

# CS1 Lecture 15

Feb. 20, 2017

- HW4 due Wed. 2/22, 9 am
  - comment on Q1 list “format”
  - Q2: find \*any\* solution. Don’t try to find the best/optimal solution or all of them
  - Q3: although you cannot use a loop in Q2, you may use a loop in Q3 (but function must also be recursive)
- Exam 1, Thursday evening, 2/23, 6:30-8:00pm, MacBride Auditorium
  - You must bring ID
  - For people with conflicts, email will be sent right after class
- No points in tomorrow’s discussion sections. Not required. Consider them just extra office hours to help with HW4. If you are in a Wed. section, you can go to any Tues. section for help (or regular office hours)
- Wednesday: exam review

## Last time

- comments/advice on use of functions
  - functions that return values, functions that have side effects or print
- Started recursion (end of Ch 5)

## Today

- A few words on Exam 1
- More recursion examples

# Recursion (end of Ch 5)

- Very important and useful concept
- Not just for programming, but math and everyday life, nature, etc.
- Has undeserved reputation among some people: “recursion is bad – recursive programs are inefficient” Yes, one can write very bad recursive programs but this is true of non-recursive programs as well. And recursion *can* be super useful.

# Recursion

- Recursive function: function that contains within its definition calls to itself
- Consider math's factorial. E.g.  $3! = 3 * 2 * 1$
- You might be used to definition like:  $n! = n * (n-1) * ... * 2 * 1$
- But more precise mathematical definition of factorial function is:  
factorial(1) = 1  
factorial(n) =  $n * \text{factorial}(n-1)$ , for all  $n > 1$
- Programming-wise, can very directly translate recursive mathematical definitions into code:

```
def factorial(n):  
    if (n == 1):  
        return 1  
    else:  
        return n * factorial(n - 1)
```

- DON'T let the function call, factorial(n-1), scare you. It's just a function call. If you draw stack frames like we did in earlier lectures, it all works out fine.
- DO need to think carefully when writing/analyzing recursive programs though ...

# Important rules for recursive functions

- When writing a recursive function:
  - MUST have base case(s), situations when code *does not* make recursive call.
  - MUST ensure that recursive calls *make progress toward base cases*. I.e. you need to convince yourself that recursive call is “closer to” base case than the original problem you are working on
  - SHOULD ensure you don’t unnecessarily repeat work. Ignoring this contributes to recursion’s bad reputation. E.g. direct recursive implementation of Fibonacci is extremely and unnecessarily inefficient

# Ch3: Stack frames

```
def countdown(n):  
    if n == 0:  
        print("Blast off!")  
    else:  
        print(n)  
        countdown(n-1)
```

```
>>> n = 100
```

```
>>> y = 2
```

```
>>> countdown(y)
```

```
2
```

```
1
```

```
Blast Off!
```

```
countdown:    n 0
```

```
countdown:    n 1
```

```
countdown:    n 2
```

```
_main_:      n 100  
             y 2
```

# Recursion examples

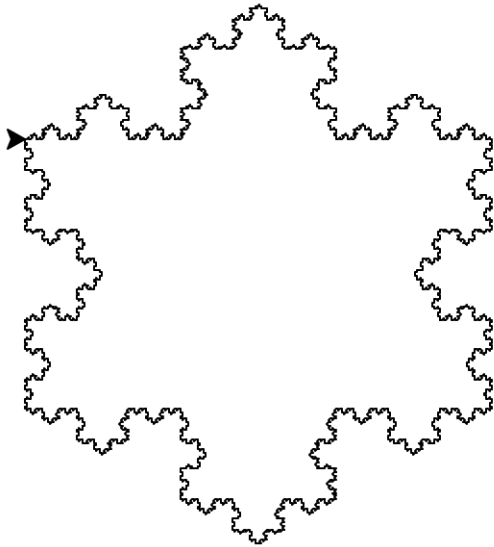
- printList, printListReverse
- sumListItems
- isPalindrome
  - “wasitacaroracatisaw” is a palindrome
  - “sitonapotatoppanotis” is a palindrome
- Fibonacci sequence: 1, 1, 2, 3, 5, 8, ...
  - $\text{fib}(0) = 1, \text{fib}(1) = 1$
  - $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2), \text{ for } n > 1$
- “flatten” a list
  - E.g. [ [[[3,[2,4]]], 0], ['a']], 23 ] -> [3,2,4,0,'a', 23]
- generate a number sequence
- Towers of Hanoi problem

# Towers of Hanoi solution

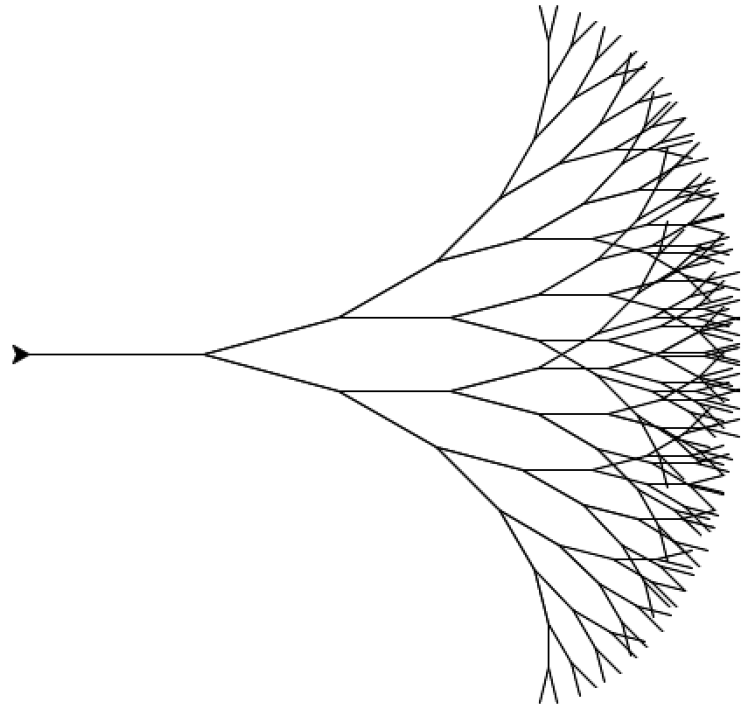
- Rods/towers A, B, C. Goal is to move disks from A to C, one at a time, never allowing larger disk to be on top of smaller disk
- Algorithm:
  - Move  $n-1$  disks from A  $\rightarrow$  B (by this algorithm!)
  - Move final disk from A  $\rightarrow$  C
  - Move  $n-1$  disks from B  $\rightarrow$  C (by this algorithm!)



# Recursion Examples



koch.py



pic.py

## Next Time

- Review for exam, and sample problems