

Linux Kernel Data Structure

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Kernel Data Structure

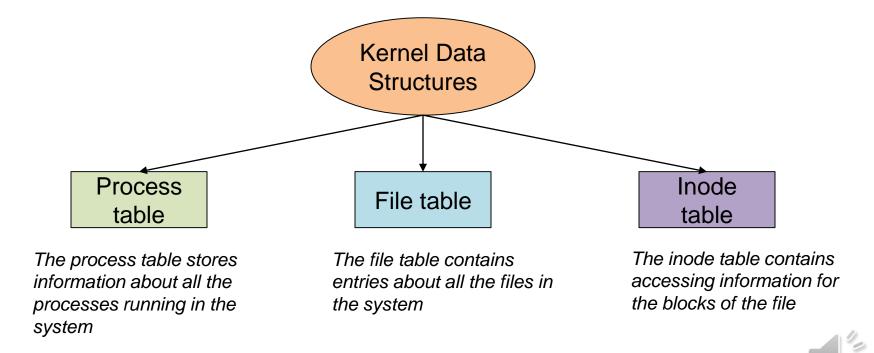
- The kernel data structures are very important as they store data about the current state of the system
 - For example, if a new process is created in the system, a kernel data structure is created that contains the detail information or state about the process
 - Kernel data structures are only accessible by the kernel and its subsystems
 - They contain data as well as pointers to other data structures





Kernel Data Structure

- The kernel data structure stores and organizes a lot of information
 - For example, it has data about which processes are running in the system, their memory requirements, files in use, etc
 - To handle all this, three important structures are used
 - These are process table, file table, and inode table





Kernel Data Structures

- This chapter introduces several built-in data structures for use in Linux kernel code
 - As with any large software project, the Linux kernel provides these generic data structures to encourage code reuse
 - Kernel developers can use these data structures





Types of Linux Kernel Data structure

- Linked list
- Red Black tree
- Hash table

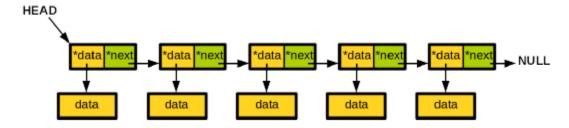




Singly linked list

- Singly linked list is a basic linked list type
- Singly linked list is a collection of nodes linked together in a sequential way
- Each node of singly linked list contains a data field and an address field which contains the reference of the next node

```
struct my_list_element{
    void *data; /* void pointer to point on a generic data */
    struct my_list_element *next; /* pointer to a next element */
};
```



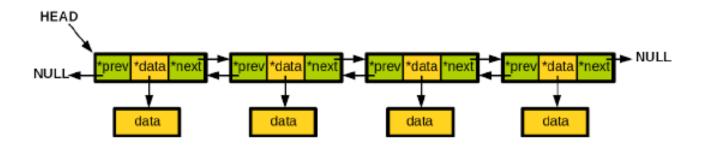
- Starts from HEAD and terminates at NULL
- Traverses forward only
- When empty, **HEAD** is **NULL**





Doubly linked list

 Each node contains three fields: two link fields (references to the previous and to the next node in the sequence of nodes) and one data field



```
struct my_list_element{
    void *data; /* void pointer to point on a generic data */
    struct my_list_element *prev; /* pointer to a previous element */
    struct my_list_element *next; /* pointer to a next element */
};
```

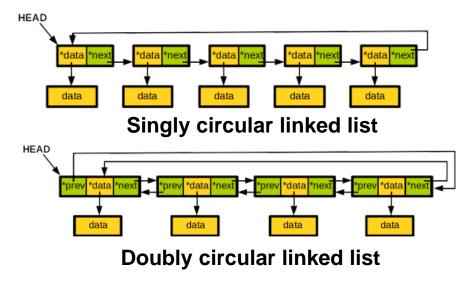
- Starts from HEAD and terminates at NULL
- Traverses forward and backward
- When empty, **HEAD** is **NULL**





Circular linked list

- Circular linked list is a linked list where all nodes are connected to form a circle
 - There is no NULL at the end
 - A circular linked list can be a singly circular linked list or doubly circular linked list

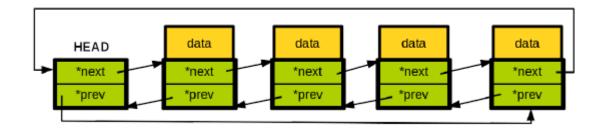


- Starts from HEAD and terminates at HEAD
- When empty, **HEAD** is **NULL**





Linux kernel linked list

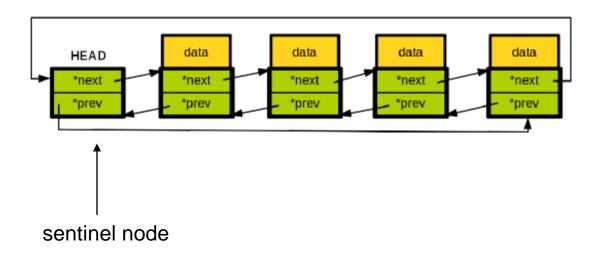


- Starts from HEAD and terminates at HEAD
- When empty, HEAD is not NULL
 - Prev and next of HEAD points HEAD
 - HEAD is a sentinel node
- Easy to insert a new element at the end of a list
- There is no exceptional case to handle NULL





- Linux kernel linked list is a circular doubly linked list
- Two differences from the typical design
 - Embedding a linked list node in the structure
 - Using a sentinel node as a list header
- linux/include/linux/list.h







- struct list_head is the key data structure
- list_head is embedded in the data structure
- Start of a list is also list_head my_car_list → sentinel node





- Getting a data from list_head
 - use list_entry (ptr, type, member)

struct car *amazing_car = list_entry(car_list_ptr, struct car, list)

list_entry – get the struct for this entry

- @ptr: the &struct list_head pointer
- @type: the type of the struct this is embedded in
- @member: the name of the list_head within the struct
- #define list_entry(ptr, type, member) container_of(ptr, type, member)





Linux kernel linked list APIs

```
/* Insert a new entry after the specified head */
void list_add(struct list_head *new, struct list_head *head);
/* Insert a new entry before the specified head */
void list_add_tail(struct list_head *new, struct list_head *head);
/* Delete a list entry
* NOTE: You still have to take care of the memory deallocation if
needed */
void list_del(struct list_head *entry);
/* Delete from one list and add as another's head */
void list_move(struct list_head *list, struct list_head *head);
/* Delete from one list and add as another's tail */
void list move tail(struct list head *list, struct list head *head);
/* Join two lists (merge a list to the specified head) */
void list_splice(const struct list_head *list, struct list_head *head);
```





Linux kernel linked list APIs

```
/**
* list for each - iterate over a list
* @pos: the &struct list_head to use as a loop cursor.
* @head: the head for your list.
#define list_for_each(pos, head) \
     for (pos = (head)->next; pos != (head); pos = pos->next)
/**
* list_for_each_entry - iterate over list of given type
* @pos: the type * to use as a loop cursor.
* @head: the head for your list.
* @member: the name of the list_head within the struct.
#define list_for_each_entry(pos, head, member) \
     for (pos = list_first_entry(head, typeof(*pos), member); \
         &pos->member != (head); \
         pos = list_next_entry(pos, member))
```





An example of Linux kernel linked list

```
#include #include
```

```
int __init hello_module_init(void)
   struct_exmaple();
   printk("module init\n");
   return 0;
void __exit hello_module_cleanup(void)
   printk("Bye Module\n");
module_init(hello_module_init);
module_exit(hello_module_cleanup);
```





Linux linked list

Linux linked list example

```
struct my_node {
                                                                     struct list head list;
                                                                     int data:
void struct exmaple(void)
             struct list_head my_list;
                                                                 #define list for each(pos, head) \
                                                                      for (pos = (head)->next; pos != (head); pos = pos->next)
            /* initialize list */
                                                                 #define list for each entry(pos, head, member) \
             INIT LIST HEAD(&my list);
                                                                      for (pos = list_first_entry(head, typeof(*pos), member); \
                                                                         &pos->member != (head); \
            /* add list element */
                                                                         pos = list_next_entry(pos, member))
            int i:
            for (i = 0; i < 10; i++)
                         struct my node *new = kmalloc(sizeof(struct my node), GFP KERNEL);
                         new->data = i:
                         list_add(&new->list, &my_list);
            struct my_node *current_node; /* This will point on the actual data structures during the iteration */
             sturct list_head *p; /* Temporary variable needed to iterate; */
             list_for_each(p, &my_list){
                         current node = list entry(p, struct my node, list);
                         printk("current value : %d\n", current_node->data);
             list_for_each_entry(current_node, &my_list, list){
                         printk("current value : %d\n", current node->data):
```



Linux linked list

```
/* iterate over a list reversely*/
list_for_each_entry_reverse(current_node, &my_list, list){
           printk("current value : %d\n", current_node->data);
/* delete list element */
list_for_each_entry(current_node, &my_list, list){
          if(current_node->data == 2){
                     printk("current node value : %d\n", current_node->data);
                     list_del(&current_node->list);
                     kfree(current node);
/* iterate over a list */
list_for_each_entry(current_node, &my_list, list){
          printk("current value : %d\n", current_node->data);
```



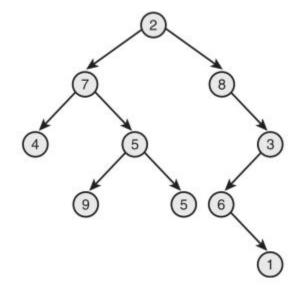


- Usage of linked list in the kernel
 - Kernel code makes extensive use of linked lists
 - a list of threads under the same parent PID
 - a list of superblocks of a file system
 - and many more





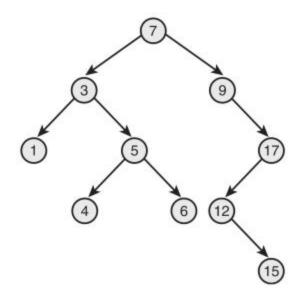
- Tree basics : binary tree
- Properties
 - Nodes have zero, one, or two children
 - Root has no parent, other nodes have one







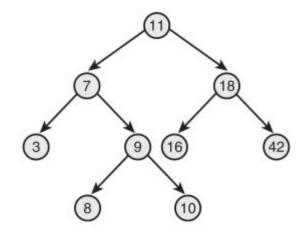
- Tree basics : binary search tree
- Properties
 - Left children < parent
 - Right children > parent
 - Search and ordered traversal are efficient







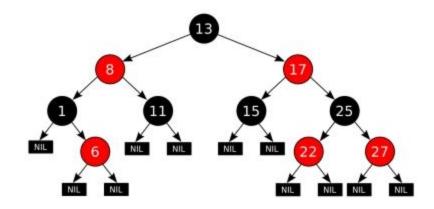
- Tree basics : balanced binary search tree
- Properties
 - Depth of all leaves differs by at most one







- Tree basics : red-black tree
- Self-balancing binary search tree
- Properties
 - A type of self-balancing binary search tree
 - Nodes: red or black
 - Leaves: black, no data







Self-balancing binary search tree

```
/* Rbtree node, which is embedded to your data structure like

* list_head and hlist node */

struct rb_node {

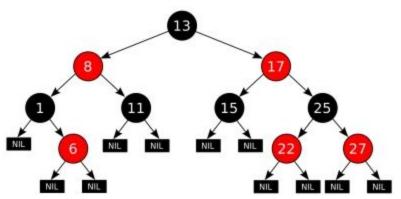
    unsigned long __rb_parent_color;
    struct rb_node *rb_right;
    struct rb_node *rb_left;

};

/* Root of a rbtree */

struct rb_root {

    struct rb_node *rb_node;
};
```







red-black tree APIs

```
/* Find logical next and previous nodes in a tree */
struct rb_node *rb_next(const struct rb_node *);
struct rb_node *rb_prev(const struct rb_node *);
struct rb_node *rb_first(const struct rb_root *);
struct rb_node *rb_last(const struct rb_root *);
/* Insert a new node under a parent connected via rb link */
void rb_link_node(struct rb_node *node, struct rb_node *parent,
                     struct rb_node **rb_link);
/* Re-balance an rbtree after inserting a node if necessary */
void rb_insert_color(struct rb_node *, struct rb_root *);
/* Delete a node */
void rb_erase(struct rb_node *, struct rb_root *);
```





Red-Black Tree example

```
#include linux/kernel.h>
#include linux/module.h>
#include linux/init.h>
#include linux/rbtree.h>
#include linux/slab.h> // for kmalloc
#define FALSE 0
#define TRUE 1
struct my_node {
          struct rb node node:
          int key;
          int value;
};
```

```
int init hello module init(void)
          RB_exmaple();
          printk("module init\n");
          return 0;
void __exit hello_module_cleanup(void)
          printk("Bye Module\n");
module_init(hello_module_init);
module_exit(hello_module_cleanup);
```





```
void RB_exmaple(void)
           struct rb_root root_node = RB_ROOT;
           int i;
         /* rb node create and insert */
           for(i=0;i<20;i++)
                       struct my_node *new_node = kmalloc(sizeof(struct my_type),GFP_KERNEL);
                       if(!new node)
                                  return NULL:
                       new node->value = i*10;
                      new_node->key = i;
                       ret = rb insert color(new node, &root node);
           /* rb tree traversal using iterator */
           struct rb_node *iter_node;
           for (iter_node = rb_first(&root_node); iter_node; iter_node = rb_next(iter_node))
                      printk("(key,value)=(%d.%d)\n", \
                                  rb_entry(iter_node, struct my_node, node)->key,
                                  rb_entry(iter_node, struct my_node, node)->value);
           /* rb tree delete node */
           rb erase(new node, &root node);
```

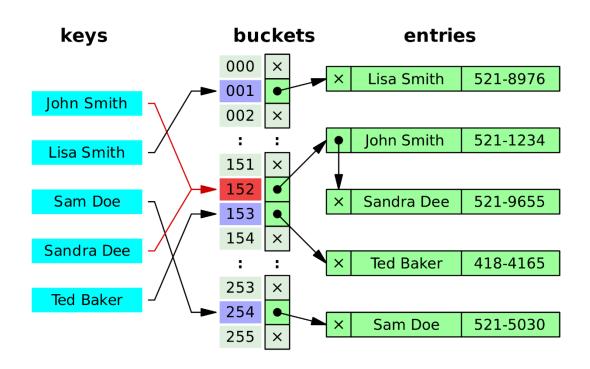


- Usage of Red-Black tree in the kernel
 - Completely Fair Scheduling (CFS)
 - Default task scheduler in Linux
 - Each task has vruntime, which presents how much time a task has run
 - CFS always picks a process with the smallest vruntime for fairness
 - Per-task vruntime structure is maintained in a rbtree





- A simple fixed-size chained hash table
 - The size of bucket array is fixed at initialization as a 2^N
 - Each bucket has a singly linked list or doubly linked list to resolve hash collision







Hash table APIs

```
* Define a hashtable with 2^bits buckets
#define DEFINE_HASHTABLE(name, bits) ...
/**
* hash_init - initialize a hash table
* @hashtable: hashtable to be initialized
#define hash_init(hashtable) ...
/**
* hash_add - add an object to a hashtable
* @hashtable: hashtable to add to
* @node: the &struct hlist_node of the object to be added
* @key: the key of the object to be added
*/
#define hash_add(hashtable, node, key) ...
```





Hash table APIs

```
* hash for each - iterate over a hashtable
* @name: hashtable to iterate
* @bkt: integer to use as bucket loop cursor
* @obj: the type * to use as a loop cursor for each entry
* @member: the name of the hlist_node within the struct
*/
#define hash_for_each(name, bkt, obj, member) ...
/**
* hash del - remove an object from a hashtable
* @node: &struct hlist node of the object to remove
*/
void hash_del(struct hlist_node *node);
```





Hash table example

```
#include #include
```

```
int __init hello_module_init(void)
          hash_exmaple();
          printk("module init\n");
          return 0;
void exit hello module cleanup(void)
          printk("Bye Module\n");
module_init(hello_module_init);
module_exit(hello_module_cleanup);
```





```
void hash_exmaple(void)
           DEFINE_HASHTABLE(my_hash, MY_HASH_BITS);
           hash_init(my_hash);
           /* (key,value) insert */
           int i;
           for(i=0;i<10;i++){
                      struct my_node *new = kmalloc(sizeof(struct my_node), GFP_KERNEL);
                      new->value = i * 10;
                      new->key = i;
                      memset(&new->hnode, 0, sizeof(struct hlist_node));
                      hash_add(my_hash,&new->hnode,new->key);
           /* You can complete hash_for_each() function!!! */
           hash_for_each(....){
           /* You can complete hash_del() function!!! */
           hash_del(....){
```



Usage of hash table in the kernel

- Hugepage
 - finds physically consecutive 4KB pages
 - remaps consecutive 4KB pages to a 2MB page (huge page)
 - saves TLB entries and improves memory access performance by reduing TLB miss
 - maintains per-process memory structure, struct mm_struct



