



SWINBURNE
UNIVERSITY OF
TECHNOLOGY

COS10004 Computer Systems

Lecture 8.3 LED Flash (part 3) – a better busy wait (the Timer Register)

CRICOS provider 00111D

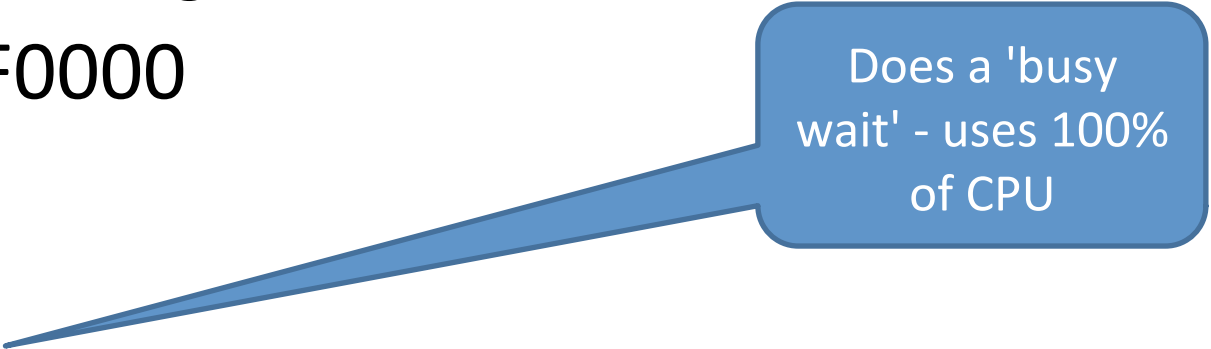
Dr Chris McCarthy

INSERT DELAY WITH BUSY WAIT TIMER

- Pseudocode:
 - Program GPIO18 LED for output
 - Loop1:
 - Turn LED on (pull GPIO18 high)
 - Busy wait
 - Turn LED off (pull GPIO18 low)
 - Busy wait
 - branch to loop1

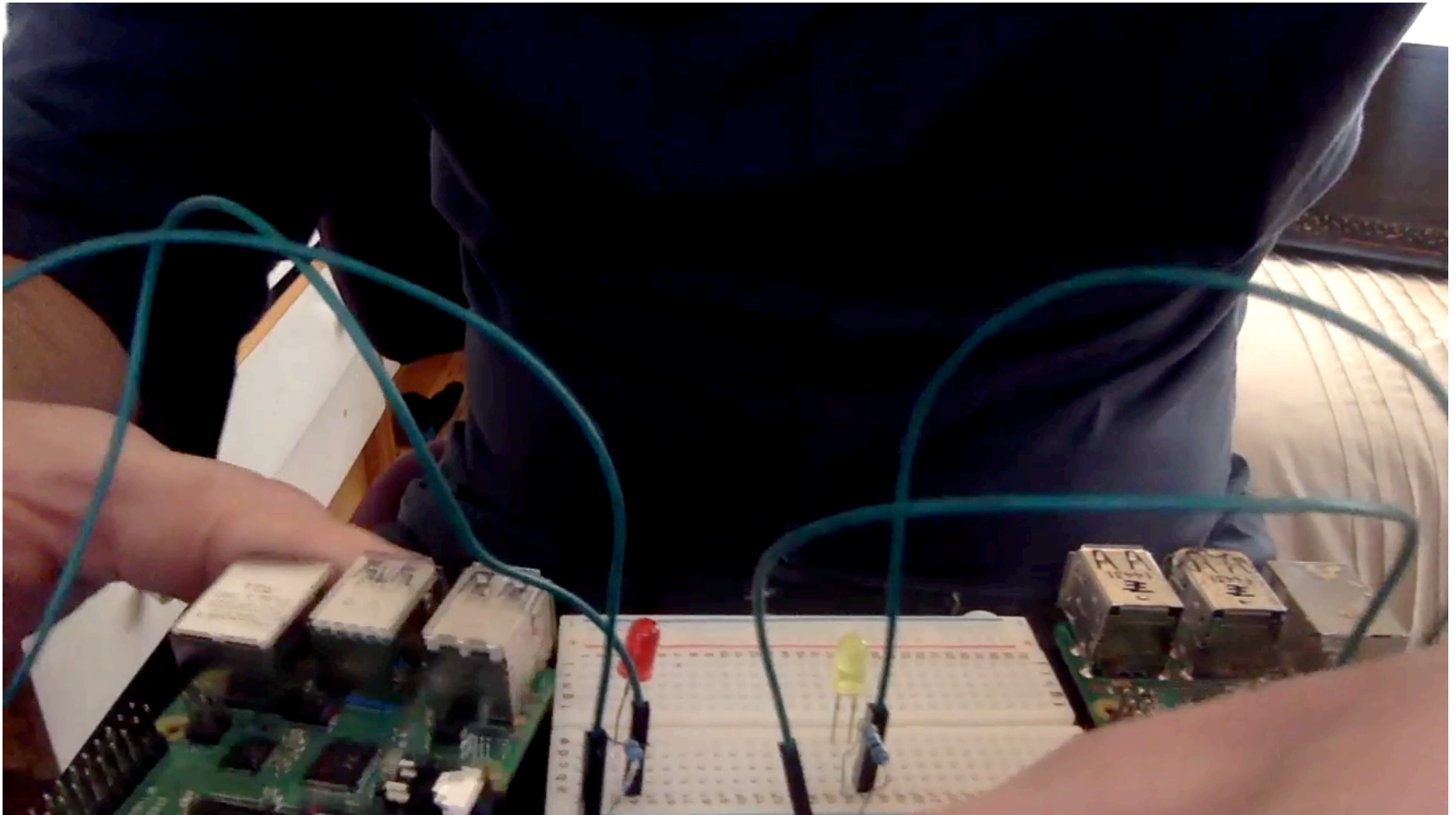
A DUMB TIMER

- Variables:
 - r0 = GPIO base address
 - r1 = working memory (for setting bits, registers)
 - use r2 for timing
- ```
mov r2,$3F0000
loop1:
 sub r2,#1
 cmp r2,#0
 bne loop1
```



Does a 'busy wait' - uses 100% of CPU

# OK2 RUNNING ON PI 2B AND PI 4B



# A BETTER TIMER

- RPi has a dedicated timer register:
  - independent of clock speed.
  - housed inside the same chip as the ARM CPU, the GPIO, GPU, RAM and most other things.
  - This chip is called the SoC (System on a Chip).
- The Timer registers start at  
BASE address + 0x3000
  - Timer counts 1 microsecond intervals ( $10^6$  per second)

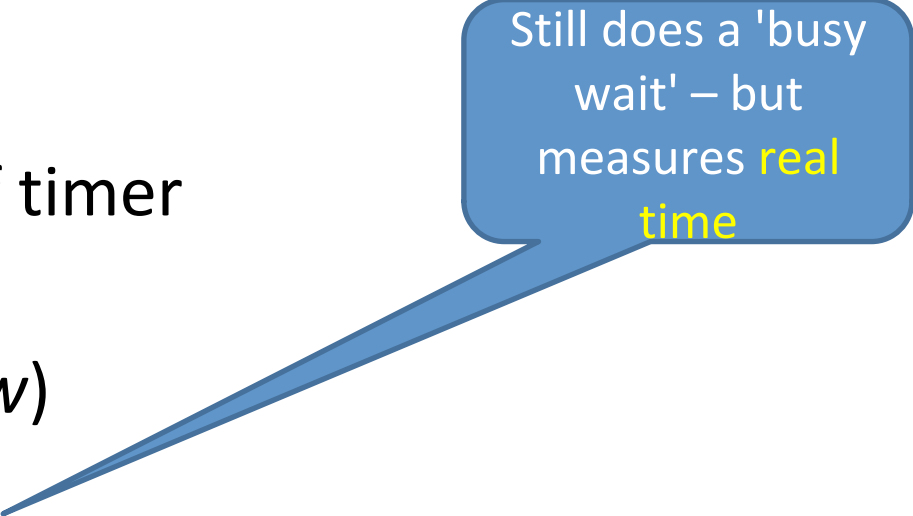
# A BETTER TIMER

- The RPi timer registers:

| Byte offset<br>(from BASE) | Size /<br>Bytes | Name             | Description                                                          | Read or<br>Write |
|----------------------------|-----------------|------------------|----------------------------------------------------------------------|------------------|
| 0x3000                     | 4               | Control / Status | Register used to control and clear timer channel comparator matches. | RW               |
| 0x3004                     | 8               | Counter          | A counter that increments at 1MHz.                                   | R                |
| 0x300C                     | 4               | Compare 0        | 0th Comparison register.                                             | RW               |
| 0x3010                     | 4               | Compare 1        | 1st Comparison register.                                             | RW               |
| 0x3014                     | 4               | Compare 2        | 2nd Comparison register.                                             | RW               |
| 0x3018                     | 4               | Compare 3        | 3rd Comparison register.                                             | RW               |

# PSEUDOCODE


- pseudocode:
  - store base address of timer
  - store delay
  - mov start\_time (=now)
  - loop:
    - read *now*
    - remaining\_time = (*now* – starttime)
    - compare remaining\_time, delay
    - loop if remaining\_time <= delay



Still does a 'busy wait' – but measures **real time**

# READING THE TIME (TIMESTAMP)

- Variables:
- r3 - base address of timer
- r4 - desired delay
- r6, r7 - where the current time will be stored
- r5 - start time
- r8 - elapsed time (r6-r5)



Timer returns  
64-bit time, but  
we only need  
the lower  
register (r6) -  
**must be an even  
register**



# WITH REGISTERS

store base address of timer (r3)

store delay (r4)

mov start\_time (r5)(=current\_time (r6))

timerloop:

    read current\_time (r6)

    remaining\_time (r8)= current\_time (r6) –  
    start\_time (r5)

    compare remaining\_time (r8), delay (r4)

    loop if LE (remaining\_time <= delay)

# RECALL LOAD REGISTER ?

delimiter

ldr r0,[r1]

register 0

r1 must contain a  
pointer to a value in  
memory

- load register 0 with the value pointed to by r1

# A BETTER TIMER

- BASE = \$FE000000  
TIMER\_OFFSET = \$3000  
mov r3,BASE  
orr r3,TIMER\_OFFSET ;base address of timer
- ldr<sup>d</sup> r6,r7,[r3,#4] ;copy timestamp [4 bytes after r3 (low byte), r7(high byte)].
  - d suffix means double-word (64 bit)
  - <http://infocenter.arm.com/help/index.jsp?topic=/com.arm.doc.ddi0338g/Chddeedh.html>
- mov r4,\$800000 ;0.524sec

<sup>d</sup> suffix  
means  
double-word  
(64 bits)

## WITH CODE (4B)

```
BASE = $FE000000
TIMER_OFFSET = $3000
mov r3,BASE
orr r3,TIMER_OFFSET ;store as
mov r4,$80000 ;store delay (r4)
ldrd r6,r7,[r3,#4]
mov r5,r6 ;mov starttime (r5)=(currenttime (r6))
timerloop:
```

- $r3 = \text{BASE} + \text{TIMER\_OFFSET} + 4$   
(0x3F003004)
  - $[r3, \#4]$  means  
value at ((address in r3) + 4)

can't re-use loop label - each  
must be unique: loop1,  
loop2, loop3...

```
ldrd r6,r7,[r3,#4] ;read currenttime (r6)
sub r8,r6,r5 ;remainingtime (8)= currenttime (r6) - starttime (r5)
cmp r8,r4 ;compare remainingtime (r8), delay (r4)
bls timerloop ;loop if LE (remainingtime <= delay)
```

BASE = \$20000000 ;B B+. Use \$3F000000 for 2  
GPIO\_OFFSET = \$200000

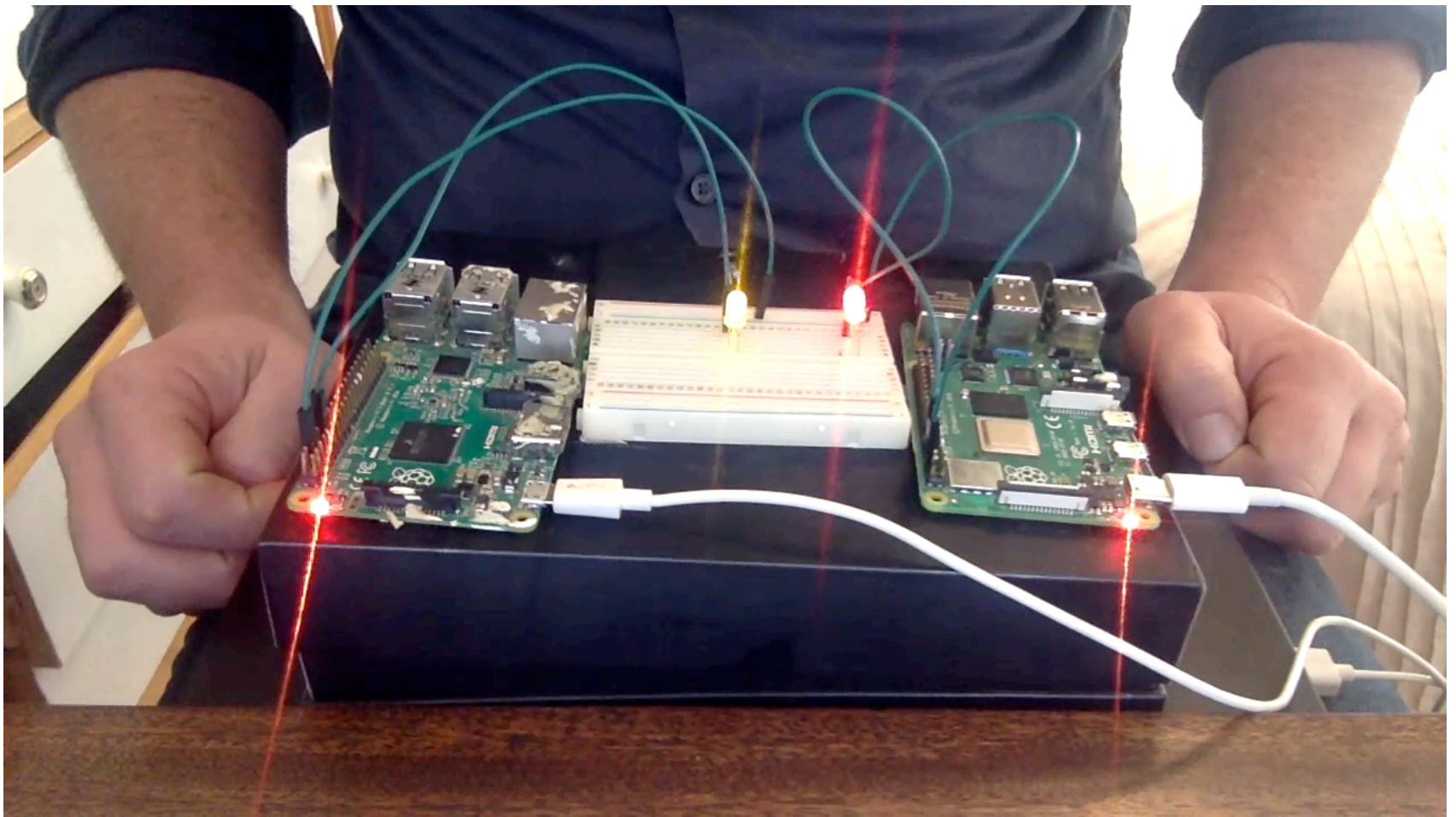
# ALL THE CODE (B)

```
mov r0,BASE
orr r0,GPIO_OFFSET ;Base address of GPIO

mov r1,#1
lsl r1,#18 ;B
str r1,[r0,#4] ;enable output
mov r1,#1
lsl r1,#16
loop$:
 str r1,[r0,#40] ;Turn on LED
 ;new timer
TIMER_OFFSET = $3000
mov r3,BASE
orr r3,TIMER_OFFSET ;store base address of timer (r3)
 mov r4,524288;store delay (r4)
 ldrd r6,r7,[r3,#4]
 mov r5,r6 ;store starttime (r5)(=currenttime (r6))
loop1:
 ldrd r6,r7,[r3,#4] ;read currenttime (r6)
 sub r8,r6,r5 ;remainingtime (8)= currenttime (r6) - starttime (r5)
 cmp r8,r4 ;compare remainingtime (r8), delay (r4)
 bls loop1 ;loop if LE (remainingtime <= delay)
 str r1,[r0,#28] ;turn off LED
 ;re-use timer
 ldrd r6,r7,[r3,#4]
 mov r5,r6 ;store starttime (r5)(=currenttime (r6))
loop2:
 ldrd r6,r7,[r3,#4] ;read currenttime (r6)
 sub r8,r6,r5 ;remainingtime (8)= currenttime (r6) - starttime (r5)
 cmp r8,r4 ;compare remainingtime (r8), delay (r4)
 bls loop2 ;loop if LE (remainingtime <= delay)
b loop$
```

# ALL THE CODE (2)

```
BASE = $FE000000 ;4B
GPIO_OFFSET = $200000
mov r0,BASE
orr r0,GPIO_OFFSET ;Base address of GPIO
mov r1,#1
lsl r1,#24
str r1,[r0,#4] ;enable output
mov r1,#1
lsl r1,#18
loop$:
 str r1,[r0,#28] ;Turn on LED
 ;new timer
TIMER_OFFSET = $3000
mov r3,BASE
orr r3,TIMER_OFFSET ;store base address of timer (r3)
mov r4, 524288 ;store delay (r4)
ldrd r6,r7,[r3,#4]
mov r5,r6 ;store starttime (r5) (=currenttime (r6))
loopt1:
 ldrd r6,r7,[r3,#4] ;read currenttime (r6)
 sub r8,r6,r5 ;remainingtime (8)= currenttime (r6) - starttime (r5)
 cmp r8,r4 ;compare remainingtime (r8), delay (r4)
 bls loopt1 ;loop if LE (remainingtime <= delay)
 str r1,[r0,#40] ;turn off LED
 ;re-use timer
 ldrd r6,r7,[r3,#4]
 mov r5,r6 ;store starttime (r5) (=currenttime (r6))
 loopt2:
 ldrd r6,r7,[r3,#4] ;read currenttime (r6)
 sub r8,r6,r5 ;remainingtime (8)= currenttime (r6) - starttime (r5)
 cmp r8,r4 ;compare remainingtime (r8), delay (r4)
 bls loopt2 ;loop if LE (remainingtime <= delay)
 b loop$
```



# USING THE TIMER REGISTER

- An absolute timer (no longer dependent on processor speed!)
- Our approach is still a form of busy wait though !
  - CPU still occupied for duration of delay
  - Essentially this is polling the timer
- Is there a potentially even better approach ?
  - Well yeah! We could implement an interrupt based timer
  - More complicated though (and beyond)



# SUMMARY

- The timer register provides an accurate timer that can easily be accessed
- We read it using `ldrd`
  - a version of load register that reads 64 bit words into two specified 32 bit registers
- Our timer is better than before, but is still a “busy wait” timer
  - Interrupt based timing would solve this