

Embedded OS Implementation, Fall 2025
Project #1 (due at November 3rd, 2025 (Monday) 08:00)

[PART I] Task Control Block Linked List

Objective:

Following the previous homework (HW1), please add some code to the μ C/OS-II scheduler in the kernel level to observe the operations of the task control block (TCB) and TCB linked list.

- ※ The TCB address is dynamic.
- ※ This part will be included in the subsequent output and does not require separate code submission.

The example output results are shown below:

```
OSTick    created, Thread ID 11760
Task[ 63] created, TCB Address  92e9e0
-----After TCB[63] begin linked-----
Previous TCB point to adress      0
Current  TCB point to adress      92e9e0
Next     TCB point to adress      0

The file 'TaskSet.txt' was opened
Task[  1] created, TCB Address  92ea54
-----After TCB[ 1] begin linked-----
Previous TCB point to adress      0
Current  TCB point to adress      92ea54
Next     TCB point to adress      92e9e0

Task[  2] created, TCB Address  92eac8
-----After TCB[ 2] begin linked-----
Previous TCB point to adress      0
Current  TCB point to adress      92eac8
Next     TCB point to adress      92ea54

=====TCB linked list=====
Task    Prev_TCB_addr    TCB_addr    Next_TCB_addr
  2             0         92eac8         92ea54
  1        92eac8         92ea54         92e9e0
 63        92ea54         92e9e0             0
```

[PART II] RM Scheduler Implementation

Objective:

To implement the Rate Monotonic (RM) scheduler for periodic tasks and observe the scheduling behaviors.

Problem Definition:

Implement the following three task sets of periodic tasks. Add necessary code to the $\mu\text{C}/\text{OS-II}$ scheduler **in the kernel level** to observe how the task suffers from the scheduler. We give the files for the parameter of the task.

Periodic Task Set = $\{\tau_{\text{ID}}(\text{ID}, \text{arrival time}, \text{execution time}, \text{period})\}$

Example Task Set 1 = $\{\tau_1(1, 0, 1, 3), \tau_2(2, 0, 3, 5)\}$

Example Task Set 2 = $\{\tau_1(1, 0, 1, 3), \tau_2(2, 1, 1, 4), \tau_3(3, 2, 1, 5)\}$

Example Task Set 3 = $\{\tau_1(1, 0, 3, 8), \tau_2(2, 1, 2, 6), \tau_3(3, 0, 4, 15)\}$

※ If tasks have same period, the task with the **smaller TaskID** will be executed first.

※ The priority of the task is set according to the RM scheduling rules.

The input file format:

Task ID	Arrive Time	Execution Time	Task Period
##	##	##	##

Example of task set file 1:

```
1 0 1 3
2 0 3 5
```

Evaluation:

The output format:

Tick	Event	CurrentTask ID	NextTask ID	Response Time	Preemption Time	OSTimeDly
##	Preemption	task(ID)(job number)	task(ID)(job number)			
##	Completion	task(ID)(job number)	task(ID)(job number)	##	##	##
##	MissDeadline	task(ID)(job number)	-----			
##	task(ID) is running					

※ If the task is Idle Task, print “*task(priority)*”.

Response Time: the duration between the task's arrival time and the time it is completed.

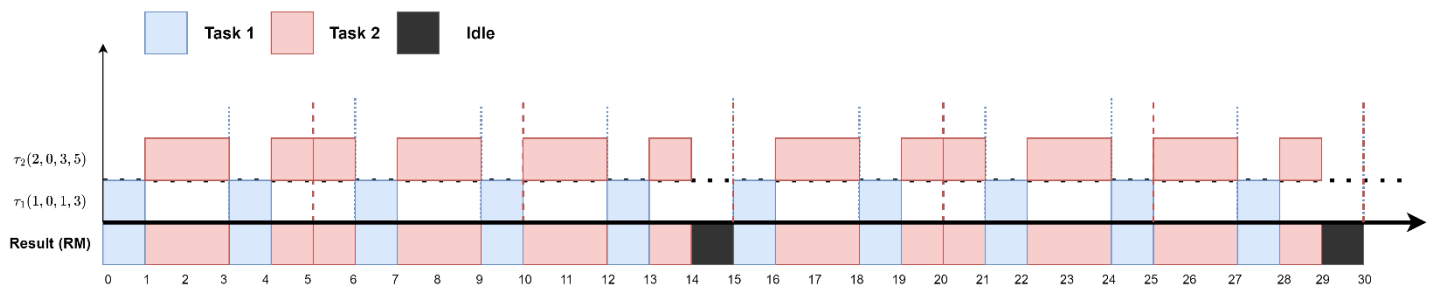
Preemption Time: the time this task is preempted by higher-priority tasks.

OSTimeDly: the remaining delay time for this task

※ **You have to print “task(ID) is running” message in task() function.**

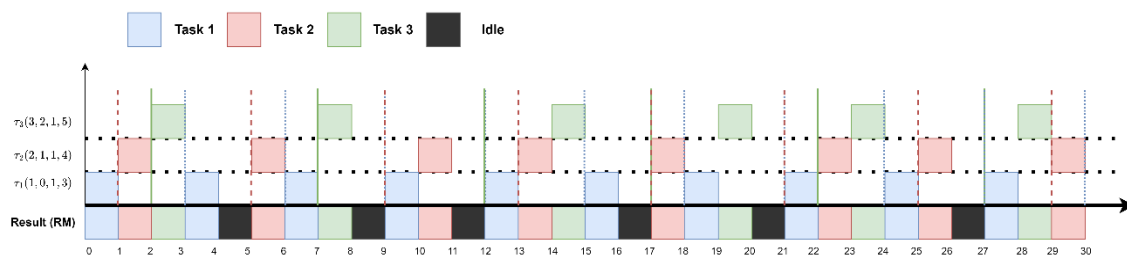
※ **Output priority: MissDeadline > Preemption > Completion > task(ID) is running .**

The scheduled results of **Task Set 1** = $\{\tau_1(1, 0, 1, 3), \tau_2(2, 0, 3, 5)\}$:



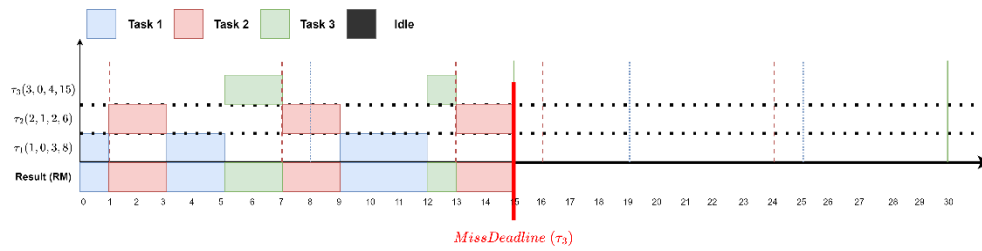
0	task (1) is running				
1	Completion task(1)(0)	task(2)(0)	1	0	2
1	task (2) is running				
2	task (2) is running				
3	Preemption task(2)(0)	task(1)(1)			
3	task (1) is running				
4	Completion task(1)(1)	task(2)(0)	1	0	2
4	task (2) is running				
5	Completion task(2)(0)	task(2)(1)	5	2	0
5	task (2) is running				
6	Preemption task(2)(1)	task(1)(2)			
6	task (1) is running				
7	Completion task(1)(2)	task(2)(1)	1	0	2
7	task (2) is running				
8	task (2) is running				
9	Completion task(2)(1)	task(1)(3)	4	1	1
9	task (1) is running				
10	Completion task(1)(3)	task(2)(2)	1	0	2
10	task (2) is running				
11	task (2) is running				
12	Preemption task(2)(2)	task(1)(4)			
12	task (1) is running				
13	Completion task(1)(4)	task(2)(2)	1	0	2
13	task (2) is running				
14	Completion task(2)(2)	task(63)	4	1	1
15	Preemption task(63)	task(1)(5)			
15	task (1) is running				

The output results of **Task Set 2** = $\{\tau_1(1, 0, 1, 3), \tau_2(2, 1, 1, 4), \tau_3(3, 2, 1, 5)\}$:



0	task (1) is running					
1	Completion task(1)(0)	task(2)(0)	1	0		2
1	task (2) is running					
2	Completion task(2)(0)	task(3)(0)	1	0		3
2	task (3) is running					
3	Completion task(3)(0)	task(1)(1)	1	0		4
3	task (1) is running					
4	Completion task(1)(1)	task(63)	1	0		2
5	Preemption task(63)	task(2)(1)				
5	task (2) is running					
6	Completion task(2)(1)	task(1)(2)	1	0		3
6	task (1) is running					
7	Completion task(1)(2)	task(3)(1)	1	0		2
7	task (3) is running					
8	Completion task(3)(1)	task(63)	1	0		4
9	Preemption task(63)	task(1)(3)				
9	task (1) is running					
10	Completion task(1)(3)	task(2)(2)	1	0		2
10	task (2) is running					
11	Completion task(2)(2)	task(63)	2	1		2
12	Preemption task(63)	task(1)(4)				
12	task (1) is running					
13	Completion task(1)(4)	task(2)(3)	1	0		2
13	task (2) is running					
14	Completion task(2)(3)	task(3)(2)	1	0		3
14	task (3) is running					
15	Completion task(3)(2)	task(1)(5)	3	2		2
15	task (1) is running					

The output results of **Task Set 3** = $\{\tau_1(1, 0, 3, 8), \tau_2(2, 1, 2, 6), \tau_3(3, 0, 4, 15)\}$:



```

0 task ( 1) is running
1 Preemption task( 1)( 0) task( 2)( 0)
1 task ( 2) is running
2 task ( 2) is running
3 Completion task( 2)( 0) task( 1)( 0) 2 0 4
3 task ( 1) is running
4 task ( 1) is running
5 Completion task( 1)( 0) task( 3)( 0) 5 2 3
5 task ( 3) is running
6 task ( 3) is running
7 Preemption task( 3)( 0) task( 2)( 1)
7 task ( 2) is running
8 task ( 2) is running
9 Completion task( 2)( 1) task( 1)( 1) 2 0 4
9 task ( 1) is running
10 task ( 1) is running
11 task ( 1) is running
12 Completion task( 1)( 1) task( 3)( 0) 4 1 4
12 task ( 3) is running
13 Preemption task( 3)( 0) task( 2)( 2)
13 task ( 2) is running
14 task ( 2) is running
15 MissDeadline task( 3)( 0) -----

```

[Part III] FIFO Scheduler Implementation

Objective:

To implement the non-preemptive First In First Out (FIFO) scheduling for periodic tasks, and handle the miss deadline behaviors.

Problem Definition:

Implement the following task set of periodic tasks. Add necessary code to the μ C/OS-II scheduler **in the kernel level** to observe how the task suffers the schedule delay.

Periodic Task Set = $\{\tau_{ID} (ID, \text{arrival time}, \text{execution time}, \text{period})\}$

Task Set 1 = $\{\tau_1 (1, 0, 1, 4), \tau_2 (2, 0, 3, 5)\}$

Task Set 2 = $\{\tau_1 (1, 0, 1, 3), \tau_2 (2, 1, 2, 7), \tau_3 (3, 4, 3, 12)\}$

※If tasks arrive simultaneously, the task with the **smaller TaskID** will be executed first.

Evaluation:

The output format:

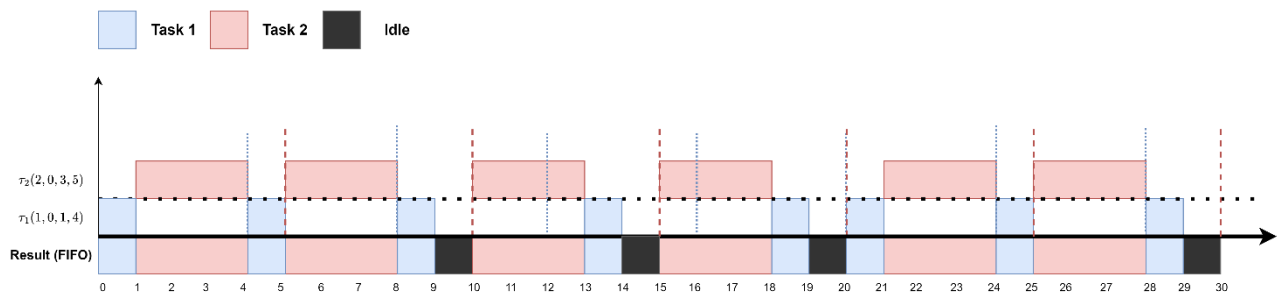
Tick	Event	CurrentTask ID	NextTask ID	Response Time	Preemption Time	OSTimeDly
##	Preemption	task(ID)(job number)	task(ID)(job number)			
##	Completion	task(ID)(job number)	task(ID)(job number)	##	##	##
##	MissDeadline	task(ID)(job number)	-----			
##	task(ID) is running					

※ If the task is Idle Task, print “*task(priority)*”.

※ **You have to print “task(ID) is running” message in task() function.**

※ **Output priority: MissDeadline > Preemption > Completion > task(ID) is running .**

The output results of **Task Set 1** = $\{\tau_1(1, 0, 1, 4), \tau_2(2, 0, 3, 5)\}$:

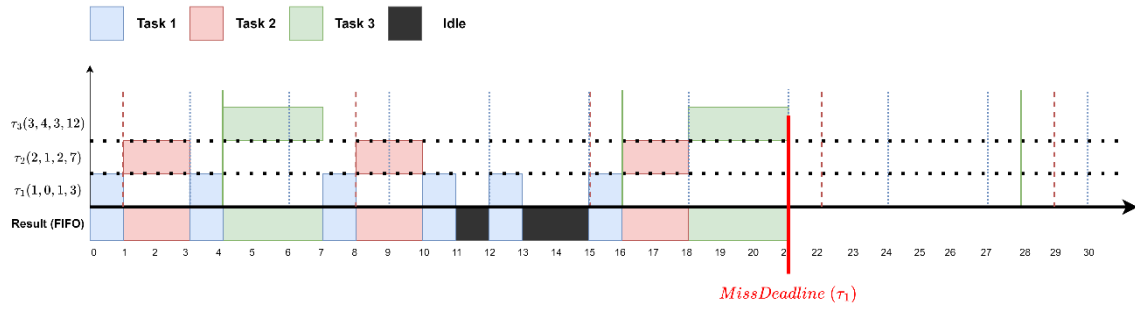


```

0 task ( 1) is running
1 Completion task( 1)( 0) task( 2)( 0) 1 0 3
1 task ( 2) is running
2 task ( 2) is running
3 task ( 2) is running
4 Completion task( 2)( 0) task( 1)( 1) 4 1 1
4 task ( 1) is running
5 Completion task( 1)( 1) task( 2)( 1) 1 0 3
5 task ( 2) is running
6 task ( 2) is running
7 task ( 2) is running
8 Completion task( 2)( 1) task( 1)( 2) 3 0 2
8 task ( 1) is running
9 Completion task( 1)( 2) task( 63) 1 0 3
10 Preemption task( 63) task( 2)( 2)
10 task ( 2) is running
11 task ( 2) is running
12 task ( 2) is running
13 Completion task( 2)( 2) task( 1)( 3) 3 0 2
13 task ( 1) is running
14 Completion task( 1)( 3) task( 63) 2 1 2
15 Preemption task( 63) task( 2)( 3)
15 task ( 2) is running

```

The output results of **Task Set 2** = $\{\tau_1(1, 0, 1, 3), \tau_2(2, 1, 2, 7), \tau_3(3, 4, 3, 12)\}$:



0	task (1) is running				
1	Completion task(1)(0)	task(2)(0)	1	0	2
1	task (2) is running				
2	task (2) is running				
3	Completion task(2)(0)	task(1)(1)	2	0	5
3	task (1) is running				
4	Completion task(1)(1)	task(3)(0)	1	0	2
4	task (3) is running				
5	task (3) is running				
6	task (3) is running				
7	Completion task(3)(0)	task(1)(2)	3	0	9
7	task (1) is running				
8	Completion task(1)(2)	task(2)(1)	2	1	1
8	task (2) is running				
9	task (2) is running				
10	Completion task(2)(1)	task(1)(3)	2	0	5
10	task (1) is running				
11	Completion task(1)(3)	task(63)	2	1	1
12	Preemption task(63)	task(1)(4)			
12	task (1) is running				
13	Completion task(1)(4)	task(63)	1	0	2
15	Preemption task(63)	task(1)(5)			
15	task (1) is running				
16	Completion task(1)(5)	task(2)(2)	1	0	2
16	task (2) is running				
17	task (2) is running				
18	Completion task(2)(2)	task(3)(1)	3	1	4
18	task (3) is running				
19	task (3) is running				
20	task (3) is running				
21	MissDeadline task(1)(6)	-----			

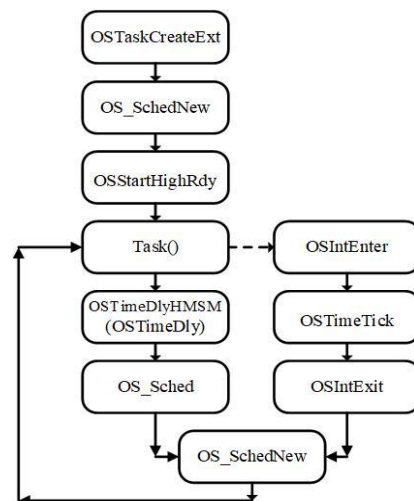
Credit:

[PART I] Task Control Block Linked List [20%]

- The screenshot results. (10%)
- A report that describes your implementation (please attach the screenshot of the code and **MARK** the modified part). (10%)

[PART II] RM Scheduler Implementation [70%]

- The correctness of schedule results of examples (**Output.txt**). Note the testing task set might not be the same as the given example task set. (25%)
- A report that describes your implementation (please attach the screenshot of the code and **MARK** the modified part). (40%)
 - Indicate the modified parts in the given example flowchart and provide reasons for their placement within the process.
- Implement and describe how to handle the deadline missing situation under RM. (5%)



Example Flowchart

[PART III] FIFO Scheduler Implementation [10%]

- The correctness of schedule results of examples (**Output.txt**). Note the testing task set might not be the same as the given example task set. (5%)
- Implement FIFO and compare the schedule results with that of RM (please attach the screenshot of the code and **MARK** the modified part). (5%)

※ You must modify the source code!

※ Standard input and output filenames in the project are necessary for the checker. Please check the file names before submitting.

```
#define INPUT_FILE_NAME "./TaskSet.txt"
```

```
#define OUTPUT_FILE_NAME "./Output.txt"
```

※ Please set the system end time as 30 seconds in this project.

```
#define SYSTEM_END_TIME 30
```

※ You must define the algorithm in `ucos_ii.h`.

```
#define RM 0
```

```
#define FIFO 1
```

```
#define ALGORITHM RM/FIFO
```

※ We will use **different task sets** to verify your code.

※ We will check your answer on `Output.txt`.

※ When the current task is completed, the completion information shall be printed even if there is one task missing its deadline.

Project submit:

Submit to Moodle.

Submit deadline: **November 3rd, 2025 (Moonday) 08:00**

File name format: RTOS_Myyyddxxx_PA1.zip

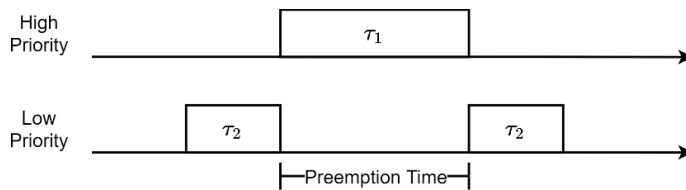
RTOS_Myyyddxxx_PA1.zip includes (The tree structure of files is shown as hints):

- The report (RTOS_Myyyddxxx_PA1.pdf).
- Folder with the executable project (RTOS_Myyyddxxx_PA1).

※ Plagiarizing is strictly prohibited.

Hints:

1. Preemption time is introduced in multiple tasking.



2. Advanced print function

```
1  #define LOG_CONSOLE 1
2  #define LOG_FILE    2
3  #define LOG_BOTH    3
4
5  void LOG_print( INT8U LOG_type, char *outfile,
6                 _In_z_ _Printf_format_string_ char const* const _Format,
7                 ...){
8      va_list _ArgList;
9      __crt_va_start(_ArgList, _Format);
10
11     if((LOG_type & LOG_CONSOLE))
12         vprintf(_Format, _ArgList);
13
14     if((LOG_type & LOG_FILE) && ((Output_err = fopen_s(&Output_fp, outfile, "a")) == 0)){
15         vfprintf(Output_fp, _Format, _ArgList);
16         fclose(Output_fp);
17     }
18     return;
19 }
```

3. RTOS_Myyyddxxx_PA1.zip include files as follows:

```

C:.\
├──RTOS_ Myyyddxxx_PA1.pdf
├──RTOS_Myyyddxxx_PA1
│   └──ReadMe.txt
├──Micrium
│   └──Software
│       ├──uC-CPU
│       │   ├──cpu_cache.h
│       │   ├──cpu_core.c
│       │   ├──cpu_core.h
│       │   └──cpu_def.h
│       ├──Win32
│       │   └──Visual_Studio
│       │       ├──cpu.h
│       │       └──cpu_c.c
│       ├──uC-LIB
│       │   ├──lib_ascii.c
│       │   ├──lib_ascii.h
│       │   ├──lib_def.h
│       │   ├──lib_math.c
│       │   ├──lib_math.h
│       │   ├──lib_mem.c
│       │   ├──lib_mem.h
│       │   ├──lib_str.c
│       │   └──lib_str.h
│       └──uCOS-II
│           ├──Ports
│           │   └──Win32
│           │       └──Visual Studio
│           │           ├──os_cpu.h
│           │           ├──os_cpu_c.c
│           │           └──os_cpu_c.c.bak
│           └──Source
│               ├──os.h
│               ├──os_cfg_r.h
│               ├──os_core.c
│               ├──os_core.c.bak
│               ├──os_dbg_r.c
│               ├──os_flag.c
│               ├──os_mbox.c
│               ├──os_mem.c
│               ├──os_mutex.c
│               ├──os_q.c
│               ├──os_sem.c
│               └──os_task.c

```

```

├──os_task.c
├──os_time.c
├──os_tmr.c
├──os_trace.h
├──ucos_ii.c
├──ucos_ii.h
└──ucos_ii.h.bak

├──Microsoft
│   ├──BSP
│   │   └──Windows
│   │       └──bsp_cpu.c
│   ├──Windows
│   │   └──Kernel
│   │       ├──app_cfg.h
│   │       ├──cpu_cfg.h
│   │       └──lib_cfg.h
│   └──OS2
│       ├──app_hooks.c
│       ├──app_hooks.c.bak
│       ├──main.c
│       ├──main.c.bak
│       ├──os_cfg.h
│       └──os_cfg.h.bak
└──VS
    ├──OS2.sln
    ├──OS2.vcxproj
    ├──OS2.vcxproj.filters
    ├──OS2.vcxproj.user
    ├──Output.txt
    └──TaskSet.txt

```