

A Note on Floor Computation

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1 Introduction

This note is an informal discussion of an algorithm I developed (and its C++ implementation) to find the set of pairs $(\xi, \zeta) \in \mathbb{R}^2$ such that the equality

$$\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$$

holds for all n in a fixed interval of integers, where x, y are arbitrary fixed real numbers. The goal is to compute $\lfloor nx + y \rfloor$ for some fixed x, y , but since x, y may be very complicated numbers, direct computation might be infeasible. Therefore, we want to replace x, y with some ξ, ζ specifically chosen for fast computation of $\lfloor n\xi + \zeta \rfloor$. In practice, preferred choices of ξ, ζ are often of the form $\xi = \frac{m}{2^k}$ and $\zeta = \frac{s}{2^k}$ for some integers k, m, s , so that

$$\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor = \left\lfloor \frac{nm + s}{2^k} \right\rfloor$$

can be computed with a multiplication, an addition, and a bit-shift of integers.

One of the motivating examples is the problem of fast computation of integer division $\lfloor \frac{n}{d} \rfloor$ by a constant divisor d , which is a special case of the above problem with $x = \frac{1}{d}$ and $y = 0$. In modern CPUs, integer division typically has much worse performance characteristics compared to other basic integer arithmetic operations. For this reason, techniques for replacing divisions with multiply-shift or multiply-add-shift operations have been explored at least since the 1990s, and some of these techniques are already widely applied in modern optimizing compilers. Some well-known earlier works on these techniques include [1], [2] and [3]. In these works, the time complexity of searching for feasible k, m, s is polynomial in $\log N$, where N is the size of the range of n , so it can be done very quickly even for extremely large N .

On the other hand, there are situations where we may want to compute $\lfloor nx + y \rfloor$ in forms more general than $\lfloor \frac{n}{d} \rfloor$. For instance, many of modern floating-point formatting algorithms, such as Ryū [4] and Dragonbox [5], are based on the idea of computing $\left\lfloor \frac{np}{q} \right\rfloor$ (i.e., $x = \frac{p}{q}$ and $y = 0$) through the identity $\left\lfloor \frac{np}{q} \right\rfloor = \left\lfloor \frac{nm}{2^k} \right\rfloor$, where $\frac{p}{q}$ is a fixed rational number with a potentially very large denominator. The algorithms used for searching for k, m again have polynomial-in- $\log N$ time complexity. In these scenarios, $\log N$ is comparable to the word size (e.g., 32, 64 or 128), so having such time complexity is crucial. If, for example, it were $O(N)$, the search would take an unreasonable amount of time.

There is also prior art that considers the case when $y \neq 0$: a work by Neri and Schneider [6] addresses the computation of $\lfloor \frac{na+b}{d} \rfloor$ (i.e., $x = \frac{a}{d}$ and $y = \frac{b}{d}$) for some fixed integers a, b, d , in the context of calendar algorithms. Note that we can still apply the aforementioned works on integer division in this scenario by separately computing $na + b$ and its division by d . However, in that case, we must perform two integer multiplications: one for computing na and another for computing $m(na + b)$ (where $\xi = \frac{m}{2^k}$). In contrast, the work by Neri and Schneider describes a search algorithm for finding k, m, s such that $\lfloor \frac{na+b}{d} \rfloor = \lfloor \frac{nm+s}{2^k} \rfloor$ holds for all n in the given range, allowing us to perform the computation with only one integer multiplication. However, the time complexity of this search algorithm is essentially $O(d)$, making it impractical for cases where x, y are not rational numbers with small denominators.

Note that the idea of computing $\lfloor nx + y \rfloor$ by replacing (x, y) with a pair $(\xi, \zeta) = (\frac{m}{2^k}, \frac{s}{2^k})$ has been known since the early days. The contribution of [6] seems to be in proving a sufficient condition for (ξ, ζ) such that the trick works and developing an algorithm to find k, m, s as a consequence, rather than proposing the idea itself. Indeed, it seems obvious that $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ must hold if (ξ, ζ) is a fine enough binary approximation of (x, y) , and it is not hard to find actual examples where this idea is exploited. For instance, the implementation of the Schubfach algorithm [7] relies on the same technique to compute $\lfloor n \log_{10} 2 - \log_{10} \frac{4}{3} \rfloor$ as one of its substeps.

However, somewhat counterintuitively, the best choice of $(\xi, \zeta) = (\frac{m}{2^k}, \frac{s}{2^k})$ to replace (x, y) , that is, the one that allows the equality $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ to hold over the widest range of n , or the one using the fewest bits for a given range of n , is *not* the one closest to the true value of (x, y) among candidates with the same number of bits. It can actually be quite far from the closest approximation, because the errors in the approximations of x and of y do not simply add; sometimes they can cancel each other out. This was already evident in the work by Neri and Schneider [6], but it remains unclear if there is a simple intuitive way of guessing what the optimal choice for m and s would then look like. It seems that until the algorithm presented in this note was developed, there was no known fast algorithm for finding k, m, s in cases where the algorithm in [6] cannot be applied, i.e., when x, y are not rational numbers with small denominators, even when the range size N is small.¹

In this note, we present a fast algorithm for finding a minimal description of the *precise set* of $(\xi, \zeta) \in \mathbb{R}^2$ satisfying $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all n in a given range of size N , with a time complexity that is (probably) polynomial in $\log N$.² Once we have such a description, it is easy to select a point of the form $(\xi, \zeta) = (\frac{m}{2^k}, \frac{s}{2^k})$ in that set. Since the algorithm finds the *precise set*, in effect we are describing a *necessary and sufficient condition* for (ξ, ζ) to satisfy the constraint $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$. Although in practice it is often enough to have only a *sufficient* condition on (ξ, ζ) , sometimes we want to ensure that our choice is absolutely the best. Our algorithm allows us to make such a guarantee because what it computes is not only a sufficient condition but also a necessary and sufficient condition.

Before delving into the details, let me point out that every algorithm described in this note can be implemented using arbitrary-precision exact rational arithmetic, even with that x, y are only assumed to be arbitrary *real numbers* (represented by their *continued fraction expansions*; see Section 3 for details). In particular, there are no floating-point operations involved, and there is no concern about rounding errors leading to imprecise results.

Even though arbitrary-precision exact rational arithmetic might run slower than alternatives (e.g. floating-point arithmetic if it can be leveraged at all), squeezing out the last bit of performance is not our primary goal, as the computation of (ξ, ζ) is done as a preparation

¹Of course, if N is small, we can exhaustively verify whether given choices of k, m, s work. However, that does not mean we know how to find the best possible k, m, s .

²We do not formally prove this statement, but it seems fairly obvious from how the algorithm works.

step (that needs to be done only once) for computing $\lfloor nx + y \rfloor$. The only performance concern is whether the implementation can return the result in a reasonable amount of time, and the proposed algorithm is more than sufficient in that regard. Indeed, as mentioned earlier, the time complexity is (probably) only polynomial in $\log N$, and more importantly than the theoretical asymptotic bound, the C++ reference implementation has no problem completing its run almost instantaneously on my laptop I bought in 2018, even in the non-optimized debug build, for $N \sim 2^{64}$ and almost any fairly complicated rational or irrational x, y .

I also note that, although the theory behind the main algorithm we develop does not care heavily about rationality of x and y , an actual implementation of the algorithm that accepts possibly irrational inputs is necessarily more complicated than one that only accepts rational inputs. The reference implementation does support irrational inputs, but I expect a much simpler implementation could be achieved if we restricts the inputs to rational numbers. Readers who are only interested in the rational case can skip the followings: Theorem 7.6, Section 8, Section 10.3, Section 12 and Section 13.

Now, let us begin the formal description of the proposed algorithm. Throughout this note, we let

$$[n_{\min} : n_{\max}] := [n_{\min}, n_{\max}] \cap \mathbb{Z}$$

for any $n_{\min}, n_{\max} \in \mathbb{Z}$. Recall that our goal is to find a concise description of the set of all $(\xi, \zeta) \in \mathbb{R}^2$ satisfying $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$. First, observe that having this equality for all $n \in [n_{\min} : n_{\max}]$ is equivalent to the following inequalities:

$$\begin{aligned} \max_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{>0}} \frac{\lfloor nx + y \rfloor - \zeta}{n} \leq \xi < \min_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{>0}} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n}, \\ \max_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{<0}} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} < \xi \leq \min_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{<0}} \frac{\lfloor nx + y \rfloor - \zeta}{n}, \end{aligned}$$

possibly along with $\lfloor y \rfloor = \lfloor \zeta \rfloor$ if $0 \in [n_{\min} : n_{\max}]$. Thus, the problem boils down to the task of computing the maximum and the minimum in the above inequalities. Of course, we will not obtain closed-form expressions for them, and instead our goal is to derive an efficient algorithm for computing these extrema.

A paper by Drane et al [8] provides useful geometric intuition for our problem. Note that the equality $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ can be rewritten as the inequality

$$n(x - \xi) + (y - \zeta) \leq (nx + y) - \lfloor nx + y \rfloor < n(x - \xi) + 1 + (y - \zeta).$$

Here, the left-hand side and right-hand side are affine functions of n with the same slope $x - \xi$, so what we want to have is to ensure that the “sawtooth function” in the middle lies between the two lines, as illustrated in Figure 1.

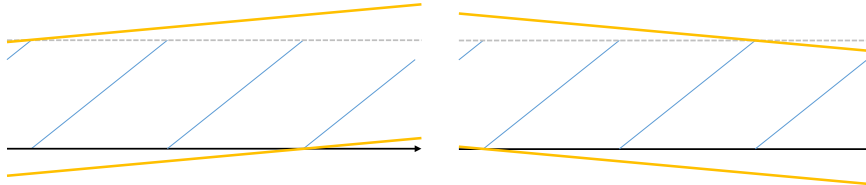


Figure 1: Left: when $x > \xi$, Right: when $x < \xi$

Of course, if n were allowed to vary continuously over all real numbers between n_{\min} and n_{\max} , then except for fairly special cases, placing the sawtooth function between the two lines is simply impossible unless $x = \xi$, since the height of the sawtooth is precisely the vertical distance between the lines. However, when n is restricted to integers, the sawtooth function may not actually reach its top or bottom corners, leaving some room for the lines to pass through. This observation leads to the insight that the maximum and the minimum of the fractional part $(nx + y) - \lfloor nx + y \rfloor$ will play a significant role, as they describe the size of these gaps. Indeed, computing the extremizers of $(nx + y) - \lfloor nx + y \rfloor$ is one of the first steps in the final algorithm we present in Section 11.

Thus, after developing some fundamental theoretical background in Section 2, Section 3 and Section 4, we will first describe an algorithm for computing the extremizers of $(nx + y) - \lfloor nx + y \rfloor$ in Section 5.

2 Best rational approximations from below and above

Before moving into the general case, let us first examine a special case when $y = 0$ and $n_{\min} = 1$, that is, we are concerned with the extremizers of $nx - \lfloor nx \rfloor$ when $n = 1, \dots, n_{\max}$. This special case in fact is the most fundamental building block for solving the general case. First, we reformulate this problem with the following terminology.

Definition 2.1 (Best rational approximations from below/above).

Let x be a real number and $\frac{p}{q}$ a rational number in its reduced form. Then we say $\frac{p}{q}$ is:

1. a *best rational approximation* of x from below (above, resp.) in the weak sense, if $\frac{p}{q} \leq x$ ($\frac{p}{q} \geq x$, resp.) holds and for any rational number $\frac{a}{b}$ with $0 < b \leq q$ and $\frac{a}{b} \leq x$ ($\frac{a}{b} \geq x$, resp.), $\frac{a}{b} \leq \frac{p}{q} \leq x$ ($\frac{a}{b} \geq \frac{p}{q} \geq x$, resp.) holds,
2. a *best rational approximation* of x from below (above, resp.) in the strong sense, if $\frac{p}{q} \leq x$ ($\frac{p}{q} \geq x$, resp.) holds and for any rational number $\frac{a}{b}$ with $0 < b \leq q$ and $\frac{a}{b} \leq x$ ($\frac{a}{b} \geq x$, resp.), $bx - a \geq qx - p \geq 0$ ($a - bx \geq p - qx \geq 0$, resp.) holds.

Remark 2.2.

Best rational approximation is a standard terminology in mathematical literatures, but they usually use this term to denote a related but different concept: $\frac{p}{q}$ is said to be a best rational approximation of x if it is a best approximation *regardless of the direction* of approximation, from below or above. For example, $\frac{1}{3}$ is a best rational approximation from below of $\frac{3}{7}$ (in the weak sense), but it is not a best rational approximation in the usual, non-directional sense, because $\frac{1}{2}$ is a better approximation with a strictly less denominator. But this concept of non-directional best rational approximations is quite irrelevant to our application, so any usage of the term *best rational approximation* in this note always means the directional ones, either from below or above.

Of course, the terminology is set in the way that “strong” implies “weak”.

Lemma 2.3.

Let x be a real number. If a rational number $\frac{p}{q}$ in its reduced form is a best rational approximation of x from below (above, resp.) in the strong sense, then it is a best rational approximation of x from below (above, resp.) in the weak sense as well.

Proof. We only show the statement about best rational approximations from below. The corresponding statement for best rational approximations from above can be done in a similar manner.

Suppose $\frac{p}{q}$ is a best rational approximation of x from below in the strong sense. It is then enough to show that for any rational number $\frac{a}{b}$ with $0 < b \leq q$ and $\frac{a}{b} \leq x$, we must have $\frac{a}{b} \leq \frac{p}{q}$. This follows directly from the assumptions:

$$x - \frac{a}{b} = \frac{bx - a}{b} \geq \frac{qx - p}{b} \geq \frac{qx - p}{q} = x - \frac{p}{q},$$

thus we conclude $\frac{a}{b} \leq \frac{p}{q}$. \square

Actually, it turns out that the converse is also true, that is, two concepts are in fact equivalent to each other. This is a very useful fact but unfortunately not entirely trivial as far as I know. We will prove it in Section 4 after developing a sufficient amount of the theory of *continued fractions* in Section 3 and Section 4. Before we prove the equivalence, we will distinguish the weak and the strong notions of best rational approximations from below/above.

Clearly, if $\frac{p}{q}$ is a best rational approximation of x from below, either in the weak or the strong sense, then we must have $p = \lfloor qx \rfloor$, because $p > \lfloor qx \rfloor$ implies $\frac{p}{q} > x$ and $p < \lfloor qx \rfloor$ implies $\frac{p}{q} < \frac{\lfloor qx \rfloor + 1}{q} \leq x$, which both contradict to the assumption. Similarly, if $\frac{p}{q}$ is a best rational approximation of x from above, then we must have $p = \lceil qx \rceil$. From this simple observation, we obtain the following result.

Theorem 2.4 (Extremizers of $nx - \lfloor nx \rfloor$).

Let x be a real number and n_{\max} a positive integer.

1. If x is rational and its reduced form $\frac{p}{q}$ satisfies $q \leq n_{\max}$, then

$$\operatorname{argmin}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor) = \{kq : k \in \mathbb{Z}_{>0}, kq \leq n_{\max}\}$$

and

$$\operatorname{argmax}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor) = \{kq - b : k \in \mathbb{Z}_{>0}, kq - b \leq n_{\max}\},$$

where $b \in [0: q - 1]$ is the modular inverse of p with respect to q .

2. If x is either irrational or is rational whose reduced denominator is strictly larger than n_{\max} , then

$$\operatorname{argmin}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor) = \{q_*\}$$

and

$$\operatorname{argmax}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor) = \{q^*\},$$

where $\frac{p_*}{q_*}$ and $\frac{p^*}{q^*}$ are best rational approximations in their reduced forms of x from below and above, respectively, in the strong sense, with the largest (reduced) denominators $q_*, q^* \leq n_{\max}$.

Proof. 1. Clearly,

$$nx - \lfloor nx \rfloor = \frac{1}{q} \left(np - \left\lfloor \frac{np}{q} \right\rfloor q \right) = \frac{(np \bmod q)}{q},$$

thus the minimizers of $nx - \lfloor nx \rfloor$ are precisely the multipliers of q . Similarly, the maximizers are precisely the n 's such that the remainder of np divided by q is $q - 1$. Those n 's are precisely the ones with $n \equiv -b \pmod{q}$, so we arrive at the desired conclusion.

2. Take any $n_0 \in \operatorname{argmin}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor)$. We claim that $\frac{\lfloor n_0 x \rfloor}{n_0}$ is a best rational approximation of x from below in the strong sense and is in its reduced form. Let $\frac{a}{b}$ be any rational number in its reduced form such that $\frac{a}{b} \leq x$ and $b \leq q$. Then since $a \leq bx$ implies $a \leq \lfloor bx \rfloor$, we have

$$bx - a \geq bx - \lfloor bx \rfloor \geq n_0 x - \lfloor n_0 x \rfloor \geq 0.$$

since n_0 is a minimizer of $nx - \lfloor nx \rfloor$. In particular, if $\frac{a}{b}$ is the reduced form of $\frac{\lfloor n_0 x \rfloor}{n_0}$, then $n_0 x - \lfloor n_0 x \rfloor = k(bx - a)$ for some positive integer k , so we must have either $n_0 x = \lfloor n_0 x \rfloor$ or $k = 1$. The first case is impossible by the assumption on x , and the claim follows.

Next, we claim $n_0 = q_*$, the largest denominator of all best rational approximations of x from below in the strong sense whose denominators are no more than n_0 . Indeed, since $\frac{p_*}{q_*}$ is a best rational approximation of x from below in the strong sense and since $n_0 \leq q_*$ by definition of q_* , we must have

$$n_0 x - \lfloor n_0 x \rfloor \geq q_* x - \lfloor q_* x \rfloor = q_* x - p_* \geq 0.$$

However, since n_0 is a minimizer of $nx - \lfloor nx \rfloor$, the first inequality in the above should be an equality, thus we obtain

$$(q_* - n_0)x = \lfloor q_* x \rfloor - \lfloor n_0 x \rfloor.$$

Since the right-hand side is an integer, the assumption on x implies $q_* = n_0$.

Next, we show the statement about the maximizer. First, note that since nx is never an integer, we have $\lceil nx \rceil = \lfloor nx \rfloor + 1$ for all $n \in [1: n_{\max}]$. Hence,

$$\begin{aligned} \operatorname{argmax}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor) &= \operatorname{argmin}_{n \in [1: n_{\max}]} (\lfloor nx \rfloor - nx) = \operatorname{argmin}_{n \in [1: n_{\max}]} (\lceil nx \rceil - nx) \\ &= \operatorname{argmin}_{n \in [1: n_{\max}]} (n(-x) - \lfloor n(-x) \rfloor), \end{aligned}$$

thus by what we have just shown, $\operatorname{argmax}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor)$ must be the singleton set whose unique element is the denominator of the best rational approximation $\frac{p}{q}$ of $-x$ from below in the strong sense with the largest denominator $q \leq n_{\max}$. Then it easily follows from the definition that $\frac{-p}{q}$ is a best rational approximation of x from above in the strong sense with the largest denominator $q \leq n_{\max}$, so we conclude that $q = q^*$ is the unique element in $\operatorname{argmax}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor)$. \square

Note that when $x \notin \mathbb{Z}$, b appearing in the first case is precisely the largest denominator of all best rational approximations of x from below in the strong sense with denominators at most $q - 1$, because clearly it is the unique minimizer of $nx - \lfloor nx \rfloor$ for $n \in [1: q - 1]$. (Similarly, $q - b$ is the largest denominator of all best rational approximations of x from above in the strong sense with denominators at most $q - 1$.) Hence, in any case, the problem of finding the extremizers of $nx - \lfloor nx \rfloor$ is equivalent to the problem of finding best rational approximations of x from below/above in the strong sense.

In Section 4, we will derive an efficient algorithm for computing the best rational approximations of x from below/above with the largest denominators bounded by given n_{\max} , using the theory of *continued fractions*.

3 Continued fractions

The problem of finding best rational approximations of x from below/above can be solved using the *continued fraction expansion* of x , so let us take a bit of detour and introduce some basic theory of continued fractions first. Many of what follows in this section and the next section are directly copied and only slightly modified from my previous paper on a floating-point formatting algorithm, *Dragonbox* [5], though there are some added materials too. Those from [5] are reproduced here to make this note more self-contained. Contents in this section specifically is standard in the literature, and can be found, for example in [9].

A *continued fraction* means either a finite or infinite sequence of the form

$$a_0, a_0 + \frac{b_1}{a_1}, a_0 + \frac{b_1}{a_1 + \frac{b_2}{a_2}}, a_0 + \frac{b_1}{a_1 + \frac{b_2}{a_2 + \frac{b_3}{a_3}}}, \dots$$

The i th term appearing in the sequence above is called the i th *convergent* of the continued fraction.

Lemma 3.1.

For each $i \geq 0$, inductively define

$$\begin{cases} p_i = b_i p_{i-2} + a_i p_{i-1}, \\ q_i = b_i q_{i-2} + a_i q_{i-1}, \end{cases}$$

where $b_0 := 1, p_{-1}, q_{-2} := 1$ and $q_{-1}, p_{-2} := 0$. Then

$$\begin{pmatrix} p_i & p_{i-1} \\ q_i & q_{i-1} \end{pmatrix} = \begin{pmatrix} a_0 & 1 \\ b_0 & 0 \end{pmatrix} \begin{pmatrix} a_1 & 1 \\ b_1 & 0 \end{pmatrix} \dots \begin{pmatrix} a_i & 1 \\ b_i & 0 \end{pmatrix}$$

holds and $\frac{p_i}{q_i}$ is equal to the i th convergent.

Proof. The recurrence relation and the initial condition immediately gives us

$$\begin{pmatrix} p_i & p_{i-1} \\ q_i & q_{i-1} \end{pmatrix} = \begin{pmatrix} p_{i-1} & p_{i-2} \\ q_{i-1} & q_{i-2} \end{pmatrix} \begin{pmatrix} a_i & 1 \\ b_i & 0 \end{pmatrix},$$

which then implies the desired matrix identity at once. To show that $\frac{p_i}{q_i}$ is equal to the i th convergent, we use induction on i . For $i = 0$, the conclusion is trivial. Suppose the result holds for some $i \geq 0$, then by applying the induction hypothesis to the continued fraction we obtain by replacing (a_j, b_j) by (a_{j+1}, b_{j+1}) for each j , we conclude that the i th convergent should be equal to

$$a_0 + \frac{b_1}{r_{i-1}/s_{i-1}} = \frac{a_0 r_{i-1} + b_1 s_{i-1}}{r_{i-1}}$$

where $((r_j, s_j))_j$ is the sequence defined by the recurrence relation

$$\begin{cases} r_j = b'_{j+1} r_{j-2} + a_{j+1} r_{j-1}, \\ s_j = b'_{j+1} s_{j-2} + a_{j+1} s_{j-1} \end{cases}$$

with the same initial condition as $((p_j, q_j))_j$, where $b'_j := b_j$ for all $j > 1$ and $b'_1 := 1$. Then we have

$$\begin{pmatrix} r_{i-1} & r_{i-2} \\ s_{i-1} & s_{i-2} \end{pmatrix} = \begin{pmatrix} a_1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} a_2 & 1 \\ b_2 & 0 \end{pmatrix} \dots \begin{pmatrix} a_i & 1 \\ b_i & 0 \end{pmatrix}.$$

Hence, multiplying

$$\begin{pmatrix} a_0 & 1 \\ b_0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & b_1 \end{pmatrix} = \begin{pmatrix} a_0 & b_1 \\ 1 & 0 \end{pmatrix}$$

to both sides from the left shows that

$$\begin{aligned} \begin{pmatrix} a_0 & b_1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} r_{i-1} & r_{i-2} \\ s_{i-1} & s_{i-2} \end{pmatrix} &= \begin{pmatrix} a_0 & 1 \\ b_0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & b_1 \end{pmatrix} \begin{pmatrix} a_1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} a_2 & 1 \\ b_2 & 0 \end{pmatrix} \cdots \begin{pmatrix} a_i & 1 \\ b_i & 0 \end{pmatrix} \\ &= \begin{pmatrix} a_0 & 1 \\ b_0 & 0 \end{pmatrix} \begin{pmatrix} a_1 & 1 \\ b_1 & 0 \end{pmatrix} \begin{pmatrix} a_2 & 1 \\ b_2 & 0 \end{pmatrix} \cdots \begin{pmatrix} a_i & 1 \\ b_i & 0 \end{pmatrix} = \begin{pmatrix} p_i & p_{i-1} \\ q_i & q_{i-1} \end{pmatrix}, \end{aligned}$$

which gives the desired conclusions $p_i = a_0 r_{i-1} + b_1 s_{i-1}$ and $q_i = r_{i-1}$. \square

Corollary 3.2.

With the notations from Lemma 3.1,

$$p_i q_{i-1} - p_{i-1} q_i = \prod_{j=0}^i (-b_j) \quad (1)$$

and

$$\frac{p_i}{q_i} - \frac{p_{i-1}}{q_{i-1}} = \frac{\prod_{j=0}^i (-b_j)}{q_{i-1} q_i}$$

hold for each i .

Proof. The first claim follows from the matrix identity of Lemma 3.1 and $\det \begin{pmatrix} a_i & 1 \\ b_i & 0 \end{pmatrix} = -b_i$. The second claim is a trivial consequence of the first claim. \square

In general, we can consider the case when a_i 's and b_i 's are whatever objects that can make sense of the entries of the continued fraction, but we are mostly concerned with the case when following further assumptions hold:

- all b_i 's are equal to 1, and
- all a_i 's except for a_0 are positive integers, and
- a_0 is an integer.

In such a case, let us call the continued fraction *regular*.³

For any given finite or infinite sequence a_0, a_1, a_2, \dots of integers with $a_i \geq 1$ for all $i > 0$, there uniquely exists a regular continued fraction which we formally denote as

$$[a_0; a_1, a_2, \dots] := a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{\dots}}}$$

Here, by *formal* we mean that the right-hand side is just a fancy way of referring to the sequence of convergents, and there is no issue of actual convergence involved here.

³It seems terminologies for these concepts are somewhat hazy in the literature. Some refers to regular continued fractions as *simple* continued fractions, and some other refers to them as *standard* continued fractions. In fact, in many of the literatures they just reserve the term *continued fractions* only for the regular ones. In this note, non-regular continued fractions do appear, but only in the context of computing logarithms in Section 13, so we will use the term *continued fractions* to refer to the general case, but we will also mostly only consider the regular ones.

However, we will show below that the sequence actually does converge to a real number, thus we make no distinction between the sequence $[a_0; a_1, a_2, \dots]$ with its limit, although this is strictly speaking an abuse of notation. In particular, for each i such that a_0, \dots, a_i are well-defined, the i th convergent is denoted as $[a_0; a_1, \dots, a_i]$. For example, the 5th convergent of $[0; 1, 2, 3, 4, \dots]$ is denoted as $[0; 1, 2, 3, 4, 5]$ which is equal to

$$0 + \frac{1}{1 + \frac{1}{2 + \frac{1}{3 + \frac{1}{4 + \frac{1}{5}}}}} = \frac{157}{225}.$$

Let us now define the sequence $\left(\frac{p_i}{q_i}\right)_i$ as in Lemma 3.1, that is,

$$\begin{cases} p_i = p_{i-2} + a_i p_{i-1}, \\ q_i = q_{i-2} + a_i q_{i-1} \end{cases}$$

with $p_{-1}, q_{-2} := 1$ and $q_{-1}, p_{-2} := 0$, so that $\frac{p_i}{q_i}$ is equal to the i th convergent. Specializing Corollary 3.2 then gives

$$p_i q_{i-1} - p_{i-1} q_i = (-1)^{i+1} \quad (2)$$

and

$$\frac{p_i}{q_i} - \frac{p_{i-1}}{q_{i-1}} = \frac{(-1)^{i+1}}{q_i q_{i-1}}.$$

From (2), we deduce that p_i and q_i are coprime, so $\frac{p_i}{q_i}$ is in fact the reduced form of the convergent. (Note that q_i is always strictly positive because $q_{-1} = 0$ and $a_i > 0$ for all $i > 0$.) Also, inductively applying the above formula gives

$$\frac{p_i}{q_i} = a_0 + \sum_{j=0}^{i-1} \frac{(-1)^j}{q_j q_{j+1}}.$$

Since the recurrence relation and positivity of a_i 's guarantee at least exponential growth of q_i 's, we obtain:

Corollary 3.3.

We have

$$\frac{p_0}{q_0} < \frac{p_2}{q_2} < \dots < \frac{p_3}{q_3} < \frac{p_1}{q_1}.$$

Also, if the sequence $(a_i)_i$ is infinite, then the sequence $\left(\frac{p_i}{q_i}\right)_{i=0}^\infty$ of convergents converges to a real number.

Conversely, for a given real number x , there is a well-known algorithm for obtaining a continued fraction converging to x . First, take the integer part $a_0 := \lfloor x \rfloor$. Then $x - a_0 \in [0, 1)$. If it is zero, then that means x is an integer, and we terminate the procedure and returns the sequence (a_0) consisting of the single term a_0 . If not, then consider the reciprocal $x_1 := \frac{1}{x - a_0} \in (1, \infty)$. Take the integer part $a_1 := \lfloor x_1 \rfloor$, then again we have $x_1 - a_1 \in [0, 1)$. If it is zero, then stop and returns the sequence $(a_0, \lfloor a_0; a_1 \rfloor)$ consisting of two terms. Otherwise, take the reciprocal and let $x_2 := \frac{1}{x_1 - a_1} \in (1, \infty)$. Again, take the integer part $a_2 := \lfloor x_2 \rfloor$,

and if $x_2 - a_2 = 0$, then stop and returns the sequence $(a_0, [a_0; a_1], [a_0; a_1, a_2])$ consisting of three terms. Otherwise, take the reciprocal $x_3 := \frac{1}{x_2 - a_2}$ and continue.

To see why the resulting continued fractions converge to x , it is sufficient to see $\frac{p_{2i}}{q_{2i}} \leq x$ and $x \leq \frac{p_{2i+1}}{q_{2i+1}}$ always holds, because of Corollary 3.3. Note that by induction one can easily see

$$x = [a_0; a_1, \dots, a_i, x_{i+1}]$$

always holds (here we are allowing non-integers to appear in continued fractions), so it is enough to show that for given positive real numbers $a < b$,

$$[a_0; a_1, \dots, a_{2i}, a] < [a_0; a_1, \dots, a_{2i}, b]$$

and

$$[a_0; a_1, \dots, a_{2i+1}, a] > [a_0; a_1, \dots, a_{2i+1}, b]$$

holds. These can be easily shown using induction.

Note that if this algorithm terminates in a finite step, then that means the final term in the resulting sequence is precisely the number x we are given with. In particular, x must be a rational number in this case, so if x is irrational then the algorithm produces an infinite sequence.

On the other hand, if $x = \frac{p}{q}$ is rational in its reduced form⁴, then what the algorithm does is basically what the Euclid algorithm does to p and q to eventually produce $\gcd(p, q) = 1$. (Details are omitted.) With this correspondence, one can see that whenever x is rational the algorithm must terminate in a finite step. In other words, the algorithm terminates in a finite step if and only if x is rational.

Now, suppose we have $x = [a_0; a_1, \dots]$. Then, it can be easily seen that

$$a_0 \leq x < a_0 + 1$$

must hold unless a_1 is the last term and $a_1 = 1$, in which case we have $x = a_0 + 1$. Also, since $[a_0; a_1, \dots] = [a_0; b_1, \dots]$ implies $[a_1; a_2, \dots] = [b_1; b_2, \dots]$, we get that two continued fractions $[a_0; a_1, \dots]$ and $[b_0; b_1, \dots]$ yield the same real number if and only if:

1. Either they are identical sequences, or
2. Both are finite sequences with their lengths differ by 1, and assuming $[a_0; a_1, \dots] = [a_0; a_1, \dots, a_i]$ is the shorter one, $a_j = b_j$ holds for all $j < i$ and $a_i = b_i + 1$, $b_{i+1} = 1$ hold.

In other words,

1. When x is irrational, then there uniquely exists an infinite continued fraction converging to x , which is generated by the algorithm described above, and
2. When x is rational, then along with the one generated by the algorithm, there is another continued fraction which is obtained by replacing the last coefficient a_i with $[a_i - 1; 1]$. Note that whenever the algorithm terminates in a finite step, the last coefficient must be at least 2.

From now on, we will call the continued fraction $[a_0; a_1, \dots]$ obtained from the algorithm as *the continued fraction expansion of x* , and the i th convergent of it as *the i th convergent of x* . We call this $[a_0; a_1, \dots]$ the *standard* continued fraction expansion of x , if we need

⁴That is, q is a positive integer and p is an integer with $\gcd(p, q) = 1$.

to distinguish it from the other expansion, which we call the *alternative* continued fraction expansion.

Given a continued fraction expansion $x = [a_0; a_1, \dots]$, whenever we say “the i th coefficient” or “the i th convergent” or anything like that, we are always implicitly assuming that those are well-defined. In particular, the index i will be always assumed to be at least 0. However, sometimes it is convenient to talk about “the -1 st convergent” which is conventionally defined as ∞ . In a similar spirit, we often view “the coefficient next to the last coefficient” to be ∞ , and the one next to it to be 0, and the one next to it to be ∞ again, and so on. Then it is natural to conceptually think that, if $\frac{p_i}{q_i}$ is the last convergent (so that $x = \frac{p_i}{q_i}$), then p_{i+1} and q_{i+1} are both ∞ , while at the same time the quotient $\frac{p_{i+1}}{q_{i+1}}$ is considered to be equal to $\frac{p_i}{q_i}$. Then, we have $p_{i+2} = p_i$ and $q_{i+2} = q_i$, and then $p_{i+3} = p_{i+1} = \infty$ and $q_{i+3} = q_{i+1} = \infty$ and $\frac{p_{i+3}}{q_{i+3}} = \frac{p_i}{q_i}$, and then again $p_{i+4} = p_i$ and $q_{i+4} = q_i$, and so on.

4 Semiconvergents

In this section, we will show that best rational approximations of x from below/above in the weak sense and in the strong sense are same, by showing that they are equivalent to *proper semiconvergents* of x , which we define below. This immediately gives an efficient algorithm for computing the best rational approximations from below/above.

Definition 4.1 (Semiconvergents).

Let $\frac{p_i}{q_i}$ denote the i th convergent of a given regular continued fraction $x = [a_0; a_1, \dots]$. Then a rational number $\frac{p}{q}$ in its reduced form is said to be a *semiconvergent* of the continued fraction $[a_0; a_1, \dots]$, if either

1. $\frac{p}{q} = x$, or
2. $\frac{p}{q} = \frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$ for some consecutive convergents $\frac{p_{i-1}}{q_{i-1}}, \frac{p_i}{q_i}$ with $i \geq 0$ and an integer $s \in [0, a_{i+1}]$, with the convention $a_{i+1} := \infty$ if $x = \frac{p_i}{q_i}$.

A semiconvergent $\frac{p}{q}$ is said to be *proper* if either x is irrational or is rational and its reduced denominator is strictly larger than q .⁵

Remark 4.2.

Semiconvergents are meant to be a monotonic sequence interpolating between two consecutive even/odd convergents (see below). However, there seems to be some ambiguity in how exactly they should be defined, regarding various corner cases occurring for finite continued fractions. For instance, the Wikipedia page on continued fractions (retrieved 2024-07-06) defines semiconvergents as any numbers of the form $\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$ for consecutive convergents $\frac{p_{i-1}}{q_{i-1}}, \frac{p_i}{q_i}$ and an integer $s \in [0, a_{i+1}]$. However, it is unclear if they intend to include the cases $i = 0$, or when $\frac{p_i}{q_i}$ is the last convergent (there is no mention at all of how they would define a_{i+1} in this case), or if they allow $i = 0$ then whether they want to include ∞ , or things like that. Note that, when x is an integer and the standard expansion is used, then the only possible interpretation of their definition is that, either we count every number of the form $x + \frac{1}{s}$ as semiconvergent but *not* x itself, or no number is counted as a semiconvergent. Or maybe they allow $s = \infty$ for the case of last convergent? When x is an integer and the alternative expansion is used, then x is included if and only if they allow $i = 0$, and $x - 1$ is included if and only if they allow the case of last convergent. They do not clarify

⁵The term *proper semiconvergent* is not standard.

which exact option they want to take for this case, though it is not even clear if the alternative expansion is allowed at all. When x is not an integer and $\frac{p_1}{q_1}$ is the last convergent (the standard expansion is used), then x itself is only counted as a semiconvergent if they allow $i = 0$, or maybe $s = \infty$. Again it is hard to know which option they wanted to take. Therefore, I tried my best to clarify these ambiguities as much as possible (because these corner cases do impact the correctness of the algorithm we will describe), and arrived at the definition written above. A few notable special cases are:

- If x is rational, then we *always* count x itself as a semiconvergent.
- We *do allow* the case $i = 0$ and the case when $\frac{p_i}{q_i}$ is the last convergent. For the latter case, s can be any nonnegative integer. We explicitly do not include the case $s = \infty$, but $x = \frac{p_i}{q_i}$ is always counted as a semiconvergent anyway, which is effectively equivalent to allowing $s = \infty$.
- We *do not* count $\frac{p_{-1}}{q_{-1}} = \infty$ as a semiconvergent, because semiconvergents are a priori defined to be rational numbers. Hence, when $i = 0$, s should be a positive integer.
- On the other hand, we *do count* every convergent $\frac{p_i}{q_i}$ with $i \geq 0$ as a semiconvergent. Indeed, if $\frac{p_i}{q_i}$ is not the last convergent, then $\frac{p_i}{q_i} = \frac{p_i + 0 \cdot p_{i+1}}{q_i + 0 \cdot q_{i+1}}$ is a semiconvergent, and if it is the last convergent, then $x = \frac{p_i}{q_i}$ so it is still counted as a semiconvergent.
- It may seem like we disallow the case $s = a_{i+1}$, but this has no impact because $\frac{p_{i-1} + a_{i+1}p_i}{q_{i-1} + a_{i+1}q_i} = \frac{p_{i+1}}{q_{i+1}}$ is a semiconvergent whenever $\frac{p_i}{q_i}$ is not the last convergent, and if it is the last convergent then $a_{i+1} = \infty$.
- There always are exactly two integer semiconvergents, namely a_0 and $a_0 + 1$. Indeed, $a_0 + 1 = \frac{p_{-1} + 1 \cdot p_0}{q_{-1} + 1 \cdot q_0}$ is always a semiconvergent. When there are at least two convergents, then $a_0 = \frac{p_0 + 0 \cdot p_1}{q_0 + 0 \cdot q_1}$ is a semiconvergent, and when there is only one convergent, then $a_0 = \frac{p_0}{q_0} = x$ is anyway a semiconvergent. Conversely, if $\frac{p}{1}$ is a semiconvergent, then either $x = \frac{p}{1}$ or $1 = q_{i-1} + s q_i$ for some i and $s \in [0, a_{i+1})$. For the first case, if the standard expansion is used, then $a_0 = x = \frac{p}{1}$, and if the alternative expansion is used, then $a_0 + 1 = x = \frac{p}{1}$. For the second case, we either have $q_{i-1} = 0$ and $s = q_i = 1$ or $q_{i-1} = 1$ and $s = 0$. The first case can happen only when $i = 0$, in which case we have $\frac{p}{1} = \frac{1 + p_0}{1} = a_0 + 1$. For the second case, since $(q_j)_j$ is a strictly increasing sequence, we must have $i = 1$, so $\frac{p}{1} = \frac{p_{i-1}}{q_{i-1}} = \frac{p_0}{q_0} = a_0$.
- Hence, if x is not an integer, then $\lfloor x \rfloor = a_0$ and $\lceil x \rceil = a_0 + 1$ are the only integer semiconvergents, and they are both *proper* semiconvergents.
- If x is an integer, then x and $x + 1$ are the only integer semiconvergents if the standard expansion is used, while $x - 1$ and x are the only integer semiconvergents if the alternative expansion is used. There is no proper semiconvergent in this case.

When i is even, we call $\frac{p_i}{q_i}$ an *even convergent*, and when i is odd, we call it an *odd convergent*. Similarly, by an *even semiconvergent* we mean either x itself when x is an even convergent of itself, or any semiconvergent of the form

$$\frac{p_{2i} + s p_{2i+1}}{q_{2i} + s q_{2i+1}},$$

and by an *odd semiconvergent* we mean either x itself when x is an odd convergent of itself, or any semiconvergent of the form

$$\frac{p_{2i-1} + sp_{2i}}{q_{2i-1} + sq_{2i}}.$$

By *consecutive semiconvergents* we mean two semiconvergents of the form $\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$ and $\frac{p_{i-1} + (s+1)p_i}{q_{i-1} + (s+1)q_i}$ with $s < a_{i+1}$. Note that $s + 1 = a_{i+1}$ is allowed, in which case the second semiconvergent becomes $\frac{p_{i+1}}{q_{i+1}}$.

Lemma 4.3.

Let $\frac{p_i}{q_i}$ denote the i th convergent of a given regular continued fraction $x = [a_0; a_1, \dots]$. Then for a given two consecutive convergents $\frac{p_{i-1}}{q_{i-1}}, \frac{p_i}{q_i}$, as s varies from 0 to a_{i+1} , the semiconvergent

$$\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$$

strictly increases from $\frac{p_{i-1}}{q_{i-1}}$ to $\frac{p_{i+1}}{q_{i+1}}$ when i is odd, and it strictly decreases from $\frac{p_{i-1}}{q_{i-1}}$ to $\frac{p_{i+1}}{q_{i+1}}$ when i is even, where we consider $\frac{p_{i+1}}{q_{i+1}}$ to be equal to x if $\frac{p_i}{q_i}$ is the last convergent.

Proof. Consider the function

$$f(s) := \frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$$

of real variable s , then from the identity

$$p_i q_{i-1} - p_{i-1} q_i = (-1)^{i+1}$$

it can be easily verified that $f'(s) > 0$ when i is odd and $f'(s) < 0$ when i is even, on $s > 0$. Hence, the claim follows. \square

In fact, $s = a_{i+1}$ (assuming $a_{i+1} < \infty$) is the last s such that the order between x and $\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$ is retained; that is, we have

$$x \leq \frac{p_{i-1} + (a_{i+1} + 1)p_i}{q_{i-1} + (a_{i+1} + 1)q_i} = \frac{p_i + p_{i+1}}{q_i + q_{i+1}}$$

when i is odd, and

$$x \geq \frac{p_{i-1} + (a_{i+1} + 1)p_i}{q_{i-1} + (a_{i+1} + 1)q_i} = \frac{p_i + p_{i+1}}{q_i + q_{i+1}}$$

when i is even. To see why, note that if $a_{i+2} = \infty$, then $x = \frac{p_{i+1}}{q_{i+1}}$, and (2) shows

$$\frac{p_i + p_{i+1}}{q_i + q_{i+1}} - \frac{p_{i+1}}{q_{i+1}} = \frac{p_i q_{i+1} - p_{i+1} q_i}{(q_i + q_{i+1})q_{i+1}} = \frac{(-1)^{i+1}}{(q_i + q_{i+1})q_{i+1}},$$

so we indeed get the desired inequalities and they are always strict. If $a_{i+2} < \infty$, then

$$\begin{aligned} \frac{p_i + p_{i+1}}{q_i + q_{i+1}} - \frac{p_{i+2}}{q_{i+2}} &= \frac{(p_i q_{i+2} - p_{i+2} q_i) + (p_{i+1} q_{i+2} - p_{i+1} q_{i+2})}{(q_i + q_{i+1})q_{i+2}} \\ &= \frac{a_{i+2}(p_i q_{i+1} - p_{i+1} q_i) + (-1)^i}{(q_i + q_{i+1})q_{i+2}} = \frac{(-1)^{i+1}(a_{i+2} - 1)}{(q_i + q_{i+1})q_{i+2}}. \end{aligned}$$

Since $x \leq \frac{p_{i+2}}{q_{i+2}}$ when i is odd and $x \geq \frac{p_{i+2}}{q_{i+2}}$ when i is even, we get

$$x \leq \frac{p_i + p_{i+1}}{q_i + q_{i+1}} - \frac{(-1)^{i+1}(a_{i+2} - 1)}{(q_i + q_{i+1})q_{i+2}} \leq \frac{p_i + p_{i+1}}{q_i + q_{i+1}}$$

when i is odd, and

$$x \geq \frac{p_i + p_{i+1}}{q_i + q_{i+1}} - \frac{(-1)^{i+1}(a_{i+2} - 1)}{(q_i + q_{i+1})q_{i+2}} \geq \frac{p_i + p_{i+1}}{q_i + q_{i+1}}$$

when i is even, with the equalities if and only if $x = \frac{p_{i+2}}{q_{i+2}}$ and $a_{i+2} = 1$.

Note that this equality condition means that x is rational and the continued fraction $[a_0; a_1, \dots]$ ended with 1, which can never happen if we use the standard expansion of x . Hence, we get:

Lemma 4.4.

Let $\frac{p_i}{q_i}$ denote the i th convergent of the standard continued fraction expansion $x = [a_0; a_1, \dots]$. Then

$$a_{i+1} = \left\lfloor \frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} \right\rfloor$$

holds for each i .

Proof. Follows directly from the discussion above together with that

$$\frac{p_{2i} + sp_{2i+1}}{q_{2i} + sq_{2i+1}} \leq x$$

is equivalent to

$$s \leq \frac{q_{2i}x - p_{2i}}{p_{2i+1} - q_{2i+1}x}$$

and

$$\frac{p_{2i-1} + sp_{2i}}{q_{2i-1} + sq_{2i}} \geq x$$

is equivalent to

$$s \leq \frac{p_{2i-1} - q_{2i-1}x}{q_{2i}x - p_{2i}}.$$

□

From now on, we will build a theory about how semiconvergents divide the approximation intervals and how rational numbers in between two consecutive semiconvergents should look like. These will lead to the main theorem of this section, which is the characterization of best rational approximations from below/above in terms of semiconvergents.

Lemma 4.5.

Let $\frac{p_i}{q_i}$ denote the i th convergent of a given regular continued fraction $x = [a_0; a_1, \dots]$. Then for any $y \in [a_0, x)$, there uniquely exist two consecutive even semiconvergents $\frac{p_{2i} + sp_{2i+1}}{q_{2i} + sq_{2i+1}}$ and $\frac{p_{2i} + (s+1)p_{2i+1}}{q_{2i} + (s+1)q_{2i+1}}$ such that

$$\frac{p_{2i} + sp_{2i+1}}{q_{2i} + sq_{2i+1}} \leq y < \frac{p_{2i} + (s+1)p_{2i+1}}{q_{2i} + (s+1)q_{2i+1}}.$$

Similarly, for any $z \in (x, a_0 + 1]$, there uniquely exist two consecutive odd semiconvergents $\frac{p_{2i-1} + sp_{2i}}{q_{2i-1} + sq_{2i}}$ and $\frac{p_{2i-1} + (s+1)p_{2i}}{q_{2i-1} + (s+1)q_{2i}}$ such that

$$\frac{p_{2i-1} + (s+1)p_{2i}}{q_{2i-1} + (s+1)q_{2i}} < z \leq \frac{p_{2i-1} + sp_{2i}}{q_{2i-1} + sq_{2i}}.$$

Proof. By Corollary 3.3, either y lies in between two consecutive even convergents $\frac{p_{2i}}{q_{2i}}$ and $\frac{p_{2i+2}}{q_{2i+2}}$, that is,

$$\frac{p_{2i}}{q_{2i}} \leq y < \frac{p_{2i+2}}{q_{2i+2}},$$

or it lies in between the last even convergent $\frac{p_{2i}}{q_{2i}}$ and $x = \frac{p_{2i+1}}{q_{2i+1}}$, that is,

$$\frac{p_{2i}}{q_{2i}} \leq y < \frac{p_{2i+1}}{q_{2i+1}} = x.$$

Then for either case, Lemma 4.3 shows that there uniquely exists an integer $s \in [0, a_{2i+2})$ such that

$$\frac{p_{2i} + sp_{2i+1}}{q_{2i} + sq_{2i+1}} \leq y < \frac{p_{2i} + (s+1)p_{2i+1}}{q_{2i} + (s+1)q_{2i+1}}$$

holds.

Similarly, with the convention $\frac{p_{-1}}{q_{-1}} = \infty$, z must satisfy either

$$\frac{p_{2i+1}}{q_{2i+1}} < z \leq \frac{p_{2i-1}}{q_{2i-1}}$$

or

$$x = \frac{p_{2i}}{q_{2i}} < z \leq \frac{p_{2i-1}}{q_{2i-1}},$$

and for either case we can find a unique integer $s \in [0, a_{2i+2})$ such that

$$\frac{p_{2i-1} + (s+1)p_{2i}}{q_{2i-1} + (s+1)q_{2i}} < z \leq \frac{p_{2i-1} + sp_{2i}}{q_{2i-1} + sq_{2i}}.$$

□

Lemma 4.6.

Let $\frac{a}{b} < \frac{c}{d}$ be rational numbers in their reduced form. Then any rational number $\frac{x}{y} \in (\frac{a}{b}, \frac{c}{d})$ in its reduced form can be uniquely written as

$$\frac{x}{y} = \frac{qa + pc}{qb + pd}$$

for some coprime positive integers p, q . If $bc - ad = 1$, then we have

$$\begin{cases} p = bx - ay, \\ q = cy - dx, \end{cases} \quad \text{and} \quad \begin{cases} x = qa + pc, \\ y = qb + pd. \end{cases}$$

Hence, $\frac{qa+pc}{qb+pd}$ is in its reduced form in this case. The exact same result is true even when we allow $\frac{a}{b} = \frac{0}{1} = 0$ or $\frac{c}{d} = \frac{1}{0} = \infty$ or both.

Proof. Consider the function

$$f(t) := \frac{a + tc}{b + td}$$

of real variable t , then

$$f'(t) = \frac{bc - ad}{(b + td)^2} > 0,$$

so f is strictly increasing on $t > 0$. Hence, if we have

$$\frac{q_1 a + p_1 c}{q_1 b + p_1 d} = \frac{q_2 a + p_2 c}{q_2 b + p_2 d}$$

for some positive integers p_1, q_1, p_2, q_2 , then we must have

$$\frac{p_1}{q_1} = \frac{p_2}{q_2}.$$

In particular, if $\gcd(p_1, q_1) = \gcd(p_2, q_2) = 1$, then we must have $p_1 = p_2$ and $q_1 = q_2$. This shows the uniqueness part. Also, note that

$$\frac{x}{y} = \frac{(cy - dx)a + (bx - ay)c}{(cy - dx)b + (bx - ay)d}$$

always holds. Indeed, the numerator is equal to $(bc - ad)x$ while the denominator is equal to $(bc - ad)y$. Hence, let $g := \gcd(cy - dx, bx - ay)$, then we have $p = (bx - ay)/g$ and $q = (cy - dx)/g$. Then it only remains to show that $g = 1$ if $bc - ad = 1$.

Since x, y are coprime, we can find integers z, w such that $xz - yw = 1$. Then, the determinant of the matrix

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} y & -z \\ -x & w \end{pmatrix} = \begin{pmatrix} ay - bx & bw - az \\ cy - dx & dw - cz \end{pmatrix}$$

must be equal to $bc - ad$. Hence, when $bc - ad = 1$, then $bx - ay$ and $cy - dx$ must be indeed coprime to each other. \square

Corollary 4.7.

Let $\frac{p_i}{q_i}$ denote the i th convergent of a given regular continued fraction $[a_0; a_1, \dots]$. Then any rational number non-strictly between two consecutive semiconvergents $\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$ and $\frac{p_{i-1} + (s+1)p_i}{q_{i-1} + (s+1)q_i}$ is of the form

$$\frac{bp_{i-1} + (a + bs)p_i}{bq_{i-1} + (a + bs)q_i}$$

for some rational number $\frac{a}{b} \in [0, 1]$. Furthermore, the above expression is always in its reduced form if $\frac{a}{b}$ is in its reduced form.

Proof. It immediately follows from the previous lemma that any rational number between those two semiconvergents must be of the form

$$\frac{qp_{i-1} + pp_i}{qq_{i-1} + pq_i}$$

for some coprime nonnegative integers p, q , because such a rational number must be in between $\frac{p_{i-1}}{q_{i-1}}$ and $\frac{p_i}{q_i}$. Furthermore, the above expression must be of the reduced form. Note that

$$\begin{aligned} & (qp_{i-1} + pp_i)(q_{i-1} + sq_i) - (qq_{i-1} + pq_i)(p_{i-1} + sp_i) \\ &= (p - qs)(p_i q_{i-1} - p_{i-1} q_i) = (-1)^{i-1}(p - qs), \end{aligned}$$

and similarly the same computation with s replaced by $s + 1$ gives $(-1)^{i-1}(p - q(s + 1))$. Therefore, $\frac{qp_{i-1} + pp_i}{qq_{i-1} + pq_i}$ lies in between $\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$ and $\frac{p_{i-1} + (s+1)p_i}{q_{i-1} + (s+1)q_i}$ if and only if $\frac{p}{q} \in [s, s + 1]$. Hence, let $a = p - qs$ and $b = q$, then we get the desired result. \square

Corollary 4.8.

Let $\frac{p_i}{q_i}$ denote the i th convergent of a given regular continued fraction $x = [a_0; a_1, \dots]$. Then for any rational number $\frac{p}{q}$ in its reduced form between two consecutive semiconvergents, in either $\left[\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}, \frac{p_{i-1} + (s+1)p_i}{q_{i-1} + (s+1)q_i} \right)$ or $\left(\frac{p_{i-1} + (s+1)p_i}{q_{i-1} + (s+1)q_i}, \frac{p_{i-1} + sp_i}{q_{i-1} + sq_i} \right]$ depending on the parity of i , we always have

$$|qx - p| \geq |(q_{i-1} + sq_i)x - (p_{i-1} + sp_i)|.$$

The equality condition is:

1. If $s + 1 = a_{i+1}$ and $x = \frac{p_{i+1}}{q_{i+1}}$ (thus in particular $\frac{p_i}{q_i}$ is not the last convergent and $x \neq \frac{p_i}{q_i}$), then the equality holds if and only if

$$\begin{cases} p = bp_{i-1} + (bs + (b - 1))p_i, \\ q = bq_{i-1} + (bs + (b - 1))q_i \end{cases}$$

for a positive integer b .

2. Otherwise, the equality holds if and only if

$$\begin{cases} p = p_{i-1} + sp_i, \\ q = q_{i-1} + sq_i. \end{cases}$$

Proof. By Corollary 4.7, we know

$$\begin{cases} p = bp_{i-1} + (a + bs)p_i, \\ q = bq_{i-1} + (a + bs)q_i \end{cases}$$

for some rational number $\frac{a}{b} \in [0, 1)$ in its reduced form. Assume $p_i - q_i x \neq 0$ first, then we get

$$|qx - p| = |b(q_{i-1}x - p_{i-1}) - (a + bs)(p_i - q_i x)| = b|p_i - q_i x| \left| \frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} - \left(s + \frac{a}{b}\right) \right|,$$

and similarly

$$|(q_{i-1} + sq_i)x - (p_{i-1} + sp_i)| = |p_i - q_i x| \left| \frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} - s \right|.$$

Recall from Lemma 4.4 that $\left\lfloor \frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} \right\rfloor = a_{i+1}$ if $[a_0; a_1, \dots]$ is the standard continued fraction expansion of x . Otherwise, if $\frac{p_{i+1}}{q_{i+1}}$ is the last convergent, then $\frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} = a_{i+1} = 1$ holds because we know $x = \frac{p_{i-1} + p_i}{q_{i-1} + q_i}$ which shows that the left-hand side is equal to

$$\frac{q_{i-1}(p_{i-1} + p_i) - p_{i-1}(q_{i-1} + q_i)}{p_i(q_{i-1} + q_i) - q_i(p_{i-1} + p_i)} = 1,$$

and if $\frac{p_{i+2}}{q_{i+2}}$ is the last convergent, then $\left\lfloor \frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} \right\rfloor = a_{i+1} + 1$ holds, and for any other case, $\left\lfloor \frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} \right\rfloor = a_{i+1}$ should hold. Therefore, in any case, we always have $\frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} \geq a_{i+1} \geq s + 1$. This shows

$$\frac{|qx - p| - |(q_{i-1} + sq_i)x - (p_{i-1} + sp_i)|}{|p_i - q_i x|} = (b - 1) \left(\frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} - \left(s + \frac{a}{b}\right) \right) - \frac{a}{b}.$$

Then since

$$\frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} \geq s + 1 \geq s + \frac{a + 1}{b},$$

we get

$$\frac{|qx - p| - |(q_{i-1} + sq_i)x - (p_{i-1} + sp_i)|}{|p_i - q_i x|} \geq \frac{b - 1}{b} - \frac{a}{b} \geq 0,$$

as desired. Also, the equality holds if and only if either $\frac{a}{b} = 0$ (i.e., $a = 0$ and $b = 1$) or $a = b - 1$ and

$$\frac{q_{i-1}x - p_{i-1}}{p_i - q_i x} = s + 1,$$

which means $x = \frac{p_{i+1}}{q_{i+1}}$. Hence, the claimed equality condition follows.

The case $p_i - q_i x = 0$ is simpler; in this case, we have

$$|qx - p| = b |q_{i-1}x - p_{i-1}|$$

and

$$|(q_{i-1} + sq_i)x - (p_{i-1} + sp_i)| = |q_{i-1}x - p_{i-1}|,$$

so we get the desired inequality and the equality condition is $b = 1$, which implies $a = 0$ and $\frac{p}{q} = \frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$. \square

Lemma 4.9.

Let $\frac{p_i}{q_i}$ denote the i th convergent of a given regular continued fraction $x = [a_0; a_1, \dots]$.

Then for any two consecutive semiconvergents $\frac{p_{i-1} + sp_i}{q_{i-1} + sq_i}$ and $\frac{p_{i-1} + (s+1)p_i}{q_{i-1} + (s+1)q_i}$,

$$|(q_{i-1} + sq_i)x - (p_{i-1} + sp_i)| \geq |(q_{i-1} + (s+1)q_i)x - (p_{i-1} + (s+1)p_i)|.$$

The equality holds if and only if $x = \frac{p_i}{q_i}$.

Proof. This follows from direct computation:

$$((q_{i-1} + sq_i)x - (p_{i-1} + sp_i)) - ((q_{i-1} + (s+1)q_i)x - (p_{i-1} + (s+1)p_i)) = p_i - q_i x.$$

If i odd, then the right-hand side is nonnegative and $(q_{i-1} + sq_i)x - (p_{i-1} + sp_i)$, $(q_{i-1} + (s+1)q_i)x - (p_{i-1} + (s+1)p_i)$ are both nonnegative, so we get the desired inequality. If i is even, all three are nonpositive, so we still get the desired inequality. The equality condition is precisely $x = \frac{p_i}{q_i}$. \square

Hence, we get:

Corollary 4.10.

Let $\left(\frac{r_j}{s_j}\right)_j$ be the sequence of all even (odd, resp.) semiconvergents of a regular continued fraction $x = [a_0; a_1, \dots]$, all in their reduced form, such that $r_j < r_{j+1}$ holds for all j .⁶

1. If x is rational and is an even (odd, resp.) convergent of itself, then the sequence $(|s_j x - r_j|)_j$ is finite and it strictly decreases until it reaches zero at the final term.
2. If x is rational and is an odd (even, resp.) convergent of itself, then the sequence $(|s_j x - r_j|)_j$ is infinite and it strictly decreases until $\frac{r_j}{s_j}$ becomes the last even convergent of x , and it stays constant after that.
3. If x is irrational, then the sequence $(|s_j x - r_j|)_j$ is infinite and it strictly decreases forever.

Finally, we prove our main theorem in this section.

Theorem 4.11 (Characterization of best rational approximations from below/above).

Let $x = [a_0; a_1, \dots]$ be a regular continued fraction. Then for a rational number $\frac{p}{q} \neq x$ in its reduced form, the followings are equivalent.

1. $\frac{p}{q}$ is an even (odd, resp.) proper semiconvergent of x .
2. $\frac{p}{q}$ is a best rational approximation of x from below (above, resp.) in the strong sense.
3. $\frac{p}{q}$ is a best rational approximation of x from below (above, resp.) in the weak sense.

Proof. Let $\frac{p_i}{q_i}$ denote the i th convergent of $x = [a_0; a_1, \dots]$.

($1 \Rightarrow 2$) Let $\frac{p}{q}$ be a proper even semiconvergent. Then we can write $\frac{p}{q} = \frac{p_{2i} + s p_{2i+1}}{q_{2i} + s q_{2i+1}}$ for some consecutive convergents $\frac{p_{2i}}{q_{2i}}, \frac{p_{2i+1}}{q_{2i+1}}$ and an integer $s \in [0, a_{2i+2})$. Clearly, $\frac{p}{q} \leq x$ holds. Let $\frac{a}{b}$ be any rational number with $0 < b \leq q$ and $\frac{a}{b} \leq x$. Since $\frac{p}{q}$ is a proper semiconvergent and $0 < b \leq q$, we should have $\frac{a}{b} \neq x$. First, suppose $\frac{a}{b} < a_0$. Then Corollary 4.10 shows

$$bx - a > b(x - a_0) = b(q_0 x - p_0) \geq b(qx - p) \geq qx - p.$$

Next, suppose $\frac{a}{b} \geq a_0$, then by Lemma 4.5, there uniquely exist two consecutive even semiconvergents $\frac{p_{2j} + t p_{2j+1}}{q_{2j} + t q_{2j+1}}$ and $\frac{p_{2j} + (t+1)p_{2j+1}}{q_{2j} + (t+1)q_{2j+1}}$ such that

$$\frac{p_{2j} + t p_{2j+1}}{q_{2j} + t q_{2j+1}} \leq \frac{a}{b} < \frac{p_{2j} + (t+1)p_{2j+1}}{q_{2j} + (t+1)q_{2j+1}}.$$

By Corollary 4.7, we know

$$q_{2j} + t q_{2j+1} \leq b \leq q = q_{2i} + s q_{2i+1},$$

thus Corollary 4.10 shows

$$(q_{2j} + t q_{2j+1})x - (p_{2j} + t p_{2j+1}) \geq qx - p.$$

Hence, from Corollary 4.8, we get

$$bx - a \geq (q_{2j} + t q_{2j+1})x - (p_{2j} + t p_{2j+1}) \geq qx - p,$$

⁶Note that no two distinct even (odd, resp.) semiconvergents share their denominators.

thus we conclude that $\frac{p}{q}$ is a best rational approximation of x from below in the strong sense. In the same way, we can show that if $\frac{p}{q}$ is a proper odd semiconvergent, then it is a best rational approximation of x from above in the strong sense.

($2 \Rightarrow 3$) Done in Lemma 2.3.

($3 \Rightarrow 1$) Let $\frac{p}{q}$ be a best rational approximation of x from below in the weak sense. Since $\frac{p_0}{q_0} = a_0 \leq x$ and $q_0 = 1 \leq q$ holds, the assumption on $\frac{p}{q}$ implies $a_0 \leq \frac{p}{q} < x$. Hence, by Lemma 4.5, there uniquely exist two consecutive semiconvergents $\frac{p_{2i}+sp_{2i+1}}{q_{2i}+sq_{2i+1}}$ and $\frac{p_{2i}+(s+1)p_{2i+1}}{q_{2i}+(s+1)q_{2i+1}}$ such that

$$\frac{p_{2i}+sp_{2i+1}}{q_{2i}+sq_{2i+1}} \leq \frac{p}{q} < \frac{p_{2i}+(s+1)p_{2i+1}}{q_{2i}+(s+1)q_{2i+1}}.$$

Hence, Corollary 4.7 shows that there exist a rational number $\frac{a}{b} \in [0, 1)$ in its reduced form such that

$$\frac{p}{q} = \frac{bp_{2i} + (a + bs)p_{2i+1}}{bq_{2i} + (a + bs)q_{2i+1}}.$$

We claim that $\frac{a}{b} = \frac{0}{1}$. Suppose otherwise, then $a, b \geq 1$, so $q \geq bq_{2i} + (a + bs)q_{2i+1}$ holds, but since $\frac{p}{q}$ is a best rational approximation of x from below in the weak sense, this implies $\frac{p_{2i}+(s+1)p_{2i+1}}{q_{2i}+(s+1)q_{2i+1}} \leq \frac{p}{q}$, which is a contradiction. Therefore, we conclude that $\frac{p}{q} = \frac{p_{2i}+sp_{2i+1}}{q_{2i}+sq_{2i+1}}$ is an even semiconvergent. Since $\frac{p_{2i}+(s+1)p_{2i+1}}{q_{2i}+(s+1)q_{2i+1}}$ is also a semiconvergent with a strictly larger denominator, it follows that $\frac{p}{q}$ is a proper semiconvergent. \square

This characterization of best rational approximations from below/above immediately leads us to the following conclusion: for a given real number x and a positive integer n_{\max} , the best rational approximations $\frac{p_*}{q_*}$ from below and $\frac{p^*}{q^*}$ from above of x with the largest denominators $q_*, q^* \leq n_{\max}$, are either both equal to x , or the even/odd semiconvergents of largest possible denominators bounded by n_{\max} (which are necessarily proper). Therefore, we get the following simple algorithm for computing $\frac{p_*}{q_*}$ and $\frac{p^*}{q^*}$, for any given x of which we can compute the continued fraction expansion:

Algorithm 4.12 (Finding best rational approximations from below and above).

1. Find the last convergent $\frac{p_i}{q_i}$ of x such that $q_i \leq n_{\max}$.
2. If $x = \frac{p_i}{q_i}$, then return $\left(\frac{p_*}{q_*}, \frac{p^*}{q^*}\right) = \left(\frac{p_i}{q_i}, \frac{p_i}{q_i}\right)$.
3. Otherwise, find the last semiconvergent $\frac{p_{i-1}+sp_i}{q_{i-1}+sq_i}$ such that $q_{i-1}+sq_i \leq n_{\max}$.
4. If i is even, then return $\left(\frac{p_*}{q_*}, \frac{p^*}{q^*}\right) = \left(\frac{p_i}{q_i}, \frac{p_{i-1}+sp_i}{q_{i-1}+sq_i}\right)$.
5. If i is odd, then return $\left(\frac{p_*}{q_*}, \frac{p^*}{q^*}\right) = \left(\frac{p_{i-1}+sp_i}{q_{i-1}+sq_i}, \frac{p_i}{q_i}\right)$.

Note that all we need to know about x is how to compute its continued fraction expansion. Because of this, the reference implementation is written in a way that all real numbers are supposed to be encoded in terms of their continued fraction expansions. This viewpoint of thinking real numbers and their continued fraction expansions as identical will be retained throughout the whole note, as it is particularly convenient for our purpose.

5 Extremizers of $(nx + y) - \lfloor nx + y \rfloor$

In this section, we now obtain the algorithm for efficiently computing extremizers of $(nx + y) - \lfloor nx + y \rfloor$ with full generality. The algorithm we derive in fact is a simple extension of the classic Euclid's algorithm.

We will actually find the *smallest minimizer* and the *largest maximizer* of $(nx + y) - \lfloor nx + y \rfloor$ for reasons that will become apparent later. Throughout this section, let $x, y \in \mathbb{R}$ be any real numbers and $n_{\min} \leq n_{\max}$ be any integers.

We start with the algorithm for finding the smallest minimizer. In fact, the algorithm follows at once by examining how exactly the fractional part of $nx + y$ would change when n increases. Let us say we start with $n_0 = n_{\min}$ and we are trying to compare the fractional part $(n_0x + y) - \lfloor n_0x + y \rfloor$ of $n_0x + y$ with the fractional part $((n_0 + n)x + y) - \lfloor (n_0 + n)x + y \rfloor$ of $(n_0 + n)x + y$. Clearly, if the fractional part of nx is small, then adding it to $n_0x + y$ will just increase the fractional part. Specifically, if the fractional part of nx is strictly less than current the margin of the fractional part of $n_0x + y$, then the fractional parts of them will just get added together. On the other hand, if the fractional part of nx is at least as large as the said margin, then the fractional part will “wrap around” to zero, and the result should be strictly less than the fractional part of $n_0x + y$. Or equivalently, if the fractional part of $n_0x + y$ is greater than or equal to the margin of nx , then adding them together will make the fractional part to strictly decrease, precisely by the amount of said margin. The lemma below describes this rigorously.

Lemma 5.1.

Let $n_0 \in \mathbb{Z}$ and $n \in \mathbb{Z}_{\geq 0}$. Then

$$\begin{aligned} & ((n_0 + n)x + y) - \lfloor (n_0 + n)x + y \rfloor \\ &= \begin{cases} ((n_0x + y) - \lfloor n_0x + y \rfloor) - (\lfloor nx \rfloor + 1 - nx) & \text{if } \lfloor nx \rfloor + 1 - nx \leq (n_0x + y) - \lfloor n_0x + y \rfloor, \\ ((n_0x + y) - \lfloor n_0x + y \rfloor) + (nx - \lfloor nx \rfloor) & \text{if } nx - \lfloor nx \rfloor < \lfloor n_0x + y \rfloor + 1 - (n_0x + y). \end{cases} \end{aligned}$$

In particular,

$$((n_0 + n)x + y) - \lfloor (n_0 + n)x + y \rfloor < (n_0x + y) - \lfloor n_0x + y \rfloor$$

holds if and only if

$$\lfloor nx \rfloor + 1 - nx \leq (n_0x + y) - \lfloor n_0x + y \rfloor.$$

Proof. If

$$\lfloor nx \rfloor + 1 - nx \leq (n_0x + y) - \lfloor n_0x + y \rfloor$$

holds, then

$$(n_0 + n)x + y \geq \lfloor n_0x + y \rfloor + \lfloor nx \rfloor + 1$$

and if it does not hold, then

$$(n_0 + n)x + y < \lfloor n_0x + y \rfloor + \lfloor nx \rfloor + 1.$$

Since $\lfloor (n_0 + n)x + y \rfloor$ is equal to either $\lfloor n_0x + y \rfloor + \lfloor nx \rfloor$ or $\lfloor n_0x + y \rfloor + \lfloor nx \rfloor + 1$, we should have

$$\lfloor (n + n_0)x + y \rfloor = \begin{cases} \lfloor n_0x + y \rfloor + \lfloor nx \rfloor + 1 & \text{if } \lfloor nx \rfloor + 1 - nx \leq (n_0x + y) - \lfloor n_0x + y \rfloor, \\ \lfloor n_0x + y \rfloor + \lfloor nx \rfloor & \text{if } \lfloor nx \rfloor + 1 - nx > (n_0x + y) - \lfloor n_0x + y \rfloor. \end{cases}$$

For the first case,

$$\begin{aligned} ((n_0 + n)x + y) - \lfloor (n_0 + n)x + y \rfloor &= ((n_0x + y) - \lfloor n_0x + y \rfloor) - (\lfloor nx \rfloor + 1 - nx) \\ &< (n_0x + y) - \lfloor n_0x + y \rfloor, \end{aligned}$$

and for the second case,

$$\begin{aligned} ((n_0 + n)x + y) - \lfloor (n_0 + n)x + y \rfloor &= ((n_0x + y) - \lfloor n_0x + y \rfloor) + (nx - \lfloor nx \rfloor) \\ &\geq (n_0x + y) - \lfloor n_0x + y \rfloor. \end{aligned}$$

□

Therefore, the first $n \in \mathbb{Z}_{>0}$ that makes the fractional part of $(n_0 + n)x + y$ strictly smaller than that of $n_0x + y$ is the smallest $n \in \mathbb{Z}_{>0}$ such that the $\lfloor nx \rfloor + 1 - nx$ is smaller than or equal to the current fractional part of $n_0x + y$. Let us call such n as n_1 . If we add n_1 to n_0 , then the fractional part decreases by $\lfloor n_1x \rfloor + 1 - n_1x$. Since it has decreased, no $n \in [1: n_1 - 1]$ can ever satisfy the condition

$$\lfloor nx \rfloor + 1 - nx \leq ((n_0 + n_1)x + y) - \lfloor (n_0 + n_1)x + y \rfloor,$$

because n_1 was chosen to be the smallest that satisfies a less strict inequality $\lfloor nx \rfloor + 1 - nx \leq (n_0x + y) - \lfloor n_0x + y \rfloor$. On the other hand, $n = n_1$ might still work out, that is, we may still have

$$\lfloor n_1x \rfloor + 1 - n_1x \leq ((n_0 + n_1)x + y) - \lfloor (n_0 + n_1)x + y \rfloor.$$

If that happens, then still n_1 is the smallest that satisfies the above, so $n = n_0 + 2n_1$ is the next smallest after $n_0 + n_1$ where $(nx + y) - \lfloor nx + y \rfloor$ decreases further than its previous minimum over $[n_0: n - 1]$. Now, note that adding n_1 to $n_0 + n_1$ decreases the fractional part by the exact same amount $\lfloor n_1x \rfloor + 1 - n_1x$. Hence, this trend will continue, that is, n_1 continues to be the smallest increment that further decreases the fractional part, until the fractional part finally becomes strictly smaller than $\lfloor n_1x \rfloor + 1 - n_1x$. Therefore, the exact number of times that we can reuse n_1 is given as

$$a = \left\lfloor \frac{(n_0x + y) - \lfloor n_0x + y \rfloor}{\lfloor n_1x \rfloor + 1 - n_1x} \right\rfloor,$$

and after adding n_1 for a times, finally we will need a different n to make the fractional part even smaller. So we replace n_0 by $n_0 + an_1$ and find the smallest n_1 that satisfies

$$\lfloor n_1x \rfloor + 1 - n_1x \leq (n_0x + y) - \lfloor n_0x + y \rfloor.$$

Then we can repeat this procedure until n_0 becomes bigger than n_{\max} .

Now, so far this feels like a kind of a greedy algorithm, and it is not immediately clear if this approach will really give us a true minimizer, let alone the smallest minimizer. To see it indeed does, let $m \in [n_{\min}: n_{\max}]$ be the smallest minimizer of $(nx + y) - \lfloor nx + y \rfloor$ over $[n_{\min}: n_{\max}]$. Note that because of Lemma 5.1, we must have either $m = n_0 = n_{\min}$ or

$$\lfloor (m - n_0)x + y \rfloor + 1 - ((m - n_0)x + y) \leq (n_0x + y) - \lfloor n_0x + y \rfloor.$$

If $m = n_0$ is the case, then we should have failed to find n_1 (because otherwise the fractional part will strictly decrease when we add n_1 to m), so we would get the correct minimizer n_0 . If $m > n_0$, then because of the minimality of n_1 , we must have $m - n_0 \geq n_1$, so m is

still greater than or equal to $n_0 + n_1$. In this case, since m is a minimizer, Lemma 5.1 still implies either $m = n_0 + n_1$ or

$$\lfloor (m - n_0 - n_1)x + y \rfloor + 1 - ((m - n_0 - n_1)x + y) \leq ((n_0 + n_1)x + y) - \lfloor (n_0 + n_1)x + y \rfloor.$$

Therefore, $m - n_0 - n_1$ should be greater than or equal to the smallest n such that $\lfloor nx \rfloor + 1 - nx$ is smaller than or equal to the right-hand side of the above inequality. Therefore, our iteration procedure should eventually exhaust the entire gap between m and n_0 , and thus reaches m .⁷ Consequently, we obtain the following result:

Corollary 5.2 (The smallest minimizer of the fractional part).

Let $n_0^{(0)} := n_{\min}$. For $i \in \mathbb{Z}_{>0}$, inductively define $n_0^{(i)} := n_0^{(i-1)} + a^{(i)}n_1^{(i)}$,

$$n_1^{(i)} := \min \left\{ n \in \mathbb{Z}_{>0} : \lfloor nx \rfloor + 1 - nx \leq (n_0^{(i-1)}x + y) - \lfloor n_0^{(i-1)}x + y \rfloor \right\}$$

and

$$a^{(i)} := \left\lfloor \frac{(n_0^{(i-1)}x + y) - \lfloor n_0^{(i-1)}x + y \rfloor}{\lfloor n_1^{(i)}x \rfloor + 1 - n_1^{(i)}x} \right\rfloor$$

when the set defining $n_1^{(i)}$ is not empty, and otherwise $n_0^{(i)} := n_1^{(i)} := \infty$. Let k be the largest nonnegative integer such that $n_0^{(k)} \leq n_{\max}$. Then

$$\min_{n \in [n_{\min} : n_{\max}]} \operatorname{argmin} ((nx + y) - \lfloor nx + y \rfloor) = n_0^{(k)} + \left\lfloor \frac{n_{\max} - n_0^{(k)}}{n_1^{(k+1)}} \right\rfloor n_1^{(k+1)}$$

with the convention $0 \cdot \infty = 0$.

Finding the maximizer can be done in a similar manner. In fact, Lemma 5.1 already gives all needed information. Starting with $n_0 = n_{\min}$, we find the smallest $n_1 \in \mathbb{Z}_{>0}$ satisfying

$$n_1x - \lfloor n_1x \rfloor < \lfloor n_0x + y \rfloor + 1 - (n_0x + y),$$

and add it to n_0 for a times, where

$$a = \left\lceil \frac{\lfloor n_0x + y \rfloor + 1 - (n_0x + y)}{n_1x - \lfloor n_1x \rfloor} \right\rceil - 1$$

is the smallest integer such that

$$n_1x - \lfloor n_1x \rfloor \geq \lfloor (n_0 + an_1)x + y \rfloor + 1 - ((n_0 + an_1)x + y)$$

holds. Indeed, the right-hand side must be equal to

$$1 - ((n_0x + y) - \lfloor n_0x + y \rfloor + a(n_1x - \lfloor n_1x \rfloor)),$$

so a is the smallest integer satisfying

$$a + 1 \geq \frac{\lfloor n_0x + y \rfloor + 1 - (n_0x + y)}{n_1x - \lfloor n_1x \rfloor}.$$

⁷Note that this exhaustion happens precisely because m is the *smallest* minimizer. Indeed, if m is not the smallest minimizer, then once our iteration procedure reaches to the smallest minimizer, then we can no longer find n that satisfies $\lfloor nx \rfloor + 1 - nx \leq (n_0x + y) - \lfloor n_0x + y \rfloor$, because if that were possible, then we can further decrease the fractional part.

Then, we replace n_0 by $n_0 + an_1$ and repeat this procedure until n_0 becomes larger than n_{\max} .

Of course, this breaks down if n_1x happens to be an integer. We claim that in that case, we have already reached the smallest maximizer. If the claim holds, then we can simply add as many multiple of n_1 as needed to n_0 to find the largest maximizer. (Note that in this case n_1 must be the smallest positive integer such that n_1x is an integer, and adding n_1 to n_0 should not change the fractional part of $n_0x + y$.)

To see why the claim holds, let $m \in [n_{\min} : n_{\max}]$ be the smallest maximizer as in the case of minimizer. Then since m is the smallest, there should be no $n \in [1 : m - n_{\min}]$ such that nx is an integer, because once we find such n , then $m - n$ must be a strictly smaller maximizer than m . Hence, the iteration procedure must exhaust the gap between n_{\min} and m before we find n_1 making n_1x into an integer. Consequently, we obtain the following result:

Corollary 5.3 (The largest maximizer of the fractional part).

Let $n_0^{(0)} := n_{\min}$. For $i \in \mathbb{Z}_{>0}$, inductively define $n_0^{(i)} := n_0^{(i-1)} + a^{(i)}n_1^{(i)}$,

$$n_1^{(i)} := \min \left\{ n \in \mathbb{Z}_{>0} : nx - \lfloor nx \rfloor < \left\lfloor \frac{n_0^{(i-1)}x + y}{n_1^{(i-1)}} \right\rfloor + 1 - (n_0^{(i-1)}x + y) \right\}$$

and

$$a^{(i)} := \left\lceil \frac{\left\lfloor \frac{n_0^{(i-1)}x + y}{n_1^{(i-1)}} \right\rfloor + 1 - (n_0^{(i-1)}x + y)}{n_1^{(i)}x - \lfloor n_1^{(i)}x \rfloor} \right\rceil - 1$$

when $n_1^{(i)}x$ is not an integer, and otherwise $n_0^{(i)} := \infty$. Let k be the largest nonnegative integer such that $n_0^{(k)} \leq n_{\max}$. Then

$$\max_{n \in [n_{\min} : n_{\max}]} \operatorname{argmax} ((nx + y) - \lfloor nx + y \rfloor) = n_0^{(k)} + \left\lfloor \frac{n_{\max} - n_0^{(k)}}{n_1^{(k+1)}} \right\rfloor n_1^{(k+1)}$$

with the convention $0 \cdot \infty = 0$.

Therefore, the problems of finding the smallest minimizer and the largest maximizer of $(nx + y) - \lfloor nx + y \rfloor$ becomes the problems of finding the smallest $n \in \mathbb{Z}_{>0}$ such that

$$\lfloor nx \rfloor + 1 - nx \leq \tau \quad \text{with} \quad \tau := (n_0x + y) - \lfloor n_0x + y \rfloor \in [0, 1)$$

and

$$nx - \lfloor nx \rfloor < \tau \quad \text{with} \quad \tau := \left\lfloor \frac{n_0x + y}{n_1} \right\rfloor + 1 - (n_0x + y) \in (0, 1],$$

respectively.

Recall from Theorem 2.4 that finding extremizers of $nx - \lfloor nx \rfloor$ is equivalent to finding best rational approximations of x from below/above, which, by Theorem 4.11, can be done by enumerating semiconvergents of x with increasingly larger denominators.

Let us give some more details here. As in Section 3 and Section 4, let $x = [a_0; a_1, \dots]$ and $\frac{p_i}{q_i}$ be the i th convergent. For the problem $\lfloor nx \rfloor + 1 - nx \leq \tau$, let $n = n_1$ be the smallest positive integer satisfying the inequality, if any. Then n_1 is the unique maximizer of $nx - \lfloor nx \rfloor$ over $n \in [1 : n_1]$. By Theorem 2.4 and Theorem 4.11, n_1 must be the denominator of a proper odd semiconvergent.⁸

⁸In particular, n_1 cannot be the denominator of x , if x is rational. Of course, in this case $\lfloor n_1x \rfloor + 1 - n_1x = 1 > \tau$.

Hence, we find the first odd convergent $\frac{p_i}{q_i}$ such that $p_i - q_i x \leq \tau$ holds. Note that $p_i = \lceil q_i x \rceil = \lfloor q_i x \rfloor + 1$ holds unless $x = \frac{p_i}{q_i}$, thus except for that case, $p_i - q_i x \leq \tau$ is equivalent to $\lfloor q_i x \rfloor + 1 - q_i x \leq \tau$.

Suppose that we found such an odd convergent $\frac{p_i}{q_i}$ that is not equal to x . Then this means that the smallest n we are looking for must be of the form $n = q_i - sq_{i-1}$ for some integer $s \in [0, a_i)$. By construction, we know

$$(p_i - q_i x) + a_i(q_{i-1}x - p_{i-1}) = p_{i-2} - q_{i-2}x > \tau \geq p_i - q_i x.$$

In this case, s must be the largest integer such that

$$\tau \geq (p_i - q_i x) + s(q_{i-1}x - p_{i-1}),$$

that is,

$$s = \left\lfloor \frac{\tau - (p_i - q_i x)}{q_{i-1}x - p_{i-1}} \right\rfloor.$$

Then $n = q_i - sq_{i-1}$ is the desired output, and additionally we know

$$\lfloor nx \rfloor + 1 - nx = (p_i - q_i x) + s(q_{i-1}x - p_{i-1}).$$

Next, suppose that there is no odd convergent $\frac{p_i}{q_i}$ such that $p_i - q_i x \leq \tau$. This can happen if $x = \frac{p_j}{q_j}$ is its own even convergent. In this case, $\frac{p_{j-1}}{q_{j-1}}$ is the proper odd semiconvergent of x with the largest denominator, and since it failed to satisfy $p_{j-1} - q_{j-1}x \leq \tau$, we conclude that there is no n satisfying $\lfloor nx \rfloor + 1 - nx \leq \tau$.

Next, suppose that $x = \frac{p_i}{q_i}$ itself is the first convergent satisfying $p_i - q_i x \leq \tau$. As noted earlier, this does not mean $\lfloor q_i x \rfloor + 1 - q_i x \leq \tau$; rather, the left-hand side is equal to 1. In this case, we may need to look at proper odd semiconvergents of the form $\frac{p_i - sp_{i-1}}{q_i - sq_{i-1}}$ for $s \in (0, a_i)$, which are better rational approximations of x from above than the last odd convergent $\frac{p_{i-2}}{q_{i-2}}$. Hence, we try the same formula as before:

$$s = \left\lfloor \frac{\tau - (p_i - q_i x)}{q_{i-1}x - p_{i-1}} \right\rfloor.$$

In this case, if we get $s = 0$, then we conclude there is no n satisfying $\lfloor nx \rfloor + 1 - nx \leq \tau$, and otherwise, $n = q_i - sq_{i-1}$ is the desired output and

$$\lfloor nx \rfloor + 1 - nx = (p_i - q_i x) + s(q_{i-1}x - p_{i-1}) = s(q_{i-1}x - p_{i-1}).$$

Now we look at the problem $nx - \lfloor nx \rfloor < \tau$. Similarly, the smallest $n = n_1$ satisfying the inequality should be the unique minimizer of $nx - \lfloor nx \rfloor$ over $n \in [1 : n_1]$, which, by Theorem 2.4 and Theorem 4.11, is either the denominator of x (if x is rational) or the denominator of a proper even semiconvergent of x .

Hence, we find the first even convergent $\frac{p_i}{q_i}$ such that $q_i x - p_i < \tau$ holds. If there is no such even convergent, then that means $x = \frac{p_j}{q_j}$ is an odd convergent of itself. In this case, the last proper even semiconvergent is $\frac{p_{j-1}}{q_{j-1}}$ which did not satisfy $q_{j-1}x - p_{j-1} < \tau$, thus $n = q_j$ is the smallest n satisfying $nx - \lfloor nx \rfloor < \tau$.

Next, suppose that we found the first even convergent $\frac{p_i}{q_i}$ satisfying $q_i x - p_i < \tau$. If $i = 0$, then $n = 1$ is the desired output. Otherwise, we look at even semiconvergents of the form $\frac{p_i - sp_{i-1}}{q_i - sq_{i-1}}$ for $s \in [0, a_i)$. Note that there is a possibility of $x = \frac{p_i}{q_i}$ so that $\frac{p_i}{q_i}$ is not a proper

semiconvergent, but that is okay because $n = q_i$ is the unique minimizer even in that case anyway. Now, we have

$$(q_i x - p_i) + a_i(p_{i-1} - q_{i-1}x) = q_{i-2}x - p_{i-2} \geq \tau > q_i x - p_i,$$

so s must be the largest integer satisfying

$$(q_i x - p_i) + s(p_{i-1} - q_{i-1}x) < \tau,$$

that is,

$$s = \left\lceil \frac{\tau - (q_i x - p_i)}{p_{i-1} - q_{i-1}x} \right\rceil - 1.$$

Then $n = q_i - s q_{i-1}$ is the desired output, and additionally we know

$$n x - \lfloor n x \rfloor = (q_i x - p_i) + s(p_{i-1} - q_{i-1}x).$$

Note that these formulas also work for the case $i = 0$, so in fact we do not need to make a special branch for $i = 0$.

In practice, since one problem only cares about even semiconvergents and the other problem only cares about odd semiconvergents, we can solve both at the same time in a single run. More specifically, we have the following algorithm:

Algorithm 5.4 (Finding extremizers of $(n x + y) - \lfloor n x + y \rfloor$).

1. Set $n_{0,\min} \leftarrow n_{\min}$, $n_{0,\max} \leftarrow n_{\min}$, $\text{found}_{\min} \leftarrow \text{false}$, $\text{found}_{\max} \leftarrow \text{false}$, and $\frac{p}{q} \leftarrow \frac{1}{0}$, the -1 st convergent of x .
2. Set $\text{terminated} \leftarrow \text{false}$.
3. Set $\tau_{\min} \leftarrow (n_{0,\min} x + y) - \lfloor n_{0,\min} x + y \rfloor$ and $\tau_{\max} \leftarrow \lfloor n_{0,\max} x + y \rfloor + 1 - (n_{0,\max} x + y)$.
4. Check if there are further convergents of x . If so, let $\frac{p}{q}$ be the next convergent of x ; in this case, $\frac{p}{q}$ must be an even convergent of x . If not, set $\text{terminated} \leftarrow \text{true}$; in this case, x is equal to the last convergent $\frac{p}{q}$ and it is an odd convergent.
5. If found_{\max} is false,
 - (a) If terminated is false,
 - i. Check if $q x - p < \tau_{\max}$ holds. If it doesn't, we move to the next convergent, so go to Step 6.
 - ii. If it does, then we find the first semiconvergent between the current and the previous even convergents that satisfies the inequality. Set

$$s \leftarrow \left\lceil \frac{\tau_{\max} - (q x - p)}{p' - q' x} \right\rceil - 1, \quad n_1 \leftarrow q - s q' \quad \text{and} \quad \delta \leftarrow (q x - p) + s(p' - q' x),$$

where $\frac{p'}{q'}$ is the previous convergent.

- iii. Check if $\delta = 0$. If that is the case, then go to Step 5(b)-i.
- iv. Otherwise, set $a \leftarrow \lceil \tau_{\max} / \delta \rceil - 1$.
- v. If $n_{0,\max} + a n_1 \leq n_{\max}$, then set $n_{0,\max} \leftarrow n_{0,\max} + a n_1$, recompute τ_{\max} accordingly, and then go back to Step 5(a)-i.
- vi. Otherwise, go to Step 5(b)-i.

- (b) If terminated is true, set $n_1 \leftarrow q$; in this case, n_1 is the smallest n satisfying $nx - \lfloor nx \rfloor < \tau$, and we have $n_1x - \lfloor n_1x \rfloor = 0$.

i. Set

$$n_{0,\max} \leftarrow n_{0,\max} + \left\lfloor \frac{n_{\max} - n_{0,\max}}{n_1} \right\rfloor n_1.$$

ii. Set $\text{found}_{\max} \leftarrow \text{true}$ and go to Step 6.

6. Check if there are further convergents of x . If so, let $\frac{p}{q}$ be the next convergent of x ; in this case, $\frac{p}{q}$ must be an odd convergent of x . If not, set $\text{terminated} \leftarrow \text{true}$; in this case, x is equal to the last convergent $\frac{p}{q}$ and it is an even convergent.

7. If found_{\min} is false,

(a) If terminated is false,

i. Check if $p - qx \leq \tau_{\min}$ holds. If it doesn't, we move to the next convergent. Go to the Step 8.

ii. If it does, then we find the first semiconvergent between the current and the previous odd convergents that satisfies the inequality. Set

$$s \leftarrow \left\lfloor \frac{\tau_{\min} - (p - qx)}{q'x - p'} \right\rfloor, \quad n_1 \leftarrow q - sq' \quad \text{and} \quad \delta \leftarrow (p - qx) + s(q'x - p')$$

where $\frac{p'}{q'}$ is the previous convergent.

- iii. Check if $\delta = 0$. If that is the case, then we conclude there is no n satisfying $\lfloor nx \rfloor + 1 - nx \leq \tau_{\min}$, so go to Step 6(b)-i.
- iv. Otherwise, set $a \leftarrow \lfloor \tau_{\min} / \delta \rfloor$.
- v. If $n_{0,\min} + an_1 \leq n_{\max}$, set $n_{0,\min} \leftarrow n_{0,\min} + an_1$, recompute τ_{\min} accordingly, and then go back to Step 7(a)-i.
- vi. Otherwise, set

$$n_{0,\min} \leftarrow n_{0,\min} + \left\lfloor \frac{n_{\max} - n_{0,\min}}{n_1} \right\rfloor n_1$$

and go to Step 6(b)-i.

- (b) If terminated is true, then there is no n satisfying $\lfloor nx \rfloor + 1 - nx \leq \tau_{\min}$.

i. Set $\text{found}_{\min} \leftarrow \text{true}$, and then go to Step 8.

8. If either of found_{\min} or found_{\max} is false, then go back to Step 4. Otherwise, $n_{0,\min}$ is the smallest minimizer and $n_{0,\max}$ is the largest maximizer.

At the end of Section 4, I remarked that we will retain the viewpoint of identifying all real numbers with their continued fraction expansions. However, note that in this viewpoint, it is not immediately clear how we can actually compute s and a in the above algorithm unless x is a rational number (so that the computations can be directly done in an arbitrary-precision exact rational arithmetic).

Note that it is not a good idea to use floating-point approximations of real numbers either, because the kind of computations we do are exactly the ones for which working with floating-point numbers can be tricky: subtracting and dividing numbers of comparable sizes. At least for me, it sounds very challenging to find out the exact error bound that will ensure the correct computation, for any given $[n_{\min} : n_{\max}]$.

To resolve this issue, we will develop two algorithms. The first algorithm, called *Gosper's algorithm*, which we develop in Section 12 after giving the explanation of the main algorithm of this whole note for establishing $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$, is an algorithm for computing the continued fraction expansion of any number obtained by combining two numbers x and y , with known continued fraction expansions, in a certain way. The way we combine x and y is general enough to cover all computations we need in Algorithm 5.4, so that those can be done with arbitrary-precision exact rational arithmetic.

However, executing this algorithm everytime in all of the iterations in Algorithm 5.4 would be very inefficient, which is where the desire for having the second algorithm arises: it is an algorithm that computes good enough binary approximations of x and y of the form $\frac{m}{2^k}$ and $\frac{s}{2^k}$, where m, s are integers and k is a nonnegative integer, such that Algorithm 5.4 is guaranteed to give the same output if we replace (x, y) by $(\frac{m}{2^k}, \frac{s}{2^k})$. If we can come up with such k, m, s , then arbitrary-precision exact rational arithmetic is enough to execute all the needed computations. This second algorithm is actually relying on the first algorithm, but the first algorithm only needs to be called just several, a bounded amount of times.

We will develop this binary approximation algorithm in Section 8 and Section 9. Before we get there, we have to prepare ourselves with some more background works which we do in Section 6 and Section 7.

6 Lemmas for expanding the floor function

We begin with several simple lemmas about extremizers of $(nx + y) - \lfloor nx + y \rfloor$. Throughout this section, let x, y be any real numbers and let $n_{\min} \leq n_{\max}$ be any integers.

Lemma 6.1.

Let $n_0 \in \operatorname{argmin}_{n \in [n_{\min} : n_{\max}]} ((nx + y) - \lfloor nx + y \rfloor)$. Then for any $n \in [n_{\min} - n_0 : n_{\max} - n_0]$,

$$\lfloor (n_0 + n)x + y \rfloor = \lfloor n_0x + y \rfloor + \lfloor nx \rfloor.$$

Proof. By the assumption on n_0 , we have

$$((n_0 + n)x + y) - \lfloor (n_0 + n)x + y \rfloor \geq (n_0x + y) - \lfloor n_0x + y \rfloor,$$

thus

$$\lfloor (n_0 + n)x + y \rfloor \leq \lfloor n_0x + y \rfloor + nx < \lfloor n_0x + y \rfloor + \lfloor nx \rfloor + 1.$$

Since $\lfloor (n_0 + n)x + y \rfloor$ should be equal to either $\lfloor n_0x + y \rfloor + \lfloor nx \rfloor$ or $\lfloor n_0x + y \rfloor + \lfloor nx \rfloor + 1$, the only possibility is

$$\lfloor (n_0 + n)x + y \rfloor = \lfloor n_0x + y \rfloor + \lfloor nx \rfloor.$$

□

The intuition behind the above lemma is that, since n_0 achieves the minimum of the fractional part of $nx + y$, adding any nx into $n_0x + y$ should not make its fractional part to “wrap around”, because at the exact moment of the wrapping, the fractional part will become strictly smaller than that of $n_0x + y$.

Lemma 6.2.

Let $n_0 \in \operatorname{argmax}_{n \in [n_{\min} : n_{\max}]} ((nx + y) - \lfloor nx + y \rfloor)$. Then for any $n \in [n_0 - n_{\max} : n_0 - n_{\min}]$,

$$\lfloor (n_0 - n)x + y \rfloor = \lfloor n_0x + y \rfloor - \lfloor nx \rfloor.$$

Proof. By the assumption on n_0 , we have

$$((n_0 - n)x + y) - \lfloor (n_0 - n)x + y \rfloor \leq (n_0x + y) - \lfloor n_0x + y \rfloor,$$

thus

$$\lfloor (n_0 - n)x + y \rfloor \geq \lfloor n_0x + y \rfloor - nx > \lfloor n_0x + y \rfloor - \lfloor nx \rfloor - 1.$$

Since $\lfloor (n_0 - n)x + y \rfloor$ should be equal to either $\lfloor n_0x + y \rfloor - \lfloor nx \rfloor$ or $\lfloor n_0x + y \rfloor - \lfloor nx \rfloor - 1$, the only possibility is

$$\lfloor (n_0 - n)x + y \rfloor = \lfloor n_0x + y \rfloor - \lfloor nx \rfloor.$$

□

Similarly to the previous lemma, subtracting any nx from $n_0x + y$ should not make its fractional part to wrap around, because at the exact moment of the wrapping, the fractional part will become strictly larger than that of $n_0x + y$.

Lemma 6.3.

Let $n_0 = \min \arg \min_{n \in [n_{\min} : n_{\max}]} ((nx + y) - \lfloor nx + y \rfloor)$. Then for any $n \in [1 : n_0 - n_{\min}]$,

$$\lfloor (n_0 - n)x + y \rfloor = \lfloor n_0x + y \rfloor - \lfloor nx \rfloor - 1.$$

Proof. By Lemma 6.1, we know

$$\lfloor (n_0 - n)x + y \rfloor = \lfloor n_0x + y \rfloor + \lfloor (-n)x \rfloor = \lfloor n_0x + y \rfloor - \lceil nx \rceil.$$

Note that if nx is integer, then we must have

$$((n_0 - n)x + y) - \lfloor (n_0 - n)x + y \rfloor = (n_0x + y) - \lfloor n_0x + y \rfloor,$$

so by the minimality of n_0 , nx is never an integer. Therefore, $\lceil nx \rceil = \lfloor nx \rfloor + 1$, so we get the desired conclusion. □

The intuition for the above lemma is that, if n_0 minimizes the fractional part of $nx + y$, then subtracting any nx from $n_0x + y$ should always make the fractional part wrap around, because otherwise the fractional part of $(n_0 - n)x + y$ will become strictly smaller than that of $n_0x + y$. Except when nx is an integer, in which case the fractional part stays the same, but since we have chosen the *smallest* minimizer, there is no n such that nx is an integer and $n_0 - n$ is still in the range $[n_{\min} : n_{\max}]$.

Lemma 6.4.

Let $n_0 = \max \arg \max_{n \in [n_{\min} : n_{\max}]} ((nx + y) - \lfloor nx + y \rfloor)$. Then for any $n \in [1 : n_{\max} - n_0]$,

$$\lfloor (n_0 + n)x + y \rfloor = \lfloor n_0x + y \rfloor + \lfloor nx \rfloor + 1.$$

Proof. By Lemma 6.2, we know

$$\lfloor (n_0 + n)x + y \rfloor = \lfloor n_0x + y \rfloor - \lfloor (-n)x \rfloor = \lfloor n_0x + y \rfloor + \lceil nx \rceil.$$

Note that if nx is integer, then we must have

$$((n_0 + n)x + y) - \lfloor (n_0 + n)x + y \rfloor = (n_0x + y) - \lfloor n_0x + y \rfloor,$$

so by the maximality of n_0 , nx is never an integer. Therefore, $\lceil nx \rceil = \lfloor nx \rfloor + 1$, so we get the desired conclusion. □

Similarly to the previous lemma, adding any nx to n_0x+y should make the fractional part wrap around (unless nx is an integer), because otherwise the fractional part of $(n_0+n)x+y$ should be larger than that of n_0x+y .

The following special cases of Lemma 6.1 and Lemma 6.4 will be also useful for our purpose.

Lemma 6.5.

Let $n_0 \in \operatorname{argmin}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor)$. Then for any $k, n \in \mathbb{Z}_{\geq 0}$ such that $kn_0 + n \leq n_{\max}$,

$$\lfloor (kn_0 + n)x \rfloor = k \lfloor n_0x \rfloor + \lfloor nx \rfloor.$$

Proof. We use induction on k . The base case $k = 0$ is trivial, so we prove the induction step. By Lemma 6.1, we have

$$\lfloor (kn_0 + n)x \rfloor = \lfloor n_0x \rfloor + \lfloor ((k-1)n_0 + n)x \rfloor,$$

thus by applying the induction hypothesis to expand $\lfloor ((k-1)n_0 + n)x \rfloor$, we get the desired conclusion. \square

Lemma 6.6.

Let $n_0 = \max \operatorname{argmax}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor)$. Then for any $k, n \in \mathbb{Z}_{\geq 0}$ such that $kn_0 + n \leq n_{\max}$ and $n \neq 0$,

$$\lfloor (kn_0 + n)x \rfloor = k \lfloor n_0x \rfloor + 1 + \lfloor nx \rfloor.$$

Proof. We use induction on k . The base case $k = 0$ is trivial, so we prove the induction step. By Lemma 6.4, we have

$$\lfloor (kn_0 + n)x \rfloor = \lfloor n_0x \rfloor + \lfloor ((k-1)n_0 + n)x \rfloor + 1,$$

thus by applying the induction hypothesis to expand $\lfloor ((k-1)n_0 + n)x \rfloor$, we get the desired conclusion. \square

7 Extremizers of $\frac{\lfloor nx \rfloor}{n}$ and $\frac{\lfloor nx \rfloor + 1}{n}$

In this section, we solve a special case of our main problem: we establish the necessary and sufficient condition for ξ to have $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ for all $n \in [1: n_{\max}]$. This is also one of the fundamental building blocks for solving our main problem.

In fact, I already described a solution for this special case in detail in [5], but here I reproduced it for convenience of both the readers and myself. Actually, some of the theorems stated here are slight improvements of the corresponding theorems in [5].

Theorem 7.1 (Extremizers of $\frac{\lfloor nx \rfloor}{n}$ and $\frac{\lfloor nx \rfloor + 1}{n}$).

Let x be a real number and n_{\max} a positive integer. Then we have the following relations between extremizers of $nx - \lfloor nx \rfloor$, $\frac{\lfloor nx \rfloor}{n}$, and $\frac{\lfloor nx \rfloor + 1}{n}$:

1. Let $n_0 := \min \operatorname{argmin}_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor)$, then

$$\operatorname{argmax}_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor}{n} = \{kn_0 : k \in \mathbb{Z}_{>0}, kn_0 \leq n_{\max}\}$$

and $\frac{\lfloor n_0x \rfloor}{n_0}$ is a best rational approximation of x from below in its reduced form.

2. Let $n_0 := \max \arg\max_{n \in [1: n_{\max}]} (nx - \lfloor nx \rfloor)$, then

$$\arg\min_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor + 1}{n} = \{kn_0 : k \in \mathbb{Z}_{>0}, kn_0 \leq n_{\max}\}$$

and $\frac{\lfloor n_0 x \rfloor + 1}{n_0}$ is in its reduced form. If $nx \notin \mathbb{Z}$ for all $n \in [1: n_{\max}]$, then $\frac{\lfloor n_0 x \rfloor + 1}{n_0}$ is a best rational approximation of x from above.

Proof. 1. Theorem 2.4 immediately shows that $\frac{\lfloor n_0 x \rfloor}{n_0}$ is a best rational approximation of x from below in its reduced form. Choose any

$$n_1 \in \arg\max_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor}{n}$$

and let $\frac{p}{q}$ be the reduced form of $\frac{\lfloor n_1 x \rfloor}{n_1}$. Then for any rational number $\frac{a}{b}$ with $0 < b \leq q$ and $\frac{a}{b} \leq x$, we have

$$\frac{a}{b} \leq \frac{\lfloor bx \rfloor}{b} \leq \frac{\lfloor n_1 x \rfloor}{n_1} = \frac{p}{q} \leq x,$$

thus $\frac{p}{q}$ is a best rational approximation of x from below.

Hence, if $n_0 \leq q$, then

$$qx - p \leq n_0 x - \lfloor n_0 x \rfloor$$

holds, but by definition of n_0 , this must be an equality. Then by Theorem 2.4, q should be a multiple of n_0 .

On the other hand, if $n_0 \geq q$, then

$$\frac{\lfloor n_0 x \rfloor}{n_0} \geq \frac{p}{q}$$

holds since $\frac{\lfloor n_0 x \rfloor}{n_0}$ is a best rational approximation of x from below in its reduced form, thus by definition of $\frac{p}{q}$, this must be an equality. Since both sides are in their reduced form, we conclude $q = n_0$.

Therefore, for any cases, we conclude that n_1 is a multiple of n_0 . Conversely, any multiple of n_0 bounded above by n_{\max} is a maximizer, because Lemma 6.5 shows that for any $k \in \mathbb{Z}_{>0}$ such that $kn_0 \leq n_{\max}$, $\frac{\lfloor kn_0 x \rfloor}{kn_0} = \frac{\lfloor n_0 x \rfloor}{n_0}$.

2. We claim that $\frac{\lfloor n_0 x \rfloor + 1}{n_0}$ is in its reduced form and any minimizer of $\frac{\lfloor nx \rfloor + 1}{n}$ is a multiple of n_0 .

First, suppose that nx is an integer for some $n \in [1: n_{\max}]$, then Theorem 2.4 shows that n_0 is the largest integer in $[1: n_{\max}]$ with $n_0 p \equiv -1 \pmod{q}$. Choose any

$$n_1 \in \arg\min_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor + 1}{n}.$$

Then by rearranging the inequality

$$\frac{\lfloor n_1 x \rfloor + 1}{n_1} \leq \frac{\lfloor n_0 x \rfloor + 1}{n_0},$$

we get

$$n_0 ((\lfloor n_1 x \rfloor + 1)q - n_1 p) \leq n_1 ((\lfloor n_0 x \rfloor + 1)q - n_0 p).$$

Since

$$n_0 (([n_1x] + 1)q - n_1p) \equiv n_1 (([n_0x] + 1)q - n_0p) \pmod{q},$$

there exists a nonnegative integer e such that

$$n_1 (([n_0x] + 1)q - n_0p) = n_0 (([n_1x] + 1)q - n_1p) + eq.$$

Now, note that

$$([n_0x] + 1)q - n_0p = \left(\frac{n_0p - (q - 1)}{q} + 1 \right) q - n_0p = 1,$$

which in particular shows that $[n_0x] + 1$ and n_0 are coprime. Similarly, we have

$$([n_1x] + 1)q - n_1p = \left(\frac{n_1p - (n_1p \bmod q)}{q} + 1 \right) q - n_1p = q - (n_1p \bmod q) \geq 1,$$

thus we obtain

$$n_1 = (q - (n_1p \bmod q)) n_0 + eq \geq n_0 + eq,$$

but by definition of n_0 , we must have $n_0 + q > n_{\max}$. Hence, $e = 0$ follows since $n_1 \leq n_{\max}$, which shows $n_1 = (q - (n_1p \bmod q)) n_0$.

Next, suppose that nx is not an integer for all $n \in [1 : n_{\max}]$, then Theorem 2.4 shows that $\frac{[n_0x] + 1}{n_0}$ is a best rational approximation of x from above in its reduced form. Choose any

$$n_1 \in \operatorname{argmin}_{n \in [1 : n_{\max}]} \frac{[nx] + 1}{n}$$

and let $\frac{p}{q}$ be the reduced form of $\frac{[n_1x] + 1}{n_1}$. Then for any rational number $\frac{a}{b}$ with $0 < b \leq q$ and $\frac{a}{b} \geq x$, we have

$$\frac{a}{b} \geq \frac{[bx]}{b} = \frac{[bx] + 1}{b} \geq \frac{[n_1x] + 1}{n_1} = \frac{p}{q} \geq x,$$

thus $\frac{p}{q}$ is a best rational approximation of x from above.

Hence, if $n_0 \leq q$, then

$$p - qx \leq [n_0x] + 1 - n_0x$$

holds, but since Theorem 2.4 shows that n_0 is the unique minimizer of $[nx] - nx = [nx] + 1 - nx = 1 - (nx - [nx])$, we conclude that the above is an equality and $q = n_0$.

On the other hand, if $n_0 \geq q$, then

$$\frac{[n_0x] + 1}{n_0} \leq \frac{p}{q}$$

holds since $\frac{[n_0x] + 1}{n_0}$ is a best rational approximation of x from above in its reduced form, thus by definition of $\frac{p}{q}$, this must be an equality. Since both sides are in their reduced form, we conclude $q = n_0$.

Therefore, for any cases, we conclude that $\frac{[n_0x] + 1}{n_0}$ is in its reduced form and n_1 is a multiple of n_0 , as claimed. Finally, any multiple of n_0 bounded above by n_{\max} is a minimizer, because Lemma 6.6 shows that for any $k \in \mathbb{Z}_{>0}$ such that $kn_0 \leq n_{\max}$,

$$\frac{[kn_0x] + 1}{kn_0} = \frac{(k - 1)([n_0x] + 1) + [n_0x] + 1}{kn_0} = \frac{[n_0x] + 1}{n_0}.$$

□

Remark 7.2.

Somewhat counterintuitively, the set

$$\operatorname{argmin}_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor + 1}{n} = \{kn_0 : k \in \mathbb{Z}_{>0}, kn_0 \leq n_{\max}\}$$

does not need to be consisting of only one element, regardless of if nx can be an integer or not. For instance, if $x = \frac{1}{7}$ and $n_{\max} = 12$, then the set above is $\{6, 12\}$, because 6 is the only element in $[1: n_{\max}]$ whose remainder is 6 when divided by 7. Indeed, both 6 and 12 minimizes $\frac{\lfloor nx \rfloor + 1}{n}$, with the minimum value $\frac{1}{6}$. For another instance, let $x = \frac{2}{7}$ and $n_{\max} = 6$, then the set above is $\{3, 6\}$. Indeed, both 3 and 6 minimizes $\frac{\lfloor nx \rfloor + 1}{n}$, with the minimum value $\frac{1}{3}$.

As a corollary of the above theorem, we get the following result which we use in Section 10.

Corollary 7.3.

Let x be a real number and n_{\max} be a positive integer. Then we have the followings.

1. Let $n_0 \in \operatorname{argmax}_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor}{n}$, then for any $n \in [0: n_{\max} - n_0]$,
$$\lfloor (n_0 + n)x \rfloor = \lfloor n_0 x \rfloor + \lfloor nx \rfloor.$$
2. Let $n_0 \in \operatorname{argmin}_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor + 1}{n}$, then for any $n \in [1: n_{\max} - n_0]$,
$$\lfloor (n_0 + n)x \rfloor = \lfloor n_0 x \rfloor + \lfloor nx \rfloor + 1.$$

Proof. 1. By Theorem 7.1, we know $n_0 = kn'_0$ for some $k \in \mathbb{Z}_{>0}$ where n'_0 is a minimizer of $nx - \lfloor nx \rfloor$ for $n \in [1: n_{\max}]$. Hence, for any $n \in [0: n_{\max} - n_0]$, Lemma 6.5 shows

$$\lfloor (n_0 + n)x \rfloor = \lfloor (kn'_0 + n)x \rfloor = k \lfloor n'_0 x \rfloor + \lfloor nx \rfloor = \lfloor kn'_0 x \rfloor + \lfloor nx \rfloor = \lfloor n_0 x \rfloor + \lfloor nx \rfloor.$$

2. By Theorem 7.1, we know $n_0 = kn'_0$ for some $k \in \mathbb{Z}_{>0}$ where n'_0 is the largest maximizer of $nx - \lfloor nx \rfloor$ for $n \in [1: n_{\max}]$. Hence, for any $n \in [1: n_{\max} - n_0]$, Lemma 6.6 shows

$$\begin{aligned} \lfloor (n_0 + n)x \rfloor &= \lfloor (kn'_0 + n)x \rfloor \\ &= k(\lfloor n'_0 x \rfloor + 1) + \lfloor nx \rfloor \\ &= (k-1)(\lfloor n'_0 x \rfloor + 1) + \lfloor n'_0 x \rfloor + \lfloor nx \rfloor + 1 \\ &= \lfloor kn'_0 x \rfloor + \lfloor nx \rfloor + 1 \\ &= \lfloor n_0 x \rfloor + \lfloor nx \rfloor + 1. \end{aligned}$$

□

Theorem 7.1 immediately solves the aforementioned special case of our main problem:

Theorem 7.4.

Let x be a real number and n_{\max} a positive integer. Then for a real number ξ , we have the followings.

1. If x is rational and its reduced form $\frac{p}{q}$ satisfies $q \leq n_{\max}$, then we have

$$\lfloor nx \rfloor = \lfloor n\xi \rfloor$$

for all $n \in [1: n_{\max}]$ if and only if

$$x \leq \xi < x + \frac{1}{vq}$$

holds, where v is the largest integer such that $vp \equiv -1 \pmod{q}$ and $v \leq n_{\max}$.

2. If x is either irrational or rational whose reduced denominator is strictly larger than n_{\max} , then we have

$$\lfloor nx \rfloor = \lfloor n\xi \rfloor$$

for all $n \in [1: n_{\max}]$ if and only if

$$\frac{p_*}{q_*} \leq \xi < \frac{p^*}{q^*}$$

holds, where $\frac{p_*}{q_*}, \frac{p^*}{q^*}$ are the best rational approximations of x from below and above, respectively, with the largest denominators $q_*, q^* \leq n_{\max}$.

Proof. Since having the equality $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ for all $n \in [1: n_{\max}]$ is equivalent to having the inequality

$$\max_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor}{n} \leq \xi < \min_{n \in [1: n_{\max}]} \frac{\lfloor nx \rfloor + 1}{n},$$

the conclusion immediately follow from Theorem 7.1 and Theorem 2.4. \square

Remark 7.5.

When x is rational and its reduced form $\frac{p}{q}$ satisfies $q \leq n_{\max}$, we have

$$x + \frac{1}{vq} = \frac{p}{q} + \frac{1}{vq} = \frac{vp + 1}{vq} = \frac{(vp + 1)/q}{v}.$$

By the definition of v , $(vp + 1)/q$ is an integer, and it is coprime to v since $\left(\frac{vp+1}{q}\right)q - vp = 1$, which means that $\frac{(vp+1)/q}{v}$ is the reduced form of the rational number $x + \frac{1}{vq}$. On the other hand, when either x is irrational or rational whose reduced denominator is strictly larger than n_{\max} , then q^* is the largest integer such that $q^*p_* \equiv -1 \pmod{q_*}$ and $q^* \leq n_{\max}$, because $\frac{p_*}{q_*}$ must appear in the right-hand side of the range of ξ given by the first case of Theorem 7.4 when we apply the theorem to $x \leftarrow \frac{p_*}{q_*}$, as we know by Theorem 7.4 that $\lfloor nx \rfloor = \left\lfloor \frac{np_*}{q_*} \right\rfloor$ holds for all $n \in [1: n_{\max}]$. Furthermore, applying the identity $\left(\frac{vp+1}{q}\right)q - vp = 1$ gives us $p^*q_* - q^*p_* = 1$, or equivalently, $\frac{p^*}{q^*} - \frac{p_*}{q_*} = \frac{1}{q^*q_*}$.

One of the important consequences of having $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ for all $n \in [1: n_{\max}]$ is that the fractional parts of nx and $n\xi$ should be ordered in the same way. This is also one of the results which I already have shown in [5]. The algorithm for finding a good enough approximation to use in Algorithm 5.4, which we will develop in Section 8, is crucially relying on this result.

Theorem 7.6 (Correspondence of fractional part ordering).

Let x, ξ be real numbers and n_{\max} a positive integer such that

$$\lfloor nx \rfloor = \lfloor n\xi \rfloor$$

holds for all $n \in [1: n_{\max}]$. Then we have the followings.

1. If x is rational and its reduced form $\frac{p}{q}$ satisfies $q \leq n_{\max}$, then for any $n_1, n_2 \in [1: n_{\max}]$,

$$n_1x - \lfloor n_1x \rfloor < n_2x - \lfloor n_2x \rfloor$$

or

$$n_1x - \lfloor n_1x \rfloor = n_2x - \lfloor n_2x \rfloor \quad \text{and} \quad n_1 < n_2$$

implies

$$n_1\xi - \lfloor n_1\xi \rfloor \leq n_2\xi - \lfloor n_2\xi \rfloor,$$

with possibly the equality only when $\xi = x$.

2. If x is either irrational or is rational whose reduced denominator is strictly larger than n_{\max} , then for any $n_1, n_2 \in [1: n_{\max}]$,

$$n_1x - \lfloor n_1x \rfloor < n_2x - \lfloor n_2x \rfloor$$

implies

$$n_1\xi - \lfloor n_1\xi \rfloor \leq n_2\xi - \lfloor n_2\xi \rfloor,$$

with possibly the equality only when $n_1 < n_2$ and $\xi = \frac{p_*}{q_*}$ is the best rational approximation of x from below with the largest denominator $q_* \leq n_{\max}$.

Note that when x is either irrational or rational with the denominator strictly greater than n_{\max} , then the mapping $n \mapsto nx - \lfloor nx \rfloor$ is one-to-one. Indeed, if $n_1x - \lfloor n_1x \rfloor = n_2x - \lfloor n_2x \rfloor$ holds with $n_1 \neq n_2$, then

$$x = \frac{\lfloor n_2x \rfloor - \lfloor n_1x \rfloor}{n_2 - n_1},$$

which is absurd.

Therefore, the theorem establishes the equivalence of two orderings on $[1: n_{\max}]$ induced from the embeddings

$$\begin{aligned} n &\mapsto (nx - \lfloor nx \rfloor, n) \quad \text{and} \\ n &\mapsto (n\xi - \lfloor n\xi \rfloor, n) \end{aligned}$$

into $[0, 1) \times [1: n_{\max}]$ endowed with the lexicographic ordering.

Proof. 1. We want to show the inequality

$$(n_1 - n_2)\xi \leq \lfloor n_1\xi \rfloor - \lfloor n_2\xi \rfloor = \lfloor n_1x \rfloor - \lfloor n_2x \rfloor,$$

which can be rewritten as

$$(n_1 - n_2)(\xi - x) \leq (n_2x - \lfloor n_2x \rfloor) - (n_1x - \lfloor n_1x \rfloor).$$

Since $x = \frac{p}{q}$ and $q \leq n_{\max}$, Theorem 7.4 implies

$$x \leq \xi < x + \frac{1}{vq}$$

where v is the largest integer such that $vp \equiv -1 \pmod{q}$ and $v \leq n_{\max}$. In particular, the inequality we want to show is trivial if $n_1 < n_2$ or $x = \xi$, so assume $n_1 > n_2$ and $x \neq \xi$, then we have

$$(n_1 - n_2)(\xi - x) < \frac{n_1 - n_2}{vq}.$$

Let r_1, r_2 be the remainders of n_1p divided by q and n_2p divided by q , respectively, then it suffices to show

$$(n_2x - \lfloor n_2x \rfloor) - (n_1x - \lfloor n_1x \rfloor) = \frac{r_2 - r_1}{q} \geq \frac{n_1 - n_2}{vq},$$

or equivalently,

$$n_1 - n_2 \leq v(r_2 - r_1),$$

assuming $r_2 > r_1$. Suppose not, so assume $n_1 - n_2 > v(r_2 - r_1)$. Note that

$$v(r_2 - r_1)p \equiv r_1 - r_2 \equiv (n_1 - n_2)p \pmod{q},$$

and since p and q are coprime, we get $v(r_2 - r_1) \equiv n_1 - n_2 \pmod{q}$, thus there exists a positive integer e such that

$$n_1 - n_2 = v(r_2 - r_1) + eq.$$

Since $r_2 > r_1$, this implies

$$n_{\max} \geq n_1 = v(r_2 - r_1) + eq + n_2 > v + q$$

which contradicts to the definition of v . Therefore, we get the desired inequality.

Note that since $n_1 - n_2 \leq v(r_2 - r_1)$ always holds whenever $n_1 > n_2$ and $r_2 > r_1$, in that case it is impossible to achieve the equality $n_1\xi - \lfloor n_1\xi \rfloor = n_2\xi - \lfloor n_2\xi \rfloor$. Hence, the equality might hold only when $n_1 \leq n_2$. Then one can easily see that $\xi = x$ is a necessary condition for having the equality.

2. Let $\frac{p_*}{q_*}$ be the best rational approximation of x from below with the largest denominator $q_* \leq n_{\max}$. Applying Theorem 7.4 with $\xi \leftarrow \frac{p_*}{q_*}$, we immediately get that $\left\lfloor \frac{np_*}{q_*} \right\rfloor = \lfloor nx \rfloor = \lfloor n\xi \rfloor$ holds for all $n \in [1: n_{\max}]$. Note that from $n_1x - \lfloor n_1x \rfloor < n_2x - \lfloor n_2x \rfloor$, we get

$$n_1 \left(x - \frac{p_*}{q_*} \right) + \frac{n_1p_*}{q_*} - \left\lfloor \frac{n_1p_*}{q_*} \right\rfloor < n_2 \left(x - \frac{p_*}{q_*} \right) + \frac{n_2p_*}{q_*} - \left\lfloor \frac{n_2p_*}{q_*} \right\rfloor. \quad (3)$$

We claim that

$$\frac{n_1p_*}{q_*} - \left\lfloor \frac{n_1p_*}{q_*} \right\rfloor \leq \frac{n_2p_*}{q_*} - \left\lfloor \frac{n_2p_*}{q_*} \right\rfloor.$$

This actually follows directly from the first part of the theorem, because if this is not the case, then the first part applied to $x \leftarrow \frac{p_*}{q_*}$ and $\xi \leftarrow x$ implies

$$n_1x - \lfloor n_1x \rfloor \geq n_2x - \lfloor n_2x \rfloor,$$

directly contradicting to the assumption.

Now, if

$$\frac{n_1p_*}{q_*} - \left\lfloor \frac{n_1p_*}{q_*} \right\rfloor < \frac{n_2p_*}{q_*} - \left\lfloor \frac{n_2p_*}{q_*} \right\rfloor$$

holds, then again we can apply the first part with $x \leftarrow \frac{p_*}{q_*}$ and $\xi \leftarrow \xi$ to conclude

$$n_1\xi - \lfloor n_1\xi \rfloor < n_2\xi - \lfloor n_2\xi \rfloor;$$

note that the equality cannot hold in this case. Therefore, we only need to consider the case

$$\frac{n_1 p_*}{q_*} - \left\lfloor \frac{n_1 p_*}{q_*} \right\rfloor = \frac{n_2 p_*}{q_*} - \left\lfloor \frac{n_2 p_*}{q_*} \right\rfloor.$$

Clearly, in this case (3) shows that we should have $n_1 < n_2$, thus we can again apply the first part to conclude

$$n_1 \xi - \lfloor n_1 \xi \rfloor \leq n_2 \xi - \lfloor n_2 \xi \rfloor,$$

with possibly the equality when $\xi = \frac{p_*}{q_*}$.

□

8 Simultaneously satisfying $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ and $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$

In this section, we now develop an algorithm for finding an approximation (ξ, ζ) of (x, y) such that Algorithm 5.4 should return the same outputs for (x, y) and (ξ, ζ) . By how Algorithm 5.4 is derived, one can see that a sufficient condition on (ξ, ζ) guaranteeing this equivalence is that the inequality

$$nx - \lfloor nx \rfloor \geq \lfloor mx + y \rfloor + 1 - (mx + y) \quad (4)$$

holds for some $m, n \in \mathbb{Z}$ with $m \geq n_{\min}$, $n > 0$ and $m + n \leq n_{\max}$ if and only if the same inequality holds with (x, y) replaced by (ξ, ζ) . Indeed, recall that the way we iteratively approached to the smallest minimizer and the largest maximizer is to successively add to m the smallest n satisfying/violating the above inequality. Hence, if the set of (m, n) satisfying/violating the inequality is same for (x, y) and (ξ, ζ) , then Algorithm 5.4 should return the same outputs for them.

One should however note that this does not mean that the *sets* of minimizers/maximizers for (x, y) and (ξ, ζ) should be equal, rather, it merely says that the *smallest minimizer* and the *largest maximizer* should be equal. Indeed, if nx is an integer for certain n 's, then adding or subtracting n to any minimizer/maximizer should still give a minimizer/maximizer, but the corresponding $n\xi$ does not need to be an integer, because the only thing the condition that nx is an integer imposes on (ξ, ζ) is that the inequality $n\xi - \lfloor n\xi \rfloor < \lfloor m\xi + \zeta \rfloor + 1 - (m\xi + \zeta)$ is satisfied for all m , which does not necessarily mean that the left-hand side is zero. This is one of the reasons why we only care about the smallest minimizer and the largest maximizer, rather than the whole sets of minimizers/maximizers.

The inequality (4) is equivalent to

$$(m + n)x + y \geq \lfloor nx \rfloor + \lfloor mx + y \rfloor + 1,$$

which is equivalent to

$$\lfloor (m + n)x + y \rfloor \geq \lfloor nx \rfloor + \lfloor mx + y \rfloor + 1.$$

Therefore, a sufficient condition on (ξ, ζ) we are seeking for is:

1. $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ holds for all $n \in [n_{\min} : n_{\max}]$ and
2. $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ holds for all $n \in [0 : n_{\max} - n_{\min}]$.

It might be puzzling to some of the readers, because finding a condition for having $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ was our main goal from the first place. However, our main goal was to establish a necessary and sufficient condition for having *only* $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$. On the other hand, we are to find a sufficient condition (although we will in fact obtain a necessary and sufficient condition) for having *both* of the above two. They are different problems.

Note that having $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$ can be rewritten as

$$\lfloor nx + (n_{\min}x + y) \rfloor = \lfloor n\xi + (n_{\min}\xi + \zeta) \rfloor$$

for all $n \in [0 : n_{\max} - n_{\min}]$. Hence, let $y' := (n_{\min}x + y) - \lfloor n_{\min}x + y \rfloor$ and $\zeta' := (n_{\min}\xi + \zeta) - \lfloor n_{\min}\xi + \zeta \rfloor$, then we can reformulate our problem as finding $\xi \in \mathbb{R}$, $\zeta' \in [0, 1)$ which satisfy

$$\lfloor nx + y' \rfloor = \lfloor n\xi + \zeta' \rfloor$$

for all $n \in [1 : n_{\max} - n_{\min}]$ for given $x \in \mathbb{R}$, $y' \in [0, 1)$. Once we find such ξ, ζ' , then we can recover ζ from the identity

$$\zeta = \zeta' + \lfloor n_{\min}x + y \rfloor - n_{\min}\xi$$

which also recovers the condition $\lfloor n_{\min}x + y \rfloor = \lfloor n_{\min}\xi + \zeta \rfloor$.

We will solve this reformulated problem using the following result, which is again from [5] with some minor modifications.

Theorem 8.1 (Simultaneously Satisfying $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ and $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$).
Let x, ξ be real numbers and n_{\max} a positive integer such that

$$\lfloor nx \rfloor = \lfloor n\xi \rfloor$$

holds for all $n \in [1 : n_{\max}]$, and let $\frac{p_*}{q_*}, \frac{p^*}{q^*}$ be the best rational approximations of x from below and above, respectively, with the largest denominators $q_*, q^* \leq n_{\max}$. Given $y \in [0, 1]$, define

$$\begin{aligned} L &:= \{n \in [1 : n_{\max}] : nx - \lfloor nx \rfloor < 1 - y\}, \\ R &:= \{n \in [1 : n_{\max}] : nx - \lfloor nx \rfloor \geq 1 - y\}. \end{aligned}$$

Then for a real number ζ , we have the followings.

1. When $L = \emptyset$,

$$\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$$

holds for all $n \in [1 : n_{\max}]$ if and only if

$$p_* + 1 - q_*\xi \leq \zeta < \left\lfloor \frac{\mu p_*}{q_*} \right\rfloor + 2 - \mu\xi,$$

where μ is defined as follows.

- (1) If $x = \frac{p}{q}$ is a rational number with $q \leq n_{\max}$, then μ is the largest integer such that $\mu \leq n_{\max}$ and $\mu p \equiv -1 \pmod{q}$.
- (2) If x is either an irrational number or a rational number with the denominator strictly greater than n_{\max} , then $\mu = q^*$.

2. When $R = \emptyset$,

$$\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$$

holds for all $n \in [1: n_{\max}]$ if and only if

$$p_* - q_*\xi \leq \zeta < \left\lfloor \frac{\mu p_*}{q_*} \right\rfloor + 1 - \mu\xi,$$

where μ is defined as in the case $L = \emptyset$.

3. When L and R are both nonempty,

$$\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$$

holds for all $n \in [1: n_{\max}]$ if and only if

$$\left\lfloor \frac{\nu p_*}{q_*} \right\rfloor + 1 - \nu\xi \leq \zeta < \left\lfloor \frac{\mu p_*}{q_*} \right\rfloor + 1 - \mu\xi,$$

where μ and ν are defined as follows.

- (1) If x is rational and its reduced denominator is at most n_{\max} , then μ is the largest integer such that $\mu \leq n_{\max}$ and $\mu p_* \equiv -\lfloor q_* y \rfloor - 1 \pmod{q_*}$, and ν is the smallest positive integer such that $\nu p_* \equiv -\lfloor q_* y \rfloor \pmod{q_*}$.
- (2) If x is either irrational or is rational whose reduced denominator is strictly larger than n_{\max} , then μ is the largest positive integer such that $\mu \leq n_{\max}$,

$$\mu p_* \equiv r \pmod{q_*} \quad \text{and} \quad \mu < \frac{q_*(1-y) - r}{q_*x - p_*},$$

where $r \in [0: q_* - 1]$ is chosen to be the largest allowing such μ to exist. Similarly, ν is the smallest positive integer such that $\nu \leq n_{\max}$,

$$\nu p_* \equiv r \pmod{q_*} \quad \text{and} \quad \nu \geq \frac{q_*(1-y) - r}{q_*x - p_*},$$

where $r \in [0: q_* - 1]$ is chosen to be the smallest allowing such ν to exist.

Proof. Note that

$$\lfloor nx + y \rfloor = \begin{cases} \lfloor nx \rfloor & \text{if } nx - \lfloor nx \rfloor < 1 - y, \\ \lfloor nx \rfloor + 1 & \text{if } nx - \lfloor nx \rfloor \geq 1 - y, \end{cases}$$

so we have the equality

$$\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$$

for all $n \in [1: n_{\max}]$ if and only if

$$L = L' := \{n \in [1: n_{\max}]: -\zeta \leq n\xi - \lfloor n\xi \rfloor < 1 - \zeta\}$$

and

$$R = R' := \{n \in [1: n_{\max}]: 1 - \zeta \leq n\xi - \lfloor n\xi \rfloor < 2 - \zeta\}.$$

Now, as $L \cup R = [1: n_{\max}]$ and $L' \cap R' = \emptyset$, we have $L = L'$ and $R = R'$ if and only if we have $L \subseteq L'$ and $R \subseteq R'$. Note that $L \subseteq L'$ holds if and only if either $L = \emptyset$ or

$$-\min_{n \in L}(n\xi - \lfloor n\xi \rfloor) \leq \zeta < 1 - \max_{n \in L}(n\xi - \lfloor n\xi \rfloor),$$

and $R \subseteq R'$ holds if and only if either $R = \emptyset$ or

$$1 - \min_{n \in R}(n\xi - \lfloor n\xi \rfloor) \leq \zeta < 2 - \max_{n \in R}(n\xi - \lfloor n\xi \rfloor).$$

1. When $L = \emptyset$, we only need to verify

$$\max_{n \in [1: n_{\max}]} (n\xi - \lfloor n\xi \rfloor) = \mu\xi - \lfloor \mu\xi \rfloor$$

and

$$\min_{n \in [1: n_{\max}]} (n\xi - \lfloor n\xi \rfloor) = q_*\xi - \lfloor q_*\xi \rfloor,$$

since $\lfloor n\xi \rfloor = \lfloor nx \rfloor = \left\lfloor \frac{np_*}{q_*} \right\rfloor$ holds for all n by Theorem 7.4. These two both follow directly from Theorem 7.6 and Theorem 2.4, because q_* is the smallest minimizer of $nx - \lfloor nx \rfloor$ and μ is the largest maximizer of $nx - \lfloor nx \rfloor$ for $n \in [1: n_{\max}]$.

2. Similarly, when $R = \emptyset$, we only need to verify

$$\max_{n \in [1: n_{\max}]} (n\xi - \lfloor n\xi \rfloor) = \mu\xi - \lfloor \mu\xi \rfloor$$

and

$$\min_{n \in [1: n_{\max}]} (n\xi - \lfloor n\xi \rfloor) = q_*\xi - \lfloor q_*\xi \rfloor,$$

which we have already done in the case $L = \emptyset$.

3. Now, suppose L and R are both not empty. Then it suffices to show that

$$\mu\xi - \lfloor \mu\xi \rfloor = \max_{n \in L} (n\xi - \lfloor n\xi \rfloor)$$

and

$$\nu\xi - \lfloor \nu\xi \rfloor = \min_{n \in R} (n\xi - \lfloor n\xi \rfloor).$$

By Theorem 7.6, it suffices to show that μ is the largest maximizer of $nx - \lfloor nx \rfloor$ for $n \in L$ and ν is the smallest minimizer of $nx - \lfloor nx \rfloor$ for $n \in R$.

When x is rational whose reduced denominator is at most n_{\max} , then we have $x = \frac{p_*}{q_*}$. Then since $n \in R$ if and only if $(np_* \bmod q_*) \geq q_* - \lfloor q_*y \rfloor$, the conclusion about ν follows immediately. For μ , note that $\lfloor q_*y \rfloor < q_*$ always holds unless $y = 1$, but $y = 1$ is impossible since it implies $L = \emptyset$. Hence, $q_* - \lfloor q_*y \rfloor - 1 \geq 0$ holds, so we have $n \in L$ if and only if $(np_* \bmod q_*) \leq q_* - \lfloor q_*y \rfloor - 1$, so the conclusion about μ also follows.

Next, assume x is either irrational or rational whose reduced denominator strictly larger than n_{\max} . In this case, we can apply Theorem 7.6 to see that the largest maximizer of $nx - \lfloor nx \rfloor$ for $n \in L$ is precisely the largest maximizer of $\frac{np_*}{q_*} - \left\lfloor \frac{np_*}{q_*} \right\rfloor$ for $n \in L$. Since

$$\frac{np_*}{q_*} - \left\lfloor \frac{np_*}{q_*} \right\rfloor = \frac{(np_* \bmod q_*)}{q_*},$$

the maximum value of $\frac{np_*}{q_*} - \left\lfloor \frac{np_*}{q_*} \right\rfloor$ is obtained at $n \in L$ such that np_* has the largest possible remainder when divided by q_* . In general, for $n \in [1: n_{\max}]$, if we write $r := (np_* \bmod q_*)$, then

$$nx - \lfloor nx \rfloor = n \left(x - \frac{p_*}{q_*} \right) + \left(\frac{np_*}{q_*} - \left\lfloor \frac{np_*}{q_*} \right\rfloor \right) = n \left(x - \frac{p_*}{q_*} \right) + \frac{r}{q_*},$$

so $x \in L$ if and only if

$$n \left(x - \frac{p_*}{q_*} \right) + \frac{r}{q_*} < 1 - y.$$

or equivalently,

$$n < \frac{q_*(1-y) - r}{q_*x - p_*},$$

which gives the desired description of μ . We can obtain the description of ν in a similar way.

□

Note that Theorem 7.4 gives us a criterion for having $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ that is easy to verify, so using the above theorem in conjunction with Theorem 7.4 can give us a *necessary and sufficient* condition for having $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ and $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ for all $n \in [1 : n_{\max}]$ given $y, \zeta \in [0, 1)$, even though we only needed a sufficient condition.

When either of L or R is empty, or when x is rational whose reduced denominator is at most n_{\max} , then we can easily compute the complete set of (ξ, ζ) using Theorem 7.4 and Theorem 8.1, given that we know how to compute $\lfloor q_*y \rfloor$. This can be indeed done by using the algorithm we develop in Section 12.

However, when both L and R are not empty and when x is either irrational or is rational but the reduced denominator is strictly larger than n_{\max} , then finding μ, ν appearing in the statement of Theorem 8.1 by successively checking all r can be very inefficient, as q_* can be in general as large as n_{\max} .⁹ Hence, we describe an algorithm for efficiently finding out the optimal r with a bounded number of trials.

Recall from Remark 7.5 that $p^*q_* - q^*p_* = 1$, in particular $q^*p_* \equiv -1 \pmod{q_*}$ holds. Consider the set of $(k, r) \in \mathbb{Z} \times [0 : q_* - 1]$ such that

$$0 < q_*k - q^*r \leq n_{\max},$$

that is, we decompose $[1 : n_{\max}]$ into set of n 's such that $np_* \equiv r \pmod{q_*}$ for each r , which means $n \equiv -q^*r \pmod{q_*}$, and then let $k := \frac{n+q^*r}{q_*}$. Then the constraint

$$n < \frac{q_*(1-y) - r}{q_*x - p_*}$$

appearing in the description of μ can be equivalently written as

$$(1 - q^*(q_*x - p_*))r < q_*(1-y) - q_*(q_*x - p_*)k$$

by substituting $n \leftarrow q_*k - q^*r$. Since

$$q^*(q_*x - p_*) = q^*q_* \left(x - \frac{p_*}{q_*} \right) < q^*q_* \left(\frac{p^*}{q^*} - \frac{p_*}{q_*} \right) = 1,$$

the inequality can be further rewritten as

$$r < \frac{q_*(1-y) - q_*(q_*x - p_*)k}{1 - q^*(q_*x - p_*)}.$$

⁹In practice, I observed that the optimal r in most cases were around $q_* - \lfloor q_*y \rfloor$, often off by small numbers. However, I was not able to prove that this should always be the case, I am suspecting that this was merely because the continued fraction coefficient of x I experimented with were mostly just small numbers. In the description of the algorithm that follows, it becomes apparent that having a big continued fraction coefficient can result in making q_* a lot larger than q , which can shift the optimal r a lot from its obvious rough guess.

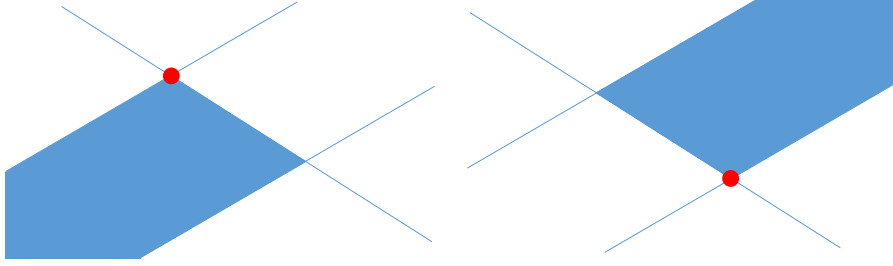


Figure 2: Left: (k, r) -region for μ , Right: (k, r) -region for ν

Next, let us ignore the constraint $(k, r) \in \mathbb{Z} \times [0: q_* - 1]$ for a while and consider the set of $(k, r) \in \mathbb{R} \times \mathbb{R}$ satisfying three inequalities

$$0 < q_* k - q^* r \leq n_{\max} \quad \text{and} \quad r < \frac{q_*(1-y) - q_*(q_* x - p_*)k}{1 - q^*(q_* x - p_*)} \quad (5)$$

as shown in the left-hand side of Figure 2.

The point marked in red is the highest possible value of r we can have in this region, once we ignore the constraint $(k, r) \in \mathbb{Z} \times [0: q_* - 1]$. We first evaluate the r -coordinate r_T of this point. This can be done by solving the equation

$$\frac{q_*(1-y) - q_*(q_* x - p_*)k}{1 - q^*(q_* x - p_*)} = \frac{q_* k}{q^*},$$

for k . The solution is

$$k_T := q^*(1-y)$$

and as a result we obtain

$$r_T = q_*(1-y).$$

Therefore, the maximum possible value of r is at most $\lceil r_T \rceil - 1$. We start from there and go further down if there is no $k \in \mathbb{Z}$ such that $(k, \lceil r_T \rceil - 1)$ is in the region. Note that r_T is clearly strictly smaller than q_* unless $y = 0$, but $y = 0$ is impossible as it yields $L = \emptyset$ contradicting to our assumption. Also, if we cannot find any $(k, r) \in \mathbb{Z} \times \mathbb{Z}$ until we further go down below $r = 0$, then this means that the actual set of (k, r) is empty, and that means $L = \emptyset$ which is again impossible. Hence, we only need to take account of the integer constraint $(k, r) \in \mathbb{Z} \times \mathbb{Z}$, and can ignore the range constraint $0 \leq r \leq q_* - 1$.

We now describe this search procedure for (k, r) . We first decompose the region defined by three inequalities in (5) into two, the top triangular region and the bottom parallelogrammatic region extending indefinitely into the left-bottom direction. These two regions are separated by the horizontal line

$$r = r_B = q_*(1-y) - (q_* x - p_*)n_{\max},$$

where r_B is the r -coordinate of the intersection between the lines

$$q_* k - q^* r = 0 \quad \text{and} \quad r = \frac{q_*(1-y) - q_*(q_* x - p_*)k}{1 - q^*(q_* x - p_*)}.$$

As a convention, we include the horizontal line $r = r_B$ into the triangular region and exclude it from the parallelogrammatic region.

We note that the parallelogrammatic region is horizontally wide enough to carry a point of integer k -coordinate for every possible r -coordinate. Indeed, for any given $r = r_0$, the horizontal gap between the two bounding lines is

$$\frac{n_{\max} + q^* r_0}{q_*} - \frac{q^* r_0}{q_*} = \frac{n_{\max}}{q_*}. \quad (6)$$

By the assumption on x , n_{\max} is strictly larger than q_* , so the right-hand side is strictly larger than 1. Hence, there must be a point of the form (k, r_0) in this region where $k \in \mathbb{Z}$, and the largest integer k -coordinate of such points can be always obtained by taking the floor of the k -coordinate of the right-most point with $r = r_0$.

There are several cases we have to deal separately.

1. First, suppose $\lceil r_T \rceil - 1 < r_B$, that is, $\lceil r_T \rceil - 1 < \lceil r_B \rceil$. Hence, the horizontal line $r = \lceil r_T \rceil - 1$ crosses the parallelogrammatic region. The k -coordinate of the right-most point in the intersection is

$$k_{PR} := \frac{n_{\max} + q^*(\lceil r_T \rceil - 1)}{q_*} = \frac{n_{\max} + q^*(\lceil r_B \rceil - 1)}{q_*}$$

and recall that by (6) the intersection is guaranteed to have a point with an integer k -coordinate. Since μ is the largest n satisfying $np_* \equiv r \pmod{q_*}$, we return $(\lfloor k_{PR} \rfloor, \lceil r_B \rceil - 1)$.

2. Next, suppose $\lceil r_T \rceil - 1 \geq \lceil r_B \rceil$, that is, the horizontal line $r = \lceil r_T \rceil - 1$ crosses the triangular region. The k -coordinate of the right-most point of the intersection is

$$k_{TR} := \frac{q_*(1 - y) - (1 - q^*(q_*x - p_*))(\lceil r_T \rceil - 1)}{q_*(q_*x - p_*)}$$

and the k -coordinate of the left-most point is

$$k_{TL} := \frac{q^*(\lceil r_T \rceil - 1)}{q_*},$$

both exclusive. If there is a point in the intersection with an integer k -coordinate, then the one with the largest k -coordinate must be the one whose k -coordinate is equal to $\lceil k_{TR} \rceil - 1$. Hence, if $\lceil k_{TR} \rceil - 1 > k_{TL}$, that is, $\lceil k_{TR} \rceil - 1 > \lfloor k_{TL} \rfloor$, then we return $(\lceil k_{TR} \rceil - 1, \lceil r_T \rceil - 1)$.

3. If $\lceil k_{TR} \rceil - 1 \leq \lfloor k_{TL} \rfloor$, then there is no point in the intersection with an integer k -coordinate. In this case we claim that, if there exists any point (k, r) in the triangular region with integer coordinates, then at least one of $(\lceil k_{TR} \rceil, r) = (\lfloor k_{TL} \rfloor + 1, r)$ or $(\lfloor k_{TL} \rfloor, r)$ must be in the triangular region. Indeed, assume $k \geq \lceil k_{TR} \rceil$, then clearly we have

$$r < \frac{q_*(1 - y) - q_*(q_*x - p_*)k}{1 - q^*(q_*x - p_*)} \leq \frac{q_*(1 - y) - q_*(q_*x - p_*)\lceil k_{TR} \rceil}{1 - q^*(q_*x - p_*)}.$$

On the other hand, since we know $k_{TR} > k_T$ from $r_T > \lceil r_T \rceil - 1$ (recall k_T is the k -coordinate of the point marked in red in the left-hand side of Figure 2), thus

$$q_* \lceil k_{TR} \rceil - q^* r > q_* k_T - q^* r > q_* k_T - q^* r_T = 0,$$

which shows $(\lceil k_{TR} \rceil, r)$ is in the triangular region. Next, assume $k \leq \lfloor k_{TL} \rfloor$, then clearly we have

$$0 < q_*k - q^*r \leq q_* \lfloor k_{TL} \rfloor - q^*r.$$

On the other hand, since we know $k_{TL} < k_T$ from $r_T > \lceil r_T \rceil - 1$, thus

$$\frac{q_*(1-y) - q_*(q_*x - p_*) \lfloor k_{TL} \rfloor}{1 - q^*(q_*x - p_*)} > \frac{q_*(1-y) - q_*(q_*x - p_*)k_T}{1 - q^*(q_*x - p_*)} = r_T > r,$$

which shows $(\lfloor k_{TL} \rfloor, r)$ is in the triangular region. Hence, the claim is shown.

Therefore, we see if there is any point in the triangular region with integer coordinates whose k -coordinate is either $\lceil k_{TR} \rceil$ or $\lfloor k_{TL} \rfloor$. If not, then we conclude there is no point with integer coordinates in the triangular region. To do this check, let

$$r_{TR} := \frac{q_*(1-y) - q_*(q_*x - p_*) \lceil k_{TR} \rceil}{1 - q^*(q_*x - p_*)} \quad \text{and} \quad r_{TL} := \frac{q_* \lfloor k_{TL} \rfloor}{q^*},$$

the r -coordinates of the intersection of the vertical line $k = \lceil k_{TR} \rceil$ with the right-boundary of the triangular region, and the intersection of the vertical line $k = \lfloor k_{TL} \rfloor$ with the left-boundary of the triangular region, respectively.

If $\lceil r_{TR} \rceil - 1 < \lceil r_B \rceil$ and $\lceil r_{TL} \rceil - 1 < \lceil r_B \rceil$, then we conclude there is no point with integer coordinates in the triangular region. In this case, we simply choose the largest integer r -coordinate $\lceil r_B \rceil - 1$ in the parallelogrammatic region. That is, let

$$k_{PR} := \frac{n_{\max} + q^*(\lceil r_B \rceil - 1)}{q_*},$$

and we return $(\lfloor k_{PR} \rfloor, \lceil r_B \rceil - 1)$.

4. If $\lceil r_{TR} \rceil - 1 \geq \lceil r_B \rceil$ and $\lceil r_{TL} \rceil - 1 < \lceil r_B \rceil$, then there exists at least one point in the triangular region with integer coordinates, and there must exist a point whose k -coordinate is equal to $\lceil k_{TR} \rceil$ among the ones whose integer r -coordinate is the largest. Hence, we simply choose the one with the largest integer k -coordinate among the points on the intersection of the triangular region and the horizontal line $r = \lceil r_{TR} \rceil - 1$. That is, let

$$k_{TRR} := \frac{q_*(1-y) - (1 - q^*(q_*x - p_*))(\lceil r_{TR} \rceil - 1)}{q_*(q_*x - p_*)}$$

and return $(\lceil k_{TRR} \rceil - 1, \lceil r_{TR} \rceil - 1)$.

5. If $\lceil r_{TR} \rceil - 1 < \lceil r_B \rceil$ and $\lceil r_{TL} \rceil - 1 \geq \lceil r_B \rceil$, then there exists at least one point in the triangular region with integer coordinates, and there must exist a point whose k -coordinate is equal to $\lfloor k_{TL} \rfloor$ among the ones whose integer r -coordinate is the largest. Hence, we simply choose the one with the largest integer k -coordinate among the points on the intersection of the triangular region and the horizontal line $r = \lceil r_{TL} \rceil - 1$. That is, let

$$k_{TLR} := \frac{q_*(1-y) - (1 - q^*(q_*x - p_*))(\lceil r_{TL} \rceil - 1)}{q_*(q_*x - p_*)}$$

and return $(\lceil k_{TLR} \rceil - 1, \lceil r_{TL} \rceil - 1)$.

6. If $\lceil r_{TR} \rceil - 1 \geq \lceil r_B \rceil$ and $\lceil r_{TL} \rceil - 1 \geq \lceil r_B \rceil$, then we choose the larger one between $\lceil r_{TR} \rceil - 1$ and $\lceil r_{TL} \rceil - 1$ and do the same. That is, if $\lceil r_{TR} \rceil - 1 \geq \lceil r_{TL} \rceil - 1$ then return $(\lceil k_{TRR} \rceil - 1, \lceil r_{TR} \rceil - 1)$, and otherwise return $(\lceil k_{TLR} \rceil - 1, \lceil r_{TL} \rceil - 1)$.

After following this procedure, we get a point $(k, r) \in \mathbb{Z} \times [0: q_* - 1]$, and we conclude $\mu = q_*k - q^*r$.

Next, we consider the case of ν . The only difference in this case is that, now the region we are looking at is defined by

$$0 < q_*k - q^*r \leq n_{\max} \quad \text{and} \quad r \geq \frac{q_*(1-y) - q_*(q_*x - p_*)k}{1 - q^*(q_*x - p_*)} \quad (7)$$

and we are interested in finding the bottom-most and left-most point (prioritizing the first) in the region with integer coordinates. Again, we start from the point marked in red in the right-hand side of Figure 2. However, in this case, this point, whose coordinates (k_B, r_B) are given as

$$k_B = q^*(1-y) + \frac{(1 - q^*(q_*x - p_*))n_{\max}}{q_*} \quad \text{and} \quad r_B = q_*(1-y) - (q_*x - p_*)n_{\max},$$

may have negative r -coordinate. Hence, instead of $\lceil r_B \rceil$ we start with

$$r'_B := \max\{\lceil r_B \rceil, 0\}.$$

On the other hand, once the point (k, r) with the smallest-coordinate in the region defined by these constraints is found, then r must be no more than $q_* - 1$, because otherwise we conclude $R = \emptyset$, contradicting the assumption.

Except for that inclusivity of boundary lines are different and now we are seeking for the left-most rather than the right-most point, the procedure of finding (k, r) is largely same as in the case of μ :

1. First, suppose $r'_B \geq r_T$, that is, $r'_B \geq \lceil r_T \rceil$ which actually means $r'_B = \lceil r_T \rceil$ since $r_T = q_*(1-y) \geq 0$. Hence, the horizontal line $r = r'_B$ crosses the parallelogrammatic region. The k -coordinate of the left-most point of the intersection is

$$k_{PL} := \frac{qr'_B}{q_*} = \frac{q^* \lceil r_T \rceil}{q_*},$$

and we return $(\lfloor k_{PL} \rfloor + 1, \lceil r_T \rceil)$.

2. Next, suppose $r'_B < \lceil r_T \rceil$, that is, the horizontal line $r = r'_B$ crosses the triangular region. The k -coordinate of the left-most point of the intersection is

$$k_{BL} := \frac{q_*(1-y) - (1 - q^*(q_*x - p_*))r'_B}{q_*(q_*x - p_*)}$$

and the k -coordinate of the right-most point is

$$k_{BR} := \frac{n_{\max} + q^*r'_B}{q_*},$$

both inclusive. If there is a point in the intersection with an integer k -coordinate, then the one with the smallest k -coordinate must be the one whose k -coordinate is equal to $\lceil k_{BL} \rceil$. Hence, if $\lceil k_{BL} \rceil \leq k_{BR}$, that is, $\lceil k_{BL} \rceil \leq \lfloor k_{BR} \rfloor$, then we return $(\lceil k_{BL} \rceil, r'_B)$.

3. If $\lceil k_{BL} \rceil > \lfloor k_{BR} \rfloor$, then there is no point in the intersection with an integer k -coordinate. By the same argument as in the case of μ , we conclude that if there

exists any point (k, r) in the triangular region with integer coordinates, then at least one of $(\lceil k_{BL} \rceil - 1, r) = (\lfloor k_{BR} \rfloor, r)$ or $(\lfloor k_{BR} \rfloor + 1, r)$ must be in the triangular region.

Therefore, we see if there is any point in the triangular region with integer coordinates whose k -coordinate is either $\lfloor k_{BR} \rfloor$ or $\lceil k_{BL} \rceil$. If not, then we conclude there is no point with integer coordinates in the triangular region. To do this check, let

$$r_{BL} := \frac{q_*(1-y) - q_*(q_*x - p_*) \lfloor k_{BR} \rfloor}{1 - q_*(q_*x - p_*)} \quad \text{and} \quad r_{BR} := \frac{q_* \lceil k_{BL} \rceil - n_{\max}}{q_*},$$

the r -coordinates of the intersection of the vertical line $k = \lfloor k_{BR} \rfloor$ with the left-boundary of the triangular region, and the intersection of the vertical line $k = \lceil k_{BL} \rceil$ with the right-boundary of the triangular region, respectively (note that “BL” is from “BR” and “BR” is from “BL”).

If $\lceil r_{BL} \rceil \geq \lceil r_T \rceil$ and $\lceil r_{BR} \rceil \geq \lceil r_T \rceil$, then we conclude there is no point with integer coordinates in the triangular region. In this case, we simply choose the smallest integer r -coordinate $\lceil r_T \rceil$ in the parallelogrammatic region. That is, let

$$k_{PL} := \frac{q_* \lceil r_T \rceil}{q_*},$$

and we return $(\lfloor k_{PL} \rfloor + 1, \lceil r_T \rceil)$.

4. If $\lceil r_{BL} \rceil < \lceil r_T \rceil$ and $\lceil r_{BR} \rceil \geq \lceil r_T \rceil$, then there exists at least one point in the triangular region with integer coordinates, and there must exist a point whose k -coordinate is equal to $\lfloor k_{BR} \rfloor$ among the ones whose integer r -coordinate is the smallest. Hence, we simply choose the one with the smallest integer k -coordinate among the points on the intersection of the triangular region and the horizontal line $r = \lceil r_{BL} \rceil$. That is, let

$$k_{BLL} := \frac{q_*(1-y) - (1 - q_*(q_*x - p_*)) \lceil r_{BL} \rceil}{q_*(q_*x - p_*)}$$

and return $(\lceil k_{BLL} \rceil, \lceil r_{BL} \rceil)$.

5. If $\lceil r_{BL} \rceil \geq \lceil r_T \rceil$ and $\lceil r_{BR} \rceil < \lceil r_T \rceil$, then there exists at least one point in the triangular region with integer coordinates, and there must exist a point whose k -coordinate is equal to $\lfloor k_{BR} \rfloor$ among the ones whose integer r -coordinate is the smallest. Hence, we simply choose the one with the smallest integer k -coordinate among the points on the intersection of the triangular region and the horizontal line $r = \lceil r_{BR} \rceil$. That is, let

$$k_{BRL} := \frac{q_*(1-y) - (1 - q_*(q_*x - p_*)) \lceil r_{BR} \rceil}{q_*(q_*x - p_*)}$$

and return $(\lceil k_{BRL} \rceil, \lceil r_{BR} \rceil)$.

6. If $\lceil r_{BL} \rceil < \lceil r_T \rceil$ and $\lceil r_{BR} \rceil < \lceil r_T \rceil$, then we choose the smaller one between $\lceil r_{BL} \rceil$ and $\lceil r_{BR} \rceil$ and do the same. That is, if $\lceil r_{BL} \rceil \leq \lceil r_{BR} \rceil$ then return $(\lceil k_{BLL} \rceil, \lceil r_{BL} \rceil)$, and otherwise return $(\lceil k_{BRL} \rceil, \lceil r_{BR} \rceil)$.

After following this procedure, we get a point $(k, r) \in \mathbb{Z} \times [0: q_* - 1]$, and we conclude $\nu = q_*k - q_*r$.

Here we give a summary of all quantities we need, after substituting $y \leftarrow y' = (n_{\min}x + y) - \lfloor n_{\min}x + y \rfloor$ and $n_{\max} \leftarrow n_{\max} - n_{\min}$. Some of them directly involve x, y in their

representations, but the algorithm we develop in Section 12 is able to compute all of these within arbitrary-precision exact rational arithmetic. Those quantities are rewritten here in a way that the aforementioned algorithm can be directly applied:

$$\begin{aligned}
\lceil r_T \rceil &= -\lfloor q_* y' - q_* \rfloor = -\lfloor q_* n_{\min} x + q_* y \rfloor + q_* (\lfloor n_{\min} x + y \rfloor + 1), \\
\lceil r_B \rceil &= -\lfloor q_* (n_{\max} - n_{\min}) x + q_* y' - (q_* + p_* (n_{\max} - n_{\min})) \rfloor \\
&= -\lfloor q_* n_{\max} x + q_* y \rfloor + (q_* (\lfloor n_{\min} x + y \rfloor + 1) + p_* (n_{\max} - n_{\min})), \\
r'_B &= \max \{ \lceil r_B \rceil, 0 \}, \\
\lfloor k_{PR} \rfloor &= \left\lfloor \frac{(n_{\max} - n_{\min}) + q^* (\lceil r_B \rceil - 1)}{q_*} \right\rfloor, \\
\lceil k_{TR} \rceil &= k_{\text{common}}(\lceil r_T \rceil - 1), \\
\lfloor k_{TL} \rfloor &= \left\lfloor \frac{q^* (\lceil r_T \rceil - 1)}{q_*} \right\rfloor, \\
\lceil r_{TR} \rceil &= r_{\text{common}}(\lceil k_{TR} \rceil), \\
\lceil r_{TL} \rceil &= \left\lceil \frac{q_* \lfloor k_{TL} \rfloor}{q^*} \right\rceil, \\
\lceil k_{TRR} \rceil &= k_{\text{common}}(\lceil r_{TR} \rceil - 1), \\
\lceil k_{TLR} \rceil &= k_{\text{common}}(\lceil r_{TL} \rceil - 1), \\
\lfloor k_{PL} \rfloor &= \left\lfloor \frac{q^* \lceil r_T \rceil}{q_*} \right\rfloor, \\
\lceil k_{BL} \rceil &= k_{\text{common}}(r'_B), \\
\lfloor k_{BR} \rfloor &= \left\lfloor \frac{(n_{\max} - n_{\min}) + q^* r'_B}{q_*} \right\rfloor, \\
\lceil r_{BL} \rceil &= r_{\text{common}}(\lfloor k_{BR} \rfloor), \\
\lceil r_{BR} \rceil &= \left\lceil \frac{q_* \lfloor k_{BL} \rfloor - (n_{\max} - n_{\min})}{q^*} \right\rceil, \\
\lceil k_{BLL} \rceil &= k_{\text{common}}(\lceil r_{BL} \rceil), \quad \text{and} \\
\lceil k_{BRL} \rceil &= k_{\text{common}}(\lceil r_{BR} \rceil)
\end{aligned}$$

where

$$\begin{aligned}
k_{\text{common}} : r \mapsto & - \left\lfloor \frac{-q^* q_* r x + q_* y' + ((q^* p_* + 1)r - q_*)}{q_*^2 x - q_* p_*} \right\rfloor \\
&= - \left\lfloor \frac{q_* (n_{\min} - q^* r) x + q_* y + ((q^* p_* + 1)r - q_* (\lfloor n_{\min} x + y \rfloor + 1))}{q_*^2 x - q_* p_*} \right\rfloor
\end{aligned}$$

and

$$\begin{aligned}
r_{\text{common}} : k \mapsto & - \left\lceil \frac{q_*^2 k x + q_* y' - (q_* p_* k + q_*)}{-q^* q_* x + (q^* p_* + 1)} \right\rceil \\
&= - \left\lceil \frac{q_* (q_* k + n_{\min}) x + q_* y - q_* (p_* k + \lfloor n_{\min} x + y \rfloor + 1)}{-q^* q_* x + (q^* p_* + 1)} \right\rceil.
\end{aligned}$$

In addition to these, checking for emptiness of L and R also involves quantities with x and y directly in their representations. That is, we have to compare the minimum and the

maximum of $nx - \lfloor nx \rfloor$ for $n \in [1 : n_{\max}]$ with $1 - y$. By Theorem 2.4, we know

$$\min_{n \in [1 : n_{\max}]} (nx - \lfloor nx \rfloor) = q_*x - p_*,$$

and

$$\max_{n \in [1 : n_{\max}]} (nx - \lfloor nx \rfloor) = \frac{q_* - 1}{q_*}$$

if $x = \frac{p_*}{q_*}$, and otherwise

$$\max_{n \in [1 : n_{\max}]} (nx - \lfloor nx \rfloor) = q^*x - p^* + 1.$$

Therefore, L is empty if and only if

$$q_*x - p_* \geq 1 - y,$$

that is,

$$\lfloor q_*x + y \rfloor > p_*.$$

Similarly, if $x = \frac{p_*}{q_*}$, then R is empty if and only if

$$\frac{q_* - 1}{q_*} < 1 - y,$$

that is,

$$\lfloor q_*y \rfloor = 0,$$

and otherwise,

$$q^*x - p^* + 1 < 1 - y,$$

that is,

$$\lfloor q^*x + y \rfloor < p^*.$$

Again, computation of $\lfloor q_*x + y \rfloor$ and $\lfloor q^*x + y \rfloor$ can be done within arbitrary-precision exact rational arithmetic using the algorithm developed in Section 12. Finally, we obtain the following algorithm for computing a quadruple of integers (a, b, c, d) such that $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ holds for all $n \in [n_{\min} : n_{\max}]$ if and only if

$$a - b\xi \leq \zeta < c - d\xi,$$

provided that we know $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ holds for all $n \in [0 : n_{\max} - n_{\min}]$. Recall that Theorem 8.1 is supposed to compute the range of $\zeta' := (n_{\min}\xi + \zeta) - \lfloor n_{\min}x + y \rfloor$ since y is replaced by $y' := (n_{\min}x + y) - \lfloor n_{\min}x + y \rfloor$. Hence, when the condition dictated by Theorem 8.1 is of the form

$$a - b\xi \leq \zeta < c - d\xi,$$

then it actually means

$$(a + \lfloor n_{\min}x + y \rfloor) - (b + n_{\min})\xi \leq \zeta < (c + \lfloor n_{\min}x + y \rfloor) - (d + n_{\min})\xi.$$

Also, we must ensure $\zeta' = \zeta - \lfloor n_{\min}x + y \rfloor + n_{\min}\xi$ belongs to the interval $[0, 1)$, because otherwise we fail to have $\lfloor n_{\min}x + y \rfloor = \lfloor n_{\min}\xi + \zeta \rfloor$. When $L = \emptyset$, the right-hand side

$$\left\lfloor \frac{\mu p_*}{q_*} \right\rfloor + 2 - \mu\xi$$

appearing in Theorem 8.1 for the case $L = \emptyset$ needs to be replaced by 1, and the left-hand side

$$p_* - q_*\xi$$

for the case $R = \emptyset$ needs to be replaced by 0. Indeed, by Theorem 7.4, we have

$$\lfloor n\xi \rfloor = \left\lfloor \frac{np_*}{q_*} \right\rfloor$$

for all $n \in [1 : n_{\max} - n_{\min}]$, thus the right-hand side for the case $L = \emptyset$ is always strictly greater than 1 and the left-hand side for the case $R = \emptyset$ is always less than or equal to 0. All the other bounds are at least as strict as $\zeta' \in [0, 1)$, so should not be touched.

Algorithm 8.2 (Simultaneously having $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ and $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$).

1. If $n_{\min} = n_{\max}$, we only need to ensure $\lfloor n_{\min}x + y \rfloor = \lfloor n_{\min}\xi + \zeta \rfloor$, which is equivalent to

$$\lfloor n_{\min}x + y \rfloor - n_{\min}\xi \leq \zeta < \lfloor n_{\min}x + y \rfloor + 1 - n_{\min}\xi,$$

thus return $(\lfloor n_{\min}x + y \rfloor, n_{\min}, \lfloor n_{\min}x + y \rfloor + 1, n_{\min})$.

2. Otherwise, compute the best rational approximations $\frac{p_*}{q_*}$ from below and $\frac{p^*}{q^*}$ from above of x with the largest reduced denominators $q_*, q^* \leq n_{\max} - n_{\min}$.
3. Compute the reduced denominator v of the right-hand side appearing in Theorem 7.4; that is, v is the largest integer such that $vp_* \equiv -1 \pmod{q_*}$ and $v \leq n_{\max}$.
4. If $\lfloor (q_* + n_{\min})x + y \rfloor > p_* + \lfloor n_{\min}x + y \rfloor$ holds, then $L = \emptyset$. In this case, return

$$(p_* + \lfloor n_{\min}x + y \rfloor + 1, q_* + n_{\min}, \lfloor n_{\min}x + y \rfloor + 1, n_{\min}).$$

5. Otherwise, if $q_* = q^*$ and $\lfloor q_*n_{\min}x + q_*y \rfloor - q_* \lfloor n_{\min}x + y \rfloor = 0$, or $q_* \neq q^*$ and $\lfloor (q^* + n_{\min})x + y \rfloor < p^* + \lfloor n_{\min}x + y \rfloor$, then $R = \emptyset$. In this case, return

$$\left(\lfloor n_{\min}x + y \rfloor, n_{\min}, \left\lfloor \frac{vp_*}{q_*} \right\rfloor + \lfloor n_{\min}x + y \rfloor + 1, v + n_{\min} \right).$$

6. Otherwise, if $q_* = q^*$, then let

$$\begin{aligned} \mu &\leftarrow ((\lfloor q_*n_{\min}x + q_*y \rfloor - q_* \lfloor n_{\min}x + y \rfloor + 1)v \bmod q_*), \\ \nu &\leftarrow ((\lfloor q_*n_{\min}x + q_*y \rfloor - q_* \lfloor n_{\min}x + y \rfloor)v \bmod q_*) \end{aligned}$$

and return

$$\left(\left\lfloor \frac{\nu p_*}{q_*} \right\rfloor + \lfloor n_{\min}x + y \rfloor + 1, \nu + n_{\min}, \left\lfloor \frac{\mu p_*}{q_*} \right\rfloor + \lfloor n_{\min}x + y \rfloor + 1, \mu + n_{\min} \right).$$

7. Otherwise, compute μ as follows:

- (a) If $\lceil r_T \rceil - 1 < \lceil r_B \rceil$, then let $(k, r) \leftarrow (\lfloor k_{PR} \rfloor, \lceil r_B \rceil - 1)$.
- (b) Otherwise, check if $\lceil k_{TR} \rceil - 1 > \lfloor k_{TL} \rfloor$ holds. If that is the case, then let $(k, r) \leftarrow (\lceil k_{TR} \rceil - 1, \lceil r_T \rceil - 1)$,
- (c) Otherwise, check if $\lceil r_{TR} \rceil - 1 < \lceil r_B \rceil$ and $\lceil r_{TL} \rceil - 1 < \lceil r_B \rceil$ hold. If that is the case, then let $(k, r) \leftarrow (\lfloor k_{PR} \rfloor, \lceil r_B \rceil - 1)$.

- (d) Otherwise, check if $\lceil r_{TR} \rceil - 1 \geq \lceil r_B \rceil$ and $\lceil r_{TL} \rceil - 1 < \lceil r_B \rceil$ hold. If that is the case, then let $(k, r) \leftarrow (\lceil k_{TRR} \rceil - 1, \lceil r_{TR} \rceil - 1)$.
 - (e) Otherwise, check if $\lceil r_{TR} \rceil - 1 < \lceil r_B \rceil$ and $\lceil r_{TL} \rceil - 1 \geq \lceil r_B \rceil$ hold. If that is the case, then let $(k, r) \leftarrow (\lceil k_{TLR} \rceil - 1, \lceil r_{TL} \rceil - 1)$.
 - (f) Otherwise,
 - i. If $\lceil r_{TR} \rceil - 1 \geq \lceil r_{TL} \rceil - 1$, then let $(k, r) \leftarrow (\lceil k_{TRR} \rceil - 1, \lceil r_{TR} \rceil - 1)$.
 - ii. Otherwise, let $(k, r) \leftarrow (\lceil k_{TLR} \rceil - 1, \lceil r_{TL} \rceil - 1)$.
 - (g) Then, let $\mu \leftarrow q_* k - q^* r$.
8. Then, compute ν as follows:
- (a) If $r'_B \geq \lceil r_T \rceil$, then let $(k, r) \leftarrow (\lfloor k_{PL} \rfloor + 1, \lceil r_T \rceil)$.
 - (b) Otherwise, check if $\lceil k_{BL} \rceil \leq \lfloor k_{BR} \rfloor$ holds. If that is the case, then let $(k, r) \leftarrow (\lceil k_{BL} \rceil, r'_B)$.
 - (c) Otherwise, check if $\lceil r_{BL} \rceil \geq \lceil r_T \rceil$ and $\lceil r_{BR} \rceil \geq \lceil r_T \rceil$ hold. If that is the case, then let $(k, r) \leftarrow (\lfloor k_{PL} \rfloor + 1, \lceil r_T \rceil)$.
 - (d) Otherwise, check if $\lceil r_{BL} \rceil < \lceil r_T \rceil$ and $\lceil r_{BR} \rceil \geq \lceil r_T \rceil$ hold. If that is the case, then let $(k, r) \leftarrow (\lceil k_{BLL} \rceil, \lceil r_{BL} \rceil)$.
 - (e) Otherwise, check if $\lceil r_{BL} \rceil \geq \lceil r_T \rceil$ and $\lceil r_{BR} \rceil < \lceil r_T \rceil$ hold. If that is the case, then let $(k, r) \leftarrow (\lceil k_{BRL} \rceil, \lceil r_{BR} \rceil)$.
 - (f) Otherwise,
 - i. If $\lceil r_{BL} \rceil \leq \lceil r_{BR} \rceil$, then let $(k, r) \leftarrow (\lceil k_{BLL} \rceil, \lceil r_{BL} \rceil)$.
 - ii. Otherwise, let $(k, r) \leftarrow (\lceil k_{BRL} \rceil, \lceil r_{BR} \rceil)$.
 - (g) Then, let $\nu \leftarrow q_* k - q^* r$.
9. Return $\left(\left\lfloor \frac{\nu p_*}{q_*} \right\rfloor + \lfloor n_{\min} x + y \rfloor + 1, \nu + n_{\min}, \left\lfloor \frac{\mu p_*}{q_*} \right\rfloor + \lfloor n_{\min} x + y \rfloor + 1, \mu + n_{\min} \right)$.

9 Lattice points in vertically parallel trapezoidal regions

Combining Theorem 7.4 with Algorithm 8.2, we obtain a complete description of the full subset of \mathbb{R}^2 consisting of all (ξ, ζ) 's satisfying

1. $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$ and
2. $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ for all $n \in [0 : n_{\max} - n_{\min}]$,

and this subset is always of the form

$$T := \{(\xi, \zeta) \in \mathbb{R}^2 : \xi_{\min} \leq \xi < \xi_{\max}, \quad a - b\xi \leq \zeta < c - d\xi\}$$

for some $\xi_{\min}, \xi_{\max} \in \mathbb{Q}$ with $\xi_{\min} < \xi_{\max}$ and $a, b, c, d \in \mathbb{Z}$, which, if drawn in the (ξ, ζ) -plane, is a trapezoid whose two edges are parallel to the vertical axis $\xi = 0$. For use in Algorithm 5.4, we want to find a point in this region T of the form $(\frac{m}{2^k}, \frac{s}{2^k})$ where $m, s \in \mathbb{Z}$ and $k \in \mathbb{Z}_{\geq 0}$.

In this section, we describe how to find such a point. More specifically, we aim to find the smallest k such that a point of the form $(\frac{m}{2^k}, \frac{s}{2^k})$ exists in T , and when such k is found, we want to enumerate all points in T of that form.

As the first step, we find the smallest $k \in \mathbb{Z}_{\geq 0}$ such that there exists some $m \in \mathbb{Z}$ satisfying

$$\xi_{\min} \leq \frac{m}{2^k} < \xi_{\max}.$$

This can be done as follows. We start with computing $k_0 := \max \{-\lfloor \log_2(\xi_{\max} - \xi_{\min}) \rfloor, 0\}$. Computation of $\max \left\{ \left\lfloor \log_2 \frac{p}{q} \right\rfloor, 0 \right\}$ for a rational number $\frac{p}{q}$ can be easily done in integer arithmetic as it is basically just computing the difference of bit-widths between p and q . Then since

$$\xi_{\max} - \xi_{\min} \geq 2^{\lfloor \log_2(\xi_{\max} - \xi_{\min}) \rfloor} \geq \frac{1}{2^{k_0}},$$

there must exist a number of the form $\frac{m_0}{2^{k_0}}$ in $[\xi_{\min}, \xi_{\max})$, and specifically $m_0 = \lceil 2^{k_0} \xi_{\min} \rceil$ satisfies that. On the other hand, if $k_0 > 0$, then since

$$\xi_{\max} - \xi_{\min} < 2^{\lfloor \log_2(\xi_{\max} - \xi_{\min}) \rfloor + 1} = \frac{1}{2^{k_0 - 1}},$$

there are at most 2 numbers of the form $\frac{m_0}{2^{k_0}}$ in $[\xi_{\min}, \xi_{\max})$. If there is only one element of such form and it is odd, then $k = k_0$ should be the smallest integer we are looking for. Otherwise, there should uniquely an even integer m_0 such that $\frac{m_0}{2^{k_0}} \in [\xi_{\min}, \xi_{\max})$. If we write $m_0 = 2^\tau \sigma$, then $k = \max \{k_0 - \tau, 0\}$ is the smallest integer we are looking for.

Then for any integer $m \in [\lceil 2^k \xi_{\min} \rceil : \lceil 2^k \xi_{\max} \rceil - 1]$ (in other words, for any integer m such that $\frac{m}{2^k} \in [\xi_{\min}, \xi_{\max})$), if we set $\xi = \frac{m}{2^k}$, then the corresponding bound on ζ becomes

$$\frac{2^k a - bm}{2^k} \leq \zeta < \frac{2^k c - dm}{2^k}.$$

Therefore, for an integer s , we have $(\frac{m}{2^k}, \frac{s}{2^k}) \in T$ if and only if $2^k a - bm \leq s < 2^k c - dm$.

However, that it is possible that the range of s we get here is empty. When that happens for every m , then we increase k sufficiently to make sure there is at least one m yielding a nonempty range for s . Note that by Theorem 8.1, the integers a, b, c, d are determined independently to ξ , and Theorem 7.6 and the construction of a, b, c, d ensure that $a - b\xi \leq c - d\xi$ always holds for any $\xi \in [\xi_{\min}, \xi_{\max})$. Hence, the only case when the range of s becomes empty is when $2^k a - bm = 2^k c - dm$ holds.

We claim that this linear equation $2^k a - bm = 2^k c - dm$ on m can have at most 1 solution. Indeed, it has more than 1 solutions if and only if $a = c$ and $b = d$ both hold. However, this means that $a - b\xi = c - d\xi$ holds for *every* ξ , which is clearly impossible because applying Theorem 8.1 to $\xi \leftarrow x$ and $\zeta \leftarrow y$ shows that $a - bx \leq y < c - dx$ holds.

Therefore, all we need to do is to find the smallest k' larger than the current k such that there are at least 2 distinct numbers of the form $\frac{m'}{2^{k'}}$ in $[\xi_{\min}, \xi_{\max})$. In particular, if there already were 2 or more distinct numbers of the form $\frac{m}{2^k}$ in $[\xi_{\min}, \xi_{\max})$, then at least one of them should have yielded a nonempty range for s . Hence, let m be the unique integer such that $\frac{m}{2^k} \in [\xi_{\min}, \xi_{\max})$. If $\frac{m}{2^k} - \xi_{\min} \geq \frac{m}{2^k} - \xi_{\max}$, then we find the smallest $k' > k$ such that

$$\frac{1}{2^{k'}} \leq \frac{m}{2^k} - \xi_{\min},$$

that is,

$$k' := - \left\lfloor \log_2 \left(\frac{m}{2^k} - \xi_{\min} \right) \right\rfloor,$$

so that $\frac{2^{k'} - k - 1}{2^{k'}} \in [\xi_{\min}, \xi_{\max})$ holds. Otherwise, we find the smallest $k' > k$ such that

$$\frac{1}{2^{k'}} < \xi_{\max} - \frac{m}{2^k},$$

that is,

$$k' := 1 - \left\lceil \log_2 \left(\xi_{\max} - \frac{m}{2^k} \right) \right\rceil,$$

so that $\frac{2^{k'-k}m+1}{2^{k'}} \in [\xi_{\min}, \xi_{\max})$ holds. Formally writing all of these in details then gives us the following algorithm.

Algorithm 9.1 (Finding lattice points in vertically parallel trapezoidal regions).

1. Let $k \leftarrow \max \{ -\lfloor \log_2(\xi_{\max} - \xi_{\min}) \rfloor, 0 \}$ and $m \leftarrow \lceil 2^k \xi_{\min} \rceil$.
2. If m is even, then factorize it into $m = 2^\tau \sigma$, and then let $m \leftarrow m/2^{\min\{k, \tau\}}$ and $k \leftarrow \max \{k - \tau, 0\}$.
3. Otherwise, if $m + 1 < \lceil 2^k \xi_{\max} \rceil$, then factorize $m + 1$ into $m + 1 = 2^\tau \sigma$, and then let $m \leftarrow (m + 1)/2^{\min\{k, \tau\}}$ and $k \leftarrow \max \{k - \tau, 0\}$.
4. Until m reaches $\lceil 2^k \xi_{\max} \rceil$,

(a) Let

$$s_{\min} \leftarrow 2^k a - mb \quad \text{and} \quad s_{\max} \leftarrow 2^k c - md - 1.$$

(b) If $s_{\min} \leq s_{\max}$, then record $(k, m, [s_{\min} : s_{\max}])$.

(c) Let $m \leftarrow m + 1$.

5. If there was at least one m that yielded $s_{\min} \leq s_{\max}$, then we are done, so return.
6. Otherwise, let $m \leftarrow m - 1$ (which must be equal to $\lceil 2^k \xi_{\max} \rceil - 1 = \lceil 2^k \xi_{\min} \rceil$).
7. If $m - 2^k \xi_{\min} \geq 2^k \xi_{\max} - m$, then let $k \leftarrow k - \lfloor \log_2(m - 2^k \xi_{\min}) \rfloor$. Otherwise, let $k \leftarrow k + 1 - \lceil \log_2(2^k \xi_{\max} - m) \rceil$.
8. Let $m \leftarrow \lceil 2^k \xi_{\min} \rceil$ and go back to Step 4. This time Step 5 must succeed to return.

10 Extremizers of $\frac{\lfloor nx+y \rfloor - \zeta}{n}$

So far, we have completed our explanation on how to find the smallest minimizer and the largest maximizer of the fractional part $(nx + y) - \lfloor nx + y \rfloor$ (except for the explanation of the Gosper's algorithm that will be developed in Section 12).

Taking it as granted, in this section we now describe an algorithm for finding extremizers of $\frac{\lfloor nx+y \rfloor - \zeta}{n}$ over an arbitrary interval $[n_{\min} : n_{\max}]$ minus 0, provided that ζ is an arbitrary fixed number in addition to x, y . Recall that having $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$ is equivalent to the inequalities

$$\max_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{>0}} \frac{\lfloor nx + y \rfloor - \zeta}{n} \leq \xi < \min_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{>0}} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n}$$

and

$$\max_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{<0}} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} < \xi \leq \min_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{<0}} \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

possibly along with $\lfloor y \rfloor = \lfloor \zeta \rfloor$ if $0 \in [n_{\min} : n_{\max}]$, so once we now how to find extremizers of $\frac{\lfloor nx+y \rfloor - \zeta}{n}$, we can immediately find a necessary and sufficient condition on ξ for having $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$, when x, y, ζ are all given real numbers.

In addition to that this problem of finding ξ when ζ is given has its own applications (TODO: elaborate), our main algorithm for finding a necessary and sufficient condition on both ξ and ζ , which we develop in Section 11, is based on it.

Throughout this section, let x, y, ζ be real numbers and n_{\min}, n_{\max} be integers such that $n_{\min} \leq n_{\max}$. Until Section 10.4, we will also assume $0 < n_{\min}$.

10.1 The maximization problem

We consider the maximization problem

$$\max_{n \in [n_{\min} : n_{\max}]} \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

first. Since

$$\frac{\lfloor nx + y \rfloor - \zeta}{n} = x - \frac{(nx + y) - \lfloor nx + y \rfloor}{n} - \frac{\zeta - y}{n},$$

our maximization problem can be equivalently written as the minimization problem

$$\min_{n \in [n_{\min} : n_{\max}]} \left(\frac{(nx + y) - \lfloor nx + y \rfloor}{n} + \frac{\zeta - y}{n} \right).$$

Since $\zeta - y$ is supposed to be a small constant, it is a reasonable guess that the *smallest minimizer* of $(nx + y) - \lfloor nx + y \rfloor$ will be a good approximation of the maximizer of $\frac{\lfloor nx + y \rfloor - \zeta}{n}$. Hence, we start from there.

Let n_0 be the smallest minimizer of $(nx + y) - \lfloor nx + y \rfloor$ for $n \in [n_{\min} : n_{\max}]$. Then, the problem of maximizing $\frac{\lfloor nx + y \rfloor - \zeta}{n}$ can be solved by separately solving two maximization problems

$$\max_{n \in [n_{\min} : n_0]} \frac{\lfloor nx + y \rfloor - \zeta}{n} \tag{8}$$

and

$$\max_{n \in [n_0 : n_{\max}]} \frac{\lfloor nx + y \rfloor - \zeta}{n}. \tag{9}$$

The following two theorems, when combined with Theorem 7.1, Theorem 2.4 and Algorithm 4.12, give complete solutions to these two problems, respectively.

Theorem 10.1 (The smallest maximizer on the left).

Let $n_0^{(0)} := n_0$. For $i \in \mathbb{Z}_{>0}$, inductively define $n_0^{(i)} := n_0^{(i-1)} - n_1^{(i)}$ and

$$n_1^{(i)} := \begin{cases} \max_{n \in [1 : n_0^{(i-1)} - n_{\min}]} \operatorname{argmin} \frac{\lfloor nx \rfloor + 1}{n} & \text{if } n_0^{(i-1)} > n_{\min}, \\ 0 & \text{otherwise.} \end{cases}$$

Let k be the smallest nonnegative integer such that either $n_0^{(k)} = n_{\min}$ or

$$\frac{\lfloor n_1^{(k+1)} x \rfloor + 1}{n_1^{(k+1)}} > \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}}$$

holds. Then

$$n_0^{(k)} = \min_{n \in [n_{\min} : n_0]} \operatorname{argmax} \frac{\lfloor nx + y \rfloor - \zeta}{n}.$$

Proof. First, we claim that

$$\left\lfloor (n_0^{(i)} - n)x + y \right\rfloor = \left\lfloor n_0^{(i)}x + y \right\rfloor - \lfloor nx \rfloor - 1$$

holds for all $i \in \mathbb{Z}_{\geq 0}$ and $n \in [1 : n_0^{(i)} - n_{\min}]$. We use induction on i . The base case $i = 0$ is precisely Lemma 6.3, so we only prove the induction step. Since $n_0^{(i)} = n_0^{(i-1)} - n_1^{(i)}$, the induction hypothesis implies

$$\left\lfloor (n_0^{(i)} - n)x + y \right\rfloor = \left\lfloor n_0^{(i-1)}x + y \right\rfloor - \left\lfloor (n_1^{(i)} + n)x \right\rfloor - 1.$$

On the other hand, since

$$n_1^{(i)} + n \leq n_1^{(i)} + n_0^{(i)} - n_{\min} = n_0^{(i-1)} - n_{\min},$$

Corollary 7.3 implies

$$\left\lfloor (n_1^{(i)} + n)x \right\rfloor = \left\lfloor n_1^{(i)}x \right\rfloor + \lfloor nx \rfloor + 1.$$

Again by the induction hypothesis, we know

$$\left\lfloor n_0^{(i)}x + y \right\rfloor = \left\lfloor n_0^{(i-1)}x + y \right\rfloor - \left\lfloor n_1^{(i)}x \right\rfloor - 1, \quad (10)$$

so combining these completes the induction step. Hence, the claim is proved.

Next, we claim that for each $i = 0, \dots, k$,

$$n_0^{(i)} = \min_{n \in [n_0^{(i)} : n_0]} \operatorname{argmax} \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

holds. We use induction on i . The base case $i = 0$ is trivial, so we show the induction step. By the induction hypothesis, we have

$$\frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} \geq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0^{(i-1)} : n_0]$. Also, note that

$$\begin{aligned} & \frac{\left\lfloor n_0^{(i)}x + y \right\rfloor - \zeta}{n_0^{(i)}} - \frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} \\ &= \frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \left\lfloor n_1^{(i)}x \right\rfloor - 1 - \zeta}{n_0^{(i)}} - \frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} \\ &= \frac{n_1^{(i)}}{n_0^{(i)}} \left(\frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} - \frac{\left\lfloor n_1^{(i)}x \right\rfloor + 1}{n_1^{(i)}} \right), \end{aligned}$$

and by the definition of k , the right-hand side is nonnegative. Therefore,

$$\frac{\left\lfloor n_0^{(i)}x + y \right\rfloor - \zeta}{n_0^{(i)}} \geq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

holds for all $n \in [n_0^{(i-1)} : n_0]$.

On the other hand, the previous claim shows that for any $n \in [1 : n_1^{(i)}]$,

$$\begin{aligned} & \frac{\lfloor n_0^{(i)} x + y \rfloor - \zeta}{n_0^{(i)}} - \frac{\lfloor (n_0^{(i-1)} - n)x + y \rfloor - \zeta}{n_0^{(i-1)} - n} \\ &= \frac{\lfloor n_0^{(i-1)} x + y \rfloor - \lfloor n_1^{(i)} x \rfloor - 1 - \zeta}{n_0^{(i)}} - \frac{\lfloor n_0^{(i-1)} x + y \rfloor - \lfloor nx \rfloor - 1 - \zeta}{n_0^{(i-1)} - n} \\ &= \frac{n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} - n)} \cdot \frac{\lfloor n_0^{(i-1)} x + y \rfloor - \zeta}{n_0^{(i-1)}} + \frac{\lfloor nx \rfloor + 1}{n_0^{(i-1)} - n} - \frac{\lfloor n_1^{(i)} x \rfloor + 1}{n_0^{(i)}}. \end{aligned}$$

Recall that by the definition of k , we have

$$\frac{\lfloor n_0^{(i-1)} x + y \rfloor - \zeta}{n_0^{(i-1)}} \geq \frac{\lfloor n_1^{(i)} x \rfloor + 1}{n_1^{(i)}},$$

thus we get

$$\begin{aligned} & \frac{\lfloor n_0^{(i)} x + y \rfloor - \zeta}{n_0^{(i)}} - \frac{\lfloor (n_0^{(i-1)} - n)x + y \rfloor - \zeta}{n_0^{(i-1)} - n} \\ & \geq \frac{\lfloor nx \rfloor + 1}{n_0^{(i-1)} - n} - \frac{n_1^{(i)}(n_0^{(i-1)} - n) - n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} - n)} \cdot \frac{\lfloor n_1^{(i)} x \rfloor + 1}{n_1^{(i)}} \\ & = \frac{n}{n_0^{(i-1)} - n} \left(\frac{\lfloor nx \rfloor + 1}{n} - \frac{\lfloor n_1^{(i)} x \rfloor + 1}{n_1^{(i)}} \right) \geq 0 \end{aligned}$$

by the definition of $n_1^{(i)}$. As a result, we conclude

$$\frac{\lfloor n_0^{(i)} x + y \rfloor - \zeta}{n_0^{(i)}} \geq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0^{(i)} : n_0]$, completing the induction step. Therefore, the claim is proved.

Consequently, it is enough to show that for any $n \in [1 : n_0^{(k)} - n_{\min}]$, we have

$$\frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} > \frac{\lfloor (n_0^{(k)} - n)x + y \rfloor - \zeta}{n_0^{(k)} - n}.$$

This is vacuously true if $n_0^{(k)} = n_{\min}$, so suppose otherwise, thus

$$\frac{\lfloor n_1^{(k+1)} x \rfloor + 1}{n_1^{(k+1)}} > \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}}$$

holds. By the first claim, for any $n \in [1: n_0^{(k)} - n_{\min}]$, we have

$$\begin{aligned} & \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} - \frac{\lfloor (n_0^{(k)} - n)x + y \rfloor - \zeta}{n_0^{(k)} - n} \\ &= \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} - \frac{\lfloor n_0^{(k)} x + y \rfloor - \lfloor nx \rfloor - 1 - \zeta}{n_0^{(k)} - n} \\ &= \frac{n}{n_0^{(k)} - n} \left(\frac{\lfloor nx \rfloor + 1}{n} - \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} \right), \end{aligned}$$

and by the definition of $n_1^{(k+1)}$, the right-hand side is strictly positive. Therefore, we are done. \square

Theorem 10.2 (The largest maximizer on the right).

Let $n_0^{(0)} := n_0$. For $i \in \mathbb{Z}_{>0}$, inductively define $n_0^{(i)} := n_0^{(i-1)} + n_1^{(i)}$ and

$$n_1^{(i)} := \begin{cases} \max_{n \in [1: n_{\max} - n_0^{(i-1)}]} \operatorname{argmax} \frac{\lfloor nx \rfloor}{n} & \text{if } n_0^{(i-1)} < n_{\max}, \\ 0 & \text{otherwise.} \end{cases}$$

Let k be the smallest nonnegative integer such that either $n_0^{(k)} = n_{\max}$ or

$$\frac{\lfloor n_1^{(k+1)} x \rfloor}{n_1^{(k+1)}} < \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}}$$

holds. Then

$$n_0^{(k)} = \max_{n \in [n_0: n_{\max}]} \operatorname{argmax} \frac{\lfloor nx + y \rfloor - \zeta}{n}.$$

Proof. First, we claim that

$$\lfloor (n_0^{(i)} + n)x + y \rfloor = \lfloor n_0^{(i)} x + y \rfloor + \lfloor nx \rfloor$$

holds for all $i \in \mathbb{Z}_{\geq 0}$ and $n \in [0: n_{\max} - n_0^{(i)}]$. We use induction on i . The base case $i = 0$ directly follows from Lemma 6.1, so we only prove the induction step. Since $n_0^{(i)} = n_0^{(i-1)} + n_1^{(i)}$, the induction hypothesis implies

$$\lfloor (n_0^{(i)} + n)x + y \rfloor = \lfloor n_0^{(i-1)} x + y \rfloor + \lfloor (n_1^{(i)} + n)x \rfloor.$$

On the other hand, since

$$n_1^{(i)} + n \leq n_1^{(i)} + n_{\max} - n_0^{(i)} = n_{\max} - n_0^{(i-1)},$$

Corollary 7.3 implies

$$\lfloor (n_1^{(i)} + n)x \rfloor = \lfloor n_1^{(i)} x \rfloor + \lfloor nx \rfloor.$$

Again by the induction hypothesis, we know

$$\lfloor n_0^{(i)}x + y \rfloor = \lfloor n_0^{(i-1)}x + y \rfloor + \lfloor n_1^{(i)}x \rfloor, \quad (11)$$

so combining these completes the induction step. Hence, the claim is proved.

Next, we claim that for each $i = 0, \dots, k$,

$$n_0^{(i)} = \max_{n \in [n_0: n_0^{(i)}]} \operatorname{argmax} \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

holds. We use induction on i . The base case $i = 0$ is trivial, so we show the induction step. By the induction hypothesis, we have

$$\frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}} \geq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0: n_0^{(i-1)}]$. On the other hand, the previous claim shows that for any $n \in [0: n_1^{(i)}]$,

$$\begin{aligned} & \frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} - \frac{\lfloor (n_0^{(i-1)} + n)x + y \rfloor - \zeta}{n_0^{(i-1)} + n} \\ &= \frac{\lfloor n_0^{(i-1)}x + y \rfloor + \lfloor n_1^{(i)}x \rfloor - \zeta}{n_0^{(i)}} - \frac{\lfloor n_0^{(i-1)}x + y \rfloor + \lfloor nx \rfloor - \zeta}{n_0^{(i-1)} + n} \\ &= \frac{\lfloor n_1^{(i)}x \rfloor}{n_0^{(i)}} - \frac{\lfloor nx \rfloor}{n_0^{(i-1)} + n} - \frac{n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} + n)} \cdot \frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}}. \end{aligned}$$

Recall that by the definition of k , we have

$$\frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}} \leq \frac{\lfloor n_1^{(i)}x \rfloor}{n_1^{(i)}},$$

thus we get

$$\begin{aligned} & \frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} - \frac{\lfloor (n_0^{(i-1)} + n)x + y \rfloor - \zeta}{n_0^{(i-1)} + n} \\ & \geq \frac{n_1^{(i)}(n_0^{(i-1)} + n) - n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} + n)} \cdot \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_1^{(i)}} - \frac{\lfloor nx \rfloor}{n_0^{(i-1)} + n} \\ & = \frac{n \lfloor n_1^{(i)}x \rfloor}{n_1^{(i)}(n_0^{(i-1)} + n)} - \frac{\lfloor nx \rfloor}{n_0^{(i-1)} + n}. \end{aligned}$$

If $n = 0$, then the above is equal to zero, and if $n \neq 0$, then it is equal to

$$\frac{n}{n_0^{(i-1)} + n} \left(\frac{\lfloor n_1^{(i)}x \rfloor}{n_1^{(i)}} - \frac{\lfloor nx \rfloor}{n} \right) \geq 0$$

by the definition of $n_1^{(i)}$. As a result, we conclude

$$\frac{\lfloor n_0^{(i)} x + y \rfloor - \zeta}{n_0^{(i)}} \geq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0 : n_0^{(i)}]$, completing the induction step. Therefore, the claim is proved.

Consequently, it is enough to show that for any $n \in [1 : n_{\max} - n_0^{(k)}]$, we have

$$\frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} > \frac{\lfloor (n_0^{(k)} + n)x + y \rfloor - \zeta}{n_0^{(k)} + n}.$$

This is vacuously true if $n_0^{(k)} = n_{\max}$, so suppose otherwise, thus

$$\frac{\lfloor n_1^{(k+1)} x \rfloor}{n_1^{(k+1)}} < \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}}$$

holds. By the first claim, for any $n \in [1 : n_{\max} - n_0^{(k)}]$, we have

$$\begin{aligned} & \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} - \frac{\lfloor (n_0^{(k)} + n)x + y \rfloor - \zeta}{n_0^{(k)} + n} \\ &= \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} - \frac{\lfloor n_0^{(k)} x + y \rfloor + \lfloor nx \rfloor - \zeta}{n_0^{(k)} + n} \\ &= \frac{n}{n_0^{(k)} + n} \left(\frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}} - \frac{\lfloor nx \rfloor}{n} \right), \end{aligned}$$

and by the definition of $n_1^{(k+1)}$, the right-hand side is strictly positive. Therefore, we are done. \square

10.2 The minimization problem

Next, we consider the minimization problem

$$\min_{n \in [n_{\min} : n_{\max}]} \frac{\lfloor nx + y \rfloor - \zeta}{n}.$$

Since

$$\frac{\lfloor nx + y \rfloor - \zeta}{n} = x - \frac{(nx + y) - \lfloor nx + y \rfloor}{n} - \frac{\zeta - y}{n},$$

our minimization problem can be equivalently written as the maximization problem

$$\max_{n \in [n_{\min} : n_{\max}]} \left(\frac{(nx + y) - \lfloor nx + y \rfloor}{n} + \frac{\zeta - y}{n} \right).$$

Again, since $\zeta - y$ is supposed to be a small constant, it is a reasonable guess that the n the *largest maximizer* of $(nx + y) - \lfloor nx + y \rfloor$ will be a good approximation of the minimizer of $\frac{\lfloor nx + y \rfloor - \zeta}{n}$. Hence, we start from there.

Let n_0 be the largest maximizer of $(nx + y) - \lfloor nx + y \rfloor$ for $n \in [n_{\min} : n_{\max}]$. Then, the problem of minimizing $\frac{\lfloor nx + y \rfloor - \zeta}{n}$ can be solved by separately solving two minimization problems

$$\min_{n \in [n_{\min} : n_0]} \frac{\lfloor nx + y \rfloor - \zeta}{n} \quad (12)$$

and

$$\min_{n \in [n_0 : n_{\max}]} \frac{\lfloor nx + y \rfloor - \zeta}{n}. \quad (13)$$

The following two theorems, when combined with Theorem 7.1, Theorem 2.4 and Algorithm 4.12, give complete solutions to these two problems, respectively.

Theorem 10.3 (The smallest minimizer on the left).

Let $n_0^{(0)} := n_0$. For $i \in \mathbb{Z}_{>0}$, inductively define $n_0^{(i)} := n_0^{(i-1)} - n_1^{(i)}$ and

$$n_1^{(i)} := \begin{cases} \max_{n \in [1 : n_0^{(i-1)} - n_{\min}]} \operatorname{argmax} \frac{\lfloor nx \rfloor}{n} & \text{if } n_0^{(i-1)} > n_{\min}, \\ 0 & \text{otherwise.} \end{cases}$$

Let k be the smallest nonnegative integer such that either $n_0^{(k)} = n_{\min}$ or

$$\frac{\lfloor n_1^{(k+1)} x \rfloor}{n_1^{(k+1)}} < \frac{\lfloor n_0^{(k)} x + y \rfloor - \zeta}{n_0^{(k)}}$$

holds. Then

$$n_0^{(k)} = \min_{n \in [n_{\min} : n_0]} \operatorname{argmin} \frac{\lfloor nx + y \rfloor - \zeta}{n}.$$

Proof. First, we claim that

$$\lfloor (n_0^{(i)} - n)x + y \rfloor = \lfloor n_0^{(i)} x + y \rfloor - \lfloor nx \rfloor$$

holds for all $i \in \mathbb{Z}_{\geq 0}$ and $n \in [0 : n_0^{(i)} - n_{\min}]$. We use induction on i . The base case $i = 0$ directly follows from Lemma 6.2, so we only prove the induction step. Since $n_0^{(i)} = n_0^{(i-1)} - n_1^{(i)}$, the induction hypothesis implies

$$\lfloor (n_0^{(i)} - n)x + y \rfloor = \lfloor n_0^{(i-1)} x + y \rfloor - \lfloor (n_1^{(i)} + n)x \rfloor.$$

On the other hand, since

$$n_1^{(i)} + n \leq n_1^{(i)} + n_0^{(i)} - n_{\min} = n_0^{(i-1)} - n_{\min},$$

Corollary 7.3 implies

$$\lfloor (n_1^{(i)} + n)x \rfloor = \lfloor n_1^{(i)} x \rfloor + \lfloor nx \rfloor.$$

Again by the induction hypothesis, we know

$$\lfloor n_0^{(i)} x + y \rfloor = \lfloor n_0^{(i-1)} x + y \rfloor - \lfloor n_1^{(i)} x \rfloor, \quad (14)$$

so combining these completes the induction step. Hence, the claim is proved.

Next, we claim that for each $i = 0, \dots, k$,

$$n_0^{(i)} = \min_{n \in [n_0^{(i)} : n_0]} \operatorname{argmin} \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

holds. We use induction on i . The base case $i = 0$ is trivial, so we show the induction step. By the induction hypothesis, we have

$$\frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}} \leq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0^{(i-1)} : n_0]$. On the other hand, the previous claim shows that for any $n \in [0 : n_1^{(i)}]$,

$$\begin{aligned} & \frac{\lfloor (n_0^{(i-1)} - n)x + y \rfloor - \zeta}{n_0^{(i-1)} - n} - \frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} \\ &= \frac{\lfloor n_0^{(i-1)}x + y \rfloor - \lfloor nx \rfloor - \zeta}{n_0^{(i-1)} - n} - \frac{\lfloor n_0^{(i-1)}x + y \rfloor - \lfloor n_1^{(i)}x \rfloor - \zeta}{n_0^{(i)}} \\ &= \frac{\lfloor n_1^{(i)}x \rfloor}{n_0^{(i)}} - \frac{\lfloor nx \rfloor}{n_0^{(i-1)} - n} - \frac{n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} - n)} \cdot \frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}}. \end{aligned}$$

Recall that by the definition of k , we have

$$\frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}} \leq \frac{\lfloor n_1^{(i)}x \rfloor}{n_1^{(i)}},$$

thus we get

$$\begin{aligned} & \frac{\lfloor (n_0^{(i-1)} - n)x + y \rfloor - \zeta}{n_0^{(i-1)} - n} - \frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} \\ & \geq \frac{n_1^{(i)}(n_0^{(i-1)} - n) - n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} - n)} \cdot \frac{\lfloor n_1^{(i)}x \rfloor}{n_1^{(i)}} - \frac{\lfloor nx \rfloor}{n_0^{(i-1)} - n} \\ & = \frac{n \lfloor n_1^{(i)}x \rfloor}{n_1^{(i)}(n_0^{(i-1)} - n)} - \frac{\lfloor nx \rfloor}{n_0^{(i-1)} - n}. \end{aligned}$$

If $n = 0$, then the above is equal to zero, and if $n \neq 0$, then it is equal to

$$\frac{n}{n_0^{(i-1)} - n} \left(\frac{\lfloor n_1^{(i)}x \rfloor}{n_1^{(i)}} - \frac{\lfloor nx \rfloor}{n} \right) \geq 0$$

by the definition of $n_1^{(i)}$. As a result, we conclude

$$\frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} \leq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0^{(i)} : n_0]$, completing the induction step. Therefore, the claim is proved.

Consequently, it is enough to show that for any $n \in [1 : n_0^{(k)} - n_{\min}]$, we have

$$\frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} < \frac{\lfloor (n_0^{(k)} - n)x + y \rfloor - \zeta}{n_0^{(k)} - n}.$$

This is vacuously true if $n_0^{(k)} = n_{\min}$, so suppose otherwise, thus

$$\frac{\lfloor n_1^{(k+1)}x \rfloor}{n_1^{(k+1)}} < \frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}}$$

holds. By the first claim, for any $n \in [1 : n_0^{(k)} - n_{\min}]$, we have

$$\begin{aligned} & \frac{\lfloor (n_0^{(k)} - n)x + y \rfloor - \zeta}{n_0^{(k)} - n} - \frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} \\ &= \frac{\lfloor n_0^{(k)}x + y \rfloor - \lfloor nx \rfloor - \zeta}{n_0^{(k)} - n} - \frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} \\ &= \frac{n}{n_0^{(k)} - n} \left(\frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} - \frac{\lfloor nx \rfloor}{n} \right), \end{aligned}$$

and by the definition of $n_1^{(k+1)}$, the right-hand side is strictly positive. Therefore, we are done. \square

Theorem 10.4 (The largest minimizer on the right).

Let $n_0^{(0)} := n_0$. For $i \in \mathbb{Z}_{>0}$, inductively define $n_0^{(i)} := n_0^{(i-1)} + n_1^{(i)}$ and

$$n_1^{(i)} := \begin{cases} \max_{n \in [1 : n_{\max} - n_0^{(i-1)}]} \operatorname{argmin} \frac{\lfloor nx \rfloor + 1}{n} & \text{if } n_0^{(i-1)} < n_{\max}, \\ 0 & \text{otherwise.} \end{cases}$$

Let k be the smallest nonnegative integer such that either $n_0^{(k)} = n_{\max}$ or

$$\frac{\lfloor n_1^{(k+1)}x \rfloor + 1}{n_1^{(k+1)}} > \frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}}$$

holds. Then

$$n_0^{(k)} = \max_{n \in [n_0 : n_{\max}]} \operatorname{argmin} \frac{\lfloor nx + y \rfloor - \zeta}{n}.$$

Proof. First, we claim that

$$\left\lfloor (n_0^{(i)} + n)x + y \right\rfloor = \left\lfloor n_0^{(i)}x + y \right\rfloor + \lfloor nx \rfloor + 1$$

holds for all $i \in \mathbb{Z}_{\geq 0}$ and $n \in [1 : n_{\max} - n_0^{(i)}]$. We use induction on i . The base case $i = 0$ is precisely Lemma 6.4, so we only prove the induction step. Since $n_0^{(i)} = n_0^{(i-1)} + n_1^{(i)}$, the induction hypothesis implies

$$\left\lfloor (n_0^{(i)} + n)x + y \right\rfloor = \left\lfloor n_0^{(i-1)}x + y \right\rfloor + \left\lfloor (n_1^{(i)} + n)x \right\rfloor + 1.$$

On the other hand, since

$$n_1^{(i)} + n \leq n_1^{(i)} + n_{\max} - n_0^{(i)} = n_{\max} - n_0^{(i-1)},$$

Corollary 7.3 implies

$$\left\lfloor (n_1^{(i)} + n)x \right\rfloor = \left\lfloor n_1^{(i)}x \right\rfloor + \lfloor nx \rfloor + 1.$$

Again by the induction hypothesis, we know

$$\left\lfloor n_0^{(i)}x + y \right\rfloor = \left\lfloor n_0^{(i-1)}x + y \right\rfloor + \left\lfloor n_1^{(i)}x \right\rfloor + 1, \quad (15)$$

so combining these completes the induction step. Hence, the claim is proved.

Next, we claim that for each $i = 0, \dots, k$,

$$n_0^{(i)} = \max_{n \in [n_0 : n_0^{(i)}]} \operatorname{argmin}_n \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

holds. We use induction on i . The base case $i = 0$ is trivial, so we show the induction step. By the induction hypothesis, we have

$$\frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} \leq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0 : n_0^{(i-1)}]$. Also, note that

$$\begin{aligned} & \frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} - \frac{\left\lfloor n_0^{(i)}x + y \right\rfloor - \zeta}{n_0^{(i)}} \\ &= \frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} - \frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor + \left\lfloor n_1^{(i)}x \right\rfloor + 1 - \zeta}{n_0^{(i)}} \\ &= \frac{n_1^{(i)}}{n_0^{(i)}} \left(\frac{\left\lfloor n_0^{(i-1)}x + y \right\rfloor - \zeta}{n_0^{(i-1)}} - \frac{\left\lfloor n_1^{(i)}x \right\rfloor + 1}{n_1^{(i)}} \right), \end{aligned}$$

and by the definition of k , the right-hand side is nonnegative. Therefore,

$$\frac{\left\lfloor n_0^{(i)}x + y \right\rfloor - \zeta}{n_0^{(i)}} \leq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

holds for all $n \in [n_0 : n_0^{(i-1)}]$.

On the other hand, the previous claim shows that for any $n \in [1 : n_1^{(i)}]$,

$$\begin{aligned} & \frac{\lfloor (n_0^{(i-1)} + n)x + y \rfloor - \zeta}{n_0^{(i-1)} + n} - \frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} \\ &= \frac{\lfloor n_0^{(i-1)}x + y \rfloor + \lfloor nx \rfloor + 1 - \zeta}{n_0^{(i-1)} + n} - \frac{\lfloor n_0^{(i-1)}x + y \rfloor + \lfloor n_1^{(i)}x \rfloor + 1 - \zeta}{n_0^{(i)}} \\ &= \frac{n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} + n)} \cdot \frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}} + \frac{\lfloor nx \rfloor + 1}{n_0^{(i-1)} + n} - \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_0^{(i)}}. \end{aligned}$$

Recall that by the definition of k , we have

$$\frac{\lfloor n_0^{(i-1)}x + y \rfloor - \zeta}{n_0^{(i-1)}} \geq \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_1^{(i)}},$$

thus we get

$$\begin{aligned} & \frac{\lfloor (n_0^{(i-1)} + n)x + y \rfloor - \zeta}{n_0^{(i-1)} + n} - \frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} \\ & \geq \frac{\lfloor nx \rfloor + 1}{n_0^{(i-1)} + n} - \frac{n_1^{(i)}(n_0^{(i-1)} + n) - n_0^{(i-1)}(n_1^{(i)} - n)}{n_0^{(i)}(n_0^{(i-1)} + n)} \cdot \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_1^{(i)}} \\ & = \frac{n}{n_0^{(i-1)} + n} \left(\frac{\lfloor nx \rfloor + 1}{n} - \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_1^{(i)}} \right) \geq 0 \end{aligned}$$

by the definition of $n_1^{(i)}$. As a result, we conclude

$$\frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} \leq \frac{\lfloor nx + y \rfloor - \zeta}{n}$$

for all $n \in [n_0 : n_0^{(i)}]$, completing the induction step. Therefore, the claim is proved.

Consequently, it is enough to show that for any $n \in [1 : n_{\max} - n_0^{(k)}]$, we have

$$\frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} < \frac{\lfloor (n_0^{(k)} + n)x + y \rfloor - \zeta}{n_0^{(k)} + n}.$$

This is vacuously true if $n_0^{(k)} = n_{\max}$, so suppose otherwise, thus

$$\frac{\lfloor n_1^{(k+1)}x \rfloor + 1}{n_1^{(k+1)}} > \frac{\lfloor n_0^{(k)}x + y \rfloor + 1 - \zeta}{n_0^{(k)}}$$

holds. By the first claim, for any $n \in [1: n_{\max} - n_0^{(k)}]$, we have

$$\begin{aligned} & \frac{\lfloor (n_0^{(k)} + n)x + y \rfloor - \zeta}{n_0^{(k)} + n} - \frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} \\ &= \frac{\lfloor n_0^{(k)}x + y \rfloor + \lfloor nx \rfloor + 1 - \zeta}{n_0^{(k)} + n} - \frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} \\ &= \frac{n}{n_0^{(k)} + n} \left(\frac{\lfloor nx \rfloor + 1}{n} - \frac{\lfloor n_0^{(k)}x + y \rfloor - \zeta}{n_0^{(k)}} \right), \end{aligned}$$

and by the definition of $n_1^{(k+1)}$, the right-hand side is strictly positive. Therefore, we are done. \square

10.3 Evaluating the inequalities on ζ

In order to apply the aforementioned theorems for actually computing the extremizers, we should evaluate inequalities of the form

$$\frac{a}{n_1} > \frac{b - \zeta}{c} \quad \text{or} \quad \frac{a}{n_1} < \frac{b - \zeta}{c}$$

where $a, b \in \mathbb{Z}$, $c \in \mathbb{Z}_{>0}$, and $n_1 \in [1: n_{\max} - n_{\min}]$. Assuming that ζ is an arbitrary real number, evaluating these inequalities is not completely trivial.

To explain how to do these computations, we first rewrite them into

$$n_1\zeta > n_1b - ac \quad \text{and} \quad n_1\zeta < n_1b - ac.$$

Since the right-hand sides are integers, these inequalities are equivalent to

$$\lceil n_1\zeta \rceil > n_1b - ac \quad \text{and} \quad \lfloor n_1\zeta \rfloor < n_1b - ac.$$

Since the only terms involving ζ are of the form $\lfloor n_1\zeta \rfloor$ or $\lceil n_1\zeta \rceil$, we can use Theorem 7.4 here. Let $\frac{p_*}{q_*}, \frac{p^*}{q^*}$ be the best rational approximations of ζ from below and from above, respectively, with the largest denominators $q_*, q^* \leq n_{\max} - n_{\min}$. Then by Theorem 7.4, we know $\lfloor n_1\zeta \rfloor = \left\lfloor \frac{n_1p_*}{q_*} \right\rfloor$, so we know

$$n_1\zeta < n_1b - ac \quad \text{if and only if} \quad \left\lfloor \frac{n_1p_*}{q_*} \right\rfloor < n_1b - ac.$$

Similarly, since $\frac{-p^*}{q^*}$ is the best rational approximation of $-\zeta$ from below with the largest denominator $q^* \leq n_{\max} - n_{\min}$, Theorem 7.4 tells us that

$$\lceil n_1\zeta \rceil = -\lfloor n_1(-\zeta) \rfloor = -\left\lfloor \frac{-np^*}{q^*} \right\rfloor = \left\lceil \frac{np^*}{q^*} \right\rceil.$$

Hence, we know

$$n_1\zeta > n_1b - ac \quad \text{if and only if} \quad \left\lceil \frac{n_1p^*}{q^*} \right\rceil > n_1b - ac.$$

Therefore, evaluations of these inequalities can be done in arbitrary-precision exact rational arithmetic.

In fact, there is another inequality involving ζ we need to evaluate: we have to choose between the maximizers/minimizers we get from Theorem 10.1 and Theorem 10.2, and Theorem 10.3 and Theorem 10.4, respectively. These are simply the matter of evaluating an inequality of the form

$$\frac{\lfloor n_L x + y \rfloor - \zeta}{n_L} \leq \frac{\lfloor n_R x + y \rfloor - \zeta}{n_R}$$

for some $n_L, n_R \in [n_{\min} : n_{\max}]$ with $n_L \leq n_R$. Here we are assuming that we can already evaluate $\lfloor n_L x + y \rfloor$ and $\lfloor n_R x + y \rfloor$ using Algorithm 8.2. Note that the above inequality is equivalent to

$$n_R \lfloor n_L x + y \rfloor - n_L \lfloor n_R x + y \rfloor \leq (n_R - n_L) \zeta,$$

which is equivalent to

$$n_R \lfloor n_L x + y \rfloor - n_L \lfloor n_R x + y \rfloor \leq \lfloor (n_R - n_L) \zeta \rfloor.$$

Since $n_R - n_L \in [0 : n_{\max} - n_{\min}]$, again evaluating the above inequality can be done by evaluating the inequality

$$n_R \lfloor n_L x + y \rfloor - n_L \lfloor n_R x + y \rfloor \leq \left\lfloor \frac{(n_R - n_L) p_*}{q_*} \right\rfloor$$

instead.

10.4 The case $n_{\min} \leq n_{\max} < 0$ and other variations

We have not yet talked about the case $n_{\min} \leq n_{\max} < 0$, but this case can be actually covered by what we have already covered. The idea is simple: we consider $-x$ instead of x . Indeed, if $n_{\min} \leq n_{\max} < 0$, then

$$\max_{n \in [n_{\min} : n_{\max}]} \frac{\lfloor nx + y \rfloor - \zeta}{n} = - \min_{n \in [-n_{\max} : -n_{\min}]} \frac{\lfloor n(-x) + y \rfloor - \zeta}{n},$$

and similarly

$$\min_{n \in [n_{\min} : n_{\max}]} \frac{\lfloor nx + y \rfloor - \zeta}{n} = - \max_{n \in [-n_{\max} : -n_{\min}]} \frac{\lfloor n(-x) + y \rfloor - \zeta}{n}.$$

Since we have not put any restrictions on the sign of x , we can apply previous arguments without any problem.

In a similar way, we can also solve the optimization problems

$$\max_{n \in [n_{\min} : n_{\max}]} \frac{\lceil nx + y \rceil - \zeta}{n} \quad \text{and} \quad \min_{n \in [n_{\min} : n_{\max}]} \frac{\lceil nx + y \rceil - \zeta}{n}.$$

In this case, we note

$$\frac{\lceil nx + y \rceil - \zeta}{n} = \frac{-\lfloor n(-x) - y \rfloor - \zeta}{n} = - \frac{\lfloor n(-x) - y \rfloor + \zeta}{n},$$

thus

$$\max_{n \in [n_{\min} : n_{\max}]} \frac{\lceil nx + y \rceil - \zeta}{n} = - \min_{n \in [n_{\min} : n_{\max}]} \frac{\lfloor n(-x) - y \rfloor + \zeta}{n},$$

and similarly

$$\min_{n \in [n_{\min} : n_{\max}]} \frac{\lceil nx + y \rceil - \zeta}{n} = - \max_{n \in [n_{\min} : n_{\max}]} \frac{\lfloor n(-x) - y \rfloor + \zeta}{n}.$$

11 The main algorithm

In this section, we finally develop our main algorithm, which finds the complete set of $(\xi, \zeta) \in \mathbb{R}^2$ such that

$$\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$$

holds for all $n \in [n_{\min} : n_{\max}]$.

Recall that having $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$ is equivalent to having

$$\begin{aligned} \max_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{>0}} \frac{\lfloor nx + y \rfloor - \zeta}{n} \leq \xi < \min_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{>0}} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n}, \\ \max_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{<0}} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} < \xi \leq \min_{n \in [n_{\min} : n_{\max}] \cap \mathbb{Z}_{<0}} \frac{\lfloor nx + y \rfloor - \zeta}{n}, \end{aligned}$$

possibly along with $\lfloor y \rfloor = \lfloor \zeta \rfloor$ if $0 \in [n_{\min} : n_{\max}]$. In Section 10 we described how to find extremizers appearing in the above inequalities. However, as noted earlier, that was provided that ζ is a given fixed number just like x, y . Yet, the way the actual extremizers depend on ζ is very special, and it allows us to obtain a fairly simple description of how the bounds appearing in the above inequalities change as ζ changes.

First, note that when ζ is fixed, the admissible interval for ξ can be described as the intersection of the intervals we obtain by separately solving some optimization problems to which one of the four theorems explained in Section 10 can be directly applied. Let us call such an optimization problem, which is one of the forms

$$\begin{aligned} \pm \max_{n \in [n_0 : m]} \frac{\lfloor n(\pm x) + y \rfloor - \zeta}{n}, \quad \pm \max_{n \in [m : n_0]} \frac{\lfloor n(\pm x) + y \rfloor - \zeta}{n}, \\ \pm \min_{n \in [n_0 : m]} \frac{\lfloor n(\pm x) + y \rfloor + 1 - \zeta}{n}, \quad \pm \min_{n \in [m : n_0]} \frac{\lfloor n(\pm x) + y \rfloor + 1 - \zeta}{n}, \end{aligned}$$

where n_0 is a maximizer/minimizer of the fractional part of $n(\pm x) + y$, as an *elementary optimization problem*. (The case $n = 0$ is special; it only imposes the constraint $\lfloor y \rfloor = \lfloor \zeta \rfloor$, and it does not impose any condition on ξ . For simplicity of the discussion, we will ignore it for the moment.)

Now, an important observation is that, according to the four theorems we proved in Section 10, the way ζ affects the optimizers to these elementary optimization problems is *only through* determining the iteration number k ; that is, specifying x, n_0 and m in the above already completely determines the whole sequence $(n_0^{(i)})_i$ we see in the aforementioned theorems, and the only thing ζ does is to merely “select” one of the entries of this sequence as the optimizer. More precisely, for each i , there exists an interval J_i of ζ such that the iteration number k equals i if and only if $\zeta \in J_i$. Clearly, these J_i ’s form a partition of the real line \mathbb{R} . In fact, the boundary between J_i and J_{i+1} is solely determined from the inequality at the i th step involving $n_0^{(i)}$ and $n_1^{(i+1)}$. More precisely, the threshold value determined from that inequality is monotone in i .

For instance, let us look at the following inequality from Theorem 10.1 (maximization problem on the left):

$$\frac{\lfloor n_1^{(i+1)} x \rfloor + 1}{n_1^{(i+1)}} > \frac{\lfloor n_0^{(i)} x + y \rfloor - \zeta}{n_0^{(i)}}.$$

The threshold value $\zeta^{(i)}$ is

$$\zeta^{(i)} := \frac{n_1^{(i+1)} \lfloor n_0^{(i)} x + y \rfloor - n_0^{(i)} (\lfloor n_1^{(i+1)} x \rfloor + 1)}{n_1^{(i+1)}}$$

so that the inequality above holds if and only if $\zeta > \zeta^{(i)}$. Then the claim is that $\zeta^{(i-1)} > \zeta^{(i)}$ holds for each $i = 1, \dots, i_0 - 1$, where i_0 is the first index such that $n_0^{(i_0)} = n_{\min}$ holds. To prove the claim, we simply compute $\zeta^{(i-1)} - \zeta^{(i)}$. In the proof of Theorem 10.1, we saw in (10) that

$$\lfloor n_0^{(i-1)} x + y \rfloor - \lfloor n_0^{(i)} x + y \rfloor = \lfloor n_1^{(i)} x \rfloor + 1$$

holds for all i , regardless of the actual number of iterations determined by ζ . Hence, we get

$$\begin{aligned} \zeta^{(i-1)} - \zeta^{(i)} &= \frac{n_1^{(i)} \lfloor n_0^{(i-1)} x + y \rfloor - n_0^{(i-1)} (\lfloor n_1^{(i)} x \rfloor + 1)}{n_1^{(i)}} \\ &\quad - \frac{n_1^{(i+1)} \lfloor n_0^{(i)} x + y \rfloor - n_0^{(i)} (\lfloor n_1^{(i+1)} x \rfloor + 1)}{n_1^{(i+1)}} \\ &= \left(\lfloor n_0^{(i-1)} x + y \rfloor - \lfloor n_0^{(i)} x + y \rfloor \right) + n_0^{(i)} \left(\frac{\lfloor n_1^{(i+1)} x \rfloor + 1}{n_1^{(i+1)}} - \frac{\lfloor n_1^{(i)} x \rfloor + 1}{n_1^{(i)}} \right) \\ &\quad - \left(n_0^{(i-1)} - n_0^{(i)} \right) \frac{\lfloor n_1^{(i)} x \rfloor + 1}{n_1^{(i)}} \\ &= n_0^{(i)} \left(\frac{\lfloor n_1^{(i+1)} x \rfloor + 1}{n_1^{(i+1)}} - \frac{\lfloor n_1^{(i)} x \rfloor + 1}{n_1^{(i)}} \right). \end{aligned}$$

Now, since $n_1^{(i)}$ is a minimizer of $\frac{\lfloor nx \rfloor + 1}{n}$ for $n \in [1: n_0^{(i-1)} - n_{\min}]$ and $n_1^{(i+1)}$ is in this range, the right-hand side of the above must be nonnegative. In fact, it must be strictly positive since $n_1^{(i)}$ is the *largest* minimizer. Indeed, from Theorem 7.1 we know that the minimizers of $\frac{\lfloor nx \rfloor + 1}{n}$ are precisely all the multiples of some $m \in [1: n_0^{(i-1)} - n_{\min}]$ contained in $[1: n_0^{(i-1)} - n_{\min}]$. However, since

$$n_1^{(i)} + n_1^{(i+1)} \leq n_1^{(i)} + n_0^{(i)} - n_{\min} = n_1^{(i-1)} - n_{\min}$$

still holds, if both $n_1^{(i)}$ and $n_1^{(i+1)}$ were minimizers of $\frac{\lfloor nx \rfloor + 1}{n}$, then $n_1^{(i)} + n_1^{(i+1)}$ must be a minimizer as well, contradicting to the maximality of $n_1^{(i)}$. Therefore, the claim is proved.

We can apply the same analysis onto other theorems as well. For Theorem 10.2 (maximization problem on the right), we have the inequality

$$\frac{\lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}} < \frac{\lfloor n_0^{(i)} x + y \rfloor - \zeta}{n_0^{(i)}}$$

if and only if $\zeta < \zeta^{(i)}$, where

$$\zeta^{(i)} := \frac{n_1^{(i+1)} \lfloor n_0^{(i)} x + y \rfloor - n_0^{(i)} \lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}}.$$

In this case, we must have $\zeta^{(i-1)} < \zeta^{(i)}$ for each $i = 1, \dots, i_0 - 1$ where i_0 is the first index such that $n_0^{(i_0)} = n_{\max}$ holds. To show that, recall (11) from the proof of Theorem 10.2 that

$$\lfloor n_0^{(i)} x + y \rfloor = \lfloor n_0^{(i-1)} x + y \rfloor + \lfloor n_1^{(i)} x \rfloor$$

holds for all i , which shows

$$\begin{aligned} \zeta^{(i)} - \zeta^{(i-1)} &= \lfloor n_1^{(i)} x \rfloor + n_0^{(i)} \left(\frac{\lfloor n_1^{(i)} x \rfloor}{n_1^{(i)}} - \frac{\lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}} \right) - (n_0^{(i)} - n_0^{(i-1)}) \frac{\lfloor n_1^{(i)} x \rfloor}{n_1^{(i)}} \\ &= n_0^{(i)} \left(\frac{\lfloor n_1^{(i)} x \rfloor}{n_1^{(i)}} - \frac{\lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}} \right). \end{aligned}$$

In the same way as before, the characterization of maximizers of $\frac{\lfloor nx \rfloor}{n}$ given in Theorem 7.1 shows that the above must be strictly positive.

For Theorem 10.3 (minimization on the left) with $\zeta \leftarrow \zeta - 1$, we have the inequality

$$\frac{\lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}} < \frac{\lfloor n_0^{(i)} x + y \rfloor + 1 - \zeta}{n_0^{(i)}}$$

if and only if $\zeta < \zeta^{(i)}$, where

$$\zeta^{(i)} := \frac{n_1^{(i+1)} \left(\lfloor n_0^{(i)} x + y \rfloor + 1 \right) - n_0^{(i)} \lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}}.$$

In this case, we must have $\zeta^{(i-1)} < \zeta^{(i)}$ for each $i = 1, \dots, i_0 - 1$ where i_0 is the first index such that $n_0^{(i_0)} = n_{\min}$ holds. Again, recall that (14) from the proof of Theorem 10.3 tells us

$$\lfloor n_0^{(i)} x + y \rfloor = \lfloor n_0^{(i-1)} x + y \rfloor - \lfloor n_1^{(i)} x \rfloor$$

holds for all i , which shows

$$\begin{aligned} \zeta^{(i)} - \zeta^{(i-1)} &= -\lfloor n_1^{(i)} x \rfloor + n_0^{(i)} \left(\frac{\lfloor n_1^{(i)} x \rfloor}{n_1^{(i)}} - \frac{\lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}} \right) + (n_0^{(i-1)} - n_0^{(i)}) \frac{\lfloor n_1^{(i)} x \rfloor}{n_1^{(i)}} \\ &= n_0^{(i)} \left(\frac{\lfloor n_1^{(i)} x \rfloor}{n_1^{(i)}} - \frac{\lfloor n_1^{(i+1)} x \rfloor}{n_1^{(i+1)}} \right), \end{aligned}$$

and again since $n_1^{(i)}$ is the *largest* maximizer, Theorem 7.1 shows that the above must be strictly positive.

Finally, for Theorem 10.4 (minimization on the right) with $\zeta \leftarrow \zeta - 1$, we have the inequality

$$\frac{\lfloor n_1^{(i+1)}x \rfloor + 1}{n_1^{(i+1)}} > \frac{\lfloor n_0^{(i)}x + y \rfloor + 1 - \zeta}{n_0^{(i)}}$$

if and only if $\zeta > \zeta^{(i)}$, where

$$\zeta^{(i)} := \frac{n_1^{(i+1)} \left(\lfloor n_0^{(i)}x + y \rfloor + 1 \right) - n_0^{(i)} \left(\lfloor n_1^{(i+1)}x \rfloor + 1 \right)}{n_1^{(i+1)}}.$$

In this case, we must have $\zeta^{(i-1)} > \zeta^{(i)}$ for each $i = 1, \dots, i_0 - 1$ where i_0 is the first index such that $n_0^{(i_0)} = n_{\max}$ holds. Recall that (15) from the proof of Theorem 10.4 tells us

$$\lfloor n_0^{(i)}x + y \rfloor = \lfloor n_0^{(i-1)}x + y \rfloor + \lfloor n_1^{(i)}x \rfloor + 1$$

holds for all i , which shows

$$\begin{aligned} \zeta^{(i-1)} - \zeta^{(i)} &= - \left(\lfloor n_1^{(i)}x \rfloor + 1 \right) + n_0^{(i)} \left(\frac{\lfloor n_1^{(i+1)}x \rfloor + 1}{n_1^{(i+1)}} - \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_1^{(i)}} \right) \\ &\quad + \left(n_0^{(i)} - n_0^{(i-1)} \right) \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_1^{(i)}} \\ &= n_0^{(i)} \left(\frac{\lfloor n_1^{(i+1)}x \rfloor + 1}{n_1^{(i+1)}} - \frac{\lfloor n_1^{(i)}x \rfloor + 1}{n_1^{(i)}} \right), \end{aligned}$$

and again since $n_1^{(i)}$ is the *largest* minimizer, Theorem 7.1 shows that the above must be strictly positive.

Therefore, for given elementary optimization problem, we can completely characterize the set of (ξ, ζ) satisfying all the constraints given by the problem. More explicitly, we have the following conclusions:

1. For a maximization problem on the left with positive n 's,

$$\begin{aligned} &\left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \geq \max_{n \in [m, n_0]} \frac{\lfloor nx + y \rfloor - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \geq \frac{\lfloor n_0^{(i)}x + y \rfloor - \zeta}{n_0^{(i)}} \text{ and } \zeta \in \left(\zeta^{(i)}, \zeta^{(i-1)} \right] \right\} \end{aligned}$$

where $\zeta^{(-1)} := +\infty$ and $\zeta^{(i_0)} := -\infty$.

2. For a maximization problem on the left with negative n 's,

$$\begin{aligned} &\left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \leq - \max_{n \in [m, n_0]} \frac{\lfloor n(-x) + y \rfloor - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \leq \frac{\lfloor -n_0^{(i)}x + y \rfloor - \zeta}{-n_0^{(i)}} \text{ and } \zeta \in \left(\zeta^{(i)}, \zeta^{(i-1)} \right] \right\} \end{aligned}$$

where $\zeta^{(-1)} := +\infty$ and $\zeta^{(i_0)} := -\infty$.

3. For a maximization problem on the right with positive n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \geq \max_{n \in [n_0, m]} \frac{\lfloor nx + y \rfloor - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \geq \frac{\lfloor n_0^{(i)} x + y \rfloor - \zeta}{n_0^{(i)}} \text{ and } \zeta \in [\zeta^{(i-1)}, \zeta^{(i)}] \right\} \end{aligned}$$

where $\zeta^{(-1)} := -\infty$ and $\zeta^{(i_0)} := +\infty$.

4. For a maximization problem on the right with negative n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \leq -\max_{n \in [n_0, m]} \frac{\lfloor n(-x) + y \rfloor - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \leq \frac{\lfloor -n_0^{(i)} x + y \rfloor - \zeta}{-n_0^{(i)}} \text{ and } \zeta \in [\zeta^{(i-1)}, \zeta^{(i)}] \right\} \end{aligned}$$

where $\zeta^{(-1)} := -\infty$ and $\zeta^{(i_0)} := +\infty$.

5. For a minimization problem on the left with positive n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi < \min_{n \in [m, n_0]} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi < \frac{\lfloor n_0^{(i)} x + y \rfloor + 1 - \zeta}{n_0^{(i)}} \text{ and } \zeta \in [\zeta^{(i-1)}, \zeta^{(i)}] \right\} \end{aligned}$$

where $\zeta^{(-1)} := -\infty$ and $\zeta^{(i_0)} := +\infty$.

6. For a minimization problem on the left with negative n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi > -\min_{n \in [m, n_0]} \frac{\lfloor n(-x) + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi > \frac{\lfloor -n_0^{(i)} x + y \rfloor + 1 - \zeta}{-n_0^{(i)}} \text{ and } \zeta \in [\zeta^{(i-1)}, \zeta^{(i)}] \right\} \end{aligned}$$

where $\zeta^{(-1)} := -\infty$ and $\zeta^{(i_0)} := +\infty$.

7. For a minimization problem on the right with positive n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi < \min_{n \in [n_0, m]} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi < \frac{\lfloor n_0^{(i)} x + y \rfloor + 1 - \zeta}{n_0^{(i)}} \text{ and } \zeta \in (\zeta^{(i)}, \zeta^{(i-1)}] \right\} \end{aligned}$$

where $\zeta^{(-1)} := +\infty$ and $\zeta^{(i_0)} := -\infty$.

8. For a minimization problem on the right with negative n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi > - \min_{n \in [n_0, m]} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcup_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi > \frac{\lfloor -n_0^{(i)} x + y \rfloor + 1 - \zeta}{-n_0^{(i)}} \text{ and } \zeta \in \left(\zeta^{(i)}, \zeta^{(i-1)} \right] \right\} \end{aligned}$$

where $\zeta^{(-1)} := +\infty$ and $\zeta^{(i_0)} := -\infty$.

We claim that the sets on the right-hand sides are equal to the one we obtain by removing the condition on ζ and replacing the union by the intersection:

1. For a maximization problem on the left with positive n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \geq \max_{n \in [m, n_0]} \frac{\lfloor nx + y \rfloor - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : n_0^{(i)} \xi + \zeta \geq \lfloor n_0^{(i)} x + y \rfloor \right\}. \end{aligned}$$

2. For a maximization problem on the left with negative n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \leq - \max_{n \in [m, n_0]} \frac{\lfloor n(-x) + y \rfloor - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : -n_0^{(i)} \xi + \zeta \geq \lfloor -n_0^{(i)} x + y \rfloor \right\}. \end{aligned}$$

3. For a maximization problem on the right with positive n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \geq \max_{n \in [n_0, m]} \frac{\lfloor nx + y \rfloor - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : n_0^{(i)} \xi + \zeta \geq \lfloor n_0^{(i)} x + y \rfloor \right\}. \end{aligned}$$

4. For a maximization problem on the right with negative n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi \leq - \max_{n \in [n_0, m]} \frac{\lfloor n(-x) + y \rfloor - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : -n_0^{(i)} \xi + \zeta \geq \lfloor -n_0^{(i)} x + y \rfloor \right\}. \end{aligned}$$

5. For a minimization problem on the left with positive n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi < \min_{n \in [m, n_0]} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : n_0^{(i)} \xi + \zeta < \lfloor n_0^{(i)} x + y \rfloor + 1 \right\}. \end{aligned}$$

6. For a minimization problem on the left with negative n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi > - \min_{n \in [m, n_0]} \frac{\lfloor n(-x) + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : -n_0^{(i)} \xi + \zeta < \lfloor -n_0^{(i)} x + y \rfloor + 1 \right\}. \end{aligned}$$

7. For a minimization problem on the right with positive n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi < \min_{n \in [n_0, m]} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : n_0^{(i)} \xi + \zeta < \lfloor n_0^{(i)} x + y \rfloor + 1 \right\}. \end{aligned}$$

8. For a minimization problem on the right with negative n 's,

$$\begin{aligned} & \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \xi > - \min_{n \in [n_0, m]} \frac{\lfloor nx + y \rfloor + 1 - \zeta}{n} \right\} \\ &= \bigcap_{i=0}^{i_0} \left\{ (\xi, \zeta) \in \mathbb{R}^2 : -n_0^{(i)} \xi + \zeta < \lfloor -n_0^{(i)} x + y \rfloor + 1 \right\}. \end{aligned}$$

For simplicity, let us denote the one with the union as U and one with the intersection as I . Also, let U_i be the set inside the union so that $U = \bigcup_{i=0}^{i_0} U_i$. Since I clearly contains the left-hand sides of the above list, $U \subseteq I$ follows. On the other hand, for any $(\xi, \zeta) \in I$, ζ must be in one of the intervals between $\zeta^{(i-1)}$ and $\zeta^{(i)}$, so clearly we must have $(\xi, \zeta) \in U_i$. Therefore, $I \subseteq U$ also follows.

All we have done is to find subsets L and U of $[n_{\min} : n_{\max}]$ such that $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ for all $n \in [n_{\min} : n_{\max}]$ if and only if

$$n\xi + \zeta \geq \lfloor nx + y \rfloor \text{ for all } n \in L \quad \text{and} \quad n\xi + \zeta < \lfloor nx + y \rfloor + 1 \text{ for all } n \in U.$$

The point here is that the number of elements in L and U is in general much smaller than that of $[n_{\min} : n_{\max}]$. By construction, it appears to me that $|L \cup U|$ is probably of $O(\log N)$ where N is the size of the interval $[n_{\min} : n_{\max}]$, but I have not tried to show this rigorously. Also, note that we can now include the case $n = 0$ here by adding 0 to both L and U .

We can even obtain the following generalization without any further complication: given a finite collection $\{(A_i, b_i, [n_{\min, i} : n_{\max, i}])\}_{i=1}^r$ of triples of a linear map $A_i : \mathbb{R}^2 \rightarrow \mathbb{R}^2$, a vector $b_i \in \mathbb{R}^2$ and an integer interval $[n_{\min, i} : n_{\max, i}]$, the set of (ξ, ζ) such that

$$\left| \left\langle \begin{pmatrix} n \\ 1 \end{pmatrix}, A_i \begin{pmatrix} x \\ y \end{pmatrix} + b_i \right\rangle \right| = \left| \left\langle \begin{pmatrix} n \\ 1 \end{pmatrix}, A_i \begin{pmatrix} \xi \\ \zeta \end{pmatrix} + b_i \right\rangle \right|$$

holds for all $n \in [n_{\min, i} : n_{\max, i}]$ for all $i = 1, \dots, r$, can be described as the intersection of finitely many closed half-planes of the form

$$\left\{ (\xi, \zeta) \in \mathbb{R}^2 : \left\langle A_i^T \begin{pmatrix} n \\ 1 \end{pmatrix}, \begin{pmatrix} \xi \\ \zeta \end{pmatrix} \right\rangle \geq \left\lfloor \left\langle \begin{pmatrix} n \\ 1 \end{pmatrix}, A_i \begin{pmatrix} x \\ y \end{pmatrix} + b_i \right\rangle \right\rfloor - \left\langle \begin{pmatrix} n \\ 1 \end{pmatrix}, b_i \right\rangle \right\}$$

or open half-planes of the form

$$\left\{ (\xi, \zeta) \in \mathbb{R}^2 : \left\langle A_i^T \begin{pmatrix} n \\ 1 \end{pmatrix}, \begin{pmatrix} \xi \\ \zeta \end{pmatrix} \right\rangle < \left\lfloor \left\langle \begin{pmatrix} n \\ 1 \end{pmatrix}, A_i \begin{pmatrix} x \\ y \end{pmatrix} + b_i \right\rangle \right\rfloor + 1 - \left\langle \begin{pmatrix} n \\ 1 \end{pmatrix}, b_i \right\rangle \right\}.$$

(Here, $\langle u, v \rangle$ means the inner product of the vectors u and v .) In this case, the $n_0^{(i)}$'s we considered are obtained by applying the previous arguments to $A_j \begin{pmatrix} x \\ y \end{pmatrix} + b_j$ instead of $\begin{pmatrix} x \\ y \end{pmatrix}$.

It is fair to say that at this point we have obtained a complete description of the set of (ξ, ζ) we are seeking for: it is the intersection of some fairly small finite collection of closed or open half-planes which we know the complete description of. Nevertheless, let us try to obtain a simpler description of the intersection, because in general many of the half-planes would be redundant and does not contribute to the intersection. Finding the intersection of closed half-planes is a well-studied problem in computational geometry, and possibility of having open half-planes is just a simple (though quite annoying) extension of it. Among others, one typical strategy is to use duality to turn the problem into the problem of finding the convex hull.

11.1 General outline for computing the intersection

We are now given with a collection $\{(\xi_i^*, \zeta_i^*, \eta_i^*)\}_{i=1}^k \subseteq \mathbb{R}^3$, and our next task is to find the intersection of the sets

$$H_i := \{(\xi, \zeta) \in \mathbb{R}^2 : \xi_i^* \xi + \zeta_i^* \zeta + \eta_i^* \geq 0\}.$$

The case of possibly having open half-planes is explained in the next subsection. Before explaining our approach in detail, let us first note what kinds of shapes are possible the intersection can be of and clarify what exactly we want to obtain as a description of the intersection.

- If $k = 0$, of course the intersection is the entire plane \mathbb{R}^2 .
- If the boundary of every half-plane is parallel to each other, then the intersection is either empty, an infinite parallelogram, an infinite straight line, or a half-plane. In our case, the intersection cannot be empty because we know at least (x, y) belongs to it. It also cannot be a half-plane, because the set we are ultimately looking for is the intersection of infinite parallelograms of the form

$$\{(\xi, \zeta) \in \mathbb{R}^2 : n\xi + \zeta \geq \lfloor nx + y \rfloor\} \cap \{(\xi, \zeta) \in \mathbb{R}^2 : n\xi + \zeta < \lfloor nx + y \rfloor + 1\}$$

(with an appropriate modification if we consider the generalization taking the affine transform $\begin{pmatrix} x \\ y \end{pmatrix} \mapsto A \begin{pmatrix} x \\ y \end{pmatrix} + b$ into account; note that such a generalization may turn an infinite parallelogram into the entire \mathbb{R}^2 , but it can never turn it into a half-plane), although we have separately compactified the list of n 's we need to consider for the lower bounds and the upper bounds so that for each i there does not need to exist a j such that $(\xi_i^*, \zeta_i^*) = -(\xi_j^*, \zeta_j^*)$.

- Otherwise, the intersection must be a (possibly degenerate) bounded convex polygon, because the intersection of two non-parallel parallelograms is already bounded, and the intersection of convex sets is convex. The polygon may have nonempty interior, or may degenerate into either a point or a bounded line segment.

Remark 11.1.

The degenerate cases, that the intersection becomes a region of dimension strictly less than 2, can only happen for the generalization involving affine transforms, and it is not possible to happen in the original problem, because as noted in Section 9, the set of (ξ, ζ) *simultaneously* satisfying $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ and $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ already has a nonempty interior.

For the trivial case of the entire plane \mathbb{R}^2 , of course there is no need of any further description. For the case of single point, of course specifying the point itself is what we would do, and for the case of line segment, we can for example specify a base point b and a vector v so that $\{b + tv \in \mathbb{R}^2 : t \in [0, 1]\}$ is the intersection.

For the case of infinite parallelogram, the intersection can be described by specifying a normal vector to the boundary and an interval of values that the inner product between the normal vector and (ξ, ζ) can take. The infinite straight line case then becomes a special case when the interval is a single point.

For the case of non-degenerate bounded polygon, there can be many representations (specifying the list of supporting hyperplanes being one of) but we will obtain the list of “vertical slices” of the polygon which are vertically parallel trapezoids, where the slice happens exactly at the vertices of the polygon. This feels to me like a generally useful representation for further processings, for instance, to find lattice points.

Having defined our goal more precisely, now let us see how to achieve it. Our strategy is to divide $(\xi_i^*, \zeta_i^*, \eta_i^*)$ ’s into those with positive ζ_i^* , negative ζ_i^* , and zero ζ_i^* , compute the intersection of them separately, and then compute the intersection of the resulting intersections. We assume that none of $(\xi_i^*, \zeta_i^*) \in \mathbb{R}^2$ is zero, because in our original problem such a case can happen only if the corresponding H_i is the whole plane \mathbb{R}^2 (in other words, it is never empty).

First, let us look at the case of positive ζ_i^* ’s. For notational simplicity, suppose that all ζ_i^* ’s are positive. In this case, the intersection

$$\bigcap_{i=1}^k H_i = \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \zeta \geq -\frac{\xi_i^*}{\zeta_i^*} \xi - \frac{\eta_i^*}{\zeta_i^*} \right\}$$

is nothing but the *epigraph* of the continuous convex function

$$f : \mathbb{R} \rightarrow \mathbb{R}$$

$$\xi \mapsto \max_{i \in [1:k]} \left(-\frac{\xi_i^*}{\zeta_i^*} \xi - \frac{\eta_i^*}{\zeta_i^*} \right).$$

By separating hyperplane theorem, one can show that f is equal to its double *convex conjugate*, that is,

$$f(\xi) = \sup_{\xi^* \in \mathbb{R}} (\xi^* \xi - f^*(\xi^*))$$

holds for all $\xi \in \mathbb{R}$ where we define

$$f^* : \mathbb{R} \rightarrow (-\infty, +\infty]$$

$$\xi^* \mapsto \sup_{\xi \in \mathbb{R}} (\xi^* \xi - f(\xi)).$$

More generally, given a function $f : \mathbb{R} \rightarrow (-\infty, +\infty]$, *Fenchel-Moreau theorem* says that f^{**} is equal to the *convex envelope* of f , which is the supremum of every convex lower semicontinuous function bounded above by f . This is a consequence of the separating hyperplane theorem and is a standard result in the literature, which can be found for example in [10].

In general, if $\{f_i : \mathbb{R} \rightarrow [-\infty, +\infty]\}_{i \in I}$ is any collection of convex lower semicontinuous functions, then Fenchel-Moreau theorem shows that

$$\left(\inf_{i \in I} f_i^* \right)^* (\xi) = \sup_{\xi^*, i} (\xi^* \xi - f_i^*(\xi^*)) = \sup_{i \in I} f_i^{**}(\xi) = \sup_{i \in I} f_i(\xi),$$

thus

$$\left(\sup_{i \in I} f_i \right)^* = \left(\inf_{i \in I} f_i^* \right)^{**}$$

is the convex envelope of $\inf_{i \in I} f_i^*$.

In our case, let $f_i: \xi \mapsto -\frac{\xi_i^*}{\zeta_i^*} \xi - \frac{\eta_i^*}{\zeta_i^*}$, then

$$f_i^*(\xi^*) = \sup_{\xi \in \mathbb{R}} \left(\left(\xi^* + \frac{\xi_i^*}{\zeta_i^*} \right) \xi + \frac{\eta_i^*}{\zeta_i^*} \right) = \begin{cases} \frac{\eta_i^*}{\zeta_i^*} & \text{if } \xi = -\frac{\xi_i^*}{\zeta_i^*}, \\ \infty & \text{otherwise} \end{cases},$$

thus f^* is the convex envelope of the function

$$\xi^* \mapsto \min_{i \in [1: k]} f_i^*(\xi^*) = \begin{cases} \min_{j: \frac{\xi_i^*}{\zeta_i^*} = \frac{\xi_j^*}{\zeta_j^*}} \frac{\eta_j^*}{\zeta_j^*} & \text{if } \xi = -\frac{\xi_i^*}{\zeta_i^*} \text{ for some } i \in [1: k], \\ \infty & \text{otherwise.} \end{cases}$$

The convex envelope g^{**} of a function $g: \mathbb{R} \rightarrow (-\infty, \infty]$ is the unique function having the closed convex hull of the epigraph of g as its epigraph, so any convex hull algorithm can be utilized for finding the convex envelope. In the reference implementation, I used a variant of *Andrew's monotone chain algorithm*. Roughly speaking, we first sort the points $\left\{ \left(-\frac{\xi_i^*}{\zeta_i^*}, \frac{\eta_i^*}{\zeta_i^*} \right) \right\}_{i=1}^k$ in ascending order according to the first coordinate, and while looping over all points from left to right, add each new point to the stack after removing enough previously added points to make a counterclockwise turn. The original algorithm by A. M. Andrew is two-phased because it has to march to right from left and then come back to left from right, but in our case the second phase is not needed. We omit details here.

After running the algorithm, we find a subset of $\left\{ \left(-\frac{\xi_i^*}{\zeta_i^*}, \frac{\eta_i^*}{\zeta_i^*} \right) \right\}_{i=1}^k$ so that f^* is a piecewise affine function interpolating between those selected points. To ease notational overhead, let us again write this subset as just $\left\{ \left(-\frac{\xi_i^*}{\zeta_i^*}, \frac{\eta_i^*}{\zeta_i^*} \right) \right\}_{i=1}^k$, and furthermore assume $\left(-\frac{\xi_i^*}{\zeta_i^*} \right)_{i=1}^k$ is strictly increasing. Then $f^*(\xi^*)$ turns out to be equal to

$$h_i(\xi^*) := \frac{\eta_i^* \zeta_{i+1}^*}{\xi_{i+1}^* \zeta_i^* - \xi_i^* \zeta_{i+1}^*} \left(\xi^* + \frac{\xi_{i+1}^*}{\zeta_{i+1}^*} \right) - \frac{\eta_{i+1}^* \zeta_i^*}{\xi_{i+1}^* \zeta_i^* - \xi_i^* \zeta_{i+1}^*} \left(\xi^* + \frac{\xi_i^*}{\zeta_i^*} \right)$$

whenever $\xi^* \in \left[-\frac{\xi_{i+1}^*}{\zeta_{i+1}^*}, -\frac{\xi_i^*}{\zeta_i^*} \right]$ for some $i = 1, \dots, k-1$, and otherwise $f^*(\xi^*) = \infty$. Or, we can write f^* as the minimum of functions $g_i: \mathbb{R} \rightarrow (-\infty, +\infty]$ defined as

$$g_i: \xi^* \mapsto \begin{cases} h_i(\xi^*) & \text{if } \xi^* \in \left[-\frac{\xi_{i+1}^*}{\zeta_{i+1}^*}, -\frac{\xi_i^*}{\zeta_i^*} \right], \\ \infty & \text{otherwise,} \end{cases}$$

thus $f = (f^*)^* = \max_{i \in [1: k]} g_i^*$. Note that $g_i \left(-\frac{\xi_i^*}{\zeta_i^*} \right) = -\frac{\eta_i^*}{\zeta_i^*}$ and $g_i \left(-\frac{\xi_{i+1}^*}{\zeta_{i+1}^*} \right) = -\frac{\eta_{i+1}^*}{\zeta_{i+1}^*}$, so for each $i = 1, \dots, k$, let

$$\xi_i := \frac{\eta_i^* \zeta_{i+1}^* - \eta_{i+1}^* \zeta_i^*}{\xi_{i+1}^* \zeta_i^* - \xi_i^* \zeta_{i+1}^*},$$

then it follows that

$$\begin{aligned} g_i^* : \xi &\mapsto \sup_{\xi^* \in \left[-\frac{\xi_{i+1}^*}{\zeta_{i+1}^*}, -\frac{\xi_i^*}{\zeta_i^*}\right]} \left(\xi^* (\xi - \xi_i) + \frac{\eta_{i+1}^* \xi_i^* - \eta_i^* \xi_{i+1}^*}{\xi_{i+1}^* \zeta_i^* - \xi_i^* \zeta_{i+1}^*} \right) \\ &= \begin{cases} f_i(\xi) & \text{if } \xi \geq \xi_i, \\ f_{i+1}(\xi) & \text{if } \xi \leq \xi_i. \end{cases} \end{aligned}$$

Since f^* is convex and $\left(-\frac{\xi_i^*}{\zeta_i^*}\right)_{i=1}^k$ is strictly increasing, the slope ξ_i of h_i should be increasing in i . By merging neighboring h_i 's with the same slope if necessary, we can even assume that $(\xi_i)_{i=1}^k$ is strictly increasing. Then it follows that $f(\xi) = f_i(\xi)$ if $\xi \in [\xi_{i-1}, \xi_i]$ for some $i = 1, \dots, k$, where $\xi_0 := -\infty$ and $\xi_n := +\infty$. This gives a direct piecewise affine description of f , which then immediately gives a similar description of the epigraph $\bigcap_{i=1}^k H_i$ of f . Let us call these ξ_i 's the *turning points* of f .

Computing the intersection for the case of negative ζ_i^* 's can be done similarly. In this case, the intersection

$$\bigcap_{i=1}^k H_i = \left\{ (\xi, \zeta) \in \mathbb{R}^2 : \zeta \leq -\frac{\xi_i^*}{\zeta_i^*} \xi - \frac{\eta_i^*}{\zeta_i^*} \right\}$$

is the *hypograph* of the concave function

$$\begin{aligned} f : \mathbb{R} &\rightarrow \mathbb{R} \\ \xi &\mapsto \min_{i \in [1:k]} \left(-\frac{\xi_i^*}{\zeta_i^*} \xi - \frac{\eta_i^*}{\zeta_i^*} \right), \end{aligned}$$

or equivalently, inversion of the epigraph of the convex function

$$\begin{aligned} -f : \mathbb{R} &\rightarrow \mathbb{R} \\ \xi &\mapsto \max_{i \in [1:k]} \left(\frac{\xi_i^*}{\zeta_i^*} \xi + \frac{\eta_i^*}{\zeta_i^*} \right). \end{aligned}$$

Then all the arguments we made for the positive case apply here too without any problem.

The case of zero ζ_i^* 's is the simplest. In this case, the intersection

$$\bigcap_{i=1}^k H_i = \{ (\xi, \zeta) \in \mathbb{R}^2 : \xi_i^* \xi \geq -\eta_i^* \}$$

can be found by finding the largest $\frac{-\eta_i^*}{\xi_i^*}$ among i 's such that $\xi_i^* > 0$, and finding the smallest $\frac{-\eta_i^*}{\xi_i^*}$ among i 's such that $\xi_i^* < 0$. (Recall that we already have excluded the case $\xi_i^* = 0$.)

Next, once we find the intersections for $\zeta_i^* > 0$, $\zeta_i^* < 0$, and $\zeta_i^* = 0$ separately, each denoted as P , N , and Z , we have to find the intersection $P \cap N \cap Z$. Note that computing the intersection with Z is simply a matter of cutting the region vertically, so we will focus only on how to compute the intersection of P and N , and omit details about taking the intersection with Z at the final stage.

Recall that P is the epigraph of a continuous convex piecewise affine function, and N is the hypograph of a continuous concave piecewise affine function. Let f, g denote these functions, respectively, then the region we are looking for is

$$P \cap N = \{ (\xi, \zeta) \in \mathbb{R}^2 : f(\xi) \leq \zeta \leq g(\xi) \}.$$

We will refer to each of the affine piece of f and g defined on their respective intervals as *affine segments* of f and g . For instance, if f itself is an affine function, then its unique affine segment is f itself. On the other hand, if f is given as

$$f: \xi \mapsto \begin{cases} 0 & \text{if } \xi \leq 0, \\ \xi & \text{if } \xi \geq 0, \end{cases}$$

then its affine segments are functions $\xi \mapsto 0$ and $\xi \mapsto \xi$ defined on $(-\infty, 0]$ and $[0, \infty)$, respectively.

Consider the projection $\pi_\xi[P \cap N]$ of $P \cap N$ onto the ξ -axis. Note that this is precisely the set of $\xi \in \mathbb{R}$ such that $(f - g)(\xi) \leq 0$. Since $f - g$ is convex and continuous, this set must be closed and convex, that is, a closed interval. Since we know $(x, y) \in P \cap N$, $\pi_\xi[P \cap N]$ is never empty.

Suppose that $\pi_\xi[P \cap N] = \mathbb{R}$. This means $f - g \leq 0$ holds on all of \mathbb{R} , which implies that $f - g$ is a constant. Indeed, if not, then there exists a supporting hyperplane of $f - g$ with a nonzero slope, enforcing either $\lim_{\xi \rightarrow +\infty} (f - g)(\xi) = +\infty$ or $\lim_{\xi \rightarrow -\infty} (f - g)(\xi) = +\infty$. Hence, we have $g = f + c$ for some constant $c \geq 0$, so g should be convex since f is convex. However, since g is concave, the only possibility is that g is affine. Hence, $f = g - c$ is also affine. Therefore, we conclude that $\pi_\xi[P \cap N] = \mathbb{R}$ holds if and only if f and g are parallel affine functions.

Then $\pi_\xi[P \cap N]$ can be either of the forms $[\xi_{\min}, \xi_{\max}]$, $(-\infty, \xi_{\min}]$, or $[\xi_{\max}, \infty)$. To distinguish between these cases, it is enough to compare the slopes of the left-most and the right-most segments of f and g . More precisely, we claim that the slope of the left-most affine segment of f is strictly smaller than that of g if and only if $\inf \pi_\xi[P \cap N] > -\infty$. Indeed, if the slope of the segment of f is strictly smaller, then the left-most affine segment of $f - g$ has a strictly negative slope, so we have $\lim_{\xi \rightarrow -\infty} (f - g)(\xi) = +\infty$, which shows $\inf \pi_\xi[P \cap N] > -\infty$. On the other hand, if the slope of the segment of f is strictly larger, then similarly we conclude $\lim_{\xi \rightarrow -\infty} (f - g)(\xi) = -\infty$ so $\inf \pi_\xi[P \cap N] = -\infty$. If two slopes are the same, then $(f - g)(\xi) = c$ for some constant $c \in \mathbb{R}$ whenever ξ is small enough. Note that this c must be nonpositive; otherwise, choose any $\xi < x$ small enough so that $(f - g)(\xi) = c$, then the slope of the line segment joining $(\xi, (f - g)(\xi))$ and $(x, (f - g)(x))$ is strictly negative, contradicting to the convexity of $f - g$. Hence, we get $\inf \pi_\xi[P \cap N] = -\infty$ in this case as well. In the same way, one can see that the slope of the right-most affine segment of f is strictly larger than that of g if and only if $\sup \pi_\xi[P \cap N] < \infty$.

In summary,

1. If the slope of the left-most affine segment of f is strictly smaller than that of g ,
 - (a) If the slope of the right-most affine segment of f is strictly larger than that of g , then $\pi_\xi[P \cap N]$ is of the form $[\xi_{\min}, \xi_{\max}]$ for some $\xi_{\min}, \xi_{\max} \in \mathbb{R}$ with $\xi_{\min} \leq \xi_{\max}$.
 - (b) Otherwise, $\pi_\xi[P \cap N]$ is of the form $[\xi_{\min}, +\infty)$ for some $\xi_{\min} \in \mathbb{R}$.
2. Otherwise,
 - (a) If the slope of the right-most affine segment of f is strictly larger than that of g , then $\pi_\xi[P \cap N]$ is of the form $(-\infty, \xi_{\max}]$ for some $\xi_{\max} \in \mathbb{R}$.
 - (b) Otherwise, f and g are parallel affine functions.

We can detect and return early for the case 2(b). In this case, if the intersection Z is not the entire \mathbb{R}^2 , then we get a bounded parallelogram, and otherwise we get an infinite parallelogram.

For all other cases, the graph of f and the graph of g should meet at some point, so we want to find their first intersection. To do so, we start from the left-most segments of f and g , and see if they intersect. If not, then compare the right endpoint of their domains, and choose the one that is smaller, and move to the next segment. If the domains have the same right endpoint, then we proceed both f and g .

When the intersection of the graphs is found for the first time, either the segments of f and g we are looking at can be parallel (that is, they are the same functions except for their domains) or not. In fact, they can be parallel only if they are the left-most segments of f and g , because otherwise the intersection should happen at the left endpoint of the domain of either of the segments, but in that case the intersection should have been already found from the previous segment. Hence, let us suppose that f and g share the same left-most segment except for their domains. This can happen only when $\pi_\xi[P \cap N]$ is of the form $(-\infty, \xi_{\max}]$, and ξ_{\max} is precisely where one of the left-most segments of f and g ends, because after the segment ends the next segment must have a strictly larger slope if it were for f and a strictly smaller slope if it were for g . Hence, the convex function $f - g$ is strictly positive at some $\xi > \xi_{\max}$ if ξ close enough to ξ_{\max} , and by convexity $f - g$ can never hit zero again. Therefore, the shared part of the left-most segments of f and g is precisely the intersection $P \cap N$. In this case, Z must be nontrivial and the intersection $P \cap N \cap Z$ is either a line segment or a point, and for either case it is easy to get the desired description of the resulting set.

Next, suppose that the first intersection happens on non-parallel segments. In this case, what matters is how the graphs of f and g will continue immediately after the intersection. There are three possibilities: f goes strictly above g , f goes strictly below g , or f and g proceed together. These possibilities can be distinguished by comparing the slopes of the affine segments of f and g . Since we are interested in what happens *after* the intersection, we have to choose the next segment instead if the intersection happens at the right endpoint of the domain of the current segment.

If the slope for f is strictly larger, then f goes strictly above g . If we know $\inf \pi_\xi[P \cap N]$ is finite, then since this is the first intersection we found, it should coincide with $\inf \pi_\xi[P \cap N]$. By convexity, $f - g$ can never hit zero again, thus in this case $P \cap N$ should be a single point and we know exactly where the point is. On the other hand, if $\inf \pi_\xi[P \cap N]$ is infinite, then this firstly found intersection must be exactly $\xi_{\max} := \sup \pi_\xi[P \cap N] < \infty$ because again convexity ensures that $f - g$ can never hit zero again. Then by vertically slicing the domains of the segments of f and g at all turning points of f and g before ξ_{\max} , we obtain a description of $P \cap N$ in terms of a union of essentially disjoint vertically parallel trapezoids. Since the set we get is not bounded, in this case Z must be not trivial, and by taking the intersection with Z , we get either a single point (if the left boundary of Z passes through $P \cap N$ exactly at ξ_{\max}) or a non-degenerate bounded polygon with the desired description.

If the slope of f is either same as that of g or is strictly smaller, then f either proceeds with g together or goes strictly below g , at least for a while. Note that in this case we can never have $\inf \pi_\xi[P \cap N] = -\infty$, because if that were the case, then $(f - g)(\xi)$ is nonpositive for all sufficiently small ξ while $f - g$ becomes zero at some point and is nonpositive on a nontrivial interval immediately following that point, which by convexity is only possible if $(f - g)(\xi) = 0$ for all small enough ξ . However, that means f and g share their left-most segments, which cannot be the case as we are looking at the first intersection of their graphs. Therefore, $\xi_{\min} := \inf \pi_\xi[P \cap N]$ is finite, so the intersection we found must be exactly at ξ_{\min} .

Specializing further, if the slope of f is same as that of g , then they remain to be same immediately after the intersection at least for a while. Since the convexity guarantees that

f can never go below g after that, either f and g separates and never meet again, or they continue to be together forever. For both cases, after taking the intersection with Z , we get either a single point or a line segment, and it is easy to obtain their desired descriptions.

If the slope for f is strictly smaller, then f goes strictly below g . If $\sup \pi_\xi[P \cap N] = \infty$, then it must be of the form $[\xi_{\min}, \infty)$. Hence, by slicing the domains of the segments of f and g at all turning points of f and g after ξ_{\min} , we obtain a description of $P \cap N$ in terms of a union of essentially disjoint vertically parallel trapezoids. Since the set we get is not bounded, in this case Z must be not trivial, and by taking the intersection with Z , we get either a single point (if the right boundary of Z passes through $P \cap N$ exactly at ξ_{\min}) or a non-degenerate bounded polygon with the desired description.

Finally, suppose $\xi_{\max} := \sup \pi_\xi[P \cap N] < \infty$. In this case, we must find the second intersection which again can be done by successively moving onto the next segments of f and g while checking for the intersection. Meanwhile, we record the turning points of f and g and the corresponding vertically parallel trapezoidal slices. Once the second intersection is found, it must be ξ_{\max} , because $f - g$ is strictly negative right before the intersection and we assumed that its slope at ξ_{\min} is strictly negative. Hence, the slope of $f - g$ at the intersection should be strictly positive and $f - g$ never hits zero again. Hence, we found $[\xi_{\min}, \xi_{\max}]$, and $P \cap N$ is a non-degenerate bounded polygon. Again, taking the intersection with Z can result in a single point (if either the left boundary of Z is at ξ_{\max} or the right boundary of Z is at ξ_{\min}) or a bounded polygon.

11.2 Consideration of open half-planes

The only issue remaining is that some of the half-planes we consider are open rather than closed. What we need to do in this case is mostly just keeping track of whether an affine segment is coming from a closed half-plane or an open half-plane, and we will not go into all details about where this can matter.

Instead, I will point out one particular issue that might be considered nontrivial. Recall that when we obtained the piecewise affine representation of the function f by computing the convex envelope of a function, we assumed that the turning points $(\xi_i)_{i=1}^k$ are strictly increasing, and this was done by merging neighboring affine segments of f^* with the same slope. In terms of how the monotone chain algorithm works, such a merging happens if we detect that points we are looking at are colinear. If all half-planes are closed, we can simply drop intermediate points and only take the extreme points. However, if some half-planes are open, then we may lose some information from it. Indeed, three points in the dual domain being colinear means that the corresponding three affine functions in the original domain having a common intersection. If, in the dual domain the two outer points are closed but the inner one is open, then this common intersection point should not be included into the intersection even though both of the visible affine segments around it are from closed half-planes. Hence, when colinear points are detected, we have to keep track of whether some of them are from open half-planes.

12 Gosper's algorithm

I completed explaining all of the main algorithm of this note, except for the algorithm for computing the continued fraction expansion of numbers we obtain by combining other numbers with known continued fraction expansions. This section is for that missing piece, called *Gosper's algorithm*, invented by Bill Gosper around 1970's.

Apparently Gosper did not publish his work into any academic journal (probably he felt it is not that significant), instead he wrote some informal notes [11] about the algorithm. Unfortunately these notes are somewhat “hidden” from the Internet behind lots of so-called “tutorials” with *wrong* explanations of the algorithm. I uploaded the original notes written by Gosper I found on the Internet into the Github repository for my C++ implementation of our main algorithm. Interested readers can take a look at it. (There probably are more notes by Gosper on this topic, but I did not feel obliged to collect all of them.)

In fact, the algorithm explained in this section is not exactly what Gosper invented. Rather, it is a somewhat more complete, but also more complicated form of it. This itself is probably a novel work, though I did not put a lot of effort to try to find existing literatures. I felt some need of this “completed” version for a fully automated implementation of our main algorithm we saw in Section 11 when I first wrote it, but probably the simpler, original form by Gosper may just work fine. Since the algorithm explained in this section is fairly complicated, I recommend readers to just read Gosper’s original notes if they want a quick implementation. I also want to remind them that Gosper’s algorithm is not needed at all if all they care are rational inputs.

There are two versions of Gosper’s algorithm that we want to use: unary and binary. Binary version is strictly more general, but since it is much more complicated we begin with the unary version first.

12.1 Unary Gosper’s algorithm

Given a real number x , *unary Gosper’s algorithm* is an algorithm for computing the continued fraction expansion of $\frac{ax+b}{cx+d}$ using the continued fraction expansion of x , for any given $a, b, c, d \in \mathbb{Z}$. For the purpose of this section, it is much more convenient to work with the projective line \mathbb{RP}^1 rather than the real line \mathbb{R} , so we will assume x is an element in \mathbb{RP}^1 .¹⁰ For a moment we let a, b, c, d live in \mathbb{R} , although we eventually want them to be in \mathbb{Z} . Also, it is more convenient for our purpose to work with the matrix $A := \begin{pmatrix} a & b \\ c & d \end{pmatrix}$ rather than the individual components.

Note that the linear map $A: \mathbb{R}^2 \rightarrow \mathbb{R}^2$ induces a map $T_A: \text{dom}(T_A) \subseteq \mathbb{RP}^1 \rightarrow \mathbb{RP}^1$ given as

$$T_A: [u] \mapsto [Au],$$

where $[\cdot]: \mathbb{R}^2 \setminus \{0\} \rightarrow \mathbb{RP}^1$ is the natural map that identifies \mathbb{RP}^1 as a quotient space of $\mathbb{R}^2 \setminus \{0\}$, and $\text{dom}(T_A)$ is the set of $[u] \in \mathbb{RP}^1$ such that $Au \neq 0$; note that $\text{dom}(T_A)$ is a well-defined subset of \mathbb{RP}^1 . To be more precise,

1. If $\text{rank } A = 0$, i.e., $A = 0$, then $\text{dom}(T_A) = \emptyset$.
2. If $\text{rank } A = 1$, then $\ker A$ corresponds to a unique element in \mathbb{RP}^1 , so $\text{dom}(T_A)$ is \mathbb{RP}^1 minus the single point set $\{\ker A\}$. Note that in this domain T_A is a constant function, whose unique value is the image $\text{im } A$ of A .
3. If $\text{rank } A = 2$, i.e., A is invertible, then $\text{dom}(T_A) = \mathbb{RP}^1$.

¹⁰In case you are not familiar, \mathbb{RP}^1 is the set of all equivalence classes of $\mathbb{R}^2 \setminus \{0\}$ by the equivalence relation $x \sim y$ if and only if $y = kx$ for some $k \in \mathbb{R} \setminus \{0\}$. \mathbb{RP}^1 can also be regarded as the set of all 1-dimensional linear subspaces of \mathbb{R}^2 . We can regard \mathbb{R} as a subset of \mathbb{RP}^1 , by identifying \mathbb{R} with the affine subspace $\mathbb{R} \times \{1\} = \{(x, 1) \in \mathbb{R}^2: x \in \mathbb{R}\}$ of \mathbb{R}^2 , with the obvious mapping $x \mapsto (x, 1)$. Then each element $(x, 1)$ in $\mathbb{R} \times \{1\}$ gives a unique equivalence class in \mathbb{RP}^1 namely the linear subspace spanned by $(x, 1)$. On the other hand, \mathbb{RP}^1 contains one more element than \mathbb{R} , the subspace spanned by $(1, 0)$. One can regard \mathbb{RP}^1 as the real line with this “infinity” added.

For the case $\text{rank } A = 1$, since T_A is constant on all of its domain, it is natural to think of its extension onto the whole \mathbb{RP}^1 defined to take the unique constant value even at the singularity $\ker A \in \mathbb{RP}^1$. In fancier terms, it admits the unique *continuous extension* onto \mathbb{RP}^1 . Let us denote this unique extension as \bar{T}_A from now on. For all practical purposes, it is this extension \bar{T}_A that matters so we can safely forget about the original T_A , but for the sake of developing the theory, let us be clear about this distinction.

In general, for linear maps A, B on \mathbb{R}^2 , we have

$$\bar{T}_{AB} = \bar{T}_A \circ \bar{T}_B$$

provided that $AB \neq 0$. Indeed, it immediately follows from the definition of T that $T_{AB} = T_A \circ T_B$ holds on the intersection of their domains. If A is not invertible, then the right-hand side is just a constant function whose unique value is the image of A . At the same time, the left-hand side is not invertible, and since $AB \neq 0$, AB is of rank 1, so $\text{im}(AB) = \text{im } A$ should hold, which means that the left-hand side is also the constant function whose unique value is $\text{im } A$. If B is not invertible, then since $AB \neq 0$, $\text{im}(AB) = A[\text{im } B]$ is not zero, so it should be a 1-dimensional subspace which must be the unique value of both sides.

With the usual identification $\mathbb{RP}^1 = \mathbb{R} \sqcup \{\infty\}$ where $(u_1 : u_2) := [(u_1, u_2)] \in \mathbb{RP}^1$ belongs to \mathbb{R} if and only if $u_2 \neq 0$, in which case it corresponds to $\frac{u_1}{u_2} \in \mathbb{R}$, one can easily see that for $x = (x_1 : x_2) \in \text{dom}(T_A)$,

$$T_A(x) = \frac{ax_1 + bx_2}{cx_1 + dx_2},$$

so in particular when $x \in \mathbb{R}$,

$$T_A(x) = \frac{ax + b}{cx + d}.$$

Of course, we define $T_A(x) = \infty$ if $x \in \text{dom}(T_A)$ and $cx + d = 0$. If $\text{rank } A = 1$, the unique constant value that T_A takes is $\frac{a}{c}$ if one of a, c is not zero, or $\frac{b}{d}$ otherwise.

The reason for considering this particular class of functions is that the transforms we perform for computing the continued fraction expansion do not leave this class of functions, and as a result the whole problem of computing the continued fraction expansion only boils down to the computing the floor of $\bar{T}_A(x)$. Indeed, suppose that we know how to compute $[\bar{T}_A(x)]$ for any $A \neq 0$, assuming $\bar{T}_A(x) \in \mathbb{R}$. When $\bar{T}_A(x) = \infty$, we can regard its continued fraction is “empty”, that is, the -1 st convergent, which is always defined to be $\frac{1}{0} = \infty$ by convention, is the last convergent.

Then, we consider a new matrix

$$A' := \begin{pmatrix} c & d \\ a - [\bar{T}_A(x)]c & b - [\bar{T}_A(x)]d \end{pmatrix} = SA \quad (16)$$

where

$$S := \begin{pmatrix} 0 & 1 \\ 1 - [\bar{T}_A(x)] & \end{pmatrix}.$$

Then $\bar{T}_{A'} = \bar{T}_S \circ \bar{T}_A$ holds (even for the case $A = 0$ as both sides are just empty function in that case), thus

$$\bar{T}_{A'}(x) = \bar{T}_S(\bar{T}_A(x)) = \frac{1}{\bar{T}_A(x) - [\bar{T}_A(x)]}.$$

Hence, the 0th continued fraction coefficient of $\bar{T}_A(x)$ is $[\bar{T}_A(x)]$, the 1st coefficient of it is $[\bar{T}_{A'}(x)]$, and the 2nd coefficient can be obtained by repeating the same procedure with $A \leftarrow A'$, and so on. Note that since S is always invertible, we never get $A' = 0$ unless $A = 0$.

The main idea for computing $\lfloor \overline{T}_A(x) \rfloor$ is to use a good enough rational approximation of x . If our approximation is good enough, then we may be able to narrow down the possible range of $\lfloor \overline{T}_A(x) \rfloor$ into a single point so that we know what $\lfloor \overline{T}_A(x) \rfloor$ is. If the approximation is not good enough for that, then we refine the approximation and retry until we get a good enough one.

Since the viewpoint we want to adapt here is that it is more natural to put the problem into the framework of \mathbb{RP}^1 rather than of \mathbb{R} , we need a notion of “good approximation” in \mathbb{RP}^1 . To incorporate that idea, we adapt the notion of *cyclic order*, which is a natural replacement for \mathbb{RP}^1 of the usual linear order on \mathbb{R} .

Definition 12.1 (Cyclic order).

Given a set X , a *cyclic order* on X is a ternary relation $[\cdot, \cdot, \cdot]$ satisfying the following axioms:

1. *Cyclicity*: if $[a, b, c]$, then $[b, c, a]$.
2. *Asymmetry*: if $[a, b, c]$, then not $[c, b, a]$.
3. *Transitivity*: if $[a, b, c]$, and $[a, c, d]$, then $[a, b, d]$.
4. *Connectedness*: if a, b, c are distinct, then either $[a, b, c]$ or $[c, b, a]$.

One can think a cyclically ordered set as anything that can be arranged on a circle. Then $[a, b, c]$ means a, b, c are arranged on the circle in the counterclockwise direction. Note that $[a, b, c]$ is always false if a, b, c are not distinct. Indeed, suppose $[a, a, b]$ holds, then by the cyclicity axiom we get $[a, b, a]$, but then the asymmetry axiom contradicts itself by saying that $[a, b, a]$ should not hold. Similarly, $[a, b, b]$ is also impossible.

Identifying \mathbb{RP}^1 with $\mathbb{R} \sqcup \{\infty\}$, i.e., we “wrap” \mathbb{R} onto a circle exactly once, with $-\infty$ and $+\infty$ being attached at a single point on the circle, then it is natural to define a cyclic order on \mathbb{RP}^1 as for $a, b, c \in \mathbb{RP}^1$, $[a, b, c]$ holds if and only if:

1. $a, b, c \in \mathbb{R}$ and one of $a < b < c$, $b < c < a$, and $c < a < b$ holds, or
2. $a = \infty$, $b, c \in \mathbb{R}$ and $b < c$ holds, or
3. $b = \infty$, $c, a \in \mathbb{R}$ and $c < a$ holds, or
4. $c = \infty$, $a, b \in \mathbb{R}$ and $a < b$ holds.

In fact, there is a more natural way of writing this in terms of the representation of \mathbb{RP}^1 as a quotient space of $\mathbb{R}^2 \setminus \{0\}$ ¹¹: for $a, b, c \in \mathbb{RP}^1$ with the representation $a = (a_1 : a_2)$, $b = (b_1 : b_2)$, $c = (c_1 : c_2)$, we have

$$[a, b, c] \text{ if and only if } (a_1 b_2 - a_2 b_1)(b_1 c_2 - b_2 c_1)(c_1 a_2 - c_2 a_1) > 0. \quad (17)$$

(Note that the cyclic order on \mathbb{RP}^1 depends on how we embed \mathbb{R} into \mathbb{RP}^1 , i.e., whether we identify $(a_1 : a_2)$ with $\frac{a_1}{a_2}$ or $\frac{a_2}{a_1}$. The above equivalence is only true assuming the first convention. We should invert the inequality if we use the other convention.)

Definition 12.2 (Cyclic intervals).

Let X be a set and $[\cdot, \cdot, \cdot]$ a cyclic order on X . A *cyclic open interval* in X is a subset of X which is either \emptyset , X , the complement of a singleton set, or of the form

$$(s, t) := \{x \in X : [s, x, t]\}$$

¹¹Seung uk Jang taught me about this.

for some distinct $s, t \in X$. Similarly, a *cyclic closed interval* in X is a subset of X which is either \emptyset , X , a singleton set, or of the form

$$[s, t] := (s, t) \cup \{s, t\}$$

for some distinct $s, t \in X$. We similarly define the notion of *cyclic half-open intervals* and the notations $(s, t]$, $[s, t)$ in the obvious way. (Due to ambiguity, we never use the notation (s, t) , $(s, t]$, $[s, t)$, or $[s, t]$ when $s = t$.)

For the cyclic order on \mathbb{RP}^1 , the cyclic open interval (s, t) is given as follows:

1. If $s, t \in \mathbb{R}$ and $s < t$, then (s, t) is the usual open interval (s, t) .
2. If $s, t \in \mathbb{R}$ and $t < s$, then (s, t) is the union of $(s, \infty) \cup (-\infty, t)$ and $\{\infty\}$.
3. If $s = \infty$ and $t \in \mathbb{R}$, then (s, t) is the usual open interval $(-\infty, t)$.
4. If $s \in \mathbb{R}$ and $t = \infty$, then (s, t) is the usual open interval (s, ∞) .

Note that for \mathbb{RP}^1 , cyclic open intervals are precisely connected open subsets, while cyclic closed intervals are precisely connected closed subsets.¹²

Proposition 12.3.

Let $A \in \mathbb{R}^{2 \times 2} \setminus \{0\}$ and $[s, t] \subseteq \mathbb{RP}^1$ be any cyclic closed interval. Then:

$$\overline{T}_A [[s, t]] = \begin{cases} [\overline{T}_A(s), \overline{T}_A(t)] & \text{if } \det A > 0, \\ [\overline{T}_A(t), \overline{T}_A(s)] & \text{if } \det A < 0, \\ \{\overline{T}_A(s)\} & \text{if } \det A = 0. \end{cases}$$

Furthermore, for the first two cases, the endpoints are mapped into the endpoints.

Proof. If $\text{rank } A = 1$ then \overline{T}_A is a constant function, so we may assume that A is invertible.

Then $T_A = \overline{T}_A$ is a homeomorphism of \mathbb{RP}^1 , so $T_A [(s, t)]$ must be a nonempty connected open proper subset of \mathbb{RP}^1 , thus is a cyclic open interval (u, v) . Similarly, $T_A [[s, t]]$ is a cyclic closed interval, and by continuity this must be $[u, v]$. Since T_A is injective, we also have $\{T_A(s), T_A(t)\} = \{u, v\}$. Hence, it is enough to show that $T_A(x) \in (T_A(s), T_A(t))$ holds for all $x \in (s, t)$ if and only if $\det A > 0$.

In fact, note that for any $(u_1 : u_2)$ and $(v_1 : v_2)$ in \mathbb{RP}^1 ,

$$u_1 v_2 - u_2 v_1 = \det \begin{pmatrix} u_1 & v_1 \\ u_2 & v_2 \end{pmatrix},$$

so if $T_A((u_1 : u_2)) = (u'_1 : u'_2)$ and $T_A((v_1 : v_2)) = (v'_1 : v'_2)$, then

$$u'_1 v'_2 - u'_2 v'_1 = (\det A)(u_1 v_2 - u_2 v_1).$$

Therefore, the sign of $u_1 v_2 - u_2 v_1$ is preserved by T_A if and only if $\det A > 0$, so (17) gives the desired conclusion. \square

Therefore, if we know in advance that our input $x \in \mathbb{RP}^1$ belongs to a nonempty cyclic interval I , then we can compute $\overline{T}_A[I]$ by knowing the sign of $\det A$ and evaluating \overline{T}_A at the endpoints of I . Based on this idea, we now present Gosper's algorithm in detail.

¹²This can be easily proven by using the fact that \mathbb{RP}^1 minus any point is homeomorphic to \mathbb{R} .

The algorithm is based on the sole assumption that we have a way to obtain a decreasing sequence of cyclic intervals $(I_n)_{n=1}^\infty$ in \mathbb{RP}^1 such that $x \in I_n$ holds for all n and $\{x\} = \bigcap_{n=1}^\infty I_n$. (Recall that we do allow $x = \infty$.) In other words, we will only rely on a sequence of approximations of x and that is the only information available to us.

Suppose we know that x is in a cyclic interval $I \subseteq \mathbb{RP}^1$. Then $\bar{T}_A[I]$ should be a cyclic interval as well. If it is a singleton set, then we obtain the precise value of $\bar{T}_A(x)$, so in particular we can compute $\lfloor \bar{T}_A(x) \rfloor$. If it is not a singleton set but its closure contains ∞ , then we cannot compute $\lfloor \bar{T}_A(x) \rfloor$, so refine I and retry.

Otherwise, $\bar{T}_A[I]$ must be a bounded interval strictly contained in \mathbb{R} . Let $u, v \in \mathbb{R}$ be the left and the right endpoints of it. Then we compute $\lfloor u \rfloor$ and see if $\lfloor u \rfloor + 1 \in I$. If $\lfloor u \rfloor + 1 \notin I$, then $\lfloor \bar{T}_A(x) \rfloor$ must be equal to $\lfloor u \rfloor$. If not, then we refine the interval I and retry.

Now, note that when $\bar{T}_A(x)$ happens to be an integer, then this procedure may not succeed no matter how refined I is, unless $I = \{x\}$. Since \bar{T}_A is a continuous function, for small enough I , $\bar{T}_A[I]$ should be contained in any small neighborhood around $\bar{T}_A(x)$, but whenever $\bar{T}_A[I]$ includes an element that is strictly below $\bar{T}_A(x)$, then $\lfloor u \rfloor = \bar{v}$ can never happen, so in practice the computation procedure will trap inside an infinite loop. Similarly, when $\bar{T}_A(x) = \infty$, then that fact is discoverable through this procedure only if $\bar{T}_A(x)$ eventually becomes $\{\infty\}$.

Of course, these can happen if and only if:

1. Either $I_n = \{x\}$ for all large enough n , or
2. $\det A = 0$, or
3. $\det A > 0$ and x is the left endpoint of I_n for all large enough n , or
4. $\det A < 0$ and x is the right endpoint of I_n for all large enough n .

Assuming that coefficients a, b, c, d of A are all integers, when we repeatedly apply the transform $A \mapsto A'$ given in (16), $\bar{T}_A(x)$ eventually becomes an integer or ∞ if and only if $\bar{T}_A(x) \in \mathbb{Q} \cup \{\infty\} = \mathbb{QP}^1$. Actually, it can become ∞ only after $\bar{T}_A(x)$ becomes an integer, unless $\bar{T}_A(x) = \infty$ initially holds. Hence, $\bar{T}_A(x) \in \mathbb{QP}^1$ is the only potentially problematic case. Note that this implies there exists $(p : q) \in \mathbb{QP}^1$ such that

$$q(ax_1 + bx_2) = p(cx_1 + dx_2),$$

so

$$(qa - pc)x_1 + (qb - pd)x_2 = 0,$$

where $x = (x_1 : x_2)$. Since we assume $a, b, c, d \in \mathbb{Z}$, this is possibly the case only when $\det A = 0$ or $x \in \mathbb{QP}^1$. As pointed out before, the case $\det A = 0$ is actually fine as it always yield $\bar{T}_A[I_n] = \{\bar{T}_A(x)\}$.

Recall that our interval estimate I_n would be in practice based on rational approximations of x , and we may assume that I_n is always with rational endpoints, and $I_n = \{x\}$ holds for large enough n if and only if x is rational. Hence, we always end up with either $\det A = 0$ or $I_n = \{x\}$ whenever $\bar{T}_A(x) \in \mathbb{QP}^1$, so this infinite loop problem does not actually exist. Nevertheless, I want to point out that other than this potential infinite loop problem, everything should work out without any problem even if we do not impose any further assumptions on I_n 's other than that it is decreasing and $\bigcap_{n=1}^\infty I_n = \{x\}$.

In summary, if we assume that I_n is always with rational endpoints and $I_n = \{x\}$ holds for large enough n if and only if x is rational, then the described algorithm always succeeds in computing the continued fraction expansion of $\bar{T}_A(x)$ indefinitely.

Remark 12.4.

In Gosper's note, he claims that the interval estimate for \bar{T}_A can be done a little bit better by observing that, except for the first coefficient matrix A , we always have $\bar{T}_A(x) \in (1, \infty]$ because $\bar{T}_{A'}(x) = (\bar{T}_A(x) - \lfloor \bar{T}_A(x) \rfloor)^{-1}$. Therefore, we can further impose the condition $\bar{T}_A(x) \in (1, \infty]$ when computing $\lfloor \bar{T}_A(x) \rfloor$ which may narrow down the list of possible values of $\lfloor \bar{T}_A(x) \rfloor$ a little bit. However, this actually does not improve the interval estimate. Indeed, note that $\bar{T}_{A'}[I]$ is necessarily equal to $\bar{T}_S[\bar{T}_A[I]]$ where

$$S := \begin{pmatrix} 0 & 1 \\ 1 - \lfloor \bar{T}_A(x) \rfloor & \end{pmatrix}.$$

That we were able to find $\lfloor \bar{T}_A(x) \rfloor$ means that we found out that $\bar{T}_A[I]$ is contained in the interval $[\lfloor \bar{T}_A(x) \rfloor, \lfloor \bar{T}_A(x) \rfloor + 1)$, which is equivalent to say that $\bar{T}_{A'}[I]$ is contained in $(1, \infty]$. Therefore, imposing $\bar{T}_{A'}(x) \in (1, \infty]$ does not help.

12.2 A special case of unary Gosper's algorithm

The special case of the algorithm explained in the previous section, when the cyclic closed interval I_n is chosen to be the interval formed by the $(n-1)$ th and the n th convergents of x , is often called *Gosper's algorithm*. I am not sure if this special case was his original proposal before he ended up writing the notes I cited [11], but it seems like this version is more widely recognized as *Gosper's algorithm* than the one actually written in the cited notes. This seemingly wider exposure is the sole reason why I included this special case here even though I never needed it in the reference implementation.

This version of Gosper's algorithm is somewhat special in that we do not actually need to compute the endpoints of the cyclic intervals. Rather, we just need to update the coefficients a, b, c, d in an appropriate way when a new convergent is supposed to be evaluated.

Suppose that x admits the continued fraction expansion

$$x = a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \dots}}}$$

Recall the recurrence relation

$$P_n := \begin{pmatrix} p_n & p_{n-1} \\ q_n & q_{n-1} \end{pmatrix} = \begin{pmatrix} p_{n-1} & p_{n-2} \\ q_{n-1} & q_{n-2} \end{pmatrix} \begin{pmatrix} a_n & 1 \\ 1 & 0 \end{pmatrix} = \dots = \begin{pmatrix} p_{-1} & p_{-2} \\ q_{-1} & q_{-2} \end{pmatrix} \prod_{i=0}^n \begin{pmatrix} a_i & 1 \\ 1 & 0 \end{pmatrix} \quad (18)$$

where we define

$$\begin{pmatrix} p_{-1} & p_{-2} \\ q_{-1} & q_{-2} \end{pmatrix} := \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}.$$

Let $C_n := \begin{pmatrix} a_n & 1 \\ 1 & 0 \end{pmatrix}$ so that $P_n = P_{n-1}C_n$. Since x is in between $\frac{p_{n-1}}{q_{n-1}}$ and $\frac{p_n}{q_n}$, we want to evaluate \bar{T}_A at those two rational numbers. However, since $P_n = \begin{pmatrix} p_n & p_{n-1} \\ q_n & q_{n-1} \end{pmatrix}$, we see that $\bar{T}_A(p_n/q_n)$ and $\bar{T}_A(p_{n-1}/q_{n-1})$ are nothing but (the equivalence classes in \mathbb{RP}^1 of) two columns of the matrix AP_n . Hence, all we need to do is to keep updating the coefficient matrix AP_n by successively multiplying C_n 's.

The procedure of keep multiplying C_n to the previous coefficient matrix at each iteration can be interpreted in another way too. Note that if we define

$$x_0 := x, \quad x_{i+1} := T_{C_i^{-1}}(x_i)$$

for each i , then

$$\bar{T}_A(x) = \bar{T}_{AC_0}(x_1) = \bar{T}_{AC_0C_1}(x_2) = \cdots = \bar{T}_{AP_n}(x_{n+1}).$$

Since

$$C_i^{-1} = \begin{pmatrix} 0 & 1 \\ 1 & -a_i \end{pmatrix},$$

in fact we have

$$x_{i+1} = \frac{1}{x_i - a_i},$$

so inductively we get $a_{i+1} = \lfloor x_{i+1} \rfloor$ for each i . Note that this is exactly the procedure of getting the continued fraction expansion of x . Then x_n is nothing but the n th complete fraction of x , so we can interpret the computation of AP_n from AP_{n-1} as the replacement of x_{n-1} by its one-level deeper nested fraction of the expansion of x .

Since each x_n for $n > 0$ should be in $(1, \infty]$, one can interpret the procedure of comparing $T_A(p_n/q_n)$ and $T_A(p_{n-1}/q_{n-1})$ as the estimation of $\frac{at+b}{ct+d}$ from $\frac{a}{c}$ and $\frac{b}{d}$. Note that for $t \in (0, \infty)$, $\frac{at+b}{ct+d}$ would be in between $\frac{a}{c}$ and $\frac{b}{d}$, so we may say $\lfloor \frac{at+b}{ct+d} \rfloor$ is equal to $\lfloor \frac{a}{c} \rfloor$ or $\lfloor \frac{b}{d} \rfloor$ when the latter two are equal. However, this naïve idea subtly fails if we allow negative values for x, c, d , because then $\frac{at+b}{ct+d}$ no longer needs to be in between $\frac{a}{c}$ and $\frac{b}{d}$. The core reason of the failure is because the denominator of $t \mapsto \frac{at+b}{ct+d}$ can pass through 0, which means its value abruptly changes from $\pm\infty$ to $\mp\infty$ when t passes that. However, this is a natural phenomena if one works with \mathbb{RP}^1 instead of \mathbb{R} .

It seems that many sources available freely on the Internet are not so careful about this issue and suggest that inspecting the equality $\lfloor \frac{a}{c} \rfloor = \lfloor \frac{b}{d} \rfloor$ alone would be sufficient. And indeed that is wrong. For example, consider

$$T_A(x) = \frac{1-3x}{1-2x}$$

with $x = \frac{1}{4}$. Then we have $\lfloor \frac{a}{c} \rfloor = \lfloor \frac{b}{d} \rfloor = 1$ but $T_A(x) = \frac{1/4}{1-1/2} = \frac{1}{2}$ so $\lfloor T_A(x) \rfloor = 0$. Note that this failure persists even if we proceed one step further by replacing x by $x_1 = \frac{1}{1-\lfloor x \rfloor} = \frac{1}{x} = 4$, so that

$$T_{AC_0}(x_1) = \frac{x_1-3}{x_1-2}.$$

To avoid this trap, one must take into account of the orientation of the cyclic interval between $\bar{T}_A(p_n/q_n)$ and $\bar{T}_A(p_{n-1}/q_{n-1})$. Note that this cannot be done by just looking at these two values, because regardless of which one is bigger, there always exist two possible orientations: the one that avoids ∞ and the one that includes ∞ . It is only the former one that allows us to do the floor analysis. Note that, nevertheless, the correct tracking of the orientation is in fact very simple: one just needs to remember the determinant of A . If it is positive, then the direction is preserved. If it is negative, then the direction is reversed.

12.3 Binary Gosper's algorithm

We wish to extend our argument from the previous sections into the case when there are two unknowns x, y involved. It turns out, the natural class of combinations of x, y we may consider are numbers of the following form:

$$F(x, y) := \frac{axy + bx + cy + d}{exy + fx + gy + h},$$

where $a, b, c, d, e, f, g, h \in \mathbb{Z}$. Our goal is thus to find a way to compute $\lfloor F(x, y) \rfloor$. Once we are able to do so for generic a, b, c, d, e, f, g, h , then by replacing the above f by

$$F'(x, y) = \frac{exy + fx + gy + h}{(a - ke)xy + (b - kf)x + (c - kg)y + (d - kh)} \quad (19)$$

with $k := \lfloor F(x, y) \rfloor$ and keep repeating the same procedure, we can compute the continued fraction expansion of $F(x, y)$.

Again, it is much more convenient to think x, y as elements in \mathbb{RP}^1 rather than \mathbb{R} . Let

$$A := \begin{pmatrix} a & b \\ c & d \end{pmatrix}, \quad B := \begin{pmatrix} e & f \\ g & h \end{pmatrix}$$

and view \mathbb{RP}^1 as a quotient space of $\mathbb{R}^2 \setminus \{0\}$. Let $u, v \in \mathbb{R}^2 \setminus \{0\}$ and write $u = (u_1, u_2)$, $v = (v_1, v_2)$, then we can see that

$$F(u_1/u_2, v_1/v_2) = \frac{\langle u, Av \rangle}{\langle u, Bv \rangle}$$

provided that the denominators are not zero. Hence, we can naturally regard F as a function

$$F(u, v) = (\langle u, Av \rangle : \langle u, Bv \rangle) \in \mathbb{RP}^1,$$

provided that $\langle u, Av \rangle$ and $\langle u, Bv \rangle$ are not simultaneously zero. (This representation motivates us to call such a function F a *bilinear fractional mapping*.)

To be more precise, define

$$\text{dom}(F) := \{([u], [v]) \in \mathbb{RP}^1 \times \mathbb{RP}^1 : \langle u, Av \rangle \neq 0 \text{ or } \langle u, Bv \rangle \neq 0\},$$

and

$$\begin{aligned} F: \text{dom}(F) &\rightarrow \mathbb{RP}^1 \\ ([u], [v]) &\mapsto (\langle u, Av \rangle : \langle u, Bv \rangle). \end{aligned}$$

As in the unary case, we want to build an algorithm for computing the range of possible values of $F(x, y)$ assuming we know x, y are from some cyclic intervals $I, J \subseteq \mathbb{RP}^1$, respectively, and based on that, come up with a way to compute $\lfloor F(x, y) \rfloor$. More precisely, as in the unary case we assume we have decreasing sequences $\{I_n\}_{n=1}^\infty, \{J_n\}_{n=1}^\infty$ of cyclic intervals in \mathbb{RP}^1 such that $\{x\} = \bigcap_{n=1}^\infty I_n$ and $\{y\} = \bigcap_{n=1}^\infty J_n$, and given n , we compute the range $F[I_n \times J_n]$. Since F is continuous and $I \times J$ is connected, $F[I \times J]$ is a connected subset of \mathbb{RP}^1 , that is, a cyclic interval. Once we figure out this cyclic interval, if that allows only one possible value of $\lfloor F(x, y) \rfloor$, then we output that value.

However, we have to be careful about the fact that $\text{dom}(F)$ may be a proper subset of $\mathbb{RP}^1 \times \mathbb{RP}^1$. It does not make sense to compute the continued fraction expansion of $F(x, y)$

when $F(x, y)$ is not even well-defined. Recall that in the unary case, this was not a big issue, because the domain of the linear fractional mapping T_A is a proper subset of \mathbb{RP}^1 if and only if the matrix A is singular, and for such a case the mapping T_A is nothing but a constant function, unless A is entirely zero (which is an obvious degenerate case that can and should be rejected upfront). Therefore, it was not at all a problem to continuously extend T_A onto \mathbb{RP}^1 and deal with the extension instead of the partial function T_A itself.

Similarly, we may wish to see if we can continuously extend F to a larger domain, hopefully $\mathbb{RP}^1 \times \mathbb{RP}^1$. However, unlike the unary case, fully extending F onto $\mathbb{RP}^1 \times \mathbb{RP}^1$ is impossible in general. A prototypical counterexample (aside from the obvious degenerate case $A = B = 0$) is $F(x, y) = xy$: of course, there is no meaningful way of defining $F(0, \infty) = 0 \cdot \infty$.

Note that in the binary case, the function in general is much likely to possess a singularity than the unary case. In the unary case, $x \in \mathbb{RP}^1$ hits the singularity only if it is the solution to *two* equations. Since \mathbb{RP}^1 is only 1-dimensional, this is very unlikely. However, in the binary case, the domain is now 2-dimensional, so having at least one solution to *two* equations is quite likely. Indeed, for almost every practically meaningful examples of A, B , the domain of F , even after trying our best to obtain an extension, is still a proper subset of $\mathbb{RP}^1 \times \mathbb{RP}^1$.

Nevertheless, let us try to find the largest continuous extension of F . Note that $\text{dom}(F)$ is already a “large” subset; it is a *dense open subset*:

Proposition 12.5.

Let $\tilde{F}: (\mathbb{R}^2 \setminus \{0\})^2 \rightarrow \mathbb{R}^2$ be the bilinear mapping defined as

$$\tilde{F}: (u, v) \mapsto (\langle u, Av \rangle, \langle u, Bv \rangle)$$

for some $A, B \in \mathbb{R}^{2 \times 2}$. Then the projection onto $\mathbb{RP}^1 \times \mathbb{RP}^1$ of the zero set $\tilde{F}^{-1}[\{0\}]$ is compact. Therefore, $\text{dom}(F)$ is an open subset of $\mathbb{RP}^1 \times \mathbb{RP}^1$. Furthermore, if A and B are not both zero, then $\text{dom}(F)$ is dense in $\mathbb{RP}^1 \times \mathbb{RP}^1$.

Proof. Clearly, \tilde{F} is continuous so $\tilde{F}^{-1}[\{0\}]$ is closed. Note that whenever $(u, v) \in \tilde{F}^{-1}[\{0\}]$, we also have $\left(\frac{u}{|u|}, \frac{v}{|v|}\right) \in \tilde{F}^{-1}[\{0\}]$. Therefore, let S^1 be the unit circle in \mathbb{R}^2 , then the projection of $\tilde{F}^{-1}[\{0\}]$ onto $\mathbb{RP}^1 \times \mathbb{RP}^1$ and that of $\tilde{F}^{-1}[\{0\}] \cap (S^1 \times S^1)$ are same. Clearly, $\tilde{F}^{-1}[\{0\}] \cap (S^1 \times S^1)$ is compact, thus its projection onto $\mathbb{RP}^1 \times \mathbb{RP}^1$ is compact, thus is closed.

Now, suppose $\text{dom}(F)$ is not dense, that is, this projection of $\tilde{F}^{-1}[\{0\}]$ has a nonempty interior. By continuity, this implies that there is an open set $U \subseteq (\mathbb{R}^2 \setminus \{0\})^2$ such that $\tilde{F}[U]$ is contained in $\{0\}$. Hence, for given $(u, v) \in U$, the kernel of the linear functional $x \mapsto \langle u, Ax \rangle$ contains an open neighborhood of v in \mathbb{R}^2 , which means $A^T u = 0$. On the other hand, the kernel of A^T also contains an open neighborhood of u , thus we get $A = 0$. By the same reason, $B = 0$ holds as well, as desired. \square

To construct the extension of F defined on the largest possible domain, we use the following theorem by J. L. Kelley [12]: a continuous function $f: A \rightarrow Y$ from a dense subspace $A \subseteq X$ of a topological space X into a regular Hausdorff space Y admits a unique continuous extension onto X , if and only if for each $x \in X$, there exists $y \in Y$ such that for any net $(x_\alpha)_{\alpha \in D}$ in A converging to x , the corresponding net $(f(x_\alpha))_{\alpha \in D}$ in Y converges to y . Since \mathbb{RP}^1 is clearly regular and Hausdorff, we can apply this theorem to find the largest continuous extension of $F: \text{dom}(F) \rightarrow \mathbb{RP}^1$.

Lemma 12.6.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and define $F: (x, y) \mapsto (\langle x, Ay \rangle : \langle x, By \rangle)$ as before.

Suppose $([x], [y]) \in (\mathbb{RP}^1 \times \mathbb{RP}^1) \setminus \text{dom}(F)$. Then F admits a unique continuous extension onto $\text{dom}(F) \cup \{([x], [y])\}$ if and only if one of the following conditions hold:

1. A and B are linearly dependent, or
2. $Ay = By = 0$, or
3. $A^T x = B^T x = 0$.

Proof. Let $R := \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$ be the counterclockwise 90° -rotation matrix. Then we can write

$$\begin{aligned} A &= axy^T + bx(Ry)^T + c(Rx)y^T + d(Rx)(Ry)^T, \\ B &= exy^T + fx(Ry)^T + g(Rx)y^T + h(Rx)(Ry)^T \end{aligned}$$

for some $a, b, c, d, e, f, g, h \in \mathbb{R}$. Since $([x], [y])$ is not in the domain of F , we know $a = e = 0$.

Now, suppose that $Ay = By = 0$, which means $c = g = 0$. Hence,

$$\begin{aligned} A &= (bx + dRx)(Ry)^T, \\ B &= (fx + hRx)(Ry)^T. \end{aligned}$$

Thus, take any net $(([x_\alpha], [y_\alpha]))_{\alpha \in D}$ in $\text{dom}(F)$ converging to $([x], [y])$, then

$$\begin{aligned} F([x_\alpha], [y_\alpha]) &= (\langle x_\alpha, bx + dRx \rangle \langle Ry, y_\alpha \rangle : \langle x_\alpha, fx + hRx \rangle \langle Ry, y_\alpha \rangle) \\ &= (\langle x_\alpha, bx + dRx \rangle : \langle x_\alpha, fx + hRx \rangle) \end{aligned}$$

since $\langle x_\alpha, bx + Rx \rangle$ and $\langle x_\alpha, fbx + hRx \rangle$ are not simultaneously zero and $\langle Ry, y_\alpha \rangle \neq 0$ by the assumption $([x_\alpha], [y_\alpha]) \in \text{dom}(F)$. If b, f are not simultaneously zero, then d, h cannot be simultaneously zero and

$$F([x_\alpha], [y_\alpha]) = (d : h)$$

for any α . On the other hand, if b, f are not simultaneously zero, then we have

$$F([x_\alpha], [y_\alpha]) \rightarrow (b : f).$$

Therefore, $(F([x_\alpha], [y_\alpha]))_{\alpha \in D}$ always converges and the limit is determined independently of the choice of the net. Similarly, if $A^T x = B^T x = 0$, then one can show that $(F([x_\alpha], [y_\alpha]))_{\alpha \in D}$ converges to either $(c : g)$ or $(d : h)$, determined independently of the choice of the net.

Next, suppose that b, f are not simultaneously zero and also c, g are not simultaneously zero. If A, B are linearly dependent, that is, if $tA = sB$ for some $(s, t) \in \mathbb{R}^2 \setminus \{0\}$, then clearly F is constantly equal to $(s : t)$, so suppose otherwise. Then there are two cases: (1) $(d : h) \in \mathbb{RP}^1$ and $(b : f) = (c : g) \neq (d : h)$, or (2) $(b : f) \neq (c : g)$.

For the first case, take a nonzero vector $v = (v_1, v_2) \in \mathbb{R}^2$ such that $(b, f) = sv$ and $(c, g) = tv$ for some $s, t \in \mathbb{R}$. For each $n \in \mathbb{Z}_{>0}$, define

$$x_n := x - \frac{s}{n}Rx, \quad y_n := y + \frac{t}{n}Ry.$$

The clearly, $(([x_n], [y_n]))_{n=1}^\infty$ converges to $([x], [y])$. Note that

$$\langle x_n, Ay_n \rangle = \frac{bt}{n} |x|^2 |y|^2 - \frac{cs}{n} |x|^2 |y|^2 - \frac{dst}{n^2} |x|^2 |y|^2 = -\frac{dst |x|^2 |y|^2}{n^2}$$

since $bt = cs$, and similarly

$$\langle x_n, By_n \rangle = -\frac{hst|x|^2|y|^2}{n^2}.$$

Therefore, $([x_n], [y_n]) \in \text{dom}(F)$ and

$$F([x_n], [y_n]) = (d : h)$$

holds for all n . On the other hand, for each $n \in \mathbb{Z}_{>0}$, define

$$z_n := x + \frac{t}{n}Rx, \quad w_n := y + \frac{s}{n}Ry.$$

The clearly, $(([z_n], [w_n]))_{n=1}^\infty$ converges to $([x], [y])$, and

$$\begin{aligned} \langle z_n, Aw_n \rangle &= \frac{bs}{n}|x|^2|y|^2 + \frac{ct}{n}|x|^2|y|^2 + \frac{dst}{n^2}|x|^2|y|^2 \\ &= \frac{(s^2 + t^2)v_1|x|^2|y|^2}{n} + \frac{dst|x|^2|y|^2}{n^2} \end{aligned}$$

and

$$\langle z_n, Bw_n \rangle = \frac{(s^2 + t^2)v_2|x|^2|y|^2}{n} + \frac{hst|x|^2|y|^2}{n^2}.$$

Hence, $([x_n], [y_n]) \in \text{dom}(F)$ whenever n is large enough and

$$F([z_n], [w_n]) = \left(v_1 + \frac{dst}{n(s^2 + t^2)} : v_2 + \frac{hst}{n(s^2 + t^2)} \right)$$

for such n . Hence, $F([z_n], [w_n]) \rightarrow [v] \neq (d : h)$. Thus, F does not admit a continuous extension onto $\text{dom}(F) \cup \{([x], [y])\}$ in this case.

For the second case, for each $n \in \mathbb{Z}_{>0}$, define

$$(x_n, y_n) := \left(x, y + \frac{1}{n}Ry \right), \quad (z_n, w_n) := \left(x + \frac{1}{n}Rx, y \right).$$

Clearly, $([x_n], [y_n]), ([z_n], [w_n]) \rightarrow ([x], [y])$ as $n \rightarrow \infty$. Note that

$$\begin{aligned} \langle x_n, Ay_n \rangle &= \frac{b}{n}|x|^2|y|^2, \\ \langle x_n, By_n \rangle &= \frac{f}{n}|x|^2|y|^2, \end{aligned}$$

and since b and f are not simultaneously zero, we get $([x_n], [y_n]) \in \text{dom}(F)$ and

$$F([x_n], [y_n]) = (b : f).$$

Similarly, one can see that

$$F([z_n], [w_n]) = (c : g) \neq (b : f),$$

thus F does not admit a continuous extension onto $\text{dom}(F) \cup \{([x], [y])\}$ in this case. \square

Note that $\langle x, Ay \rangle$ and $\langle x, By \rangle$ can simultaneously be zero only if Ay and By are parallel, which means that the linear fractional mapping

$$F(\cdot, [y]): [x] \mapsto (\langle x, Ay \rangle : \langle x, By \rangle)$$

is singular. Hence, it is expected that the determinant of $F(\cdot, [y])$, that is,

$$\det \begin{pmatrix} y^T A^T \\ y^T B^T \end{pmatrix} = \langle RAy, By \rangle = \langle y, B^T RAy \rangle = \frac{1}{2} \langle y, (B^T RA - A^T RB)y \rangle$$

will play an important role in determining the singularities of F .

Theorem 12.7 (Characterization of singularities of bilinear fractional mappings).

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and define $F: (x, y) \mapsto (\langle x, Ay \rangle : \langle x, By \rangle)$ as before. Let $R := \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$ be the 90°-counterclockwise rotation matrix and L_1 the symmetric bilinear form on \mathbb{R}^2 defined as

$$L_1: (u, v) \mapsto \langle u, (B^T RA - A^T RB)v \rangle.$$

1. If A and B are linearly dependent, then $L_1 = 0$ and F admits a unique continuous extension onto $\mathbb{RP}^1 \times \mathbb{RP}^1$, and it is a constant function. In this case, if A, B are both scalar multiples of a nonzero matrix $C \in \mathbb{R}^{2 \times 2}$, then

$$\text{dom}(F) = \{([x], [y]) \in \mathbb{RP}^1 \times \mathbb{RP}^1 : y \notin \ker C \text{ and } [x] \neq [RCy]\}.$$

2. If $\det L_1 > 0$, then $\text{dom}(F) = \mathbb{RP}^1 \times \mathbb{RP}^1$.
3. If $\det L_1 < 0$, then the complement $(\mathbb{RP}^1 \times \mathbb{RP}^1) \setminus \text{dom}(F)$ of $\text{dom}(F)$ has exactly two elements whose second coordinates are precisely the only two distinct elements $[y_1], [y_2] \in \mathbb{RP}^1$ satisfying $L_1(y_1, y_1) = L_1(y_2, y_2) = 0$. Furthermore, F does not admit any proper continuous extension onto a larger subset of $\mathbb{RP}^1 \times \mathbb{RP}^1$.
4. If $\det L_1 = 0$ and $L_1 \neq 0$, then the second coordinate of any element in the complement $(\mathbb{RP}^1 \times \mathbb{RP}^1) \setminus \text{dom}(F)$ of $\text{dom}(F)$ is the unique element $[y_0] \in \mathbb{RP}^1$ satisfying $L_1(y_0, y_0) = 0$.

- (a) If $Ay_0 = By_0 = 0$, then F admits a unique continuous extension onto $\mathbb{RP}^1 \times \mathbb{RP}^1$. In this case, F is independent to its second variable on its domain, thus its extension is actually a linear fractional mapping of only one variable. Furthermore, A, B should be of the form $A = uv^T$, $B = wv^T$ for some $u, v, w \in \mathbb{R}^2 \setminus \{0\}$ and u, w are linearly independent. Also,

$$\text{dom}(F) = \{([x], [y]) \in \mathbb{RP}^1 \times \mathbb{RP}^1 : \langle v, y \rangle \neq 0\}.$$

- (b) If $Ay_0 \neq 0$ or $By_0 \neq 0$, then $(\mathbb{RP}^1 \times \mathbb{RP}^1) \setminus \text{dom}(F)$ consists of a single element $([x_0], [y_0])$ for some $[x_0] \in \mathbb{RP}^1$. In this case, F does not admit any proper continuous extension onto a larger subset of $\mathbb{RP}^1 \times \mathbb{RP}^1$.

5. If $L_1 = 0$, then F admits a unique continuous extension onto $\mathbb{RP}^1 \times \mathbb{RP}^1$. In this case, F is independent to its first variable on its domain, thus its extension is actually a linear fractional mapping of only one variable. If A, B are linearly independent, then A, B should be of the form $A = uv^T$, $B = uw^T$ for some $u, v, w \in \mathbb{R}^2 \setminus \{0\}$ and v, w are linearly independent. Also,

$$\text{dom}(F) = \{([x], [y]) \in \mathbb{RP}^1 \times \mathbb{RP}^1 : \langle x, u \rangle \neq 0\}.$$

Proof. As noted, for any $y \in \mathbb{R}^2$, $L_1(y, y) = 0$ happens if and only if Ay and By are parallel. Now we show each of the claims one by one.

1. If A and B are linearly dependent, then clearly $L_1 = 0$. Also, in this case, F is clearly a constant function on $\text{dom}(F)$, so by density it admits a unique continuous extension onto $\mathbb{RP}^1 \times \mathbb{RP}^1$ as a constant function. The claim about $\text{dom}(F)$ is also trivial.
2. Suppose $\det L_1 > 0$. This means that L_1 is either a positive-definite form or a negative definite form, so in particular $L_1(u, u) = 0$ holds if and only if $u = 0$. Therefore, there is no $y \in \mathbb{R}^2 \setminus \{0\}$ such that Ay and By are parallel, thus there is no $x \in \mathbb{R}^2 \setminus \{0\}$ simultaneously satisfying $\langle x, Ay \rangle = 0$ and $\langle x, By \rangle = 0$. This shows $\text{dom}(F) = \mathbb{RP}^1 \times \mathbb{RP}^1$.
3. Suppose $\det L_1 < 0$. This means that L_1 has two distinct eigenvalues with different signs. Hence, if $\lambda_1 > 0 > -\lambda_2$ are eigenvalues with associated normalized eigenvectors $u_1, u_2 \in \mathbb{R}^2 \setminus \{0\}$, then for any $a, b \in \mathbb{R}$,

$$L_1(au_1 + bu_2, au_1 + bu_2) = \lambda_1 a^2 - \lambda_2 b^2 = 0$$

if and only if $\sqrt{\lambda_1}a = \pm\sqrt{\lambda_2}b$. Therefore, there are exactly two $[y_1], [y_2] \in \mathbb{RP}^1$ such that $L_1(y_1, y_1) = L_1(y_2, y_2) = 0$, which are

$$y_1 := \sqrt{\lambda_2}u_1 + \sqrt{\lambda_1}u_2 \quad \text{and} \quad y_2 := \sqrt{\lambda_2}u_1 - \sqrt{\lambda_1}u_2.$$

Now, since Ay_1 and By_1 are parallel to each other, there exists a nonzero vector $u \in \mathbb{R}^2 \setminus \{0\}$ and $s, t \in \mathbb{R}$ such that $Ay_1 = su$ and $By_1 = tu$. Note that s, t cannot be both zero because that implies

$$(B^T RA - A^T RB)y_1 = sB^T Ru - tA^T Ru = 0,$$

but since $\det L_1 \neq 0$, the left-hand side cannot be zero. Therefore, define $x_1 := Ru$, then $[x_1]$ is the unique element in \mathbb{RP}^1 such that $([x_1], [y_1]) \notin \text{dom}(F)$. Note that we have

$$(B^T RA - A^T RB)y_1 = sB^T Ru - tA^T Ru = sB^T x_1 - tA^T x_1,$$

but again since the left-hand side is not zero, we cannot have $A^T x_1 = B^T x_1 = 0$. Therefore, the lemma shows that F cannot be continuously extended to $\text{dom}(F) \cup \{([x_1], [y_1])\}$. The same argument for y_2 shows that there uniquely exists $[x_2] \in \mathbb{RP}^1$ such that $([x_2], [y_2]) \notin \text{dom}(F)$ and F cannot be continuously extended to $\text{dom}(F) \cup \{([x_2], [y_2])\}$.

4. Suppose $\det L_1 = 0$ and $L_1 \neq 0$. This means that L_1 has two distinct eigenvalues and one of them is zero. Then by representing L_1 in the eigenbasis, it becomes immediately clear that the unique element $[y_0] \in \mathbb{RP}^1$ such that $L_1(y_0, y_0) = 0$ is the equivalence class of eigenvectors associated to the eigenvalue 0. Hence, we not only have $L_1(y_0, y_0) = 0$ but also have $B^T RAy_0 = A^T RB y_0$.

- (a) Suppose $Ay_0 = By_0 = 0$. Then there exist $u, v \in \mathbb{R}^2$ such that $A = u(Ry_0)^T$ and $B = v(Ry_0)^T$. Since $L_1 \neq 0$, A, B are linearly independent, thus u, v are linearly independent. Now, let $x, y \in \mathbb{R}^2$ be arbitrary. Then

$$\langle x, Ay \rangle = \langle x, u \rangle \langle Ry_0, y \rangle \quad \text{and} \quad \langle x, By \rangle = \langle x, v \rangle \langle Ry_0, y \rangle,$$

thus in particular if $([x], [y]) \in (\mathbb{RP}^1 \times \mathbb{RP}^1) \setminus \text{dom}(F)$, then

$$\langle x, u \rangle \langle Ry_0, y \rangle = \langle x, v \rangle \langle Ry_0, y \rangle = 0.$$

Since u, v are linearly independent, $\langle x, u \rangle$ and $\langle x, v \rangle$ cannot be simultaneously zero, thus we must have $\langle Ry_0, y \rangle = 0$. This shows the claim about the forms of A , B , and $\text{dom } F$. Also, if $\langle Ry_0, y \rangle = 0$, then $A^T x_0 = B^T x_0 = 0$ follows, thus the lemma shows that F admits a unique continuous extension onto $\text{dom}(F) \cup \{([x], [y])\}$. Since $([x], [y])$ is arbitrary, we conclude that F admits a unique continuous extension onto all of $\mathbb{RP}^1 \times \mathbb{RP}^1$. Also, since

$$F([x], [y]) = (\langle x, u \rangle \langle Ry_0, y \rangle : \langle x, v \rangle \langle Ry_0, y \rangle) = (\langle x, u \rangle : \langle x, v \rangle)$$

whenever $([x], [y]) \in \text{dom}(F)$, the unique extension \bar{F} should be given as

$$\bar{F}: ([x], [y]) \mapsto (\langle x, u \rangle : \langle x, v \rangle)$$

which is a linear fractional mapping of $[x]$.

- (b) Next, suppose $Ay_0 \neq 0$ or $By_0 \neq 0$. Without loss of generality, let us assume $Ay_0 \neq 0$. Note that $([x], [y]) \notin \text{dom}(F)$ if and only if $\langle x, Ay \rangle = \langle x, By \rangle = 0$ and it implies that Ay and By are parallel, thus we must have $L_1(y, y) = 0$. Since $[y_0]$ is the unique element such that $L_1(y_0, y_0) = 0$, $[y] = [y_0]$ follows. Since $Ay_0 \neq 0$, $[x] = [RAy_0]$ also follows. Hence, define $x_0 := RAy_0$, then the complement of $\text{dom}(F)$ is the singleton set $\{([x_0], [y_0])\}$.

Since $Ay_0 \neq 0$ and A, B are linearly independent, the lemma shows that F admits a continuous extension onto $\mathbb{RP}^1 \times \mathbb{RP}^1$ if and only if $A^T x_0 = B^T x_0 = 0$. Note that

$$RA^T x_0 = RA^T RAy_0,$$

so using the 2×2 -matrix identity

$$RA^T R = -(\text{adj } A),$$

we get

$$R(A^T x_0) = -(\text{adj } A)Ay_0 = -(\det A)y_0.$$

Therefore, $A^T x_0 = 0$ holds if and only if $\det A = 0$, that is, A is not invertible.

If $By_0 \neq 0$ in addition to $Ay_0 \neq 0$, then we must have $[Ay_0] = [By_0]$, thus we can perform the same computation to verify that $B^T x_0 = 0$ holds if and only if B is not invertible. On the other hand, if $By_0 = 0$, then we always have

$$B^T x_0 = (B^T RA)y_0 = (A^T RB)y_0 = 0.$$

Of course, in this case B must be not invertible. Therefore, F admits a unique continuous extension onto $\mathbb{RP}^1 \times \mathbb{RP}^1$ if and only if A, B are both not invertible. Since $A = 0$ or $B = 0$ yields $L_1 = 0$, this means $\text{rank } A = \text{rank } B = 1$, so there exist $u_1, v_1, u_2, v_2 \in \mathbb{R}^2 \setminus \{0\}$ such that $A = u_1 v_1^T$ and $B = u_2 v_2^T$. However, in this case we have

$$L_1 = v_1 u_1^T R u_2 v_2^T - v_2 u_2^T R u_1 v_1^T = \langle u_1, R u_2 \rangle (v_1 v_2^T + v_2 v_1^T).$$

By the 2×2 determinant sum formula $\det(A+B) = \det(A) + \det(B) + \operatorname{tr}(A) \operatorname{tr}(B) - \operatorname{tr}(AB)$, we have

$$\det(v_1 v_2^T + v_2 v_1^T) = \langle v_1, v_2 \rangle^2 - |v_1|^2 |v_2|^2,$$

so

$$\det L_1 = \langle u_1, Ru_2 \rangle^2 \left(\langle v_1, v_2 \rangle^2 - |v_1|^2 |v_2|^2 \right).$$

Since we assumed $\det L_1 = 0$ and $L_1 \neq 0$, we conclude that u_1, u_2 are not parallel and v_1, v_2 are parallel. Then since $L_1(y_0, y_0) = 0$, the only possibility is that v_1, v_2 are scalar multiples of Ry_0 . However, this contradicts to the assumption that either of Ay_0 or By_0 is nonzero.

5. Suppose $L_1 = 0$. This means that Ay and By are parallel for every $y \in \mathbb{R}^2$. Suppose first that there exists $y_0 \in \mathbb{R}^2 \setminus \{0\}$ such that $Ay_0 = By_0 = 0$. Then there exist $u, v \in \mathbb{R}^2$ such that $A = u(Ry_0)^T$ and $B = v(Ry_0)^T$. Since $A(Ry_0)$ and $B(Ry_0)$ should be parallel as well, it follows that u and v are parallel. Therefore, A and B are linearly dependent in this case.

Next, suppose there exists $y_0 \in \mathbb{R}^2 \setminus \{0\}$ such that $Ay_0 \neq 0$ and $By_0 = 0$. Again, there exists $v \in \mathbb{R}^2$ such that $B = v(Ry_0)^T$. Let $x_0 := RAy_0$, then

$$B^T x_0 = B^T RAy_0 = A^T RBy_0 = 0.$$

Since $B^T x_0 = \langle v, x_0 \rangle Ry_0$, we get $\langle v, x_0 \rangle = \langle v, RAy_0 \rangle = 0$, thus v is parallel to Ay_0 . Hence, there exists $\lambda \in \mathbb{R}$ such that $B = \lambda(Ay_0)(Ry_0)^T$. Clearly, if $\lambda = 0$ then A and B are linearly dependent, so we may assume $\lambda \neq 0$. Since By is parallel to Ay for all $y \in \mathbb{R}^2$, and By is parallel to Ay_0 for all $y \in \mathbb{R}^2$, and By is not zero for all y not parallel to y_0 , it follows that $A = (Ay_0)u^T$ for some $u \in \mathbb{R}^2$. If u and Ry_0 are linearly dependent, then A and B are linearly dependent, so suppose they are linearly independent.

Now, take arbitrary $x, y \in \mathbb{R}^2$. Then by the representations of A, B we got, we know

$$\langle x, Ay \rangle = \langle x, Ay_0 \rangle \langle u, y \rangle \quad \text{and} \quad \langle x, By \rangle = \langle x, Ay_0 \rangle \langle Ry_0, y \rangle,$$

thus in particular if $([x], [y]) \in (\mathbb{RP}^1 \times \mathbb{RP}^1) \setminus \operatorname{dom}(F)$, then

$$\langle x, Ay_0 \rangle \langle u, y \rangle = \langle x, Ay_0 \rangle \langle Ry_0, y \rangle = 0.$$

Since u and Ry_0 are linearly independent, $\langle u, y \rangle$ and $\langle Ry_0, y \rangle$ cannot be simultaneously zero. Therefore, we must have $\langle x, Ay_0 \rangle = 0$. Note that this means $A^T x = \langle x, Ay_0 \rangle u$ and $B^T x = \langle x, Ay_0 \rangle Ry_0$ are both zero, thus by the lemma we conclude that F extends to $\operatorname{dom}(F) \cup \{([x], [y])\}$. Since $([x], [y])$ is arbitrary, it follows that F admits a unique continuous extension to all of $\mathbb{RP}^1 \times \mathbb{RP}^1$. Also, since

$$F([x], [y]) = (\langle x, Ay_0 \rangle \langle u, y \rangle : \langle x, Ay_0 \rangle \langle Ry_0, y \rangle) = (\langle u, y \rangle : \langle Ry_0, y \rangle)$$

whenever $([x], [y]) \in \operatorname{dom}(F)$, the unique extension \overline{F} should be given as

$$\overline{F}: ([x], [y]) \mapsto (\langle u, y \rangle : \langle Ry_0, y \rangle),$$

which is a linear fractional mapping of $[y]$. By symmetry, we get the same conclusion if we assume there exists $y_0 \in \mathbb{R}^2 \setminus \{0\}$ such that $Ay_0 = 0$ and $By_0 \neq 0$.

Finally, assume Ay, By are both nonzero for all $y \in \mathbb{R}^2 \setminus \{0\}$. This means that A and B are both invertible. Since $A^{-1}By$ is parallel to y for all $y \in \mathbb{R}^2$, every vector in \mathbb{R}^2 is an eigenvector of $A^{-1}B$. This is only possible if $A^{-1}B$ is a scalar matrix, so A and B are linearly dependent. Hence, if A and B are linearly independent, then there exists $y_0 \in \mathbb{R}^2 \setminus \{0\}$ with $Ay_0 \neq 0$ and $By_0 = 0$, or $Ay_0 = 0$ and $By_0 \neq 0$. We have already shown that in this case A, B and $\text{dom}(F)$ are of the claimed forms.

□

Definition 12.8 (Extended bilinear fractional mappings).

We denote the unique maximal continuous extension of $F: \text{dom}(F) \rightarrow \mathbb{RP}^1$ appearing in the above theorem as $\overline{F}: \text{dom}(\overline{F}) \rightarrow \mathbb{RP}^1$. The complement of $\text{dom}(\overline{F})$ in $\mathbb{RP}^1 \times \mathbb{RP}^1$ is called the *indeterminacy locus* of F .

Corollary 12.9.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero.

1. If $A = sC, B = tC$ for some $C \in \mathbb{R}^{2 \times 2} \setminus \{0\}$, then \overline{F} is constantly equal to $(s : t)$.
2. If A, B are linearly independent and there exist $u, v, w \in \mathbb{R}^2 \setminus \{0\}$ such that $A = uv^T, B = uw^T$, then

$$\overline{F}(\cdot, \cdot): (x, y) \mapsto (\langle v, y \rangle : \langle w, y \rangle)$$

and is defined everywhere on $\mathbb{RP}^1 \times \mathbb{RP}^1$.

3. If A, B are linearly independent and there exists $u, v, w \in \mathbb{R}^2 \setminus \{0\}$ such that $A = uv^T, B = vw^T$, then

$$\overline{F}(\cdot, \cdot): (x, y) \mapsto (\langle x, u \rangle : \langle x, w \rangle)$$

and is defined everywhere on $\mathbb{RP}^1 \times \mathbb{RP}^1$.

4. Otherwise, $\overline{F} = F$. In this case,

- (a) $\det L_1 > 0$ if and only if the indeterminacy locus is empty,
- (b) $\det L_1 = 0$ if and only if the indeterminacy locus has one element, and
- (c) $\det L_1 < 0$ if and only if the indeterminacy locus has two elements,

where we define $L_1: (x, y) \mapsto \langle x, (B^T R A - A^T R B)y \rangle$ as in Theorem 12.7.

For the first three cases, \overline{F} is a proper extension of F .

Proof. The first three cases follow easily from the definition of \overline{F} as a continuous extension of F . Suppose those are not the case, then Theorem 12.7 shows that we must have $F = \overline{F}$ and the number of elements in the indeterminacy locus is determined according to the determinant of L_1 as claimed. □

Consequently, \overline{F} either:

1. is a constant function, or
2. is a linear fractional mapping of just one variable, or
3. does not add any new point to the domain of F .

So far, we have established the exact function \overline{F} we want to work with. Next, our goal is to figure out how to compute the image $\overline{F}[(I \times J) \cap \text{dom}(\overline{F})]$ of $I \times J$ under \overline{F} , where I, J are arbitrary nonempty cyclic intervals in \mathbb{RP}^1 . Obviously, we can just do what we did for the unary case if \overline{F} is either a constant function or a linear fractional mapping of one variable, so from now on let us suppose that $\overline{F} = F$.

We need some preliminary works. First, we have to figure out how F behaves around the indeterminacy locus. Let $([x_0], [y_0])$ be a point in the indeterminacy locus. Clearly, since Ay_0 and By_0 are parallel and not both zero, the function $F(\cdot, [y_0])$ is a rank 1 linear fractional mapping. In the same reason, $F([x_0], \cdot)$ is a rank 1 linear fractional mapping. In other words, F is constant on $(\mathbb{RP}^1 \times \{[y_0]\}) \cap \text{dom}(F)$ and on $(\{[x_0]\} \times \mathbb{RP}^1) \cap \text{dom}(F)$, while these two constants may be different. Let us first see what precisely is the set of points in $\text{dom}(F)$ where F achieves these constants.

Proposition 12.10.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and suppose $\overline{F} = F$. Then for given $(s : t) \in \mathbb{RP}^1$, $S := F^{-1}[\{(s : t)\}] \subseteq \text{dom}(F)$ can be described as follows.

1. If $tA - sB$ is invertible, then

$$\begin{aligned} S &= \{([x], [(tA - sB)^{-1}Rx]) \in \text{dom}(F) : [x] \in \mathbb{RP}^1\} \\ &= \{([(tA^T - sB^T)^{-1}Ry], [y]) \in \text{dom}(F) : [y] \in \mathbb{RP}^1\} \end{aligned}$$

where $R := \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$ is the 90° -counterclockwise rotation matrix.

2. If $tA - sB$ is singular, then it is of rank 1 and

$$S = ((\{[x_0]\} \times \mathbb{RP}^1) \cup (\mathbb{RP}^1 \times \{[y_0]\})) \cap \text{dom}(F)$$

where $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $y_0 \in \ker(tA - sB) \setminus \{0\}$. In this case, $([x_0], [y_0])$ is in the indeterminacy locus of F if and only if it is the unique point there.

Proof. Note that for $([x], [y]) \in \text{dom}(F)$, $F([x], [y]) = (\langle x, Ay \rangle : \langle x, By \rangle) = (s : t)$ holds if and only if $\langle x, (tA - sB)y \rangle = 0$. If $tA - sB$ is invertible, then each $[x] \in \mathbb{RP}^1$ uniquely corresponds to some $[y] \in \mathbb{RP}^1$ such that $\langle x, (tA - sB)y \rangle = 0$, namely $[y] = [(tA - sB)^{-1}Rx]$, and similarly each $[y] \in \mathbb{RP}^1$ uniquely corresponds to some $[x] \in \mathbb{RP}^1$ such that $\langle x, (tA - sB)y \rangle = 0$, namely $[x] = [(tA^T - sB^T)^{-1}Ry]$.

Next, suppose $tA - sB$ is singular. Clearly, $tA - sB = 0$ implies that A, B are linearly dependent, but since we assumed $\overline{F} = F$, it cannot be the case. Hence, $tA - sB$ is of rank 1. Then we can write $tA - sB = uv^T$ for some nonzero $u, v \in \mathbb{R}^2$. Hence, clearly $\langle x, (tA - sB)y \rangle = 0$ holds if and only if either $[x] = [Ru]$ or $[y] = [Rv]$. Then it can be easily seen that $[x_0] := [Ru]$ is the unique element in \mathbb{RP}^1 such that $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $[y_0] := [Rv]$ is the unique element in \mathbb{RP}^1 such that $y_0 \in \ker(tA - sB) \setminus \{0\}$. This shows the claimed representation of S .

Now, suppose that $([x_0], [y_0])$ is in the indeterminacy locus of F . Then either $Ay_0 = 0$ or $[x_0] = [RAy_0]$ holds, and similarly either $By_0 = 0$ or $[x_0] = [RBy_0]$ holds, but not both of $Ay_0 = By_0 = 0$ hold. Suppose $Ay_0 = 0$, then since $tAy_0 = sBy_0$, we obtain $s = 0 \neq t$. Then $x_0 \in \ker(tA^T - sB^T)$ implies $A^T x_0 = 0$, thus $A^T RBy_0 = 0$. Therefore, $(A^T RB - B^T RA)y_0 = 0$, which shows $\det L = 0$. As noted in Corollary 12.9, this implies that $([x_0], [y_0])$ is the unique element of the indeterminacy locus. Next, suppose $Ay_0 \neq 0$, then $tAy_0 = sBy_0$ shows that $s \neq 0$. Since $[x] = [RAy_0]$ holds, $x \in \ker(tA^T - sB^T)$ shows

$$tA^T RAy_0 = sB^T RAy_0,$$

and substituting $tAy_0 = sBy_0$ on the left-hand side gives

$$sA^T RBy_0 = sB^T RAy_0,$$

concluding $(B^T RA - A^T RB)y_0 = 0$. Again, this shows that $([x_0], [y_0])$ is the unique element of the indeterminacy locus. \square

Corollary 12.11.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and suppose $\overline{F} = F$. For a point (x_0, y_0) in the indeterminacy locus, let $(s_1 : t_1)$ be the constant taken by $F(\cdot, y_0)$ and $(s_2 : t_2)$ be the constant taken by $F(x_0, \cdot)$, then

$$\begin{aligned} F^{-1}[\{(s_1 : t_1)\}] &= ((\{[x'_0]\} \times \mathbb{RP}^1) \cup (\mathbb{RP}^1 \times \{[y_0]\})) \cap \text{dom}(F) \text{ and} \\ F^{-1}[\{(s_2 : t_2)\}] &= ((\{[x_0]\} \times \mathbb{RP}^1) \cup (\mathbb{RP}^1 \times \{[y'_0]\})) \cap \text{dom}(F) \end{aligned}$$

where $x'_0 \in \ker(t_1 A^T - s_1 B^T) \setminus \{0\}$ and $y'_0 \in \ker(t_2 A - s_2 B) \setminus \{0\}$. Furthermore, we have

$$(\mathbb{RP}^1 \times \mathbb{RP}^1) \setminus \text{dom}(F) = \{([x_0], [y_0]), ([x'_0], [y'_0])\},$$

and $(s_1 : t_1) = (s_2 : t_2)$ holds if and only if $x_0 = x'_0$ holds if and only if $y_0 = y'_0$ holds.

Proof. Follows easily from Proposition 12.10; details are omitted. \square

Note that, if $F(\cdot, [y_0])$ is constant on $(\mathbb{RP}^1 \times \{[y_0]\}) \cap \text{dom}(F)$, then Ay_0 and By_0 are parallel, thus by Theorem 12.7, there should exist $[x_0] \in \mathbb{RP}^1$ such that $([x_0], [y_0])$ is in the indeterminacy locus of F , and this $[x_0]$ is precisely the unique element such that $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ where $(s : t) \in \mathbb{RP}^1$ is the unique element such that $tAy_0 = sBy_0$.

In particular, whenever F is constant on $(\mathbb{RP}^1 \times \{[y_0]\}) \cap \text{dom}(F)$, there exists a neighborhood $V = ([u], [v])$ of $[y_0] \in \mathbb{RP}^1$ such that F is not constant on $(\mathbb{RP}^1 \times \{[y]\}) \cap \text{dom}(F)$ for all $[y] \in V \setminus \{[y_0]\}$. This in fact also shows $\mathbb{RP}^1 \times \{[y]\} \subseteq \text{dom}(F)$, so $F(\cdot, [y])$ is a homeomorphism from \mathbb{RP}^1 onto \mathbb{RP}^1 for all such $[y]$. Hence, provided that V is small enough, the sign of the determinant of $F(\cdot, [y])$ is constant on the open cyclic intervals $([u], [y_0])$ and $([y_0], [v])$. Let us denote these signs as *the determinant sign of $F(\cdot, [y_0]^-)$* and *the determinant sign of $F(\cdot, [y_0]^+)$* , respectively. We can similarly define the determinant signs of $F([x_0]^-, \cdot)$ and $F([x_0]^+, \cdot)$ as well. Formally:

Definition 12.12.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and suppose $\overline{F} = F$. Given $(x_0, y_0) \in \mathbb{RP}^1 \times \mathbb{RP}^1$, the limits

$$\sigma(F(x^\pm, \cdot)) := \lim_{x \rightarrow x_0^\pm} \frac{\det F(x, \cdot)}{|\det F(x, \cdot)|} \quad \text{and} \quad \sigma(F(\cdot, y^\pm)) := \lim_{y \rightarrow y_0^\pm} \frac{\det F(\cdot, y)}{|\det F(\cdot, y)|},$$

all four of which should always exist, are called *the determinant signs of $F(x^\pm, \cdot)$* and *the determinant signs of $F(\cdot, y^\pm)$* , respectively.

These determinant signs play an important role in determining the behavior of F near the indeterminacy locus.

Proposition 12.13.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and suppose $\overline{F} = F$. Let $([x_0], [y_0])$ be a point in the

indeterminacy locus of F . Let $R := \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$ be the 90° -counterclockwise rotation matrix and $L_1, L_2: \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$ be bilinear forms defined as

$$\begin{aligned} L_1: (x, y) &\mapsto \langle x, (B^T R A - A^T R B)y \rangle, \\ L_2: (x, y) &\mapsto \langle x, (B R A^T - A R B^T)y \rangle. \end{aligned}$$

1. If $([x_0], [y_0])$ is the unique point in the indeterminacy locus of F , then $L_1(Ry_0, Ry_0)$, $L_2(Rx_0, Rx_0)$ are both nonzero and

$$\sigma(F(\cdot, [y_0]^\pm)) = \frac{L_1(Ry_0, Ry_0)}{|L_1(Ry_0, Ry_0)|}, \quad \sigma(F([x_0]^\pm, \cdot)) = \frac{L_2(Rx_0, Rx_0)}{|L_2(Rx_0, Rx_0)|}.$$

2. Otherwise, $L_1(y_0, Ry_0)$, $L_2(x_0, Rx_0)$ are both nonzero and

$$\sigma(F(\cdot, [y_0]^\pm)) = \frac{\mp L_1(y_0, Ry_0)}{|L_1(y_0, Ry_0)|} = \frac{\pm L_2(x_0, Rx_0)}{|L_2(x_0, Rx_0)|} = \sigma(F([x_0]^\mp, \cdot)).$$

Proof. For given $[y] \in \mathbb{RP}^1$, recall that

$$\det F(\cdot, [y]) = \det \begin{pmatrix} y^T A^T \\ y^T B^T \end{pmatrix} = \frac{1}{2} L_1(y, y),$$

and similarly for given $[x] \in \mathbb{RP}^1$,

$$\det F([x], \cdot) = \det \begin{pmatrix} x^T A \\ x^T B \end{pmatrix} = \frac{1}{2} L_2(x, x).$$

Hence, to see the sign of $\det F(\cdot, [y])$ near $[y_0]$, it is enough to see how the sign of the function

$$\sigma_1: \epsilon \mapsto L_1(y_0 - \epsilon Ry_0, y_0 - \epsilon Ry_0) = -2\epsilon L_1(y_0, Ry_0) + \epsilon^2 L_1(Ry_0, Ry_0),$$

defined near 0 changes around 0, where the equality follows from $L_1(y_0, y_0) = 0$. Similarly, for $F([x_0]^\pm, \cdot)$, we consider the function

$$\sigma_2: \epsilon \mapsto -2\epsilon L_2(x_0, Rx_0) + \epsilon^2 L_2(Rx_0, Rx_0).$$

Now, let

$$P := \begin{pmatrix} \frac{x_0}{|x_0|} & \frac{Rx_0}{|x_0|} \end{pmatrix}, \quad Q := \begin{pmatrix} \frac{y_0}{|y_0|} & \frac{Ry_0}{|y_0|} \end{pmatrix},$$

then

$$\begin{aligned} (P^T B Q)^T R (P^T A Q) - (P^T A Q)^T R (P^T B Q) &= Q^T (B^T (P R P^T) A - A^T (P R P^T) B) Q \\ &= Q^T (B^T R A - A^T R B) Q, \end{aligned}$$

where the second equality is because P, R are both rotation matrices, which implies $PR = RP$ and $P^T = P^{-1}$, so

$$P R P^T = R P P^T = R.$$

Let us write

$$P^T A Q = \begin{pmatrix} a & b \\ c & d \end{pmatrix} =: A', \quad P^T B Q = \begin{pmatrix} e & f \\ g & h \end{pmatrix} =: B',$$

and define a bilinear form $L'_1: (x, y) \mapsto L_1(Qx, Qy) = \langle x, (B'^T RA' - A'^T RB')y \rangle$. By direct computation, one can see

$$B'^T RA' - A'^T RB' = \begin{pmatrix} 0 & bg - cf \\ bg - cf & 2(bh - df) \end{pmatrix}$$

Similarly, we can define a bilinear form $L'_2: (x, y) \mapsto L_2(Px, Py)$ then the matrix entries of L_2 are given as

$$A'RB'^T - B'RA'^T = \begin{pmatrix} 0 & cf - bg \\ cf - bg & 2(ch - dg) \end{pmatrix}.$$

Then one can easily see that

$$\begin{aligned} L_1(y_0, Ry_0) &= |y_0|^2 (bg - cf), \\ L_1(Ry_0, Ry_0) &= 2|y_0|^2 (bh - df), \\ L_2(x_0, Rx_0) &= -|x_0|^2 (bg - cf), \text{ and} \\ L_2(Rx_0, Rx_0) &= 2|x_0|^2 (ch - dg) \end{aligned}$$

hold.

1. By Corollary 12.9, we know $\det L_1 = \det(B'^T RA' - A'^T RB') = 0$, that is, $bg = cf$. Thus, we get $L_1(y_0, Ry_0) = L_2(x_0, Rx_0) = 0$. Also, $L_1(Ry_0, Ry_0) = 0$ yields $bh = df$, so $L_1 = 0$, which is impossible by Theorem 12.7 as we assumed $\bar{F} = F$. Similarly, $L_2(Rx_0, Rx_0) = 0$ yields $L_2 = 0$, so applying Theorem 12.7 after swapping the roles of x, y yields a contradiction. Hence, we conclude that $L_1(Ry_0, Ry_0)$ and $L_2(Rx_0, Rx_0)$ are both nonzero. Therefore,

$$\sigma(F(\cdot, [y_0]^\pm)) = \lim_{\epsilon \rightarrow 0^\pm} \frac{\sigma_1(\epsilon)}{|\sigma_1(\epsilon)|} = \frac{L_1(Ry_0, Ry_0)}{|L_1(Ry_0, Ry_0)|}$$

and similarly

$$\sigma(F([x_0]^\pm, \cdot)) = \lim_{\epsilon \rightarrow 0^\pm} \frac{\sigma_2(\epsilon)}{|\sigma_2(\epsilon)|} = \frac{L_2(Rx_0, Rx_0)}{|L_2(Rx_0, Rx_0)|}.$$

2. In this case, Corollary 12.9 shows that $\det L_1 \neq 0$, so $L_1(y_0, Ry_0)$ and $L_2(x_0, Rx_0)$ are both nonzero and of opposite signs. Then,

$$\sigma(F(\cdot, [y_0]^\pm)) = \lim_{\epsilon \rightarrow 0^\pm} \frac{\sigma_1(\epsilon)}{|\sigma_1(\epsilon)|} = \frac{\mp L_1(y_0, Ry_0)}{|L_1(y_0, Ry_0)|} = \frac{\pm L_2(x_0, Rx_0)}{|L_2(x_0, Rx_0)|} = \sigma(F([x_0]^\mp, \cdot))$$

follows.

□

Remark 12.14.

Since $\det F(\cdot, y) = 0$ can hold only when y is associated to a point in the indeterminacy locus, Proposition 12.13 shows that when there is only one point (x_0, y_0) in the indeterminacy locus, then every $\det F(\cdot, y)$ for $y \neq y_0$ has the identical sign to $\sigma(F(\cdot, y_0^\pm))$, because otherwise there must exist some $y \in \mathbb{RP}^1$ such that $\det F(\cdot, y) = 0$. By the same reason, every $\det F(x, \cdot)$ for $x \neq x_0$ has the identical sign to $\sigma(F(x_0^\pm, \cdot))$.

Using these, we can precisely describe local images of F near the indeterminacy locus.

Proposition 12.15.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and suppose $\bar{F} = F$.

1. Suppose (x_0, y_0) is the unique point in the indeterminacy locus of F and $k \in \mathbb{RP}^1$ the common value of F along the horizontal and the vertical lines passing through (x_0, y_0) minus (x_0, y_0) . Then we have the followings:

- If $\sigma(F(\cdot, y_0^\pm)) = \sigma(F(x_0^\pm, \cdot)) = +1$, then

$$\begin{aligned} F[(\{u\} \times (v, y_0)) \cup ((u, x_0) \times \{v\})] &= (F(u, v), k), \\ F[((x_0, u) \times \{v\}) \cup (\{u\} \times (v, y_0))] &= \mathbb{RP}^1 \setminus \{F(u, v), k\}, \\ F[(\{u\} \times (y_0, v)) \cup ((x_0, u) \times \{v\})] &= (k, F(u, v)), \\ F[((u, x_0) \times \{v\}) \cup (\{u\} \times (y_0, v))] &= \mathbb{RP}^1 \setminus \{F(u, v), k\} \end{aligned}$$

for any $u, v \in \mathbb{RP}^1$ such that $u \neq x_0, v \neq y_0$.

- If $\sigma(F(\cdot, y_0^\pm)) = \sigma(F(x_0^\pm, \cdot)) = -1$, then

$$\begin{aligned} F[(\{u\} \times (v, y_0)) \cup ((u, x_0) \times \{v\})] &= (k, F(u, v)), \\ F[((x_0, u) \times \{v\}) \cup (\{u\} \times (v, y_0))] &= \mathbb{RP}^1 \setminus \{F(u, v), k\}, \\ F[(\{u\} \times (y_0, v)) \cup ((x_0, u) \times \{v\})] &= (F(u, v), k), \\ F[((u, x_0) \times \{v\}) \cup (\{u\} \times (y_0, v))] &= \mathbb{RP}^1 \setminus \{F(u, v), k\} \end{aligned}$$

for any $u, v \in \mathbb{RP}^1$ such that $u \neq x_0, v \neq y_0$.

- If $\sigma(F(\cdot, y_0^\pm)) = +1$ and $\sigma(F(x_0^\pm, \cdot)) = -1$, then

$$\begin{aligned} F[(\{u\} \times (v, y_0)) \cup ((u, x_0) \times \{v\})] &= \mathbb{RP}^1 \setminus \{F(u, v), k\}, \\ F[((x_0, u) \times \{v\}) \cup (\{u\} \times (v, y_0))] &= (k, F(u, v)), \\ F[(\{u\} \times (y_0, v)) \cup ((x_0, u) \times \{v\})] &= \mathbb{RP}^1 \setminus \{F(u, v), k\}, \\ F[((u, x_0) \times \{v\}) \cup (\{u\} \times (y_0, v))] &= (F(u, v), k) \end{aligned}$$

for any $u, v \in \mathbb{RP}^1$ such that $u \neq x_0, v \neq y_0$.

- If $\sigma(F(\cdot, y_0^\pm)) = -1$ and $\sigma(F(x_0^\pm, \cdot)) = +1$, then

$$\begin{aligned} F[(\{u\} \times (v, y_0)) \cup ((u, x_0) \times \{v\})] &= \mathbb{RP}^1 \setminus \{F(u, v), k\}, \\ F[((x_0, u) \times \{v\}) \cup (\{u\} \times (v, y_0))] &= (F(u, v), k), \\ F[(\{u\} \times (y_0, v)) \cup ((x_0, u) \times \{v\})] &= \mathbb{RP}^1 \setminus \{F(u, v), k\}, \\ F[((u, x_0) \times \{v\}) \cup (\{u\} \times (y_0, v))] &= (k, F(u, v)) \end{aligned}$$

for any $u, v \in \mathbb{RP}^1$ such that $u \neq x_0, v \neq y_0$.

2. Suppose $(x_1, y_1), (x_2, y_2)$ are distinct points in the indeterminacy locus of F and $k_1, k_2 \in \mathbb{RP}^1$ the common values of F along the horizontal and the vertical lines passing through $(x_1, y_2), (x_2, y_1)$, respectively. Then we have the followings:

- If $\sigma(F(\cdot, y_1^+)) = \sigma(F(x_1^-, \cdot)) = +1$,

$$\begin{aligned} F[(\{u_-\} \times (v_-, y_1)) \cup ((u_-, x_1) \times \{v_-\})] &= (k_1, k_2) \setminus \{F(u_-, v_-)\}, \\ F[(\{x_1, u_+\} \times \{v_-\}) \cup (\{u_+\} \times (v_-, y_1))] &= (k_2, k_1) \setminus \{F(u_+, v_-)\}, \\ F[(\{u_+\} \times (y_1, v_+)) \cup ((x_1, u_+) \times \{v_+\})] &= (k_1, k_2) \setminus \{F(u_+, v_+)\}, \\ F[(\{u_-, x_1\} \times \{v_+\}) \cup (\{u_-\} \times (y_1, v_+))] &= (k_2, k_1) \setminus \{F(u_-, v_+)\} \end{aligned}$$

for any $u_- \in (x_2, x_1)$, $u_+ \in (x_1, x_2)$, $v_- \in (y_2, y_1)$ and $v_+ \in (y_1, y_2)$.

- If $\sigma(F(\cdot, y_1^+)) = \sigma(F(x_1^-, \cdot)) = -1$,

$$\begin{aligned} F[(\{u_-\} \times (v_-, y_1)) \cup ((u_-, x_1) \times \{v_-\})] &= (k_2, k_1) \setminus \{F(u_-, v_-)\}, \\ F[(\{x_1, u_+\} \times \{v_-\}) \cup (\{u_+\} \times (v_-, y_1))] &= (k_1, k_2) \setminus \{F(u_+, v_-)\}, \\ F[(\{u_+\} \times (y_1, v_+)) \cup ((x_1, u_+) \times \{v_+\})] &= (k_2, k_1) \setminus \{F(u_+, v_+)\}, \\ F[(\{u_-, x_1\} \times \{v_+\}) \cup (\{u_-\} \times (y_1, v_+))] &= (k_1, k_2) \setminus \{F(u_-, v_+)\} \end{aligned}$$

for any $u_- \in (x_2, x_1)$, $u_+ \in (x_1, x_2)$, $v_- \in (y_2, y_1)$ and $v_+ \in (y_1, y_2)$.

- Proof.* 1. Follows immediately from Proposition 12.13 and Proposition 12.3 by observing that $\det F(\cdot, v)$ has the same sign with $\sigma(F(\cdot, y_0^\pm))$ and similarly $\det F(u, \cdot)$ has the same sign with $\sigma(F(x_0^\pm, \cdot))$.
2. Follows immediately from Proposition 12.13 and Proposition 12.3 by observing that $\det F(\cdot, v_+)$ and $\det F(u_-, \cdot)$ have the same sign with $\sigma(F(\cdot, y_1^+))$ and similarly $\det F(\cdot, v_-)$ and $\det F(u_+, \cdot)$ have the opposite sign with $\sigma(F(y_1^+, \cdot))$. \square

Here is another useful result.

Proposition 12.16.

Suppose $A, B \in \mathbb{R}^{2 \times 2}$ are linearly independent. Then $F: \text{dom}(F) \rightarrow \mathbb{RP}^1$ is an open map.

Proof. It is enough to show that for $([x], [y]) \in \text{dom}(F)$ and a connected neighborhood $U \subseteq \text{dom}(F)$ of $([x], [y])$, $F([x], [y])$ is in the interior of $F[U]$.

Define $\tilde{F}: (\mathbb{R}^2 \setminus \{0\})^2 \rightarrow \mathbb{R}^2$ as

$$\tilde{F}: (u, v) \mapsto (\langle u, Av \rangle, \langle u, Bv \rangle)$$

as before and consider a function $G: (-\delta, \delta)^2 \rightarrow \mathbb{R}$ defined as

$$G: (\epsilon_1, \epsilon_2) \mapsto \left\langle \tilde{F}(x + \epsilon_1 Rx, y + \epsilon_2 Ry), R\tilde{F}(x, y) \right\rangle$$

where $\delta > 0$ is small enough to ensure that $([x + \epsilon_1 Rx], [y + \epsilon_2 Ry])$ is contained in U . Note that if $([x], [y])$ is not in the interior of the cyclic interval $F[U]$, then G must attain a local extrema at $(0, 0)$. We will show that this is only possible when A, B are linearly dependent.

For notational simplicity, define

$$\begin{aligned} a &:= \langle x, Ay \rangle, \quad b := \langle x, ARy \rangle, \quad c := \langle Rx, Ay \rangle, \quad d := \langle Rx, ARy \rangle, \\ e &:= \langle x, By \rangle, \quad f := \langle x, BRy \rangle, \quad g := \langle Rx, By \rangle, \quad h := \langle Rx, BRy \rangle. \end{aligned}$$

Note that

$$\tilde{F}(x + \epsilon_1 Rx, y + \epsilon_2 Ry) = (a + c\epsilon_1 + b\epsilon_2 + d\epsilon_1\epsilon_2, e + g\epsilon_1 + f\epsilon_2 + h\epsilon_1\epsilon_2)$$

thus

$$\begin{aligned} G(\epsilon_1, \epsilon_2) &= \left\langle \tilde{F}(x + \epsilon_1 Rx, y + \epsilon_2 Ry), (-e, a) \right\rangle \\ &= (ag - ce)\epsilon_1 + (af - be)\epsilon_2 + (ah - de)\epsilon_1\epsilon_2. \end{aligned}$$

If $(0, 0)$ is a critical point of G , then the first derivative of G should vanish at $(0, 0)$, which means $af = be$ and $ag = ce$ hold. In this case, the second derivative of G becomes the bilinear form

$$D^2G(0, 0): (h_1, h_2) \mapsto 2(ah - de)h_1h_2.$$

Therefore, if the first derivative of G vanishes at $(0, 0)$, then $(0, 0)$ is either a saddle point of G or all of $af = be$, $ag = ce$ and $ah = de$ hold. In other words, $(0, 0)$ is a local extrema of G only if the latter condition holds, but that implies A and B are linearly dependent, as claimed. Hence, when A, B are linearly independent, $F([x], [y])$ must be in the interior of $F[U]$, so $F[U]$ is open. \square

Next, we need some more detailed understanding of cyclic order in general. Specifically, we want to see that a cyclic order can be always extended to an n -ary relation for arbitrary positive integer n , where $[x_1, \dots, x_n]$ means x_1, \dots, x_n are arranged cyclically in that order. Here I wrote a proof of this fact in detail.

Definition 12.17 (Monotone functions).

A function $f: X \rightarrow Y$ between cyclically ordered sets X, Y is said to be *monotone* if $[f(a), f(b), f(c)]$ implies $[a, b, c]$ for any $a, b, c \in X$. We say f an *embedding* if f is monotone and injective, and an *isomorphism* if it is monotone and bijective.

Proposition 12.18.

Let X, Y be cyclically ordered sets and assume $|X| \geq 3$. Then a function $f: X \rightarrow Y$ is an embedding if and only if $[a, b, c]$ implies $[f(a), f(b), f(c)]$ for any $a, b, c \in X$.

Proof. (\Rightarrow) Let a, b, c be distinct elements in X satisfying $[a, b, c]$. Then since f is injective, we have either $[f(a), f(b), f(c)]$ or $[f(c), f(b), f(a)]$. Since f is monotone, the second case is impossible as it implies $[c, b, a]$, thus we get $[f(a), f(b), f(c)]$.

(\Leftarrow) Let $a, b, c \in X$. Suppose $[f(a), f(b), f(c)]$, then a, b, c must be distinct, so either $[a, b, c]$ or $[c, b, a]$ should hold. The latter is thus impossible because it implies $[f(c), f(b), f(a)]$, contradicting to $[f(a), f(b), f(c)]$. Hence, we obtain $[a, b, c]$, so f is monotone. To show that f is injective, let a, b be distinct elements in X . Then since $|X| \geq 3$, we can find another distinct element $c \in X$. Then either $[a, b, c]$ or $[c, b, a]$ should hold, thus by the assumption we must have either $[f(a), f(b), f(c)]$ or $[f(c), f(b), f(a)]$. For any cases, we have that $f(a) \neq f(b)$, thus concluding that f is an injection. \square

Proposition 12.19.

For a positive integer n , consider the group \mathbb{Z}/n of integers modulo n . We define a cyclic order on \mathbb{Z}/n as $[a, b, c]$ if and only if there exists $k \in \mathbb{Z}/n$ with

$$(a + k \bmod n) < (b + k \bmod n) < (c + k \bmod n).$$

Then a bijection $\sigma: \mathbb{Z}/n \rightarrow \mathbb{Z}/n$ is an isomorphism (of cyclically ordered sets) if and only if there exists $k \in \mathbb{Z}/n$ such that $\sigma(i) = i + k$ for all $i \in \mathbb{Z}/n$.

Proof. Clearly, if such k exists, then σ is an isomorphism. For the other direction, we use induction on n . The conclusion we want to draw is vacuously true for $n \leq 2$, so take $n \geq 2$ and suppose that it is true for that n . Let $\sigma: \{0, \dots, n\} \rightarrow \{0, \dots, n\}$ be an isomorphism

(when we identify $\{0, \dots, n\}$ with $\mathbb{Z}/(n+1)$), then by adding an appropriate element in $\mathbb{Z}/(n+1)$, we can assume without loss of generality that $\sigma(n) = 0$. Then since $\sigma|_{\{0, \dots, n-1\}}$ is an isomorphism, the induction hypothesis shows that there exists $k = 0, \dots, n-1$ such that $\sigma(i) + k \equiv i \pmod{n}$ holds for all $i = 0, \dots, n-1$.

Now we show k must be zero. Suppose not, then $n-k$, n , and 0 are distinct elements in $\{0, \dots, n\}$. Adding $k+1$ and then taking the modulo operation by $n+1$ then maps them into $0, k, k+1$, respectively, thus by definition of the cyclic order on $\mathbb{Z}/(n+1)$, we must have $[n-k, n, 0]$. However, since σ is an isomorphism, this implies $[\sigma(n-k), \sigma(n), \sigma(k)]$, which means $[0, n, k]$ with $k \leq n-1$. This contradicts to the definition of the cyclic order on $\mathbb{Z}/(n+1)$, so we must have $k = 0$, so σ is the identity map on $\mathbb{Z}/(n+1)$. \square

Definition 12.20.

Let X be a cyclically ordered set. Then for any positive integer n , we define an n -ary relation $[\cdot, \dots, \cdot]$ on X as $[x_0, \dots, x_{n-1}]$ (or $[x_i]_{i \in \mathbb{Z}/n}$ as an abbreviation) if and only if the function $\mathbb{Z}/n \rightarrow X$ given as $i \mapsto x_i$ is an embedding. In particular, we always have $[x_0]$, and $[x_0, x_1]$ holds if and only if $x_0 \neq x_1$.

Lemma 12.21.

Let X be a cyclically ordered set.

1. *When $n = 3$, the ternary relation defined above is precisely the original cyclic order.*
2. *For any $\{x_i\}_{i \in \mathbb{Z}/n} \subseteq X$, $[x_i]_{i \in \mathbb{Z}/n}$ implies $[x_{i+1}]_{i \in \mathbb{Z}/n}$.*
3. *For any x_1, \dots, x_{n+1} , if $[x_1, \dots, x_n]$ and $[x_n, x_{n+1}, x_1]$ hold, then $[x_1, \dots, x_n, x_{n+1}]$ holds.*
4. *Let $x_0, \dots, x_n \in X$ be distinct. Then there uniquely exists a permutation $\sigma \in S_n$ such that $[x_0, x_{\sigma(1)}, \dots, x_{\sigma(n)}]$ holds.*

Proof. 1. To distinguish the two relations, let us denote the newly defined relation as $\langle \cdot, \cdot, \cdot \rangle$ for a moment. For given $x_0, x_1, x_2 \in X$, suppose $[x_0, x_1, x_2]$. Then x_0, x_1, x_2 are necessarily distinct, so the mapping $i \mapsto x_i$ from $\mathbb{Z}/3$ is injective. Then it can be easily checked that it is monotone, thus is an embedding. Hence, we conclude $\langle x_0, x_1, x_2 \rangle$. Conversely, suppose $\langle x_0, x_1, x_2 \rangle$. Note that this implies $|X| \geq 3$, so we get $[x_0, x_1, x_2]$ from the relation $[0, 1, 2]$ on $\mathbb{Z}/3$. Therefore, the desired equivalence is established.

2. Follows from that $i \mapsto i-1$ is an isomorphism on \mathbb{Z}/n .
3. Let $i, j, k \in \{1, \dots, n+1\}$, $i < j < k$. Then it is enough to show that $[x_i, x_j, x_k]$ holds. This is clear if $k \leq n$, so let $k = n+1$. If $j = n$, then we can apply the transitivity axiom to $[x_n, x_{n+1}, x_1]$ and $[x_n, x_1, x_i]$ to conclude $[x_n, x_{n+1}, x_i]$. Inductively, if we have $[x_i, x_j, x_{n+1}]$ and $i < j-1$, then $[x_i, x_{j-1}, x_j]$ and $[x_i, x_j, x_{n+1}]$ together gives $[x_i, x_{j-1}, x_{n+1}]$. Hence, we have $[x_i, x_j, x_k]$ for all $i < j < k$, thus $[x_1, \dots, x_{n+1}]$ follows.
4. Uniqueness follows directly from Proposition 12.19. To show existence, we use induction on n . The base case $n = 1$ is trivial. Suppose the conclusion is true for some $n \geq 1$, and let $x_0, \dots, x_{n+1} \in X$ be any distinct elements. Then there exists a permutation $\sigma \in S_n$ such that $[x_0, x_{\sigma(1)}, \dots, x_{\sigma(n)}]$ holds. Without loss of generality, we may assume σ is the identity. We claim that at least one of $[x_0, x_{n+1}, x_1]$, $[x_1, x_{n+1}, x_2], \dots, [x_{n-1}, x_{n+1}, x_n], [x_n, x_{n+1}, x_0]$ holds.

Suppose not, then we have all of $[x_1, x_{n+1}, x_0]$, $[x_2, x_{n+1}, x_1]$, \dots , $[x_n, x_{n+1}, x_{n-1}]$, $[x_0, x_{n+1}, x_n]$. From $[x_{n+1}, x_0, x_1]$ and $[x_{n+1}, x_1, x_2]$, applying the transitivity axiom gives us $[x_{n+1}, x_0, x_2]$. Then again applying the transitivity axiom to it together with $[x_{n+1}, x_2, x_3]$ gives $[x_{n+1}, x_0, x_3]$. Thus, inductively, we obtain $[x_{n+1}, x_0, x_n]$, which contradicts to $[x_0, x_{n+1}, x_n]$. This shows the claim.

Therefore, there exists $j = 0, \dots, n$ such that $[x_j, x_{n+1}, x_{(j+1) \bmod (n+1)}]$ holds. Hence, applying 3 shows $[x_{(j+1) \bmod (n+1)}, \dots, x_n, x_0, \dots, x_j, x_{n+1}]$, or equivalently, $[x_0, \dots, x_j, x_{n+1}, x_{j+1}, \dots, x_n]$, thus we have found a desired permutation σ . \square

As an application, we show a simple statement about the *Hausdorff distance* on \mathbb{RP}^1 . Here, it does not matter which specific metric we endow \mathbb{RP}^1 with; what matters is that we do have one so that at each point in \mathbb{RP}^1 , we can find a neighborhood base consisting of open cyclic intervals whose diameters shrink to zero.

Lemma 12.22.

Let $a, b \in \mathbb{RP}^1$ and $(a_n)_{n=1}^\infty, (b_n)_{n=1}^\infty$ be sequences in \mathbb{RP}^1 converging to a, b , respectively, such that $a_n \neq b_n$ for each $n \in \mathbb{Z}_{>0}$.

1. If $a \neq b$, then $([a_n, b_n])_{n=1}^\infty$ converges to $[a, b]$ in the Hausdorff metric d_H .
2. If $a = b$ and there exists a sequence $(c_n)_{n=1}^\infty$ in \mathbb{RP}^1 which does not have a as an accumulation point and $[a_n, c_n, b_n]$ holds for all $n \in \mathbb{Z}_{>0}$, then $([a_n, b_n])_{n=1}^\infty$ converges to \mathbb{RP}^1 in the Hausdorff metric d_H .
3. If $a = b$ and there exists a sequence $(c_n)_{n=1}^\infty$ in \mathbb{RP}^1 which does not have a as an accumulation point and $[b_n, c_n, a_n]$ holds for all $n \in \mathbb{Z}_{>0}$, then $([a_n, b_n])_{n=1}^\infty$ converges to $\{a\}$ in the Hausdorff metric d_H .

Proof. 1. Since $a \neq b$, any small enough open neighborhoods (a_-, a_+) of a and (b_-, b_+) of b satisfy $[a_-, a_+] \cap [b_-, b_+] = \emptyset$. Then we have $[a_-, a_+, b_-]$ and $[b_-, b_+, a_-]$, thus $[a_-, a_+, b_-, b_+]$, or more precisely, $[a_-, a, a_+, b_-, b, b_+]$. In particular, $[a_+, b_-] \subseteq [a, b] \subseteq [a_-, b_+]$. Similarly, for sufficiently large n , we must have $a_n \in (a_-, a_+)$ and $b_n \in (b_-, b_+)$, so $[a_-, a_n, a_+, b_-, b_n, b_+]$ holds, which implies $[a_+, b_-] \subseteq [a_n, b_n] \subseteq [a_-, b_+]$. Since the diameters of (a_-, a_+) and (b_-, b_+) can be set to be arbitrarily small and $d_H([a_+, b_-], [a_-, b_+])$ is bounded by the maximum of those two (because $[a_-, b_+] \setminus [a_+, b_-] = [a_-, a_+] \cup (b_-, b_+]$), it can be set to be arbitrarily small too. This shows that $[a_n, b_n] \rightarrow [a, b]$ as $n \rightarrow \infty$.

2. We show that any subsequence of $([a_n, b_n])_{n=1}^\infty$ admits a further subsequence convergent to \mathbb{RP}^1 . For ease of notation, let us just write $([a_n, b_n])_{n=1}^\infty$ to denote one of its subsequences. Since a is not an accumulation point of $(c_n)_{n=1}^\infty$, any subsequence of $(c_n)_{n=1}^\infty$ does not converge to c . By compactness of \mathbb{RP}^1 , we can find a convergent further subsequence of the subsequence of $(c_n)_{n=1}^\infty$ corresponding to the subsequence we chose for $([a_n, b_n])_{n=1}^\infty$. Again for ease of notation, let us write this further subsequence and the corresponding further subsequence of $([a_n, b_n])_{n=1}^\infty$ just as $(c_n)_{n=1}^\infty$, $([a_n, b_n])_{n=1}^\infty$, respectively.

Let $c \in \mathbb{RP}^1$ be the limit of $(c_n)_{n=1}^\infty$. Then by the assumption, we know $a \neq c$, so we have $[a_-, a_+] \cap [c_-, c_+] = \emptyset$ for any small enough neighborhoods (a_-, a_+) and (c_-, c_+) of a , respectively. In particular, we have $[a_+, c, a_-]$.

Take a sufficiently large n so that $a_n, b_n \in (a_-, a_+)$ and $c_n \in (c_-, c_+)$ hold. Then since $a_n, b_n \notin [c_-, c_+]$, we obtain $[c_-, c_n, c_+]$, $[c_+, a_n, c_-]$ and $[c_+, b_n, c_-]$. Combining the

first two gives $[c_+, a_n, c_-, c_n]$, which gives $[a_n, c_-, c_n]$. Combining it with $[c_n, b_n, a_n]$ then gives $[a_n, c_-, c_n, b_n]$, in particular $[a_n, c_-, b_n]$. Combining it with $[c_-, c_+, b_n]$ then gives $[b_n, a_n, c_-, c_+]$, which together with $[c_-, c, c_+]$ implies $[c_+, b_n, a_n, c_-, c]$, in particular $[a_n, c, b_n]$.

On the other hand, since $a_n, b_n \in (a_-, a_+)$, we have $[a_-, a_n, a_+]$ and $[a_-, b_n, a_+]$. These together with $[a_+, c, a_-]$ imply $[a_-, a_n, a_+, c]$ and $[a_-, b_n, a_+, c]$. Then $[a_n, a_+, c]$ together with $[c, b_n, a_n]$ implies $[a_n, a_+, c, b_n]$, and $[b_n, a_n, a_+, c]$ together with $[c, a_-, b_n]$ implies $[b_n, a_n, a_+, c, a_-]$, or equivalently, $[a_n, a_+, c, a_-, b_n]$. Therefore, $[a_n, b_n]$ contains $[a_+, a_-]$, or equivalently, $\mathbb{RP}^1 \setminus [a_n, b_n]$ is contained in (a_-, a_+) . Since the diameter of the latter can be made arbitrarily small, we obtain $[a_n, b_n] \rightarrow \mathbb{RP}^1$ as $n \rightarrow \infty$.

3. By the same logic as in 2 but with a_n and b_n swapped, we conclude that a further subsequence of any subsequence of $([a_n, b_n])_{n=1}^\infty$ is eventually contained in an arbitrary neighborhood (a_-, a_+) of a . This shows $[a_n, b_n] \rightarrow \{a\}$ as $n \rightarrow \infty$. \square

Given these, now let us see in details how to compute $F[(I \times J) \cap \text{dom}(F)]$ for arbitrary nonempty cyclic intervals I, J .

First, suppose that at least one of I, J is a singleton set. Consider the case when J is a singleton set; the case when I is a singleton set can be done in a similar way. Let $J = \{[y_0]\}$. Then $F(\cdot, [y_0])$ is a linear fractional mapping. Then Theorem 12.7 shows that the function $F(\cdot, [y_0])$ is either a homeomorphism onto \mathbb{RP}^1 or $[y_0]$ is the second coordinate of a point in the indeterminacy locus, where the second case happens exactly when $L_1(y_0, y_0) = 0$. For the second case, Theorem 12.7 shows that Ay_0, By_0 are not both zero, so there uniquely exists $(s : t) \in \mathbb{RP}^1$ such that $tAy_0 = sBy_0$ and $F([x], [y_0]) = (s : t)$ for all $[x] \in \mathbb{RP}^1$ such that $([x], [y_0]) \in \text{dom}(F)$. Hence, in this case, $F[(I \times J) \cap \text{dom}(F)] = \{(s : t)\}$ if $I \neq \{[x_0]\}$ where $[x_0] \in \mathbb{RP}^1$ is the unique element such that $([x_0], [y_0]) \notin \text{dom}(F)$, and $F[(I \times J) \cap \text{dom}(F)] = \emptyset$ if $I = \{[x_0]\}$. If $L_1(y_0, y_0) \neq 0$, then $F(\cdot, [y_0])$ is defined everywhere and $F[(I \times J) \cap \text{dom}(F)] = F(\cdot, [y_0])[I]$ can be computed using Proposition 12.3.

Next, suppose that at least one of I, J is \mathbb{RP}^1 and both of them are not singleton sets. Then we must have $F[(I \times J) \cap \text{dom}(F)] = \mathbb{RP}^1$. Indeed, suppose $I = \mathbb{RP}^1$; the case $J = \mathbb{RP}^1$ is same. Since J is a nontrivial cyclic interval, it has infinitely many distinct elements. But since there are at most two points in the indeterminacy locus, there exists at least one $[y] \in J$ such that $F(\cdot, [y])$ is a homeomorphism onto \mathbb{RP}^1 . Thus we get $F[(I \times J) \cap \text{dom}(F)] = \mathbb{RP}^1$.

If one of I, J is the complement of a single point and the other is not singleton, then $F[(I \times J) \cap \text{dom}(F)]$ is either \mathbb{RP}^1 or the complement of a single point. Indeed, suppose that $J = \mathbb{RP}^1 \setminus \{[y_0]\}$. Since I is assumed to be a nontrivial cyclic interval, we can find two distinct $[x_1], [x_2] \in I$ such that $F([x_1], \cdot)$ and $F([x_2], \cdot)$ are both homeomorphisms onto \mathbb{RP}^1 . Thus, each of the images of $J = \mathbb{RP}^1 \setminus \{[y_0]\}$ by those maps avoids exactly one point in \mathbb{RP}^1 , namely $F([x_1], [y_0])$ and $F([x_2], [y_0])$, respectively. Therefore, if these two are distinct points, then we get $F[(I \times J) \cap \text{dom}(F)] = \mathbb{RP}^1$. Note however that they are equal if and only if $F(\cdot, [y_0])$ is not a homeomorphism, that is, it is constantly equal to some $(s : t) \in \mathbb{RP}^1$. By Proposition 12.10, there exists $([x], [y]) \in (I \times J) \cap \text{dom}(F)$ such that $F([x], [y]) = (s : t)$ if and only if the unique element $[x_0] \in \mathbb{RP}^1$ with $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ belongs to I . Therefore, we have $F[(I \times J) \cap \text{dom}(F)] = \mathbb{RP}^1$ if and only if either $L_1(y_0, y_0) \neq 0$ or $L_1(y_0, y_0) = 0$ but the unique element $[x_0] \in \mathbb{RP}^1$ with $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ belongs to I . Otherwise, we have $F[(I \times J) \cap \text{dom}(F)] = \mathbb{RP}^1 \setminus \{(s : t)\}$.

The only remaining cases are now that I, J are both nontrivial cyclic intervals bounded by two distinct endpoints. Let $u_1, u_2 \in \mathbb{RP}^1$ be the left and the right endpoints of I , and $v_1, v_2 \in \mathbb{RP}^1$ the left and the right endpoints of J . Let us first consider the case when I, J

are both closed and $I \times J \subseteq \text{dom}(F)$. In this case, the main idea is that, because of the “basically monotone” behavior of F (quoting Bill Gosper), the image $F[I \times J]$ should be the union of images of F on each of the four edges of the cyclic rectangle $I \times J$. Since the restrictions of F onto those edges are linear fractional mappings, we know how to compute the image of a cyclic interval. The proposition below formally proves this claim, allowing us to compute $F[I \times J]$ by separately computing the images of four edges of $I \times J$ using Proposition 12.3 and then joining them.

Proposition 12.23.

Let $[\lambda, \rho], [\beta, \tau]$ be cyclic closed intervals in \mathbb{RP}^1 such that $[\lambda, \rho] \times [\beta, \tau] \subseteq \text{dom}(F)$. Then

$$F[[\lambda, \rho] \times [\beta, \tau]] = F[[\lambda, \rho] \times \{\beta\}] \cup F[\{\rho\} \times [\beta, \tau]] \\ \cup F[[\lambda, \rho] \times \{\tau\}] \cup F[\{\lambda\} \times [\beta, \tau]].$$

Proof. Clearly, the right-hand side is contained in the left-hand side. To show the other direction, note that we can assume without loss of generality that $[\lambda, \rho] = [\beta, \tau] = [0, \infty]$, by performing an appropriate change of basis on A and B . Then any $(x, y) \in [\lambda, \rho] \times [\beta, \tau]$ can be uniquely written as $((1-s)s, (1-t)t)$ for $(s, t) \in [0, 1] \times [0, 1]$.

Now, consider the function $\tilde{F}: [0, 1] \times [0, 1] \rightarrow \mathbb{R}^2$ given as

$$\tilde{F}(s, t) \mapsto \left((1-s)s A \begin{pmatrix} 1-t \\ t \end{pmatrix}, (1-s)s B \begin{pmatrix} 1-t \\ t \end{pmatrix} \right)$$

so that $F((1-s)s, (1-t)t) = [\tilde{F}(s, t)]$. Clearly, we have

$$\tilde{F}(s, t) = (1-s)(1-t)\tilde{F}(0, 0) + (1-s)t\tilde{F}(0, 1) + s(1-t)\tilde{F}(1, 0) + st\tilde{F}(1, 1)$$

for all $(s, t) \in [0, 1] \times [0, 1]$, thus $F[[\lambda, \rho] \times [\beta, \tau]]$ is the image under the natural projection $[\cdot]: \mathbb{R}^2 \setminus \{0\} \rightarrow \mathbb{RP}^1$ of the convex hull K of four points $\tilde{F}(0, 0)$, $\tilde{F}(0, 1)$, $\tilde{F}(1, 0)$, and $\tilde{F}(1, 1)$ in \mathbb{R}^2 . Then it is enough to show that any straight line passing through a point in K should intersect with one of the four line segments $[\tilde{F}(0, 0), \tilde{F}(1, 0)]$, $[\tilde{F}(1, 0), \tilde{F}(1, 1)]$, $[\tilde{F}(1, 1), \tilde{F}(0, 1)]$, and $[\tilde{F}(0, 1), \tilde{F}(0, 0)]$. This is proved in the lemma below. \square

Lemma 12.24.

Let $\{x_i\}_{i \in \mathbb{Z}/n} \subseteq \mathbb{R}^2$ be any finite collection of points and let y be any point in the convex hull of $\{x_i\}_{i \in \mathbb{Z}/n}$. Then any straight line passing through y must intersect with at least one of the line segments in $\{[x_i, x_{i+1}]\}_{i \in \mathbb{Z}/n}$.

Proof. Without loss of generality, we may let $y = 0$ and the given straight line is the horizontal axis. Define

$$P := \{x_i : i \in \mathbb{Z}/n, x_{i2} > 0\}, \\ N := \{x_i : i \in \mathbb{Z}/n, x_{i2} < 0\}, \\ Z := \{x_i : i \in \mathbb{Z}/n, x_{i2} = 0\},$$

where x_{i2} is the vertical coordinate of x_i . Clearly, we cannot have $P = Z = \emptyset$ because in that case any point in the convex hull of $\{x_i\}_{i \in \mathbb{Z}/n}$ should have strictly negative vertical coordinate, contradicting to $y = 0$. In the same reason, we cannot have $N = Z = \emptyset$.

Now, if $Z \neq \emptyset$, then the horizontal axis must pass through any line segment $[x_i, x_{i+1}]$ such that either $x_i \in Z$ or $x_{i+1} \in Z$, so we may assume $Z = \emptyset$, so that P and N are both nonempty. Hence, there exists i such that $x_i \in P$ and $x_{i+1} \in N$. Then the horizontal line must pass through the line segment $[x_i, x_{i+1}]$. \square

Therefore, at least when I, J are closed and $I \times J \subseteq \text{dom}(F)$, we only need to know how to compute the union

$$F[[\lambda, \rho] \times [\beta, \tau]] = F[[\lambda, \rho] \times \{\beta\}] \cup F[\{\rho\} \times [\beta, \tau]] \\ \cup F[[\lambda, \rho] \times \{\tau\}] \cup F[\{\lambda\} \times [\beta, \tau]].$$

In fact, it turns out that basically the same approach works even when the aforementioned assumptions do not hold.

Theorem 12.25 (Cyclic rectangle mapping algorithm).

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and suppose $\bar{F} = F$. Let I and J be any cyclic intervals in \mathbb{RP}^1 such that $(\lambda, \rho) \subseteq I \subseteq [\lambda, \rho]$ and $(\beta, \tau) \subseteq J \subseteq [\beta, \tau]$ for some $\lambda \neq \rho, \beta \neq \tau$ in \mathbb{RP}^1 . Let $\{(x_1, y_1), (x_2, y_2)\}$ be the indeterminacy locus of F (where $(x_1, y_1) = (x_2, y_2)$ is allowed) and let $k_1, k_2 \in \mathbb{RP}^1$ be the common value of F assumed on the horizontal and the vertical lines passing through (x_1, y_2) and (x_2, y_1) , respectively, except on $\{(x_1, y_1), (x_2, y_2)\}$. With $(i, j) \in \{(1, 2), (2, 1)\}$, define

$$E_1 := \begin{cases} [F(\lambda, \beta), F(\rho, \beta)] & \text{if } \det F(\cdot, \beta) > 0, \\ [F(\rho, \beta), F(\lambda, \beta)] & \text{if } \det F(\cdot, \beta) < 0, \\ \{k_j\} & \text{if } y_i = \beta \text{ and } x_i \notin [\lambda, \rho], \\ \mathbb{RP}^1 & \text{if } y_i = \beta \text{ and } x_i \in (\lambda, \rho), \\ [k_i, k_j] & \text{if } y_i = \beta, x_i = \lambda, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \beta^+)) = +1, \\ [k_j, k_i] & \text{if } y_i = \beta, x_i = \lambda, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \beta^+)) = -1, \\ \{k_j\} & \text{if } y_i = \beta, x_i = \lambda, k_1 = k_2 \text{ and } \sigma(F(\cdot, \beta^+)) = \sigma(F(\lambda^+, \cdot)), \\ \mathbb{RP}^1 & \text{if } y_i = \beta, x_i = \lambda, k_1 = k_2 \text{ and } \sigma(F(\cdot, \beta^+)) \neq \sigma(F(\lambda^+, \cdot)), \\ [k_j, k_i] & \text{if } y_i = \beta, x_i = \rho, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \beta^+)) = +1, \\ [k_i, k_j] & \text{if } y_i = \beta, x_i = \rho, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \beta^+)) = -1, \\ \mathbb{RP}^1 & \text{if } y_i = \beta, x_i = \rho, k_1 = k_2 \text{ and } \sigma(F(\cdot, \beta^+)) = \sigma(F(\rho^-, \cdot)), \\ \{k_j\} & \text{if } y_i = \beta, x_i = \rho, k_1 = k_2 \text{ and } \sigma(F(\cdot, \beta^+)) \neq \sigma(F(\rho^-, \cdot)), \end{cases}$$

$$E_2 := \begin{cases} [F(\rho, \beta), F(\rho, \tau)] & \text{if } \det F(\rho, \cdot) > 0, \\ [F(\rho, \tau), F(\rho, \beta)] & \text{if } \det F(\rho, \cdot) < 0, \\ \{k_i\} & \text{if } x_i = \rho \text{ and } y_i \notin (\beta, \tau), \\ \mathbb{RP}^1 & \text{if } x_i = \rho \text{ and } y_i \in (\beta, \tau), \end{cases}$$

$$\begin{aligned}
E_3 := & \begin{cases} [F(\lambda, \tau), F(\rho, \tau)] & \text{if } \det F(\cdot, \tau) > 0, \\ [F(\rho, \tau), F(\lambda, \tau)] & \text{if } \det F(\cdot, \tau) < 0, \\ \{k_j\} & \text{if } y_i = \tau \text{ and } x_i \notin [\lambda, \rho], \\ \mathbb{RP}^1 & \text{if } y_i = \tau \text{ and } x_i \in (\lambda, \rho), \\ [k_i, k_j] & \text{if } y_i = \tau, x_i = \lambda, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \tau^-)) = +1, \\ [k_j, k_i] & \text{if } y_i = \tau, x_i = \lambda, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \tau^-)) = -1, \\ \mathbb{RP}^1 & \text{if } y_i = \tau, x_i = \lambda, k_1 = k_2 \text{ and } \sigma(F(\cdot, \tau^-)) = \sigma(F(\lambda^+, \cdot)), \\ \{k_j\} & \text{if } y_i = \tau, x_i = \lambda, k_1 = k_2 \text{ and } \sigma(F(\cdot, \tau^-)) \neq \sigma(F(\lambda^+, \cdot)), \\ [k_j, k_i] & \text{if } y_i = \tau, x_i = \rho, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \tau^-)) = +1, \\ [k_i, k_j] & \text{if } y_i = \tau, x_i = \rho, k_1 \neq k_2 \text{ and } \sigma(F(\cdot, \tau^-)) = -1, \\ \{k_j\} & \text{if } y_i = \tau, x_i = \rho, k_1 = k_2 \text{ and } \sigma(F(\cdot, \tau^-)) = \sigma(F(\rho^-, \cdot)), \\ \mathbb{RP}^1 & \text{if } y_i = \tau, x_i = \rho, k_1 = k_2 \text{ and } \sigma(F(\cdot, \tau^-)) \neq \sigma(F(\rho^-, \cdot)), \end{cases} \\
E_4 := & \begin{cases} [F(\lambda, \beta), F(\lambda, \tau)] & \text{if } \det F(\lambda, \cdot) > 0, \\ [F(\lambda, \tau), F(\lambda, \beta)] & \text{if } \det F(\lambda, \cdot) < 0, \\ \{k_i\} & \text{if } x_i = \lambda \text{ and } y_i \notin (\beta, \tau), \\ \mathbb{RP}^1 & \text{if } x_i = \lambda \text{ and } y_i \in (\beta, \tau). \end{cases}
\end{aligned}$$

Then

$$\overline{F[(I \times J) \cap \text{dom}(F)]} = E_1 \cup E_2 \cup E_3 \cup E_4.$$

Furthermore,

$$\begin{aligned}
& \overline{F[(I \times J) \cap \text{dom}(F)]} \setminus F[(I \times J) \cap \text{dom}(F)] \\
& \subseteq F[\{(\lambda, \beta), (\rho, \beta), (\rho, \tau), (\lambda, \tau)\} \cap \text{dom}(F)],
\end{aligned}$$

that is, the only possible extra points that are in $\overline{F[(I \times J) \cap \text{dom}(F)]}$ but not in $F[(I \times J) \cap \text{dom}(F)]$ are the images under F of the four corners of $I \times J$. Finally, $F[(I \times J) \cap \text{dom}(F)]$ is open if I, J are both open, and is closed if I, J are both closed.

Proof. Take increasing sequences $([\lambda_n, \rho_n])_{n=1}^\infty, ([\beta_m, \tau_m])_{m=1}^\infty$ of nontrivial closed cyclic intervals contained in $(\lambda, \rho), (\beta, \tau)$, respectively, such that

1. $(\lambda, \rho) = \bigcup_{n=1}^\infty [\lambda_n, \rho_n], (\beta, \tau) = \bigcup_{m=1}^\infty [\beta_m, \tau_m]$, and
2. none of λ_n 's, ρ_n 's are equal to x_1, x_2 , and none of β_m 's, τ_m 's are equal to y_1, y_2 .

We claim that $E_1 = \lim_{m \rightarrow \infty} \lim_{n \rightarrow \infty} F[[\lambda_n, \rho_n] \times \{\beta_m\}]$ where the limit is with respect to the Hausdorff metric d_H on closed subsets of \mathbb{RP}^1 . Note that we always have

$$F[[\lambda_n, \rho_n] \times \{\beta_m\}] = \begin{cases} [F(\lambda_n, \beta_m), F(\rho_n, \beta_m)] & \text{if } \sigma(F(\cdot, \beta^+)) = +1, \\ [F(\rho_n, \beta_m), F(\lambda_n, \beta_m)] & \text{if } \sigma(F(\cdot, \beta^+)) = -1 \end{cases}$$

for large enough m . Furthermore, since $\text{dom}(F)$ contains an open rectangle of the form $\mathbb{RP}^1 \times (\beta, \beta_{m_0})$ for some large enough m_0 , Lemma 12.22 shows that

$$\lim_{n \rightarrow \infty} F[[\lambda_n, \rho_n] \times \{\beta_m\}] = F[[\lambda, \rho] \times \{\beta_m\}]$$

holds for all large enough m . Using this, we show the claim by considering each of the cases appearing in the definition of E_1 . For simplicity, we assume $(i, j) = (1, 2)$, but the case $(i, j) = (2, 1)$ is not different.

Case I. When $\det F(\cdot, \beta) > 0$ or $\det F(\cdot, \beta) < 0$.

In this case, the claim immediately follows by continuity of F and Lemma 12.22.

Case II. When $y_1 = \beta$ and $x_1 \notin [\lambda, \rho]$.

In this case, we know that there exists an open rectangle R contained in $\text{dom}(F)$ such that R contains $[\lambda, \rho] \times \{\beta\}$ and F is uniformly continuous on R . Since F is constantly equal to k_2 on $[\lambda, \rho] \times \{\beta\}$, it follows that $F[[\lambda, \rho] \times \{\beta_m\}]$ must be contained in an arbitrarily small neighborhood around k_2 whenever m is large enough. Therefore, $E_1 = \{k_2\}$ follows.

Case III. When $y_1 = \beta$ and $x_1 \in (\lambda, \rho)$.

We separate the cases when $k_1 \neq k_2$ and when $k_1 = k_2$. When $k_1 \neq k_2$, since $F(\cdot, \beta_m)$ is not a constant and $F(x_1, \beta_m) = k_1$, we conclude that $F[[\lambda, \rho] \times \{\beta_m\}]$ is a nontrivial closed cyclic interval containing k_1 in its interior. Since both of the endpoints of $F[[\lambda, \rho] \times \{\beta_m\}]$ converge to $F(\lambda, \beta) = F(\rho, \beta) = k_2 \neq k_1$, Lemma 12.22 shows that $F[[\lambda, \rho] \times \{\beta_m\}] \rightarrow \mathbb{RP}^1$ as $m \rightarrow \infty$.

Now, suppose $k_1 = k_2$. We only consider the case $\sigma(F(\cdot, y_1^\pm)) = +1$ as the other case can be done similarly. Fix any $c \in \mathbb{RP}^1 \setminus \{k_2\}$. We will show that $[F(\lambda, \beta_m), c, F(\rho, \beta_m)]$ always holds for large enough m so that we can apply Lemma 12.22 to conclude $F[[\lambda, \rho] \times \{\beta_m\}] \rightarrow \mathbb{RP}^1$ as $m \rightarrow \infty$. Take a small neighborhood (k_-, k_+) of k_2 such that $c \notin [k_-, k_+]$. In particular, $c \notin [k_2, k_+]$, that is, $[k_+, c, k_2]$, and $c \notin [k_-, k_2]$, that is, $[k_2, c, k_-]$.

First, suppose $\sigma(F(x_1^\pm, \cdot)) = +1$. By the assumption, $\det F(\lambda, \cdot) > 0$ holds, so by continuity of $F(\lambda, \cdot)$ at $y_1 = \beta$, we get that $F(\lambda, \cdot)[[\beta, \beta_m]] \subseteq (k_-, k_+)$ holds for any large enough m . Hence, we have $[F(\lambda, \beta), F(\lambda, \beta_m), k_+]$, or equivalently, $[k_2, F(\lambda, \beta_m), k_+]$ for such m , which together with $[k_+, c, k_2]$ shows $[k_2, F(\lambda, \beta_m), k_+, c]$. On the other hand, since $\det F(\cdot, \beta_m) > 0$ holds from the assumption, we have $[F(\lambda, \beta_m), F(x_1, \beta_m), F(\rho, \beta_m)]$, or equivalently, $[k_2, F(\rho, \beta_m), F(\lambda, \beta_m)]$. Hence, combining this with $[k_2, F(\lambda, \beta_m), k_+, c]$ yields

$$[k_2, F(\rho, \beta_m), F(\lambda, \beta_m), k_+, c],$$

in particular $[F(\lambda, \beta_m), c, F(\rho, \beta_m)]$ as claimed, for all large enough m .

Next, suppose $\sigma(F(x_1^\pm, \cdot)) = -1$. By the assumption, $\det F(\rho, \cdot) < 0$ holds, so by continuity of $F(\rho, \cdot)$ at $y_1 = \beta$, we get that $F(\rho, \cdot)[[\beta, \beta_m]] \subseteq (k_-, k_+)$ holds for any large enough m . Hence, we have $[k_-, F(\rho, \beta_m), F(\rho, \beta)]$, or equivalently, $[k_-, F(\rho, \beta_m), k_2]$ for such m , which together with $[k_2, c, k_-]$ shows $[k_-, F(\rho, \beta_m), k_2, c]$. On the other hand, we saw in the case $\sigma(F(x_1^\pm, \cdot)) = +1$ that $[k_2, F(\rho, \beta_m), F(\lambda, \beta_m)]$ holds, which when combined with $[k_-, F(\rho, \beta_m), k_2, c]$ yields

$$[k_2, c, k_-, F(\rho, \beta_m), F(\lambda, \beta_m)],$$

in particular $[F(\lambda, \beta_m), c, F(\rho, \beta_m)]$ as claimed, for all large enough m .

Case IV. When $y_1 = \beta$, $x_1 = \lambda$ and $k_1 \neq k_2$.

In this case, we have

$$F[[\lambda, \rho] \times \{\beta_m\}] = \begin{cases} [k_1, F(\rho, \beta_m)] & \text{if } \sigma(F(\cdot, \beta^+)) = +1, \\ [F(\rho, \beta_m), k_1] & \text{if } \sigma(F(\cdot, \beta^+)) = -1, \end{cases}$$

for m large enough and that $F(\rho, \beta_m) \rightarrow k_2$ as $m \rightarrow \infty$, so Lemma 12.22 shows $F[[\lambda, \rho] \times \{\beta_m\}] \rightarrow [k_1, k_2]$ if $\sigma(F(\cdot, \beta^+)) = +1$ and $F[[\lambda, \rho] \times \{\beta_m\}] \rightarrow [k_2, k_1]$ if $\sigma(F(\cdot, \beta^+)) = -1$, respectively, as $m \rightarrow \infty$.

Case V. When $y_1 = \beta$, $x_1 = \lambda$ and $k_1 = k_2$.

In this case, we have

$$F[[\lambda, \rho] \times \{\beta_m\}] = \begin{cases} [k_2, F(\rho, \beta_m)] & \text{if } \sigma(F(\cdot, \beta^+)) = +1, \\ [F(\rho, \beta_m), k_2] & \text{if } \sigma(F(\cdot, \beta^+)) = -1, \end{cases}$$

for m large enough and that $F(\rho, \beta_m) \rightarrow k_2$ as $m \rightarrow \infty$. Fix any $c \in \mathbb{RP}^1 \setminus \{k_2\}$. We will show that $[F(\lambda, \beta_m), c, F(\rho, \beta_m)]$, or equivalently $[k_2, c, F(\rho, \beta_m)]$ always holds for large enough m if $\sigma(F(\cdot, \beta^+)) = \sigma(F(\lambda^+, \cdot))$, while $[F(\rho, \beta_m), c, k_2]$ always holds for large enough m if $\sigma(F(\cdot, \beta^+)) \neq \sigma(F(\lambda^+, \cdot))$, so that Lemma 12.22 shows $F[[\lambda, \rho] \times \{\beta_m\}] \rightarrow \{k_2\}$ for the first case, and $F[[\lambda, \rho] \times \{\beta_m\}] \rightarrow \mathbb{RP}^1$ for the second case, respectively, as $m \rightarrow \infty$. We only consider the case $\sigma(F(\cdot, y_1^\pm)) = \sigma(F(x_1^\pm, \cdot)) = +1$ and the case $\sigma(F(\cdot, y_1^\pm)) = +1$ and $\sigma(F(x_1^\pm, \cdot)) = -1$, and the other two cases can be treated similarly.

Suppose $\sigma(F(\cdot, y_1^\pm)) = \sigma(F(x_1^\pm, \cdot)) = +1$. Take a small neighborhood (k_-, k_+) of k_2 such that $c \notin [k_-, k_+]$. In particular, $c \notin [k_2, k_+]$, that is, $[k_+, c, k_2]$. As in Case III, the assumption together with the continuity of F gives us $[k_2, F(\rho, \beta_m), k_+]$, which together with $[k_+, c, k_2]$ implies $[k_2, F(\rho, \beta_m), k_+, c]$. In particular, this shows $[F(\rho, \beta_m), c, k_2]$ which is the desired conclusion.

Next, suppose $\sigma(F(\cdot, y_1^\pm)) = +1$ and $\sigma(F(x_1^\pm, \cdot)) = -1$. Take a small neighborhood (k_-, k_+) of k_2 such that $c \notin [k_-, k_+]$. In particular, $c \notin [k_-, k_2]$, that is, $[k_2, c, k_-]$. As in Case III, the assumption together with the continuity of F gives us $[k_-, F(\rho, \beta_m), k_2]$, which together with $[k_2, c, k_-]$ implies $[k_2, c, k_-, F(\rho, \beta_m)]$. In particular, this shows $[k_2, c, F(\rho, \beta_m)]$ which is the desired conclusion.

Case VI. When $y_1 = \beta$ and $x_1 = \rho$.

The same logic from Case III, IV and V applies; details are omitted.

Similarly, we claim that $E_2 = \lim_{m \rightarrow \infty} \lim_{n \rightarrow \infty} F[\{\rho_n\} \times [\beta_m, \tau_m]]$. Again we always have

$$F[\{\rho_n\} \times [\beta_m, \tau_m]] = \begin{cases} [F(\rho_n, \beta_m), F(\rho_n, \tau_m)] & \text{if } \sigma(F(\rho^-, \cdot, \cdot)) = +1, \\ [F(\rho_n, \tau_m), F(\rho_n, \beta_m)] & \text{if } \sigma(F(\rho^-, \cdot, \cdot)) = -1 \end{cases}$$

for large enough m . Again, let $(i, j) = (1, 2)$.

Case I. When $\det F(\rho, \cdot) > 0$ or $\det F(\rho, \cdot) < 0$.

In this case, the claim immediately follows by continuity of F and Lemma 12.22.

Case II. When $x_1 = \rho$ and $y_1 \notin (\beta, \tau)$.

In this case, for fixed m , we know that there exists an open rectangle R contained in $\text{dom}(F)$ such that R contains $\{\rho\} \times [\beta_m, \tau_m]$ and F is uniformly continuous on R . Since F is constantly equal to k_1 on $\{\rho\} \times [\beta_m, \tau_m]$, it follows that $F[\{\rho_n\} \times [\beta_m, \tau_m]]$ must be contained in an arbitrarily small neighborhood around k_1 whenever n is large enough. Therefore, $\lim_{n \rightarrow \infty} F[\{\rho_n\} \times [\beta_m, \tau_m]] = \{k_1\}$, and since this holds for all m , we get $E_2 = \{k_1\}$.

Case III. When $x_1 = \rho$ and $y_1 \in (\beta, \tau)$.

We separate the cases when $k_1 \neq k_2$ and when $k_1 = k_2$. When $k_1 \neq k_2$, since $F(\rho_n, \cdot)$ is not a constant and $F(\rho_n, y_1) = k_2$, we conclude that $F[\{\rho_n\} \times [\beta_m, \tau_m]]$ is a nontrivial closed cyclic interval containing k_2 in its interior. Since both of the endpoints of $F[\{\rho_n\} \times [\beta_m, \tau_m]]$ converge to $F(\rho, \beta_m) = F(\rho, \tau_m) = k_1 \neq k_2$, Lemma 12.22 shows that $F[\{\rho_n\} \times [\beta_m, \tau_m]] \rightarrow \mathbb{RP}^1$ as $n \rightarrow \infty$. Since this holds for all m , we get $E_2 = \mathbb{RP}^1$.

Now, suppose $k_1 = k_2$. We only consider the case $\sigma(F(x_1^\pm, \cdot)) = +1$ as the other case can be done similarly. Fix any $c \in \mathbb{RP}^1 \setminus \{k_1\}$. We will show that $[F(\rho_n, \beta_m), c, F(\rho, \tau_m)]$ always holds for large enough n so that we can apply Lemma 12.22 to conclude $F[\{\rho_n\} \times [\beta_m, \tau_m]] \rightarrow \mathbb{RP}^1$ as $n \rightarrow \infty$, thus $E_2 = \mathbb{RP}^1$ follows as well since m is arbitrary. Take a small neighborhood (k_-, k_+) of k_1 such that $c \notin [k_-, k_+]$. In particular, $c \notin [k_1, k_+]$, that is, $[k_+, c, k_1]$, and $c \notin [k_-, k_1]$, that is, $[k_1, c, k_-]$.

First, suppose $\sigma(F(\cdot, y_1^\pm)) = +1$. By the assumption, $\det F(\cdot, \tau_m) > 0$ holds, so by continuity of $F(\cdot, \tau_m)$ at $x_1 = \rho$, we get that $F(\cdot, \tau_m)[[\rho_n, \rho]] \subseteq (k_-, k_+)$ holds for any large enough n . Hence, we have $[k_-, F(\rho_n, \tau_m), F(\rho, \tau_m)]$, or equivalently, $[k_-, F(\rho_n, \tau_m), k_1]$ for such n , which together with $[k_1, c, k_-]$ shows $[k_-, F(\rho_n, \tau_m), k_1, c]$. On the other hand, since $\det F(\rho_n, \cdot) > 0$ holds from the assumption, we have $[F(\rho_n, \beta_m), F(\rho_n, y_1), F(\rho_n, \tau_m)]$, or equivalently, $[F(\rho_n, \beta_m), k_1, F(\rho_n, \tau_m)]$. Hence, combining this with $[k_-, F(\rho_n, \tau_m), k_1, c]$, we obtain

$$[k_1, c, k_-, F(\rho_n, \tau_m), F(\rho_n, \beta_m)],$$

in particular $[F(\rho_n, \beta_m), c, F(\rho_n, \tau_m)]$ as claimed, for all large enough n .

Next, suppose $\sigma(F(\cdot, y_1^\pm)) = -1$. By the assumption, $\det F(\cdot, \beta_m) < 0$ holds, so by continuity of $F(\cdot, \beta_m)$ at $x_1 = \rho$, we get that $F(\cdot, \beta_m)[[\rho_n, \rho]] \subseteq (k_-, k_+)$ holds for any large enough n . Hence, we have $[F(\rho, \beta_m), F(\rho_n, \beta_m), k_+]$, or equivalently, $[k_1, F(\rho_n, \beta_m), k_+]$ for such n , which together with $[k_+, c, k_1]$ shows $[k_1, F(\rho_n, \beta_m), k_+, c]$. On the other hand, we saw in the case $\sigma(F(\cdot, y_1^\pm)) = +1$ that $[F(\rho_n, \beta_m), k_1, F(\rho_n, \tau_m)]$, which when combined with $[k_1, F(\rho_n, \beta_m), k_+, c]$ yields

$$[k_1, F(\rho_n, \tau_m), F(\rho_n, \beta_m), k_+, c],$$

in particular $[F(\rho_n, \beta_m), c, F(\rho_n, \tau_m)]$ as claimed, for all large enough n .

Similarly, we can show

$$\begin{aligned} E_3 &= \lim_{m \rightarrow \infty} \lim_{n \rightarrow \infty} F[[\lambda_n, \rho_n] \times \{\tau_m\}], \\ E_4 &= \lim_{m \rightarrow \infty} \lim_{n \rightarrow \infty} F[\{\lambda_n\} \times [\beta_m, \tau_m]]. \end{aligned}$$

Recall that finite unions are continuous with respect to the Hausdorff metric¹³, thus we conclude

$$\begin{aligned} E_1 \cup E_2 \cup E_3 \cup E_4 &= \lim_{m \rightarrow \infty} \lim_{n \rightarrow \infty} \left(F[[\lambda_n, \rho_n] \times \{\beta_m\}] \cup F[\{\rho_n\} \times [\beta_m, \tau_m]] \right. \\ &\quad \left. \cup F[[\lambda_n, \rho_n] \times \{\tau_m\}] \cup F[\{\lambda_n\} \times [\beta_m, \tau_m]] \right). \end{aligned}$$

¹³Because $d_H(A_1 \cup A_2, B_1 \cup B_2) \leq \max\{d_H(A_1, B_1), d_H(A_2, B_2)\}$.

In particular, when $(x_1, y_1), (x_2, y_2)$ are not in the interior of $I \times J$, then Proposition 12.23 shows

$$\begin{aligned}
E_1 \cup E_2 \cup E_3 \cup E_4 &= \lim_{m \rightarrow \infty} \lim_{n \rightarrow \infty} F[[\lambda_n, \rho_n] \times [\beta_m, \tau_m]] \\
&= F \left[\overline{\bigcup_{m=1}^{\infty} \bigcup_{n=1}^{\infty} ([\lambda_n, \rho_n] \times [\beta_m, \tau_m])} \right] \\
&= \overline{F[\text{int}(I \times J)]} \\
&= \overline{F[(I \times J) \cap \text{dom}(F)]}.
\end{aligned}$$

because in Hausdorff metric any increasing sequence converges to the closure of its union.

Next, we claim that if the interior of $I \times J$ intersects with the indeterminacy locus, then both $E_1 \cup E_2 \cup E_3 \cup E_4$ and $\overline{F[(I \times J) \cap \text{dom}(F)]}$ are equal to \mathbb{RP}^1 . Clearly, Proposition 12.15 shows that the latter is \mathbb{RP}^1 , so it is enough to show that $E_1 \cup E_2 \cup E_3 \cup E_4 = \mathbb{RP}^1$.

Suppose that we know this claim is true when there is no point in the indeterminacy locus that is located on the boundary of $I \times J$. Then, if $(x_1, y_1) \neq (x_2, y_2)$, one of them is located in the interior of $I \times J$ and the other is on the boundary, then we still have $E_1 \cup E_2 \cup E_3 \cup E_4 = \mathbb{RP}^1$. Indeed, we already know that

$$\begin{aligned}
E_1 \cup E_2 \cup E_3 \cup E_4 &= \lim_{m \rightarrow \infty} \lim_{n \rightarrow \infty} \left(F[[\lambda_n, \rho_n] \times \{\beta_m\}] \cup F[\{\rho_n\} \times [\beta_m, \tau_m]] \right. \\
&\quad \left. \cup F[[\lambda_n, \rho_n] \times \{\tau_m\}] \cup F[\{\lambda_n\} \times [\beta_m, \tau_m]] \right)
\end{aligned}$$

holds, but by the claim we supposed to be true, the right-hand side is \mathbb{RP}^1 . Therefore, there are three different cases remaining to be proved: (1) when $(x_1, y_1) = (x_2, y_2)$ and it is in the interior of $I \times J$, (2) when $(x_1, y_1) \neq (x_2, y_2)$, (x_1, y_1) is in the interior of $I \times J$ and (x_2, y_2) is in the exterior of $I \times J$, and (3) when $(x_1, y_1) \neq (x_2, y_2)$ and both (x_1, y_1) and (x_2, y_2) are in the interior of $I \times J$. For all three cases, E_1, E_2, E_3, E_4 are just images of four edges of $I \times J$ under F .

When $(x_1, y_1) = (x_2, y_2)$ and it is in the interior of $I \times J$, Proposition 12.15 shows that at least two of $E_1 \cup E_2$, $E_2 \cup E_3$, $E_3 \cup E_4$ and $E_4 \cup E_1$ are \mathbb{RP}^1 . For the other two cases, Proposition 12.15 shows that any of $E_1 \cup E_2 \cup E_3$, $E_2 \cup E_3 \cup E_4$, $E_3 \cup E_4 \cup E_1$ and $E_4 \cup E_1 \cup E_2$ is \mathbb{RP}^1 . Therefore, the claim is proved.

Since F is an open map by Proposition 12.16, we know $F[(I \times J) \cap \text{dom}(F)]$ is open if I, J are both open. Next, we show that $F[(I \times J) \cap \text{dom}(F)]$ is closed if I, J are both closed. The conclusion is trivial if $I \times J \subseteq \text{dom}(F)$ by compactness, so suppose otherwise. If some (x_i, y_i) is in the interior of $I \times J$, then we know from Proposition 12.15 that $F[(I \times J) \cap \text{dom}(F)] = \mathbb{RP}^1$, so is closed. If some (x_i, y_i) is on the boundary of $I \times J$ but not on one of the four corners of $I \times J$, then still Proposition 12.15 shows $F[(I \times J) \cap \text{dom}(F)] = \mathbb{RP}^1$, so is closed. If some (x_i, y_i) is on one of the four corners of $I \times J$, then Proposition 12.15 shows that the image of $(I \times J) \cap R \cap \text{dom}(F)$ by F is closed, where R is a small enough closed cyclic rectangle containing that corner in its interior. To see why, the only nontrivial case is when $k_1 = k_2$ and $F[(I \times J) \cap R \cap \text{dom}(F)]$ is either given as a union of intervals of the form $[F(u, v), k_1]$, or that of the form $[k_1, F(u, v)]$, for $(u, v) \in (I \times J) \cap R$. However, in that case, by inspecting the determinant signs of F along the relevant edges, one can easily see that all such $[F(u, v), k_1]$ or $[k_1, F(u, v)]$ are contained in a single $[F(u, v), k_1]$ or $[k_1, F(u, v)]$, namely the one we obtain when (u, v) is the corner of R located in the interior of $I \times J$; details are omitted. Hence,

$$F[(I \times J) \cap \text{dom}(F)] = F[(I \times J) \setminus \text{int } R] \cup F[(I \times J) \cap R \cap \text{dom}(F)]$$

is closed.

Finally, we show

$$\begin{aligned} C &:= \overline{F[(I \times J) \cap \text{dom}(F)]} \setminus F[(I \times J) \cap \text{dom}(F)] \\ &\subseteq F[\{(\lambda, \beta), (\rho, \beta), (\rho, \tau), (\lambda, \tau)\} \cap \text{dom}(F)]. \end{aligned}$$

By continuity, clearly $F[\text{int}(I \times J) \cap \text{dom}(F)]$ is dense in $F[(\overline{I \times J}) \cap \text{dom}(F)]$ and we already showed that the latter is closed, so it must coincide with $\overline{F[(I \times J) \cap \text{dom}(F)]}$. Hence, any point in C must be the image of a point on the boundary of $I \times J$. Suppose that $(u, v) \in (\overline{I \times J} \setminus \text{int}(I \times J)) \cap \text{dom}(F)$ be such that (u, v) is not one of the corners of $I \times J$ and $(s : t) := F(u, v) \in C$. Then by Proposition 12.10, we have

$$\begin{aligned} F^{-1}[(s : t)] &= \{([x], [(tA - sB)^{-1}Rx]) \in \text{dom}(F) : [x] \in \mathbb{RP}^1\} \\ &= \{([(tA^T - sB^T)^{-1}Ry], [y]) \in \text{dom}(F) : [y] \in \mathbb{RP}^1\} \end{aligned}$$

when $tA - sB$ is invertible, and

$$F^{-1}[(s : t)] = ((\{[x_0]\} \times \mathbb{RP}^1) \cup (\mathbb{RP}^1 \times \{[y_0]\})) \cap \text{dom}(F)$$

otherwise, where $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $y_0 \in \ker(tA - sB) \setminus \{0\}$. For the first case, $F^{-1}[(s : t)]$ must intersect $\text{int}(I \times J)$ since $[x] \mapsto [(tA - sB)^{-1}Rx]$ and $[y] \mapsto [(tA^T - sB^T)^{-1}Ry]$ are continuous functions and (u, v) is in the interior of one of the four edges of $I \times J$.

On the other hand, for the second case, obviously either $F^{-1}[(s : t)]$ intersects $\text{int}(I \times J)$ or it contains the entire edge of $I \times J$ where (u, v) belongs to, in which case $(s : t)$ is equal to the image of F at one of the corners of $I \times J$. This completes the proof of $C \subseteq F[\{(\lambda, \beta), (\rho, \beta), (\rho, \tau), (\lambda, \tau)\} \cap \text{dom}(F)]$. \square

When I, J are not both closed, then we have to figure out the exact set

$$\begin{aligned} C &:= \overline{F[(I \times J) \cap \text{dom}(F)]} \setminus F[(I \times J) \cap \text{dom}(F)] \\ &\subseteq F[\{(\lambda, \beta), (\rho, \beta), (\rho, \tau), (\lambda, \tau)\} \cap \text{dom}(F)]. \end{aligned}$$

For instance, suppose $(\lambda, \beta) \in \text{dom}(F)$ and we want to see if $(s : t) := F(\lambda, \beta)$ belongs to $F[(I \times J) \cap \text{dom}(F)]$. Again, Proposition 12.10 shows that $F^{-1}[(s : t)]$ is given as

$$\begin{aligned} F^{-1}[(s : t)] &= \{([x], [(tA - sB)^{-1}Rx]) \in \text{dom}(F) : [x] \in \mathbb{RP}^1\} \\ &= \{([(tA^T - sB^T)^{-1}Ry], [y]) \in \text{dom}(F) : [y] \in \mathbb{RP}^1\} \end{aligned}$$

when $tA - sB$ is invertible, and

$$F^{-1}[(s : t)] = ((\{[x_0]\} \times \mathbb{RP}^1) \cup (\mathbb{RP}^1 \times \{[y_0]\})) \cap \text{dom}(F)$$

otherwise, where $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $y_0 \in \ker(tA - sB) \setminus \{0\}$. Clearly, for the second case $I \times J$ intersects with $F^{-1}[(s : t)]$ if and only if either $[x_0] \in I$ or $[y_0] \in J$ holds.

For the first case, note that the linear fractional mapping

$$T_{(tA-sB)^{-1}R} : [x] \mapsto [(tA - sB)^{-1}Rx]$$

should satisfy $\beta = T_{(tA-sB)^{-1}R}(\lambda)$ and it maps $[\lambda, \rho]$ onto either $[\beta, T_{(tA-sB)^{-1}R}(\rho)]$ or $[T_{(tA-sB)^{-1}R}(\rho), \beta]$, where which one is the case is determined by the determinant sign of $(tA - sB)^{-1}R$, or equivalently, that of $tA - sB$. If the sign is positive, then it follows that

$F^{-1}[(s : t)]$ must intersect with the interior of $I \times J$. If the sign is negative, then $F^{-1}[(s : t)]$ intersects with the interior of $I \times J$ if and only if $[\beta, T_{(tA-sB)^{-1}R}(\rho), \tau]$ holds, and it does not intersect the interior but touches the corner (ρ, τ) if and only if $T_{(tA-sB)^{-1}R}(\rho) = \tau$. Hence, if either the sign is positive or is negative and $[\beta, T_{(tA-sB)^{-1}R}(\rho), \tau]$ holds, then we conclude $(s : t) \notin C$, and otherwise we conclude $(s : t) \notin C$ if and only if either $(\lambda, \beta) \in I \times J$ holds or $T_{(tA-sB)^{-1}R}(\rho) = \tau$ (or equivalently $\rho = T_{R(tA-sB)}(\tau)$) and $(\rho, \tau) \in (I \times J) \cap \text{dom}(F)$ hold. And we can apply a similar check for other corners (ρ, β) , (ρ, τ) and (λ, τ) as well.

Now, the final missing piece is how to compute the union $E_1 \cup E_2 \cup E_3 \cup E_4$. This is in fact simple. First, we compute E_1 . If $E_1 = \mathbb{RP}^1$, then there is no need to do any further computation because the union is just \mathbb{RP}^1 .

Otherwise, we compute E_2 . Again, we may suppose $E_2 \neq \mathbb{RP}^1$. By how they are defined, E_2 must “start” at where E_1 “ended”. This is the right endpoint of E_1 if $\sigma(F(\cdot, \beta^+)) = +1$, or is the left endpoint of E_1 if $\sigma(F(\cdot, \beta^+)) = -1$, and similarly it is the left endpoint of E_2 if $\det F(\rho, \cdot) \geq 0$, or is the right endpoint of E_2 if $\det F(\rho, \cdot) \leq 0$. When the right endpoint of E_1 and the left endpoint of E_2 are the matching endpoints, then $E_1 \cup E_2$ is either \mathbb{RP}^1 or the closed cyclic interval bounded between the left endpoint of E_1 and the right endpoint of E_2 , where the first case happens if and only if the right endpoint of E_2 is in E_1 .

On the other hand, if the right endpoint of E_1 and the right endpoint of E_2 are the matching endpoints, then $E_1 \cup E_2$ is either E_1 or E_2 , where the first case happens if and only if the left endpoint of E_2 is in E_1 . The other two cases can be treated similarly.

If $E_1 \cup E_2 \neq \mathbb{RP}^1$, then we compute E_3 . Again, suppose $E_3 \neq \mathbb{RP}^1$. Now the other endpoint of E_2 that was not the matching endpoint with E_1 is to be matched with an endpoint of E_3 , the right endpoint if $\sigma(F(\cdot, \tau^-)) = +1$, or the left endpoint if $\sigma(F(\cdot, \tau^-)) = -1$.

Similarly as before, if the E_2 -side matching endpoint is the right endpoint of $E_1 \cup E_2$ and the E_3 -side matching endpoint is the left endpoint of E_3 , then we simply stitch them together in a row, which will either result in \mathbb{RP}^1 when the right endpoint of E_3 is in $E_1 \cup E_2$, or the closed cyclic interval bounded between the left endpoint of $E_1 \cup E_2$ and the right endpoint of E_3 .

If the matching is between the right endpoint of $E_1 \cup E_2$ and the right endpoint of E_3 , then the union is either $E_1 \cup E_2$ or E_3 , where the first case happens if and only if the left endpoint of E_3 is in $E_1 \cup E_2$.

If the matching is between an interior point of $E_1 \cup E_2$ and the left endpoint of E_3 , then there are three possible cyclic orders between endpoints of $E_1 \cup E_2$ and E_3 when they are all distinct, each yielding $E_1 \cup E_2$, the closed cyclic interval bounded between the left endpoint of $E_1 \cup E_2$ and the right endpoint of E_3 , and \mathbb{RP}^1 , respectively. Or, the right endpoint of E_3 may coincide with either the left or the right endpoint of $E_1 \cup E_2$, each yielding \mathbb{RP}^1 and $E_1 \cup E_2$.

More formally, this *cyclic interval appending algorithm* can be described as follows:

Proposition 12.26 (Cyclic interval appending algorithm).

Let X be a cyclically ordered set and I a closed cyclic interval in X . Suppose I is a proper subset of X , and let l, r be the left and the right endpoints of I . (Here, I can be a singleton set, in which case $l = r$.) Then for $c \in I$ and $c' \in X \setminus \{c\}$, we have

$$I \cup [c, c'] = \begin{cases} [l, c'] & \text{if } l = r \text{ or } [r, c', l], \\ X & \text{if } l \neq r, \text{ not } [r, c', l], \text{ and } r = c \text{ or } [c, r, c'], \\ I & \text{otherwise.} \end{cases}$$

Similarly,

$$I \cup [c', c] = \begin{cases} [c', r] & \text{if } l = r \text{ or } [r, c', l], \\ X & \text{if } l \neq r, \text{ not } [r, c', l], \text{ and } l = c \text{ or } [c', l, c], \\ I & \text{otherwise.} \end{cases}$$

Proof. Clearly, if $l = r$ then we have $I \cup [c, c'] = [l, c']$ and $I \cup [c', c] = [c', r]$. Next, suppose $[r, c', l]$. Then we have three cases: $l = c$, $[l, c, r]$, $r = c$. For the first case, we have $I \cup [c, c'] = [l, c']$ and $I \cup [c', c] = [c', r]$, and for the second case, we have $[l, c, r, c']$ or equivalently $[c', l, c, r]$, so $I \cup [c, c'] = [l, c']$ and $I \cup [c', c] = [c', r]$, and for the third case, we have $I \cup [c, c'] = [l, c']$ and $I \cup [c', c] = [c', r]$.

Next, suppose $l \neq r$ and not $[r, c', l]$, that is, either $r = c'$ or $l = c'$ or $[l, c', r]$. For the first case $r = c'$, neither $r = c$ nor $[c, r, c']$ is possible, while both $l = c$ and $[c', l, c]$ yield $I \cup [c', c] = X$. For the second case $l = c'$, neither $l = c$ nor $[c', l, c]$ is possible, while both $r = c$ and $[c, r, c']$ yield $I \cup [c, c'] = X$. For the third case $[l, c', r]$, $r = c$ yields $I \cup [c, c'] = X$ and $l = c$ yields $I \cup [c', c] = X$. On the other hand, $[r, l, c']$ and $[c', c, r]$ together imply $[r, l, c', c]$ or equivalently, $[l, c', c, r]$, which yields $I \cup [c, c'] = X$. Similarly, $[l, c, c']$ and $[c', r, l]$ together imply $[l, c, c', r]$, which yields $I \cup [c', c] = X$.

Next, suppose $l \neq r$ and not $[r, c', l]$ and $r \neq c$ and not $[c, r, c']$. In other words, we assume $l \neq r$, $r \neq c$, and either $r = c'$ or $l = c'$ or $[l, c', r]$, and either $r = c'$ or $[c', r, c]$. Thus we have three cases: $r = c'$, or $l = c'$ and $[c', r, c]$, or $[l, c', r]$ and $[c', r, c]$. For the first case, clearly we have $I \cup [c, c'] = I$. The second case is impossible because $c \in [l, r]$ and $c \neq c'$ imply $c \in (l, r)$, that is, $[l, c, r]$ which contradicts to $[c', r, c]$. For the third case, since $c \in [l, r]$, one of $c = l$ and $[l, c, c', r]$ should hold. Then it is clear that both possibilities lead to $I \cup [c, c'] = I$.

Finally, suppose $l \neq r$ and not $[r, c', l]$ and $l \neq c$ and not $[c', l, c]$. In other words, we assume $l \neq r$, $l \neq c$, and either $r = c'$ or $l = c'$ or $[l, c', r]$, and either $l = c'$ or $[c, l, c']$. Thus we have three cases: $l = c'$, or $r = c'$ and $[c, l, c']$, or $[l, c', r]$ and $[c, l, c']$. For the first case, clearly we have $I \cup [c', c] = I$. The second case is impossible because $c \in [l, r]$ and $c \neq c'$ imply $c \in (l, r)$, that is, $[l, c, r]$ which contradicts to $[c, l, c']$. For the third case, since $c \in [l, r]$, one of $c = r$ and $[l, c', c, r]$ should hold. Then it is clear that both possibilities lead to $I \cup [c', c] = I$. \square

Combining all of these, we finally obtain an algorithm for computing $F[(I \times J) \cap \text{dom}(F)]$:

Algorithm 12.27 (Cyclic rectangle mapping algorithm).

Given $A, B \in \mathbb{R}^{2 \times 2}$ not both zero such that $\bar{F} = F$ and nonempty cyclic intervals $I, J \subseteq \mathbb{RP}^1$, define $L_1, L_2: \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$ as

$$\begin{aligned} L_1: (x, y) &\mapsto \langle x, (B^T R A - A^T R B)y \rangle, \\ L_2: (x, y) &\mapsto \langle x, (B R A^T - A R B^T)y \rangle. \end{aligned}$$

Then the following procedure is supposed to return $F[(I \times J) \cap \text{dom}(F)]$.

1. When $J = \{[y_0]\}$, compute $L_1(y_0, y_0)$.
 - (a) If $L_1(y_0, y_0) \neq 0$, then compute $F(\cdot, [y_0])[I]$ using Proposition 12.3 and return it.
 - (b) If $L_1(y_0, y_0) = 0$, then compute the unique $(s : t) \in \mathbb{RP}^1$ such that $tAy_0 = sBy_0$, and let $x_0 \leftarrow RAy_0$ if $s \neq 0$, and $x_0 \leftarrow RB y_0$ if $t \neq 0$.
 - If $I = \{[x_0]\}$, then return \emptyset .
 - Otherwise, return $\{(s : t)\}$.

2. When $I = \{[x_0]\}$ and J is not a singleton set, compute $\overline{F([y_0], \cdot)}[J]$ using Proposition 12.3 and return it.
3. When both I and J are not singleton sets and one of them is \mathbb{RP}^1 , just return \mathbb{RP}^1 .
4. When $J = \mathbb{RP}^1 \setminus \{[y_0]\}$ and I is not a singleton set nor \mathbb{RP}^1 , compute $L_1(y_0, y_0)$.
 - (a) If $L_1(y_0, y_0) \neq 0$, then return \mathbb{RP}^1 .
 - (b) If $L_1(y_0, y_0) = 0$, then compute the unique $(s : t) \in \mathbb{RP}^1$ such that $tAy_0 = sBy_0$, and compute the unique element $[x_0] \in \mathbb{RP}^1$ such that $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$.
 - If I contains $[x_0]$, then return \mathbb{RP}^1 .
 - Otherwise, return $\mathbb{RP}^1 \setminus \{(s : t)\}$.
5. When $I = \mathbb{RP}^1 \setminus \{[x_0]\}$ and J is not a singleton set nor \mathbb{RP}^1 nor the complement of a singleton set, compute $L_2(x_0, x_0)$.
 - (a) If $L_2(x_0, x_0) \neq 0$, then return \mathbb{RP}^1 .
 - (b) If $L_2(x_0, x_0) = 0$, then compute the unique $(s : t) \in \mathbb{RP}^1$ such that $tA^Tx_0 = sB^Tx_0$, and compute the unique element $[y_0] \in \mathbb{RP}^1$ such that $y_0 \in \ker(tA - sB) \setminus \{0\}$.
 - i. If J contains $[y_0]$, then return \mathbb{RP}^1 .
 - ii. Otherwise, return $\mathbb{RP}^1 \setminus \{(s : t)\}$.
6. Otherwise, there exist $[\lambda], [\rho], [\beta], [\tau] \in \mathbb{RP}^1$ such that $[\lambda] \neq [\rho], [\beta] \neq [\tau]$ and $([\lambda], [\rho]) \subseteq I \subseteq [[\lambda], [\rho]], ([\beta], [\tau]) \subseteq J \subseteq [[\beta], [\tau]]$.
 - (a) Compute $c_{\text{start}}, c_{\text{end}}$ and E , the bottom edge, as follows.
 - i. Compute $L_1(\beta, \beta)$. If it is nonzero, then set $c_{\text{start}} \leftarrow F([\lambda], [\beta])$ and $c_{\text{end}} \leftarrow F([\rho], [\beta])$, and in this case set $E \leftarrow [c_{\text{start}}, c_{\text{end}}]$ if $L_1(\beta, \beta) > 0$, and set $E \leftarrow [c_{\text{end}}, c_{\text{start}}]$ otherwise. Proceed to 6(b).
 - ii. If $L_1(\beta, \beta) = 0$, compute $k_1 := (s_1 : t_1)$, $k_2 := (s_2 : t_2)$ and $x_0 \in \mathbb{R}^2 \setminus \{0\}$ as follows.
 - Compute the unique element $(s_2 : t_2) \in \mathbb{RP}^1$ such that $t_2A\beta = s_2B\beta$. If $s_2 \neq 0$, let $x_0 \leftarrow RA\beta$, and otherwise, $x_0 \leftarrow RB\beta$. Then find the unique element $(s_1 : t_1) \in \mathbb{RP}^1$ such that $t_1A^Tx_0 = s_1B^Tx_0$.
 - iii. If $[[\lambda], [x_0], [\rho]]$, then set $E \leftarrow \mathbb{RP}^1$ and proceed to 6(e).
 - iv. Otherwise, if $[[\rho], [x_0], [\lambda]]$, then set $c_{\text{start}} \leftarrow k_2$, $c_{\text{end}} \leftarrow k_2$ and $E \leftarrow \{k_2\}$. Proceed to 6(b).
 - v. Otherwise, check if $k_1 = k_2$ holds.
 - When $k_1 \neq k_2$, set $c_{\text{start}} \leftarrow k_1$ and $c_{\text{end}} \leftarrow k_2$ if $[x_0] = [\lambda]$, or set $c_{\text{start}} \leftarrow k_2$ and $c_{\text{end}} \leftarrow k_1$ otherwise, and then compute the sign of $L_1(\beta, R\beta)$.
 - If it is positive, then set $E \leftarrow [c_{\text{end}}, c_{\text{start}}]$.
 - If it is negative, then set $E \leftarrow [c_{\text{start}}, c_{\text{end}}]$.
 - When $k_1 = k_2$, compare the signs of $L_1(R\beta, R\beta)$ and $L_2(Rx_0, Rx_0)$.
 - If they are different and $[x_0] = [\lambda]$, or if they are same and $[x_0] = [\rho]$, then set $E \leftarrow \mathbb{RP}^1$ and proceed to 6(e).
 - Otherwise, set $c_{\text{start}} \leftarrow k_2$, $c_{\text{end}} \leftarrow k_2$ and $E \leftarrow \{k_2\}$. Proceed to 6(b).

- (b) Update c_{end} and E by appending the right edge, as described as follows.
- i. Compute $L_2(\rho, \rho)$. If $L_2(\rho, \rho) = 0$, compute $k_1 := (s_1 : t_1)$ and $y_0 \in \mathbb{R}^2 \setminus \{0\}$ as follows.
 - Compute the unique element $(s_1 : t_1) \in \mathbb{RP}^1$ such that $t_1 A^T \rho = s_1 B^T \rho$. If $s_1 \neq 0$, let $y_0 \leftarrow R A^T \rho$, and otherwise, $y_0 \leftarrow R B^T \rho$.
 If $[[\beta], [y_0], [\tau]]$, then set $E \leftarrow \mathbb{RP}^1$ and proceed to 6(e). Otherwise, do not update anything and just proceed to 6(c).
 - ii. If $L_2(\rho, \rho) \neq 0$, then set $c_{\text{new}} \leftarrow F([\rho], [\tau])$.
 - If $L_2(\rho, \rho) > 0$, then use Proposition 12.26 to replace E by $E \cup [c_{\text{end}}, c_{\text{new}}]$.
 - If $L_2(\rho, \rho) < 0$, then use Proposition 12.26 to replace E by $E \cup [c_{\text{new}}, c_{\text{end}}]$.
 If $E = \mathbb{RP}^1$, then proceed to 6(e).
- (c) Update c_{end} and E by appending the top edge, as described as follows.
- i. Compute $L_1(\tau, \tau)$. If $L_1(\tau, \tau) \neq 0$, then set $c_{\text{new}} \leftarrow F([\lambda], [\tau])$ and set σ to be the sign of $L_1(\tau, \tau)$. Then, proceed to 6(c)-vi.
 - ii. If $L_1(\tau, \tau) = 0$, compute $k_1 := (s_1 : t_1)$, $k_2 := (s_2 : t_2)$ and $x_0 \in \mathbb{R}^2 \setminus \{0\}$ as done in 6(a)-ii, except that we use τ instead of β .
 - iii. If $[[\lambda], [x_0], [\rho]]$, then set $E \leftarrow \mathbb{RP}^1$ and proceed to 6(e).
 - iv. Otherwise, if $[[\rho], [x_0], [\lambda]]$, then do not update anything and just proceed to 6(d).
 - v. Otherwise, check if $k_1 = k_2$ holds.
 - When $k_1 \neq k_2$, set $c_{\text{new}} \leftarrow k_2$ if $[x_0] = [\rho]$, or set $c_{\text{new}} \leftarrow k_1$ otherwise, and then set σ to be the sign opposite to the sign of $L_1(\tau, R\tau)$. Then, proceed to 6(c)-vi.
 - When $k_1 = k_2$, compare the signs of $L_1(R\tau, R\tau)$ and $L_2(Rx_0, Rx_0)$.
 - If they are different and $[x_0] = [\rho]$, or if they are same and $[x_0] = [\lambda]$, then set $E \leftarrow \mathbb{RP}^1$ and proceed to 6(e).
 - Otherwise, do not update anything and just proceed to 6(d).
 - vi. Append a new cyclic interval into E as follows.
 - If $\sigma < 0$, then use Proposition 12.26 to replace E by $E \cup [c_{\text{end}}, c_{\text{new}}]$.
 - If $\sigma > 0$, then use Proposition 12.26 to replace E by $E \cup [c_{\text{new}}, c_{\text{end}}]$.
 If $E = \mathbb{RP}^1$, then proceed to 6(e).
- (d) Update c_{end} and E by appending the left edge, as described as follows.
- i. Compute $L_2(\lambda, \lambda)$. If $L_2(\lambda, \lambda) = 0$, compute $k_1 := (s_1 : t_1)$ and $y_0 \in \mathbb{R}^2 \setminus \{0\}$ as in 6(b)-i. If $[[\beta], [y_0], [\tau]]$, then set $E \leftarrow \mathbb{RP}^1$ and proceed to 6(e). Otherwise, just proceed to 6(e).
 - ii. If $L_2(\lambda, \lambda) < 0$, then use Proposition 12.26 to compute $E \cup [c_{\text{end}}, c_{\text{start}}]$ and proceed to 6(e).
 - iii. If $L_2(\lambda, \lambda) > 0$, then use Proposition 12.26 to compute $E \cup [c_{\text{start}}, c_{\text{end}}]$ and proceed to 6(e).
- (e) Remove boundary points from E that are not included in $F[(I \times J) \cap \text{dom}(F)]$ and then return. If E is a singleton set, we can skip this and return immediately. If $E = \mathbb{RP}^1$, there are at most one point that might be removed, so we can return immediately if one of the corners is detected to be not in $F[(I \times J) \cap \text{dom}(F)]$. If $E = [l, r]$ for some $l \neq r$, then a corner needs to be inspected only if it is in the domain of F and its image is one of l and r . In any case, the inspection is only needed if the corner does not belong to $I \times J$.

- i. Check if $F([\lambda], [\beta])$ belongs to $F[(I \times J) \cap \text{dom}(F)]$ if the condition described above is satisfied.
 - Find the unique element $(s : t) \in \mathbb{RP}^1$ such that $F([\lambda], [\beta]) = (s : t)$, and check if $tA - sB$ is invertible or not.
 - If $tA - sB$ is not invertible, check if either $[x_0] \in I$ or $[y_0] \in J$ holds, where $[x_0], [y_0]$ are unique elements in \mathbb{RP}^1 satisfying $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $y_0 \in \ker(tA - sB) \setminus \{0\}$. If those two both fail to hold, then remove $(s : t)$ from E .
 - If $tA - sB$ is invertible, compute its determinant. If the determinant is negative, check if either $[[\beta], [(tA - sB)^{-1}R\rho], [\tau]]$ holds or $([\lambda], [\beta]) \in I \times J$ holds or $[(tA - sB)^{-1}R\rho] = [\tau]$ and $([\rho], [\tau]) \in (I \times J) \cap \text{dom}(F)$ hold. If those three all fail to hold, then remove $(s : t)$ from E .
- ii. Check if $F([\rho], [\beta])$ belongs to $F[(I \times J) \cap \text{dom}(F)]$ if the condition described above is satisfied.
 - Find the unique element $(s : t) \in \mathbb{RP}^1$ such that $F([\rho], [\beta]) = (s : t)$, and check if $tA - sB$ is invertible or not.
 - If $tA - sB$ is not invertible, check if either $[x_0] \in I$ or $[y_0] \in J$ holds, where $[x_0], [y_0]$ are unique elements in \mathbb{RP}^1 satisfying $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $y_0 \in \ker(tA - sB) \setminus \{0\}$. If those two both fail to hold, then remove $(s : t)$ from E .
 - If $tA - sB$ is invertible, compute its determinant. If the determinant is positive, check if either $[[\beta], [(tA - sB)^{-1}R\lambda], [\tau]]$ holds or $([\rho], [\beta]) \in I \times J$ holds or $[(tA - sB)^{-1}R\lambda] = [\tau]$ and $([\lambda], [\tau]) \in (I \times J) \cap \text{dom}(F)$ hold. If those three all fail to hold, then remove $(s : t)$ from E .
- iii. Check if $F([\rho], [\tau])$ belongs to $F[(I \times J) \cap \text{dom}(F)]$ if the condition described above is satisfied.
 - Find the unique element $(s : t) \in \mathbb{RP}^1$ such that $F([\rho], [\tau]) = (s : t)$, and check if $tA - sB$ is invertible or not.
 - If $tA - sB$ is not invertible, check if either $[x_0] \in I$ or $[y_0] \in J$ holds, where $[x_0], [y_0]$ are unique elements in \mathbb{RP}^1 satisfying $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $y_0 \in \ker(tA - sB) \setminus \{0\}$. If those two both fail to hold, then remove $(s : t)$ from E .
 - If $tA - sB$ is invertible, compute its determinant. If the determinant is negative, check if either $[[\beta], [(tA - sB)^{-1}R\lambda], [\tau]]$ holds or $([\rho], [\tau]) \in I \times J$ holds or $[(tA - sB)^{-1}R\lambda] = [\beta]$ and $([\lambda], [\beta]) \in (I \times J) \cap \text{dom}(F)$ hold. If those three all fail to hold, then remove $(s : t)$ from E .
- iv. Check if $F([\lambda], [\tau])$ belongs to $F[(I \times J) \cap \text{dom}(F)]$ if the condition described above is satisfied.
 - Find the unique element $(s : t) \in \mathbb{RP}^1$ such that $F([\lambda], [\tau]) = (s : t)$, and check if $tA - sB$ is invertible or not.
 - If $tA - sB$ is not invertible, check if either $[x_0] \in I$ or $[y_0] \in J$ holds, where $[x_0], [y_0]$ are unique elements in \mathbb{RP}^1 satisfying $x_0 \in \ker(tA^T - sB^T) \setminus \{0\}$ and $y_0 \in \ker(tA - sB) \setminus \{0\}$. If those two both fail to hold, then remove $(s : t)$ from E .
 - If $tA - sB$ is invertible, compute its determinant. If the determinant is positive, check if either $[[\beta], [(tA - sB)^{-1}R\rho], [\tau]]$ holds or $([\lambda], [\tau]) \in I \times J$ holds or $[(tA - sB)^{-1}R\rho] = [\beta]$ and $([\rho], [\beta]) \in (I \times J) \cap \text{dom}(F)$ hold. If those three all fail to hold, then remove $(s : t)$ from E .

Given we know how to compute $\overline{F}[(I \times J) \cap \text{dom}(\overline{F})]$, computing the continued fraction expansion of $\overline{F}(x, y)$ is now just a breeze: we just repeat the exactly the same procedure as we did for the unary case.

However, we need to decide what we would do if (x, y) is in the indeterminacy locus of F , because it does not make sense to compute the continued fraction expansion, or really just anything about $\overline{F}(x, y)$ when it is not even well-defined. There are several options we can choose from. Suppose that we implement the algorithm for generating the continued fraction expansion of $\overline{F}(x, y)$ using the cyclic rectangle mapping algorithm described above.

First, we could just assume (x, y) is never in the indeterminacy locus. Or, in other words, it is the caller's responsibility to provide a valid (x, y) , and the callee does not care about the validity of the answer when the input was wrong from the first place. In this case, what would happen if the input was indeed in the indeterminacy locus is that, the procedure may produce some meaningless numbers which it believes to be the continued fraction coefficients, and at some point it may either just stop the computation after I, J both become singleton sets so that it concludes $\overline{F}[(I \times J) \cap \text{dom}(F)] = \emptyset$, or be trapped into an infinite loop because $\overline{F}[(I \times J) \cap \text{dom}(F)]$ never shrinks down into either $\{\infty\}$ or a compact subinterval of \mathbb{R} , no matter how refined the provided cyclic intervals I, J are.

Another approach is to reject upfront any (x, y) that are in the indeterminacy locus before running Algorithm 12.27. That is, we detect the possibility $(x, y) \notin \text{dom}(\overline{F})$ at the very beginning and produces an error instead of starting the computation. However, this is almost impossible to be done effectively because all we are supposed to know is a cyclic rectangle $I \times J$ where (x, y) belongs to, and unless $I \times J$ is a singleton set, the best we can do with it is to merely checking if it intersects the indeterminacy locus or not. Hence, unless the sequence $(I_n \times J_n)_{n=1}^\infty$ shrinks down to a singleton set in a finite step, we can never confidently reject an invalid input without any false positive.

Yet, rejection with possible false positives can be still valuable. Hence, let us see how to check if $I \times J$ intersects with the indeterminacy locus or not. Note that this rejection can be also leveraged as a preparation step for computing the continued fraction expansion, because once we know $\overline{I \times J}$ is disjoint from the indeterminacy locus, we can simply use Proposition 12.23 instead of Theorem 12.25 which can substantially simplify Algorithm 12.27. Also, since replacing

$$F: (x, y) \mapsto (\langle x, Ay \rangle : \langle x, By \rangle)$$

by

$$F': (x, y) \mapsto (\langle x, By \rangle : \langle x, (A - kB)y \rangle)$$

for some $k \in \mathbb{Z}$ is nothing but taking the composition with the homeomorphism $T_S: \mathbb{RP}^1 \rightarrow \mathbb{RP}^1$ for

$$S := \begin{pmatrix} 0 & 1 \\ 1 & -k \end{pmatrix},$$

the continuous extension \overline{F}' of F' is nothing but the composition $T_S \circ \overline{F}$, and \overline{F} and \overline{F}' have the same indeterminacy locus. Therefore, once we shrink I, J sufficiently so that $\overline{I \times J}$ does not contain any point in the indeterminacy locus of F , then it will stay to be like that in the procedure of computing the continued fraction expansion of $F(x, y)$.

In this reason, let us explore how to detect nonempty intersection between $I \times J$ and the indeterminacy locus of F . Note that the number of elements in the indeterminacy locus of F is either zero or one or two, and they are given as nothing but the solutions to the quadratic equation $L_1(y, y) = 0$. Hence, we could just evaluate them and see if they belong to $I \times J$. However, since we want to confine our computational technology into arbitrary-precision exact rational arithmetic, this will be quite cumbersome if the solutions to $L_1(y, y) = 0$ turn

out to be irrational. Note that since $L_1(y, y) = 0$ is a *quadratic* equation, this can surely happen even if all coefficients of A, B are integers, as the following example demonstrates.

Example 12.28.

Consider $A = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ and $B = \begin{pmatrix} 0 & 3 \\ 1 & 0 \end{pmatrix}$, that is,

$$F(x, y) = \frac{xy + 1}{3x + y}.$$

In this case,

$$L_1 = \begin{pmatrix} 0 & 1 \\ 3 & 0 \end{pmatrix} \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} - \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 3 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 2 & 0 \\ 0 & -6 \end{pmatrix}.$$

Hence, $\det L_1 < 0$, and there are two distinct y such that $L_1(y, y) = 0$, namely

$$y = \begin{pmatrix} \sqrt{3} \\ \pm 1 \end{pmatrix},$$

and Theorem 12.7 tells us that these are indeed the y -coordinates of the points in the indeterminacy locus. The corresponding x 's are given as

$$x = \begin{pmatrix} \mp 1 \\ \sqrt{3} \end{pmatrix},$$

that is, $(x, y) = \left(\mp \frac{1}{\sqrt{3}}, \pm \sqrt{3} \right)$. We can easily see from manual computation that these (x, y) 's indeed eliminate both the numerator and the denominator of $F(x, y)$.

Remark 12.29.

Under the assumption that entries of A, B are all integers, irrational indeterminacy locus can happen only when $\{xy, x, y, 1\}$ is \mathbb{Q} -linearly dependent, as we require an equation of the form $axy + bx + cy + d = 0$ to hold. (More precisely, since we are working on the projective space \mathbb{RP}^1 , the equation we have is of the form $ax_1y_1 + bx_1y_2 + cx_2y_1 + dx_2y_2 = 0$, so $\{x_1y_1, x_1y_2, x_2y_1, x_2y_2\}$ is \mathbb{Q} -linearly dependent, where $x = (x_1 : x_2)$ and $y = (y_1 : y_2)$.)

In this reason, we use the following fact instead of directly solving $L_1(y, y) = 0$:

Proposition 12.30.

Let $A, B \in \mathbb{R}^{2 \times 2}$, and $R = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$ the rotation matrix previously mentioned. As before, let $L_1: \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$ be the symmetric bilinear form on \mathbb{R}^2 defined as

$$L: (u, v) \mapsto \langle u, (B^T R A - A^T R B)v \rangle.$$

Let $u, v \in \mathbb{R}^2$ be such that $\langle u, Rv \rangle > 0$, so in particular $[u] \neq [v]$ as elements in \mathbb{RP}^1 . Then there exists $[y] \in ([u], [v])$ such that $L_1(y, y) = 0$ if and only if:

1. $\det L_1 \leq 0$ and
2. either $L_1 = 0$ or $L_1(u, u)L_1(v, v) < 0$ or $L_1(u, u)L_1(u, v) < 0$ or $L_1(v, v)L_1(u, v) < 0$.

Similarly, there exists $[y] \in [[u], [v]]$ such that $L_1(y, y) = 0$ if and only if:

1. $\det L_1 \leq 0$ and

2. either $L_1(u, u)L_1(v, v) \leq 0$ or $L_1(u, u)L_1(u, v) < 0$.

Proof. As noted in the proof of Theorem 12.7, the equation $L_1(y, y) = 0$ does not have a solution if $\det L_1 > 0$, so suppose $\det L_1 \leq 0$. We want to see if $L_1(y, y) = 0$ admits a solution in $([u], [v])$. Consider the matrix

$$P = \begin{pmatrix} v_1 & u_1 \\ v_2 & u_2 \end{pmatrix}$$

so that P sends $(0, 1) \in \mathbb{R}^2$ to u and $(1, 0) \in \mathbb{R}^2$ to v . By linear independence of u, v , P must be invertible. Also, by replacing v by $-v$ if necessary, we can assume $\det P > 0$. Hence, by Proposition 12.3, T_P maps the interval $[(0 : 1), (1 : 0)] = [0, \infty]$ onto $[[u], [v]]$. Using (17), it can be easily seen that the interval $[(0 : 1), (1 : 0)]$ is precisely the line segment $t \mapsto (t : 1 - t)$. Define a function $h : \mathbb{R} \rightarrow \mathbb{R}$ as

$$h : t \mapsto L_1 \left(P \begin{pmatrix} t \\ 1 - t \end{pmatrix}, P \begin{pmatrix} t \\ 1 - t \end{pmatrix} \right) = t^2 L_1(v, v) + 2t(1 - t)L_1(u, v) + (1 - t)^2 L_1(u, u)$$

so that $h(0) = L_1(u, u)$ and $h(1) = L_1(v, v)$. Clearly, if $h(0)h(1) < 0$, then there exists $t \in (0, 1)$ such that $h(t) = 0$, so suppose otherwise, that is, $L_1(u, u)L_1(v, v) \geq 0$.

Assume first that $L_1(u, u) = L_1(v, v) = 0$. Since h is a quadratic function and $t = 0, 1$ are roots of h , we conclude that there exists $t \in (0, 1)$ with $h(t) = 0$ if and only if h is identically zero. This means $L_1(y, y) = 0$ for all $y \in \mathbb{R}^2$ except possibly on the 1-dimensional subspace generated by $(1, -1)$, but since L_1 is continuous, we in fact have $L_1(y, y) = 0$ even on that subspace. This means $L_1(x, y) = 0$ for all $x, y \in \mathbb{R}^2$ because $0 = L_1(x + y, x + y) = L_1(x, x) + L_1(y, y) + 2L_1(x, y) = 2L_1(x, y)$, so $L_1 = 0$. Hence, when $L_1(u, u) = L_1(v, v) = 0$, there exists $[y] \in ([u], [v])$ such that $L_1(y, y) = 0$ if and only if $L_1 = 0$.

Next, assume $L_1(v, v) \neq 0$. Consider the change of variable $\tau := \frac{t}{1-t}$ which maps $[0, 1)$ bijectively onto $[0, \infty)$, and consider a new function

$$g(\tau) := \frac{h(t)}{(1-t)^2} = \tau^2 L_1(v, v) + 2\tau L_1(u, v) + L_1(u, u).$$

Then $g(\tau) = 0$ if and only if

$$\tau = \frac{-L_1(u, v) \pm \sqrt{L_1(u, v)^2 - L_1(u, u)L_1(v, v)}}{L_1(v, v)}$$

and since we assumed $L_1(u, u)L_1(v, v) \geq 0$, both roots should be in between 0 and $-\frac{2L_1(u, v)}{L_1(v, v)}$. If $L_1(u, u) \neq 0$, then this inclusion is strict, so there exists a strictly positive root of g if and only if $-\frac{2L_1(u, v)}{L_1(v, v)} > 0$, or equivalently, $L_1(v, v)L_1(u, v) < 0$. If $L_1(u, u) = 0$, then two roots of g are 0 and $-\frac{2L_1(u, v)}{L_1(v, v)}$, so we have the same conclusion.

Similarly, for the case $L_1(v, v) = 0$ and $L_1(u, u) \neq 0$, we consider the change of variable $\tau := \frac{1-t}{t}$ which maps $(0, 1]$ bijectively onto $[0, \infty)$, and consider

$$g(\tau) := \frac{h(t)}{t^2} = \tau^2 L_1(u, u) + 2\tau L_1(u, v).$$

Then g has a strictly positive root if and only if $L_1(u, u)L_1(u, v) < 0$.

In conclusion, we proved the claim for the existence of $[y] \in ([u], [v])$ such that $L_1(y, y) = 0$. The claim for $[[u], [v]]$ then follows easily. \square

Using Theorem 12.7 and Proposition 12.30, we can derive an algorithm that checks existence of an indeterminacy locus in $I \times J$ for given cyclic interval estimates I, J of x, y , assuming $F = \overline{F}$. Note, however, that Proposition 12.30 only gives a way to check if J contains the y -coordinate of the indeterminacy locus. We need to check if the associated x -coordinate also is contained in I . To rescue ourselves from all possible weird corner cases, let us first form the intersection of J with I after appropriately transforming it. Then by applying Proposition 12.30 to the intersection, we can check the condition for I and J at the same time.

Still, forming the intersection of two cyclic intervals in \mathbb{RP}^1 is a little more complicated than taking the intersection of two intervals in \mathbb{R} : given two cyclic closed intervals $[a, b], [c, d]$ with a, b, c, d all distinct, there are exactly six possibilities: $[a, b, c, d]$, $[a, b, d, c]$, $[a, c, b, d]$, $[a, c, d, b]$, $[a, d, b, c]$, and $[a, d, c, b]$. Then one can show the following:

Proposition 12.31 (Intersection of cyclic intervals).

Let $[a, b], [c, d]$ be cyclic closed intervals in a cyclically ordered set X with $a \neq b, c \neq d$. Then

$$[a, b] \cap [c, d] = \begin{cases} \emptyset & \text{if } [a, b, c, d], \\ [a, b] & \text{if } [a, b, d, c], \\ [c, b] & \text{if } [a, c, b, d], \\ [c, d] & \text{if } [a, c, d, b], \\ [a, d] & \text{if } [a, d, b, c], \\ [a, d] \sqcup [c, b] & \text{if } [a, d, c, b], \\ [a, b] & \text{if } a = c \text{ and } [a, b, d], \\ [a, d] & \text{if } a = c \text{ and } [a, d, b], \\ \{a\} & \text{if } a = d \text{ and } [a, b, c], \\ \{a\} \sqcup [c, b] & \text{if } a = d \text{ and } [a, c, b], \\ \{b\} & \text{if } b = c \text{ and } [a, b, d], \\ [a, d] \sqcup \{b\} & \text{if } b = c \text{ and } [a, d, b], \\ [a, b] & \text{if } b = d \text{ and } [a, b, c], \\ [c, b] & \text{if } b = d \text{ and } [a, c, b], \\ [a, b] & \text{if } a = c \text{ and } b = d, \\ \{a\} \sqcup \{b\} & \text{if } a = d \text{ and } b = c. \end{cases}$$

In a similar way, we can also make a table of intersection of three intervals $[a, b], [c, d], [e, f]$. In that case, the maximum possible number of components is 3. For example, if $[a, f, e, d, c, b]$ holds, then the intersections becomes three disjoint intervals $[a, f]$, $[e, d]$, and $[c, b]$. Three is indeed the maximum number, because there are only 6 possible endpoints and to have disjoint intervals each interval should have at least two distinct points as its endpoints. This trend continues when the number of intervals increases.

Taking the procedure of computing the intersection of two cyclic intervals as granted, we obtain the following algorithm for checking for the existence of the indeterminacy locus in $I \times J$. In the below we assume that I, J are closed cyclic intervals. This is because, as pointed out earlier, knowing that $\overline{I \times J}$ does not intersect the indeterminacy locus is generally more useful than knowing that for $I \times J$. Nevertheless, it is not tremendously difficult to extend the algorithm below to the general case; see the remark after it.

Algorithm 12.32 (Checking for the existence of a point in the indeterminacy locus).

1. Compute $L_1 := B^T R A - A^T R B$.

2. If one of the following conditions holds, then we conclude that the indeterminacy locus is empty, so return immediately.
 - $\overline{F} \neq F$ (checked by using Corollary 12.9), or
 - $\det L_1 > 0$.
3. For any other cases, any $[y_0] \in \mathbb{RP}^1$ such that $L_1(y_0, y_0) = 0$ is associated to a point in the indeterminacy locus.
 - (a) If A and B are both of rank 1, let $[y_1], [y_2] \in \mathbb{RP}^1$ be unique elements such that $y_1 \in \ker A \setminus \{0\}$ and $y_2 \in \ker B \setminus \{0\}$. Then $L_1(y_1, y_1) = L_1(y_2, y_2) = 0$ holds. In this case, y_1 and y_2 should be linearly independent because otherwise we run into one of the three conditions for having $\overline{F} \neq F$ listed in Corollary 12.9. Then indeterminacy locus is consisting of exactly two points, and their y -coordinates are precisely y_1, y_2 . Hence, there is an indeterminacy locus in $I \times J$ if and only if either of $([RB y_1], [y_1]) \in I \times J$ or $([RA y_2], [y_2]) \in I \times J$ holds.
 - (b) Otherwise, at least one of A, B is invertible. If A is invertible, then let $K := T_{A^{-1}R}[I] \cap J$, and otherwise, let $K := T_{B^{-1}R}[I] \cap J$.
 - (c) If $K = \mathbb{RP}^1$, then we conclude there is an indeterminacy locus in $I \times J$.
 - (d) If $K = \emptyset$, then we conclude there is no indeterminacy locus in $I \times J$.
 - (e) If K is a singleton set, then we conclude there is an indeterminacy locus in $I \times J$ if and only if the bilinear form $y \mapsto L_1(y, y)$ evaluated at the unique element of K is zero.
 - (f) If $K = [[u], [v]]$ for some $u, v \in \mathbb{R}^2 \setminus \{0\}$, then compute $L_1(u, u)L_1(v, v)$ and $L_1(u, u)L_1(u, v)$. Then we conclude there is an indeterminacy locus in $I \times J$ if and only if $L_1(u, u)L_1(v, v) \leq 0$ and $L_1(u, u)L_1(u, v) > 0$ hold.
 - (g) If K is a union of two disjoint cyclic intervals, then we check both of the components as above.

Remark 12.33.

Proposition 12.30 already gives all necessary ingredients for extending the above to general I, J . The only slightly nontrivial cases are when any of I, J are the complements of points. For instance, K can be the complement of one or two points. For the first case, we conclude there is an indeterminacy locus in $I \times J$ if and only if there are more than one points in the indeterminacy locus or the bilinear form $y \mapsto L_1(y, y)$ evaluated at the unique element of $\mathbb{RP}^1 \setminus K$ is not zero. According to Theorem 12.7 and Corollary 12.9, the indeterminacy locus is a singleton set if and only if $\overline{F} = F$ and $\det L_1 = 0$. When K is the complement of two points, then we can see if those excluded points are in the indeterminacy locus by evaluating the bilinear form L_1 at those points. Since we know how many points are there in the indeterminacy locus by evaluating $\det L_1$, we can compare this number with how many excluded points are in the indeterminacy locus. Then $I \times J$ contains a point in the indeterminacy locus if and only if this latter number is strictly less than the total number of points in the indeterminacy locus.

Just like the unary case, extracting $\lfloor \overline{F}(x, y) \rfloor$ from $\overline{F}[(I \times J) \cap \text{dom}(\overline{F})]$ may fail forever no matter how refined the cyclic intervals I, J are, if $\overline{F}(x, y)$ happens to be in $\mathbb{Z} \cup \{\infty\}$. We want to characterize exactly when this can happen. As a preparation, we first state some definition.

Definition 12.34 (Supporting vectors).

Let $(I_n \times J_n)_{n=1}^\infty$ be a decreasing sequence of sets such that I_n, J_n 's are cyclic intervals satisfying $(x, y) \in I_n \times J_n$ for all n and

$$\{(x, y)\} = \bigcap_{n=1}^\infty I_n \times J_n.^{14}$$

Then we call a nonzero tangent vector $X \in T_{(x,y)}(\mathbb{RP}^1 \times \mathbb{RP}^1)$ a *supporting vector* if for all n , there exists a smooth curve $\gamma: (-\epsilon, \epsilon) \rightarrow \mathbb{RP}^1 \times \mathbb{RP}^1$ such that $\gamma(0) = (x, y)$, $\gamma'(0) = X$, and $\gamma[[0, \epsilon]] \subseteq I_n \times J_n$.

In other words, we call X a supporting vector if every $I_n \times J_n$ contains a small line segment starting from (x, y) into the direction X . Due to the simple geometry of $I_n \times J_n$, there are only a few possibilities:

1. If $I_n \times J_n$ shrinks down to a single point (x, y) in a finite step, then there is no supporting vector.
2. If I_n shrinks down to a single point in a finite step while J_n stays to be nontrivial, then there are three possibilities:
 - (a) y eventually becomes the “bottom” endpoint of J_n . In this case, supporting vectors are the vectors along the positive vertical direction.
 - (b) y eventually becomes the “top” endpoint of J_n . In this case, supporting vectors are the vectors along the negative vertical direction.
 - (c) y stays to be inside the interior of J_n . In this case, supporting vectors are the vectors along the vertical direction.
3. If J_n shrinks down to a single point in a finite step while I_n stays to be nontrivial, then there are three possibilities:
 - (a) x eventually becomes the “left” endpoint of I_n . In this case, supporting vectors are the vectors along the positive horizontal direction.
 - (b) x eventually becomes the “right” endpoint of I_n . In this case, supporting vectors are the vectors along the negative horizontal direction.
 - (c) x stays to be inside the interior of I_n . In this case, supporting vectors are the vectors along the horizontal direction.
4. If both I_n and J_n stays to be nontrivial, there are nine possibilities:
 - (a) (x, y) eventually becomes the “left-bottom” corner of $I_n \times J_n$. In this case, supporting vectors are the vectors in the first quadrant.
 - (b) (x, y) eventually becomes the “right-bottom” corner of $I_n \times J_n$. In this case, supporting vectors are the vectors in the second quadrant.
 - (c) (x, y) eventually becomes the “right-top” corner of $I_n \times J_n$. In this case, supporting vectors are the vectors in the third quadrant.

¹⁴Note that this condition is equivalent to that $I_n \times J_n$ is eventually contained in any neighborhood of (x, y) : if not, then for each n we can find $(x_n, y_n) \in (I_n \times J_n) \setminus U$ for some fixed neighborhood U of (x, y) . By compactness, $((x_n, y_n))_{n=1}^\infty$ admits a convergent subsequence, but by the construction this subsequence must be eventually contained in every $I_n \times J_n$, thus is convergent to x , which is a contradiction.

- (d) (x, y) eventually becomes the “left-top” corner of $I_n \times J_n$. In this case, supporting vectors are the vectors in the fourth quadrant.
- (e) (x, y) eventually stays in the interior of the “bottom” edge of $I_n \times J_n$. In this case, supporting vectors are the vectors with positive vertical coordinate.
- (f) (x, y) eventually stays in the interior of the “right” edge of $I_n \times J_n$. In this case, supporting vectors are the vectors with negative horizontal coordinate.
- (g) (x, y) eventually stays in the interior of the “top” edge of $I_n \times J_n$. In this case, supporting vectors are the vectors with negative vertical coordinate.
- (h) (x, y) eventually stays in the interior of the “left” edge of $I_n \times J_n$. In this case, supporting vectors are the vectors with positive horizontal coordinate.
- (i) (x, y) stays in the interior of $I_n \times J_n$. In this case, every nonzero tangent vector is supporting.

Proposition 12.35.

Let $A, B \in \mathbb{R}^{2 \times 2}$ be not both zero and suppose $\overline{F} = F$ and $([x], [y]) \in \text{dom}(F)$. Let $L_1, L_2: \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}$ be symmetric bilinear forms given as

$$\begin{aligned} L_1: (u, v) &\mapsto \langle u, (B^T R A - A^T R B)v \rangle, \\ L_2: (u, v) &\mapsto \langle u, (B R A^T - A R B^T)v \rangle. \end{aligned}$$

Let $(I_n \times J_n)_{n=1}^\infty$ be a decreasing sequence of subsets of $\text{dom}(F)$ such that I_n ’s, J_n ’s are cyclic intervals satisfying $([x], [y]) \in I_n \times J_n \subseteq \text{dom}(F)$ for all n and

$$\{([x], [y])\} = \bigcap_{n=1}^\infty (I_n \times J_n).$$

Let $z \in \mathbb{RP}^1 \setminus \{F([x], [y])\}$. Then

$$F[I_n \times J_n] \subseteq [F([x], [y]), z] \quad (20)$$

holds for all large enough n if and only if one of the followings hold:

1. $I_n \times J_n = \{[x]\} \times \{[y]\}$ for all large enough n , or
2. $I_n = \{[x]\}$ for all large enough n and $L_2(x, x) = 0$, or
3. $I_n = \{[x]\}$, y is one of the endpoints of J_n for all large enough n , and $L_2(x, x) > 0$ if y is the bottom endpoint and $L_2(x, x) < 0$ if y is the top endpoint, or
4. $J_n = \{[y]\}$ for all large enough n and $L_1(y, y) = 0$, or
5. $J_n = \{[y]\}$, x is one of the endpoints of I_n for all large enough n , and $L_1(y, y) > 0$ if x is the left endpoint and $L_1(y, y) < 0$ if x is the right endpoint, or
6. $([x], [y])$ is one of the corners of $I_n \times J_n$ for all large enough n , $\sigma_h L_1(y, y) \geq 0$ and $\sigma_v L_2(x, x) \geq 0$ hold with at least one of the inequalities being strict, where

$$(\sigma_h, \sigma_v) = \begin{cases} (+1, +1) & \text{if } ([x], [y]) \text{ is the left-bottom corner,} \\ (-1, +1) & \text{if } ([x], [y]) \text{ is the right-bottom corner,} \\ (-1, -1) & \text{if } ([x], [y]) \text{ is the right-top corner,} \\ (+1, -1) & \text{if } ([x], [y]) \text{ is the left-top corner, or} \end{cases}$$

7. $([x], [y])$ is one of the corners of $I_n \times J_n$ for all large enough n , $L_1(y, y) = L_2(x, x) = 0$ and $\sigma(\langle Rx, ARy \rangle \langle x, By \rangle - \langle Rx, BRy \rangle \langle x, Ay \rangle) < 0$ where

$$\sigma = \begin{cases} +1 & \text{if } ([x], [y]) \text{ is the left-bottom or the right-top corner,} \\ -1 & \text{if } ([x], [y]) \text{ is the right-bottom or the left-top corner, or} \end{cases}$$

8. $([x], [y])$ is on one of the edges of $I_n \times J_n$ for all large enough n and

$$\begin{cases} L_1(y, y) = 0, L_2(x, x) > 0 & \text{if } ([x], [y]) \text{ is on the bottom edge,} \\ L_1(y, y) < 0, L_2(x, x) = 0 & \text{if } ([x], [y]) \text{ is on the right edge,} \\ L_1(y, y) = 0, L_2(x, x) < 0 & \text{if } ([x], [y]) \text{ is on the top edge,} \\ L_1(y, y) > 0, L_2(x, x) = 0 & \text{if } ([x], [y]) \text{ is on the left edge} \end{cases}$$

holds.

Similarly,

$$F[I_n \times J_n] = \{F(x, y)\} \quad (21)$$

holds for large enough n if and only if one of 1, 2, 4 holds.

Proof. Consider the bilinear mapping

$$\begin{aligned} \tilde{F}: \mathbb{R}^2 \times \mathbb{R}^2 &\rightarrow \mathbb{R}^2 \\ (u, v) &\mapsto (\langle u, Av \rangle, \langle u, Bv \rangle). \end{aligned}$$

Let $\pi: (\mathbb{R}^2 \setminus \{0\}) \rightarrow \mathbb{RP}^1$ be the canonical projection. Take a unit vector $(a, b) \in \mathbb{R}^2 \times \mathbb{R}^2$ such that $d(\pi, \pi)_{(x, y)}(a, b)$ is a supporting vector. Note that for any $w \in \mathbb{R}^2 \setminus \{0\}$, we have $\ker d\pi_w = \langle w \rangle$, so we can identify the tangent space $T_w \mathbb{RP}^1$ with $\langle Rw \rangle \subseteq \mathbb{R}^2$. Note that in this identification, the direction corresponding to the positive direction on \mathbb{R} , with the usual embedding $r \mapsto [r: 1]$ of \mathbb{R} into \mathbb{RP}^1 , is the direction along $-Rw$. Therefore, the condition of having

$$F[I_n \times J_n] \subseteq [F([x], [y]), z)$$

for all large enough n is equivalent to having

$$\tau(t) := \left\langle \tilde{F}((x, y) + t(a, b)), R\tilde{F}(x, y) \right\rangle \leq 0$$

for all $t \geq 0$ small enough and for all (a, b) . Note that τ is a quadratic function of t , whose constant coefficient is zero and the linear coefficient is equal to

$$\begin{aligned} &\left\langle (\langle x, Ab \rangle + \langle a, Ay \rangle, \langle x, Bb \rangle + \langle a, By \rangle), R\tilde{F}(x, y) \right\rangle \\ &= \det \left(\begin{pmatrix} a^T \\ x^T \end{pmatrix} (Ay \ B y) \right) + \det \left(\begin{pmatrix} b^T \\ y \end{pmatrix} (A^T x \ B^T x) \right), \end{aligned}$$

which by applying the identity $\det \begin{pmatrix} u_1 & u_2 \end{pmatrix} = \langle Ru_1, u_2 \rangle$ can be shown to be equal to

$$\frac{1}{2} (\langle a, Rx \rangle L_1(y, y) + \langle b, Ry \rangle L_2(x, x)).$$

Finally, the leading coefficient is equal to

$$\left\langle (\langle a, Ab \rangle, \langle a, Bb \rangle), R\tilde{F}(x, y) \right\rangle = \langle a, Ab \rangle \langle x, By \rangle - \langle a, Bb \rangle \langle x, Ay \rangle.$$

Hence, $F[I_n \times J_n] \subseteq F([x], [y], z)$ holds for all large enough n if and only if for any (a, b) such that $d(\pi, \pi)_{(x,y)}(a, b)$ is a supporting vector, either

$$\langle a, Rx \rangle L_1(y, y) + \langle b, Ry \rangle L_2(x, x) < 0$$

or

$$\begin{aligned} \langle a, Rx \rangle L_1(y, y) + \langle b, Ry \rangle L_2(x, x) &= 0 \quad \text{and} \\ \langle a, Ab \rangle \langle x, By \rangle - \langle a, Bb \rangle \langle x, Ay \rangle &\leq 0 \end{aligned}$$

holds.

Note that the “horizontal coordinate” of the supporting vector $d(\pi, \pi)_{(x,y)}(a, b)$ is equal to $-\langle a, Rx \rangle$ while the “vertical coordinate” is equal to $-\langle b, Ry \rangle$, according to our identification of $T_x \mathbb{RP}^1$ with $\langle Rx \rangle$. Then considering how the set of supporting vectors look like, it can be easily seen that $L_1(y, y) \geq 0$ should hold if there is a supporting vector of nonnegative vertical coordinate, $L_1(y, y) \leq 0$ should hold if there is a supporting vector of nonpositive vertical coordinate, and similarly $L_2(x, x) \geq 0$ or $L_2(x, x) \leq 0$ should hold according to existence of a supporting vector with positive or negative horizontal coordinate.

Similarly, $F[I_n \times J_n] = \{F([x], [y])\}$ holds for all large enough n if and only if for any (a, b) such that $d(\pi, \pi)_{(x,y)}(a, b)$ is a supporting vector,

$$\langle a, Rx \rangle L_1(y, y) + \langle b, Ry \rangle L_2(x, x) = \langle a, Ab \rangle \langle x, By \rangle - \langle a, Bb \rangle \langle x, Ay \rangle = 0$$

holds. In particular, we must have $L_1(y, y) = 0$ if there is a supporting vector with a nonzero horizontal coordinate while we must have $L_2(x, x) = 0$ if there is one with a nonzero vertical coordinate. Based on this, we consider each possible location of $([x], [y])$ in $I_n \times J_n$.

- When $I_n \times J_n = \{[x]\} \times \{[y]\}$ for large enough n .

In this case, we always have $F[I_n \times J_n] = \{F([x], [y])\}$, so there is nothing to prove.

- When $I_n = \{[x]\}$ for large enough n , but J_n is not a single point for all n .

Note that $L_2(x, x) = 0$ is a necessary condition for having (21). To see sufficiency, note that $\tau(t)$ has no quadratic term in t if $a = 0$ or $b = 0$, so the linear term being zero is same as it being identically zero. Hence, whenever J_n is small enough, $F[I_n \times J_n]$ should be a singleton set.

On the other hand, (20) certainly implies $L_2(x, x) = 0$ if $[y]$ is in the interior of J_n , $L_2(x, x) \geq 0$ if it is the bottom endpoint, or $L_2(x, x) \leq 0$ if it is the top endpoint. As noted above, $L_2(x, x) = 0$ is equivalent to $\tau \equiv 0$, so $L_2(x, x) = 0$ is certainly a sufficient condition for all cases. When $L_2(x, x) \neq 0$, the sign of $\langle \tilde{F}((x, y) + t(0, b)), R\tilde{F}(x, y) \rangle$ for small enough t is determined by the sign of its first derivative at $t = 0$, so $L_2(x, x) > 0$ or $L_2(x, x) < 0$ should be a sufficient condition as well, respectively for when $[y]$ is the bottom or the top endpoint of J_n .

- When $J_n = \{[y]\}$ for large enough n , but I_n is not a single point for all n .

Exactly the same analysis applies, so we omit this case.

- When neither I_n nor J_n is a singleton set for all n .

To analyze (20), we further divide the cases.

- When $([x], [y])$ is one of the corners of $I_n \times J_n$ for large enough n .
Clearly, $\sigma_h L_1(y, y) \geq 0$ and $\sigma_v L_2(x, x) \geq 0$ together are necessary conditions for (20). Then it is enough to show three claims: (1) if $L_1(y, y) = L_2(x, x) = 0$, then (20) is equivalent to

$$\sigma (\langle Rx, ARy \rangle \langle x, By \rangle - \langle Rx, BRy \rangle \langle x, Ay \rangle) \leq 0, \quad (22)$$

(2) the above cannot be zero in addition to $L_1(y, y) = L_2(x, x) = 0$, and (3) having a strict inequality for one of $\sigma_h L_1(y, y) \geq 0$ and $\sigma_v L_2(x, x) \geq 0$ is enough to conclude (20).

For the first claim, suppose $L_1(y, y) = L_2(x, x) = 0$. Then (20) is equivalent to having

$$\langle a, Ab \rangle \langle x, By \rangle - \langle a, Bb \rangle \langle x, Ay \rangle \leq 0$$

for all (a, b) such that $d(\pi, \pi)_{(x, y)}(a, b)$ is a supporting vector. Since $L_1(y, y) = 0$ means that Ay and By are parallel, and $L_2(x, x) = 0$ means that $A^T x$ and $B^T x$ are parallel, decomposing a into x -component and Rx -component and b into y -component and Ry -component shows that the above equality is equivalent to

$$\langle a, Rx \rangle \langle b, Ry \rangle (\langle x, ARy \rangle \langle x, By \rangle - \langle Rx, BRy \rangle \langle x, Ay \rangle) \leq 0.$$

Then it is easy to see that σ is precisely the sign of $\langle a, Rx \rangle \langle b, Ry \rangle$, so the first claim is proved.

For the second claim, note that having all three being zero means that the function τ is identically zero. Therefore, $F|_{I_n \times J_n}$ should be a constant in a neighborhood of $([x], [y])$, contradicting to Proposition 12.16.

The third claim directly follows from that the sign of $\langle a, Rx \rangle$ and $\langle b, Ry \rangle$ is the negative of σ_h and σ_v , respectively.

- When $([x], [y])$ is on one of the edges of $I_n \times J_n$ for large enough n .

In this case, sufficiency of

$$\begin{cases} L_1(y, y) = 0, L_2(x, x) > 0 & \text{if } ([x], [y]) \text{ is on the bottom edge,} \\ L_1(y, y) < 0, L_2(x, x) = 0 & \text{if } ([x], [y]) \text{ is on the right edge,} \\ L_1(y, y) = 0, L_2(x, x) < 0 & \text{if } ([x], [y]) \text{ is on the top edge,} \\ L_1(y, y) > 0, L_2(x, x) = 0 & \text{if } ([x], [y]) \text{ is on the left edge} \end{cases}$$

for having (20) is clear. Then it is enough to show that $L_1(y, y) = L_2(x, x) = 0$ cannot be true. This follows from that (22) should hold for both positive and negative σ , because $\langle a, Rx \rangle \langle b, Ry \rangle$ can have two different signs. Therefore, τ is identically zero, thus $F|_{I_n \times J_n}$ is a constant in a neighborhood of $([x], [y])$, contradicting to Proposition 12.16.

- When (x, y) stays inside the interior of $I_n \times J_n$.

By Proposition 12.16, in this case (20) cannot hold.

□

As in the unary case, when we repeatedly apply the transform $F \mapsto F'$ given in (19), $\bar{F}(x, y) \in \mathbb{Z} \cup \{\infty\}$ eventually happens if and only if $\bar{F}(x, y) \in \mathbb{Q} \cup \{\infty\} = \mathbb{QP}^1$. (Actually, it can become ∞ only after $\bar{F}(x, y) \in \mathbb{Z}$ holds, unless $\bar{F}(x, y) = \infty$ were the case from the beginning.) Hence, the problem may exist only when $\bar{F}(x, y) \in \mathbb{QP}^1$.

Note that if we assume $\overline{F} = F$, all coefficients of A, B are integers, and also that the interval estimates I_n, J_n are always with rational endpoints and they converge down to single points if and only if x, y are rational, respectively, then this infinite loop problem can only happen when x, y are both irrational numbers, in which case none of the items from the proposition above applies since (x, y) is always in the interior of $I_n \times J_n$. Indeed, let $x = [u]$, $y = [v]$ for $u, v \in \mathbb{R}^2 \setminus \{0\}$ and suppose x is irrational while y is rational (i.e., $y \in \mathbb{QP}^1$), then

$$F(x, y) = (\langle u, Av \rangle : \langle u, Bv \rangle) \in \mathbb{QP}^1$$

implies that there exists $(p : q) \in \mathbb{QP}^1$ such that

$$\langle u, (qA - pB)v \rangle = 0,$$

and since $[u]$ is irrational and $[(qA - pB)v]$ is rational, we must have

$$qAv = pBv.$$

Therefore, as noted in Theorem 12.7, $L_1(v, v) = 0$ holds, thus the condition 4 from Proposition 12.35 is satisfied.

Similarly, if $F(x, y) \in \mathbb{QP}^1$ then $\{xy, x, y, 1\}$ must be \mathbb{Q} -linearly dependent. Indeed, from the equation

$$\langle u, (qA - pB)v \rangle = 0,$$

assuming \mathbb{Q} -linear independence of $\{xy, x, y, 1\}$ (which by definition means \mathbb{Q} -linear independence of $\{u_1v_1, u_1v_2, u_2v_1, u_2v_2\}$) shows $qA - pB = 0$, contradicting to $\overline{F} = F$.

Finally, note that if (x, y) stays inside the interior, then (20) holds if and only if A, B are linearly dependent. In other words, if A, B are linearly independent and x, y are both irrational, then having $\overline{F}(x, y) \in \mathbb{QP}^1$ guarantees the procedure to fall into an infinite loop.

In summary, assuming $A, B \in \mathbb{Z}^{2 \times 2}$ and that I_n 's and J_n 's are always with rational endpoints and they converge down to single points in a finite step if and only if x, y are rational, respectively, then the described procedure either succeeds in computing the continued fraction expansion of $\overline{F}(x, y)$ indefinitely, or rejects the input (x, y) for successfully identifying $(x, y) \notin \text{dom}(\overline{F})$, if and only if one of the following conditions hold:

1. A, B are linearly dependent, or
2. \overline{F} depends only on x or y , or
3. $\overline{F} = F$ and at least one of x, y is rational, or
4. $(x, y) \in \text{dom}(\overline{F})$ and $\overline{F}(x, y) \notin \mathbb{QP}^1$.

Also, assuming that A, B are \mathbb{Q} -linearly independent, the last condition automatically holds whenever $\{xy, x, y, 1\}$ is \mathbb{Q} -linearly independent.

Unfortunately, generically speaking, if x, y are both irrational but $\overline{F}(x, y)$ is rational, then the procedure is almost guaranteed to trap into an infinite loop, but this is rather a fundamentally unsolvable issue.

Remark 12.36.

There are two notable differences between the explained algorithm and what Gosper originally proposed.

1. There is no mention of singularity in Gosper's original proposal. In particular, it does not mention the possible continuous extension. Thus, the implementation may run into divisions of the form $\frac{0}{0}$ even when (x, y) is technically a "valid" input. For truly invalid inputs, i.e., those from the indeterminacy locus, not only $\frac{0}{0}$ but also other kinds of pathologies can happen, for instance spitting out meaningless numbers which it believes to be the continued fraction coefficient.
2. Computation of $F[I \times J]$ is much simpler in Gosper's original proposal. Aside from not distinguishing various types of cyclic intervals, a big factor is that it just gives up and try to refine I, J immediately whenever any of the images of four edges by F contains ∞ . Since $F[I \times J]$ is the union of those four images and we cannot compute $\lfloor F(x, y) \rfloor$ anyway if $F[I \times J]$ contains ∞ , this is totally fine, and it significantly simplifies the procedure of computing the union because all we need to do is to simply find the minimum and the maximum of endpoints once every interval is contained in \mathbb{R} . However, I considered the unconditional computation of $F[I \times J]$ might be useful in certain scenarios, for example when we want to feed $F(x, y)$ as another input to Gosper's algorithm, e.g. to compute $F(F(x, y), z)$. My viewpoint is that Gosper's algorithm is, in essence, an algorithm for finding the corresponding interval estimate from given interval estimates for the input, rather than an algorithm for computing the continued fraction expansion. The latter just immediately follows from the former as a trivial application.

13 Continued fractions for logarithms

Some of the (in fact, *the only*) motivating examples I had for computing $\lfloor nx + y \rfloor$ that involved irrational x, y were when x, y were some general logarithms with base 2 or 10. Hence, I devoted a section for explaining how to obtain continued fraction expansions of logarithms.

When we restrict ourselves to the case when the base and the input are both rational, there is a very simple algorithm for computing the continued fraction expansion. Let us say that a, b are positive rational numbers with $a \neq 1$, and we want to compute the continued fraction expansion of $\log_a b$. Note that, once we know $k := \lfloor \log_a b \rfloor$, then

$$(\log_a b - k)^{-1} = (\log_a(b/a^k))^{-1} = \log_{(b/a^k)} a$$

is still a general logarithm with rational base and input. Therefore, it is enough to see how to compute $\lfloor \log_a b \rfloor$. For that, we can use the definition of logarithms: we find the largest integer k such that

$$a^k \leq b < a^{k+1}$$

if $a > 1$, or

$$a^k \geq b > a^{k+1}$$

if $a < 1$.

However, no matter what specific algorithm we use for finding k , this way of computing the continued fraction expansion of $\log_a b$ generally does not work in practice, not because it is wrong, rather because the required precision of involved numbers grows exponentially, so the computation very quickly becomes totally infeasible.

Instead, we compute the continued fraction expansion of $\log_a b$ using the identity $\log_a b = \frac{\ln b}{\ln a}$: once we know how to compute the continued fraction expansions of $\ln a$ and $\ln b$, we can apply Gosper's algorithm. In fact, we do not really need to compute the continued fraction

expansions for $\ln a$ and $\ln b$. What we really need are the interval estimates $(I_n)_{n=1}^\infty, (J_n)_{n=1}^\infty$ of $\ln a$ and $\ln b$, because that is all we need to run Gosper's algorithm.

Therefore, any method of computing successively better approximations of the natural logarithm should work. Nevertheless, when I started writing the implementation, I did not really realize this, so was trying to find a way to compute the continued fraction expansions of natural logarithms. Thus, I searched for existing literatures, and found two papers by Isao Makino and Takeshi Aoyama on computing natural logarithm using its continued fraction expansion [13][14].¹⁵ After I realized there probably are other methods that fit my purpose better, I already had implemented the method explained in these papers and it did its job good enough, so I just moved on. The following subsections explain some details of this method I implemented.

13.1 Continued fractions for natural logarithm

As far as I am aware, there is no known formula for the regular continued fraction expansion of natural logarithms. The method I will explain in this subsection does not compute such; rather, it computes a *generalized* continued fraction that converges to the natural logarithm. However, this is fine because all we need is just a sequence of successively better interval estimates. In the end of the day, coming up with an error bound is the only thing we need to do.

We start from the following expansion of the inverse hyperbolic tangent:

$$\log\left(\frac{1+z}{1-z}\right) = 2 \tanh^{-1} z = \frac{2z}{1 - \frac{z^2}{3 - \frac{4z^2}{5 - \frac{9z^2}{7 - \dots}}}}$$

In particular, when $z = \frac{p}{q}$, we can write it as

$$2 \tanh^{-1}\left(\frac{p}{q}\right) = \frac{2p}{q - \frac{p^2}{3q - \frac{4p^2}{5q - \frac{9p^2}{7q - \dots}}}} \quad (23)$$

Recall from Section 3 that given a generalized continued fraction

$$a_0 + \frac{b_1}{a_1 + \frac{b_2}{a_2 + \frac{b_3}{a_3 + \frac{b_4}{a_4 + \dots}}}} \quad (24)$$

the recurrence relation for the convergents is given as

$$\begin{cases} p_{n+1} = a_{n+1}p_n + b_{n+1}p_{n-1} \\ q_{n+1} = a_{n+1}q_n + b_{n+1}q_{n-1} \end{cases}, \quad \begin{cases} (p_{-1}, p_0) = (1, a_0) \\ (q_{-1}, q_0) = (0, 1) \end{cases}$$

¹⁵These papers are quite poorly written and especially contain a lot of typos.

where $\frac{p_n}{q_n}$ is the n th convergent. Specializing this for (23), we obtain

$$\begin{cases} p_{n+1} = (2n+1)qp_n - n^2p^2p_{n-1} \\ q_{n+1} = (2n+1)qq_n - n^2p^2q_{n-1} \end{cases}, \quad \begin{cases} (p_0, p_1) = (0, 2p) \\ (q_0, q_1) = (1, q) \end{cases}$$

for $n \geq 1$. At this point, it is not clear if the sequence $\left(\frac{p_n}{q_n}\right)_{n=1}^{\infty}$ even converges.¹⁶

Recall from Corollary 3.2 the identity

$$\frac{p_{n-1}}{q_{n-1}} - \frac{p_n}{q_n} = \frac{\prod_{i=1}^n (-b_i)}{q_{n-1}q_n},$$

so in the case of (23), this becomes

$$\frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n} = \frac{2p^{2n+1}(n!)^2}{q_n q_{n+1}} \quad (25)$$

for $n \geq 0$. Therefore, if q_n grows sufficiently faster than $p^n n!$, then $\left(\frac{p_n}{q_n}\right)_{n=1}^{\infty}$ converges to some real number α and the convergence is monotone.

Now, recall that \tanh^{-1} is not well-defined for real numbers of modulus greater than or equal to 1, so our $z = \frac{p}{q}$ should be such that $|p| < q$. Note that, because of this, if we had only the first term in the recurrence relation

$$q_{n+1} = (2n+1)qq_n - n^2p^2q_{n-1},$$

then we indeed have an enough growth of q_n needed to conclude the convergence. Of course, the problem is that the second term is of the opposite sign so it is not obvious if we can still ensure a similar rate of growth. In this reason, Makino and Aoyama tried to compare a smaller multiple of $(2n+1)qq_n$ with $n^2p^2q_{n-1}$, and obtained the result below. To be precise, they claimed they have a more general result (which they called the “generalized error evaluation equation” or *GEEE* in short), but I seriously doubt the credibility of that result.

Theorem 13.1 (Makino-Aoyama, 2000).

Let $z = \frac{p}{q}$ be any rational number with $|z| < 1$. Then

$$qq_n > np^2q_{n-1}$$

holds for all $n \geq 1$, where $\frac{p_n}{q_n}$ is the n th convergent from the generalized continued fraction (23).

Proof. We use induction on n . When $n = 1$,

$$qq_1 - p^2q_0 = q^2 - p^2 > 0$$

by the assumption on z . Now, suppose

$$qq_n > np^2q_{n-1}$$

¹⁶It seems that (23) is a specific instance of the so called *Gauss’ continued fractions*, which seem to have a general theory of convergence already developed. However, I think that is not extremely relevant here. I will probably look into it later.

holds for some $n \geq 1$, then

$$\begin{aligned} qq_{n+1} - (n+1)p^2q_n &= q((2n+1)qq_n - n^2p^2q_{n-1}) - (n+1)p^2q_n \\ &= ((2n+1)q^2 - (n+1)p^2)q_n - qn^2p^2q_{n-1} \\ &> (n+1)(p^2 - q^2)q_n \end{aligned}$$

where the last inequality follows from the induction hypothesis. By the assumption on z , the right-hand side is positive, so we get the conclusion. \square

Corollary 13.2.

In the setting of Theorem 13.1, we have followings:

1. $q_{n+1} \geq (n+1)qq_n$ holds for all $n \geq 0$ where the equality holds only when $n = 0$, and
2. The sequence $\left(\frac{p_n}{q_n}\right)_{n=0}^{\infty}$ converges to some $\alpha \in \mathbb{R}$ and

$$\left|\frac{p_n}{q_n} - \alpha\right| \leq \frac{2|z|^{2n+1}}{(n+1)(1-|z|^2)} \quad (26)$$

holds for all $n \geq 0$. Furthermore, the sequence $\left(\frac{p_n}{q_n}\right)_{n=0}^{\infty}$ strictly increases to the limit α if $z > 0$ and it strictly decreases to α if $z < 0$.

Proof. 1. Directly follows from the recurrence relation $q_{n+1} = (2n+1)qq_n - n^2p^2q_{n-1}$ and Theorem 13.1.

2. Since $q_n \geq (n!)q^n$ follows from 1, (25) yields

$$\left|\frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n}\right| \leq \frac{2|z|^{2n+1}}{n+1}.$$

Hence, for any $m > n$,

$$\left|\frac{p_m}{q_m} - \frac{p_n}{q_n}\right| \leq \frac{2|z|^{2n+1}}{n+1} (1 + |z|^2 + \dots + |z|^{2(m-n-1)}) \leq \frac{2|z|^{2n+1}}{(n+1)(1-|z|^2)}$$

holds. This shows the convergence and the inequality (26). Note that in (25), the sign of the right-hand side is equal to the sign of z , so the last statement also follows. \square

Therefore, we can use the convergent $\frac{p_n}{q_n}$ to approximate $\log\left(\frac{1+z}{1-z}\right)$ as precise as desired, where the error bound can be estimated using (26). However, the convergence speed estimated by (26) is not very great, so we derive a better bound.

Proposition 13.3.

In the setting of Theorem 13.1, let $\alpha := \lim_{n \rightarrow \infty} \frac{p_n}{q_n}$, then

$$\left|\frac{p_0}{q_0} - \alpha\right| \leq \frac{2|p|q}{q^2 - p^2}$$

and

$$\left|\frac{p_n}{q_n} - \alpha\right| \leq \frac{2|p|^{2n+1}(n-1)!n!}{(n+1)(q^2 - p^2)q_{n-1}q_n}$$

holds for all $n \geq 1$.

Proof. Recall that for all $n \geq 0$, we have $q_{n+1} \geq (n+1)q_n$. Hence,

$$\begin{aligned} \left| \frac{p_{n+2}}{q_{n+2}} - \frac{p_{n+1}}{q_{n+1}} \right| &= \frac{2|p|^{2n+3}((n+1)!)^2}{q_{n+1}q_{n+2}} \leq \frac{|p|^2(n+1)}{q^2(n+2)} \cdot \frac{2|p|^{2n+1}(n!)^2}{q_nq_{n+1}} \\ &= \frac{n+1}{n+2} |z|^2 \left| \frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n} \right| \end{aligned} \quad (27)$$

holds for all n . Therefore, inductively we get

$$\begin{aligned} \left| \frac{p_{n+k+1}}{q_{n+k+1}} - \frac{p_{n+k}}{q_{n+k}} \right| &\leq \frac{n+k}{n+k+1} \cdot \frac{n+k-1}{n+k} \cdot \dots \cdot \frac{n+1}{n+2} |z|^{2k} \left| \frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n} \right| \\ &= \frac{n+1}{n+k+1} |z|^{2k} \left| \frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n} \right| \end{aligned}$$

for any $k \geq 0$. Thus,

$$\left| \frac{p_{n+k+1}}{q_{n+k+1}} - \frac{p_n}{q_n} \right| \leq \frac{1}{1-|z|^2} \left| \frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n} \right|$$

holds for all $k \geq 0$, so taking the limit $k \rightarrow \infty$ gives

$$\left| \frac{p_n}{q_n} - \alpha \right| \leq \frac{1}{1-|z|^2} \left| \frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n} \right|.$$

Specializing this with $n = 0$ then gives

$$\left| \frac{p_0}{q_0} - \alpha \right| \leq \frac{q^2}{q^2 - p^2} \cdot \frac{2|p|}{q} = \frac{2|p|q}{q^2 - p^2}.$$

On the other hand, if $n \geq 1$, then (27) shows

$$\left| \frac{p_{n+1}}{q_{n+1}} - \frac{p_n}{q_n} \right| \leq \frac{n}{n+1} |z|^2 \left| \frac{p_n}{q_n} - \frac{p_{n-1}}{q_{n-1}} \right|,$$

thus we get

$$\left| \frac{p_n}{q_n} - \alpha \right| \leq \frac{|z|^2}{1-|z|^2} \cdot \frac{n}{n+1} \cdot \frac{2|p|^{2n-1}((n-1)!)^2}{q_{n-1}q_n} = \frac{2|p|^{2n+1}(n!)^2}{n(n+1)(q^2 - p^2)q_{n-1}q_n}$$

as desired. \square

Having established an error bound formula for $2 \tanh^{-1} \left(\frac{p}{q} \right)$, we compute that for $\ln x$ by doing the change of variable $z = \frac{x-1}{x+1}$. Note that $z \in (-1, 1)$ holds if and only if $x \in (0, \infty)$ holds. Then when x is rational, we can find p, q by finding the reduced form of $\frac{x-1}{x+1}$.

13.2 When the logarithm is rational

Once we have a way to compute successively better interval estimates for natural logarithm of positive rational numbers, we can immediately compute the continued fraction expansion of $\log_a b$ for any positive rational numbers a, b (even including the case $a = 1$, if we interpret $\log_a b$ as ∞ in that case) by applying Gosper's algorithm to $\ln b / \ln a$. However, recall that

Gosper's algorithm traps in an infinite loop if $\log_a b$ is rational.¹⁷ Therefore, we have to detect this case upfront, and if it turns out to be the case, then use the continued fraction expansion algorithm for rational numbers.

Essentially, this can be in fact done by going back to our first, naïve attempt of computing the continued fraction expansion of $\log_a b$. Roughly speaking, what we do is basically to just follow the naïve algorithm, and conclude that $\log_a b$ should be irrational if the numbers involved grow at all.

To be precise, suppose if we want to see $\log_a b = \frac{m}{n}$ holds where $\frac{m}{n}$ is a rational number in its reduced form, which means $a^m = b^n$. If $a = \prod_{i=1}^d p_i^{e_i}$ is the prime factorization, then $a^m = \prod_{i=1}^d p_i^{me_i}$ must be the prime factorization of b^n , but since m and n are coprime, it follows that n divides every e_i . Therefore, there exists a rational number c such that $a = c^n$, which shows $b = c^m$.

Now, let $a = \frac{p}{q}$, $b = \frac{r}{s}$, and $c = \frac{t}{u}$ be their reduced forms. Then since $\frac{p}{q}$ and $\frac{t^n}{u^n}$ are both of their reduced forms, we conclude $p = t^n$ and $q = u^n$. Similarly, we have $r = t^m$ and $s = u^m$ if $m \geq 0$, or $r = u^{-m}$ and $s = t^{-m}$ if $m < 0$.

For simplicity, so we first detect the sign of m and replace b by $1/b$ to make m nonnegative if necessary. Since $\log_a b \geq 0$ if and only if either $a > 1$ and $b \geq 1$ or $a < 1$ and $b \leq 1$, this can be easily done by inspecting these inequalities. Hence, without loss of generality, we can assume $m \geq 0$, so $r = t^m$ and $s = u^m$.

Next, we ensure $m \geq n$ by swapping a and b if necessary. Note that $\frac{m}{n} \geq 1$ if and only if either $a > 1$ and $b \geq a$ or $a < 1$ and $b \leq a$, so by simply inspecting these inequalities, we can assume without loss of generality that $m \geq n$.

Then, now we see if the equations $p = t^n$, $r = t^m$ have a solution. To do so, note that if it ever had a solution, then p must divide r since $m \geq n$ is assumed. If this is not the case, then we immediately conclude that $\log_a b$ is irrational. Otherwise, write $r = p^d p'$ where d is the highest integer exponent such that p^d divides r . If $\log_a b$ is rational, then this d must be equal to $\lfloor \frac{m}{n} \rfloor$ and p' must be equal to $t^{(m \bmod n)}$. Hence, we run Euclid's algorithm at this point. That is, we replace (p, r) by (p', p) and repeat this procedure. If $\log_a b$ is rational, then at some point p' must become 1. Conversely, if we reach to the point where p' becomes 1, then we conclude that $p = t^n$ and $r = t^m$ have a solution.

In addition to that, by tracking the relevant coefficients, we can also figure out m, n from this procedure. Concretely, we start from

$$\begin{pmatrix} \log_t p \\ \log_t r \end{pmatrix} = \begin{pmatrix} m \\ n \end{pmatrix},$$

and note that

$$\begin{pmatrix} \log_t p' \\ \log_t p \end{pmatrix} = \begin{pmatrix} -d & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \log_t p \\ \log_t r \end{pmatrix} = \begin{pmatrix} -d & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} m \\ n \end{pmatrix},$$

so

$$\begin{pmatrix} m \\ n \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & d \end{pmatrix} \begin{pmatrix} \log_t p' \\ \log_t p \end{pmatrix}.$$

Hence, by repeating the procedure until p' becomes 1, we get

$$\begin{pmatrix} m \\ n \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & d_1 \end{pmatrix} \cdots \begin{pmatrix} 0 & 1 \\ 1 & d_k \end{pmatrix} \begin{pmatrix} 0 \\ \log_t r_k \end{pmatrix}$$

where (p_{i+1}, r_{i+1}) is the pair (p', p) obtained from (p_i, r_i) , $(p_0, r_0) = (p, r)$ and d_{i+1} is the highest integer exponent such that $p_i^{d_{i+1}}$ divides r_i . Since m and n are coprime, r_k should

¹⁷Note that $\ln x$ is never rational if x is rational. Indeed, if $e^{p/q} = x$ holds, then e is in a finite field extension of $\mathbb{Q}(x)$, but since e is transcendental, this is impossible if x is rational.

be equal to t , so the second entry of the last matrix in the right-hand side is actually 1. Therefore, by keeping track of the product of $\begin{pmatrix} 0 & 1 \\ 1 & d_i \end{pmatrix}$'s, we can compute m, n .

Note that this only solves the half of the problem: we also need to see the existence of a solution to the equations $q = u^n$ and $s = u^m$. However, the exactly the same numbers should appear for d_i 's, and if not, then we immediately conclude that $\log_a b$ is irrational. When p' reaches 1, the corresponding q' also must be 1 if $\log_a b$ is rational. Otherwise, we conclude $\log_a b$ is irrational. If they are both 1, then we conclude $\log_a b$ is rational.

14 TODO List

- Think of an algorithm that determines the exact range of $n \geq 0$ such that $\lfloor nx + y \rfloor = \lfloor n\xi + \zeta \rfloor$ holds, when x, y, ξ, ζ are all given. Will it be possible to get rid of the assumption $n_{\min} = 0$?
- Applications to James Anhalt's algorithm: find ξ, ζ such that $n = \left\lfloor \frac{\lfloor n\xi + \zeta \rfloor}{x} \right\rfloor$ holds for given x and $n \in [n_{\min} : n_{\max}]$, or equivalently, $\lfloor nx \rfloor \leq n\xi + \zeta < \lfloor nx + x \rfloor$.
- Application to log computations in Dragonbox.
- Think of an algorithm for solving $\lfloor nx \rfloor \equiv \lfloor n\xi \rfloor \pmod{D}$; i.e., for given $x \in \mathbb{R}$ and $D, n_{\max} \in \mathbb{Z}_{>0}$, what is the set of $\xi \in \mathbb{R}$ such that $\lfloor nx \rfloor \equiv \lfloor n\xi \rfloor \pmod{D}$ holds for all $n \in [1 : n_{\max}]$? Obviously, if $\lfloor nx \rfloor = \lfloor n\xi \rfloor$ holds then we do have the equivalence, so the question is how much the set of ξ would become bigger if we relax the equality into modular equivalence. When we set $x = 1$, this will give a generalized version of Lemire's remainder computation algorithm described in [3]. We can maybe also apply this into fixed-precision floating-point formatting problem I described in my blog post, though I do not expect a lot of improvement in this regard. Anyway, my current rough idea goes like the following. First, we can assume without loss of generality that $0 \leq x, \xi < D$. Then we are looking for the set

$$\bigcap_{n=1}^{n_{\max}} \bigcup_{k \in \mathbb{Z}} \left[\frac{\lfloor nx \rfloor + kD}{n}, \frac{\lfloor nx \rfloor + 1 + kD}{n} \right) \cap [0, D).$$

The size of each of the inner interval is $\frac{1}{n}$, and the gap between intervals with consecutive k 's is $\frac{D}{n}$. In particular, the union inside the intersection may have multiple components which complicates the problem. However, at least when $n = 1$, there is only one or two components in the union. Then taking the intersection for small enough n will not increase the number of components, because many of the components coming from the union will not intersect with the intersection we have for smaller n 's. Hence, probably we can run the algorithm described in Section 10 on each component for a range of n where the number of component does not change. When any of the components splits into several pieces at some n , then we record that and do this iteratively. Or at some n the number of components may decrease. What matters here is that we have to iteratively figure out a (possibly conservative) bound on n to guarantee no splitting of the component we are looking at. My intuition tells me that the number of components should not grow too rapidly, and probably will just stay small, or may increase linearly at worst, but I do not have any formal nor experimental proof.

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