

Merge Sort

①

→ First we need to know merging.

Combining 2 list in 1

e.g

List A	List B
2	5
8	9
15	12
18	17

The list can be of same & different sizes. But, they are sorted.

For merging, we ~~may~~ need another list c

These list can be array / linked list

A	B	C
$i \rightarrow 2$	$5 \leftarrow j$	k
8	9	
15	12	
18	17	

Compare i with j if i is smaller then copy in k

(2)

A	B	C
2	5	2
→ 8	9	14
15	12	
18	17	

A	B	C
2	5	2
i → 8	9 ← j	5
15	12	← 14
18	17	

Continue until ~~the~~ all sorted.
in the end,

A	B	C
2	5	2
8	9	5
15	12	8
i → 18	17	9
	← j	12
		15
		17
		← 14

if one list is finished, copy
remaining items of the list

so C list is with sorted elements

This was merging

Assume

A has m elements

B has n elements

so total elements sorted are

$$\underline{\underline{O(m+n)}}$$

Now Algorithm of merge.

Algorithm merge ($\overbrace{A, B}^{\text{List}}, \overbrace{m, n}^{\text{size}}$)

{
 $i = 1, j = 1, k = 1$ \therefore Assuming index 1
 while ($i \leq m$ and $j \leq n$) so far pseudocode.
 { if ($A[i] < B[j]$) \rightarrow until one list finishes.
 $C[k] = A[i]$ this will happen
 $k++ ; i++$ } $\rightarrow C[k++] + A[i++]$

else .

$C[k++] = B[j++]$
 $\}$

1 Once while is done, there might be a possibility that list have remaining item.

We don't know which one so we need to copy all remaining elements.
 (Assume List A have 5 more elements)

for (; $i \leq m$; $i++$)

$C[k++] = A[i];$

for (; $j \leq n$; $j++$)

$C[k++] = B[j];$
 $\}$

==

⇒ What if we have more than 2 list?

4

A	B	C	D
4	3	8	2
6	5	10	4
12	9	16	18

→ After comparison 1

Sorted List

2

3

4

So this one is

4 way merging
or generalize as

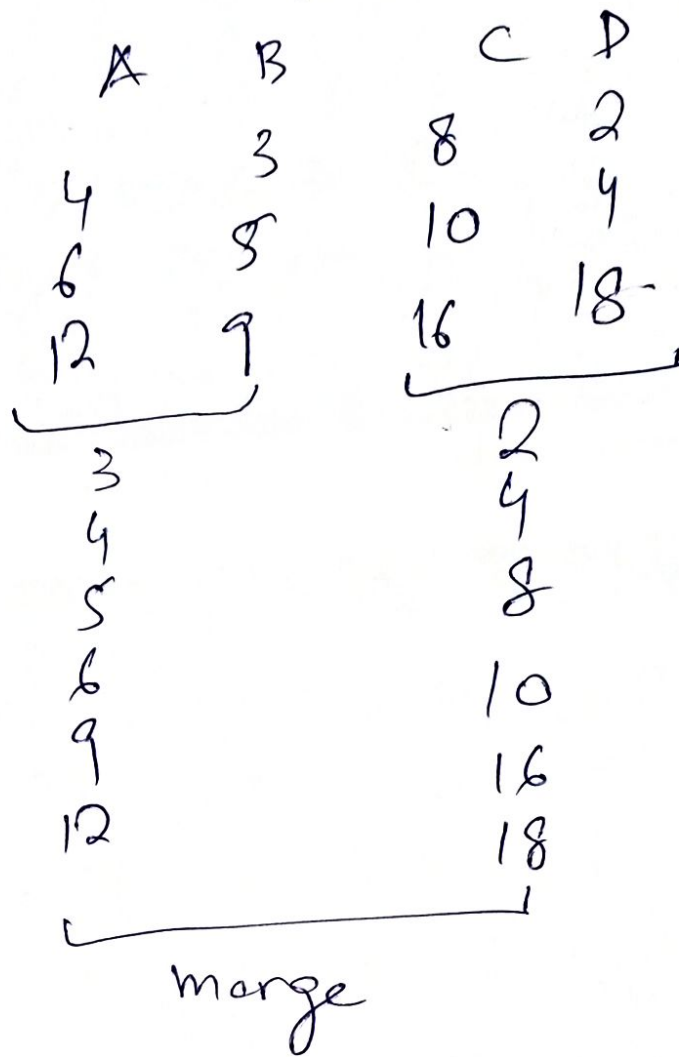
m - way merging

= mostly 2 way merging.

Now we will sort above one
with 2 way merging.

merge.

5



So take pair & do it again

Can we do it another way?

merge (AB) then (C) then (D)

or (CD) then B then A

Merge Sort

(6)

A

1	2	3	4	5	6	7	8
9	3	7	5	6	4	8	2

Done by divide & conquer / recursive

if greater than 1 element, we merge sort.

Algorithm Merge Sort ($\overset{\text{start}}{\uparrow} 1, \overset{\text{end}}{\uparrow} h$)

{ if ($1 < h$)

{

mid = $(1+h)/2$

merge Sort ($1, \text{mid}$)

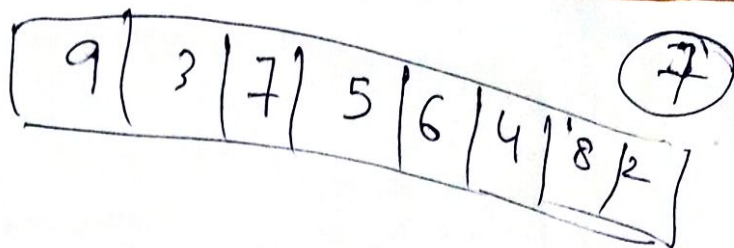
merge Sort ($\text{mid}+1, h$):

merge ($\overset{A}{1}, \text{mid}, \overset{B}{h}$)

}

}

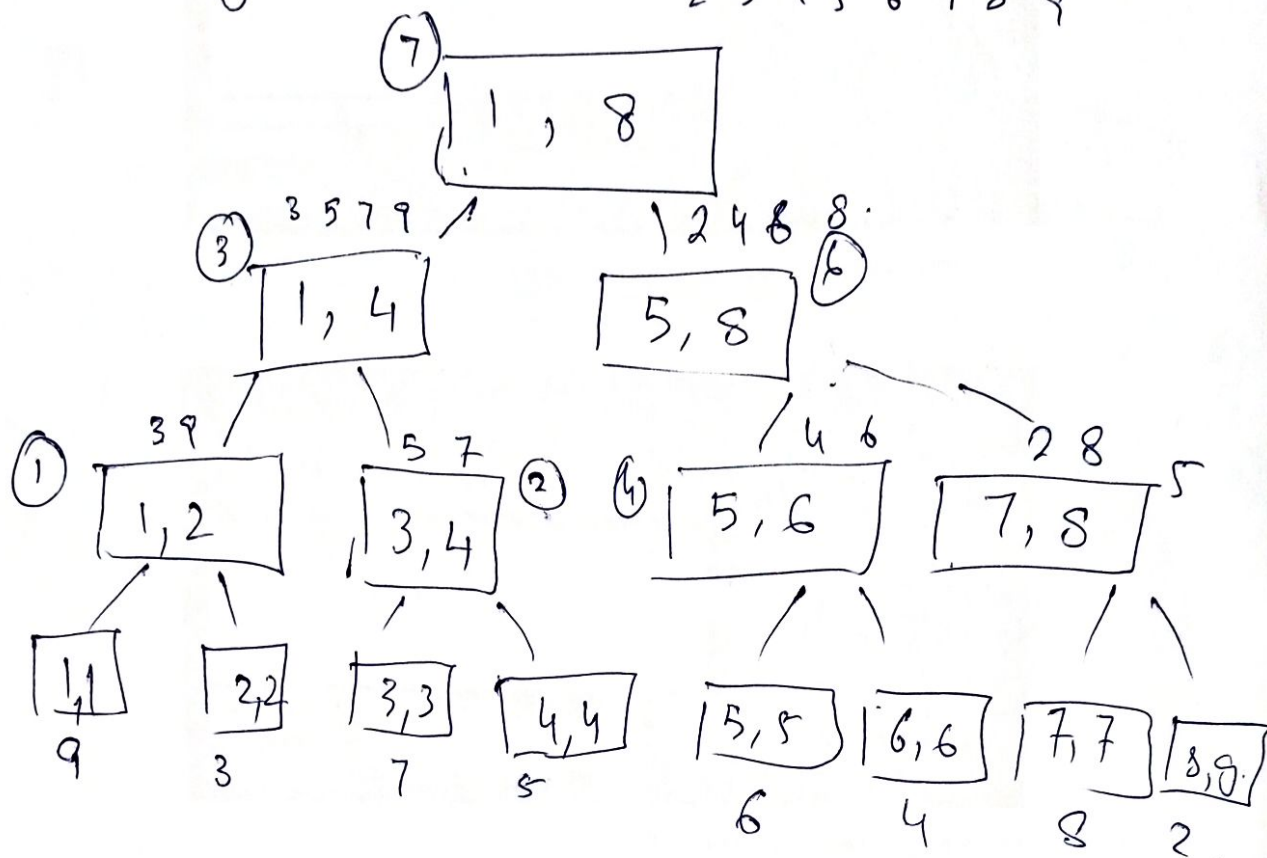
So now



low = 1

high = 8

2 3 4 5 6 7 8 9



① → First merge

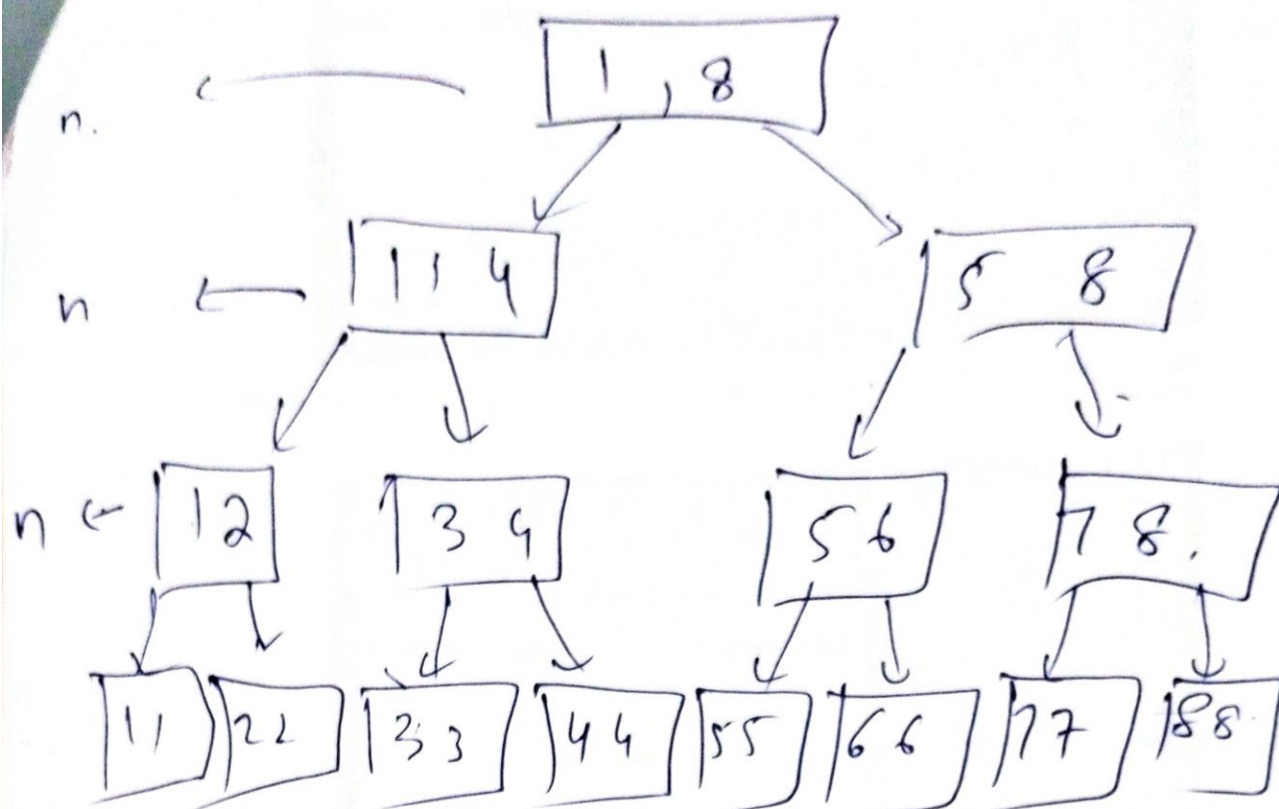
② merge

③

So divide & conquer

Now Time Analysis

(2)



How many elements merged.
Last layer is only split element
How many leads ~~of~~ for 8 elements
= n so, ~~$n \log n$~~ $(n \log n)$

~~Algo~~ Algo Complexity 2

$T(n)$ — Algo merge Sort(l, h)
if ($l < h$)

mid ($l < h$)

~~$T(h)$~~ mid = $\frac{l+h}{2}$

$T(\frac{n}{2})$ merge Sort (l, mid)
 $T(\frac{n}{2})^n$ ($mid+1, h$)

merge (l mid h)

$$T(n) = \begin{cases} 2 \left(\frac{T(n)}{2} \right) + n & n = 1 \\ & n > 1 \end{cases}$$

master theorem.

(10)

$$a = 2$$

$$b = 2$$

$$f(n) = 2$$

$$k = 1$$

$$p = 0$$

$$\log_b a = \log_2 2 = 1$$

~~is~~

$$\log_2 2 = k$$

so

$$= O(n \log n)$$