Coflow Scheduling of Multi-stage Jobs with Isolation Guarantee

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TABLE I
KEY TERMS AND DESCRIPTIONS

Terms	Description
M	The number of total jobs.
N	The number of total coflows.
K	The number of machines.
$F_i = \langle f_i^1, \ldots, f_i^{2K} \rangle$	Demand vector of coflow-i.
$d_i = \langle d_i^1, \ldots, d_i^{2K} \rangle$	Correlation vector of coflow-i.
$c_i = \langle c_i^1, \ldots, c_i^{2K} \rangle$	Bandwidth allocation of coflow-i.
$\overline{f_i} = \max_k f_i^k$	Bottleneck demand of coflow-i.
P_i	Progress of coflow-i.
Γ_m	Progress of job-m.

***Abstract*—This document is a model and instructions for LaTeX. This and the IEEEtran.cls file define the components of your paper [title, text, heads, etc.]. *CRITICAL: Do Not Use Symbols, Special Characters, Footnotes, or Math in Paper Title or Abstract.**

***Index Terms*—keyword, keyword, keyword**

I. INTRODUCTION

II. MODEL AND OBJECTIVE

In this section, we firstly delineate the model of datacenter networks and coflow and then discuss two objectives. To simplify the discussion, key terms used in our model are summarized in Table 1.

A. Model

Given the full bisection bandwidth, which has been well developed in modern datacenter [1], we treat the datacenter network as a big non-blocking switch connecting K machines. Each machine has one uplink port and one downlink port, thus the whole fabric has $2K$ ports. In this simplified model, the ports are the only congestion points. Hence we focus sorely on bandwidth of each port. In our analysis, all links are assumed of equal capacity normalized to one.

The coflow abstraction presents the communication demand within stages of parallel computing model. A coflow is composed of a collection of flows across a group of machines sharing a common performance requirement. The completion time of the latest flow defines the completion time of this coflow. In many data-parallel frameworks like MapReduce/Hadoop, the coflow properties, such as source, destination, amount of data transferred of each flow, are known as a priori [2]–[4].

Specifically, the coflow *demand vector* $F_i = \langle f_i^1, \ldots, f_i^{2K} \rangle$ captures the data demand of coflow-i, where f_i^k denotes the amount of data transferred on port k. Additionally, f_i^{hl} denotes a flow transferring data from port h to port l. Among all flows in coflow-i, we name the port with the largest traffic as bottleneck port. Let the data demand on this port be the *bottleneck demand*, defined as $\overline{f_i} = \max_k f_i^k$. To simplify our analysis, the *correlation vector* $d_i = \langle d_i^1, \ldots, d_i^{2K} \rangle$ is engaged to describe the demand correlation across ports, where d_i^k is the normalized data demand on port k by the bottleneck demand, i.e., $d_i^k = f_i^k / \overline{f_i}$. This vector indicates that for every byte coflow-i sends on bottleneck port, at least d_i^k bytes should be transferred on port k.

Coflows have elastic bandwidth demands on multiple ports, which means the the *bandwidth allocation vector* $c_i = \langle c_i^1, \ldots, c_i^{2K} \rangle$ of coflow-i, where c_i^k is the bandwidth share on port k, is not necessarily in the same ratio of demand vector d_i. Given the bandwidth allocation vector c_i for each coflow-i calculated by coflow scheduler given the demand vectors, the coflow progress is restricted by the worst-case port. Formally, *progress* of coflow-i is measured as the minimum demand-normalized rate allocation across ports, i.e.,

$$P_i = \min_{i:d_i^k>0} \frac{c_i^k}{d_i^k}, \tag{1}$$

Intuitively, progress of coflow-i means the smallest demand satisfaction ratio across all ports, which determines the CCT of coflow-i.

Assume a multi-stage job-m has N active coflows, i.e., $\{F_{m,1}, \ldots, F_{m,N}\}$. Given the bottleneck demand $\{\overline{f}_{m,1}, \ldots, \overline{f}_{m,N}\}$ and progress $\{P_{m,1}, \ldots, P_{m,N}\}$ of each coflow, the progress of job-m can be computed as the weighted average progress of active coflows in job-m, i.e.,

$$\Gamma_m = \frac{\sum_{n=1}^{N} \overline{f}_{m,n} P_{m,n}}{\sum_{n=1}^{N} \overline{f}_{m,n}}. \tag{2}$$

where $\overline{f}_{m,n}$ is the weight of coflow $F_{m,n}$. Like above, progress of job-m indicates the collectivity smallest demand satisfaction ratio of all coflows belonging to it, which has significant effect on the JCT of job-m.

B. Objective

In common consensus [5]–[7], a coflow scheduler focuses primarily on two objectives, average CCT and isolation guarantee. Under the multi-stage coflow scheduling problem, we should concern the average JCT instead.

1) *Average JCT*: To speed up data-parallel application completion time, as many jobs as possible should be finished in their fastest possible ways. Therefore minimizing the average JCT is settled as a critical objective for an efficient coflow scheduler.
2) *Isolation Guarantee*: In a shared datacenter network, all tenants expect *performance isolation guarantees*. Existing work has define such guarantee as the *minimum progress* across coflows [8], i.e., $\max_i P_i$. For multi-stage jobs, we define the isolation guarantee as the minimum progress across active jobs, i.e., $\max_m \Gamma_m$. To optimize the isolation guarantee, a coflow scheduler should look for an allocation to maximize the minimum progress.

However, on the application level, the effect of isolation guarantee cannot be perceived until the job is finished. If we take the fair scheme (e.g., DRF [9] or HUG [8]) as a baseline, as long as an application observes its jobs finish no later than they would have finished in the baseline algorithm, the isolation guarantee is provided in long run. Thus we introduce the job *long-term isolation guarantee* to our model.

Definition 1 (Long-term Isolation Guarantee): Consider a multi-stage job-i, let T_i be its JCT by coflow scheduler S. T_i^* is its JCT by a fair scheduler which enforce a minimum instantaneous progress of all active coflows. We call the scheduler S provides the job long-term isolation guarantee if all jobs complete no later than $T_i^* + D$, where D is a constant delay, i.e.,

$$T_i \leq T_i^* + D \tag{3}$$

In this paper, our objective is to gain the best of both targets in a long run, which is formalized as

$$\begin{aligned} \text{minimize} \quad & \sum_i T_i \\ \text{s.t.} \quad & T_i \leq T_i^* + D, \text{ for all job-}i \end{aligned} \tag{4}$$

III. Algorithm and Analysis

In this section, we designed a coflow scheduling algorithm of multi-stage jobs with isolation guarantee. We first elaborate how to sort coflows based on their priorities. Second we promote our bandwidth allocation algorithm.

A. Coflow Sorting

Before we allocate bandwidth to coflows, the priorities of each coflow should be determined. Consider all coflows are scheduled by the fair scheduler DRF, which enforces the equal progress across coflows [8], [9]. On account of the completion time of each coflow under such scheduling, we can obtain a priority order for all coflows. If a coflow completes faster in DRF scheduling, it has a higher priority than the lowers.

However, directly performing DRF is not suitable in the context of multi-stage jobs, because some coflows that have dependencies are not released at the start. At first, we can only calculate the completion times of coflows F' with no dependencies, which are released at time 0. Then, the minimum progress of these coflows $P^* = \min_i P_i'$ is produced, where P_i' is the progress of F_i'. It is noticed that in popular parallel frameworks, the number of coflows in different stages stays similar [10], [11]. We assume that P^* keeps across stages, then the completion time of coflow $F_i \notin F'$ is estimated as $t_i = P^* \overline{f_i}$.

we assume the job J_i is the i-th job to complete data transferring, i.e., when J_i completes, $J_1, \ldots, J_{i-1}$ have already completed. Let $(F_{i,1}, \ldots, F_{i,N_i})$ be the coflows constituting job J_i in the topological order and let super-coflow $S_{i,j}$ be the bonding super-coflow of $F_{i,j}$ when it finishes. To simplify our explanation,

Besides, the completion time of single coflow is not necessarily the priority of the job containing it. In multi-stage jobs, coflow pipelines are developed using directed acyclic graph(DAG) pattern to show the dependencies between coflows. In job-i, we set the weights of coflows in DAG as their completion times. Then we can get the *estimate completion time* of coflow-j $F_{i,j}$ as the maximum total weight along DAG edges from roots. A coflow with shorter estimate completion time has a higher priority. Additionally, the estimate completion time of job-m is the maximum estimate completion time of all coflows it contains. Like previous, if a job completes faster, it has a higher priority. Within a job, the priority order of coflows is the ascending order of estimate completion tiome. Therefore we get a prioritized queue of all coflows $\mathbf{F} = (F_{1,1}, \ldots, F_{1,N_1}, \ldots, F_{M,1}, \ldots, F_{M,N_M})$. The entire procedure is summarized in Algorithm 1.

Algorithm 1 Coflow Sorting Algorithm

Input: Data demand set of all coflows F and the job set J.
Output: An ordered queue of all coflows $\mathbf{F}$.

```
Sort F in the topological order of DAG.
F' ← coflows have no dependencies in F.
P* ← the minimum progress of F' under DRF.
for i = 1 → |F| do
    if F_i ∈ F' then
        t_i ← the completion time of F_i under DRF.
    else
        t_i ← P* f̄_i.
    if F_i has dependencies then
        t_i ← t_i + max{t_j | F_i depends on F_j}
T_m ← the maximum total time along DAG edges of J_m.
Sort J in the ascending order of T.
Initialize F* as an empty queue.
for m = 1 → |J| do
    F ← F + F̄_m
```

B. Bandwidth Allocation

After obtaining the prioritized queue of coflows, we can allocated bandwidth to the active coflows one by one. Specifically, we introduce the *super-coflow* conception which was proposed in [7] for our solution.

Super-coflow. Given a queue of coflows $\mathbf{F}$, the super-coflow S_i is defined as the sequential aggregation of the first i

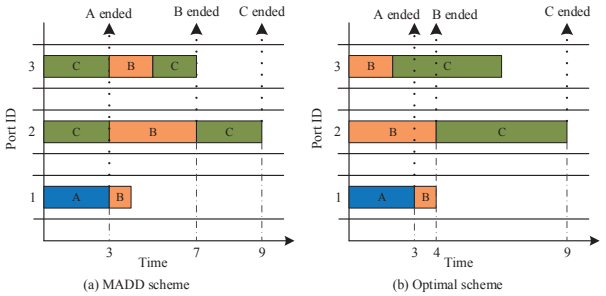

Fig. 1.

coflows $F_1, \cdots, F_i$. Additionally, we call S_i is the *bonding super-coflow* of F_i. Formally, the demand vector $D_i = \langle D_i^1, \ldots, D_i^{2K} \rangle$ of S_i is the accumulation of the first i coflows, i.e., $D_i = \sum_{j=1}^{i} d_j$. Super-coflow S_1 trivially degrades into coflow F_1 as a special case.

In previous works, coflows are mostly scheduled individually. We take the minimum-allocation-for-desired-duration (MADD) algorithm, which is a common approach for sequential coflow bandwidth allocation, as an example [2], [12]. Under MADD, the least amount of bandwidth was allocated to a coflow to obtain the maximum possible progress. Let R_k be the remaining bandwidth on port k and $d_i = \langle d_i^1, \ldots, d_i^{2K} \rangle$ be the correlation vector of coflow-i. Thus the maximum possible progress P' of coflow-i can be calculated as

$$P' = \min_{1 \le k \le 2K} \frac{R_k}{d_i^k}. \tag{5}$$

To acquiring the maximum possible progress, the bandwidth allocation of coflow-i on port k is at least $P' d_i^k$. Under MADD, the allocation is exactly $P' d_i^k$.

However, MADD may lead to the priority inversion problem. If a coflow is not able to get any progress, then it gets no bandwidth, even if allocating it some bandwidth will accelerate its completion.

On the contrary, our super-coflow conception eliminates the priority inversion situation as much as possible. When computing the bandwidth allocation, coflow F_i is justified if it contributes to super-coflow S_i. As long as F_i affects the progress of S_i, a bandwidth allocation to F_i is always guaranteed, even when it gains no progress alone, which easily triggers a priority inversion under MADD.

Assume the active coflow priority queue is $F' = (F'_1, \ldots, F'_{N'})$ and the super-coflow queue is $S = (S_1, \ldots, S_{N'})$. Our bandwidth allocation algorithm runs in turns. In the i-th turn, we will allocate the least bandwidth to F_i, while the allocations of earlier turns stay unchanged, in order to achieve the minimum possible completion time of S_i. To easily describe our algorithm, we next focus on the bandwidth allocation of F_i.

To achieve the minimum possible completion time of S_i, we first calculate the bottleneck data demand of S_i, i.e.,

$$\overline{D}_i = \max_{1 \le k \le 2K} D_i^k. \tag{6}$$

Let D_i^{hl} be the data amount transferred from port h to port l in S_i. Then at least $D_i^{hl}/\overline{D}_i$ bandwidth should be allocated for the flows between port h and port l. It is noted that the first $i-1$ coflows have received some bandwidth in the earlier turns. We denote the bandwidth received by flow f_k^{hl} in coflow-k as u_k^{hl} for $k < i$, thus we can distribute at most $(D_i^{hl}/\overline{D}_i - \sum_{k=1}^{i-1} u_k^{hl})^+$ bandwidth to f_i^{hl}, where $(x)^+ = \max(0, x)$. Given the remaining bandwidth R_h on port h and R_l on port l, which limit the actual bandwidth available for f_i^{hl}, we finally get u_i^{hl} as

$$u_i^{hl} = \min[(D_i^{hl}/\overline{D}_i - \sum_{k=1}^{i-1} u_k^{hl})^+, R_h, R_l]. \tag{7}$$

After then, the remaining bandwidth on these two ports are updated. When all turns have finished, unused bandwidth is distributed to coflows. For each ingress port h, remaining bandwidth R_h is distributed to corresponding flows according to ratio of their current received bandwidth u_i^{hl}, restricted by the corresponding egress port l. We summarize the whole procedure as SuperflowAllocation(F) in Algorithm 2.

Algorithm 2 Bandwidth Allocation Algorithm

procedure SUPERFLOWALLOCATION(Coflows F)
  Initialize unused bandwidth $R_k \leftarrow 1$ on port-k
  for $i = 1 \rightarrow |F|$ **do**
    Assemble demands $D_i = \sum_{k=1}^{i} d_k$.
    $\overline{D}_i \leftarrow \max_k D_i^k$
    for all flow $f_i^{hl} \in F_i$ **do**
      $u_i^{hl} \leftarrow \min[(D_i^{hl}/\overline{D}_i - \sum_{k=1}^{i-1} u_k^{hl})^+, R_h, R_l]$.
      $R_h \leftarrow R_h - u_i^{hl}$.
      $R_l \leftarrow R_l - u_i^{hl}$.
  Allocate remaining bandwidth to coflows in the order of F by FIFO.
procedure BANDWIDTHALLOCATION(Coflows F)
  F $\leftarrow$ sorted F by Algorithm 1.
  Initialize active coflows $F' \leftarrow \emptyset$.
  while True **do**
    $F' \leftarrow$ released coflows in the order of **F**.
    if F' is changed **then**
      Update remaining transfer demand of F'.
      SuperflowAllocation(F').
    if $F' == \emptyset$ **then**
      break.

When an old coflow is finished or a new coflow is released, the active coflow set F' is changed and the bandwidth allocations will be rescheduled. All the active coflows are sorted as the order we get by Algorithm 1 to maintain the priorities of all jobs through our scheduling. The main procedure is summarized as BandwidthAllocation(F) in Algorithm 2.

C. Long-term Isolation Guarantee

Our algorithm provides the job long-term isolation guarantee of Definition 1. For each job-i, the completion time T_i is guaranteed not to exceed a constant time over its completion time in DRF. We formalize this theorem as following:

Theorem 1 (Long-term Isolation Guarantee): Assume that jobs J are released at time 0. For all job $J_i \in J$, let T_i be the JCT of J_i in Algorithm 2, and let T_i^* be the completion time of J_i in DRF. The completion time delay is bounded as

$$T_i \leq T^* + \overline{d}_i. \tag{8}$$

Proof: To simplify our proof, we abbreviate F_{i,N_i} and S_{i,N_i} to F_i and S_i through this proof and let $F(i)$ be the coflows before F_i in our priority order. Since F_i is the last completed coflow of J_i, T_i, the JCT of J_i, equals to the completion time of F_i. Additionally we let $F_0(i)$ be the coflows that has shorter estimate completion time than F_i, i.e., coflows in $F_0(i)$ are the ones finish before F_i under DRF. We have $F(i) \subseteq F_0(i)$ because all coflows in $\{J_1, \ldots, J_i\}$ except F_i have shorter estimate completion time than F_i. To discuss T_i, we consider the following two cases.

Case 1. The bottleneck port of super-coflow S_i is kept using full bandwidth during the data transmission. In this case, the completion time of S_i is simply the time of transferring data on bottleneck port, i.e., $\overline{D}_i$. When S_i completes, F_i must have completed. Therefore we have $T_i \leq \overline{D}_i$. Then we turn to DRF. According to our algorithm, all jobs are sorted by their JCT under DRF. When J_i completes, previous $i-1$ jobs $J_1, \ldots, J_{i-1}$ must have all completed, so does S_i. It is noticed that $\overline{D}_i$ is the minimum possible completion time of S_i, i.e., $\overline{D}_i \leq T_i^*$, thus we finally have

$$T_i \leq \overline{D}_i \leq T_i^*. \tag{9}$$

Case 2. The bottleneck port of super-coflow S_i is not fully used at some time during the data transmission. Recall that the bandwidth allocation on bottleneck port is constrained by the available bandwidth on the coupled ports due to line 7 in Algorithm 2. In particular, we define port k is a coupled port of port l if there are flows in S_i transferring data between these two ports but getting less bandwidth allocation than $D_i^{kl}/\overline{D}_i$. Let B_i be the bottleneck port of S_i and let $C(B_i)$ be the set of coupled ports of B_i. Additionally, let $C_0(B_i)$ be the ports that have data transmission with B_i in S_i.

Let t_B be the time when port B_i starts to get full bandwidth allocation until the completion of S_i. Specifically, let t_k be the time when port $k \in C(B_i)$ finished data transmission of all the coflows in $F(i)$. Under our algorithm, the lacking of bandwidth utilization on B_i can only occurs when bandwidth of ports in $C(B_i)$ are all fully used, otherwise our algorithm will tempt to allocate the spare bandwidth to a flow between such coupled port and B_i. Therefore we have

$$t_B \leq \max_{k \in C(B_i)} t_k. \tag{10}$$

Since the bandwidth allocation of each port varies throughout the data transmission because of completions of coflows, we denote $a_i^{kl}(t)$ and $b_i^{kl}(t)$ as the expected and actual bandwidth allocation of the flow transferring between port k and port l at time t. In particular,

$$a_i^{kl}(t) = D_i^{kl}(t)/\overline{D}_i(t); \tag{11}$$

$$b_i^{kl}(t) = u_i^{kl}(t) + \sum_{j \in F(i)} u_j^{kl}(t). \tag{12}$$

We notice that during the transmission of S_i, the actual bandwidth allocation on some ports may continues being less than its expected bandwidth when S_i is running alone under DRF. In some extreme situations, the bottleneck port of S_i may change from B_i to another port k if port k keeps getting less bandwidth allocation than expected, which we call the bottleneck port of S_i *shifts*. To bound our JCT of J_i, we next discuss two sub-cases, differentiated by whether the bottleneck port of S_i shifts.

Sub-Case 1. The bottleneck port B_i of S_i does not *shifts* during the data transmission. In the worst case, bottleneck port B_i can only get full bandwidth allocation when all previous coflows finish their data transmissions on the coupled ports of B_i, which means $t_B = \max_{k \in C(B_i)} t_k$. We have the JCT constraint of J_i as

$$T_i \leq \overline{D}_i + \sum_{k \in C_0(B_i)} \int_0^{t_k} (a_i^{kB_i}(t) - b_i^{kB_i}(t))dt. \tag{13}$$

Since S_i is the aggregation of $F(i)$ and F_i, we can split $\overline{D}_i$ into two parts as

$$\overline{D}_i = d_i^{B_i} + \mathbf{D}_i^{B_i}, \tag{14}$$

where $d_i^{B_i}$ is the data demand of F_i on port B_i and $\mathbf{D}_i^{B_i}$ is the sum of data demands of coflows in $F(i)$ on B_i. Thus we have

$$\mathbf{D}_i^{B_i} = \sum_{k \in C_0(B_i)} \int_0^{t_k} (b_i^{kB_i}(t) - u_i^{kB_i}(t))dt. \tag{15}$$

Additionally we have

$$u_i^{kB_i}(t) \geq 0. \tag{16}$$

By combining Eq.(14)(15)(16), we can transform Eq.(13) into

$$T_i \leq d_i^{B_i} + \sum_{k \in C_0(B_i)} \int_0^{t_k} a_i^{kB_i}(t)dt. \tag{17}$$

Then we first focus on the ports $k \in C(B_i)$. Under DRF, when coflow F_i finishes, all coflows in $F(i)$ must have finished. Since port k keeps using full bandwidth, F_i will not finish before t_k. Thus we have

$$\max_{k \in C(B_i)} t_k \leq T_i^*. \tag{18}$$

Second we discuss the ports $k \in C_0(B_i) - C(B_i)$. On these ports, the actual bandwidth allocations are not greater than the expected allocations, i.e.,

$$a_i^{kB_i}(t) \leq b_i^{kB_i}(t). \tag{19}$$

Combining Eq.(9)(11)(19),we have

$$\max_{k \in C_0(B_i) - C(B_i)} t_k \leq \overline{D}_i \leq T_i^*. \tag{20}$$

Combining Eq.(18)(20), we can bound the $t_i^{\max}$ as

$$t_i^{\max} = \max_{k \in C_0(B_i)} t_k \leq T_i^*. \tag{21}$$

Note that the bandwidth allocation on B_i cannot exceed 1, i.e.,

$$\sum_{j \in C_0(B_i)} a_i^{kB_i}(t) \leq 1. \tag{22}$$

According to Eq.(21)(22) we transform Eq.(17) into

$$\begin{aligned} T_i &\leq d_i^{B_i} + T_i^* \\ &\leq \overline{d}_i + T_i^*. \end{aligned} \tag{23}$$

Sub-Case 2 The bottleneck port of S_i *shifts* to $B_i^1, B_i^2, \ldots, B_i$ in order. We denote the time when bottleneck port shifts to B_i as t^*. Consider the worst case that before t^* there exists no bandwidth allocation on B_i. Similar to Eq.(17), we have

$$\begin{aligned} T_i &\leq t^* + \overline{D}_i + \sum_{k \in C_0(B_i)} \int_{t^*}^{t^k} (a_i^{kB_i}(t) - b_i^{kB_i}(t)) dt \\ &\leq t^* + (d_i^{B_i} + \mathbf{D}_i^{B_i}) + \sum_{k \in C_0(B_i)} \int_{t^*}^{t^k} (a_i^{kB_i}(t) - b_i^{kB_i}(t)) dt \\ &\leq d_i^{B_i} + t^* + \sum_{k \in C_0(B_i)} \int_{t^*}^{t^k} a_i^{kB_i}(t) dt. \end{aligned} \tag{24}$$

By combining Eq.(22), similar to the previous analysis, we finally transform Eq.(24) to

$$\begin{aligned} T_i &\leq d_i^{B_i} + t^* + (t_i^{\max} - t_*) \\ &\leq d_i^{B_i} + t_i^{\max} \\ &\leq \overline{d}_i + T^*. \end{aligned} \tag{25}$$

■

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