

Embedded OS Implementation

Project 3

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[PART I] NPCS Implementation

Objective:

Implement the non-preemptible critical section (NPCS) based on the **RM scheduler** in uC/OS-II.

- ucos_ii.h:

```
66 #include <string.h>
67 #define SYSTEM_END_TIME 100
68
69 #define INPUT_FILE_NAME "./TaskSet.txt"
70 #define OUTPUT_FILE_NAME "./Output.txt"
71 #define MAX 20
72 #define INFO 10
73 #define RM 0
74 #define FIFO 1
75 #define ALGORITHM RM
76 #define NPCS 0
77 #define CPP 1
78 #define PROTOCOL NPCS
79
80 #define MAX_PRINTF_BUFFER 16
81
82 typedef struct task_para_set {
83     INT16U TaskID;
84     INT16U TaskArriveTime;
85     INT16U TaskExecutionTime;
86     INT16U TaskPeriodic;
87     INT16U TaskNumber;
88     INT16U TaskPriority;
89     INT16U InitArriveTime;
90     INT16U Ready;
91     INT16U job_no; //Job number
92     INT16U RemainTime; //任務剩餘時間
93     INT16U deadline; //任務deadline
94     INT16U ResponseTime; //紀錄response time
95     INT16U Delay; //任務需要被osTimeDly多久
96     INT16U BlockTime;
97     INT16U PreemptiveTime;
98     INT16U R1Lock;
99     INT16U R1UnLock;
100     INT16U R2Lock;
101     INT16U R2UnLock;
102 } task_para_set;
```

➤ 新增 task_para_set 的 member：

- BlockTime：任務被 Block 的時間
- PreemptiveTime：任務被 Preemp 的時間
- R1Lock、R1UnLock：Resource1 被 Lock、Unlock 時間
- R2Lock、R2UnLock：Resource2 被 Lock、Unlock 時間

➤ Define 使用的 PROTOCOL

- NPCS(non-preemptible critical section)
- CPP(Ceiling Priority Protocol)

- app_hooks.h :

```

105 while (ptr != NULL) {
106     TaskInfo[i] = atoi(ptr);
107     ptr = strtok_s(NULL, " ", &pTmp);
108     if (i == 0) {
109         TASK_NUMBER++;
110         TaskParameter[j].TaskID = TASK_NUMBER;
111     }
112     else if (i == 1) {
113         TaskParameter[j].InitArriveTime = TaskInfo[i];
114         TaskParameter[j].TaskArriveTime = TaskInfo[i];
115     }
116     else if (i == 2) {
117         TaskParameter[j].TaskExecutionTime = TaskInfo[i];
118         TaskParameter[j].RemainTime = TaskInfo[i];
119     }
120     else if (i == 3) {
121         TaskParameter[j].TaskPeriodic = TaskInfo[i];
122     }
123     else if (i == 4) {
124         TaskParameter[j].R1Lock = TaskInfo[i];
125     }
126     else if (i == 5) {
127         TaskParameter[j].R1UnLock = TaskInfo[i];
128     }
129     else if (i == 6) {
130         TaskParameter[j].R2Lock = TaskInfo[i];
131     }
132     else if (i == 7) {
133         TaskParameter[j].R2UnLock = TaskInfo[i];
134         TaskParameter[j].deadline = TaskParameter[j].TaskArriveTime + TaskParameter[j].TaskPeriodic;
135         TaskParameter[j].BlockTime = 0;
136         TaskParameter[j].PreemptiveTime = 0;
137         TaskParameter[j].ready = 0;
138         i++;
139     }
140     TaskParameter[j].TaskPriority = j;
141     j++;
142 }

```

➤ 讀取 Taskset.txt 資料、並初始化 BlockTime、PreemptiveTime

- App_hooks.h(App_TimeTickHook) :

```

369 void App_TimeTickHook (void)
370 {
371     #if (APP_CFG_PROBE_OS_PLUGIN_EN == DEF_ENABLED) && (OS_PROBE_HOOKS_EN > 0)
372         OSProbe_TickHook();
373     #endif
374
375     OS_TCB* ptcb;
376     task_para_set* pdata;
377     INT32U CurrentTime;
378     // 因為 Hook 是在 OSTime++ 之前執行的，所以我們要用 "下一刻的時間" 來判斷到達
379     CurrentTime = OSTime + 1;
380     ptcb = OSTCBLIST;
381     while (ptcb != (OS_TCB*)0) {
382         if (ptcb->OSTCBExtPtr != (void*)0) {
383             pdata = (task_para_set*)ptcb->OSTCBExtPtr;
384             if (CurrentTime >= pdata->InitArriveTime) {
385                 INT32U time_diff = CurrentTime - pdata->InitArriveTime;
386
387                 if (pdata->TaskPeriodic > 0 && (time_diff % pdata->TaskPeriodic) == 0) {
388                     pdata->RemainTime = pdata->TaskExecutionTime;
389                     pdata->Ready = 1;
390                 }
391             }
392         }
393         // 檢查下一個 TCB
394         ptcb = ptcb->OSTCBNext;
395     }
396 }

```

➤ 將任務下次 ArriveTime 抵達時，補充 RemainTime=ExecutionTime

➤ 檢查整個 OSTCBLIST

- main.c(task_function):

```

110 void task(void* p_arg) {
111     task_para_set* task_data = (task_para_set*)p_arg;
112     INT32U current_time;
113     INT32U executed_time;
114     INT8U err;
115     INT8U base_prio = task_data->TaskPriority;
116     INT8U current_prio_log;
117     INT8U target_prio_log;
118     current_time = OSTimeGet();
119     while (stopAlltask == 0) {
120         task_data->Ready = 1;
121         while (task_data->RemainTime > 0 && stopAlltask == 0) {
122             if (stopAlltask == 1) {
123                 break;
124             }
125         }
126         executed_time = task_data->TaskExecutionTime - task_data->RemainTime;
127         // 檢查 R1 Lock
128         if ((executed_time == task_data->R1Lock) && (task_data->R1UnLock != 0)) {
129             #if PROTOCOL == NPCS
130                 LOG_print(3, "./Output.txt", "%d\tLockResource\ttask(%2d)(%2d)\t%s\n",
131                     OSTime, task_data->TaskID, task_data->job_no, "R1");
132                 OSSchedLock(); // 禁止搶佔
133             }
134         }
135         // 檢查 R2 Lock
136         if ((executed_time == task_data->R2Lock) && (task_data->R2UnLock != 0)) {
137             #if PROTOCOL == NPCS
138                 LOG_print(3, "./Output.txt", "%d\tLockResource\ttask(%2d)(%2d)\t%s\n",
139                     OSTime, task_data->TaskID, task_data->job_no, "R2");
140                 OSSchedLock(); // 禁止搶佔
141             }
142         }
143         // 檢查 R1 UnLock
144         if ((executed_time == task_data->R1UnLock) && (task_data->R1UnLock != 0)) {
145             #if PROTOCOL == NPCS
146                 LOG_print(3, "./Output.txt", "%d\tUnlockResource\ttask(%2d)(%2d)\t%s\n",
147                     OSTime, task_data->TaskID, task_data->job_no, "R1");
148                 OSSchedUnlock(); // 恢復排程
149             }
150         }
151         // 檢查 R2 UnLock
152         if ((executed_time == task_data->R2UnLock) && (task_data->R2UnLock != 0)) {
153             #if PROTOCOL == NPCS
154                 LOG_print(3, "./Output.txt", "%d\tUnlockResource\ttask(%2d)(%2d)\t%s\n",
155                     OSTime, task_data->TaskID, task_data->job_no, "R2");
156                 OSSchedUnlock(); // 恢復排程
157             }
158         }
159     }
160 }

```

- 檢查 executed_time 是否為 Lock 或 UnLock 的時間

- Lock : OSSchedLock() 禁止搶斷
- UnLock : OSSchedUnlock() 恢復排程

- Print 出 Lock、UnLock 訊息

- os_core.c(OS_SchedNew):

```

1979  static void OS_SchedNew(void)
1980  {
1981  #if OS_LOWEST_PRIO <= 63u                /* See if we support up to 64 tasks */
1982      INT8U y;
1983      static INT8U first_check = 0u;
1984
1985      if (first_check == 0u) {
1986          OS_TCB* ptcb;
1987          INT32U delay_ticks;
1988          first_check = 1u;
1989          ptcb = OSTCBLIST;
1990
1991          while (ptcb != (OS_TCB*)0 && (ptcb->OSTCBEExtPtr != NULL)) {
1992              delay_ticks = ((task_para_set*)(ptcb->OSTCBEExtPtr))->InitArriveTime;
1993              if (delay_ticks > 0) {
1994                  ptcb->OSTCBDly = delay_ticks;
1995                  // 從 Ready Table 中移除該任務
1996                  if ((OSRdyTbl[ptcb->OSTCBY] &= ~ptcb->OSTCBBitX) == 0u) {
1997                      OSRdyGrp &= ~ptcb->OSTCBBitY;
1998                  }
1999                  ptcb = ptcb->OSTCBNext; /* 移至下一個任務 */
2000              }
2001          }
2002      }
2003
2004      y = OSUnMapTbl[OSRdyGrp];
2005      OSPrioHighRdy = (INT8U)((y << 3u) + OSUnMapTbl[OSRdyTbl[y]]);

```

- 原始設計為在 task_function 中檢查 InitArriveTime，但會有低優先級任務來不及移除 RdyList 導致訊息印錯的狀況
- InitArriveTime != 0 的情況，OS_Start()後第一次呼叫 OS_SchedNew()搜尋 OSTCBLIST，如果 InitArriveTime != 0，將其從 RdyList 中移除，並設定其 OSTCBDly

- os_core.c(OS_TimeTick):

```

#if PROTOCOL == NPCS
    if (ptcb->OSTCBEExtPtr != (void*)0) {
        task_para_set* p_task_data = (task_para_set*)(ptcb->OSTCBEExtPtr);
        if (ptcb->OSTCBStat == OS_STAT_RDY && ptcb != OSTCBCur && p_task_data->RemainTime > 0 && (currentTime > p_task_data->InitArriveTime))
            if (OSPrioCur != 63) {
                // Blocking Time
                if (ptcb->OSTCBPrio < OSTCBCur->OSTCBPrio) {
                    p_task_data->BlockTime++;
                    printf("kd\ttask(%2d)\tblockTime +1\n", OSTimeGet(), p_task_data->TaskID);
                }
                // Preemption Time
                else {
                    p_task_data->PreemptiveTime++;
                    printf("kd\ttask(%2d)\tpreemptiveTime +1\n", OSTimeGet(), p_task_data->TaskID);
                }
            }
    }
    ptcb = ptcb->OSTCBNext;

```

- 搜尋 OSTCBLIST 當 ptcb 指向的 task 的 Priority 小於正在執行的 taskPrio 且 RemainTime>0(代表還沒做完)
 - 任務被 Block : BlockTime++
- 搜尋 OSTCBLIST 當 ptcb 指向的 task 的 Priority 大於正在執行的 taskPrio 且 RemainTime>0(代表還沒做完)
 - 任務被 Preempt : PreemptTime++

[PART II] CPP Implementation

Objective:

Implement the ceiling- priority protocol (CPP) based on the **RM scheduler** in uC/OS-II.

- main.c(main_function) :

```
320  #elif PROTOCOL == CPP
321      INT8U R1_Prio = 255;
322      INT8U R2_Prio = 255;
323      for (int i = 0; i < TASK_NUMBER; i++) {
324          TaskParameter[i].TaskPriority = 3*(i+1);
325      }
326
327      for (int i = 0; i < TASK_NUMBER; i++) {
328          // 檢查該 Task 是否使用 R1
329          if (TaskParameter[i].R1Lock < TaskParameter[i].R1UnLock) {
330              if (TaskParameter[i].TaskPriority < R1_Prio) {
331                  R1_Prio = TaskParameter[i].TaskPriority;
332              }
333          }
334          // 檢查該 Task 是否使用 R2
335          if (TaskParameter[i].R2Lock < TaskParameter[i].R2UnLock) {
336              if (TaskParameter[i].TaskPriority < R2_Prio) {
337                  R2_Prio = TaskParameter[i].TaskPriority;
338              }
339          }
340      }
341
342      R1_Ceiling = R1_Prio - 1;
343      R2_Ceiling = R2_Prio - 2;
344
345      R1 = OSMutexCreate(R1_Ceiling, &err);
346      R2 = OSMutexCreate(R2_Ceiling, &err);
347
348      for (int n = 0; n < TASK_NUMBER; n++) {
349          Task_STK[n] = malloc(TASK_STACKSIZE * sizeof(int));
350          OSTaskCreateExt(task,
351              &TaskParameter[n],
352              &Task_STK[n][TASK_STACKSIZE - 1],
353              TaskParameter[n].TaskPriority,
354              TaskParameter[n].TaskID,
355              &Task_STK[n][0],
356              TASK_STACKSIZE,
357              &TaskParameter[n],
358              (OS_TASK_OPT_STK_CHK | OS_TASK_OPT_STK_CLR));
359      }
360  #endif
```

- 在 CPP Protocol 中，當任務由 Periodic 由低排到高後(RM)，將任務的 Priority 設為 $3*(i+1)$ ，需要留 Priority 給 R1、R2
 - Task priority : 3、6、9...
- 檢查是否有用到 R1、R2，如果有，先將 R1、R2 Prio 設為當前 Task 的 Priority
 - 將 R1、R2 的 Prio-index 即為 R1、R2 的 Ceiling Priority
- 建立 Mutex 以及 periodic 任務

- main.c(task_function):

```

119 while (stopAlltask == 0) {
120     task_data->Ready = 1;
121     while (task_data->RemainTime > 0 && stopAlltask == 0) {
122         if (stopAlltask == 1) {
123             break;
124         }
125
126         executed_time = task_data->TaskExecutionTime - task_data->RemainTime;
127         // 檢查 R1 Lock
128         if ((executed_time == task_data->R1Lock) && (task_data->R1UnLock != 0)) {
129             #if PROTOCOL == NPCS
130                 LOG_print(3, "../Output.txt", "%d\tLockResource\ttask(%2d)(%2d)\t%s\n",
131                     OSTime, task_data->TaskID, task_data->job_no, "R1");
132                 OSSchedLock(); // 禁止搶佔
133             #elif PROTOCOL == CPP
134                 current_prio_log = base_prio;
135                 if (executed_time > task_data->R2Lock && executed_time < task_data->R2UnLock) {
136                     if (R2_Ceiling < current_prio_log) current_prio_log = R2_Ceiling;
137                 }
138                 target_prio_log = (R1_Ceiling < current_prio_log) ? R1_Ceiling : current_prio_log;
139                 LOG_print(3, "../Output.txt", "%d\tLockResource\ttask(%2d)(%2d)\t%s\t%d to %d\n",
140                     OSTime, task_data->TaskID, task_data->job_no, "R1", current_prio_log, target_prio_log);
141                 OSMutexPend(R1, 0, &err);
142             #endif
143         }
144     }
145 }

```

```

146 // 檢查 R2 Lock
147 if ((executed_time == task_data->R2Lock) && (task_data->R2UnLock != 0)) {
148     #if PROTOCOL == NPCS
149         LOG_print(3, "../Output.txt", "%d\tLockResource\ttask(%2d)(%2d)\t%s\n",
150             OSTime, task_data->TaskID, task_data->job_no, "R2");
151         OSSchedLock(); // 禁止搶佔
152     #elif PROTOCOL == CPP
153         // Lock R2 前·檢查是否持有 R1
154         current_prio_log = base_prio;
155         if (executed_time > task_data->R1Lock && executed_time < task_data->R1UnLock) {
156             if (R1_Ceiling < current_prio_log) current_prio_log = R1_Ceiling;
157         }
158         // Lock R2 後·再跟 R2 Ceiling 比
159         target_prio_log = (R2_Ceiling < current_prio_log) ? R2_Ceiling : current_prio_log;
160         LOG_print(3, "../Output.txt", "%d\tLockResource\ttask(%2d)(%2d)\t%s\t%d to %d\n",
161             OSTime, task_data->TaskID, task_data->job_no, "R2", current_prio_log, target_prio_log);
162         OSMutexPend(R2, 0, &err);
163     #endif
164 }
165 }

```

```

166 // 檢查 R1 UnLock
167 if ((executed_time == task_data->R1UnLock) && (task_data->R1UnLock != 0)) {
168     #if PROTOCOL == NPCS
169         LOG_print(3, "../Output.txt", "%d\tUnlockResource\ttask(%2d)(%2d)\t%s\n",
170             OSTime, task_data->TaskID, task_data->job_no, "R1");
171         OSSchedUnlock(); // 恢復排程
172     #elif PROTOCOL == CPP
173         // 檢查是否持有 R2
174         current_prio_log = base_prio;
175         if (R1_Ceiling < current_prio_log) current_prio_log = R1_Ceiling; // 因為正在解鎖 R1·肯定持有 R2
176         if (executed_time > task_data->R2Lock && executed_time < task_data->R2UnLock) {
177             if (R2_Ceiling < current_prio_log) current_prio_log = R2_Ceiling;
178         }
179         // Unlock 後·R1 移除·檢查 R2
180         target_prio_log = base_prio;
181         if (executed_time > task_data->R2Lock && executed_time < task_data->R2UnLock) {
182             if (R2_Ceiling < target_prio_log) target_prio_log = R2_Ceiling;
183         }
184         LOG_print(3, "../Output.txt", "%d\tUnlockResource\ttask(%2d)(%2d)\t%s\t%d to %d\n",
185             OSTime, task_data->TaskID, task_data->job_no, "R1", current_prio_log, target_prio_log);
186         OSMutexPost(R1);
187     #endif
188 }
189 }

```

```

190 // 檢查 R2 UnLock
191 if ((executed_time == task_data->R2UnLock) && (task_data->R2UnLock != 0)) {
192     #if PROTOCOL == NPCS
193         LOG_print(3, "../Output.txt", "%d\tUnlockResource\ttask(%2d)(%2d)\t%s\n",
194             OSTime, task_data->TaskID, task_data->job_no, "R2");
195         OSSchedUnlock(); // 恢復排程
196     #elif PROTOCOL == CPP
197         // 檢查 R1
198         current_prio_log = base_prio;
199         if (R2_Ceiling < current_prio_log) current_prio_log = R2_Ceiling; // 肯定持有 R2
200         if (executed_time > task_data->R1Lock && executed_time < task_data->R1UnLock) {
201             if (R1_Ceiling < current_prio_log) current_prio_log = R1_Ceiling;
202         }
203         // Unlock 後·R2 移除·檢查有沒有 R1
204         target_prio_log = base_prio;
205         if (executed_time > task_data->R1Lock && executed_time < task_data->R1UnLock) {
206             if (R1_Ceiling < target_prio_log) target_prio_log = R1_Ceiling;
207         }
208         LOG_print(3, "../Output.txt", "%d\tUnlockResource\ttask(%2d)(%2d)\t%s\t%d to %d\n",
209             OSTime, task_data->TaskID, task_data->job_no, "R2", current_prio_log, target_prio_log);
210         OSMutexPost(R2);
211     #endif
212 }
213 }

```

➤ 變數定義:

- base_prio: Task 本身的原始優先權
- current_prio_log: 動作前的當下優先權
- target_prio_log: 動作後的目標優先權
- Ceiling: 資源的最高優先權上限

- Lock Resource：準備鎖定 Resource 時，優先權會從目前的狀態提升到包含 Resource Ceiling 的狀態
 - 計算 Current：檢查目前是否已經持有其他資源->若有，Current = max_prio(Base, R2_Ceiling)
 - 計算 Target：比較 Current 與 R1 的 Ceiling->Target = max_prio(Current, R1_Ceiling)
- Unlock Resource：準備釋放 R1 時，優先權會從目前的高優先權降回剩餘資源的最高 Ceiling 或是原始優先權
 - 計算 Current：因為正在持有 R1（且可能持有 R2），Current 是所有持有資源中最高的 Ceiling
 - 計算 Target：釋放 R1 後，回歸 Base，再檢查是否還持有 R2。若還持有 R2，Target = max_prio(Base, R2_Ceiling)

- os_core.c(OS_TimeTick)：

```

1110  #elif PROTOCOL == CPP
1111      if (ptcb->OSTCBExtPtr != (void*)0) {
1112          task_para_set* p_task_data = (task_para_set*)(ptcb->OSTCBExtPtr);
1113          if (ptcb != OSTCBCur && p_task_data->RemainTime > 0 &&
1114              (currentTime > p_task_data->InitArriveTime) &&
1115              p_task_data->Ready == 1 &&
1116              currentTime > p_task_data->TaskArriveTime) {
1117
1118              task_para_set* p_running_data = (task_para_set*)0;
1119              if (OSTCBCur->OSTCBExtPtr != (void*)0) {
1120                  p_running_data = (task_para_set*)(OSTCBCur->OSTCBExtPtr);
1121              }
1122
1123              if (p_running_data != (task_para_set*)0) {
1124                  INT8U my_base_prio = p_task_data->TaskPriority;
1125                  INT8U running_base_prio = p_running_data->TaskPriority;
1126                  if (ptcb->OSTCBStat == OS_STAT_RDY) {
1127                      // Blocking
1128                      if (running_base_prio > my_base_prio) {
1129                          p_task_data->BlockTime++;
1130                          printf("%d\ttask(%2d)\tblockTime +1 (Prio Inv)\n", OSTimeGet(), p_task_data->TaskID);
1131                      }
1132                      // Preemption
1133                      else {
1134                          p_task_data->PreemptiveTime++;
1135                          printf("%d\ttask(%2d)\tpreemptiveTime +1\n", OSTimeGet(), p_task_data->TaskID);
1136                      }
1137                  }
1138              }
1139          }
1140      }
1141  #endif
1142      ptcb = ptcb->OSTCBNext; /* Point at next TCB in TCB list */

```

- 判斷 Block Time 與 PreemptiveTime
 - Running_Base_Prio < My_Base_Prio (執行任務的原始優先級比我低)。代表執行任務是因為持有資源，優先權被暫時提升才擋住=>這屬於優先權反轉，計入 BlockTime。
 - Running_Base_Prio > My_Base_Prio (執行任務的原始優先級本來就比較高)。代表這是正常的排程搶佔，計入 Preemptive Time。

- os_mutex.c(OSMutexPend):

```

495     OS_ENTER_CRITICAL();
496     pcp = (INT8U)(pevent->OSEventCnt >> 8u);
497
498
499     if ((INT8U)(pevent->OSEventCnt & OS_MUTEX_KEEP_LOWER_8) == OS_MUTEX_AVAILABLE) {
500         pevent->OSEventCnt &= OS_MUTEX_KEEP_UPPER_8;
501         pevent->OSEventCnt |= OSTCBCur->OSTCBPrio;
502         pevent->OSEventPtr = (void*)OSTCBCur;
503         if (pcp != OS_PRIO_MUTEX_CEIL_DIS) {
504             if (OSTCBCur->OSTCBPrio > pcp) {
505                 if ((OSRdyTbl[OSTCBCur->OSTCBY] & OSTCBCur->OSTCBBitX) != 0) {
506                     OSRdyTbl[OSTCBCur->OSTCBY] &= (OS_PRIO)-OSTCBCur->OSTCBBitX;
507                     if (OSRdyTbl[OSTCBCur->OSTCBY] == 0) {
508                         OSRdyGrp &= (OS_PRIO)-OSTCBCur->OSTCBBitY;
509                     }
510                 }
511                 OSTCBCur->OSTCBPrio = pcp;
512             }
513         }
514         #if OS_LOWEST_PRIO <= 63u
515             OSTCBCur->OSTCBY = (INT8U)(OSTCBCur->OSTCBPrio >> 3u);
516             OSTCBCur->OSTCBX = (INT8U)(OSTCBCur->OSTCBPrio & 0x07u);
517         #else
518             OSTCBCur->OSTCBY = (INT8U)((INT8U)(OSTCBCur->OSTCBPrio >> 4u) & 0xFFu);
519             OSTCBCur->OSTCBX = (INT8U)(OSTCBCur->OSTCBPrio & 0x0Fu);
520         #endif
521         OSTCBCur->OSTCBBitY = (OS_PRIO)(1uL << OSTCBCur->OSTCBY);
522         OSTCBCur->OSTCBBitX = (OS_PRIO)(1uL << OSTCBCur->OSTCBX);
523         OSRdyGrp |= OSTCBCur->OSTCBBitY;
524         OSRdyTbl[OSTCBCur->OSTCBY] |= OSTCBCur->OSTCBBitX;
525
526         OSTCBPrioTbl[pcp] = OSTCBCur;
527     }
528
529     // if (pcp != OS_PRIO_MUTEX_CEIL_DIS) {
530     //     mprio = (INT8U)(pevent->OSEventCnt & OS_MUTEX_KEEP_LOWER_8); /* Get priority of mutex owner */
531     //     ptcb = (OS_TCB*)(pevent->OSEventPtr); /* Point to TCB of mutex owner */
532     //     if (ptcb->OSTCBPrio > pcp) { /* Need to promote prio of owner? */
533     //         if (mprio > OSTCBCur->OSTCBPrio) {
534     //             y = ptcb->OSTCBY;

```

- 獲取 Mutex 時立即提升優先級

- 成功獲取 Mutex 的當下，強制將當前任務的優先級提升至該 Mutex 定義的 Ceiling Priority，確保任務以最高優先級運行

- 原有的優先級繼承 (Priority Inheritance) 偵測。由於任務持有資源時已處於最高優先級，因此不再需要當 Mutex 發生競爭時才動態提升優先級的邏輯。

- os_mutex.c(OSMutexPost):

```

647     INT8U OSMutexPost(OS_EVENT* pevent)
648     {
649         INT8U pcp; /* Priority ceiling priority */
650         INT8U prio;
651         BOOLEAN prio_restored = OS_FALSE; /* [新增] 標記是否發生了優先權還原 */
652         #if OS_CRITICAL_METHOD == 3u /* Allocate storage for CPU status register */
653             OS_CPU_SR cpu_sr = 0u;
654         #endif
655
656         if (OSIntNesting > 0u) { /* See if called from ISR ... */
657             return (OS_ERR_POST_ISR); /* ... can't POST mutex from an ISR */
658         }
659         #if OS_ARG_CHK_EN > 0u
660             if (pevent == (OS_EVENT*)0) { /* Validate 'pevent' */
661                 return (OS_ERR_PEVENT_NULL);
662             }
663         #endif
664
665         OS_TRACE_MUTEX_POST_ENTER(pevent);
666
667         if (pevent->OSEventType != OS_EVENT_TYPE_MUTEX) { /* Validate event block type */
668             OS_TRACE_MUTEX_POST_EXIT(OS_ERR_EVENT_TYPE);
669             return (OS_ERR_EVENT_TYPE);
670         }
671         OS_ENTER_CRITICAL();
672         pcp = (INT8U)(pevent->OSEventCnt >> 8u); /* Get priority ceiling priority of mutex */
673         prio = (INT8U)(pevent->OSEventCnt & OS_MUTEX_KEEP_LOWER_8); /* Get owner's original priority */
674         if (OSTCBCur != (OS_TCB*)pevent->OSEventPtr) { /* See if posting task owns the MUTEX */
675             OS_EXIT_CRITICAL();
676             OS_TRACE_MUTEX_POST_EXIT(OS_ERR_NOT_MUTEX_OWNER);
677             return (OS_ERR_NOT_MUTEX_OWNER);
678         }
679         if (pcp != OS_PRIO_MUTEX_CEIL_DIS) {
680             if (OSTCBCur->OSTCBPrio == pcp) {
681                 OS_TRACE_MUTEX_TASK_PRIO_DISINHERIT(OSTCBCur, prio);
682                 OSMutexRdyAtPrio(OSTCBCur, prio); /* Restore the task's original priority */
683                 prio_restored = OS_TRUE; /* 降低了自己的優先權 */
684             }
685             OSTCBPrioTbl[pcp] = OS_TCB_RESERVED; /* Reserve table entry */
686         }
687         if (pevent->OSEventGrp != 0u) {

```



```

686     OSTCBPrioTbl[pcp] = OS_TCB_RESERVED; /* Reserve table entry */
687 }
688 if (pevent->OSEventGrp != 0u) {
689     prio = OS_EventTaskRdy(pevent, (void*)0, OS_STAT_MUTEX, OS_STAT_PEND_OK);
690     pevent->OSEventCnt &= OS_MUTEX_KEEP_UPPER_8;
691     pevent->OSEventPtr = prio;
692     pevent->OSEventPtr = OSTCBPrioTbl[prio];
693     if ((pcp != OS_PRIO_MUTEX_CEIL_DIS) && (prio <= pcp)) {
694         OS_EXIT_CRITICAL();
695         OS_Sched();
696         OS_TRACE_MUTEX_POST_EXIT(OS_ERR_PCP_LOWER);
697         return (OS_ERR_PCP_LOWER);
698     }
699 }
700 else {
701     OS_EXIT_CRITICAL();
702     OS_Sched(); /* Find highest priority task ready to run */
703     OS_TRACE_MUTEX_POST_EXIT(OS_ERR_NONE);
704     return (OS_ERR_NONE);
705 }
706 }
707 pevent->OSEventCnt |= OS_MUTEX_AVAILABLE;
708 pevent->OSEventPtr = (void*)0;
709 OS_EXIT_CRITICAL();
710 if (prio_restored == OS_TRUE) {
711     OS_Sched();
712 }
713
714 OS_TRACE_MUTEX_POST_EXIT(OS_ERR_NONE);
715 return (OS_ERR_NONE);
716 }

```

➤ 新增變數 `prio_restored` 來追蹤優先級變化

- 修改內容：宣告了一個 `BOOLEAN prio_restored = OS_FALSE`；當檢測到任務需要將優先級從 Ceiling Priority 還原回原始優先級時，將此旗標設為 `OS_TRUE`

➤ 釋放 Mutex 後觸發排程 (`OS_Sched`)

- 當你優先級降低時，可能有其他沒有在等待此 `Mutex`，但優先級比原始優先級高的任務（原本被 `CPP` 壓制的任務）現在應該要執行了。因此，必須加入 `OS_Sched()`，確保優先級降低後，系統能立即將 `CPU` 交給更高優先級的就緒任務。

[PART III] Performance Analysis [10%]

- Explain the difference between NPCS and CPP, and how each mechanism avoids the deadlock problem.
 - Describe how it works、Advantage、Disadvantage

➤ **CPP(Ceiling Priority Protocol):**

- ◆ 允許搶佔，但受到嚴格限制。當任務進入 `Critical Section` 時，會立即將優先級提升至該資源的 `Ceiling Priority`。這只會阻止會使用該資源或優先級較低的任務進行搶佔，但仍允許那些不使用該資源且優先級更高的任務搶佔，保留了系統的反應能力
- ◆ 如何避免 `Deadlock`：
 - 由於任務一旦取得資源，優先級就會被提升至最高潛在競爭者的等級，系統保證了當前任務在持有資源期間，不會被任何可能請求同一資源的任務搶佔，確保任務能順利執行完畢並釋放資源，不會卡在中間等待更高優先級的任務釋放資源
- ◆ **Advantage:**
 - 允許優先級高於 `Ceiling Priority`(即不使用該資源) 的任務進行搶佔。這確保了系統能對緊急且不衝突的事件做出即時反應

- 高優先級任務在執行過程中，最多只會被阻塞一次，避免了連鎖阻塞 Transitive blocking 的發生。

◆ Disadvantages :

- 動態調整優先級並操作 Ready Table，程式邏輯較為繁瑣
- 涉及優先級計算與 Ready Queue 的改變，CPU 使用週期高於 NPCS

➤ NPCS(non-preemptible critical section) :

- ◆ 將 Critical Section 視為不可分割的。透過關閉中斷或鎖住排程器，禁止任何 Context Switch。這意味著執行中的任務不能被任何其他任務搶佔，無論對方的優先級多高

◆ 如何避免 Deadlock:

- 因為持有資源的任務獨佔 CPU，其他任務無法執行，無法形成資源依賴的循環，直到當前任務釋放資源為止

◆ Advantage :

- 邏輯簡單，不需要複雜的優先級計算或修改 OS 核心
- 進入與退出僅較少的 CPU 指令，減少系統時間

◆ Disadvantage :

- 它會阻塞所有任務，包含那些完全不需要此資源的最高優先級緊急任務。在 Real-time 系統中，這可能導致危急任務無法及時執行
- 中斷延遲：若實作方式為關閉中斷，系統在 Critical Section 期間將無法響應硬體訊號。