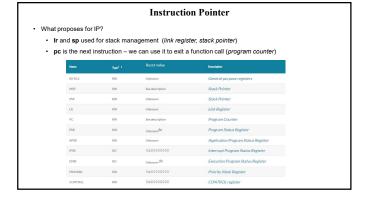


Instruction Pointer • The running a program is managed by Instruction pointer (IP) • How does the Instruction pointer work? Address int main() { 0x0001 int a, b, c; Instruction pointer (IP) stores the address of command 0x0004 a=5; Instruction pointer (IP) stores the address of command 0x0008 b=6; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction pointer (IP) stores the address of command 0x00012 c=a+b; Instruction 0

Instruction Pointer									
	unning a program is mai does the instruction poin	naged by Instruction pointer (IP) nter work?							
Store a,b,c, current IP (1)									
Address	int main()	to memory and go to the int add(int a, int b)	Address						
0x0001 0x0004 0x0008 0x00012	int a, b, c; a=5; b=6; c=a+b; add(a,b);	P	0x00A1 0x00A4 0x00A8						
	c=15; ←								

Instruction Pointer							
What p	roposes for IP?						
• Ir a	and sp used for	r stack management (link register, stack pointer)					
• pc	is the next inst	ruction – we can use it to exit a function call (program counter)					
	R0	→ Application Binary Interface (ABI) sets standard way of using ARM registers.					
	R1						
	R2						
Low registers <	R3						
	R4						
	R5	General p urpose r egisters					
	R6						
_	R7						
	R8						
	R9						
High registers <	R10						
	R11						
	R12	l J					
Stack Pointer	SP (R13)	PSP MSP					
Link Register	LR (R14)						
Program Counter	PC (R15)						
	PSR	Program Status Register					
	PRIMASK	Interrupt mask register Special registers					
	C ONTROL	Control Register					
		https://developer.arm.com/documentation/dui0497/a/CHDBIBGJ					



Instruction Pointer • What proposes for IP? • Ir (link register) contains the address of the next instruction after a function call. • We use this to tell the code what to run after a function finishes. • The current address of code to be run is stored in the program counter (pc). Setting this to the value in Ir makes the program resume after a function has finished. FunctionLabel: • Alternatively (better) FunctionLabel: • Calling function: bi FunctionLabel push {tr} do something → Maybe Ir is changed → store it in the stack pop {pc}

Instruction Pointer push {Ir} mov r3,\$3F000000 ; RPi2 and 3 orr r3,\$00003000 mov r4.\$80000 :~0.5s Idrd r6,r7,[r3,#4] mov r5,r6 ;label still has to be different from all the others loopt1: Idrd r6,r7,[r3,#4] sub r8,r6,r5 cmp r8,r4 bls loopt1 ;branch if lower or same (<=) pop {pc} ;return ;; -> safety

Delay:		
mov r3,\$3F000000	; RPi2 and 3	
orr r3,\$00003000		
mov r4,\$80000	;~0.5s	
ldrd r6,r7,[r3,#4]		
mov r5,r6		
loopt1:	;label still has to be different from all the others	
Idrd r6,r7,[r3,#4]		
sub r8,r6,r5		
cmp r8,r4		
bls loopt1	;branch if lower or same (<=)	
bx Ir	;branch to Ir without updating PC	
;; -> This way works best with	the FASMARM compiler	

NEW COMMAND

For details

B loopA ; Branch to loopA

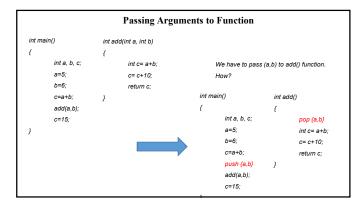
 $\bullet \quad BL \quad \text{ funC} \qquad ; Branch \ with \ link \ (Call) \ to \ function \ funC, \ return \ address \ stored \ in \ LR$

BX LR ; Return from function call

- BLX R0 ; Branch with link and exchange (Call) to a address stored in R0 $\,$

 $\bullet \ \ BEQ \quad labelD \ \ ; Conditionally \ branch \ to \ labelD \ if \ last \ flag \ setting \ instruction \ set \ the \ Z \ flag, \ else \ do \ not \ branch.$

Function ;;; test.asm mov r0,\$3F000000 mov r3,\$3F000000 orr r0,\$00200000 orr r3,\$00003000 mov r1,#1 mov r4,\$80000 Isl r1,#24 ;GPIO18 Idrd r6,r7,[r3,#4] str r1,[r0,#4] mov r5,r6 mov r1,#1 loopt1: Isl r1,#18 Idrd r6,r7,[r3,#4] loop\$: sub r8,r6,r5 str r1,[r0,#32] cmp r8,r4 bl Delay ;call Delay bls loopt1 str r1,[r0,#44] ;off bl Delay ;call Delay ;; save to timer3.asm b loop\$ include "TIMER3.asm"



v c {	oid Increment(counter[], digitIdx)	Increment: ;r0 = counter_address ;r1 = digit ;r2 = maxDigit ;r3 = radix-1 mov r4,r1 ;copy for later to a temp variable cmp r1,r3; → cmp r1,r2	
	if (digitIdx <= maxDigitIdx)	return	, a dgit maxagit
	{ if (counter[digitldx] == radix-1) //carry { counter[digitldx]=0; //carry digitldx+1); //carry digitldx+1);	beq end;; ->bls end cmp r0[r1], r3 bne continue ;carry mov r0[r1], #0 add r4,#1 push (lr)	;if this digit != radix-1 (e.g. 9) ;just add 1 (increment) ;reset this counter ;add 1 to copy of digit
	} else {	;backup Ir (we'll need it later when the next line returns)	
}	counter[digitIdx]++; //increment } }	bl Increment pop (ir) b end ;all done continue: add r0[r1], #1 end: bx Ir	;increment ;call display function here

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