

$$(\bar{\delta t})^2 = 5 (1/5)^2 S^2 = S^2/5$$

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Generally $\bar{\delta t} = \frac{S(t)}{\sqrt{N}}$ uncertainty in mean

Numerical computing 'errors'

- approximation (we cut off Taylor series after finite terms)

- roundoff
computer doesn't have infinite memory

to store $\sqrt{2}$

uses floating point representation

python floats are double precision (doubles)

$$X_{\text{double}} = (-1)^s \times 1.f \times 2^{n-1023}$$

$s \equiv \text{sign}$
1 bit

$f \equiv \text{mantissa}$
52 bits
precision

$n \equiv \text{exponent}$
11 bits
range

range of doubles $\pm 2^{-1074} \leftrightarrow \pm 2^{1024}$

$\pm 4.9e-324 \leftrightarrow \pm 1.8e308$

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We get underflow or overflow for
trying to store a smaller or larger number

We can't store 324 sig figs

Precision given by $1/2^{50} \approx 2.2e-16$

(about 16 decimal digits of precision)

[small jupyter code]

Careful comparing floats

[code]

Computing e^x (an example)

Taylor
$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!}$$

approximate as
$$e^x \approx \sum_{n=0}^{n_{\max}} x^n / n!$$

problem: as n increases (more terms)

x^n and $n!$ can become huge

even though ratio might be small

We can try to be smarter

Note that $\frac{x^n}{n!} = \frac{x}{n} \frac{x^{n-1}}{(n-1)!}$ 2-9

↙ new term ↘ old term

What should we take for n_{\max} ? How many terms?

Let's keep adding terms until the sum
doesn't change (can't do any better)

[comp exp code]

Derivatives

$$\vec{v} = d\vec{r}/dt \quad \vec{E} = -\nabla\phi - \partial\vec{A}/\partial t$$

Typical case: We have set of n discrete data points $(x_i, f(x_i))$ $i=0, 1, \dots, n-1$
 We want to calculate $f'(x)$ at specific point

Recall analytical formula

$$\frac{df(x)}{dx} = \lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h}$$

We can code this and use a 'small' value for h

→ What uncertainty is in our answer

→ Small relative to what

As h ~~gets~~ gets small $f(x+h)$ gets close to $f(x)$
 and then it gets divided by small h

Let's try to be more systematic

Nlok: we have 2 sources of error 3-3
 approximation due to cutting off Taylor series
 roundoff due to subtraction and division

$$\mathcal{E}_{app} = \frac{h}{2} |f''(x)|$$

From Chapter 2: $c = a - b$ $\Delta c \equiv$ uncertainty in c
 $|\Delta c| \leq |A| + |B|$

$$|Af(x+h)| \approx |f(x)| \epsilon_m \quad |Af(x)| \approx |f(x)| \epsilon_m$$

$$\mathcal{E}_{ro} \approx \frac{2 |f(x)| \epsilon_m}{h}$$

details not that important
 $\mathcal{E}_{ro} \approx 2/h$

$$\mathcal{E}_{tot} = \mathcal{E}_{app} + \mathcal{E}_{ro} \approx \beta h + 2/h$$

one term prefers small h for less error

one term prefers large h for less error

We can find optimal by

$$\frac{d\mathcal{E}_{tot}}{dh_*} = 0 = \beta - 2/h_*^2 \quad h_* = \sqrt{2/\beta}$$

$$\mathcal{E}_{tot} = \beta h_* + 2/h_* = \beta \sqrt{2/\beta} + 2\sqrt{\beta/2}$$

$$= 2\sqrt{2\beta} \quad \text{See Egn 3.20}$$