



# Computer Networks Lab



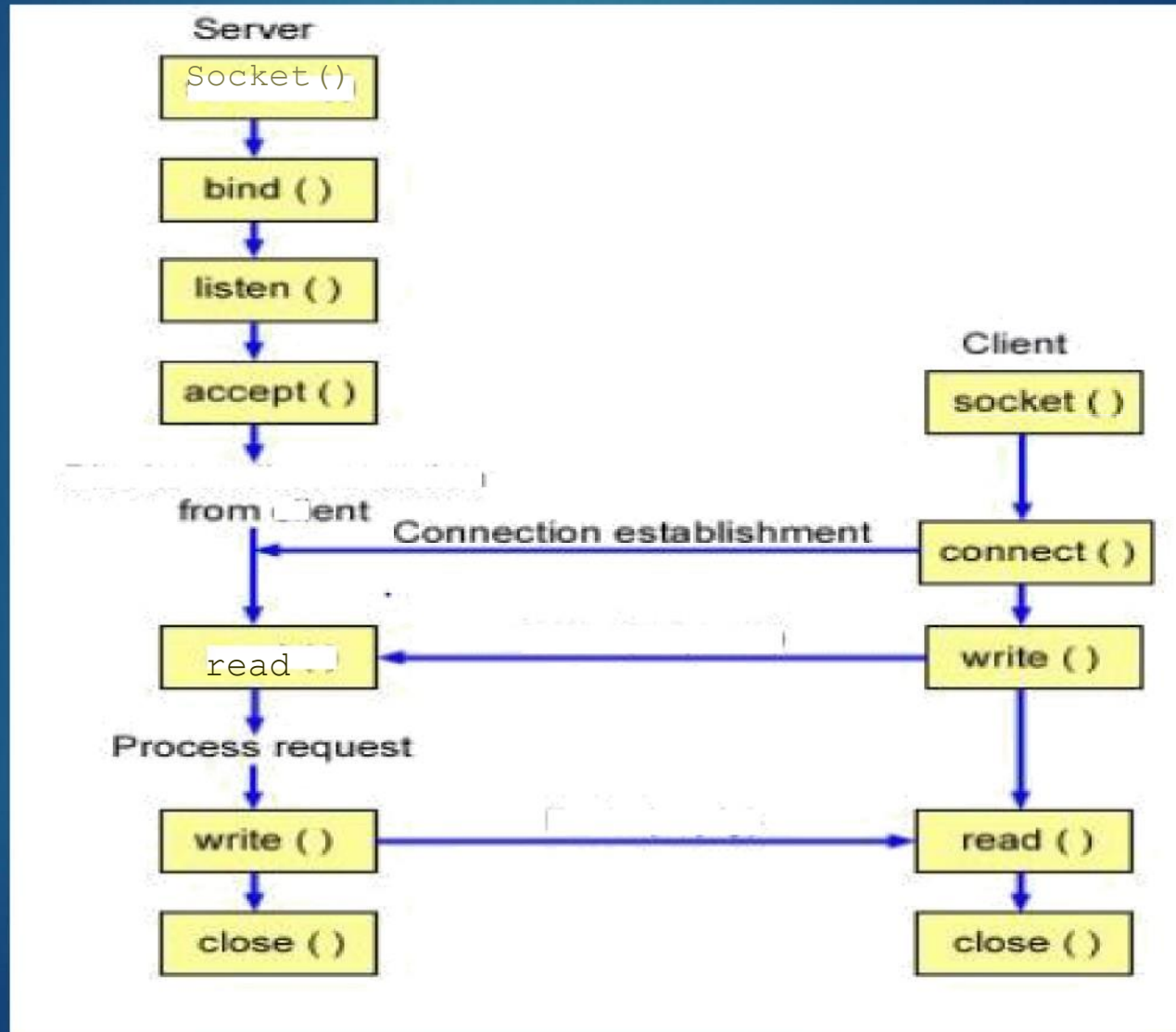
Lab 6:

Socket Programming:  
(TCP) Concurrent Servers

# Outline

1. Iterative servers
2. Fork()
3. Concurrent servers
4. Handling Child Servers

# TCP Client Server Interaction



# Iterative Server

- Server is iteratively waiting for clients. Iterative servers are fairly simple and are suitable for transactions that do not last long.
- When a client calls connect function the server gets connected with it and provides it services.
- During this period no more client can connect with the server because it is tied up with a single client.

# Iterative servers

```
create a socket- Socket()
bind to the desired port- Bind()

put the socket in passive mode - listen()
while (1)
{
    Accept the next connection-Accept()

    while (client requests)
    {
        read the request from the client - read()/recv()
        process request
        send a reply to client- write()/send()
    }

    close the client socket- close()
}
```

# Fork() System Call

When the server receives and *accepts* the client's connection, it *forks* a copy of itself called child and lets the child handle the client

The prototype of "fork" call is:

***pid\_t fork(void);***

where the definition of "pid\_t" is given in `<sys/types>` include file, and the include file `<unistd.h>` contain the declaration of "fork" system call.

When successful, "fork" returns the process ID of the child process in the parent process, and it returns a "0" in the child process.



# Fork() System Call

By checking the return value from "fork", a process can easily determine if it is a parent or child process.

A process can generate multiple child processes, but each child process has only one parent.

What happens when a fork system call is made?

- ✓ When a fork system call is made, the operating system generates a copy of the parent process which becomes the child process.
- ✓ The OS will pass to the child process most of the parent's system information (e.g. open file descriptors, environment information etc.)



# Fork() System Call

Some information, however, is unique to the child process.

Process ID

Different parent process

ID Process times

Own copy of file descriptors

Resource utilization (initialized to zero)

# Fork()

```
pid_t new_pid;  
new_pid = fork();  
switch(new_pid) {  
case -1 : /* Error */  
break;  
case 0 : /* I am child */  
break;  
default : /* I am parent */  
break;  
}
```

# But we want the child process to do something else...

```
int pid;  
int status = 0;  
  
pid = fork();  
if (pid > 0) {  
    /* parent */  
    .....  
    pid = wait(&status);  
} else {  
    /* child */  
    .....  
    exit(status);  
}
```

The **fork** syscall returns a zero to the child and the child process ID to the parent

Parent uses **wait** to sleep until child returns and

**Fork** creates an exact copy of the parent process

**Wait** variants allow wait on a specific child or

Child process passes status back to parent on **exit**, to report success/failure

# Concurrent servers

Concurrent servers can handle multiple clients at same time.

Concurrent servers fork a child to handle each client.

After the establishment of connection sever calls fork and creates child to serve the client. Then parent closes the connected socket.

# Concurrent servers

```
bind to the desired port
put the socket in listen mode (passive)
while (1)
{
    Accept the next client connection
    fork() // create a new process to handle the request if
    (child) // i.e if return value of the previous fork is zero {
        this part of the code will only be executed by the child
        process communicate with the client socket process client
        request
        close client socket
        exit // you can safely terminate the child process-> parent is still
        listening for new clients
    }
    else // i.e. for the parent process
    {
        close client socket as parent will not handle it
    }
}
```



Thank You