# Design and Analysis of Algorithms

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## **Algorithm Characteristics**

#### The necessary features of an algorithm:

- Definiteness
  - The steps of an algorithm must be precisely defined.
- Effectiveness
  - Individual steps are all do-able.
- Finiteness
  - It won't take forever to produce the desired output for any input in the specified domain.
- Output
  - Information or data that goes out.

## **Algorithm Characteristics**

#### Other important features of an algorithm:

- Input.
  - Information or data that comes in.
- Correctness.
  - Outputs correctly relate to inputs.
- Generality.
  - Works for many possible inputs.
- Efficiency.
  - Takes little time & memory to run.

# **Complexity Analysis**

## Want to achieve platform-independence

Use an abstract machine that uses *steps* of time and *units* of memory, instead of seconds or bytes

- ✓ each elementary operation takes 1 step
- ✓ each elementary instance occupies 1 unit of memory

### Standard Analysis Techniques

- Constant time statements
- Analyzing Loops
- Analyzing Nested Loops
- Analyzing Sequence of Statements
- Analyzing Conditional Statements

Simple statement sequence

```
s_1; s_2; .... ; s_k
```

- Basic Step = 1 as long as k is constant
- Simple loops

```
for (i=0; i<n; i++) { s; }
where s is Basic Step = 1</pre>
```

- Basic Steps: n
- Nested loops

```
for(i=0; i<n; i++)
for(j=0; j<n; j++) { s; }</pre>
```

• Basic Steps:  $n^2$ 

Loop index depends on outer loop index

```
for (j=0; j<=n; j++)
for (k=0; k<j; k++) {
    s;
}</pre>
```

- Inner loop executed
  - 1, 2, 3, ...., n times

$$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$

 $\therefore$  Basic Steps:  $n^2$ 

```
// Input: int A[N], array of N integers
// Output: Sum of all numbers in array A
int Sum(int A[], int N)
{
  int s=0;
  for (int i=0; i< N; i++)
    s = s + A[i];
  return s;
}</pre>
```

How should we analyse this?

```
// Input: int A[N], array of N integers
// Output: Sum of all numbers in array A
int Sum(int A[], int N){
   int [s=0]; ←
   for (int i=0; i< N; i++)
               A[i];
                                     1,2,8: Once
   return s;
                                     3,4,5,6,7: Once per each iteration
                                             of for loop, N iteration
                                     Total: 5N + 3
                                     The complexity function of the
                                     algorithm is : f(N) = 5N + 3
```

#### Constant time statements

- Simplest case: O(1) time statements
- Assignment statements of simple data types int x = y;
- Arithmetic operations:

$$x = 5 * y + 4 - z;$$

Array referencing:

$$A[j] = 5;$$

Most conditional tests:

### **Analyzing Loops**

- Any loop has two parts:
  - How many iterations are performed?
  - How many steps per iteration?

```
int sum = 0,j;
for (j=0; j < N; j++)
sum = sum +j;
```

#### **Analyzing Loops**

- Any loop has two parts:
  - How many iterations are performed?
  - How many steps per iteration?

```
int sum = 0,j;
for (j=0; j < N; j++)
sum = sum +j;
```

- Loop executes N times (0..N-1)
- O(1) steps per iteration
- Total time is N \* O(1) = O(N\*1) = O(N)

# Class Activity

```
int sum = 0;
for (i=1;i<n; i=i+2)
sum = sum +i;
```

# Class Activity

```
int sum = 0,j;
for (i=1;i<n; i=i+2)
sum = sum +i;
```

$$F(n)=n/2;$$
$$F(n)=O(n)$$

### **Analyzing Loops**

What about this for loop?

```
int sum =0, j;
for (j=0; j < 100; j++)
sum = sum +j;
```

- Loop executes 100 times
- O(1) steps per iteration
- Total time is 100 \* O(1) = O(100 \* 1) = O(100) =
   O(1)

```
int j,k;
for (j=0; j<n; j++)
for (k=0; k<n; k++)
sum += k+j;
```

```
int j,k;
for (j=0; j<n; j++)
for (k=0; k<n; k++)
sum += k+j;

n*n times
```

$$F(n)= 2n^2+2n+1$$
  
 $F(n)= O(n^2)$ 

```
int j,k;
for (j=0; j<n; j++)
for (k=0; k<n; k++)
sum += k+j;
```

- Start with outer loop:
  - How many iterations? N
  - How much time per iteration? Need to evaluate inner loop
- Inner loop uses O(N) time
- Total time is  $N * O(N) = O(N*N) = O(N^2)$

```
int j,k;
for (j=0; j<N; j++)
for (k=N; k>0; k--)
sum += k+j;
```



```
int j,k;
for (j=0; j<N; j++)
for (k=N; k>0; k--)
sum += k+j;
```

- Start with outer loop:
  - How many iterations? N
  - How much time per iteration? Need to evaluate inner loop
- Inner loop uses O(N) time
- Total time is  $N * O(N) = O(N*N) = O(N^2)$

```
int j,k;
for (j=0; j<N; j++)
for (k=N; k>0; k--)
sum += k+j;
```

- Start with outer loop:
  - How many iterations? N
  - How much time per iteration? Need to evaluate inner loop
- Inner loop uses O(N) time
- Total time is  $N * O(N) = O(N*N) = O(N^2)$

 What if the number of iterations of one loop depends on the counter of the other?

```
int j,k;
for (j=0; j < N; j++)
for (k=0; k < j; k++)
sum += k+j;
```

- Analyze inner and outer loop together:
- Number of iterations of the inner loop is:
- $0 + 1 + 2 + ... + (N-1) = O(N^2)$

# Class Activity

```
void add( int A[ ], int B[ ], int n)
  for (i=0; i<n; i++)
      for (j=0; j< n; j++)
              c[i,j] = A[i][k] + B[k][j];
```

# Class Activity

```
void multiply( int A[ ], int B[ ], int n)
  for (i=0; i<n; i++)
      for (j=0; j< n; j++)
          c[i,j]=0;
              for (k=0; k< n; ++)
                  c[i,j]+=A[i][k]*b[k][j];
```