Algoritmos Taller 2: GRAFOS, complesidad Computacional, 2018/1 Programación Dinamica.

UNIVERSIDAD NACIONAL DE COLOMBIA.

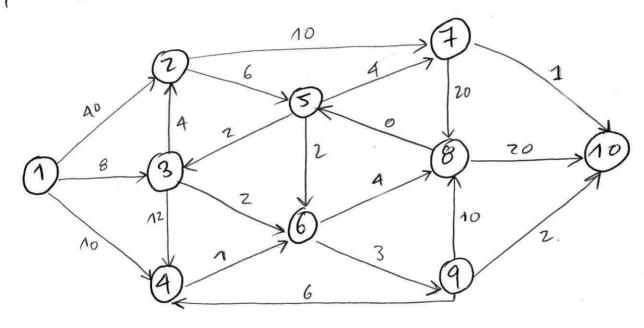
NOMBRE: JHON JAIRD MUESES Q.

Copigo: 2879355

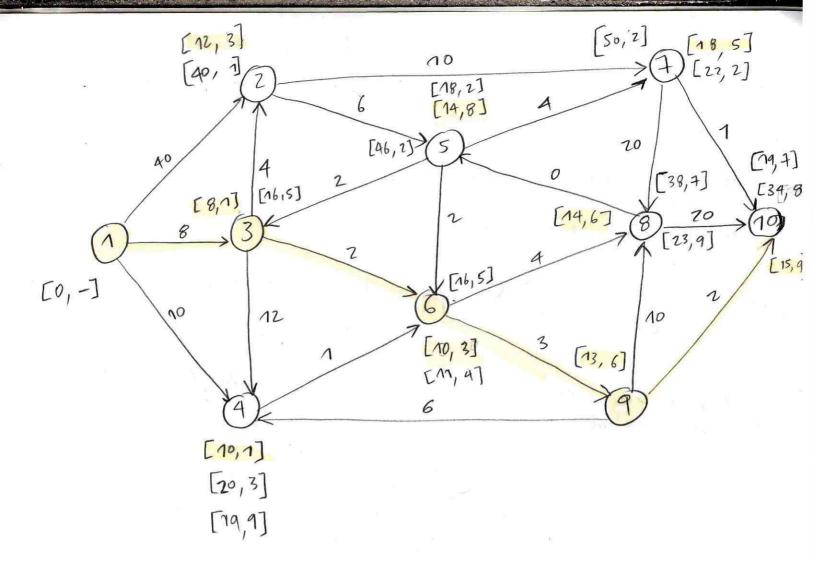
PROFESOR: GERMAN HERNANDEZ.

1) Considere el grafo de la figura 1 (solo tenga en cuenta los pesos en parintesis):

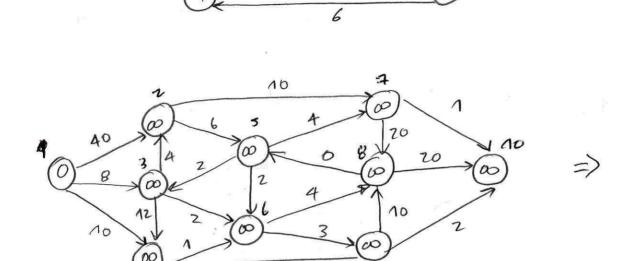
(a) Ejecute el algoritmo de Dijkstra detallando claramente los pasos ejecutados.



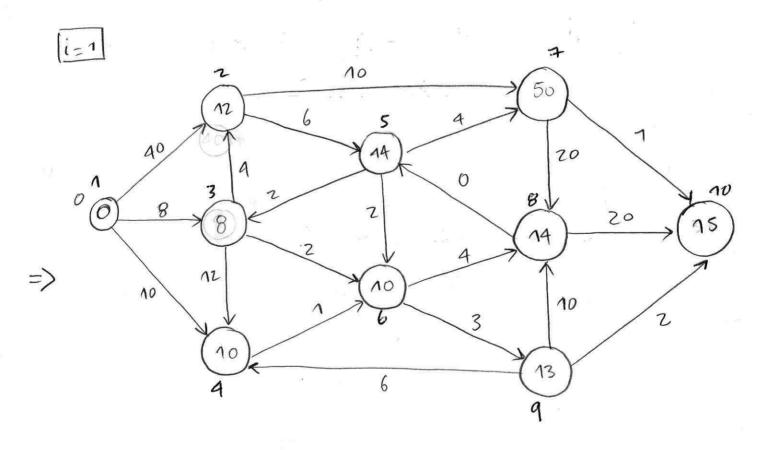
Dijkstra: Puta mas corta en ma red

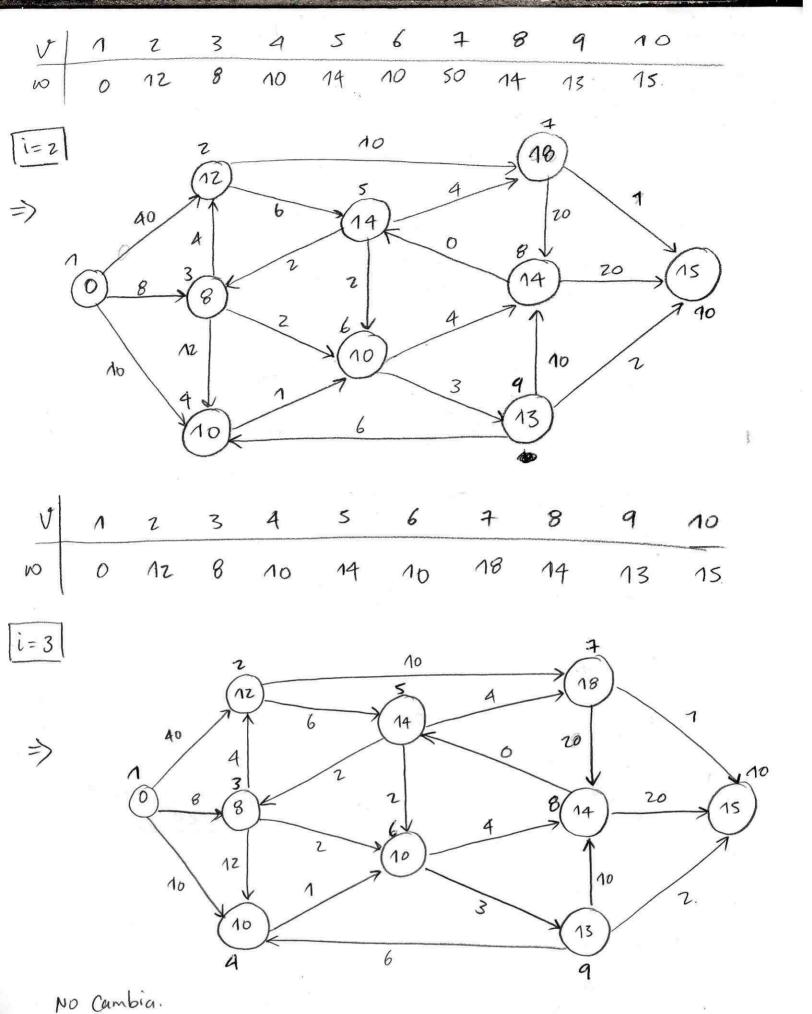


(b) Ejecute el algoritmo de Bellman-Ford detallando claramente los pasos ejecutados.



6





$$i = 4$$

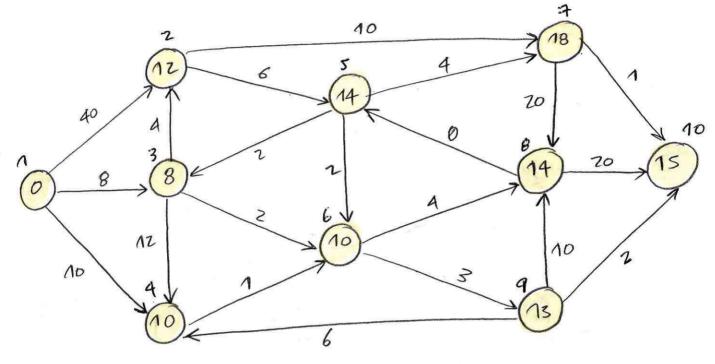
$$i = 5$$

$$i = 6$$

$$i = [6.0] - 1 = 20$$

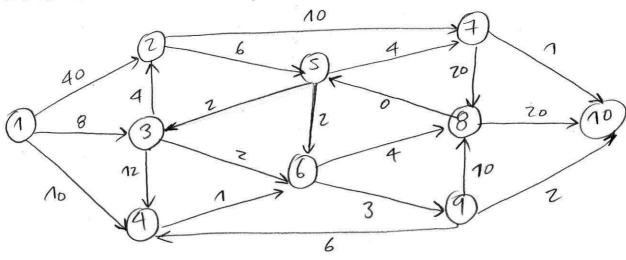
$$\begin{cases}
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i = 6
\end{cases}$$
Se obtiene el mismo grafo.

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(C) Ejecute el algorit mo de Floyd_warshal, detallando claramente los pasos ejecutados.



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