神秘模板库

Toy ASM Truck
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一切的开始

宏定义

数据结构

ST 表

二维

```
int f[maxn][maxn][10][10];
    inline int highbit(int x) { return 31 - __builtin_clz(x); }
    inline int calc(int x, int y, int xx, int yy, int p, int q) {
        return max(
            \max(f[x][y][p][q], f[xx - (1 << p) + 1][yy - (1 << q) + 1][p][q]),
            \max(f[xx - (1 << p) + 1][y][p][q], f[x][yy - (1 << q) + 1][p][q])
        );
   }
    void init() {
        FOR (x, 0, highbit(n) + 1)
        FOR (y, 0, highbit(m) + 1)
11
            FOR (i, 0, n - (1 << x) + 1)
12
            FOR (j, 0, m - (1 << y) + 1) {
13
                if (!x && !y) { f[i][j][x][y] = a[i][j]; continue; }
14
                f[i][j][x][y] = calc(
16
                    i + (1 << x) - 1, j + (1 << y) - 1,
17
                    max(x - 1, 0), max(y - 1, 0)
18
                );
19
            }
21
    inline int get_max(int x, int y, int xx, int yy) {
22
        return calc(x, y, xx, yy, highbit(xx - x + 1), highbit(yy - y + 1));
23
24
```

数学

模整数类

除法为整除, 请乘逆元。

```
template<int P>
    struct moint {
        int x;
        moint():x(0){}
        moint(int n) {x=n<0?n%P+P:n%P;}</pre>
        moint(ll n) {x=n<0?n%P+P:n%P;}</pre>
        int get()const{return (int)x;}
        moint &operator+=(moint b){x+=b.x;if(x>=P)x-=P;return *this;}
        moint &operator==(moint b) {x==b.x;if(x<0)x+=P;return *this;}</pre>
        moint &operator*=(moint b){x=1ll*x*b.x%P;return *this;}
        moint &operator/=(moint b) {x=x/b.x;return *this;}
11
        moint &operator%=(moint b){x=x%b.x;return *this;}
12
        moint operator+(moint b)const{return moint(*this)+=b;}
13
        moint operator-(moint b)const{return moint(*this)-=b;}
14
        moint operator*(moint b)const{return moint(*this)*=b;}
15
        moint operator/(moint b)const{return moint(*this)/=b;}
16
17
        moint operator%(moint b)const{return moint(*this)%=b;}
        moint operator+(int b)const{return moint(*this)+=moint(b);}
18
        moint operator-(int b)const{return moint(*this)-=moint(b);}
19
        moint operator*(int b)const{return moint(*this)*=moint(b);}
20
        moint operator/(int b)const{return moint(*this)/=moint(b);}
21
22
        moint operator%(int b)const{return moint(*this)%=moint(b);}
        bool operator==(moint b)const{return x==b.x;}
23
24
        bool operator>=(moint b)const{return x>=b.x;}
        bool operator!=(moint b)const{return x!=b.x;}
25
   };
26
    typedef moint<998244353> mint;
```

类欧几里得

```
• m = \lfloor \frac{an+b}{a} \rfloor.
               • f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor: 当 a \geq c or b \geq c 时,f(a,b,c,n) = (\frac{a}{c})n(n+1)/2 + (\frac{b}{c})(n+1) + f(a \bmod c, b \bmod c, c, n); 否则 f(a,b,c,n) = nm - f(c,c-b-1,a,m-1)。
               g(a \bmod c, b \bmod c, c, n); 否则 g(a, b, c, n) = \frac{1}{2}(n(n+1)m - f(c, c-b-1, a, m-1) - h(c, c-b-1, a, m-1))。
               • h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2: \exists a \geq c \text{ or } b \geq c \text{ ft}, \ h(a,b,c,n) = (\frac{a}{c})^2 n(n+1)(2n+1)/6 + (\frac{b}{c})^2 (n+1) + (\frac{
                     (\frac{a}{c})(\frac{b}{c})n(n+1)+h(a \mod c,b \mod c,c,n)+2(\frac{a}{c})g(a \mod c,b \mod c,c,n)+2(\frac{b}{c})f(a \mod c,b \mod c,c,n); 否则
                    h(a,b,c,n) = nm(m+1) - 2q(c,c-b-1,a,m-1) - 2f(c,c-b-1,a,m-1) - f(a,b,c,n)
        struct ans{mint f,g,h;};
        static ans calc(int a,int b,int c,int n){
                 ans ret;
                 if(!a){
                         ret.f=mint(b/c)*(n+1);
                         ret.g=mint(b/c)*n*(n+1)*iv2;
                         ret.h=mint(b/c)*(b/c)*(n+1);
                         return ret:
                 if(a>=c||b>=c){
10
                         ans to=calc(a%c,b%c,c,n);
11
                         ret.f=mint(a/c)*n*(n+1)*iv2+mint(b/c)*(n+1)+to.f;
12
                         ret.g=mint(a/c)*n*(n+1)*(n*2+1)*iv6+mint(b/c)*n*(n+1)*iv2+to.g;
13
14
                         ret.h=mint(a/c)*(a/c)*n*(n+1)*(n*2+1)*iv6+mint(b/c)*(b/c)*(n+1)+\
15
                                  mint(a/c)*(b/c)*n*(n+1)+to.h+mint(a/c)*2*to.g+mint(b/c)*2*to.f;
16
17
                }else{
                         ll m=(1ll*a*n+b)/c;
18
                         ans to=calc(c,c-b-1,a,m-1);
                         ret.f=mint(n*m%P)-to.f;
20
                         ret.g=(mint(m*n%P*(n+1)%P)-to.f-to.h)*iv2;
                         ret.h=mint(m*n%P*(m+1)%P)-to.g*2-to.f*2-ret.f;
22
23
                         return ret;
24
       }
25
        Pollard-Rho
        template < const int test_case > // set 8 usually
        struct Pollard Rho {
                 vector<long long> fac;
                 long long quick_pow(long long a, long long b, long long mod) {
                         long long ans = 1;
                         while (b) {
                                  if (b&1) ans = (__int128)ans*(__int128)a%mod;
                                  b >>= 1, a = (__int128)a*(__int128)a%mod;
                         }
10
11
                 bool Miller_Rabin(long long n) {// return if n is a prime
12
                         if (n < 3) return n == 2;
                         long long a = n-1, b = 0;
14
15
                         while (a\%2 == 0) a /= 2, ++b;
                         for (int i = 0, j; i < test_case; i++) {</pre>
16
                                  long long x = rand()\%(n-2)+2, v = quick_pow(x, a, n);
17
                                  if (v == 1 || v == n-1) continue;
                                  for (j = 0; j < b; j++) {
19
                                          v = (\_int128)v*(\_int128)v%n;
                                          if (v == n-1) break;
21
22
                                  if (j >= b) return false;
23
                         }
24
                         return true;
25
26
                 long long f(long long x, long long c, long long n) { return ((\_int128)x * x + c) % n; }
27
                 long long rho(long long x) {
28
                         long long s = 0, t = 0;
29
```

long long $c = (_int128) rand() % (x - 1) + 1;$

```
int step = 0, goal = 1;
31
32
            long long val = 1;
            for (goal = 1;; goal <<= 1, s = t, val = 1) {</pre>
33
                 for (step = 1; step <= goal; ++step) {</pre>
34
35
                     t = f(t, c, x);
                     val = (_int128)val * abs(t - s) % x;
36
37
                     if ((step % 127) == 0) {
                         long long d = __gcd(val, x);
38
                         if (d > 1) return d;
39
                     }
40
                 }
41
                 long long d = __gcd(val, x);
42
                 if (d > 1) return d;
43
44
45
        }
        void find(long long x) {
46
47
            if (x == 1) return;
            if (Miller_Rabin(x)) {
48
                 fac.push_back(x);
                 return;
50
51
52
            long long p = x;
            while (p >= x) p = rho(x);
53
            //while ((x % p) == 0) x /= p;
            find(x/p), find(p);
55
56
        vector<long long> factor(long long n) {// return the factors of n}
57
            srand((unsigned)time(NULL));
58
59
            fac.clear();
            find(n);
60
            sort(fac.begin(), fac.end());
61
            return fac;
62
63
   };
    ex-gcd
    template<typename T>
1
2
    struct ex_gcd {
        T gcd(const T a, const T b, T &x, T &y) \{// x'=x_0+b/gcd, y'=y_0-a/gcd\}
3
4
            if (b == 0) {x = 1, y = 0; return a; }
            T d = gcd(b, a\%b, x, y);
            T t = x;
            x = y;
            y = t - a/b*y;
            return d;
10
        T inv(const T a, const T m) {// return -1 if inv is not exist
11
12
            if (a == 0 || m <= 1) return -1;
            T x, y, d = gcd(a, m, x, y);
13
            if (d != 1) return -1;
14
            return (x%m+m)%m;
15
16
   } ;
17
    crt
    template<typename T>
    struct crt {
        ex_gcd<T> *exgcd = new ex_gcd<T>();
        T cal(const T *a, const T *m, const int n) \{// a[1..n], m[1..n], gcd(m_i) = 1\}
            T M = 1, ans = 0;
            for (int i = 1; i <= n; i++) M *= m[i];</pre>
            for (int i = 1; i <= n; i++)</pre>
                 (ans += (\_int128)a[i]*(M/m[i])%M*exgcd->inv(M/m[i], m[i])%M) ~\%= M;
            return ans;
        }
   };
```

ex-crt

```
template<typename T>
1
2
    struct ex_crt {
        ex_gcd<T> *exgcd = new ex_gcd<T>();
        T cal(T *a, T *m, const int n) \{// a[1..n], m[1..n], return -1 if no ans
            T x, y, gcd, lcm;
            for (int i = 2; i <= n; i++) {</pre>
                gcd = exgcd - gcd(m[1], m[i], x, y);
                 if ((a[i]-a[1])%gcd) return -1;
                lcm = (__int128)m[1]*m[i]/gcd;
                x = (_{int128})x*(a[i]-a[1])/gcd%lcm;
                gcd = m[i]/gcd;
11
                x = (x\%gcd+gcd)\%gcd;
                a[1] = ((__int128)m[1]*x%lcm+a[1])%lcm, m[1] = lcm;
13
            return a[1];
        }
16
   } ;
    Meissel-Lehmer
    求解 1e11 内的质数个数,约为 O(n^{2/3})。
   namespace pcf{
    #define chkbit(ar, i) (((ar[(i) >> 6]) & (1 << (((i) >> 1) & 31))))
    #define setbit(ar, i) (((ar[(i) >> 6]) \mid= (1 << (((i) >> 1) & 31))))
    #define isprime(x) (( (x) && ((x)&1) && (!chkbit(ar, (x)))) || ((x) == 2))
        const int MAXN=100;
        const int MAXM=10001;
        const int MAXP=40000;
        const int MAX=400000;
        long long dp[MAXN][MAXM];
        unsigned int ar[(MAX >> 6) + 5] = {0};
10
        int len = 0, primes[MAXP], counter[MAX];
11
12
        void Sieve(){
            setbit(ar, 0), setbit(ar, 1);
13
            for (int i = 3; (i * i) < MAX; i++, i++){</pre>
14
                 if (!chkbit(ar, i)){
15
                     int k = i << 1;</pre>
16
                     for (int j = (i * i); j < MAX; j += k) setbit(ar, j);
17
                }
18
            for (int i = 1; i < MAX; i++){</pre>
20
                 counter[i] = counter[i - 1];
22
                if (isprime(i)) primes[len++] = i, counter[i]++;
23
24
        void init(){
25
            Sieve();
            for (int n = 0; n < MAXN; n++){
27
                 for (int m = 0; m < MAXM; m++){</pre>
28
                     if (!n) dp[n][m] = m;
29
                     else dp[n][m] = dp[n - 1][m] - dp[n - 1][m / primes[n - 1]];
30
                }
            }
32
33
34
        long long phi(long long m, int n){
            if (n == 0) return m;
35
            if (primes[n - 1] >= m) return 1;
            if (m < MAXM && n < MAXN) return dp[n][m];</pre>
37
38
            return phi(m, n - 1) - phi(m / primes[n - 1], n - 1);
39
        long long Lehmer(long long m){
40
41
            if (m < MAX) return counter[m];</pre>
            long long w, res = 0;
42
43
            int i, a, s, c, x, y;
            s = sqrt(0.9 + m), y = c = cbrt(0.9 + m);
44
            a = counter[y], res = phi(m, a) + a - 1;
45
            for (i = a; primes[i] <= s; i++) res = res - Lehmer(m / primes[i]) + Lehmer(primes[i]) - 1;</pre>
46
            return res;
47
```

```
}
48
49
   }
    int main(){
50
51
        pcf::init();
52
        long long n;
        while (scanf("%lld", &n) != EOF){
53
54
            printf("%lld\n",pcf::Lehmer(n));
        }
55
        return 0;
56
57
   }
    图论
   LCA
       倍增
    void dfs(int u, int fa) {
        pa[u][0] = fa; dep[u] = dep[fa] + 1;
2
        FOR (i, 1, SP) pa[u][i] = pa[pa[u][i - 1]][i - 1];
3
4
        for (int& v: G[u]) {
            if (v == fa) continue;
            dfs(v, u);
        }
8
   }
    int lca(int u, int v) {
11
        if (dep[u] < dep[v]) swap(u, v);</pre>
12
        int t = dep[u] - dep[v];
13
        FOR (i, 0, SP) if (t & (1 << i)) u = pa[u][i];
14
        FORD (i, SP - 1, -1) {
15
            int uu = pa[u][i], vv = pa[v][i];
            if (uu != vv) { u = uu; v = vv; }
17
18
        return u == v ? u : pa[u][0];
19
20
    计算几何
    二维几何: 点与向量
   #define y1 yy1
   #define nxt(i) ((i + 1) % s.size())
   typedef double LD;
   const LD PI = 3.14159265358979323846;
   const LD eps = 1E-10;
   int sgn(LD x) { return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); }
    struct L;
    struct P;
    typedef P V;
    struct P {
11
        LD x, y;
        explicit P(LD x = 0, LD y = 0): x(x), y(y) {}
12
13
        explicit P(const L& l);
   };
14
    struct L {
        Ps, t;
16
        L() {}
17
        L(P s, P t): s(s), t(t) {}
18
   };
19
   P operator + (const P& a, const P& b) { return P(a.x + b.x, a.y + b.y); }
21
    P operator - (const P& a, const P& b) { return P(a.x - b.x, a.y - b.y); }
22
   P operator * (const P& a, LD k) { return P(a.x * k, a.y * k); }
23
   P operator / (const P& a, LD k) { return P(a.x / k, a.y / k); }
24
```

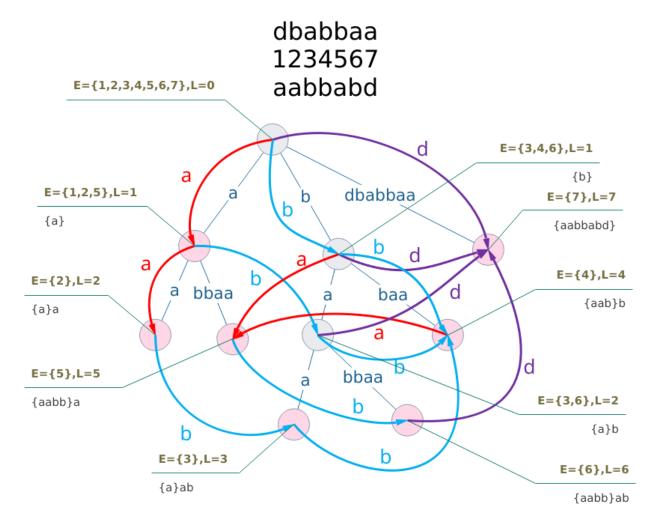
inline bool operator < (const P& a, const P& b) {</pre>

return $sgn(a.x - b.x) < 0 \mid \mid (sgn(a.x - b.x) == 0 \&\& sgn(a.y - b.y) < 0);$

```
27
28
   bool operator == (const P& a, const P& b) { return !sgn(a.x - b.x) && !sgn(a.y - b.y); }
   P::P(const L& l) { *this = l.t - l.s; }
29
   ostream &operator << (ostream &os, const P &p) {
        return (os << "(" << p.x << "," << p.y << ")");
32
33
    istream &operator >> (istream &is, P &p) {
        return (is >> p.x >> p.y);
34
35
   LD dist(const P& p) { return sqrt(p.x * p.x + p.y * p.y); }
37
   LD dot(const V& a, const V& b) { return a.x * b.x + a.y * b.y; }
   LD det(const V& a, const V& b) { return a.x * b.y - a.y * b.x; }
   LD cross(const P& s, const P& t, const P& o = P()) { return det(s - o, t - o); }
```

字符串

后缀自动机



多项式

NTT 模数

```
NTTPrimes = {1053818881, 1051721729, 1045430273, 1012924417, 1007681537, 1004535809, 998244353, 985661441,

    976224257, 975175681};

    NTTPrimitiveRoots = {7, 6, 3, 5, 3, 3, 3, 3, 17};
    FFT
    namespace FFT{
        const db pi=acos(-1);
2
        struct cp{
3
4
            db re, im;
            cp(db _re=0,db _im=0){re=_re;im=_im;}
            cp operator +(cp b){return cp(re+b.re,im+b.im);}
            cp operator -(cp b){return cp(re-b.re,im-b.im);}
            cp operator *(cp b){return cp(re*b.re-im*b.im,re*b.im+im*b.re);}
        };
        int r[N];cp c[N<<1];</pre>
10
        inline void fft(cp *a,int f,int n){
11
12
            rep(i,0,n-1) if(r[i]>i) swap(a[r[i]],a[i]);
             for(int i=1;i<n;i<<=1){</pre>
13
                 cp wn(cos(pi/i),f*sin(pi/i));
14
                 for(int j=0,p=(i<<1);j<n;j+=p){</pre>
15
                     cp w(1,0);
16
                     for(int k=0; k<i; ++k, w=w*wn){</pre>
17
                          cp x=a[j+k], y=w*a[j+k+i];
18
                          a[j+k]=x+y;a[j+k+i]=x-y;
19
20
                     }
                 }
21
22
            if(f==-1){rep(i,0,n-1) a[i].re/=n,a[i].im/=n;}
23
24
        inline int mul(db *a,db *b,int n,int m){
25
            n+=m;rep(i,0,n) c[i]=cp(a[i],b[i]);
26
             int l=0;m=n;for(n=1;n<=m;n<<=1) ++l;</pre>
27
28
            rep(i,0,n-1) r[i]=(r[i>>1]>>1)|((i&1)<<(l-1));
29
            rep(i,m+1,n) c[i]=cp(0,0);
             fft(c,1,n);rep(i,0,n-1) c[i]=c[i]*c[i];
30
             fft(c,-1,n);
31
            rep(i,0,m) a[i]=c[i].im/2;
32
33
            return n;
34
    }
35
    NTT
    namespace NTT{
1
        const int P=998244353,g=3,ig=332748118;
        inline int qpow(int a,int b){int q=1;while(b){if(b&1)q=1LL*q*a%P;a=1LL*a*a%P;b>>=1;}return q;}
        int r[N],ow[N],inv[N];
        inline void ntt(int *a,int f,int n){
             rep(i,0,n-1) if(r[i]>i) swap(a[i],a[r[i]]);
             for(int i=1;i<n;i<<=1){</pre>
                 int wn=qpow(f,(P-1)/(i<<1));</pre>
                 ow[0] = 1; rep(k,1,i-1) \ ow[k] = 1 \\ LL + ow[k-1] + wn\%P;
                 for(int j=0,p=(i<<1);j<n;j+=p){</pre>
10
11
                     for(int k=0;k<i;++k){</pre>
                          int x=a[j+k],y=1LL*ow[k]*a[j+k+i]%P;
12
                          a[j+k]=(x+y)%P;a[j+k+i]=(x+P-y)%P;
                     }
14
                 }
15
16
            if(f==ig){
17
                 int iv=qpow(n,P-2);
                 rep(i,0,n-1) a[i]=1LL*a[i]*iv%P;
19
21
        int tma[N],tmb[N];
22
```

```
inline int mul(int *a,int *b,int n,int m,int ci){
23
            int _n=n,_m=m,l=0;m+=n;for(n=1;n<=m;n<<=1) ++l;</pre>
24
            rep(i,0,n-1) r[i]=(r[i>>1]>>1)|((i&1)<<(l-1));
25
            rep(i,0,n-1) tma[i]=a[i];rep(i,0,n-1) tmb[i]=b[i];
26
            rep(i,_n+1,n) tma[i]=0;rep(i,_m+1,n) tmb[i]=0;
27
            ntt(tma,g,n);ntt(tmb,g,n);
28
29
            while(ci){
                if(ci&1) rep(i,0,n-1) tma[i]=1LL*tma[i]*tmb[i]%P;
30
                rep(i,0,n-1) tmb[i]=1LL*tmb[i]*tmb[i]%P;
31
32
                ci>>=1;
            }
33
34
            ntt(tma,ig,n);
            rep(i,0,n-1) a[i]=tma[i];
35
            return n;
36
        }
37
    }
38
    inline void prepare(){
39
        //NTT inv
40
        using NTT::inv;using NTT::P;
41
        inv[1]=1;rep(i,2,N-1) inv[i]=1LL*(P-P/i)*inv[P%i]%P;
42
   }
43
```