CMPSC 40, Section 6

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To Jiyu: remember to record the section.

1 Schedule

If you feel you need clarifications for problems in HW5, please ask your questions in Piazza. This will help other students too.

Problem from last HW (HW4) - 10 min

Inductions - 5-10 \min

Recursive Algorithms - 20-25 min

Q&A

2 Problem from HW4

Problem 6

Martin correctly multiplied two octal numbers, but forgot to write down the most significant digit d of the product. Can you help him?

$$(12345)_8 \times (54321)_8 = (d\,17743365)_8$$

Show your work.

• Most of you are using the classical multiplication algorithm you have learned in elementary school. What's the time complexity of the multiplication algorithm?

- Essentially, to multiply any two *n*-digit numbers, you'll write down *n*-rows of n/(n+1)-digit numbers.
- If addition/multiplication of two single digits (in 0-8) uses 1 step, then the time complexity is $O(n^2)$.

A better solution

- Now using modular arithmetic I can give an alternative and more "efficient" solution.
- Let's represent $(12345)_8$ as $1 \times 8^4 + 2 \times 8^3 + 3 \times 8^2 + 4 \times 8 + 5$. What happens if I compute its modulo 7?
- Since $8 \equiv 1 \mod 7$, it becomes $(1 \times 1 + 2 \times 1 + 3 \times 1 + 4 \times 1 + 5 \times 1) \mod 7 = (1 + 2 + 3 + 4 + 5) \mod 7$
- We have similar results for $(54321)_8$ and $(d17743365)_8$
- $(1+2+3+4+5) \times (5+4+3+2+1) \equiv (d+1+7+7+4+3+3+6+5) \mod 7$
- Finally we have $1 \equiv d+1 \mod 7$.
- d = 7.
- Time complexity?
- 3n + 1 = O(n).
- What about decimal multiplication? mod 9.
- In general?

3 Inductions

Find a formula for $\frac{1}{1\cdot 2} + \frac{1}{2\cdot 3} + \cdots + \frac{1}{(n-1)n}$. Prove it by induction.

Alternative proof?

4 Recursive Algorithms

4.1 Example from the textbook

Fibonacci sequence is defined recursively:

$$f(0) = 0, f(1) = 1$$

$$f(n+2) = f(n) + f(n+1)$$

Question: find f(6)

- f(2) = 1
- f(3) = 1 + 1 = 2
- f(4) = 1 + 2 = 3
- f(5) = 2 + 3 = 5
- f(6) = 3 + 5 = 8

Question: Write a computer program (algorithm) to solve f(n) (recursively).

- PROCEDURE f(n):
- Input is n, output is f(n), the n-th Fibonacci number.
- check if $n \ge 0$, if not, output "undefined"
- //Base Cases:
- if n = 0, output 0; if n = 1, output 1.
- //Recursive Steps:
- otherwise, output f(n-1) + f(n-2)

How it works?

- f(6)? Find f(5) and f(4) first!
- f(5)? Find f(4) and f(3) first.
- ...
- f(0)?, 0; f(1)?, 1.

Deficiency?

Note. One deficiency is that f(4) is computed twice! For f(6) we computed f(4), for f(5) we computed f(4).

An iterative approach is already above!

4.2 Euclidean Algorithm

Goal: Find the gcd of a, b.

Fact. if c divides a and b, then it divides a mod b. i.e. For $a = m \cdot b + n$, c divides n.

Question: suppose f(a, b) outputs the gcd of a, b, what's the recursive relation here?

Question: can you write an algorithm (in pseudo code) for finding the gcd of a, b?

WLOG, assume $a \ge b$

- PROCEDURE f(a, b)
- Input is a, b, output is f(a, b)
- What's the base cases?
- What's the recursive steps?

Exercise

Write a recursive procedure (algorithm) for finding Bezout's coefficients.

This algorithm is called the extended Euclidean algorithm.