CS3310 Data and File Structures with Java Instructor: Ajay K Gupta Lab TA: Yu Guo, Zijiang J Yang

Assignment: 4

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<u>Software Life Cycle Report – Assignment 4</u>

Phase 1: Specification

- Practice building trees and heaps
- Practice traversing trees and heaps
- Develop high-performance solutions

Program Output is at the end of this file.

Phase 2 Design:

Modules and Basic structure

Module 1: Class Main

Fields:

- 1. Int timesToRun controls the average time
- 2. ArrayList<Student> holds the original file data
- 3. Node minHeap head node to the min heap
- 4. BST binarySearch Object that holds the binary search tree
- 5. MaxHeap heap object that holds the max heap
- 6. Long start, end, total holds the run time of each search
- 7. Long[] totalArr holds the cumulative run time of each search
- 8. Int index used to copy data from the original to new data structures
- 9. String first, last the first and last name to search for
- 10. Queue queue used for breadth first traversal
- 11.Stack stack used for depth first traversal
- 12. Node found holds the data if the node is found in the search

Methods:

- Void printNodeData(Node data) prints the wanted information for the found node – parent, level, position in level, left / right child
- 2. String getSearchName() used if the user wants to input a different name other than the default
- 3. Void printAverage(long[] array, int timesRun) prints the average run time for the searches

Module 2: BinaryTree

Fields: none Methods:

- 1. Void insert(Node head, Node insert) inserts the insert Node using the head Node as the root
- 2. Void increment(Node leaf) starts at the given leaf node and increments the size of each node above it until the root is found
- 3. Void decrement(Node leaf) starts at the given leaf node and decrements the size of each node above it until the root is found
- 4. Void minHeap(Node leaf) heapify for min heaps
- 5. Void minHeapTopDown(Node head) start at the root node and heapify from the root node down to the leaves
- 6. Void printPreOrder(Node head) prints the tree in pre-order starting at the given head node
- 7. Void printPostOrder(Node head) prints the tree in post-order starting at the given head node
- 8. Void printInOrder(Node head) prints the tree in in-order starting at the given head node
- 9. Student delete(Node head) deletes the given node from the tree and rearranges the children nodes. The deleted data is returned
- 10. Void enqueue(Node head, Queue insert) inserts the tree nodes to the given queue starting at the head node as root
- 11.Void enqueue(MaxHeap maxHeap, Queue insert, int start) overloaded method to insert the nodes into the queue using the maxHeap.heap array
- 12. Void pushStack(Node head, Stack stack) pushes the tree nodes to the given stack starting at the head node as root
- 13. Node breadthFirst(Queue queue, String last, String first) performs a breadth first search using the queue comparing data to the first and last name
- 14. Node depthFirst(Stack stack, String last, String first) performs a depth first search using the stack comparing data to the first and last name
- 15. Void setPosition(Node head, HashMap<Integer, Integer> map) sets the position of the nodes in the tree using the HashMap

Module 3: BST

Fields:

- 1. Private Node head head of the BST tree
- 2. Private HashMap<Integer, Integer> map holds the positions of the nodes

Methods:

- 1. Void insert(Node head, Student insert, int level) inserts the new Student data into the tree using the head Node as the root. Level is passed recursively to set the level of the node
- 2. Student delete(Node head, String name) deletes the node with the given name and returns the Student object that was deleted
- 3. Node getHead() returns the head of the tree
- 4. Node setHead(Node head) sets the head node of the tree
- 5. Void setPosition(Node head) traverses the tree and sets the position of each node in the level that it is in starting at the head node as the root
- 6. Node searchLast(Node head, String last, String first) searches the tree for the given last and first name

Module 4: MaxHeap

Fields:

- 1. Node[] heap holds a max heap using an array
- 2. Int last keeps track of the last node index in the array

Methods

- 1. MaxHeap(ArrayList<Student> list) initializes the array using the size of the array list. Sets the last index to 0
- 2. Void insert(Student record) inserts the new student data as a node into the array
- 3. Void maxHeapify(int last) heapify the last node data
- 4. Void printlnOrder(int print) prints the nodes of the array in in-order of their tree traversal
- 5. Void printPreOrder(int print) prints the nodes of the array in pre-order of their tree traversal
- 6. Void printPostOrder(int print) prints the nodes of the array in post-order of their tree traversal
- 7. Node delete() deletes that last node in the array and returns it
- 8. Void heapify(int index) heapify starting at the given node

Module 5: Node

Fields:

- 1. Private Student data the first and last name of the student
- 2. Private int size the size of the sub trees + 1
- 3. Private Node parent Node pointer to the parent node
- 4. Private Node leftChild Node pointer to the left child
- 5. Private Node rightChild Node pointer to the right child

- 6. Int level holds the level the node was inserted at
- 7. Int position holds the position in the level relative to siblings

Methods:

- Node(String head) makes the node the root not by setting parent to null and size to 1
- 2. Getters and Setters for all the fields

Module 6: NodeQueue

Fields:

- 1. Private NodeQueue next the next Node in the queue
- 2. Private Node data the data held

Methods:

1. Getters and setters for the fields

Module 7: NodeStack

Fields:

- 1. Node data the data held
- 2. NodeStack next the next node in the stack

Methods:

1. Getters and setters for the fields

Module 8: Queue

Fields:

- 1. Private NodeQueue head the head node to the queue
- 2. Private NodeQueue tail the tail node to the queue

Methods:

- 1. Getters and setters for the fields
- 2. Void add(Node newNode) adds the new Node to the end of the queue
- 3. NodeQueue delete() deletes the head node of the queue

Module 9: ReadFile

Fields:

- 1. Static Scanner scan Used to read data
- 2. Static String fileName the file to read the data in from

Methods:

- Void readFile(Node headNode, ArrayList<Student> list) reads from the filename field and saves the data to the headNode tree and also the list Array list
- Void getFile() asks the user for a new data file if the default file is not found

Module 10: Stack

Fields:

- Private NodeStack head the head node of the stack
- 2. Private NodeStack tail the tail node of the stack

Methods:

- 1. Void insert(Node newNode) inserts the new node to the head of the stack
- 2. NodeStack delete() deletes the head node of the stack and returns it
- 3. Getters and setters for the fields

Module 11: Student

Fields:

- 1. Private String last The last name of the student
- 2. Private String first The first name of the student

Methods:

1. Getters and setters for the fields

Pseudocode

Class - Main

Main()

Make a variable to hold the number of times the program is run

Create an array list to hold the original data

Create a Node minHeap and a BST binarySearch object

Create 3 variables to hold the time data – start, end, total

Create an array of longs to hold the cumulative time for the searches

Call ReadFile.getFile() to set the data file

Call ReadFile.readFile() to read the data in and set the array list and minHeap data

Iterate the array list and set the Max heap data

Call binarySearch.setPosition to set the position of the nodes relative to their siblings

Call BinaryTree.setPosition to set the position of the minHeap relative to their siblings

Set the first and last name to search for – there is a method if this should be dynamic

Name = getSearchName()

Run a for loop for the designated number of runs

Enqueue the data from the MaxHeap

Call BinaryTree.breadthFirst to search the MaxHeap using breadthFirst search

Print the data if it is the first time the program was run

Add the total time to the array

Enqueue the minHeap

Call BinaryTree.breadthFirst to search the MinHeap using breadthFirst search

Add the total time to the array

Push the MaxHeap nodes to the stack
Call BinaryTree.depthFirst() to search the MaxHeap using depth search
Add the total time to the array
Push the MinHeap to the stack
Call BinaryTree.depthFirst() to search the MinHeap using depth search
Add the total time to the array
Search the binary search tree – binarySearch.searchLast()
Add the total time to the array

After the number of runs are done – print the average time printAverage()

printNodeData()

print the data from the given node

getSearchName

Initialize a scanner for System.in reading use a try / catch to ask the user for a name to search the trees for return the String from the user

Class – BinaryTree

Insert()

Create a temp node to iterate the list until a child is null
Use a while loop to traverse the tree until a child is null
If the left child is null — set the new node to the left child's next
If the right child is null — set the new node to the right child's next
Call increment() to add 1 to the size of all the parent nodes
Call minHeap() to move the data and maintain a MinHeap

Increment()

While the current node.parent != null Add one to the size

Decrement()

While the current node.parent != null Subtract one from the size

minHeap()

create a temp node to traverse the tree while the temp node's parent is not null and the data is less than the parent data swap the two student objects

```
minHeapTopDown()
       Create a temp node
       Create a temp student object to hold the head data
       If the head node's children are both null, then it is a leaf node
               Return
       Iterate the tree until the head node's children are null
printPreOrder()
       traverse the list
       visit the node
       if != null call printPreOrder() on the node.leftChild
       if != null call printPreOrder() on the node.rightChild
printPostOrder()
       traverse the list
       if != null call printPreOrder() on the node.leftChild
       if != null call printPreOrder() on the node.rightChild
       visit the node
printlnOrder()
       traverse the list
       if != null call printPreOrder() on the node.leftChild
       visit the node
       if != null call printPreOrder() on the node.rightChild
delete()
```

Create a temp node

Create a temp student object

Iterate the list until either of the children are null

Iterate the tree checking the size of each sub tree

Swap this node with the root node

Delete the new leaf node

Heapify from the top down by calling minHeapTopDown on the new root

Return the deleted leaf node data

Enqueue() - tree

If head == null then the tree is empty or a leaf node is found

Return

Else

Insert the head node to the queue Recursively call enqueuer on the left and right child

```
Enqueue() – array
       Add the node data to the queue
       Recursively call the enqueue() until the index is out of bounds for the array
pushStack() - tree
       if the head == null, then the tree is empty or a leaf node is found
       else
               insert the head node to the stack
               recursively call the pushStack with the left and right child
pushStack() - array
       run a for loop to push every node in the array to the stack
breadthFirst()
       if the queue.getHead() == null, then the queue is empty
               return
       else
               test the head data with the first and last name, if equal return the data
               else
                      delete the head and move to the next node in the queue
depthFirst()
       if the stack.getHead() == null, then the stack is empty
               return
       else
               test the head data with the first and last name, if equal return the data
               else
                      delete the head node and move to the next node in the stack
setPosition()
       if the current node has both children as null
               if the hash map doesn't have a key for the current level
                      set the level -> position as 1
                      set the current node position to 1
                      iterate the position
               else
                      get the position and set the current node to it
                      iterate the position
       if the left child is not null
               call setPosition on the left child node
       if the right child is not null
               call setPosition on the right child node
       set the position of the current node
```

if the map contains the key and position

get the position and set the current node to it iterate the position

else

create the key and position set the current node position to 1 iterate the position

Class: BST

Insert()

If the head node of the BST is null, then the tree is empty

Set the head node data to the input data

Else

Add one to the current node

Compare the insert data to the left child

If the insert data is >

Move to the right

If the data is null, create node and set the node to the right child

Recursively call insert() on the right child

If the insert data is <

Move to the left

If the data is null, create node and set the node to the left child Else

Recursively call insert() on the left child

Delete()

If the head data is null, then the tree is empty

Return

If the data to delete is less than the current node, move to the right

Call delete() on the left child node

Else

Call delete() on the right child node

If it is not less than or greater than, then the node to delete has been found

Check if both children of the parent are null

Create a temp student object to hold the data

Get the data

Set the parent pointer to the child to null

Call BinaryTree.decrement() on the current node to decrease the size by one

Else if the node to be deleted has only 1 child

Find which node of the parent is the node to be deleted

Set the parent's child node to the child node of the nod to be deleted

Else the node to be deleted has two children

Swap the deleted node data with a child node until it becomes a leaf

node

Call BinaryTree.decrement() to decrease the size of all the parent nodes

by one

Delete the leaf node by setting the parent node's child pointer to null

Return the deleted data

Class: MaxHeap

Insert()

Create a new node to add to the array

Save the given student data to the new node

Find where the parent should be in the array

Make sure the parent is in the bounds of the array

If the node should be set to the left child of the parent – set it

If the node should be set to the right child of the parent – set it

If the parent is out of bounds of the array

Set the parent to null

Calculate the level of the node by dividing the current index position by 2 until 0 is found Set the position in the level by subtracting the current index from the total of the inner nodes $(2 \land level - 1)$

Call maxHeapify() to move the student data until the parent node is greater

maxHeapify()

hold the parent index by taking the current (index -1) / 2

if the parent is < 0, the array will be out of bounds

return;

else compare the parent node data to the current node data

if the parent is less than the current data – swap the data

if the parent is greater than the current data

return

call maxHeapfiy() recusively until pass the array bounds

printlnOrder()

in-order traversal of the array

printPreOrder()

pre-order traversal of the array

printPostOrder()

post-order traversal of the array

delete() – delete the head node

create a temp node to hold the node to be deleted
swap the head node with the last node
delete by new last node by setting it to null
call heapfiy on the first node to keep the MaxHeap ordr
return the temp node

heapify()

create a temp node
create an integer to hold the index of the left child of the current node
if the left child index is past the array length
return as the array will be out of bounds
compare the current node to the left child
if the left child is greater than current node – swap
call heapify() on the left child index
compare the current node to the right child
if the right child is greater than the current node – swap
call heapfiy() on the right child index

Class: Node, NodeQueue, NodeStack

These classes only have standard getters and setters

Class: Stack

Insert()

Create a new node to the push to the stack with the Node data set

If the head of the stack is null – the stack is empty

Set the head and tail of the stack to the new node

Delete()

If the head node is null

Return as the stack is empty

Else

Create a temp NodeStack object to hold the head node Set the head of the stack to Head.getNext() Return the temp NodeStack object

Standard getters and setters for the fields

Class: ReadFile

readFile()

set a counter to keep track of the current level and position of the trees if the file has data to read – read the first line create a new student object to hold the data

```
set the first and last name of the new student by scan.nextLine()
               add the student object to the array list to hold the original data
               add the student object to the minHeap
               set the minHeap size to 1
               set the minHeap level to 0
               set the minHeap position to 1
       increment the position
       use a while to read through the remaining data – scan.hasNext()
               if the position is greater that what the current level can hold
                      reset the position to 0
                      increment the levels
               increment the position
               Create a new student object to hold the data
               Read and set the student first and last name – scan.next()
               Read the new line character with a blank scan.nextLine()
               Create a new node to insert into the tree
               Set the data to the created student object
               Add the node to the tree with BinaryTree.insert()
       Close the scanner
getFile()
       create a boolean to determine if the file is found
       use a while loop to prompt the user for an input file until the file is found
               use a try / catch to catch errors from opening a file that does not exist
              try {
                      Create a file object with the default file name – namelist.txt
                      Create a scanner object with the opened file
                      Set the boolean to true because if we are this far an error has not be
thrown
               catch (FileNotFoundException e) {
                      Print to the user that the file was not found and they need to enter a new
file name
                      Change the scanner to read from the System.in
                      Set the filename string to scan.nextLine()
                      Trim the filename for any leading white spaces
                      Set the goodFile to false to run the test again
              }
```

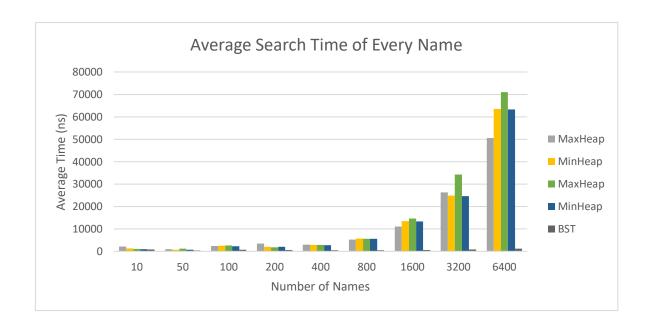
Class: ReadFile

This class only has standard getters and setters for the fields

Summary Analysis

The binary search tree performed well in the actual tests. As the number of names double for each the average search time stayed very low. The graphs of the binary search tree look nearly linear. The binary search tree had a high performance because, unlike the heaps, there is a definite decision of whether to move left or right by the comparison of the data. The issue that the binary search tree could have is if the data would come in a sorted order. This would cause the binary search tree to be a simple linked list and the searching would be linear. Assuming this doesn't happen, the binary search tree can cut off approximately half the data with the first comparison and so on through the tree.

The heaps did not perform very well with the data and name changes. The issue that the heaps have is that we cannot guarantee what the siblings of a node look like other than that they are either less than or greater than the parent depending on if a max or min heap is implemented. This causes us to search through the nodes on the level before moving to the next level. This issue can be somewhat eliminated by improving the enqueue process. If a max heap is implemented and the current node is less than the searched name, we are guaranteed by the max heap that there are only going to lesser names below it. This would allow us to not enqueue those nodes for searching.



Time Complexity

BinaryTree.depthFirst() – iterate through the stack O(n)

BinaryTree.breadthFirst() – iterate through the queue O(n)

BinaryTree.setPosition() – traverse the tree and set the position of the node O(n)

BinaryTree.pushStack() – push the item on the stack O(n)

BinaryTree.enqueue() – add the head to the queue O(n)

Print the binary tree – O(n)

BinaryTree.minHeapTopDown() – heapfiy from the root to the leaves O(h)

BinaryTree.minHeap() – start and leaf and swap until parent is lesser O(h)

BinarySearch.decrement() – traverse from leaf node to root O(h)

BinarySearch.increment() – traverse from leaf node to root O(h)

BinarySearch.insert() – traverse until a leaf node is found O(h)

BST.searchLast() – compare and move left or right O(h) *H could be linear*

BST.setPosition – set the position of the node relative to siblings O(n)

BST.insert() - navigate to a leaf, insert O(h) *H could be linear in worst case*

MaxHeap.heapfiy() – heapify from top down O(h)

Print max heap - O(n)

MaxHeap.maxHeapfiy() – compare to the parent, swap if greater O(h)

MaxHeap.insert() – insert at the end of the array O(1)

ReadFile.readFile() – read every line of the file O(n)

Node, NodeQueue, NodeStack, and Student Class – only getters and setters O(1)

Phase 3: Risk Analysis

The only risk for this application is having a file that is not found. If this happens, the user is prompted for a new file.

Phase 4: Verification

The application will search every name and output the average time to a file

Phase 5: Coding

The code of this program is included in the zip file. The code is explained with comments and line breaks to make reading easier. A Javadoc is also made to explain the use of each method.

Phase 6: Testing

Test runs have been included at the end of this file.

Phase 7: Refining the program

The tree structure of the binary search tree could be improved by finding an average name to start the root with. This will allow the data to more likely be split in half.

Phase 8: Production

A zip file including the source files from eclipse and the output of the program have been included.

Phase 9: Maintenance

Changes can be made once feedback is received.

Program Output