

Embedded OS Implementation, Fall 2025

Project #1 (due at November 3rd, 2025 (Monday) 08:00)

[PART I] Task Control Block Linked List

Objective:

Following the previous homework (HW1), please add some code to the µC/OS-II scheduler in the kernel level to observe the operations of the task control block (TCB) and TCB linked list.

- ※ The TCB address is dynamic.
- ※ This part will be included in the subsequent output and does not require separate code submission.

The example output results are shown below:

```
OSTick      created, Thread ID 11760
Task[ 63] created, TCB Address 92e9e0
-----After TCB[63] begin linked-----
Previous TCB point to adress          0
Current   TCB point to adress        92e9e0
Next      TCB point to adress        0

The file 'TaskSet.txt' was opened
Task[  1] created, TCB Address 92ea54
-----After TCB[ 1] begin linked-----
Previous TCB point to adress          0
Current   TCB point to adress        92ea54
Next      TCB point to adress        92e9e0

Task[  2] created, TCB Address 92eac8
-----After TCB[ 2] begin linked-----
Previous TCB point to adress          0
Current   TCB point to adress        92eac8
Next      TCB point to adress        92ea54

=====TCB linked list=====
Task    Prev_TCB_addr   TCB_addr   Next_TCB_addr
  2            0         92eac8       92ea54
  1           92eac8       92ea54       92e9e0
  63          92ea54       92e9e0         0
```

[PART II] RM Scheduler Implementation

Objective:

To implement the Rate Monotonic (RM) scheduler for periodic tasks and observe the scheduling behaviors.

Problem Definition:

Implement the following three task sets of periodic tasks. Add necessary code to the µC/OS-II scheduler in the kernel level to observe how the task suffers from the scheduler. We give the files for the parameter of the task.

Periodic Task Set = { τ_{ID} (ID, arrival time, execution time, period)}

Example Task Set 1 = { $\tau_1(1, 0, 1, 3), \tau_2(2, 0, 3, 5)$ }

Example Task Set 2 = { $\tau_1(1, 0, 1, 3), \tau_2(2, 1, 1, 4), \tau_3(3, 2, 1, 5)$ }

Example Task Set 3 = { $\tau_1(1, 0, 3, 8), \tau_2(2, 1, 2, 6), \tau_3(3, 0, 4, 15)$ }

※ If tasks have same period, the task with the **smaller TaskID** will be executed first.

※ The priority of the task is set according to the RM scheduling rules.

The input file format:

Task ID	Arrive Time	Execution Time	Task Period
##	##	##	##

Example of task set file 1:

1 0 1 3
2 0 3 5

Evaluation:

The output format:

Tick	Event	CurrentTask ID	NextTask ID	Response Time	Preemption Time	OSTimeDly
##	Preemption	task(ID)(job number)	task(ID)(job number)			
##	Completion	task(ID)(job number)	task(ID)(job number)	##	##	##
##	MissDeadline	task(ID)(job number)	-----			
##			task(ID) is running			

※ If the task is Idle Task, print “**task(priority)**”.

Response Time: the duration between the task's arrival time and the time it is completed.

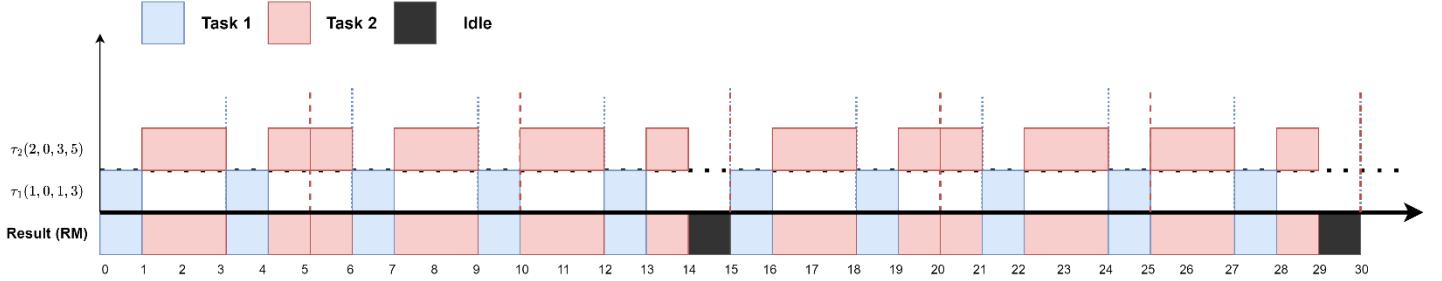
Preemption Time: the time this task is preempted by higher-priority tasks.

OSTimeDly: the remaining delay time for this task

※ You have to print “task(ID) is running” message in **task()** function.

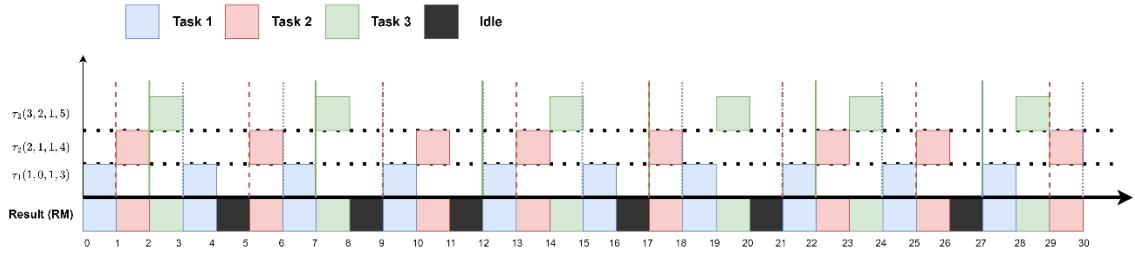
※ Output priority: MissDeadline > Preemption > Completion > task(ID) is running .

The scheduled results of Task Set 1 = { $\tau_1(1, 0, 1, 3)$, $\tau_2(2, 0, 3, 5)$ }:



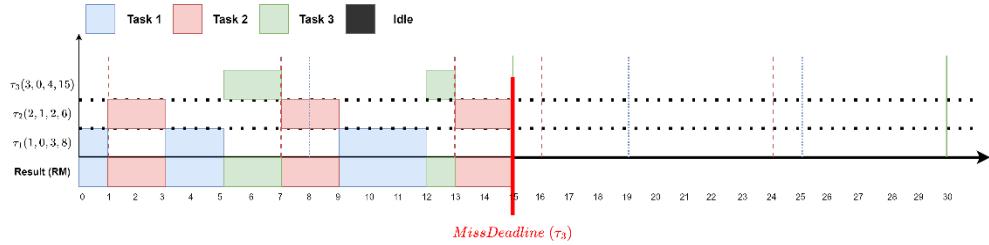
0	task (1) is running						
1	Completion	task(1)(0)	task(2)(0)	1	0	2	
1	task (2) is running						
2	task (2) is running						
3	Preemption	task(2)(0)	task(1)(1)				
3	task (1) is running						
4	Completion	task(1)(1)	task(2)(0)	1	0	2	
4	task (2) is running						
5	Completion	task(2)(0)	task(2)(1)	5	2	0	
5	task (2) is running						
6	Preemption	task(2)(1)	task(1)(2)				
6	task (1) is running						
7	Completion	task(1)(2)	task(2)(1)	1	0	2	
7	task (2) is running						
8	task (2) is running						
9	Completion	task(2)(1)	task(1)(3)	4	1	1	
9	task (1) is running						
10	Completion	task(1)(3)	task(2)(2)	1	0	2	
10	task (2) is running						
11	task (2) is running						
12	Preemption	task(2)(2)	task(1)(4)				
12	task (1) is running						
13	Completion	task(1)(4)	task(2)(2)	1	0	2	
13	task (2) is running						
14	Completion	task(2)(2)	task(63)	4	1	1	
15	Preemption	task(63)	task(1)(5)				
15	task (1) is running						

The output results of Task Set 2 = { $\tau_1(1, 0, 1, 3)$, $\tau_2(2, 1, 1, 4)$, $\tau_3(3, 2, 1, 5)$ }:



0	task (1) is running						
1	Completion task(1)(0)	task(2)(0)	1	0	2		
1	task (2) is running						
2	Completion task(2)(0)	task(3)(0)	1	0	3		
2	task (3) is running						
3	Completion task(3)(0)	task(1)(1)	1	0	4		
3	task (1) is running						
4	Completion task(1)(1)	task(63)	1	0	2		
5	Preemption task(63)	task(2)(1)					
5	task (2) is running						
6	Completion task(2)(1)	task(1)(2)	1	0	3		
6	task (1) is running						
7	Completion task(1)(2)	task(3)(1)	1	0	2		
7	task (3) is running						
8	Completion task(3)(1)	task(63)	1	0	4		
9	Preemption task(63)	task(1)(3)					
9	task (1) is running						
10	Completion task(1)(3)	task(2)(2)	1	0	2		
10	task (2) is running						
11	Completion task(2)(2)	task(63)	2	1	2		
12	Preemption task(63)	task(1)(4)					
12	task (1) is running						
13	Completion task(1)(4)	task(2)(3)	1	0	2		
13	task (2) is running						
14	Completion task(2)(3)	task(3)(2)	1	0	3		
14	task (3) is running						
15	Completion task(3)(2)	task(1)(5)	3	2	2		
15	task (1) is running						

The output results of Task Set 3 = { $\tau_1(1, 0, 3, 8)$, $\tau_2(2, 1, 2, 6)$, $\tau_3(3, 0, 4, 15)$ }:



0	task (1) is running					
1	Preemption	task(1)(0)	task(2)(0)			
1	task (2) is running					
2	task (2) is running					
3	Completion	task(2)(0)	task(1)(0)	2	0	4
3	task (1) is running					
4	task (1) is running					
5	Completion	task(1)(0)	task(3)(0)	5	2	3
5	task (3) is running					
6	task (3) is running					
7	Preemption	task(3)(0)	task(2)(1)			
7	task (2) is running					
8	task (2) is running					
9	Completion	task(2)(1)	task(1)(1)	2	0	4
9	task (1) is running					
10	task (1) is running					
11	task (1) is running					
12	Completion	task(1)(1)	task(3)(0)	4	1	4
12	task (3) is running					
13	Preemption	task(3)(0)	task(2)(2)			
13	task (2) is running					
14	task (2) is running					
15	MissDeadline	task(3)(0)	-----			

[Part III] FIFO Scheduler Implementation

Objective:

To implement the non-preemptive First In First Out (FIFO) scheduling for periodic tasks, and handle the miss deadline behaviors.

Problem Definition:

Implement the following task set of periodic tasks. Add necessary code to the µC/OS-II scheduler **in the kernel level** to observe how the task suffers the schedule delay.

Periodic Task Set = { τ_i (ID, arrival time, execution time, period)}

Task Set 1 = { τ_1 (1, 0, 1, 4), τ_2 (2, 0, 3, 5)}

Task Set 2 = { τ_1 (1, 0, 1, 3), τ_2 (2, 1, 2, 7), τ_3 (3, 4, 3, 12)}

※ If tasks arrive simultaneously, the task with the **smaller TaskID** will be executed first.

Evaluation:

The output format:

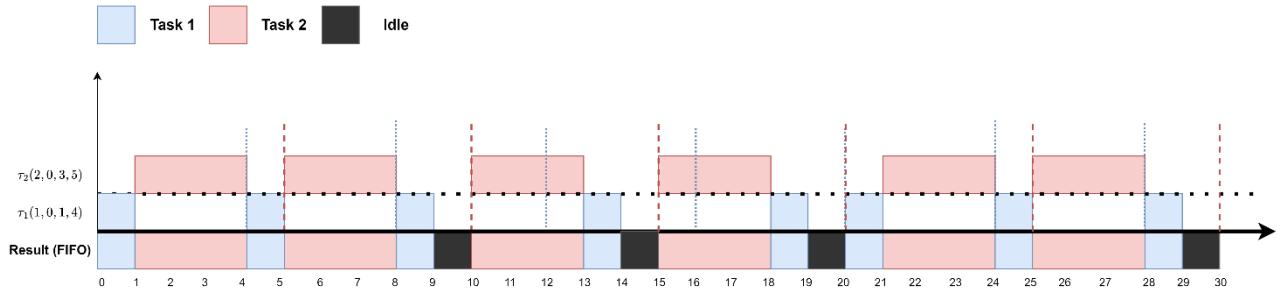
Tick	Event	CurrentTask ID	NextTask ID	Response Time	Preemption Time	OSTimeDly
##	Preemption	task(ID)(job number)	task(ID)(job number)			
##	Completion	task(ID)(job number)	task(ID)(job number)	##	##	##
##	MissDeadline	task(ID)(job number)	-----			
##		task(ID) is running				

※ If the task is Idle Task, print “*task(priority)*”.

※ You have to print “task(ID) is running” message in task() function.

※ Output priority: MissDeadline > Preemption > Completion > task(ID) is running .

The output results of Task Set 1 = { $\tau_1(1, 0, 1, 4)$, $\tau_2(2, 0, 3, 5)$ }:

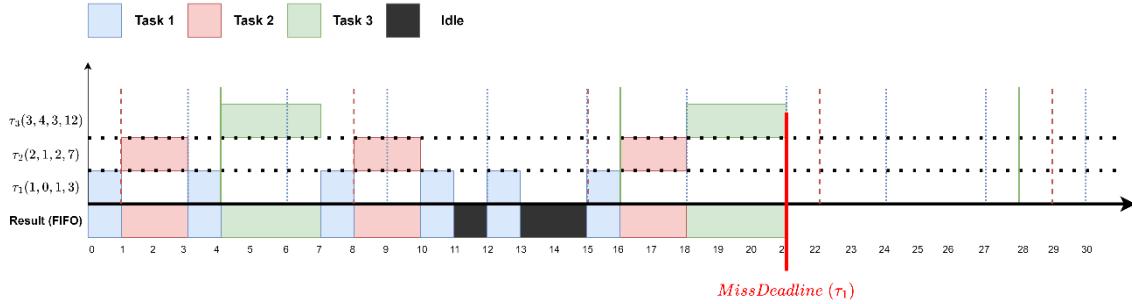


```

0  task( 1) is running
1  Completion   task( 1)( 0)    task( 2)( 0)      1          0          3
1  task( 2) is running
2  task( 2) is running
3  task( 2) is running
4  Completion   task( 2)( 0)    task( 1)( 1)      4          1          1
4  task( 1) is running
5  Completion   task( 1)( 1)    task( 2)( 1)      1          0          3
5  task( 2) is running
6  task( 2) is running
7  task( 2) is running
8  Completion   task( 2)( 1)    task( 1)( 2)      3          0          2
8  task( 1) is running
9  Completion   task( 1)( 2)    task( 63)         1          0          3
10 Preemption   task( 63)       task( 2)( 2)
10 task( 2) is running
11 task( 2) is running
12 task( 2) is running
13 Completion   task( 2)( 2)    task( 1)( 3)      3          0          2
13 task( 1) is running
14 Completion   task( 1)( 3)    task( 63)         2          1          2
15 Preemption   task( 63)       task( 2)( 3)
15 task( 2) is running

```

The output results of Task Set 2 = { $\tau_1(1, 0, 1, 3)$, $\tau_2(2, 1, 2, 7)$, $\tau_3(3, 4, 3, 12)$ } :



0	task (1) is running						
1	Completion task(1)(0)	task(2)(0)		1	0	2	
1	task (2) is running						
2	task (2) is running						
3	Completion task(2)(0)	task(1)(1)		2	0	5	
3	task (1) is running						
4	Completion task(1)(1)	task(3)(0)		1	0	2	
4	task (3) is running						
5	task (3) is running						
6	task (3) is running						
7	Completion task(3)(0)	task(1)(2)		3	0	9	
7	task (1) is running						
8	Completion task(1)(2)	task(2)(1)		2	1	1	
8	task (2) is running						
9	task (2) is running						
10	Completion task(2)(1)	task(1)(3)		2	0	5	
10	task (1) is running						
11	Completion task(1)(3)	task(63)		2	1	1	
12	Preemption task(63)	task(1)(4)					
12	task (1) is running						
13	Completion task(1)(4)	task(63)		1	0	2	
15	Preemption task(63)	task(1)(5)					
15	task (1) is running						
16	Completion task(1)(5)	task(2)(2)		1	0	2	
16	task (2) is running						
17	task (2) is running						
18	Completion task(2)(2)	task(3)(1)		3	1	4	
18	task (3) is running						
19	task (3) is running						
20	task (3) is running						
21	MissDeadline task(1)(6)	-----					

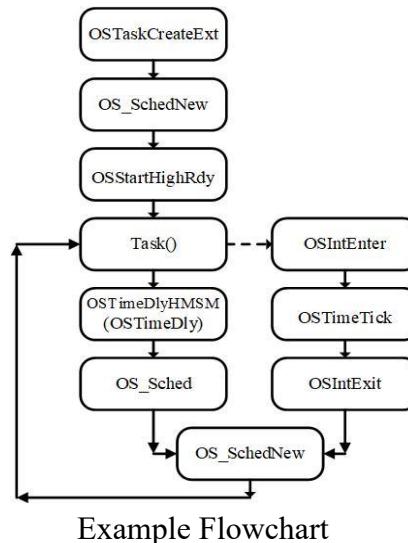
Credit:

[PART I] Task Control Block Linked List [20%]

- The screenshot results. (10%)
- A report that describes your implementation (please attach the screenshot of the code and **MARK** the modified part). (10%)

[PART II] RM Scheduler Implementation [70%]

- The correctness of schedule results of examples (**Output.txt**). Note the testing task set might not be the same as the given example task set. (25%)
- A report that describes your implementation (please attach the screenshot of the code and **MARK** the modified part). (40%)
 - Indicate the modified parts in the given example flowchart and provide reasons for their placement within the process.
- Implement and describe how to handle the deadline missing situation under RM. (5%)



[PART III] FIFO Scheduler Implementation [10%]

- The correctness of schedule results of examples (**Output.txt**). Note the testing task set might not be the same as the given example task set. (5%)
- Implement FIFO and compare the schedule results with that of RM (please attach the screenshot of the code and **MARK** the modified part). (5%)

- ※ You must modify the source code!
- ※ Standard input and output filenames in the project are necessary for the checker. Please check the file names before submitting.

```
#define INPUT_FILE_NAME "./TaskSet.txt"  
#define OUTPUT_FILE_NAME "./Output.txt"
```

- ※ Please set the system end time as 30 seconds in this project.

```
#define SYSTEM_END_TIME 30
```

- ※ You must define the algorithm in ucos_ii.h.

```
#define RM 0  
#define FIFO 1  
#define ALGORITHM RM/FIFO
```

- ※ We will use different task sets to verify your code.

- ※ We will check your answer on Output.txt.

- ※ When the current task is completed, the completion information shall be printed even if there is one task missing its deadline.

Project submit:

Submit to Moodle.

Submit deadline: **November 3rd, 2025 (Moonday) 08:00**

File name format: RTOS_Myyyddxxx_PA1.zip

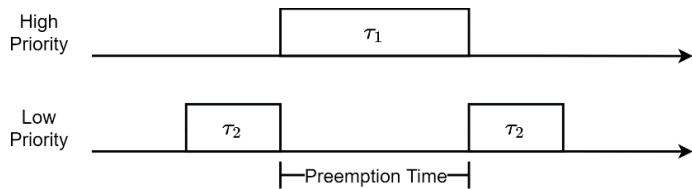
RTOS_Myyyddxxx_PA1.zip includes (The tree structure of files is shown as hints):

- The report (RTOS_Myyyddxxx_PA1.pdf).
- Folder with the executable project (RTOS_Myyyddxxx_PA1).

※ Plagiarizing is strictly prohibited.

Hints:

1. Preemption time is introduced in multiple tasking.



2. Advanced print function

```
1 #define LOG_CONSOLE 1
2 #define LOG_FILE    2
3 #define LOG_BOTH   3
4
5 void LOG_print( INT8U LOG_type, char *outfile,
6                 _In_z_ _Printf_format_string_ char const* const _Format,
7                 ...){
8     va_list _ArgList;
9     __crt_va_start(_ArgList, _Format);
10
11    if((LOG_type & LOG_CONSOLE))
12        vprintf(_Format, _ArgList);
13
14    if((LOG_type & LOG_FILE) && ((Output_err = fopen_s(&Output_fp, outfile, "a")) == 0)){
15        vfprintf(Output_fp, _Format, _ArgList);
16        fclose(Output_fp);
17    }
18    return;
19 }
```

3. RTOS_Myyddxx_PA1.zip include files as follows:

```
C:.  
  RTOS_Myyddxxx_PA1.pdf  
  RTOS_Myyddxxx_PA1  
    ReadMe.txt  
    Micrium  
      Software  
        uC-CPU  
          cpu_cache.h  
          cpu_core.c  
          cpu_core.h  
          cpu_def.h  
        Win32  
          Visual_Studio  
            cpu.h  
            cpu_c.c  
        uC-LIB  
          lib_ascii.c  
          lib_ascii.h  
          lib_def.h  
          lib_math.c  
          lib_math.h  
          lib_mem.c  
          lib_mem.h  
          lib_str.c  
          lib_str.h  
        uCOS-II  
          Ports  
            Win32  
              Visual Studio  
                os_cpu.h  
                os_cpu_c.c  
                os_cpu_c.c.bak  
          Source  
            os.h  
            os_cfg_r.h  
            os_core.c  
            os_core.c.bak  
            os_dbg_r.c  
            os_flag.c  
            os_mbox.c  
            os_mem.c  
            os_mutex.c  
            os_q.c  
            os_sem.c  
            os_task.c
```

```
  os_task.c  
  os_time.c  
  os_tmr.c  
  os_trace.h  
  ucos_ii.c  
  ucos_ii.h  
  ucos_ii.h.bak  
  Microsoft  
    BSP  
      Windows  
        bsp_cpu.c  
    Windows  
      Kernel  
        app_cfg.h  
        cpu_cfg.h  
        lib_cfg.h  
    OS2  
      app_hooks.c  
      app_hooks.c.bak  
      main.c  
      main.c.bak  
      os_cfg.h  
      os_cfg.h.bak  
    VS  
      OS2.sln  
      OS2.vcxproj  
      OS2.vcxproj.filters  
      OS2.vcxproj.user  
      Output.txt  
      TaskSet.txt
```