

Alg - 3

HW. (1) $A \sim_s B \Leftrightarrow A = TBT^{-1}$ for some invertible $T \in M_{n \times n}(R)$ and $A, B \in M_{n \times n}(R)$
 Prove that \sim_s is an equivalence rel on $M_n(R)$.

2. Define $x \sim y \Leftrightarrow x-y = n$ for some $n \in \mathbb{Z}$
 and $x, y \in R$, the set of reals
 Prove that \sim is an equivalence relation on R .

Note \mathbb{R}/\sim will be identified with

$$S' = T = \left\{ z \in \mathbb{C} : |z| = 1 \right\} \text{ where}$$

\sim is defined in ~~HW~~ at 3 (2) i.e.

$$x \sim y \Leftrightarrow x-y = n \text{ for some } n \in \mathbb{Z}.$$

Defn $p \geq 2$ integer is prime iff $p = mn$, $m, n \in \mathbb{N}$
 Then either $m=1$ or $n=1$
 (only divisors of p are $1, p$)

Thm $n \in \mathbb{N}$ is either a prime or
 \dots else $k_1 \cdots k_r$

Thm $n \in \mathbb{N}$ receives a prime factorization
 else $n = p_1 p_2 \cdots p_r$ where p_1, \dots, p_r
 are primes in \mathbb{N} , $r \geq 1$ integer.
 p_1, \dots, p_r need not be distinct.

Proof. Let $S = \{n \in \mathbb{N} : n \text{ is not a product of primes}\}$

$S \subseteq \mathbb{N}'$. To show $S = \emptyset$.

Suppose not. Then by WOP there
 is a smallest $m \in S$: $m \leq n$ $\forall n \in S$
 Since $m \in S$, m is not product of primes
 So there are $d, e \in \mathbb{N}$ st $1 < d, 1 < e$
 st. $m = de$. $d, e \notin S$ as $m > d \wedge m > e$

So $d = t_1 \cdots t_r$, $e = t'_1 \cdots t'_s$ where
 $t_1, \dots, t_r, t'_1, \dots, t'_s$ are primes

But $m = de = t_1 \cdots t_r t'_1 \cdots t'_s$, a
 product of primes $\Rightarrow m \notin S$, a
 contradiction. So $S = \emptyset$

Fundamental Theorem of arithmetic

Fundamental Theorem of arithmetic

(i) Every $z \leq n \in \mathbb{N} \Rightarrow n = p_1^{\alpha_1} \cdot \cdot \cdot p_r^{\alpha_r}$ where

$p_1 < p_2 < \dots < p_r$, primes &

$\alpha_i \geq 0$ integers $i=1 \dots r$

$$\left| \begin{array}{l} \\ m^o = 1 \end{array} \right.$$

(ii) If also $n = q_1^{\beta_1} \cdot \cdot \cdot q_s^{\beta_s}$

where $q_1 < \dots < q_s$ are primes

and $\beta_i \geq 0$ integers, $i=1 \dots s$

then $r=s$ and $p_i = q_i$, $i=1 \dots r=s$

These unique α_i , $i=1 \dots r$
are called orders of n at p_i .

$$\text{ord}_{p_i}(n) = \alpha_i$$

$$n = 8 = 2^3, \quad p_1 = 2, \quad \alpha_1 = 3$$

$$\text{ord}_2 8 = 3$$

Ex $m, n \in \mathbb{N},$

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$$m = p_1^{\alpha_1} \cdots p_r^{\alpha_r}$$

$$n = p_1^{\beta_1} \cdots p_r^{\beta_r} \quad \alpha_1, \dots, \alpha_r, \beta_1, \dots, \beta_r \geq 0$$

integers

$$(n, m) = d = p_1^{\min(\alpha_1, \beta_1)} p_2^{\min(\alpha_2, \beta_2)} \cdots p_r^{\min(\alpha_r, \beta_r)}$$

$$[n, m] = p_1^{\max(\alpha_1, \beta_1)} p_2^{\max(\alpha_2, \beta_2)} \cdots p_r^{\max(\alpha_r, \beta_r)}$$

Equivalence relations and congruences

S is a set (always non empty)

Def (1)
A relation \sim is a subset of $S \times S$
 $= \{(a, b) : a, b \in S\}$, the set of all
ordered pairs : $(a, b) = (c, d) \iff a = c$
 $b = d$

(2) A relation \sim on S is called an
equivalence relation, write $a \sim b$ if
 $(a, b) \in \sim \subseteq S \times S$ iff

(i) $a \sim a$ (reflexive prop) $\forall a \in S$
 $\therefore a \sim b \Rightarrow b \sim a$ (symmetric)

- (ii) $a \sim b \Rightarrow b \sim a$ (symmetric)
 $\wedge a, b \in S$
- (iii) $a \sim b \wedge b \sim c \stackrel{\text{and}}{\Rightarrow} a \sim c$
 (transitive prop.)

Ex | $x, y \in \mathbb{Z}$, $n \in \mathbb{N}$, fixed. Define

$$x \sim_n y \Leftrightarrow n|x-y$$

Then \sim_n is an equivalence rel
on \mathbb{Z} .

(2) $A, B \in M_{n \times n}(\mathbb{R})$. $A \sim_s B$ (A is
similar to B) iff $A = TBT^{-1}$ for
some invertible $T \in M_{n \times n}(\mathbb{R})$
then \sim_s is an equivalence rel
in $M_{n \times n}(\mathbb{R})$

(3) $A, B \in M_{n \times n}(\mathbb{C})$, $A \sim_v B$ (unitary
equivalence) iff $A = UBU^{-1}$ for
some $U \in M_{n \times n}(\mathbb{C})$ unitary: $U^{-1} = U^*$
 $= \overline{U^T}$,

$$\text{work } \propto \min_{t=1}^T \dots = \frac{\bar{V}^T}{\bar{U}} t$$

Let S be a set & \sim be an equivalence relation on S . For each $x \in S$ let $C_x = [x] = \{y \in S : y \sim x\}$, \checkmark
is called the equivalence class of x .

Thm (1) $C_x \neq \emptyset$ \checkmark

(2) $C_x \cap C_y = \emptyset$ or if $C_x \cap C_y \neq \emptyset$
then $C_x = C_y$

(3) $\{C_x\}_{x \in S}$ is a partition of S

Def. $\{A_\alpha\}_{\alpha \in I \neq \emptyset}$ of subsets of S is

a partition of S iff

(1) $A_\alpha \neq \emptyset \quad \forall \alpha \in I$

(1) $A_\alpha \neq \emptyset \quad \forall \alpha \in I$

(2) $A_\alpha \cap A_\beta = \emptyset \quad \text{for } \alpha \neq \beta$

(3) $S = \bigcup_{\alpha \in I} A_\alpha$

Proof of thm. Since $x \sim x$
 $\Rightarrow x \in C_x \Rightarrow C_x \neq \emptyset$. \therefore (1)

(2) Let $x, y \in S$. If $C_x \cap C_y = \emptyset$, done

(nothing to prove) Suppose $C_x \cap C_y \neq \emptyset$

then to show $C_x = C_y$?.

$C_x \subseteq C_y$ and $C_y \subseteq C_x$.

Since $C_x \cap C_y \neq \emptyset$, $\exists z \in C_x \cap C_y$

$z \in C_x$ and $z \in C_y$

\Downarrow

\Downarrow

$z \sim x$

$z \sim y$

\Downarrow
 $x \sim z$ by symmetry & $z \sim y$

$\Rightarrow x \sim y$ by transitivity
prop

$a \in C_x$. Want to prove $a \in C_y$

$a \in C_x \Rightarrow a \sim x$, and $x \sim y$

$\Rightarrow a \sim y$ by transitive prop

$\Rightarrow a \in C_y$. Similarly.

Show that $C_y \subseteq C_x$ · proof:

$$C_y = C_x$$

$$(3) \quad \bigcup_{x \in S} C_x = S$$

Since $C_x \subseteq S$ for each x

$\bigcup_{x \in S} C_x \subseteq S$. If $a \in S$

$$\text{II. } a \in C_x \subset \bigcup_{x \in S} C_x$$

thus $a \in C_a \subseteq \bigcup_{x \in S} C_x$

$\Rightarrow S \subseteq \bigcup_{x \in S} C_x$

proving $S = \bigcup_{x \in S} C_x$

Notation The set $S/\sim = \{[x] : x \in S\}$
 $= \{C_x : x \in S\}$

called called the quotient set.

Theorem \sim_n is an equivalence relation on \mathbb{Z}
where $x \sim_n y \Leftrightarrow n \mid x-y$

Proof (1) $x \sim_n x \Leftrightarrow n \mid x-x=0$
 $n \mid 0 = 0 \cdot n$.

(2) Suppose $x \sim_n y$. $x-y=g^n$
for $g \in \mathbb{Z}$

$$\begin{aligned}
 \text{Now } y-x &= -(x-y) \\
 &= -(g^n) \\
 &= (-g)^n \\
 \Rightarrow n &\mid y-x \cdot (\text{symmetric prop})
 \end{aligned}$$

(3) Suppose $x_n \sim y \wedge y_n \sim z$, for $x, y, z \in \mathbb{Z}$
 then $x-y = g_1^n$ and $y-z = g_2^n$
 for $g_1, g_2 \in \mathbb{Z}$

$$\begin{aligned}
 x-z &= x-y + y-z = g_1^n + g_2^n \\
 &= (g_1 + g_2)^n \\
 \Rightarrow n &\mid x-z \Rightarrow \text{transitivity}.
 \end{aligned}$$

$$[x] = [x]_n = \{y \in \mathbb{Z} : y \sim_n x\}$$

$$\begin{aligned}
 (1) \quad x_n \sim y, a_n \sim b \Rightarrow &(i) x+a \sim_n y+b \\
 &(ii) xa \sim_n yb
 \end{aligned}$$

$$\mathbb{Z}/_n = \left\{ [x]_n : x \in \mathbb{Z} \right\}$$

$$[x]_n \oplus [y]_n = [x+y]_n - \text{num}$$

$$[x]_n \otimes [y]_n = [xy]_n - \text{product}$$

x is congruent to $y \pmod n$, write
 $x \equiv y \pmod n \iff x \sim_n y$ i.e.

$$n \mid x-y.$$