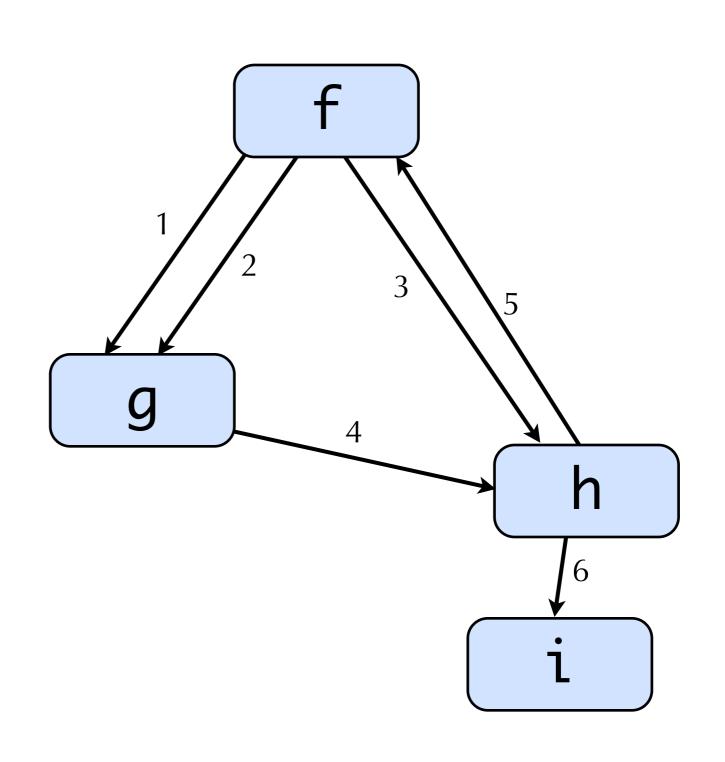


Interprocedural Analysis

CS252r Spring 2011

Call graph example

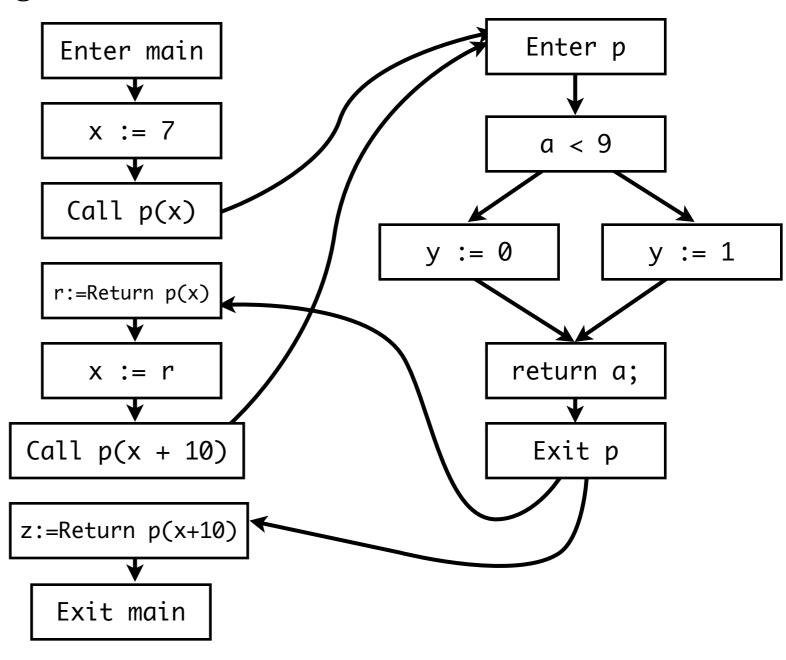
```
f() {
 1: g();
  2:
     g();
   3: h();
g() {
  4: h();
h() {
 5: f();
   6: i();
i() { ... }
```



Interprocedural dataflow analysis

- How do we deal with procedure calls?
- Obvious idea: make one big CFG

```
main() {
  x := 7;
  r := p(x);
  x := r;
  z := p(x + 10);
p(int a) {
  if (a < 9)
    y := 0;
  else
    y := 1;
  return a;
```



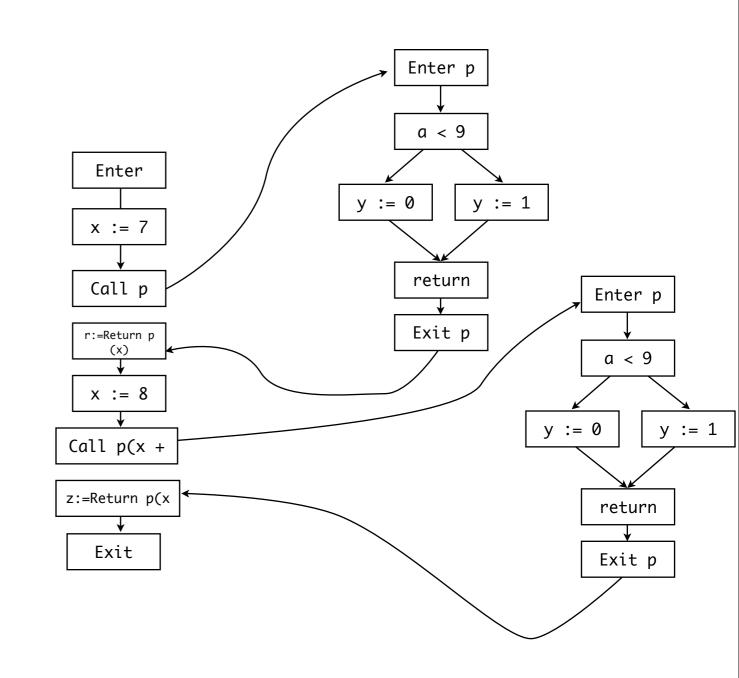
Interprocedural CFG

- CFG may have additional nodes to handle call and returns
 - Treat arguments, return values as assignments
- Note: a local program variable represents multiple locations

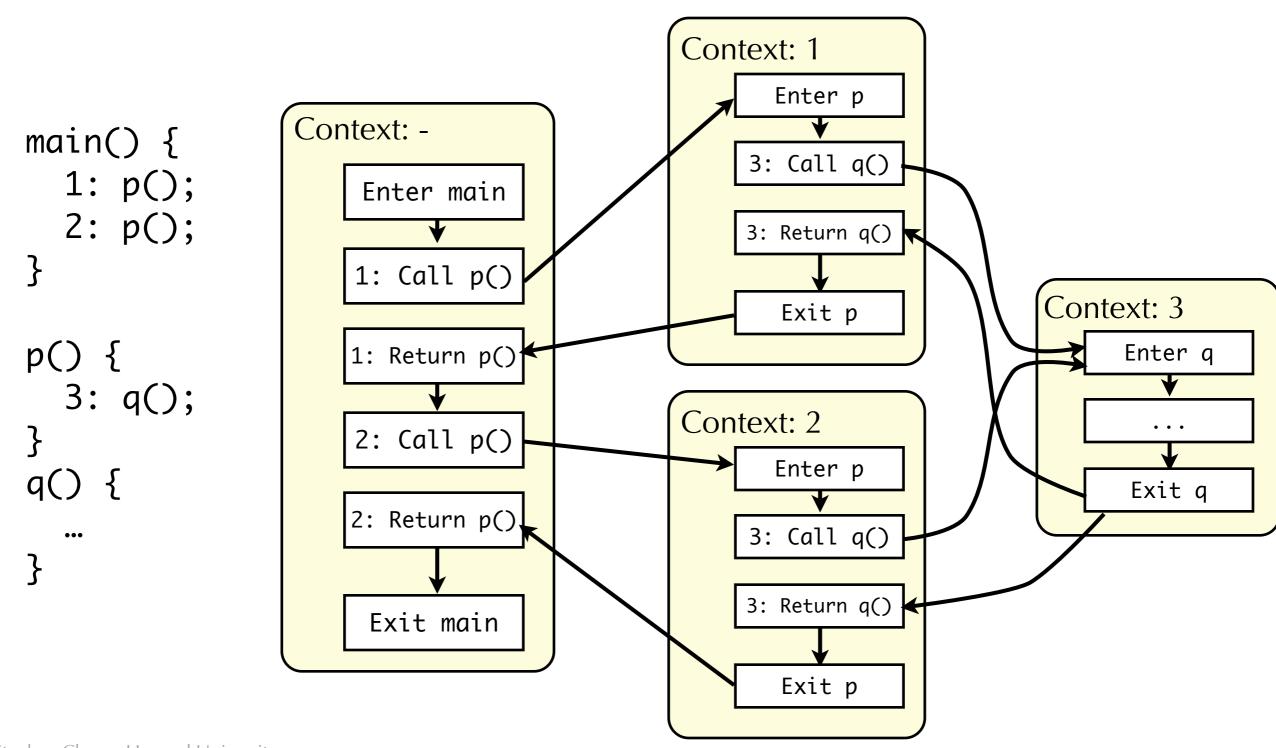
Set up environment for calling p a := x, ...Enter p Enter main x := 7 a < 9Call p(x)y := 0 y := 1 r:=Return p(x) x := rreturn a; Call p(x + 10)Exit p z:=Return p(x+10)Restore calling environment Exit main z := a

Inlining

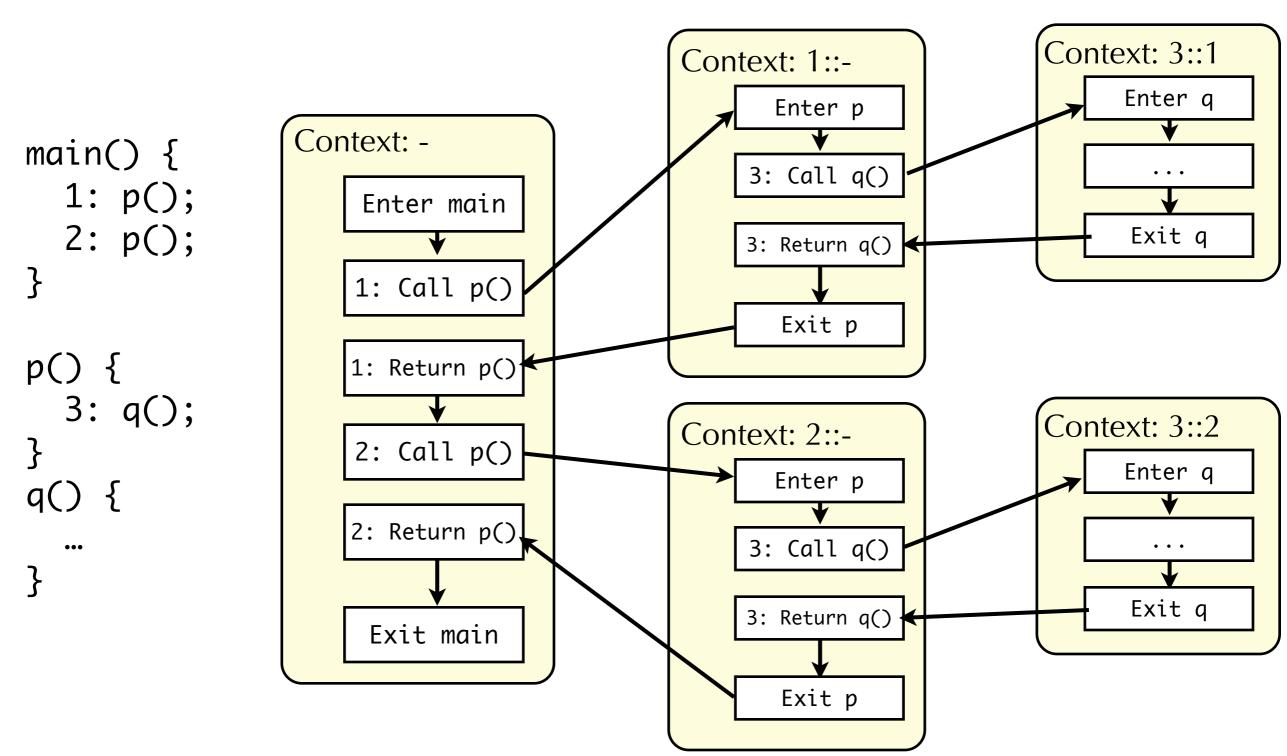
- Inlining
 - Use a new copy of a procedure's CFG at each call site
- Problems? Concerns?
 - May be expensive! Exponential increase in size of CFG
 - p() { q(); q(); } q() { r(); r() }
 r() { ... }
 - What about recursive procedures?
 - $p(int n) \{ ... p(n-1); ... \}$
 - More generally, cycles in the call graph



Context sensitivity example

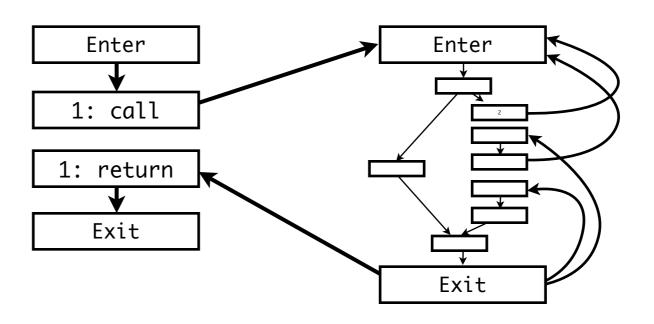


Context sensitivity example



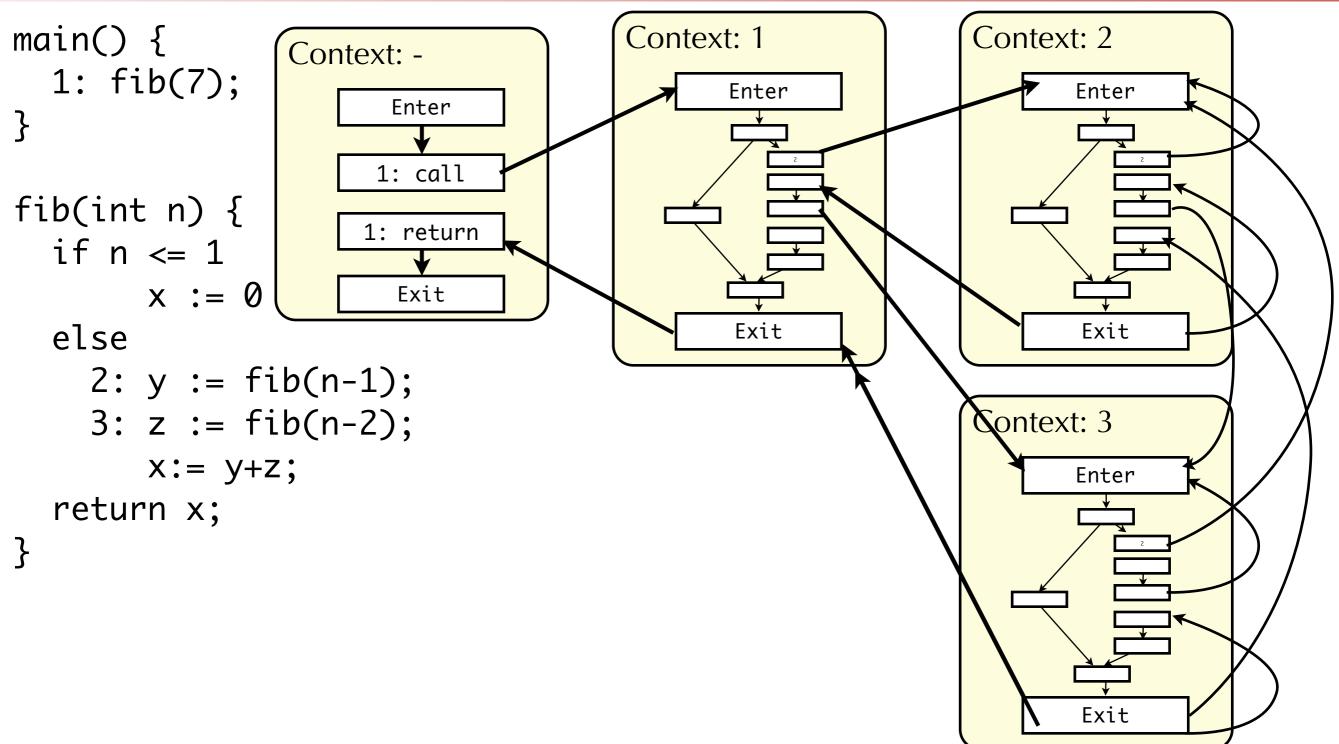
Fibonacci: context insensitive

```
main() {
  1: fib(7);
fib(int n) {
  if n <= 1
      x := 0
  else
    2: y := fib(n-1);
    3: z := fib(n-2);
       X := y+z;
  return x;
```



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Fibonacci: context sensitive, stack depth 1



Fibonacci: context sensitive, stack depth 2

