The Internet: a "Nuts And Bolts" View



Billions of connected computing *devices*:

- hosts = end systems
- running *network apps* at Internet's "edge"



Packet switches: forward packets (chunks of data)

routers, switches



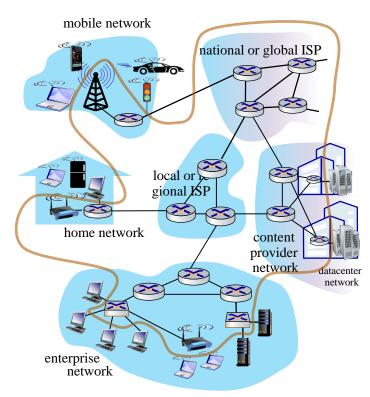
Communication links

- fiber, copper, radio, satellite
- transmission rate: bandwidth



Networks

 collection of devices, routers, links: managed by an organization

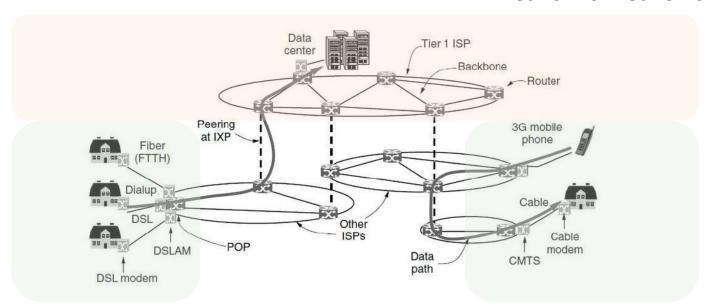




Network Core/Edge

Network core:

- interconnected routers
- network of networks



- Network edge:
 - hosts: clients and servers
 - servers often in data centers
- Access network (last-mile network)

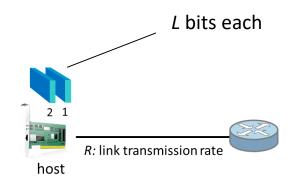




Network Core

Transmission delay:

packet transmission delay time needed to transmit L-bit packet into link L (bits)
R (bits/sec)



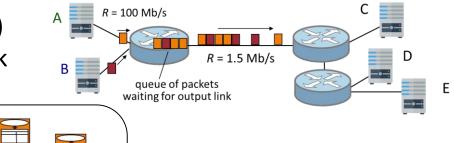
store and forward: entire packet must arrive at router before it can be transmitted on next link

L bits per packet source R bps destin

queuing and loss: if arrival rate (in bits) to link exceeds transmission rate of link for a period of time

destination address in arriving packet's header

Packet Switching & Routing









Introduction

CSE351
Computer Networks

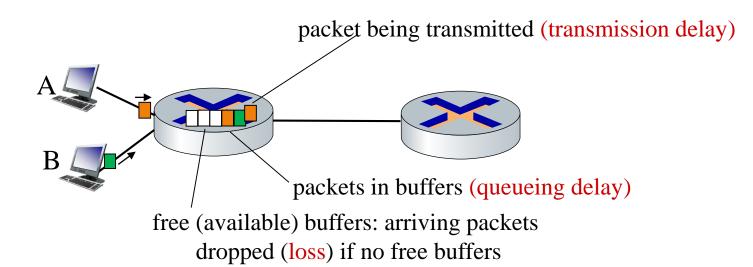




How do loss and delay occur?

packets *queue* in router buffers

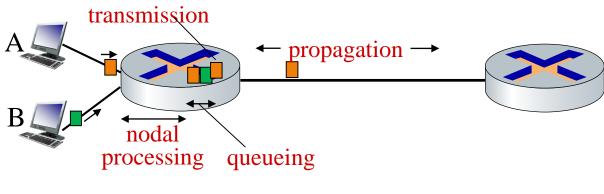
- Packet arrival rate to link (temporarily) exceeds output link capacity
- Packets queue, wait for turn
- Packets are dropped in the tail part (drop-tail queue)







Packet Delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{proc} : nodal processing

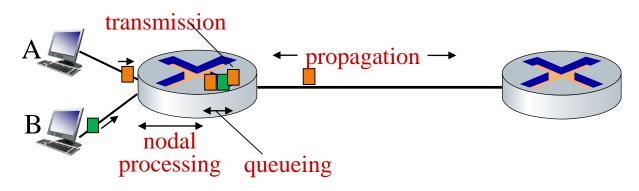
- check bit errors
- determine output link
- typically < msec

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router



Packet Delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L: packet length (bits)
- R: link transmission rate (bps)

$$\frac{d_{trans}}{d_{trans}} = L/R$$

$$\frac{d_{trans}}{d_{trans}} \text{ and } \frac{d_{prop}}{d_{prop}}$$

$$very \text{ different}$$

 d_{prop} : propagation delay:

- *d*: length of physical link
- s: propagation speed (~2x10⁸ m/sec)
- $d_{\text{prop}} = d/s$

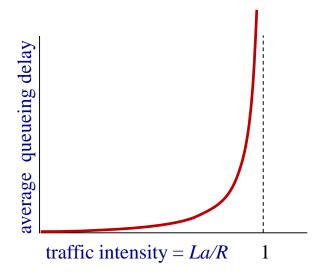
* Check out the online interactive exercises: http://gaia.cs.umass.edu/kurose ross





Queueing Delay (revisited)

- *R*: link bandwidth (bps)
- *L*: packet length (bits)
- a: average packet arrival rate
- $La/R \sim 0$: avg. queueing delay small
- *La/R* -> 1: avg. queueing delay large
- La/R > 1: more "work" arriving than can be servi ced - average delay infinite!





 $La/R \rightarrow 1$





Real Internet Delay

ping www.google.com

```
PING www.google.com (172.217.160.68): 56 data bytes
64 bytes from 172.217.160.68: icmp_seq=0 ttl=44 time=71.642 ms
64 bytes from 172.217.160.68: icmp_seq=1 ttl=44 time=82.027 ms
64 bytes from 172.217.160.68: icmp_seq=2 ttl=44 time=67.633 ms
64 bytes from 172.217.160.68: icmp_seq=3 ttl=44 time=67.492 ms
64 bytes from 172.217.160.68: icmp_seq=4 ttl=44 time=80.396 ms
64 bytes from 172.217.160.68: icmp_seq=5 ttl=44 time=83.657 ms
```

traceroute www.google.com

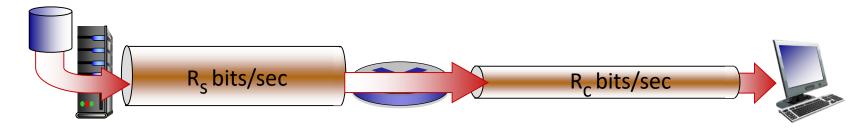
```
traceroute to www.google.com (172.217.160.68), 64 hops max, 52 byte packets
1 router.asus.com (192.168.2.1) 2.360 ms 1.030 ms 1.131 ms
2 router.asus.com (192.168.1.1) 1.624 ms 4.172 ms 1.429 ms
3 * * *
4 10.240.88.237 (10.240.88.237) 4.098 ms 3.047 ms *
5 10.204.92.105 (10.204.92.105) 2.368 ms 4.382 ms 3.328 ms
6 1.213.64.165 (1.213.64.165) 2.583 ms
   1.208.64.161 (1.208.64.161) 2.580 ms
   1.213.64.165 (1.213.64.165) 2.552 ms
7 1.214.58.177 (1.214.58.177) 9.980 ms
   1.209.58.181 (1.209.58.181) 9.874 ms
   1.214.58.205 (1.214.58.205) 9.172 ms
8 1.208.104.61 (1.208.104.61) 10.285 ms
   1.213.104.61 (1.213.104.61) 9.528 ms 10.819 ms
9 1.208.148.141 (1.208.148.141) 12.972 ms 10.580 ms 8.417 ms
10 203.233.117.81 (203.233.117.81) 10.343 ms
   203.248.208.229 (203.248.208.229) 18.325 ms
   210.120.117.113 (210.120.117.113) 9.409 ms
```



Hop-by-hop ping

Throughput and Bottleneck

- □ *Throughput:* rate (bits/time unit) at which bits transferred between sender/receiver
 - instantaneous: rate at given point in time
 - average: rate over longer period of time
 - $R_s > R_c$ What is average end-to-end throughput?



□ Bottleneck link

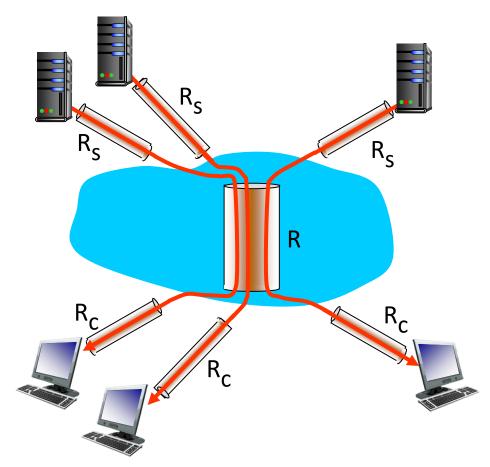
link on e2e path that constrains e2e throughput





Throughput and Bottleneck

- □ Per-connection endend throughput: $min(R_c, R_s, R/10)$
- □ In practice: R_c or R_s is often bottleneck



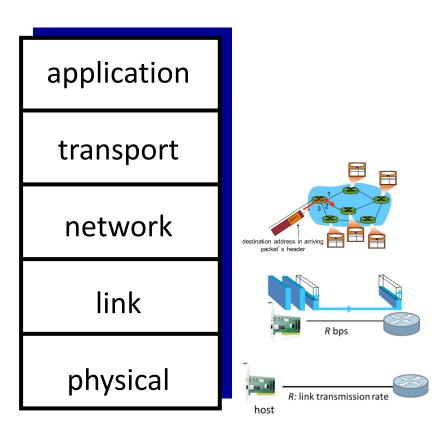
10 connections (fairly) share backbone link of *R* bits/sec





Internet Protocol Stacks (Layering)

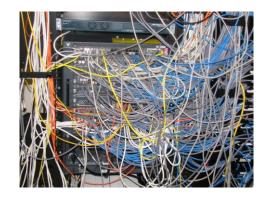
- application: supporting network applications
 - FTP, SMTP, HTTP
- transport: process-process data transfer
 - TCP, UDP
- network: routing of datagrams from source to destination
 - IP, routing protocols
- link: data transfer between neighboring network elements
 - Ethernet, 802.11 (WiFi), Bluetooth
- physical: bits "on the wire"



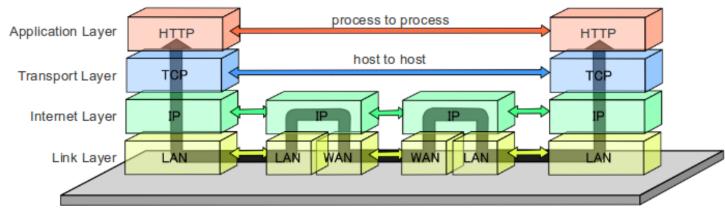




Communication over the Internet









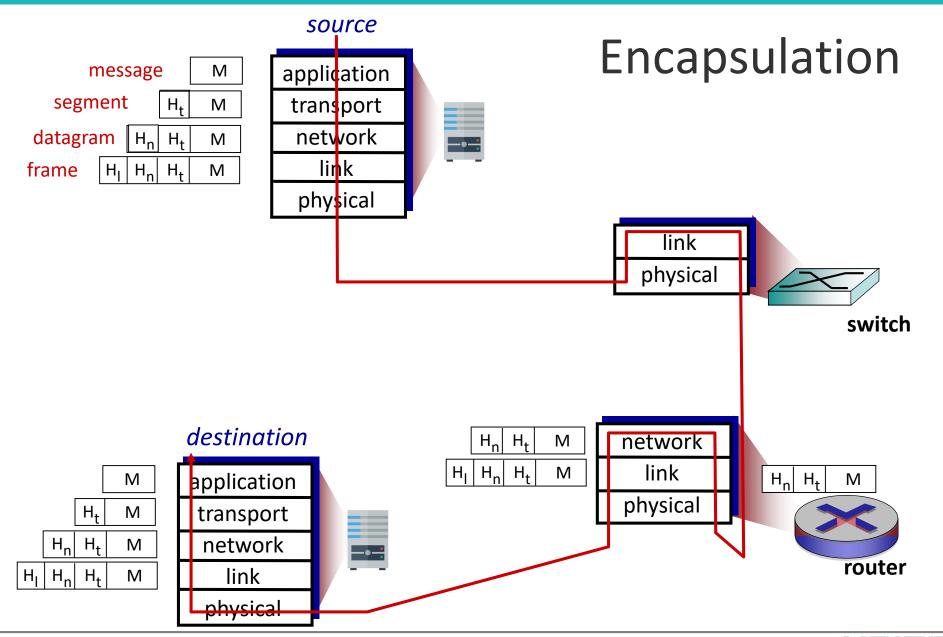








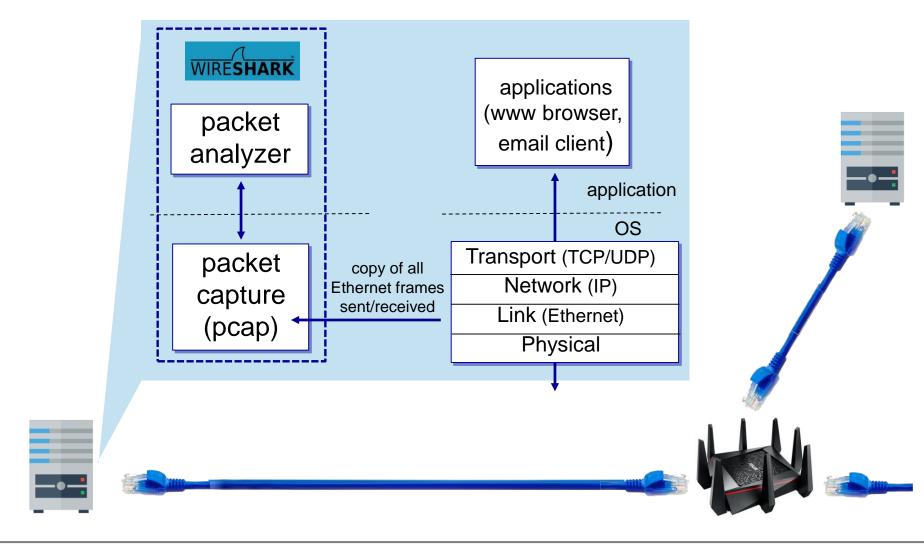








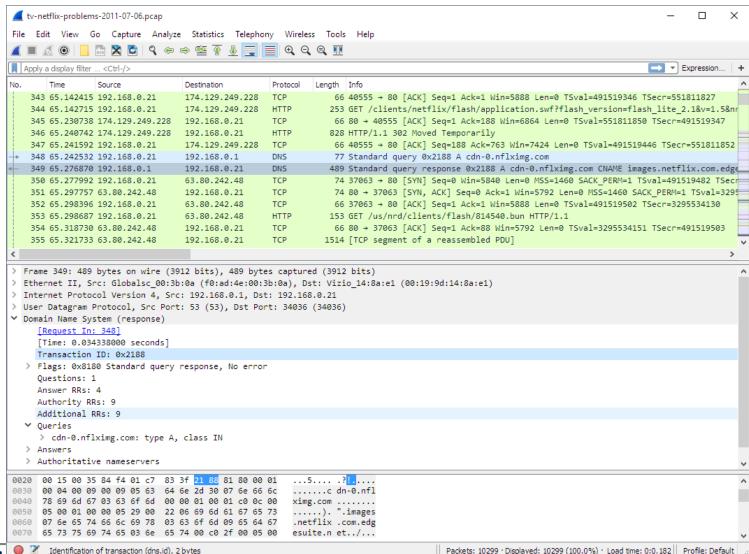
Packet Capturing







Packet Capturing





(1962-1972: Early Packet-Switching Principles)

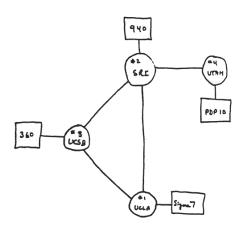
☐ 1961: Kleinrock - queueing theory shows effectiveness of packet-switching



- ☐ 1964: Baran packetswitching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

☐ 1972:

- ARPAnet public demo
- NCP (Network Control Protocol) first host-host protocol
- first e-mail program
- ARPAnet has 15 nodes



THE ARPA NETWORK





(1972-1980: Internetworking, new and proprietary nets)

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- □ late70's: proprietary architectures: DECnet, SNA, XNA
- □ late 70's: switching fixed length packets (ATM precursor)
- ☐ 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture

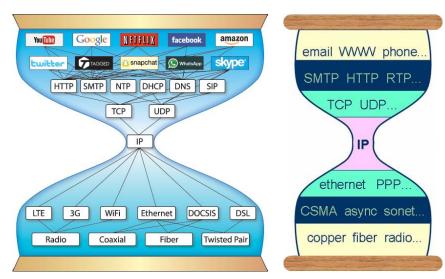




(1980-1990: new protocols, a proliferation of networks)

- ☐ 1983: deployment of TCP/IP
- ☐ 1982: smtp e-mail protocol defined
- ☐ 1983: DNS defined for nameto-IP-address translation
- 1985: ftp protocol defined
- ☐ 1988: TCP congestion control

- new national networks: CSnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks



Drawing from Johann Schleier-Smith





(1990, 2000's: commercialization, the Web, new apps)

- ☐ Early 1990's: ARPAnet decommissioned
- □ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- ☐ Early 1990's: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990's: commercialization of the Web

- ☐ Late 1990's 2000's:
 - more killer apps: instant messaging, P2P file sharing
 - network security to forefront
 - est. 50 million host, 100 million+ users
 - backbone links running at Gbps





(2005-present)

| П | ~5B devices attached to Internet (2016) |
|---|--|
| | , |
| | smartphones and tablets |
| | Aggressive deployment of broadband access |
| | Increasing ubiquity of high-speed wireless access |
| | Emergence of online social networks: |
| | ■ Facebook: ~ one billion users |
| | Service providers (Google, Microsoft) create their own networks |
| | bypass Internet, providing "instantaneous" access to search video content, email, etc. |
| | e-commerce, universities, enterprises running their services in "cloud" (e.g. Amazon FC2) |



