CHAPTER 4

LISTS

All the programs in this file are selected from

Ellis Horowitz, Sartaj Sahni, and Susan Anderson-Freed "Fundamentals of Data Structures in C", Computer Science Press, 1992.

Introduction

Array

successive items locate a fixed distance

disadvantage

- data movements during insertion and deletion
- waste space in storing n ordered lists of varying size

possible solution

linked list

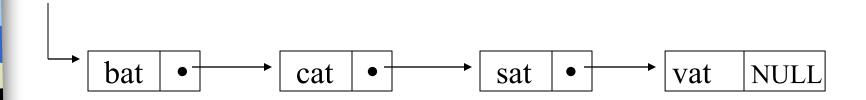
4.1.1 Pointer Can Be Dangerous

```
pointer
  int i, *pi;
  pi = &i; i=10 \text{ or } *pi=10
pi= malloc(size of(int));
  /* assign to pi a pointer to int */
pf=(float *) pi;
  /* casts an int pointer to a float pointer */
```

4.1.2 Using Dynamically Allocated Storage

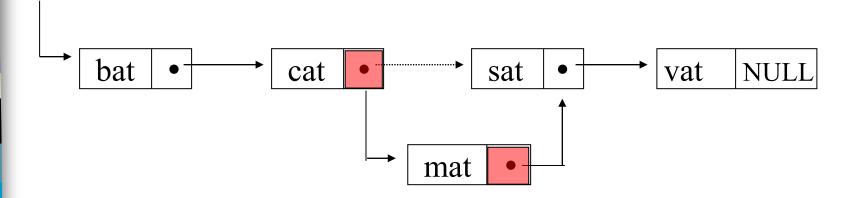
```
int i, *pi;
float f, *pf;
pi = (int *) malloc(sizeof(int));
                                           request memory
pf = (float *) malloc (sizeof(float));
*pi =1024;
*pf = 3.14;
printf("an integer = \%d, a float = \%f\n", *pi, *pf);
free(pi);
                       return memory
free(pf);
*Program4.1: Allocation and deallocation of pointers (p.138)
```

4.2 SINGLY LINKED LISTS

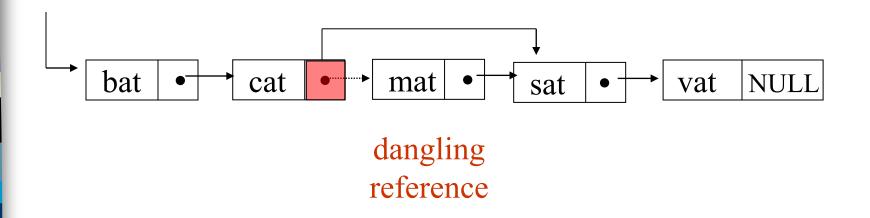


*Figure 4.1: Usual way to draw a linked list (p.139)

Insertion



*Figure 4.2: Insert mat after cat (p.140)



***Figure 4.3:** Delete *mat* from list (p.140)

Example 4.1: create a linked list of words

```
Declaration
typedef struct list node, *list pointer;
typedef struct list node {
        char data [4];
        list pointer link;
Creation
list pointer ptr =NULL;
Testing
#define IS EMPTY(ptr) (!(ptr))
Allocation
ptr=(list pointer) malloc (sizeof(list node));
```

*Figure 4.4:Referencing the fields of a node(p.142)

Example: create a two-node list

```
ptr
                                   NULL
typedef struct list node *list pointer;
typedef struct list node {
        int data;
        list pointer link;
list pointer ptr =NULL
```

Example 4.2: (p.142)

```
list pointer create2()
/* create a linked list with two nodes */
  list pointer first, second;
  first = (list pointer) malloc(sizeof(list node));
  second = ( list pointer) malloc(sizeof(list node));
  second -> link = NULL;
  second \rightarrow data = 20;
                                ptr
  first \rightarrow data = 10;
  first -> link = second;
                                       10
                                                             NULL
  return first;
                            *Program 4.2:Create a tow-node list (p.143)
```

Pointer Review (1)

```
pi = &i;

i 1000

*pi ?
```

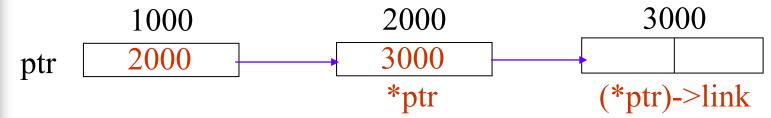
Pointer Review (2)

```
typedef struct list node *list_pointer;
typedef struct list node {
               int data;
               list pointer link;
list pointer ptr = NULL;
                                  ptr->data (*ptr).data
ptr
ptr = malloc(sizeof(list node));
                             2000
           1000
ptr
          2000
                             data
                                      link
                            CHAPTER 4
```

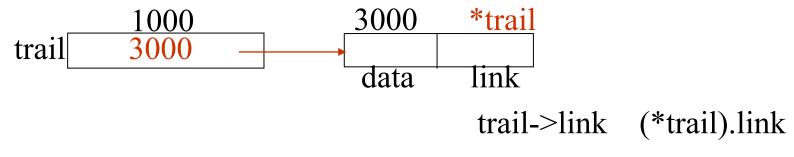
Pointer Review (3)

void delete(list_pointer *ptr, list_pointer trail, list_pinter node)

ptr: a pointer point to a pointer point to a list node



trail (node): a pointer point to a list node



Pointer Review (4)

element delete(stack pointer *top) top delete(&st) vs. delete(st) 300 st 200 top 200 300 300 200 400 400 top st 300 300 400 Does not change. **CHAPTER 4**

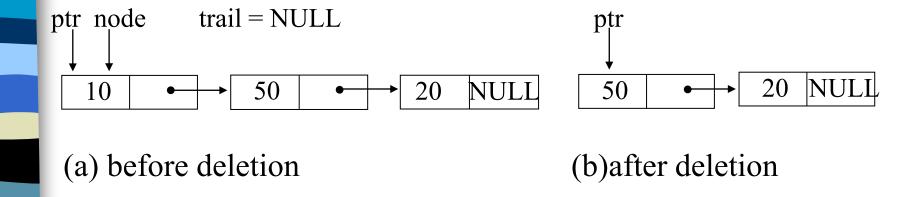
List Insertion:

Insert a node after a specific node

```
void insert(list_pointer *ptr, list_pointer node)
{
/* insert a new node with data = 50 into the list ptr after node */
    list_pointer temp;
    temp = (list_pointer) malloc(sizeof(list_node));
    if (IS_FULL(temp)) {
        fprintf(stderr, "The memory is full\n");
        exit (1);
    }
}
```

```
temp->data = 50;
if (*ptr) { noempty list
   temp->link =node ->link;
   node->link = temp;
                            pţr
                                                    20
                                                        NULL
else { empty list
                                    10
  temp->link = NULL;
                           node
  *ptr =temp;
                                           50
                                   temp
```

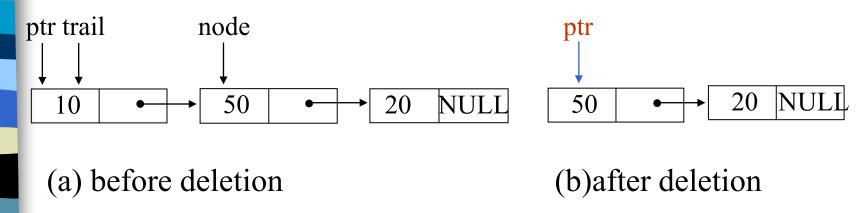
*Program 4.3: Simple insert into front of list (p.144)



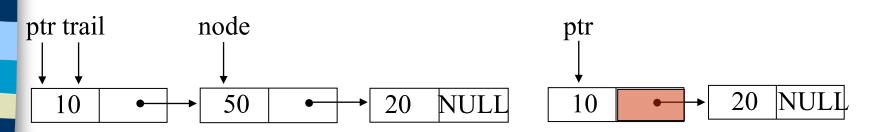
*Figure 4.7: List after the function call *Delete*(&ptr,NULL.ptr);(p.145)

List Deletion

Delete the first node.



Delete node other than the first node.



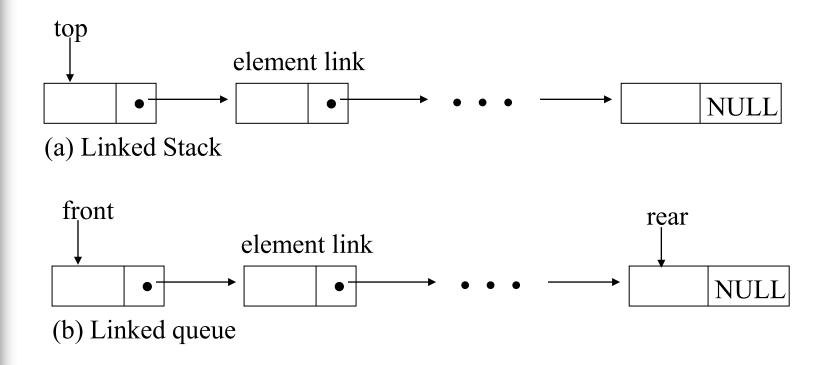
```
void delete(list pointer *ptr, list pointer trail,
                                                  list pointer node)
  delete node from the list, trail is the preceding node
                                   trail
                                            node
  ptr is the head of the list */
   if (trail)
                                              50
                               10
                                                           20
                                                                NULI
     trail->link = node->link;
   else
                                                 NULL
                                             20
                               10
     *ptr = (*ptr) ->link;
    free(node);
               node
   ptr
                                              50
                 50
                               20
                                 CHAPTER 4
                                                                      20
```

Print out a list (traverse a list)

```
void print_list(list_pointer ptr)
{
    printf("The list ocntains: ");
    for (; ptr; ptr = ptr->link)
        printf("%4d", ptr->data);
    printf("\n");
}
```

*Program 4.5: Printing a list (p.146)

4.3 DYNAMICALLY LINKED STACKS AND QUEUES



*Figure 4.10: Linked Stack and queue (p.147)

Represent n stacks

```
#define MAX STACKS 10 /* maximum number of stacks */
typedef struct {
       int key;
       /* other fields */
       } element;
typedef struct stack *stack pointer;
typedef struct stack {
        element item;
       stack pointer link;
stack pointer top[MAX STACKS];
```

Represent n queues

```
#define MAX_QUEUES 10 /* maximum number of queues */
typedef struct queue *queue_pointer;
typedef struct queue {
    element item;
    queue_pointer link;
    };
queue_pointer front[MAX_QUEUE], rear[MAX_QUEUES];
```

```
Push in the linked stack
void add(stack pointer *top, element item)
 /* add an element to the top of the stack */
  stack pointer temp =
               (stack pointer) malloc (sizeof (stack));
 if (IS FULL(temp)) {
   fprintf(stderr, "The memory is full\n");
   exit(1);
   temp->item = item;
   temp->link = *top;
    *top= temp;
                         *Program 4.6:Add to a linked stack (p.149)
```

```
pop from the linked stack
element delete(stack pointer *top) {
/* delete an element from the stack */
  stack pointer temp = *top;
  element item;
  if (IS EMPTY(temp)) {
    fprintf(stderr, "The stack is empty\n");
    exit(1);
  item = temp->item;
  *top = temp->link;
   free(temp);
   return item;
*Program 4.7: Delete from a linked stack (p.149)
```

enqueue in the linked queue

```
void addq(queue pointer *front, queue pointer *rear, element item)
{ /* add an element to the rear of the queue */
 queue pointer temp =
                (queue pointer) malloc(sizeof (queue));
 if (IS FULL(temp)) {
   fprintf(stderr, "The memory is full\n");
   exit(1);
   temp->item = item;
   temp->link = NULL;
   if (*front) (*rear) \rightarrow link = temp;
   else *front = temp;
    *rear = temp; }
                              CHAPTER 4
```

dequeue from the linked queue (similar to push)

```
element deleteq(queue pointer *front) {
/* delete an element from the queue */
  queue pointer temp = *front;
  element item;
  if (IS EMPTY(*front)) {
    fprintf(stderr, "The queue is empty\n");
    exit(1);
  item = temp->item;
  *front = temp->link;
   free(temp);
   return item;
                         CHAPTER 4
```

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Polynomials

$$A(x) = a_{m-1}x^{e_{m-1}} + a_{m-2}x^{e_{m-2}} + ... + a_0x^{e_0}$$

Representation

```
typedef struct poly_node *poly_pointer;
typedef struct poly_node {
    int coef;
    int expon;
    poly_pointer link;
};
poly pointer a, b, c;
```

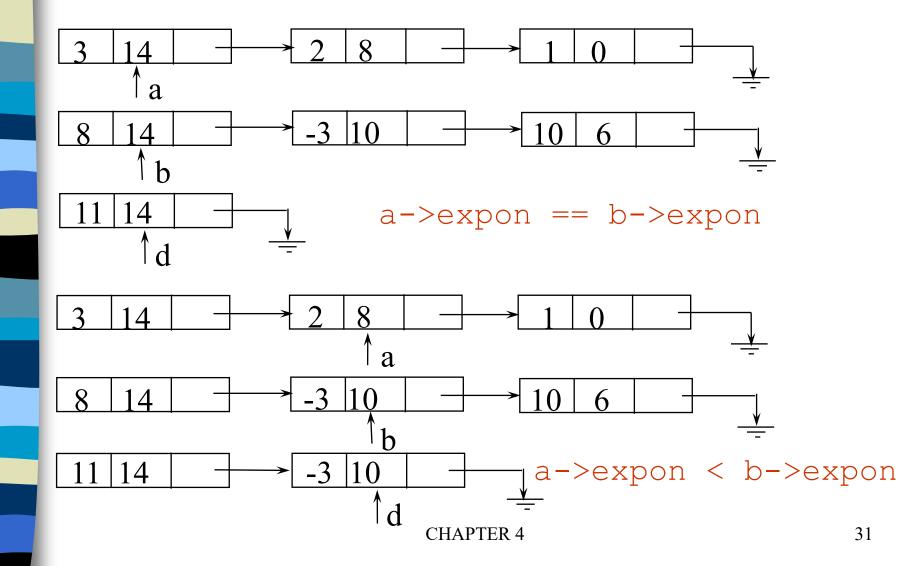
| coef expon | link |
|------------|------|
|------------|------|

Examples

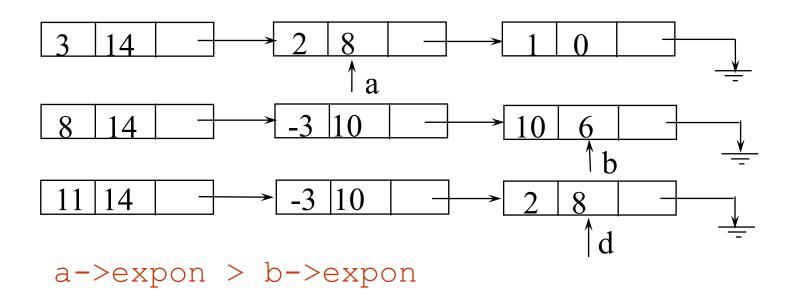
$$a = 3x^{14} + 2x^8 + 1$$

$$b = 8x^{14} - 3x^{10} + 10x^6$$

Adding Polynomials



Adding Polynomials (Continued)



Alogrithm for Adding Polynomials

```
poly_pointer padd(poly_pointer a, poly_pointer b)
{
    poly_pointer front, rear, temp;
    int sum;
    rear = (poly_pointer)malloc(sizeof(poly_node));
    if (IS_FULL(rear)) {
        fprintf(stderr, "The memory is full\n");
        exit(1);
    }
    front = rear;
    while (a && b) {
        switch (COMPARE(a->expon, b->expon)) {
```

```
case -1: /* a->expon < b->expon */
             attach(b->coef, b->expon, &rear);
             b= b->link;
             break:
        case 0: /* a->expon == b->expon */
             sum = a - coef + b - coef;
             if (sum) attach(sum,a->expon,&rear);
             a = a \rightarrow link; b = b \rightarrow link;
             break;
        case 1: /* a->expon > b->expon */
             attach(a->coef, a->expon, &rear);
             a = a - > link;
for (; a; a = a - \lambda 
    attach(a->coef, a->expon, &rear);
for (; b; b=b->link)
    attach(b->coef, b->expon, &rear);
rear->link = NULL;
temp = front; front = front->link; free(temp);
return front;
```

Delete extra initial node.

Analysis

(1) coefficient additions $0 \le \text{additions} \le \min(m, n)$

where m (n) denotes the number of terms in A (B).

(2) exponent comparisons extreme case

$$e_{m-1} > f_{m-1} > e_{m-2} > f_{m-2} > ... > e_0 > f_0$$

m+n-1 comparisons

(3) creation of new nodes extreme case m + n new nodes

Attach a Term

```
void attach (float coefficient, int exponent,
             poly pointer *ptr)
/* create a new node attaching to the node pointed to
  by ptr. ptr is updated to point to this new node. */
    poly pointer temp;
    temp = (poly pointer) malloc(sizeof(poly node));
    if (IS FULL(\overline{temp})) {
         fprintf(stderr, "The memory is full \n");
        exit(1);
    temp->coef = coefficient;
    temp->expon = exponent;
    (*ptr) - > link = temp;
    *ptr = temp;
```

A Suite for Polynomials

```
e(x) = a(x) * b(x) + d(x)

poly_pointer a, b, d, e; read_poly()

a = read_poly();
b = read_poly();
d = read_poly();
temp = pmult(a, b);
e = padd(temp, d);
print_poly()

pmult()

pmult()
```

temp is used to hold a partial result. By returning the nodes of temp, we may use it to hold other polynomials

Erase Polynomials

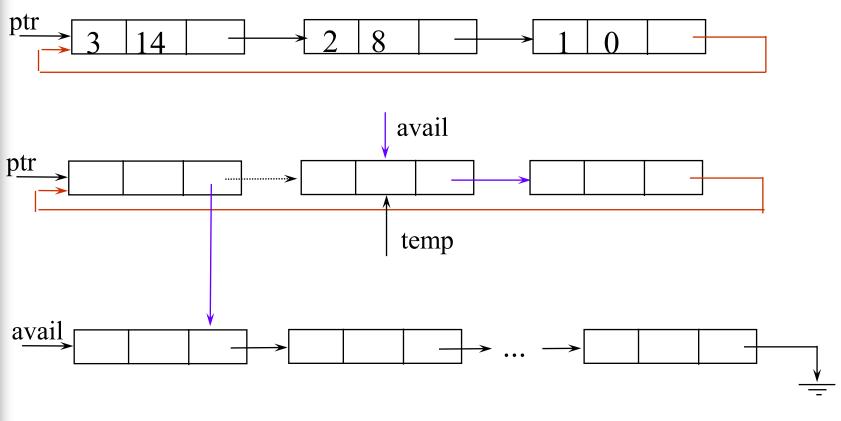
```
void earse(poly_pointer *ptr)
{
/* erase the polynomial pointed to by ptr */
    poly_pointer temp;
    while (*ptr) {
        temp = *ptr;
        *ptr = (*ptr)->link;
        free(temp);
    }
}
```

O(n)

CHAPTER 4

Circularly Linked Lists

circular list vs. chain



Maintain an Available List

```
poly_pointer get_node(void)
  poly_pointer node;
if (avail) {
      node = avail;
      avail = avail->link:
  else {
      node = (poly pointer) malloc(sizeof(poly node));
      if (IS FULL (node))
           prīntf(stderr, "The memory is full\n");
           exit(1);
  return node;
```

Maintain an Available List (Continued)

Insert ptr to the front of this list

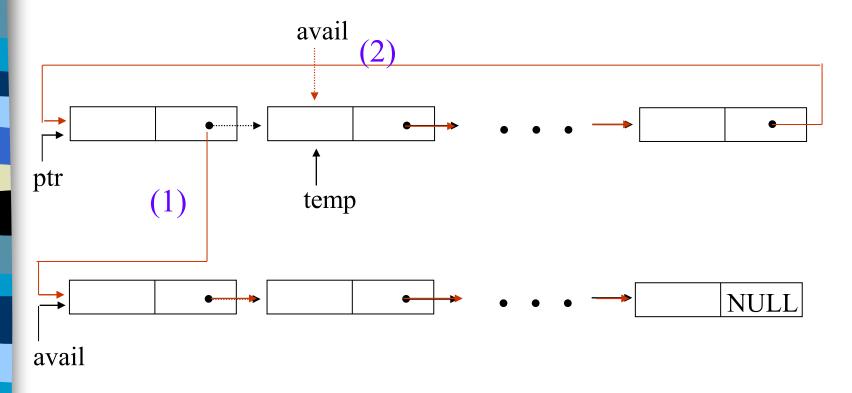
```
void ret node(poly pointer ptr)
  ptr->link = avail;
  avail = ptr;
                            Erase a circular list (see next page)
void cerase(poly pointer *ptr)
     poly pointer temp;
if (*ptr) {
          temp = (*ptr)->link;
(*ptr)->link = avail;
          avail = temp; (2)
*ptr = NULL;
```

Independent of # of nodes in a list O(1) constant time

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4.4.4 Representing Polynomials As Circularly Linked Lists

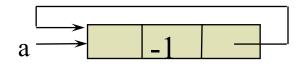


*Figure 4.14: Returning a circular list to the avail list (p.159)

Head Node

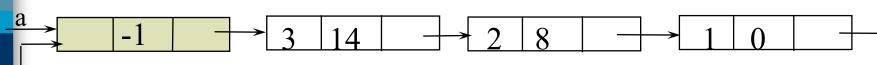
Represent polynomial as circular list.

(1) zero



Zero polynomial

(2) others



$$a = 3x^{14} + 2x^8 + 1$$

Another Padd

```
poly_pointer cpadd(poly pointer a, poly pointer b)
  poly pointer starta, d, lastd;
  int sum, done = FALSE;
  starta = a;
  a = a - > link;
                          Set expon field of head node to -1.
  b = b - > link;
  d = get node();
  d\rightarrow expon = -1; lastd = d;
  do {
    switch (COMPARE(a->expon, b->expon)) {
      case -1: attach(b->coef, b->expon, &lastd);
                b = b - > link;
                break;
```

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Another Padd (Continued)

Additional List Operations

```
typedef struct list_node *list_pointer;
typedef struct list_node {
    char data;
    list_pointer link;
};
```

Invert single linked lists
Concatenate two linked lists

Invert Single Linked Lists

Use two extra pointers: middle and trail.

```
list_pointer invert(list_pointer lead)
{
    list pointer middle, trail;
    middTe = NULL;
    while (lead) {
        trail = middle;
        middle = lead;
        lead = lead->link;
        middle->link = trail;
    }
    return middle;
}
```

0: null
1: lead → ...
≥2: lead → ...

Concatenate Two Lists

```
list pointer concatenate(list pointer
             ptr1, list pointer ptr2)
  list pointer temp;
  if (IS EMPTY(ptr1)) return ptr2;
  else {
    if (!IS EMPTY(ptr2)) {
      for (temp=ptr1;temp->link;temp=temp->link);
      temp->link = ptr2;
    return ptr1;
```

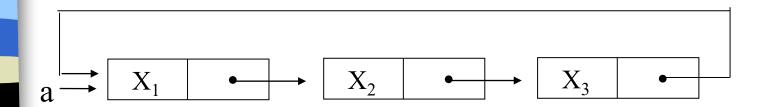
O(m) where m is # of elements in the first list

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4.5.2 Operations For Circularly Linked List

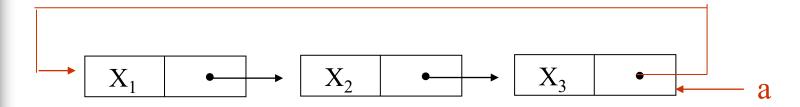
What happens when we insert a node to the front of a circular linked list?



Problem: move down the whole list.

*Figure 4.16: Example circular list (p.165)

A possible solution:

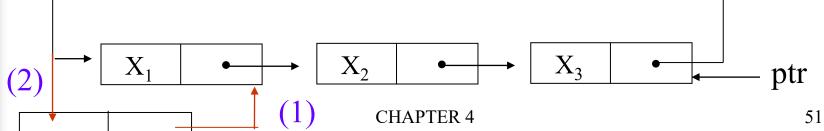


Note a pointer points to the last node.

*Figure 4.17: Pointing to the last node of a circular list (p.165)

Operations for Circular Linked Lists

```
void insert front (list pointer *ptr, list pointer
node)
    if (IS EMPTY(*ptr)) {
       *ptr= node;
       node->link = node;
    else
        node->link = (*ptr)->link;
                                      (1)
        (*ptr) -> link = node;
                                      (2)
```



Length of Linked List

```
int length(list_pointer ptr)
{
    list_pointer temp;
    int count = 0;
    if (ptr) {
        temp = ptr;
        do {
            count++;
            temp = temp->link;
        } while (temp!=ptr);
    }
    return count;
}
```

CHAPTER 4

Equivalence Relations

A relation over a set, S, is said to be an *equivalence relation* over S *iff* it is symmertric, reflexive, and transitive over S.

```
reflexive, x=x
symmetric, if x=y, then y=x
transitive, if x=y and y=z, then x=z
```

Examples

three equivalent classes {0,2,4,7,11}; {1,3,5}; {6,8,9,10}

A Rough Algorithm to Find Equivalence Classes

```
void equivalenec()
    initialize;
    while (there are more pairs) {
Phase
        read the next pair <i, j>;
        process this pair;
    initialize the output;
    do
        output a new equivalence class;
Phase
    } while (not done);
```

What kinds of data structures are adopted?

First Refinement

```
.nclude <stdio.h>
    nclude <alloc.h>
#define MAX SIZE 24
#define IS FULL(ptr) (!(ptr))
#define FATSE 0
#define TRUE 1
yoid equivalence()
        initialize seq to NULL and out to TRUE
while (there are more pairs) {
    read the next pair, <i,j>;
    put j on the seq[i] list;
    put i on the seq[j] list;
                  (i=0; i<n; i++)
if (out[i]) {
    out[i]= FALSE;
                                                                               direct equivalence
                           output this equivalence class;
                                                               Compute indirect equivalence
                                                               using transitivity
```

Lists After Pairs are input

```
[5]
                                                               [8]
                  [0]
                                   [3]
                                        [4]
                                                    [6]
0 \equiv 4
         seq
3 \equiv 1
6 \equiv 10
8 \equiv 9
                                                                                ()
7 \equiv 4
                                                                          NULL
                                                                    NULL
                       NULL
                            NULL
                                              NULL
                                                         NULL
6 \equiv 8
3 \equiv 5
2 \equiv 11
11 \equiv 0
                 NULL
                                  NULL
                                        NULL
                                                                               NULL
                                                   NULL
                                                               NULL
                 typedef struct node *node_pointer ;
                 typedef struct node {
                        int data;
                       node pointer link;
                 };
                                      CHAPTER 4
                                                                                   57
```

Final Version for Finding Equivalence Classes

```
void main(void)
  short int out[MAX SIZE];
  node pointer seq[MAX SIZE];
  node pointer x, y, top;
  int \overline{i}, j, n;
  printf ("Enter the size (<= %d) ", MAX SIZE);
  scanf("%d", &n);
  for (i=0; i< n; i++) {
      out[i] = TRUE; seq[i] = NULL;
  printf("Enter a pair of numbers (-1 -1 to quit): ");
  scanf("%d%d", &i, &j);
      Phase 1: input the equivalence pairs:
```

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```
while (i \ge 0) {
    x = (node pointer) malloc(sizeof(node));
    if (IS FULL(x))
      fprintf(stderr, "memory is full\n");
        exit(1);
    Insert x to the top of lists seq[i]
    x->data= j; x->link= seq[i]; seq[i]= x;
    if (IS FULL(x))
      fprintf(stderr, "memory is full\n");
        exit(1);
       Insert x to the top of lists seq[i]
    x->data= i; x->link= seq[j]; seq[j]= x;
    printf("Enter a pair of numbers (-1 -1 to \
          quit): ");
    scanf("%d%d", &i, &j);
```

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Phase 2: output the equivalence classes

```
for (i=0; i< n; i++) {
     if (out[i]) {
          printf("\nNew class: %5d", i);
          out[i] = FALSE;
          x = seq[i]; top = NULL;
          for (;;) {
               while (x) {
                    j = x->data;
if (out[j]) {
                         printf("%5d", j); push
out[j] = FALSE;
y = x->link; x->link = top;
                          top = x; x = y;
                    else x = x->link;
               if (!top) break;
x = seq[top->data]; top = top->link;
```

4.7 Sparse Matrices

$$\begin{bmatrix} 0 & 0 & 11 & 0 \\ 12 & 0 & 0 & 0 \\ 0 & -4 & 0 & 0 \\ 0 & 0 & 0 & -15 \end{bmatrix}$$

inadequates of sequential schemes

- (1) # of nonzero terms will vary after some matrix computation
- (2) matrix just represents intermediate results

new scheme

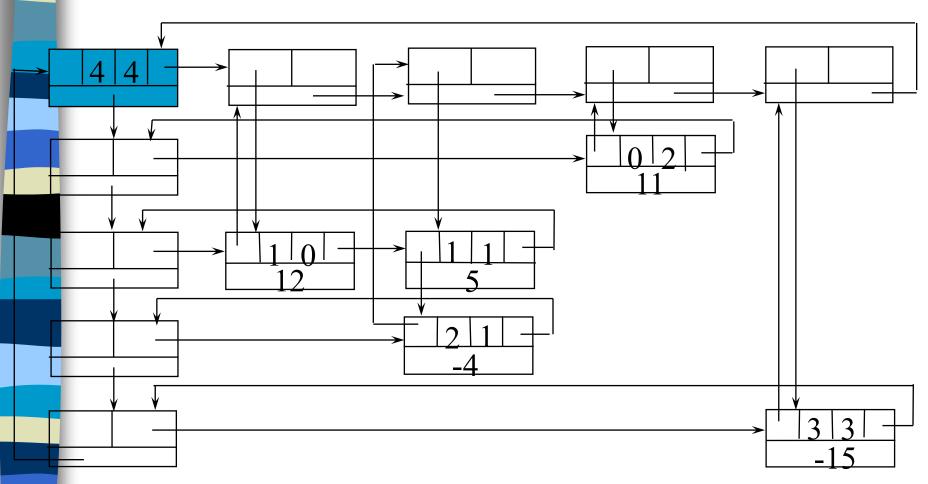
Each column (row): a circular linked list with a head node

Revisit Sparse Matrices

of head nodes = max {# of rows, # of columns} 連同一行元素 down head right 連同一列元素 head node next down col right entry row entry node value entry **i** aij aij

CHAPTER 4

Linked Representation for Matrix



```
#define MAX SIZE 50 /* size of largest matrix */
typedef enum {head, entry} tagfield;
typedef struct matrix node *matrix pointer;
typedef struct entry node {
        int row;
        int col;
        int value;
typedef struct matrix node {
        matrix pointer down;
        matrix pointer right;
        tagfield tag;
```

| | [0] | [1] | [2] |
|-----|-----|-----|-----|
| [0] | 4 | 4 | 4 |
| [1] | 0 | 2 | 11 |
| [2] | 1 | 0 | 12 |
| [3] | 2 | 1 | -4 |
| [4] | 3 | 3 | -15 |

*Figure 4.22: Sample input for sparse matrix (p.174)

Read in a Matrix

```
scanf("%d%d%d", &num rows, &num cols,
      &num terms);
 num heads =
 (num cols>num rows)? num cols : num rows;
 /* set up head node for the list of head
    nodes */
 node->u.enTry.row = num rows;
 node->u.entry.col = num-cols;
 if (!num heads) node->right = node;
 else { /* initialize the head nodes */
   for (i=0; i < num heads; i++) {
     term= new nod\overline{e}();
    hdnode[i] = temp;
                              O(\max(n,m))
     hdnode[i]->tag = head;
    hdnode[i]->right = temp;
     hdnode[i]->u.next = temp;
```

```
current row= 0; last= hdnode[0];
for (i=0; i< num terms; i++) {
 printf("Enter row, column and value:");
  scanf("%d%d%d", &row, &col, &value);
  if (row>current row) {
    last->right= The Todo [current row];
    current row= row; last=hdnode[row];
  temp = new node();
  temp->tag=entry; temp->u.entry.row=row;
  temp->u.entry.col = col;
  temp->u.entry.value = value;
  last->right = temp; /*link to row list */
  last= temp;
  /* link to column list */
  hdnode[col]->u.next->down = temp;
  hdnode[col]=>u.next = temp;
   利用next field 存放column的last node -
```

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```
/*close last row */
  last->right = hdnode[current row];
  /* close all column lists */
  for (i=0; i < num cols; i++)
    hdnode[i]->u.\overline{n}ext->down = hdnode[i];
  /* link all head nodes together */
  for (i=0; i < num heads-1; i++)
    hdnode[i] -> u.next = hdnode[i+1];
  hdnode[num heads-1]->u.next= node;
  node - > right = hdnode[0];
return node;
```

O(max{# rows, # cols}+# terms)

Write out a Matrix

```
void mwrite(matrix pointer node)
{ /* print out the matrix in row major form */
  int i;
 matrix pointer temp, head = node->right;
  printf("\n num rows = %d, num cols= %d\n",
         node->u.entry.row, node->u.entry.col);
  printf ("The matrix by row, column, and
         value: \n'); O(# rows+# terms)
  for (i=0; i< node->u.entry.row; i++) {
    for (temp=head->right; temp!=head; temp=temp->right)
      printf("%5d%5d%5d\n", temp->u.entry.row,
           temp->u.entry.col, temp->u.entry.value);
    head= head->u.next; /* next row */
```

Free the entry and head nodes by row.

Erase a Matrix

```
void merase(matrix pointer *node)
  int i, num heads;
  matrix poi\overline{n}ter x, y, head = (*node)->right; for (i\equiv 0; i<(*node)->u.entry.row; i++) {
    y=head->right;
    while (y!=head) {
       x = y; y = y - > right; free (x);
    x= head; head= head->u.next; free(x);
  y = head;
  while (y!=*node) {
    x = y; y = y->u.next; free(x);
  free(*node); *node = NULL;
  O(# rows+# cols+# terms)
```

Alternative: 利用Fig 4.14的技巧,把一列資料erase (constant time)

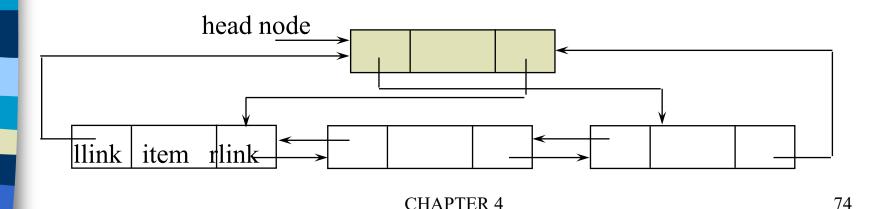
Doubly Linked List

Move in forward and backward direction.

Singly linked list (in one direction only)
How to get the preceding node during deletion or insertion?
Using 2 pointers

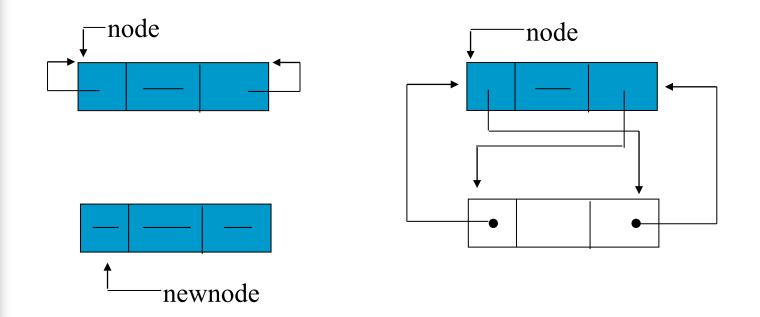
Node in doubly linked list left link field (llink) data field (item) right link field (rlink)

Doubly Linked Lists





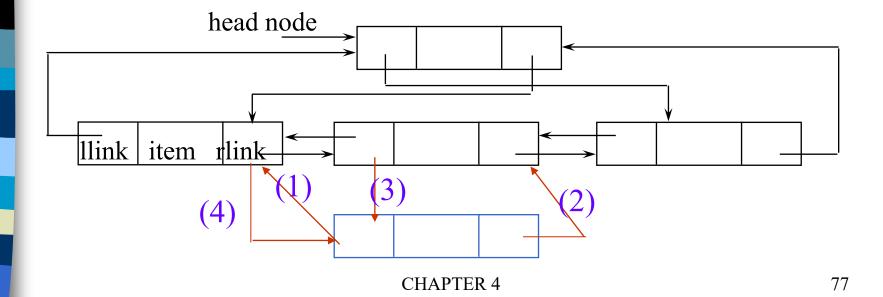
*Figure 4.24:Empty doubly linked circular list with head node (p.180)



*Figure 4.25: Insertion into an empty doubly linked circular list (p.181)

Insert

```
void dinsert(node_pointer node, node_pointer newnode)
{
    (1) newnode->llink = node;
    (2) newnode->rlink = node->rlink;
    (3) node->rlink->llink = newnode;
    (4) node->rlink = newnode;
}
```



Delete

```
void ddelete(node_pointer node, node_pointer deleted)
{
    if (node==deleted) printf("Deletion of head node not permitted.\n");
    else {
        (1) deleted->llink->rlink= deleted->rlink;
        (2) deleted->rlink->llink= deleted->llink;
        free(deleted);
}
```

