A faint, light blue network graph is visible in the background, consisting of numerous nodes (small circles) connected by thin lines, forming a complex, interconnected web.

Lecture 16: Graphs

CS 111: Intro to Computational Science
Spring 2023

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Administrative

- New homework due **Tuesday**
- Lab tomorrow
- Quiz 5 grades are on Canvas
- Remember: NO CLASS NEXT MONDAY!
- Your last quiz is next week Wednesday (Quiz 7)
 - Will be on lectures 15, 16 (PCA, Graphs)
- Re: Final Exam
 - **Wed. June 14th from 9 – 11 am → ARRIVE 10 mins EARLY!**

Networks Can Be About... Anything

Network-centric views help us look at things from a certain perspective.

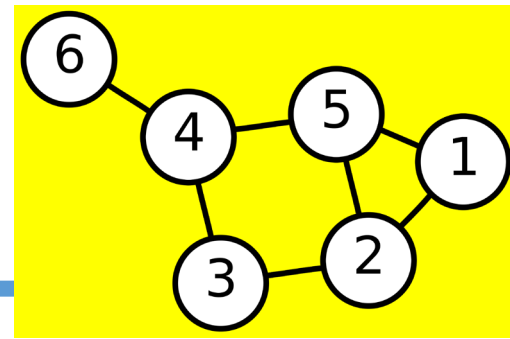
Example: **Network of Ideas**

[http://www.ted.com/talks/
eric_berlow_and_sean_gourley_mapping_ideas_worth_spreading](http://www.ted.com/talks/eric_berlow_and_sean_gourley_mapping_ideas_worth_spreading)
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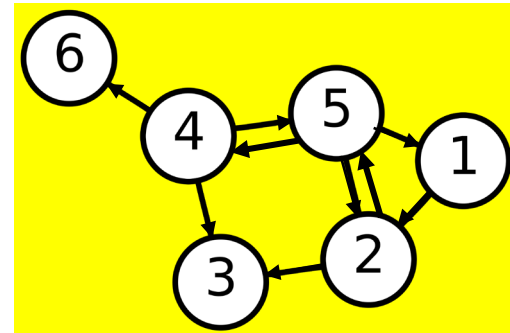
Check out this short talk!

Graph Theory

- Graphs: mathematical structures that model *pairwise relations* between objects
- Made up of **vertices** (aka **nodes**) which are connected by **edges** (aka **links**).
- Some graphs are **undirected** (use straight lines to symbolize edges)
 - where edges link two vertices symmetrically ($A \rightarrow B$ means $B \rightarrow A$ too)
- Some graphs are **directed** (use arrows to symbolize edges)
 - where edges link two vertices asymmetrically ($A \rightarrow B \neq B \rightarrow A$)



Undirected graph

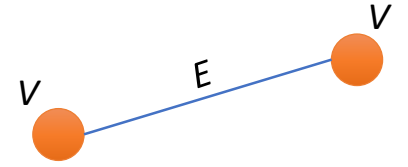


Directed graph

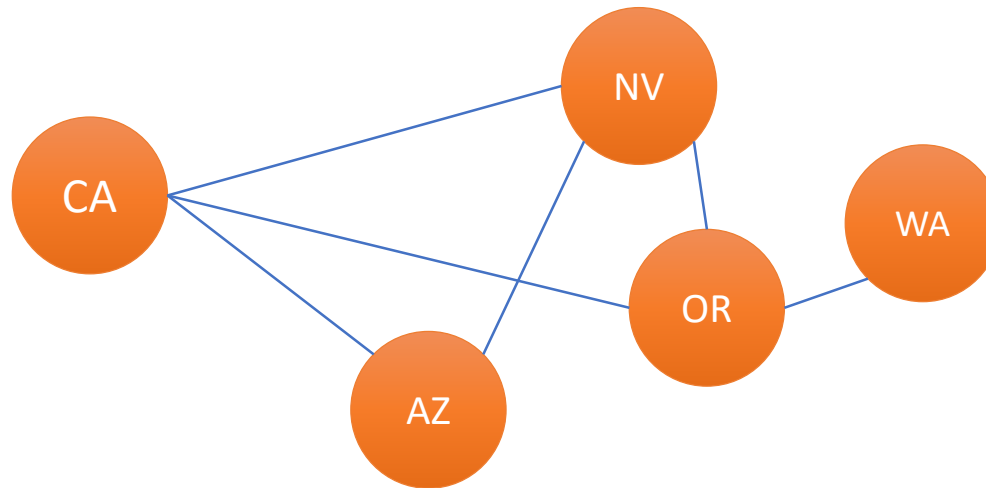
Studying Networks of... *anything*...

- Based on Graph Theory
- Network structures can describe a lot of phenomena
 - How certain cancer cell form
 - How people make social links with one another (i.e. social networks)
 - How words and their meanings are linked together (in linguistics)
 - *How documents end up being linked on the WWW*
 - *How internet protocols works*

Vertices and Edges *aka* Nodes and Links



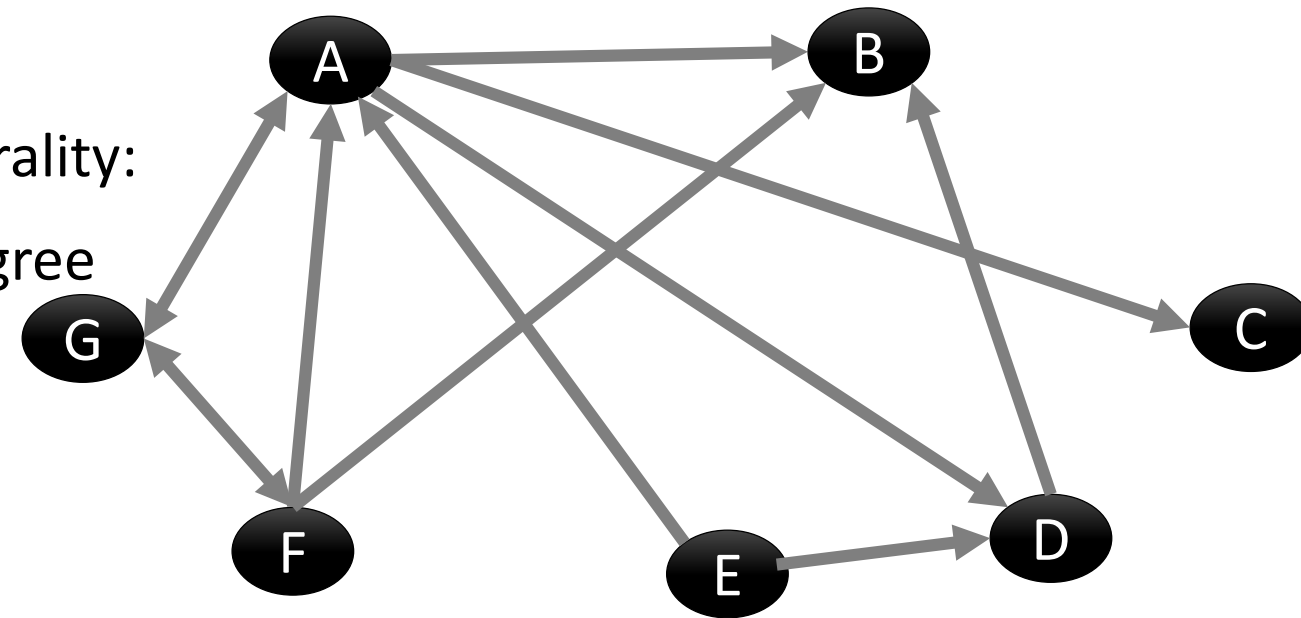
- *The only 2 components of any graph/network!*
- Vertices (nodes/entities) are connected to each other via edges (links/relationships)
- **Vertices** can be: Web docs, people, words, *etc...*
- **Edges** can be: “has URL link to”, “knows”, “found before (other words)”, *etc...*



Degree Centrality

- Measures of the “Importance” of a Vertex/Node
 - Simply this: Importance is based only on how many edges/links does a node have attached to it?

- Example: Looking at Node A’s degree centrality:
3 in-degree, 4 out-degree

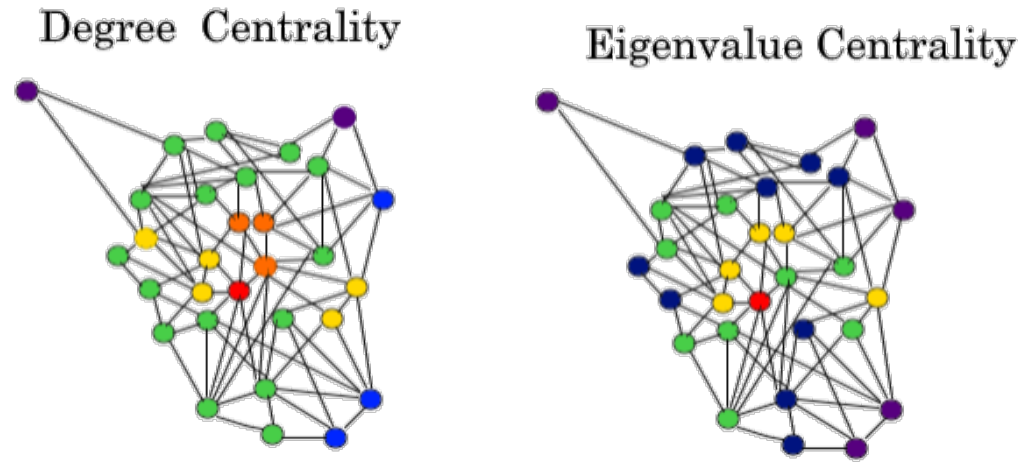


- **Pros**: It’s a simple measure of node “importance” or “connectivity”
- **Cons**: It’s too simple a measure...

Eigenvector Centrality

- A more sophisticated version of degree centrality
 - Not all “highly central” nodes have the same “importance”
 - Main Idea: A node is **important** if it is linked to **other important nodes**
- Like *degree centrality*, this one calculates number of links on a given node
 - **BUT** then each adjacent node is **weighted** by its *own* centrality
 - So, a link to a high-centrality node “**counts for**” more than one going to a less-connected node

Eigenvector Centrality

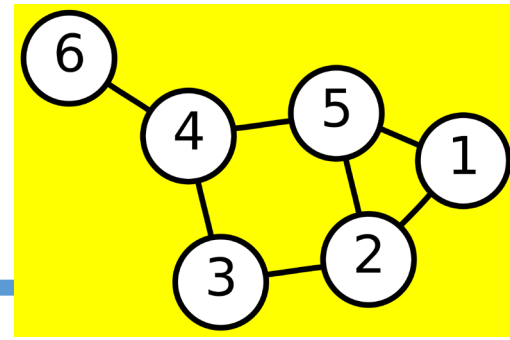


- Popular in measuring ***influence*** and ***access*** to influential nodes in a network
- The Google search algorithm (**PageRank**) uses a *variation on this measure* to decide how “important” a webpage is

Other Centrality Measures

- Example: “**Betweenness Centrality**”
- If a node has *high* betweenness centrality
 - ➔ It has a *large potential for controlling “flows”*
(like messages, influence, resources) through the network
because it “shortens” the “number of hops” between any other 2 nodes

The Degree Matrix



- Network graphs can be described using matrices
- **Degree matrix**
 - An $n \times n$ **diagonal** matrix: describes the **degree** of each node in a graph with n nodes
 - Usually used only for undirected graphs
 - For directed graphs, you need one matrix for *in-degrees* and another for *out-degrees*

• $D =$

$$\begin{pmatrix} 2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 3 & 0 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 0 & 0 \\ 0 & 0 & 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

Node 1 has degree = 2

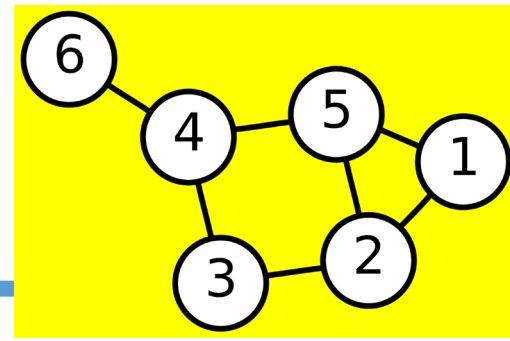
Node 2 has degree = 3

Note, that **trace(D)**, i.e. the sum of diagonals, is **twice** the number of edges for obvious reasons.

i.e. **trace(D) = 2m**, where **m** = no. of edges

In the example: $m = 7$, $\text{trace}(D) = 14$

The Adjacency Matrix

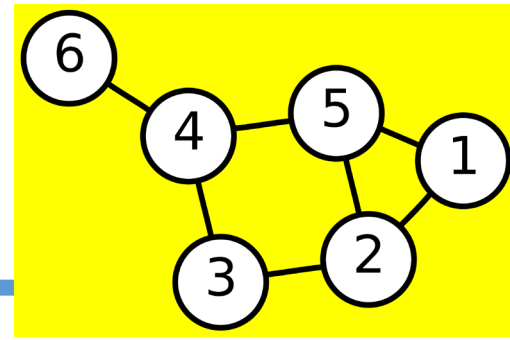


- Shows what vertices (nodes) are connected in an n sized graph
 - Is always an $n \times n$ matrix
 - In a large graph, it is typically a sparse matrix
- A very compact way of representing a graph
 - Only need **1** (for 'connected') and **0** (for 'not connected') values in it
 - Also, we can easily derive **D** from **A** *(how?)*
- In an undirected graph, **A** is always symmetrical (if node $x \rightarrow y$, then node $y \rightarrow x$ too)
- In a directed graph, **A** is *not necessarily* symmetrical

Nodes: 1 2 3 4 5 6																																											
A =	<table><tr><th>1</th><td>0</td><td>1</td><td>0</td><td>0</td><td>1</td><td>0</td></tr><tr><th>2</th><td>1</td><td>0</td><td>1</td><td>0</td><td>1</td><td>0</td></tr><tr><th>3</th><td>0</td><td>1</td><td>0</td><td>1</td><td>0</td><td>0</td></tr><tr><th>4</th><td>0</td><td>0</td><td>1</td><td>0</td><td>1</td><td>1</td></tr><tr><th>5</th><td>1</td><td>1</td><td>0</td><td>1</td><td>0</td><td>0</td></tr><tr><th>6</th><td>0</td><td>0</td><td>0</td><td>1</td><td>0</td><td>0</td></tr></table>	1	0	1	0	0	1	0	2	1	0	1	0	1	0	3	0	1	0	1	0	0	4	0	0	1	0	1	1	5	1	1	0	1	0	0	6	0	0	0	1	0	0
1	0	1	0	0	1	0																																					
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4	0	0	1	0	1	1																																					
5	1	1	0	1	0	0																																					
6	0	0	0	1	0	0																																					

Note: This shows an undirected graph

The Graph Laplacian Matrix



- Shows BOTH **degree** and **adjacency** information in an n sized graph
 - Calculated as: $L = D - A$
 - Is always an $n \times n$ matrix

$$\bullet L = \begin{pmatrix} 2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 3 & 0 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 0 & 0 \\ 0 & 0 & 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} - \begin{pmatrix} 0 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix} = \begin{pmatrix} 2 & -1 & 0 & 0 & -1 & 0 \\ -1 & 3 & -1 & 0 & -1 & 0 \\ 0 & -1 & 2 & -1 & 0 & 0 \\ 0 & 0 & -1 & 3 & -1 & -1 \\ -1 & -1 & 0 & -1 & 3 & 0 \\ 0 & 0 & 0 & -1 & 0 & 1 \end{pmatrix}$$

L has interesting properties that combine properties for D and A and other things...!

Graph Laplacian Characteristics

- Similar to **A**, vertex connections are spelled-out (here, as all the **-1** entries)
- Similar to **D**, the trace of **L** = $2m$; where m = number of edges
- In a GL matrix, every row adds to zero, so does every column

$$\begin{pmatrix} 2 & -1 & 0 & 0 & -1 & 0 \\ -1 & 3 & -1 & 0 & -1 & 0 \\ 0 & -1 & 2 & -1 & 0 & 0 \\ 0 & 0 & -1 & 3 & -1 & -1 \\ -1 & -1 & 0 & -1 & 3 & 0 \\ 0 & 0 & 0 & -1 & 0 & 1 \end{pmatrix}$$

- Let vector $\mathbf{v}_0 = [1,1,1,1,1]^T$; Note that: $\mathbf{L} \mathbf{v}_0 = \lambda_0 \mathbf{v}_0 = \mathbf{0}$

$$\begin{pmatrix} 2 & -1 & 0 & 0 & -1 & 0 \\ -1 & 3 & -1 & 0 & -1 & 0 \\ 0 & -1 & 2 & -1 & 0 & 0 \\ 0 & 0 & -1 & 3 & -1 & -1 \\ -1 & -1 & 0 & -1 & 3 & 0 \\ 0 & 0 & 0 & -1 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

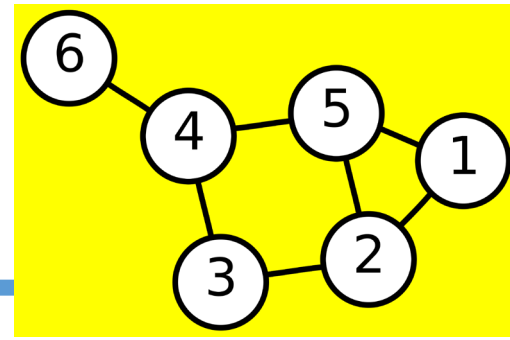
- Therefore, λ_0 of **L** is always zero!
- Therefore, **L** is a singular matrix (in the example given, the rank of this 6x6 matrix is 5)

- **L** is **symmetric** and **positive semi-definite** (i.e. it's eigenvalues are ≥ 0)

Graph Laplacian Matrix Uses

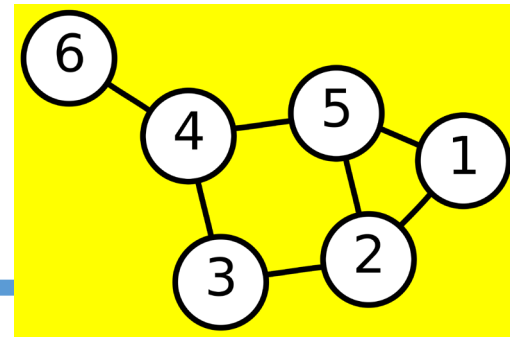
- Clustering algorithms
 - Applied in Machine Learning, task scheduling in multiprocessor computers, electronic integrated circuit design, etc...
- Kirchhoff's Theorem and Spanning Trees in Math/Theoretical C.S.
 - Spanning tree = a path in a network that includes all the nodes
 - Number of Spanning trees in a network = $\frac{1}{n} (\lambda_1 \lambda_2 \dots \lambda_{n-1})$
 - where λ_i are non-zero eigenvalues of the GL of the network
 - Helps to reduce a graph to a tree

The “Density” of a Graph/Network



- Tells us how many links/edges are “realized” in the network.
- *A fully-connected network* (den. = 100%)
is where *every* node/vertex connects to *every other* node
- *Thought exercise: If I have n vertices, how many connections can I make?*
- The **max** number of edges in a network with n nodes is: $\frac{1}{2} n(n-1)$
 - Assumes an undirected graph/network
 - **Double** that number for directed graph (i.e. it's $n(n-1)$)

The “Density” of a Graph/Network



- Finally, the **density of a network (d)** is defined as:

$$d = (\text{number of actual edges}) / (\text{max. number of edges})$$

- Example: in the above network, the density is:

$$d = 7 / (\frac{1}{2} * 6 * 5) = 7 / 15 = \underline{46.7\%}$$

- Density is a useful measure when studying the efficiency of how stuff disperses in a network
 - The most efficient networks for dispersion tend to have “medium” densities...!

Your TO DOs!

- Finish new assignment by Tuesday
- Lab tomorrow!

</LECTURE>