



Data Structures

Lists: Circular Queue and Linked List

Won Kim

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Circular Queue



Circular Queue (Ring Buffer)

- The queue forms a circle (ring).
- The queue may be implemented using a linked list or an array.
- If an array is used, once inserts reach the end of the array, new inserts will be made to the start of the array.

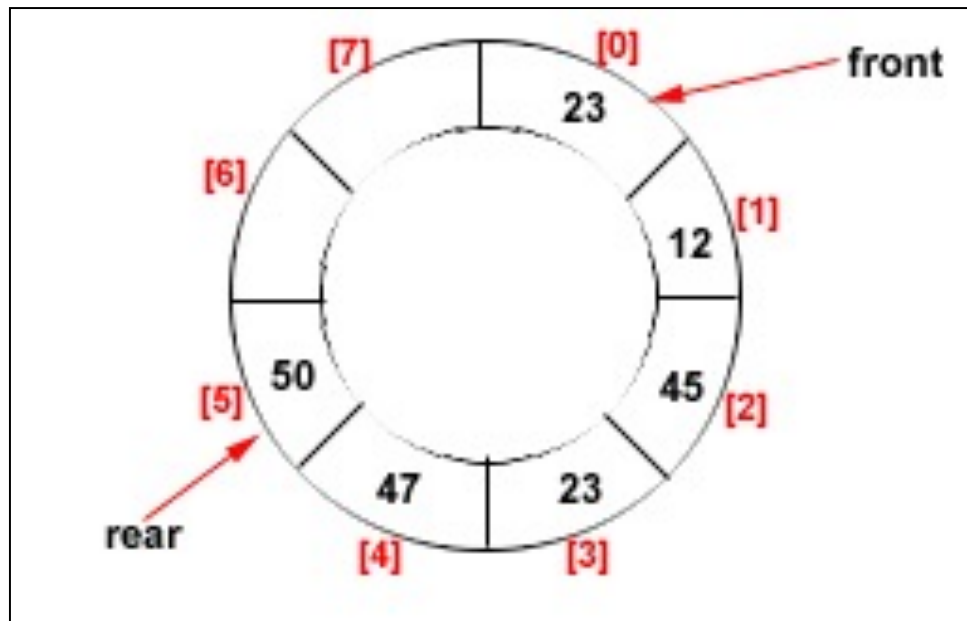


Circular Queue: Empty and Full State

- Empty state
 - $\text{front} = \text{rear}$
 - (may be anywhere on the array)
- Full state
 - $\text{front} = \text{rear}$
- Solutions
 - use all N elements of the N-element array
 - queue overwrite (i.e. force an insert)
 - queue overflow (i.e. reject insert)
 - use only N-1 elements of the N-element array

Circular Queue: Full or Overflow

- If the Queue becomes full, the Front of the Queue = the Rear, if and only if the Front has moved forward.
- Otherwise it will be "Queue overflow".





Inserting into a Circular Queue

Initially $\text{rear}=0$ and $\text{front}=0$.

- 1. If $\text{front}=0$ and $\text{rear}=0$, set $\text{front}=1$ and go to step 4.
- 2. If $\text{front}=1$ and $\text{rear}=n$ or $\text{front}=\text{rear}+1$, "circular queue overflow"(finish).
- 3. If $\text{rear}=n$, set $\text{rear}=1$ and go to step 5.
- 4. Set $\text{rear}=\text{rear}+1$
- 5. Store in rear of array.
- 6. Finish

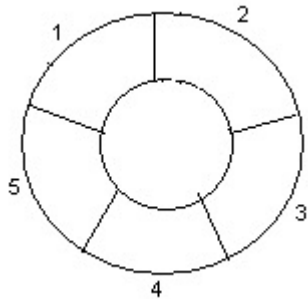


Deletion from a Circular Queue

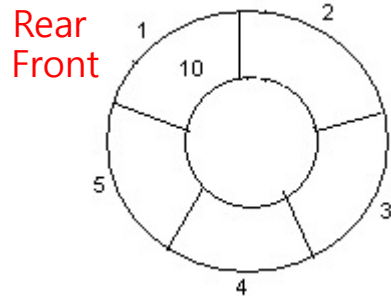
- 1. If $\text{front}=0$, "circular queue underflow" and finish.
- 2. Delete the front item
- 3. If $\text{front}=\text{rear}$, set $\text{front}=1$ and finish.
- 4. If $\text{front}=\text{rear}$, set $\text{front}=0$ and $\text{rear}=0$, and finish.
- 5. Set $\text{front}=\text{front}+1$
- 6. Finish

Example

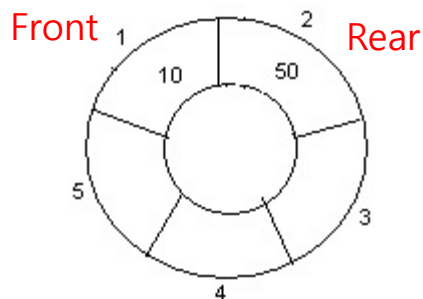
1. Initially, Rear = 0, Front = 0.



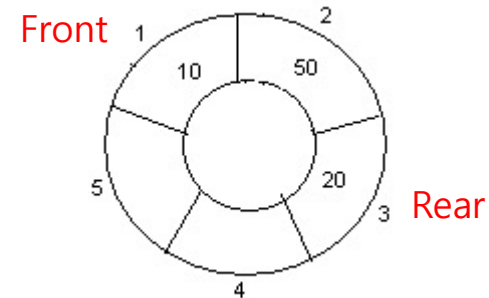
2. Insert 10, Rear = 1, Front = 1.



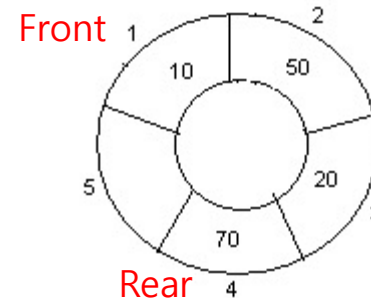
3. Insert 50, Rear = 2, Front = 1.



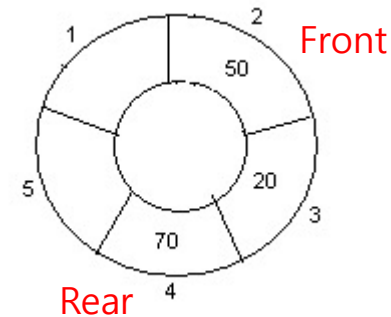
4. Insert 20, Rear = 3, Front = 1.



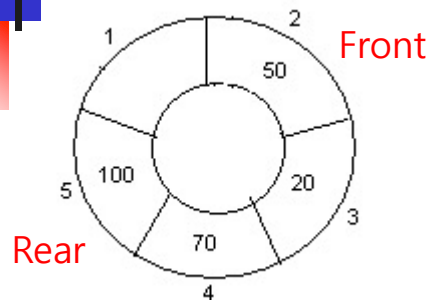
5. Insert 70, Rear = 4, Front = 1.



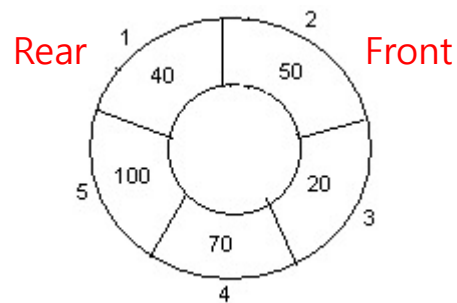
6. Delete front, Rear = 4, Front = 2.



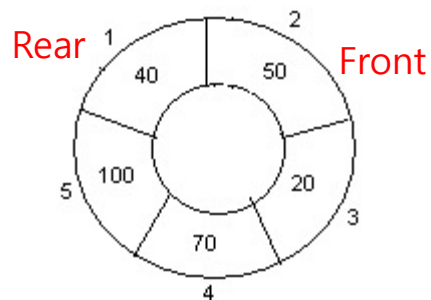
7. Insert 100, Rear = 5, Front = 2.



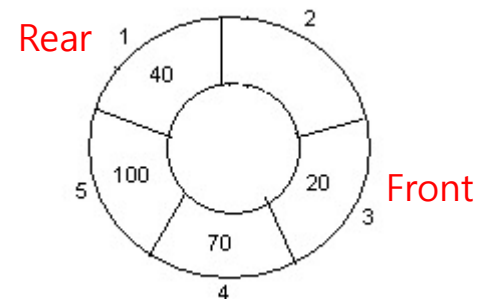
8. Insert 40, Rear = 1, Front = 2.



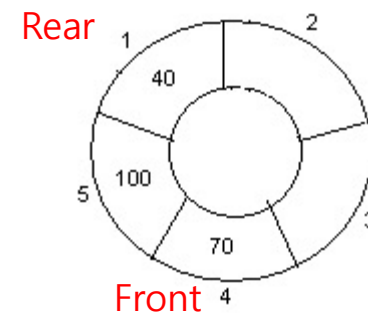
9. Insert 140, Rear = 1, Front = 2.
As $\text{Front} = \text{Rear} + 1$, so Queue overflow.



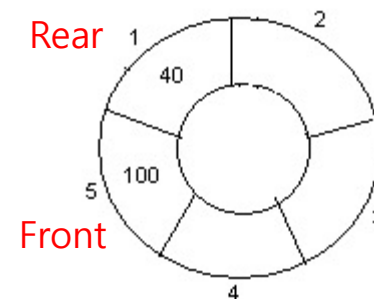
10. Delete front, Rear = 1, Front = 3.



11. Delete front, Rear = 1, Front = 4.



12. Delete front, Rear = 1, Front = 5.





Exercise

- Using buffer **overwrite**, do the same sequence of inserts and deletes shown in pages 8-9 (Draw the changes in the ring buffer).
- Using buffer **overflow**, create a ring buffer with an array of 4 elements, and do the same sequence of inserts and deletes shown in pages 8-9 (Draw the changes in the ring buffer).



Review on Linked List with struct



Programming Homework (PHW) 1



PHW 1-1 (10 points)

- Implement a ring buffer with an array of 5 elements that uses buffer overflow.
- Test the program using the sequence of inserts and deletes shown in class (pages 8-9).



PHW 1-2 (15 points)

- Implement and Test a Stack Program, Using a Singly Linked List.
 - for the same 4 functions (of Lab 1-1).
 - ** stack_full test??



PHW 1-3 (15 points)

- Implement and Test a Queue Program, Using a Singly Linked List.
 - for the same 4 functions (of Lab 1-2).
 - ** queue_full test?



PHW 1-4 (15 points)

- Implement a Virtual Integer Stack using 2 Queues.
 - Must use the Queue program (of Lab 1-2).



Solution Hints

- Create queue1 and queue2
- push: use queue1
- pop:
 - dequeue queue1, and enqueue queue2 (except the rear data)
 - dequeue queue2, and enqueue queue1



Points To Be Deducted For...

- changing the dequeue function
 - dequeue function should simply delete the **front of the queue** (and nothing else)
- directly copying data between the 2 queues
 - (e.g.) **enqueue2[k] = enqueue1[j];**
 - (should use **enqueue2(dequeue1())**)



End of Class
