

Slides to Accompany *Programming Languages and Methodologies*

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Chapter 2, Part 3: Formal Grammars

The Derivation (or Parse) Tree

A derivation or parse tree, T , has the following characteristics:

1. The root of T is the starting symbol $S \in V_N$.
2. Leaf nodes of T are terminals $\in V_T$.
3. Interior^a nodes are nonterminals $\in V_N$.
4. The children of any non-leaf node^b represent the right hand side (RHS) of some production in P , where the parent node represents the corresponding LHS of the production.

Derivation trees are important when cataloging the production of x , since (except in cases of grammatical ambiguity) they show the structure of x *independently of the possible sequences*.

^aneither root nor leaf

^brecall leaf nodes have no children.

Derivation Tree Example

Recall the productions in an earlier grammar:

$$S \rightarrow AB$$

$$S \rightarrow C$$

$$A \rightarrow C$$

$$A \rightarrow a$$

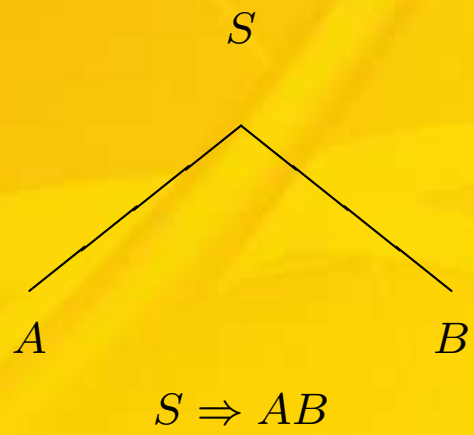
$$B \rightarrow b$$

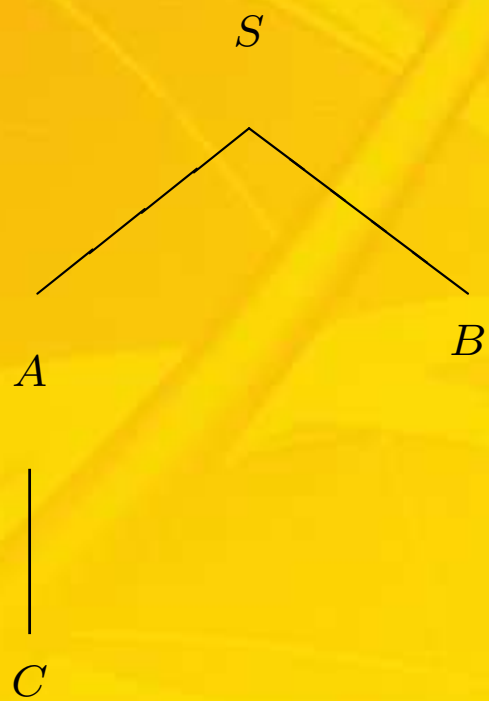
$$B \rightarrow c$$

$$C \rightarrow d$$

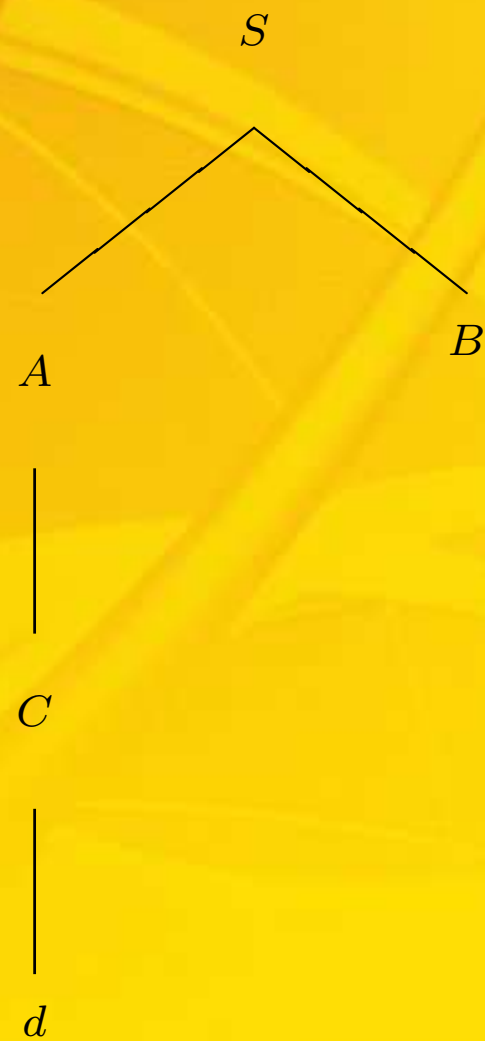
First, we show the derivation of string dc , using the replacement sequence:

$$S \Rightarrow AB \Rightarrow CB \Rightarrow dB \Rightarrow dc$$

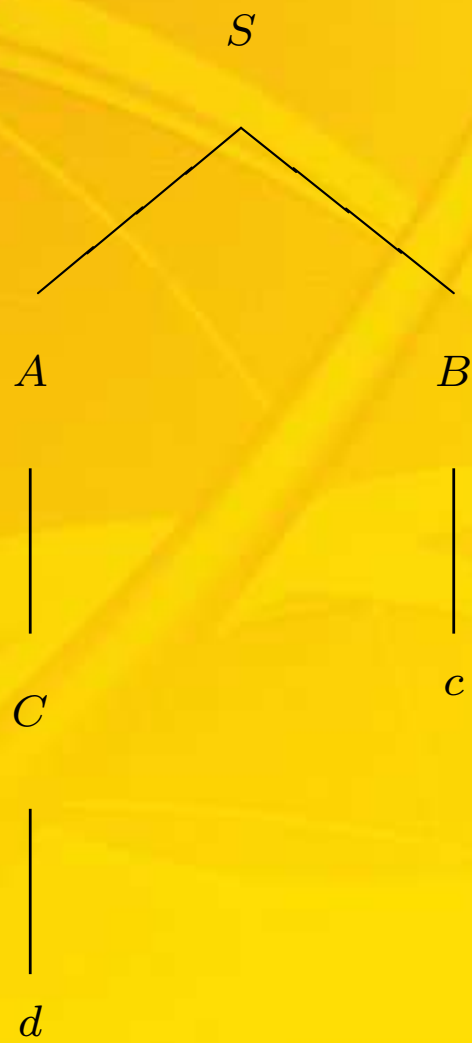




$$AB \Rightarrow CB$$



$$CB \Rightarrow dB$$



$dB \Rightarrow dc$

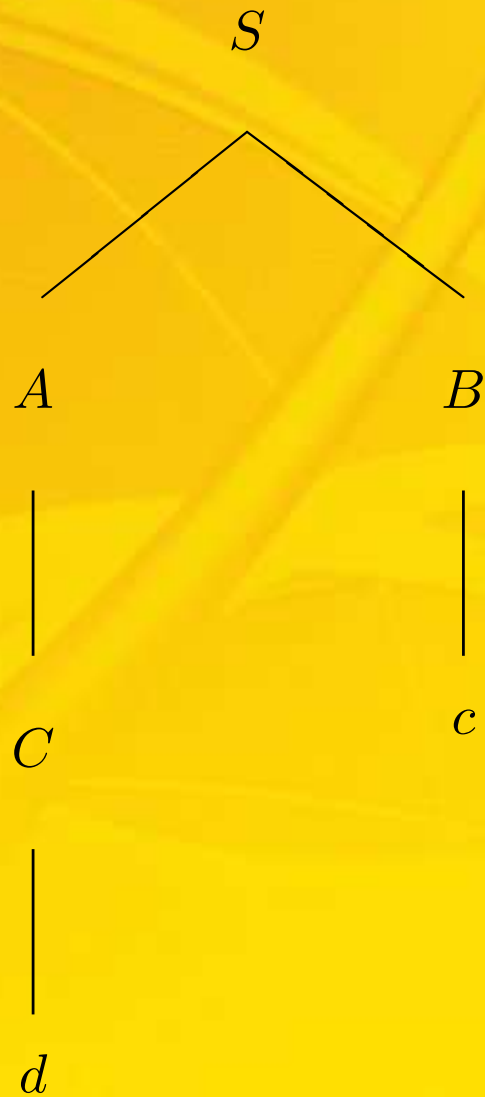
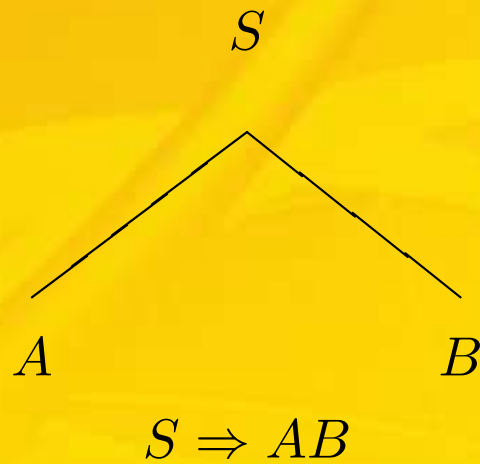


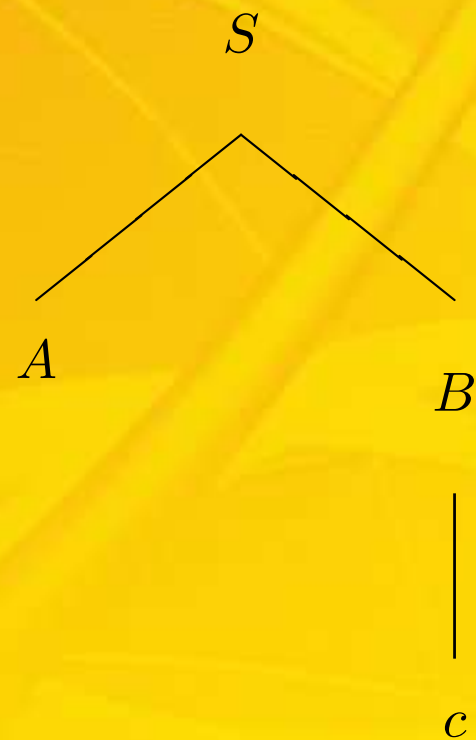
Figure 1: Derivation Tree Using the Production Sequence $S \Rightarrow AB \Rightarrow CB \Rightarrow dB \Rightarrow dc$

Now consider an alternative derivation *sequence* for the same string:

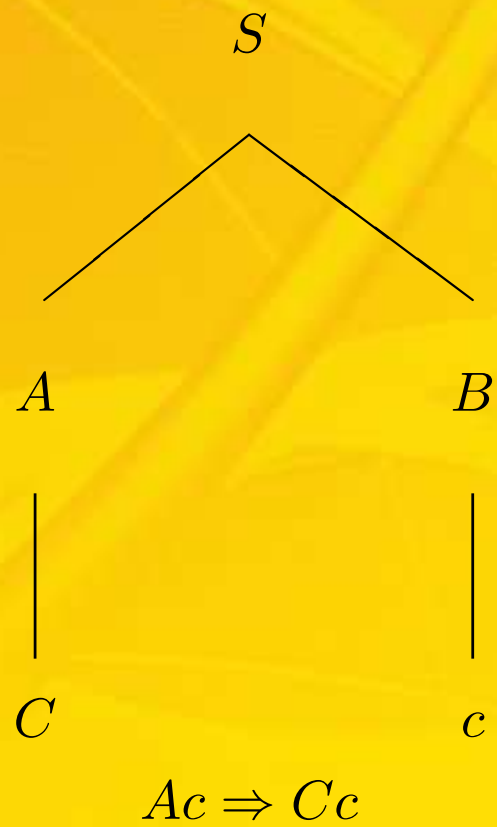
$$S \Rightarrow AB \Rightarrow Ac \Rightarrow Cc \Rightarrow dc$$

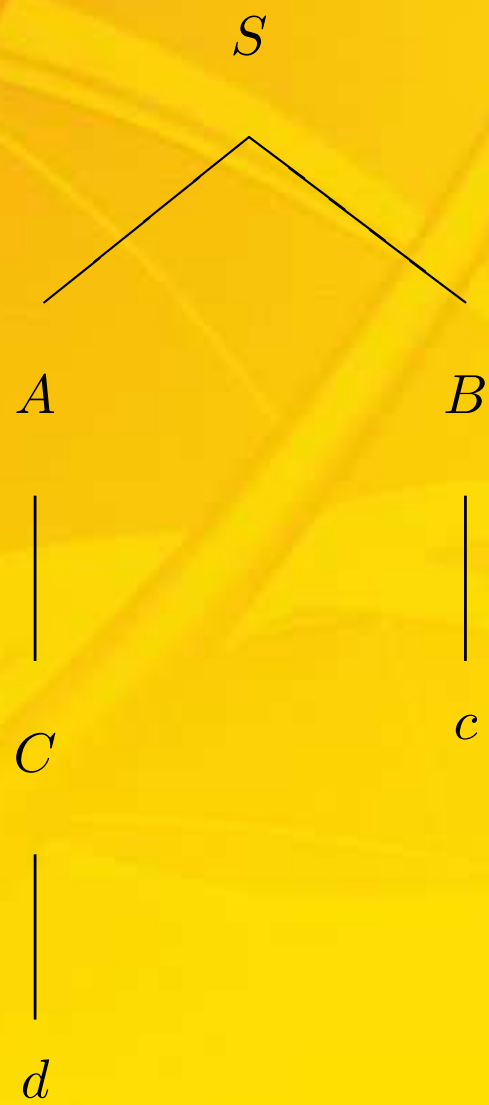
Development of the derivation tree with this sequence proceeds as follows:





$$AB \Rightarrow Ac$$





$$Cc \Rightarrow dc$$

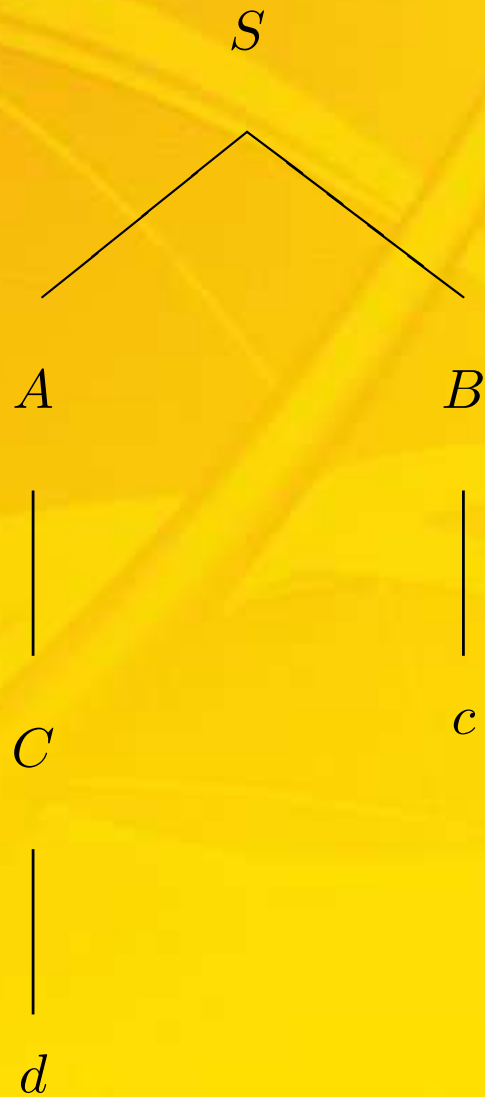


Figure 2: Derivation Tree Using the Production Sequence $S \Rightarrow AB \Rightarrow Ac \Rightarrow Cc \Rightarrow dc$

Grammatical/Syntactic Ambiguity

A grammar is ambiguous if any string in $L(G)^a$ has two or more distinct derivation trees.

^athe language generated by grammar G

Classic Examples of Ambiguity

Classic examples of potential ambiguity in programming languages include:

- Strings employing binary mathematical operators. For example, what is the underlying structure of:

$$a = b + c * a$$

Of course, specification of operator precedence together with the syntax specification may eliminate this problem.

- The use of nested "if-then-(optional) else" constructs. This is a major source of potential ambiguity, hence the syntax specification is usually augmented with a description of the corresponding interpretation algorithm; and

- Confusion between function invocation and array use. For example, how would (could) you interpret: $a(2)$? Is this function a called with argument 2 or the second (or third) element of array a ?