

## Mathematics of Computer Science. - M.I.T. opencourseware

### Abstract:

Random variable is the lecture about the way to calculate the random events probability with various method. If the possibility of events is stable, we can use indicator or distribution function or even Guessing strategy. Binomial distribution is when the possibility is  $1/2$  and all events are mutually independent, we can use it to count probability.

### For the English:

indicator, envelop, swap, unbiased, contrarian

### Lec 21. RANDOM VARIABLE:

DEF: a random variable  $R$  is a function  $R : S$  (sample space)  $\Rightarrow |R$   
ex) toss 3 coin //  $R = \#$  heads,  $M = 1$  if all 3 coins match |  $0$  otherwise

DEF: an Indicator (aka Bernoulli or characteristic)

$r1, v1$  is a  $r, v$  with range =  $[0, 1]$

{HHH, TTT}  $\sim$  check origin graph

{ $w \mid R(w) = x$ } is the event that  $R = X$

DEF:  $\Pr(R = x) = \sum \Pr(w)$

$\Pr(R=2) = \Pr(HHT, HTH, THH) = 3/8$ ,  $\Pr(m=1) = \Pr(HHH, TTT) = 1/4$

$\Pr(R \geq 2) = \sum \Pr(R=i) = 1/2$

for  $A \subseteq |R$ ,  $\Pr(R \in A) = \sum \Pr(R=a)$

$A = \{1, 3\}$ ,  $\Pr(R \in A) = 1/2$

$\Pr(R=2 \mid m=1) = 0$

DEF: two  $r1, v1$ 's  $R1, R2$  are indep- if  $\forall x1, x2 \in |R$

$\Pr(R1=x1 \mid R2=x2) = \Pr(R1=x1)$  or  $\Pr(R2=x2) = 0 \Rightarrow$  "no influence"

EQUIV: product form

$\forall x1, x2 \in |R \quad \Pr(R1=x1 \cap R2=x2) = \Pr(R1=x1) \cdot \Pr(R2=x2)$

$\Pr(R=2 \cap m=1) = 0 \neq \Pr(R=2) \cdot \Pr(m=1) \Rightarrow R \& M$  are not indep-

2 fair indep- 6side dice :  $D1, D2$

let  $S = D1 + D2$  let  $T = \{1 \text{ if } S = 7 \mid 0 = \text{otherwise}\}$

$\Pr(S=12 \cap D1=1) = 0 \neq \Pr(S=12) \cdot \Pr(D1=1) \Rightarrow S \& D \Rightarrow$  dependent

$\Pr(T=1 \mid D1=1) = 1/6 = \Pr(T=1)$  //  $\Pr(T=1 \mid D1=2) = 1/6 = \Pr(T=1)$

$\Pr(T=1 \mid D1=6) = 1/6 = \Pr(T=1)$  //  $\Pr(T=0 \mid D1=2) = 5/6 = \Pr(T=0)$  - 보충

DEF: given a  $r, v$   $R1$  the prob- (or point) distribution func (pdf) for  $R$  is  $f(x) = \Pr(R=x)$

DEF: the culmulative dist func  $F$  (cdf) for  $R$  is  $f(x) = \Pr(R \leq x) = \sum \Pr(R=y)$  ( $f(0)=p$ ,  $f(1)=1$ )

ex) for an indicator  $rv$   $f(0) = p$   $f(1) = 1-p$

For uniform  $r, v_1$  on  $[1, n]$   $f_n(k) = 1/n$  for  $1 \leq k \leq n$   
 $F_n(k) = k/n$

- Guessing strategy

0. envelopes contains  $y, z \in [0, n]$  where  $y < z$
1. player chooses  $x$  uniformly in  $\{1/2, 1|1/2, 2|1/2 \dots n|1/2\}$
2. player hopes  $y < x < z$
3. player opens random envelop to reveal  $r \in \{y, z\}$
4. player swaps if  $r < x$

check the origin tree node

$y/2n, z-y/2n, n-z/2n \Rightarrow n, n, s, n, s, s \Rightarrow l, w, w, w, w, l$   
 $\Pr(w) = y/2n + z-y/2n + z-y/2n + n-z/2n = n+z-y/2n = 1/2 + z-y/2n$   
 $\geq 1/2 + 1/2n = \text{PROFIT!}$

- Unbiased binomial dist

$f_n(k) = \binom{n}{k} 2^{-n} \quad n \geq 1 \quad 0 \leq k \leq n \quad \text{unbiased} = p = 1/2$

-General binomial dist

$f_{n,p}(k) = \binom{n}{k} p^k (1-p)^{n-k} \quad 0 < p < 1$   
ex)  $n$  components each fails mut- indep- with prob  $p$  ( $0 < p < 1$ )  
 $R = \# \text{ failures}$

THM:  $\Pr(R=k) = f_{n,p}(k)$  for any  $k$  ( $0 \leq k \leq n$ )

check the origin tree node

---- sample space  $= 2^n$  components

$\binom{n}{k}$  sample points have  $k$  failed components // each has prob-  $p^k (1-p)^{n-k}$

$\Pr(R=k) = \binom{n}{k} p^k (1-p)^{n-k} = f_{n,p}(k)$  - swap-possible

prob  $1 \quad 1/2 \quad 1/10 \quad 1/100 \quad 1/1k \quad 1/1m$

$\Pr(50H) \quad \quad \quad 1/2^{50}$

$f_{n,p}(\theta n) \approx 2^{[\theta \log p / \theta + (1-\theta) \log 1-p / 1-\theta]n} / \sqrt{2\pi \theta (1-\theta)n} \quad \text{for } 0 < \theta < 1$

max value at  $\theta = p$  for  $n = 100, p = 1/2$

$f_{n,p}(pn) \approx 1 / \sqrt{2\pi p(1-p)n}$

$\Pr(50H) = 1/\sqrt{50\pi} = 0.08...$

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