beamer named overlay specification with beanoves

Jérôme Laurens

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Abstract

This package allows the management of multiple slide lists in `beamer` documents. Slide lists are very handy both during edition and to manage complex and variable beamer overlay specifications.

Contents

1 Minimal example

The document below is a contrived example to show how the `beamer` overlay specifications have been extended.

```
1  \documentclass {beamer}
2  \RequirePackage {beanoves}
3  \begin{document}
4  \Beanoves {
5      A = 1:2,
6      B = A.next:3,
7      C = B.next,
8    }
9  \begin{frame}
10 {\Large Frame \insertframenumber}
11 {\Large Slide \insertslidenumber}
12 \visible<?(A.1)> {Only on slide 1}\\
13 \visible<?(B.1)-?(B.last)> {Only on slide 3 to 5}\\
14 \visible<?(C.1)> {Only on slide 6}\\
15 \visible<?(A.2)> {Only on slide 2}\\
16 \visible<?(B.2::B.last)> {Only on slide 4 to 5}\\
17 \visible<?(C.2)> {Only on slide 7}\\
18 \visible<?(A.3)-> {From slide 3}\\
19 \visible<?(B.3::B.last)> {Only on slide 5}\\
20 \visible<?(C.3)> {Only on slide 8}\\
21 \end{frame}
22 \end{document}
```

On line 4, we use the `\Beanoves` command to declare named slide ranges. On line 5, we declare a slide range named 'A', starting at slide 1 and with length 2. On line 12, the extended *named overlay specification* `?(A.1)` stands for 1, on line 15, `?(A.2)` stands for 2 whereas on line 18, `?(A.3)` stands for 3. On line 6, we declare a second slide range named 'B', starting after the 2 slides of 'A' namely 3. Its length is 3 meaning that its last slide number is 5, thus each `?(B.last)` is replaced by 5. The next slide number after slide range 'B' is 6 which is also the start of the third slide range due to line 7.

2 Named slide lists

2.1 Presentation

Within a beamer frame, there are different slides that appear in turn. The main slide list is a range of integers covering all the slide numbers, from one to the total amount of slides. In general, a slide list is a range of positive integers identified by a unique name. The main practical interest is that such lists may be defined relative to one another, we can even have lists of slide ranges. Finally, we can use these lists to organize beamer overlay specifications logically.

2.2 Defining named slide lists

In order to define named slide lists, we can either use the `\Beanoves` command below before a beamer frame environment, or use the `beanoves` option of this environment. The

value of the `beanoves` option is similar to the argument of the `\Beanoves` commands, but the latter takes precedence on the former. This behaviour may be useful to input the very same source code into different frames and have different combinations of slides.

`beanoves`

```
beanoves = {
  ⟨name_1⟩=⟨spec_1⟩,
  ⟨name_2⟩=⟨spec_2⟩,
  ...,
  ⟨name_n⟩=⟨spec_n⟩,
}
```

`\Beanoves`

```
\Beanoves{
  ⟨name_1⟩=⟨spec_1⟩,
  ⟨name_2⟩=⟨spec_2⟩,
  ...,
  ⟨name_n⟩=⟨spec_n⟩,
}
```

The keys $\langle name_i \rangle$ are the slide lists names, they are case sensitive and must contain no spaces nor '`/`' character. In order to avoid name conflicts with floating point functions, it is suggested to let them contain at least an uppercase letter ot an underscore. When the same key is used multiple times, only the last one is taken into account. Possible values for $\langle spec_i \rangle$ are the *slide range specifiers* $\langle first \rangle$, $\langle first \rangle$:$\langle length \rangle$, $\langle first \rangle$::$\langle last \rangle$, :$\langle length \rangle$::$\langle last \rangle$ where $\langle first \rangle$, $\langle length \rangle$ and $\langle last \rangle$ are algebraic expression possibly involving any integer valued named overlay specifications defined below.

Also possible values are *slide list specifiers* which are comma separated list of *slide range specifiers* and *slide list specifier* between square brackets. The definition

$\langle name \rangle$=[$\langle spec_1 \rangle$,$\langle spec_2 \rangle$,...,$\langle spec_n \rangle$],

is a convenient shortcut for

$\langle name \rangle$.1=$\langle spec_1 \rangle$,
$\langle name \rangle$.2=$\langle spec_2 \rangle$,
...,
$\langle name \rangle$.n=$\langle spec_n \rangle$.

The rules above can apply individually to each

$\langle name \rangle$.i=$\langle spec_i \rangle$.

Moreover we can go deeper: the definition

$\langle name \rangle$=[[$\langle spec_{1.1} \rangle$, $\langle spec_{1.2} \rangle$],[[$\langle spec_{2.1} \rangle$, $\langle spec_{2.2} \rangle$]]

happens to be a convenient shortcut for

$\langle name \rangle$.1.1=$\langle spec_{1.1} \rangle$,
$\langle name \rangle$.1.2=$\langle spec_{1.2} \rangle$,
$\langle name \rangle$.2.1=$\langle spec_{2.1} \rangle$,
$\langle name \rangle$.2.2=$\langle spec_{2.2} \rangle$

and so on.

3 Named overlay specifications

3.1 Named slide ranges

When *slide range specifications* are used, the named overlay specifications are detailled in the tables below together with their replacement meaning value as beamer standard

overlay specification.

$\langle$*name*$\rangle$ == [i, $i+1$, $i+2, \ldots$]	
syntax	**meaning**
$\langle$*name*$\rangle$.1	i
$\langle$*name*$\rangle$.2	$i+1$
$\langle$*name*$\rangle.\langle$*integer*$\rangle$	$i+\langle$*integer*$\rangle-1$

In the frame example below, we use the `\BeanovesEval` command for the demonstration. It is mainly used for debugging and testing purposes.

```
\Beanoves {
  A = 3:6,
}
\begin{frame} {Frame \insertframenumber} {Slide \insertslidenumber}
\ttfamily
\BeanovesEval(A.1) ==3,
\BeanovesEval(A.2) ==4,
\BeanovesEval(A.-1)==1,
\end{frame}
```

When the slide range has been given a length or an end, like in the frame example below, we also have

$\langle$*name*$\rangle$ == [i, $i+1, \ldots,$ j]			
syntax	**meaning**	**example**	**output**
$\langle$*name*$\rangle$.length	$j-i+1$	A.length	6
$\langle$*name*$\rangle$.last	j	A.last	8
$\langle$*name*$\rangle$.next	$j+1$	A.next	9
$\langle$*name*$\rangle$.range	i ''-'' j	A.range	3-8

```
\Beanoves {
  A = 3:6, % or equivalently A = 3::8 or A = :6::8,

}
\begin{frame} {Frame \insertframenumber} {Slide \insertslidenumber}
\ttfamily
\BeanovesEval(A.1)      == 3,
\BeanovesEval(A.length) == 6,
\BeanovesEval(A.last)   == 8,
\BeanovesEval(A.next)   == 9,
\BeanovesEval(A.range)  == 3-8,
\end{frame}
```

Using these specifications on unfinite named slide ranges is unsupported. Finally each named slide range has a dedicated counter $\langle$*name*$\rangle$.n which is some kind of variable that can be used and incremented[1].

$\langle$*name*$\rangle$.n : use the position of the counter

[1] This is actually an experimental feature.

⟨***name***⟩`.n+=`⟨***integer***⟩ : advance the counter by ⟨*integer*⟩ and use the new position

`++`⟨***name***⟩`.n` : advance the counter by 1 and use the new position

Notice that "`.n`" can generally be omitted.

3.2 Named slide lists

After the definition
⟨*name*⟩`=[`⟨$spec_1$⟩`,`⟨$spec_2$⟩`,...,`⟨$spec_n$⟩`]`
the rules of the previous section apply recursively to each individual declaration
⟨*name*⟩`.`*i*`=`⟨$spec_i$⟩.

4 `?(...)` query expressions

This is the key feature of the beanoves package, extending beamer *overlay specifications* included between pointed brackets. Before the *overlay specifications* are processed by the beamer class, the beanoves package scans them for any occurrence of '`?(`⟨***queries***⟩`)`'. Each one is then evaluated and replaced by its static counterpart. The overall result is finally forwarded to the beamer class.

The ⟨*queries*⟩ argument is a comma separated list of individual ⟨*query*⟩'s of next table. Sometimes, using ⟨***name***⟩`.range` is not allowed as it would lead to an algeabraic difference instead of a range.

query	static value	limitation
`:`	`-`	
`::`	`-`	
⟨***first expr***⟩	⟨*first*⟩	
⟨***first expr***⟩`:`	⟨*first*⟩ `-`	no ⟨***name***⟩`.range`
⟨***first expr***⟩`::`	⟨*first*⟩ `-`	no ⟨***name***⟩`.range`
⟨***first expr***⟩`:`⟨***length expr***⟩	⟨*first*⟩ `-` ⟨*last*⟩	no ⟨***name***⟩`.range`
⟨***first expr***⟩`::`⟨***end expr***⟩	⟨*first*⟩ `-` ⟨*last*⟩	no ⟨***name***⟩`.range`

Here ⟨*first expr*⟩, ⟨*length expr*⟩ and ⟨*end expr*⟩ both denote algebraic expressions possibly involving named overlay specifications and counters. As integers, they respectively evaluate to ⟨*first*⟩, ⟨*length*⟩ and ⟨*last*⟩.

For example both `?(A.next)`, `?(A.last+1)`, `?(A.1+A.length)` give the same result as soon as the slide range named '`A`' has been properly defined with a starting value and a length.

Notice that nesting `?(...)` expressions is not supported.

1 ⟨*package⟩

5 Implementation

Identify the internal prefix (LaTeX3 DocStrip convention).

2 ⟨`@@=bnvs`⟩

5.1 Package declarations

```
\NeedsTeXFormat{LaTeX2e}[2020/01/01]
\ProvidesExplPackage
  {beanoves}
  {2022/10/28}
  {1.0}
  {Named overlay specifications for beamer}
```

5.2 logging and debugging facilities

Utility message.

```
\msg_new:nnn { beanoves } { :n } { #1 }
\msg_new:nnn { beanoves } { :nn } { #1~(#2) }
\cs_set:Npn \__bnvs_DEBUG_:nn #1 #2 {
  \msg_term:nnn { beanoves } { :n } { #1~#2 }
}
\cs_new:Npn \__bnvs_DEBUG_on: {
  \cs_set:Npn \__bnvs_DEBUG:n {
    \exp_args:Nx
    \__bnvs_DEBUG_:nn
    {   \prg_replicate:nn {\l__bnvs_group_int} { } \space }
  }
}
\cs_new:Npn \__bnvs_DEBUG_off: {
  \cs_set_eq:NN \__bnvs_DEBUG:n \use_none:n
}
\__bnvs_DEBUG_off:
\cs_generate_variant:Nn \__bnvs_DEBUG:n { x, V }
\int_zero_new:N \l__bnvs_group_int
\cs_set:Npn \__bnvs_group_begin: {
  \group_begin:
  \int_incr:N \l__bnvs_group_int
}
\cs_set:Npn \__bnvs_group_end: {
  \group_end:
}
\cs_new:Npn \__bnvs_LOG:nn #1 #2 {
  \__bnvs_DEBUG:x { #1~#2 }
}
\cs_new:Npn \__bnvs_DEBUG:nn #1 {
  \exp_args:Nx
  \__bnvs_LOG:nn
  {   \prg_replicate:nn {\l__bnvs_group_int + 1} {#1} }
}
\cs_generate_variant:Nn \__bnvs_DEBUG:nn { nx, nV }
```

5.3 Local variables

We make heavy use of local variables and function scopes. Many functions are executed within a TEX group, which ensures no name collision with the caller stack. In that case, variables need not follow exactly the LATEX3 naming convention: we do not specialize with the module name. On execution, next initialization instructions declare the variables as side effect.

```
\int_if_exist:NF \l_depth_int {
  \int_new:N  \l_depth_int
}
\bool_new:N \l__bnvs_no_counter_bool
\bool_new:N \l__bnvs_no_range_bool
\bool_new:N \l__bnvs_continue_bool
\bool_new:N \l__bnvs_in_frame_bool
\bool_set_false:N \l__bnvs_in_frame_bool
```

5.4 Infinite loop management

Unending recursivity is managed here.

\g__bnvs_call_int

```
\int_zero_new:N  \g__bnvs_call_int
\int_const:Nn \c__bnvs_max_call_int { 2048 }
```

(*End definition for* `\g__bnvs_call_int`.)

__bnvs_call_reset: `\__bnvs_call_reset:`

Reset the call stack counter.

```
\cs_set:Npn  \__bnvs_call_reset: {
  \int_gset:Nn \g__bnvs_call_int { \c__bnvs_max_call_int }
}
```

__bnvs_call:*TF* `\__bnvs_call_do:TF {`⟨ *true code* ⟩`} {`⟨ *false code* ⟩`}`

Decrement the `\g__bnvs_call_int` counter globally and execute ⟨ *true code* ⟩ if we have not reached 0, ⟨ *false code* ⟩ otherwise.

```
\prg_new_conditional:Npnn  \__bnvs_call: { T, F, TF } {
  \int_gdecr:N \g__bnvs_call_int
  \int_compare:nNnTF \g__bnvs_call_int > 0 {
    \prg_return_true:
  } {
    \prg_return_false:
  }
}
```

5.5 Overlay specification

5.5.1 In slide range definitions

\g__bnvs_prop ⟨*key*⟩–⟨*value*⟩ property list to store the named slide lists. The basic keys are, assuming ⟨*id*⟩!⟨*name*⟩ is a fully qualified slide list name,

⟨***id***⟩**!**⟨***name***⟩**/A** for the first index

⟨***id***⟩**!**⟨***name***⟩**/L** for the length when provided

⟨***id***⟩**!**⟨***name***⟩**/Z** for the last index when provided

⟨***id***⟩**!**⟨***name***⟩**/C** for the counter value, when used

⟨***id***⟩**!**⟨***name***⟩**/C0** for initial value of the counter (when reset)

Other keys are eventually used to cache results when some attributes are defined from other slide ranges. They are characterized by a '`//`'.

⟨*id*⟩`!`⟨*name*⟩`//A` for the cached static value of the first index

⟨*id*⟩`!`⟨*name*⟩`//Z` for the cached static value of the last index

⟨*id*⟩`!`⟨*name*⟩`//L` for the cached static value of the length

⟨*id*⟩`!`⟨*name*⟩`//N` for the cached static value of the next index

The implementation is private, in particular, keys may change in future versions.

```
\prop_new:N \g__bnvs_prop
```

(*End definition for* `\g__bnvs_prop`.)

__bnvs_gput:nn
__bnvs_gput:nV
__bnvs_gprovide:nn
__bnvs_gprovide:nV
__bnvs_item:n
__bnvs_get:nN
__bnvs_gremove:n
__bnvs_gclear:n
__bnvs_gclear_cache:n
__bnvs_gclear:

__bnvs_gput:nn {⟨*key*⟩} {⟨*value*⟩}
__bnvs_gprovide:nn {⟨*key*⟩} {⟨*value*⟩}
__bnvs_item:n {⟨*key*⟩}
__bnvs_get:n {⟨*key*⟩} ⟨*tl variable*⟩
__bnvs_gremove:n {⟨*key*⟩}
__bnvs_gclear:n {⟨*key*⟩}
__bnvs_gclear_cache:n {⟨*key*⟩}
__bnvs_gclear:

Convenient shortcuts to manage the storage, it makes the code more concise and readable. This is a wrapper over LaTeX3 eponym functions, except `\__bnvs_gprovide:nn` which meaning is straightforward.

```
\cs_new:Npn \__bnvs_gput:nn #1 #2 {
\__bnvs_DEBUG:x {\string\__bnvs_gput:nn/key:#1/value:#2/}
  \prop_gput:Nnn \g__bnvs_prop { #1 } { #2 }
}
\cs_new:Npn \__bnvs_gprovide:nn #1 #2 {
\__bnvs_DEBUG:x {\string\__bnvs_gprovide:nn/key:#1/value:#2/}
  \prop_if_in:NnF \g__bnvs_prop { #1 } {
    \prop_gput:Nnn \g__bnvs_prop { #1 } { #2 }
  }
}
\cs_new:Npn \__bnvs_item:n {
  \prop_item:Nn \g__bnvs_prop
}
\cs_new:Npn \__bnvs_get:nN {
  \prop_get:NnN \g__bnvs_prop
}
\cs_new:Npn \__bnvs_gremove:n {
  \prop_gremove:Nn \g__bnvs_prop
}
\cs_new:Npn \__bnvs_gclear:n #1 {
  \clist_map_inline:nn { A, L, Z, C, C0, /, /A, /L, /Z, /N } {
    \__bnvs_gremove:n { #1 / ##1 }
  }
}
\cs_new:Npn \__bnvs_gclear_cache:n #1 {
  \clist_map_inline:nn { /A, /L, /Z, /N } {
    \__bnvs_gremove:n { #1 / ##1 }
  }
}
\cs_new:Npn \__bnvs_gclear: {
  \prop_gclear:N \g__bnvs_prop
}
\cs_generate_variant:Nn \__bnvs_gput:nn { nV }
\cs_generate_variant:Nn \__bnvs_gprovide:nn { nV }
```

__bnvs_if_in_p:n ⋆
__bnvs_if_in_p:V ⋆
__bnvs_if_in:nTF ⋆
__bnvs_if_in:VTF ⋆

__bnvs_if_in_p:n {⟨*key*⟩}
__bnvs_if_in:nTF {⟨*key*⟩} {⟨*true code*⟩} {⟨*false code*⟩}

Convenient shortcuts to test for the existence of some key, it makes the code more concise and readable.

```
\prg_new_conditional:Npnn \__bnvs_if_in:n #1 { p, T, F, TF } {
  \prop_if_in:NnTF \g__bnvs_prop { #1 } {
```

```
    \prg_return_true:
  } {
    \prg_return_false:
  }
}
\prg_generate_conditional_variant:Nnn \__bnvs_if_in:n {V} { p, T, F, TF }
```

__bnvs_get:nNTF
__bnvs_get:nnNTF

```
\__bnvs_get:nNTF {⟨key⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}
\__bnvs_get:nnNTF {⟨id⟩} {⟨key⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}
```

Convenient shortcuts to retrieve the value with branching, it makes the code more concise and readable. Execute ⟨*true code*⟩ when the item is found, ⟨*false code*⟩ otherwise. In the latter case, the content of the ⟨*tl variable*⟩ is undefined. NB: the predicate won't work because `\prop_get:NnNTF` is not expandable.

```
\prg_new_conditional:Npnn \__bnvs_get:nN #1 #2 { T, F, TF } {
  \prop_get:NnNTF \g__bnvs_prop { #1 } #2 {
\__bnvs_DEBUG:x { \string\__bnvs_get:nN\space TRUE/
  #1/\string#2:#2/ }
    \prg_return_true:
  } {
\__bnvs_DEBUG:x { \string\__bnvs_get:nN\space FALSE/#1/\string#2/ }
    \prg_return_false:
  }
}
```

5.5.2 Regular expressions

\c__bnvs_name_regex

The name of a slide range consists of a non void list of alphanumerical characters and underscore, but with no leading digit.

```
\regex_const:Nn \c__bnvs_name_regex {
  [[:alpha:]_][[:alnum:]_]*
}
```

(*End definition for* `\c__bnvs_name_regex`.)

\c__bnvs_id_regex

The name of a slide range consists of a non void list of alphanumerical characters and underscore, but with no leading digit.

```
\regex_const:Nn \c__bnvs_id_regex {
  (?: \ur{c__bnvs_name_regex} | [?]* ) ? !
}
```

(*End definition for* `\c__bnvs_id_regex`.)

\c__bnvs_path_regex

A sequence of .⟨*positive integer*⟩ items representing a path.

```
\regex_const:Nn \c__bnvs_path_regex {
  (?: \. [+-]? \d+ )*
}
```

(*End definition for* `\c__bnvs_path_regex`.)

\c__bnvs_key_regex
\c__bnvs_A_key_Z_regex

A key is the name of a slide range possibly followed by positive integer attributes using a dot syntax. The '`A_key_Z`' variant matches the whole string.

```
\regex_const:Nn \c__bnvs_key_regex {
  \ur{c__bnvs_id_regex} ?
  \ur{c__bnvs_name_regex}
  \ur{c__bnvs_path_regex}
}
\regex_const:Nn \c__bnvs_A_key_Z_regex {
```

2: slide ⟨*id*⟩

3: question mark, when ⟨*id*⟩ is empty

4: The range name

```
      \A ( ( \ur{c__bnvs_id_regex} ? ) \ur{c__bnvs_name_regex} )
```

5: the path, if any.

```
      ( \ur{c__bnvs_path_regex} ) \Z
    }

```

(*End definition for* `\c__bnvs_key_regex` *and* `\c__bnvs_A_key_Z_regex`.)

\c__bnvs_colons_regex

For ranges defined by a colon syntax.

```
\regex_const:Nn \c__bnvs_colons_regex { :(:+)? }
```

(*End definition for* `\c__bnvs_colons_regex`.)

\c__bnvs_list_regex

A comma separated list between square brackets.

```
\regex_const:Nn \c__bnvs_list_regex {
  \A \[ \s*
```

Capture groups:

- 2: the content between the brackets, outer spaces trimmed out

```
    ( [^\] %[---
    ]*? )
  \s* \] \Z
}
```

(*End definition for* `\c__bnvs_list_regex`.)

\c__bnvs_split_regex

Used to parse slide list overlay specifications in queries. Next are the 10 capture groups. Group numbers are 1 based because the regex is used in splitting contexts where only capture groups are considered and not the whole match.

```
\regex_const:Nn \c__bnvs_split_regex {
  \s* ( ? :
```

We start with '`++`' instrussions [2].

[2] At the same time an instruction and an expression... this is a synonym of exprection

- 1: ⟨*name*⟩ of a slide range
- 2: ⟨*id*⟩ of a slide range plus the exclamation mark

```
\+\+ ( ( \ur{c__bnvs_id_regex}? ) \ur{c__bnvs_name_regex} )
```

- 3: optionally followed by an integer path

```
( \ur{c__bnvs_path_regex} ) (?: \. n )?
```

We continue with other expressions

- 4: fully qualified ⟨*name*⟩ of a slide range,
- 5: ⟨*id*⟩ of a slide range plus the exclamation mark (to manage void ⟨*id*⟩)

```
| ( ( \ur{c__bnvs_id_regex}? ) \ur{c__bnvs_name_regex} )
```

- 6: optionally followed by an integer path

```
( \ur{c__bnvs_path_regex} )
```

Next comes another branching

```
(?:
```

- 7: the ⟨*length*⟩ attribute

```
    \. l(e)ngth
```

- 8: the ⟨*last*⟩ attribute

```
  |  \. l(a)st
```

- 9: the ⟨*next*⟩ attribute

```
  |  \. ne(x)t
```

- 10: the ⟨*range*⟩ attribute

```
  |  \. (r)ange
```

- 11: the ⟨*n*⟩ attribute

```
  |  \. (n)
```

- 12: the poor man integer expression after '`+=`', which is the longest sequence of black characters, which ends just before a space or at the very last character. This tricky definition allows quite any algebraic expression, even those involving parenthesis.

```
      (?: \s* \+= \s* ( \S+ ) )?


  )?

) \s*
}
```

(*End definition for* `\c__bnvs_split_regex`.)

5.5.3 beamer.cls interface

Work in progress.

```
\RequirePackage{keyval}
\define@key{beamerframe}{beanoves~id}[]{
  \tl_set:Nx \l__bnvs_id_tl { #1 ! }
  \__bnvs_DEBUG_on:
  \__bnvs_DEBUG:x {THIS_IS_KEY}
  \__bnvs_DEBUG_off:
}
\AddToHook{env/beamer@frameslide/before}{
  \bool_set_true:N \l__bnvs_in_frame_bool
  \__bnvs_DEBUG_on:
  \__bnvs_DEBUG:x {THIS_IS_BEFORE}
  \__bnvs_DEBUG_off:
}
\AddToHook{env/beamer@frameslide/after}{
  \bool_set_false:N \l__bnvs_in_frame_bool
  \__bnvs_DEBUG_on:
  \__bnvs_DEBUG:x {THIS_IS_BEFORE}
  \__bnvs_DEBUG_off:
}
\AddToHook{cmd/frame/before}{
  \tl_set:Nn \l__bnvs_id_tl { ?! }
  \__bnvs_DEBUG_on:
  \__bnvs_DEBUG:x {THIS_IS_FRAME}
  \__bnvs_DEBUG_off:
}
```

5.5.4 Defining named slide ranges

`\__bnvs_parse:Nnn` `\__bnvs_parse:Nnn` ⟨*command*⟩ {⟨*key*⟩} {⟨*definition*⟩}

Auxiliary function called within a group. ⟨*key*⟩ is the slide range key, including eventually a dotted integer path and a slide identifier, ⟨*definition*⟩ is the corresponding definition. ⟨*command*⟩ is `\__bnvs_range:nVVV` at runtime.

`\l_match_seq` Local storage for the match result.

(*End definition for* `\l_match_seq`. *This variable is documented on page* **??**.)

`\__bnvs_range:nnnn` `\__bnvs_range:nnnn` {⟨*key*⟩} {⟨*first*⟩} {⟨*length*⟩} {⟨*last*⟩}
`\__bnvs_range:nVVV` `\__bnvs_range_alt:nnnn` {⟨*key*⟩} {⟨*first*⟩} {⟨*length*⟩} {⟨*last*⟩}
`\__bnvs_range_alt:nnnn` `\__bnvs_range:Nnnnn` ⟨*cmd*⟩ {⟨*key*⟩} {⟨*first*⟩} {⟨*length*⟩} {⟨*last*⟩}
`\__bnvs_range_alt:nVVV`
`\__bnvs_range:Nnnnn`

Auxiliary function called within a group. Setup the model to define a range. The alt variant does not override an already existing value.

Implementation detail: the core functionality is implemented in the function `\__bnvs_range:Nnnnn` which first argument is `\__bnvs_gput:nn` for `\__bnvs_range:nnnn` and `\__bnvs_gprovide:nn` for `\__bnvs_range_alt:nnnn`.

```
\cs_new:Npn \__bnvs_range:Nnnnn #1 #2 #3 #4 #5 {
\__bnvs_DEBUG:x {\string\__bnvs_range:Nnnnn/\string#1/#2/#3/#4/#5/}
  \tl_if_empty:nTF { #3 } {
```

```
    \tl_if_empty:nTF { #4 } {
      \tl_if_empty:nTF { #5 } {
        \msg_error:nnn { beanoves } { :n } { Not~a~range:~:~#2 }
      } {
        #1 { #2/Z } { #5 }
      }
    } {
      #1 { #2/L } { #4 }
      \tl_if_empty:nF { #5 } {
        #1 { #2/Z } { #5 }
        #1 { #2/A } { #2.last - (#2.length) + 1 }
      }
    }
  } {
    #1 { #2/A } { #3 }
    \tl_if_empty:nTF { #4 } {
      \tl_if_empty:nF { #5 } {
        #1 { #2/Z } { #5 }
        #1 { #2/L } { #2.last - (#2.1) + 1 }
      }
    } {
      #1 { #2/L } { #4 }
      #1 { #2/Z } { #2.1 + #2.length - 1 }
    }
  }
}
\cs_new:Npn \__bnvs_range:nnnn #1 {
  \__bnvs_gclear:n { #1 }
  \__bnvs_range:Nnnnn \__bnvs_gput:nn { #1 }
}
\cs_generate_variant:Nn \__bnvs_range:nnnn { nVVV }
\cs_new:Npn \__bnvs_range_alt:nnnn #1 {
  \__bnvs_gclear_cache:n { #1 }
  \__bnvs_range:Nnnnn \__bnvs_gprovide:nn { #1 }
}
\cs_generate_variant:Nn \__bnvs_range_alt:nnnn { nVVV }
```

`\__bnvs_parse:Nn` `\__bnvs_parse:Nn` ⟨*command*⟩ {⟨*key*⟩}

Define a hidden range, for which slides are never shown. This is useful to conditionally show or hide a sequence of slides.

```
\cs_new:Npn \__bnvs_parse:Nn #1 #2 {
  \__bnvs_group_begin:
  \__bnvs_id_name_set:nNNTF { #2 } \l_id_tl \l_name_tl {
    \exp_args:Nx \__bnvs_gput:nn { \l_name_tl/ } { }
    \exp_args:NNNV
    \__bnvs_group_end:
    \tl_set:Nn \l__bnvs_id_tl \l__bnvs_id_tl
  } {
    \msg_error:nnn { beanoves } { :n } { Unexpected~key:~#2 }
    \__bnvs_group_end:
  }
}
```

__bnvs_do_parse:Nnn

__bnvs_do_parse:Nnn ⟨*command*⟩ {⟨*full name*⟩}

Auxiliary function for __bnvs_parse:Nn. ⟨*command*⟩ is __bnvs_range:nVVV at runtime and must have signature nVVV.

```
\cs_generate_variant:Nn \tl_if_empty:nTF { xTF }
\cs_new:Npn \__bnvs_do_parse:Nnn #1 #2 #3 {
\__bnvs_DEBUG:x {\string\__bnvs_do_parse:Nnn/\string#1/#2/#3}
```

This is not a list.

```
  \tl_clear:N \l_a_tl
  \tl_clear:N \l_b_tl
  \tl_clear:N \l_c_tl
  \regex_split:NnN \c__bnvs_colons_regex { #3 } \l_split_seq
  \seq_pop_left:NNT \l_split_seq \l_a_tl {
```

\l_a_tl may contain the ⟨*start*⟩.

```
    \seq_pop_left:NNT \l_split_seq \l_b_tl {
      \tl_if_empty:NTF \l_b_tl {
```

This is a one colon range.

```
        \seq_pop_left:NN \l_split_seq \l_b_tl
```

\l_b_tl may contain the ⟨*length*⟩.

```
        \seq_pop_left:NNT \l_split_seq \l_c_tl {
          \tl_if_empty:NTF \l_c_tl {
```

A :: was expected:

```
\msg_error:nnn { beanoves } { :n } { Invalid~range~expression(1):~#3 }
          } {
            \int_compare:nNnT { \tl_count:N \l_c_tl } > { 1 } {
\msg_error:nnn { beanoves } { :n } { Invalid~range~expression(2):~#3 }
            }
            \seq_pop_left:NN \l_split_seq \l_c_tl
```

\l_c_tl may contain the ⟨*end*⟩.

```
            \seq_if_empty:NF \l_split_seq {
\msg_error:nnn { beanoves } { :n } { Invalid~range~expression(3):~#3 }
            }
          }
        }
      } {
```

This is a two colon range.

```
        \int_compare:nNnT { \tl_count:N \l_b_tl } > { 1 } {
\msg_error:nnn { beanoves } { :n } { Invalid~range~expression(4):~#3 }
        }
        \seq_pop_left:NN \l_split_seq \l_c_tl
```

\l_c_tl contains the ⟨*end*⟩.

```
        \seq_pop_left:NNTF \l_split_seq \l_b_tl {
          \tl_if_empty:NTF \l_b_tl {
            \seq_pop_left:NN \l_split_seq \l_b_tl
```

`\l_b_tl` may contain the ⟨*length*⟩.

```
            \seq_if_empty:NF \l_split_seq {
\msg_error:nnn { beanoves } { :n } { Invalid~range~expression(5):~#3 }
            }
          } {
\msg_error:nnn { beanoves } { :n } { Invalid~range~expression(6):~#3 }
          }
        } {
          \tl_clear:N \l_b_tl
        }
      }
    }
  }
```

Providing both the ⟨*start*⟩, ⟨*length*⟩ and ⟨*end*⟩ of a range is not allowed, even if they happen to be consistent.

```
  \bool_if:nF {
    \tl_if_empty_p:N \l_a_tl
    || \tl_if_empty_p:N \l_b_tl
    || \tl_if_empty_p:N \l_c_tl
  } {
\msg_error:nnn { beanoves } { :n } { Invalid~range~expression(7):~#3 }
  }
  #1 { #2 } \l_a_tl \l_b_tl \l_c_tl
}
\cs_generate_variant:Nn \__bnvs_do_parse:Nnn { Nxn, Non }

\cs_new:Npn \__bnvs_parse_old:Nnn #1 #2 #3 {
  \__bnvs_group_begin:
  \regex_match:NnTF \c__bnvs_A_key_Z_regex { #2 } {
```

We got a valid key.

```
    \regex_extract_once:NnNTF \c__bnvs_list_regex { #3 } \l_match_seq {
```

This is a comma separated list, extract each item and go recursive.

```
      \exp_args:NNx
      \seq_set_from_clist:Nn \l_match_seq {
        \seq_item:Nn \l_match_seq { 2 }
      }
      \seq_map_indexed_inline:Nn \l_match_seq {
        \__bnvs_do_parse:Nnn #1 { #2.##1 } { ##2 }
      }
    } {
      \__bnvs_do_parse:Nnn #1 { #2 } { #3 }
    }
  } {
    \msg_error:nnn { beanoves } { :n } { Invalid~key:~#1 }
  }
  \__bnvs_group_end:
}
```

__bnvs_id_name_set:nNN*TF* `\__bnvs_id_name_set:nNNTF {⟨key⟩} ⟨id tl var⟩ ⟨full name tl var⟩ {⟨ true code⟩} {⟨ false code⟩}`

If the ⟨*key*⟩ is a key, put the name it defines into the ⟨*name tl var*⟩ with the current frame id prefix `\l__bnvs_id_tl` if none was given, then execute ⟨*true code*⟩. Otherwise execute ⟨*false code*⟩.

```
\prg_new_conditional:Npnn \__bnvs_id_name_set:nNN #1 #2 #3 { T, F, TF } {
  \__bnvs_group_begin:
  \regex_extract_once:NnNTF \c__bnvs_A_key_Z_regex { #1 } \l_match_seq {
    \tl_set:Nx #2 { \seq_item:Nn \l_match_seq 3 }
    \tl_if_empty:NTF #2 {
      \exp_args:NNNx
      \__bnvs_group_end:
      \tl_set:Nn #3 { \l__bnvs_id_tl #1 }
      \tl_set_eq:NN #2 \l__bnvs_id_tl
    } {
      \cs_set:Npn \:n ##1 {
        \__bnvs_group_end:
        \tl_set:Nn #2 { ##1 }
        \tl_set:Nn \l__bnvs_id_tl { ##1 }
      }
      \exp_args:NV
      \:n #2
      \tl_set:Nn #3 { #1 }
    }
\__bnvs_DEBUG:x { \string\__bnvs_id_name_set:nNN\space TRUE/#1/
  \string#2:#2/\string#3:#3/\string\l__bnvs_id_tl:\l__bnvs_id_tl/ }
    \prg_return_true:
  } {
    \__bnvs_group_end:
\__bnvs_DEBUG:x { \string\__bnvs_id_name_set:nNN\space FALSE/#1/
  \string#2/\string#3/ }
    \prg_return_false:
  }
}

\cs_new:Npn \__bnvs_parse:Nnn #1 #2 #3 {
\__bnvs_DEBUG:x {\string\__bnvs_parse:Nnn/\string#1/#2/#3/}
  \__bnvs_group_begin:
  \__bnvs_id_name_set:nNNTF { #2 } \l_id_tl \l_name_tl {
\__bnvs_DEBUG:x {key:#2/ID:\l_id_tl/NAME:\l_name_tl/}
    \regex_extract_once:NnNTF \c__bnvs_list_regex { #3 } \l_match_seq {
```

This is a comma separated list, extract each item and go recursive.

```
      \exp_args:NNx
      \seq_set_from_clist:Nn \l_match_seq {
        \seq_item:Nn \l_match_seq { 2 }
      }
      \seq_map_indexed_inline:Nn \l_match_seq {
        \__bnvs_do_parse:Nxn #1  { \l_name_tl.##1 } { ##2 }
      }
    } {
      \__bnvs_do_parse:Nxn #1 { \l_name_tl } { #3 }
    }
```

```
  } {
    \msg_error:nnn { beanoves } { :n } { Invalid~key:~#2 }
  }
```

We export `\l__bnvs_id_tl`:

```
  \exp_args:NNNV
  \__bnvs_group_end:
  \tl_set:Nn \l__bnvs_id_tl \l__bnvs_id_tl
}
```

`\Beanoves` `\Beanoves {⟨key--value list⟩}`

The keys are the slide range specifiers. When no value is provided, it defaults to `1`. On the contrary, ⟨*key–value*⟩ items are parsed by `\__bnvs_parse:Nnn`.

```
\NewDocumentCommand \Beanoves { sm } {
  \tl_if_eq:NnT \@currenvir { document } {
    \__bnvs_gclear:
  }
  \IfBooleanTF {#1} {
    \keyval_parse:nnn {
      \__bnvs_parse:Nn \__bnvs_range_alt:nVVV
    } {
      \__bnvs_parse:Nnn \__bnvs_range_alt:nVVV
    }
  } {
    \keyval_parse:nnn {
      \__bnvs_parse:Nn \__bnvs_range:nVVV
    } {
      \__bnvs_parse:Nnn \__bnvs_range:nVVV
    }
  }
  { #2 }
  \ignorespaces
}
```

If we use the frame `beanoves` option, we can provide default values to the various name ranges.

```
\define@key{beamerframe}{beanoves}{\Beanoves*{#1}}
```

5.5.5 Scanning named overlay specifications

Patch some beamer commands to support `?(...)` instructions in overlay specifications.

`\beamer@frame` `\beamer@frame {⟨overlay specification⟩}`
`\beamer@masterdecode` `\beamer@masterdecode {⟨overlay specification⟩}`

Preprocess ⟨*overlay specification*⟩ before beamer uses it.

`\l_ans_tl` Storage for the translated overlay specification, where `?(...)` instructions are replaced by their static counterparts.

(*End definition for* `\l_ans_tl`. *This variable is documented on page* ??.)

Save the original macro `\beamer@masterdecode` and then override it to properly preprocess the argument.

```
\cs_set_eq:NN \__bnvs_beamer@frame \beamer@frame
\cs_set:Npn \beamer@frame < #1 > {
  \__bnvs_group_begin:
  \tl_clear:N \l_ans_tl
  \__bnvs_scan:nNN { #1 } \__bnvs_eval:nN \l_ans_tl
  \exp_args:NNNV
  \__bnvs_group_end:
  \__bnvs_beamer@frame < \l_ans_tl >
}
\cs_set_eq:NN \__bnvs_beamer@masterdecode \beamer@masterdecode
\cs_set:Npn \beamer@masterdecode #1 {
  \__bnvs_group_begin:
  \tl_clear:N \l_ans_tl
  \__bnvs_scan:nNN { #1 } \__bnvs_eval:nN \l_ans_tl
  \exp_args:NNV
  \__bnvs_group_end:
  \__bnvs_beamer@masterdecode \l_ans_tl
}
```

`\__bnvs_scan:nNN` `\__bnvs_scan:nNN {`⟨*named overlay expression*⟩`}` ⟨*eval*⟩ ⟨*tl variable*⟩

Scan the ⟨*named overlay expression*⟩ argument and feed the ⟨*tl variable*⟩ replacing `?(...)` instructions by their static counterpart with help from the ⟨*eval*⟩ function, which is `\__bnvs_eval:nN`. A group is created to use local variables:

`\l_ans_tl`: is the token list that will be appended to ⟨*tl variable*⟩ on return.

`\l_depth_int` Store the depth level in parenthesis grouping used when finding the proper closing parenthesis balancing the opening parenthesis that follows immediately a question mark in a `?(...)` instruction.

(*End definition for* `\l_depth_int`. *This variable is documented on page* ??.)

`\l_query_tl` Storage for the overlay query expression to be evaluated.

(*End definition for* `\l_query_tl`. *This variable is documented on page* ??.)

`\l_token_seq` The ⟨*overlay expression*⟩ is split into the sequence of its tokens.

(*End definition for* `\l_token_seq`. *This variable is documented on page* ??.)

`\l_ask_bool` Whether a loop may continue. Controls the continuation of the main loop that scans the tokens of the ⟨*named overlay expression*⟩ looking for a question mark.

(*End definition for* `\l_ask_bool`. *This variable is documented on page* ??.)

`\l_query_bool` Whether a loop may continue. Controls the continuation of the secondary loop that scans the tokens of the ⟨*named overlay expression*⟩ looking for an opening parenthesis follow the question mark. It then controls the loop looking for the balanced closing parenthesis.

(*End definition for* `\l_query_bool`. *This variable is documented on page* ??.)

`\l_token_tl` Storage for just one token.

(*End definition for* `\l_token_tl`. *This variable is documented on page* ??.)

```
\cs_new:Npn \__bnvs_scan:nNN #1 #2 #3 {
  \__bnvs_group_begin:
  \tl_clear:N \l_ans_tl
  \int_zero:N \l_depth_int
  \seq_clear:N \l_token_seq
```

Explode the ⟨*named overlay expression*⟩ into a list of tokens:

```
  \regex_split:nnN {} { #1 } \l_token_seq
```

Run the top level loop to scan for a '`?`':

```
  \bool_set_true:N  \l_ask_bool
  \bool_while_do:Nn \l_ask_bool {
    \seq_pop_left:NN \l_token_seq \l_token_tl
    \quark_if_no_value:NTF \l_token_tl {
```

We reached the end of the sequence (and the token list), we end the loop here.

```
      \bool_set_false:N \l_ask_bool
    } {
```

`\l_token_tl` contains a 'normal' token.

```
      \tl_if_eq:NnTF \l_token_tl { ? } {
```

We found a '`?`', we first gobble tokens until the next '`(`', whatever they may be. In general, no tokens should be silently ignored.

```
        \bool_set_true:N \l_query_bool
        \bool_while_do:Nn \l_query_bool {
```

Get next token.

```
          \seq_pop_left:NN \l_token_seq \l_token_tl
          \quark_if_no_value:NTF \l_token_tl {
```

No opening parenthesis found, raise.

```
            \msg_fatal:nnx { beanoves } { :n } {Missing~'('%---)
              ~after~a~?:~#1}
          } {
            \tl_if_eq:NnT \l_token_tl { ( %)
            } {
```

We found the '`(`' after the '`?`'. Increment the parenthesis depth to 1 (on first passage).

```
              \int_incr:N \l_depth_int
```

Record the forthcomming content in the `\l_query_tl` variable, up to the next balancing '`)`'.

```
              \tl_clear:N \l_query_tl
              \bool_while_do:Nn \l_query_bool {
```

Get next token.

```
                \seq_pop_left:NN \l_token_seq \l_token_tl
                \quark_if_no_value:NTF \l_token_tl {
```

We reached the end of the sequence and the token list with no closing '`)`'. We raise and end both bool while loops. As recovery we feed `\l_query_tl` with the missing '`)`'. `\l_depth_int` is 0 whenever `\l_query_bool` is false.

```
                  \msg_error:nnx { beanoves } { :n } {Missing~%(---
                    `)':~#1 }
                  \int_do_while:nNnn \l_depth_int > 1 {
                    \int_decr:N \l_depth_int
                    \tl_put_right:Nn \l_query_tl {%(---
```

```
                    )}
                  }
                  \int_zero:N \l_depth_int
                  \bool_set_false:N \l_query_bool
                  \bool_set_false:N \l_ask_bool
                } {
                  \tl_if_eq:NnTF \l_token_tl { ( %---)
                  } {
```

We found a '(', increment the depth and append the token to `\l_query_tl`.

```
                    \int_incr:N \l_depth_int
                    \tl_put_right:NV \l_query_tl \l_token_tl
                  } {
```

This is not a '('.

```
                    \tl_if_eq:NnTF \l_token_tl { %(
                      )
                    } {
```

We found a ')', decrement the depth.

```
                      \int_decr:N \l_depth_int
                      \int_compare:nNnTF \l_depth_int = 0 {
```

The depth level has reached 0: we found our balancing parenthesis of the `?(...)` instruction. We can append the evaluated slide ranges token list to `\l_ans_tl` and stop the inner loop.

```
  \exp_args:NV #2 \l_query_tl \l_ans_tl
  \bool_set_false:N \l_query_bool
                      } {
```

The depth has not yet reached level 0. We append the ')' to `\l_query_tl` because it is not the end of sequence marker.

```
                        \tl_put_right:NV \l_query_tl \l_token_tl
                      }
```

Above ends the code for a positive depth.

```
                    } {
```

The scanned token is not a '(' nor a ')', we append it as is to `\l_query_tl`.

```
                      \tl_put_right:NV \l_query_tl \l_token_tl
                    }
                  }
                }
```

Above ends the code for Not a '('

```
              }
            }
```

Above ends the code for: Found the '(' after the '?'

```
          }
```

Above ends the code for not a no value quark.

```
        }
```

Above ends the code for the bool while loop to find the '(' after the '?'.

If we reached the end of the token list, then end both the current loop and its containing loop.

```
        \quark_if_no_value:NT \l_token_tl {
          \bool_set_false:N \l_query_bool
          \bool_set_false:N \l_ask_bool
        }
      } {
```

This is not a '?', append the token to right of `\l_ans_tl` and continue.

```
        \tl_put_right:NV \l_ans_tl \l_token_tl
      }
```

Above ends the code for the bool while loop to find a '(' after the '?'

```
    }
  }
```

Above ends the outer bool while loop to find '?' characters. We can append our result to ⟨*tl variable*⟩

```
  \exp_args:NNNV
  \__bnvs_group_end:
  \tl_put_right:Nn #3 \l_ans_tl
}
```

I

5.5.6 Resolution

Given a frame id, a name and an integer path, we resolve any intermediate standalone reference. For example, with `A=B` and `B=C`, `A` is resolved in `C`. But with `A=B+1` and `B=C`, `A` is not resolved in `C+1`. With `A=B:D` and `B=C`, `A` is not resolved in `C:D` as well.

`\__bnvs_extract_key:NNN`*TF*

`\__bnvs_extract_key:NNNTF` ⟨*id tl var*⟩ ⟨*name tl var*⟩ ⟨*path seq var*⟩ {⟨*true code*⟩} {⟨*false code*⟩}

Auxiliary function. ⟨*id tl var*⟩ contains a frame id whereas ⟨*name tl var*⟩ contains a range name. If we recognize a key, on return, ⟨*name tl var*⟩ contains the resolved name, ⟨*path seq var*⟩ is prepended with new integer path components, {⟨*true code*⟩} is executed, otherwise {⟨*false code*⟩} is executed.

```
\exp_args_generate:n { VVx }
\prg_new_conditional:Npnn \__bnvs_extract_key:NNN
    #1 #2 #3 { T, F, TF } {
\__bnvs_DEBUG:x { \string\__bnvs_extract_key:NNN/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn#3./}
  \__bnvs_group_begin:
  \exp_args:NNV
  \regex_extract_once:NnNTF \c__bnvs_A_key_Z_regex #2 \l_match_seq {
```

This is a correct key, update the path sequence accordingly

```
    \exp_args:Nx
    \tl_if_empty:nT { \seq_item:Nn \l_match_seq 3 } {
      \tl_put_left:NV #2 { #1 }
\__bnvs_DEBUG:x { VERIF~\tl_to_str:V #2 }
    }
    \exp_args:NNnx
```

```
    \seq_set_split:Nnn \l_split_seq . { \seq_item:Nn \l_match_seq 4 }
    \seq_remove_all:Nn \l_split_seq { }
    \seq_pop_left:NN \l_split_seq \l_a_tl
    \seq_if_empty:NTF \l_split_seq {
```

No new integer path component is added.

```
      \cs_set:Npn \:nn ##1 ##2 {
        \__bnvs_group_end:
        \tl_set:Nn #1 { ##1 }
        \tl_set:Nn #2 { ##2 }
      }
      \exp_args:NVV \:nn #1 #2
\__bnvs_DEBUG:x { END/\string#1:#1/\string#2:#2/ }
    } {
```

Some new integer path components are added.

```
\__bnvs_DEBUG:x { \string\__bnvs_extract_key:NNN/\string#1:#1/
  \string#2:#2/\string#3:\seq_use:Nn#3./
  \string\l_split_seq:\seq_use:Nn\l_split_seq./ }
      \cs_set:Npn \:nnn ##1 ##2 ##3 {
        \__bnvs_group_end:
        \tl_set:Nn #1 { ##1 }
        \tl_set:Nn #2 { ##2 }
        \seq_set_split:Nnn #3 . { ##3 }
        \seq_remove_all:Nn #3 { }
      }
\__bnvs_DEBUG:n{1}
      \exp_args:NVVx
      \:nnn #1 #2 {
        \seq_use:Nn \l_split_seq . . \seq_use:Nn #3 .
      }
\__bnvs_DEBUG:x { END/\string#1:#1/\string#2:#2/
  \string#3:\seq_use:Nn #3 . /
  \string\l_split_seq:\seq_use:Nn \l_split_seq . / }
    }
\__bnvs_DEBUG:x { \string\__bnvs_extract_key:NNN\space TRUE/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn #3 . /}
    \prg_return_true:
  } {
    \__bnvs_group_end:
\__bnvs_DEBUG:x { \string\__bnvs_extract_key:NNN\space FALSE/
    \string#1/\string#2/\string#3/}
    \prg_return_false:
  }
}
```

__bnvs_resolve:NNN*TF* `\__bnvs_resolve:NNNTF` ⟨*id tl var*⟩ ⟨*name tl var*⟩ ⟨*path seq var*⟩ {⟨*true code*⟩} {⟨*false code*⟩}

When too many nested calls occurred, {⟨*false code*⟩} is executed directly. ⟨*id tl var*⟩, ⟨*name tl var*⟩ and ⟨*path seq var*⟩ are meant to contain proper information. On input, {⟨*id tl var*⟩} contains a frame id, {⟨*name tl var*⟩} contains a range name and {⟨*path seq var*⟩} contains the components of an integer path, possibly empty. On return, ⟨*id tl var*⟩ contains the frame id used, ⟨*name tl var*⟩ contains the resolved range name and ⟨*path seq var*⟩ contains the sequence of integer path components that could not be resolved. To resolve a path, $\langle name_0\rangle.\langle i_1\rangle.\langle i_2\rangle...\langle i_n\rangle$ is turned into $\langle name_1\rangle.\langle i_2\rangle...\langle i_n\rangle$ where $\langle name_0\rangle.\langle i_1\rangle$ is $\langle name_1\rangle$, then $\langle name_2\rangle.\langle i_3\rangle...\langle i_n\rangle$ where $\langle name_1\rangle.\langle i_2\rangle$ is $\langle name_2\rangle$... If the above rule does not apply, $\langle name_0\rangle.\langle i_1\rangle.\langle i_2\rangle...\langle i_n\rangle$ may turn into $\langle name_2\rangle.\langle i_3\rangle...\langle i_n\rangle$ when $\langle name_0\rangle.\langle i_1\rangle.\langle i_2\rangle$ is $\langle name_2\rangle$... The algorithm is not yet more clever. The resolution algorithm is quite straightforward:

1. If ⟨*name tl var*⟩ content is the name of an unlimited range, and the first item of this range is exactly another name range with eventually a heading frame identifier or a trailing integer path, then ⟨*name tl var*⟩ is replaced by this name, the ⟨*id tl var*⟩ and `\l__bnvs_id_tl` are updates accordingly and the ⟨*path seq var*⟩ is prepended with the integer path.

2. If ⟨*path seq var*⟩ is not empty, append to the right of ⟨*name tl var*⟩ after a separating dot, all its left elements but the last one and loop. Otherwise return. None of the tl variables must be one of `\l_a_tl`, `\l_b_tl` or `\l_c_tl`. None of the seq variables must be one of `\l_a_seq`, `\l_b_seq`.

```
\prg_new_conditional:Npnn \__bnvs_resolve:NNN
    #1 #2 #3 { T, F, TF } {
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNN/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn #3./ }
 \__bnvs_group_begin:
```

Local variables:

- `\l_a_tl` contains the name with a partial index path currently resolved.
- `\l_a_seq` contains the index path components currently resolved.
- `\l_b_tl` contains the resolution.
- `\l_b_seq` contains the index path components to be resolved.

```
 \seq_set_eq:NN \l_a_seq #3
 \seq_clear:N \l_b_seq
 \cs_set:Npn \loop: {
   \__bnvs_call:TF {
     \tl_set_eq:NN \l_a_tl #2
     \seq_if_empty:NTF \l_a_seq {
       \exp_args:Nx
       \__bnvs_get:nNTF { \l_a_tl / L } \l_b_tl {
         \cs_set:Nn \loop: { \return_true: }
       } {
         \get_extract:F {
```

Unknown key ⟨\l_a_tl⟩/A or the value for key ⟨\l_a_tl⟩/A does not fit.

```
            \cs_set:Nn \loop: { \return_true: }
          }
        }
      } {
        \tl_put_right:Nx \l_a_tl { . \seq_use:Nn \l_a_seq . }
        \get_extract:F {
          \seq_pop_right:NNT \l_a_seq \l_c_tl {
            \seq_put_left:NV \l_b_seq \l_c_tl
          }
        }
      }
      \loop:
    } {
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNN\space~TOO~MANY~CALLS/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn #3./ }
      \__bnvs_group_end:
      \prg_return_false:
    }
  }
  \cs_set:Npn \get_extract:F ##1 {
    \exp_args:Nx
    \__bnvs_get:nNTF { \l_a_tl / A } \l_b_tl {
\__bnvs_DEBUG:x { RESOLUTION:~\l_a_tl / A=>\l_b_tl}
      \__bnvs_extract_key:NNNTF #1 \l_b_tl \l_b_seq {
        \tl_set_eq:NN #2 \l_b_tl
        \seq_set_eq:NN #3 \l_b_seq
        \seq_set_eq:NN \l_a_seq \l_b_seq
        \seq_clear:N \l_b_seq
      } { ##1 }
    } { ##1 }
  }
  \cs_set:Npn \return_true: {
    \cs_set:Npn \:nnn ####1 ####2 ####3 {
      \__bnvs_group_end:
      \tl_set:Nn #1 { ####1 }
      \tl_set:Nn #2 { ####2 }
      \seq_set_split:Nnn #3 . { ####3 }
      \seq_remove_all:Nn #3 { }
    }
    \exp_args:NVVx
    \:nnn #1 #2 {
      \seq_use:Nn #3 .
    }
\__bnvs_DEBUG:x { ...\string\__bnvs_resolve:NNN\space TRUE/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn #3./ }
    \prg_return_true:
  }
  \loop:
}
```

__bnvs_resolve_n:NNNTF*TF* __bnvs_resolve_n:NNNTF ⟨*id tl var*⟩ ⟨*name tl var*⟩ ⟨*path seq var*⟩ {⟨ *true code*⟩} {⟨ ⟩} false code

The difference with the function above without _n is that resolution is performed only when there is an integer path afterwards

```
\prg_new_conditional:Npnn \__bnvs_resolve_n:NNN
    #1 #2 #3 { T, F, TF } {
\__bnvs_DEBUG:x { \string\__bnvs_resolve_n:NNN/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn #3./ }
 \__bnvs_group_begin:
```

Local variables:

- \l_a_tl contains the name with a partial index path currently resolved.
- \l_a_seq contains the index path components currently resolved.
- \l_b_tl contains the resolution.
- \l_b_seq contains the index path components to be resolved.

```
 \seq_set_eq:NN \l_a_seq #3
 \seq_clear:N \l_b_seq
 \cs_set:Npn \loop: {
   \__bnvs_call:TF {
     \tl_set_eq:NN \l_a_tl #2
     \seq_if_empty:NTF \l_a_seq {
       \exp_args:Nx
       \__bnvs_get:nNTF { \l_a_tl / L } \l_b_tl {
         \cs_set:Nn \loop: { \return_true: }
       } {
         \seq_if_empty:NTF \l_b_seq {
           \cs_set:Nn \loop: { \return_true: }
         } {
           \get_extract:F {
```

Unknown key ⟨\l_a_tl⟩/A or the value for key ⟨\l_a_tl⟩/A does not fit.

```
             \cs_set:Nn \loop: { \return_true: }
           }
         }
       }
     } {
       \tl_put_right:Nx \l_a_tl { . \seq_use:Nn \l_a_seq . }
       \get_extract:F {
         \seq_pop_right:NNT \l_a_seq \l_c_tl {
           \seq_put_left:NV \l_b_seq \l_c_tl
         }
       }
     }
     \loop:
   } {
\__bnvs_DEBUG:x { \string\__bnvs_resolve_n:NNN\space~TOO~MANY~CALLS/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn #3./ }
     \__bnvs_group_end:
     \prg_return_false:
   }
```

```
  }
  \cs_set:Npn \get_extract:F ##1 {
    \exp_args:Nx
    \__bnvs_get:nNTF { \l_a_tl / A } \l_b_tl {
\__bnvs_DEBUG:x { RESOLUTION:~\l_a_tl / A=>\l_b_tl}
      \__bnvs_extract_key:NNNTF #1 \l_b_tl \l_b_seq {
        \tl_set_eq:NN #2 \l_b_tl
        \seq_set_eq:NN #3 \l_b_seq
        \seq_set_eq:NN \l_a_seq \l_b_seq
        \seq_clear:N \l_b_seq
      } { ##1 }
    } { ##1 }
  }
  \cs_set:Npn \return_true: {
    \cs_set:Npn \:nnn ####1 ####2 ####3 {
      \__bnvs_group_end:
      \tl_set:Nn #1 { ####1 }
      \tl_set:Nn #2 { ####2 }
      \seq_set_split:Nnn #3 . { ####3 }
      \seq_remove_all:Nn #3 { }
    }
    \exp_args:NVVx
    \:nnn #1 #2 {
      \seq_use:Nn #3 .
    }
\__bnvs_DEBUG:x { ...\string\__bnvs_resolve_n:NNN\space TRUE/
    \string#1:#1/\string#2:#2/\string#3:\seq_use:Nn #3./ }
    \prg_return_true:
  }
  \loop:
}
```

`\__bnvs_resolve:NNNNTF`*TF*

`\__bnvs_resolve:NNNNTF` ⟨*cs:nn*⟩ ⟨*id tl var*⟩ ⟨*name tl var*⟩ ⟨*path seq var*⟩ {⟨ *true code*⟩} {⟨ ⟩} false code

When too many nested calls occurred, {⟨*false code*⟩} is executed directly. ⟨*id tl var*⟩, ⟨*name tl var*⟩ and ⟨*path seq var*⟩ are meant to contain proper information. To resolve a path, $\langle name_0\rangle.\langle i_1\rangle.\langle i_2\rangle...\langle i_n\rangle$ is turned into $\langle name_1\rangle.\langle i_2\rangle...\langle i_n\rangle$ where $\langle name_0\rangle.\langle i_1\rangle$ is $\langle name_1\rangle$, then $\langle name_2\rangle.\langle i_3\rangle...\langle i_n\rangle$ where $\langle name_1\rangle.\langle i_2\rangle$ is $\langle name_2\rangle$... If the above rule does not apply, $\langle name_0\rangle.\langle i_1\rangle.\langle i_2\rangle...\langle i_n\rangle$ may turn into $\langle name_2\rangle.\langle i_3\rangle...\langle i_n\rangle$ when $\langle name_0\rangle.\langle i_1\rangle.\langle i_2\rangle$ is $\langle name_2\rangle$... We try to match the longest sequence of components first. The algorithm is not yet more clever. In general, ⟨*cs:nn*⟩ is just `\use_i:nn` but for in place incrementation, we must resolve only when there is an integer path. See the implementation of the `\__bnvs_if_append:...` conditionals.

```
\prg_new_conditional:Npnn \__bnvs_resolve:NNNN
    #1 #2 #3 #4 { T, F, TF } {
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNNN/
    \string#1/\string#2:#2/\string#3:#3/\string#4:\seq_use:Nn #4./ }
  #1 {
    \__bnvs_group_begin:
```

`\l_a_tl` contains the name with a partial index path currently resolved. `\l_a_seq` contains the remaining index path components to be resolved. `\l_b_seq` contains the current index path components to be resolved.

```
    \tl_set_eq:NN \l_a_tl #3
    \seq_set_eq:NN \l_a_seq #4
    \tl_clear:N \l_b_tl
    \seq_clear:N \l_b_seq
    \cs_set:Npn \return_true: {
      \cs_set:Npn \:nnn ####1 ####2 ####3 {
        \__bnvs_group_end:
        \tl_set:Nn #2 { ####1 }
        \tl_set:Nn #3 { ####2 }
        \seq_set_split:Nnn #4 . { ####3 }
        \seq_remove_all:Nn #4 { }
      }
      \exp_args:NVVx
      \:nnn #2 #3 {
        \seq_use:Nn #4 .
      }
\__bnvs_DEBUG:x { ...\string\__bnvs_resolve:NNNN\space TRUE/
  \string#1/\string#2:#2/\string#3:#3/\string#4:\seq_use:Nn #4./ }
      \prg_return_true:
    }
    \cs_set:Npn \branch:n ##1 {
      \seq_pop_right:NNTF \l_a_seq \l_b_tl {
        \seq_put_left:NV \l_b_seq \l_b_tl
\__bnvs_DEBUG:x {\string\__bnvs_resolve:NNNN\space POP~TRUE~##1}
\__bnvs_DEBUG:x {\string\l_b_tl:\l_b_tl }
\__bnvs_DEBUG:x {\string\l_a_seq:\seq_count:N\l_a_seq/\seq_use:Nn \l_a_seq ./ }
\__bnvs_DEBUG:x {\string\l_b_seq:\seq_count:N\l_b_seq/\seq_use:Nn \l_b_seq ./ }
        \tl_set:Nn \l_a_tl { #3 . }
        \tl_put_right:Nx \l_a_tl { \seq_use:Nn \l_a_seq . }
      } {
        \cs_set_eq:NN \loop: \return_true:
      }
    }
    \cs_set:Npn \branch:FF ##1 ##2 {
      \exp_args:Nx
      \__bnvs_get:nNTF { \l_a_tl / A } \l_b_tl {
        \__bnvs_extract_key:NNNTF #2 \l_b_tl \l_b_seq {
          \tl_set_eq:NN #3 \l_b_tl
          \seq_set_eq:NN #4 \l_b_seq
          \seq_set_eq:NN \l_a_seq \l_b_seq
        } { ##1 }
      } { ##2 }
    }
    \cs_set:Npn \extract_key:F {
      \__bnvs_extract_key:NNNTF #2 \l_b_tl \l_b_seq {
        \tl_set_eq:NN #3 \l_b_tl
        \seq_set_eq:NN #4 \l_b_seq
        \seq_set_eq:NN \l_a_seq \l_b_seq
      }
    }
    \cs_set:Npn \loop: {
      \__bnvs_call:TF {
        \exp_args:Nx
        \__bnvs_get:nNTF { \l_a_tl / L } \l_b_tl {
```

If there is a length, no resolution occurs.

```
          \branch:n { 1 }
        } {
          \seq_pop_right:NNTF \l_a_seq \l_c_tl {
            \seq_clear:N \l_b_seq
            \tl_set:Nn \l_a_tl { #3 . }
            \tl_put_right:Nx \l_a_tl { \seq_use:Nn \l_a_seq . . }
            \tl_put_right:NV \l_a_tl \l_c_tl
            \branch:FF {
```

The value for key ⟨\l_a_tl⟩/L is not just a (qualified) name.

```
\seq_put_left:NV \l_b_seq \l_c_tl
            } {
```

Unknown key ⟨\l_a_tl⟩/L.

```
\seq_put_left:NV \l_b_seq \l_c_tl
            }
          } {
            \branch:FF {
              \cs_set_eq:NN \loop: \return_true:
            } {
              \cs_set:Npn \loop: {
                \__bnvs_group_end:
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNNN\space FALSE/
  \string#1/\string#2:#2/\string#3:#3/\string#4:\seq_use:Nn #4./
  \g__bnvs_call_int :\int_use:N\g__bnvs_call_int/
}
                \prg_return_false:
              }
            }
          }
        }
      } {
        \cs_set:Npn \loop: {
          \__bnvs_group_end:
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNNN\space FALSE/
  \string#1/\string#2:#2/\string#3:#3/\string#4:\seq_use:Nn #4./
  \g__bnvs_call_int :\int_use:N\g__bnvs_call_int/
}
          \prg_return_false:
        }
      }
      \loop:
    }
    \loop:
  } {
    \prg_return_true:
  }
}
\prg_new_conditional:Npnn \__bnvs_resolve_OLD:NNNN
    #1 #2 #3 #4 { T, F, TF } {
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNNN/
    \string#1/\string#2:#2/\string#3:#3/\string#4:\seq_use:Nn #4./ }
  #1 {
    \__bnvs_group_begin:
```

`\l_a_tl` contains the name with a partial index path to be resolved. `\l_a_seq` contains the remaining index path components to be resolved.

```
    \tl_set_eq:NN \l_a_tl #3
    \seq_set_eq:NN \l_a_seq #4
    \cs_set:Npn \return_true: {
      \cs_set:Npn \:nnn ####1 ####2 ####3 {
        \__bnvs_group_end:
        \tl_set:Nn #2 { ####1 }
        \tl_set:Nn #3 { ####2 }
        \seq_set_split:Nnn #4 . { ####3 }
        \seq_remove_all:Nn #4 { }
      }
      \exp_args:NVVx
      \:nnn #2 #3 {
        \seq_use:Nn #4 .
      }
\__bnvs_DEBUG:x { ...\string\__bnvs_resolve:NNNN\space TRUE/
  \string#1/\string#2:#2/\string#3:#3/\string#4:\seq_use:Nn #4./ }
      \prg_return_true:
    }
    \cs_set:Npn \branch:n ##1 {
      \seq_pop_left:NNTF \l_a_seq \l_b_tl {
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNNN\space POP~TRUE~##1/
  \string\l_b_tl:\l_b_tl/\string\l_a_seq:\seq_count:N\l_a_seq/
  \seq_use:Nn \l_a_seq ./ }
        \tl_put_right:Nn \l_a_tl { . }
        \tl_put_right:NV \l_a_tl \l_b_tl
      } {
        \cs_set_eq:NN \loop: \return_true:
      }
    }
    \cs_set:Npn \loop: {
      \__bnvs_call:TF {
        \exp_args:Nx
        \__bnvs_get:nNTF { \l_a_tl / L } \l_b_tl {
          \branch:n { 1 }
        } {
          \exp_args:Nx
          \__bnvs_get:nNTF { \l_a_tl / A } \l_b_tl {
            \__bnvs_extract_key:NNNTF #2 \l_b_tl \l_a_seq {
              \tl_set_eq:NN \l_a_tl \l_b_tl
              \tl_set_eq:NN #3 \l_b_tl
              \seq_set_eq:NN #4 \l_a_seq
            } {
              \branch:n { 2 }
            }
          } {
            \branch:n { 3 }
          }
        }
      } {
        \cs_set:Npn \loop: {
          \__bnvs_group_end:
\__bnvs_DEBUG:x { \string\__bnvs_resolve:NNNN\space FALSE/
```

```
  \string#1/\string#2:#2/\string#3:#3/\string#4:\seq_use:Nn #4./
  \g__bnvs_call_int :\int_use:N\g__bnvs_call_int/
}
          \prg_return_false:
        }
      }
      \loop:
    }
    \loop:
  } {
    \prg_return_true:
  }
}
```

5.5.7 Evaluation bricks

`\__bnvs_fp_round:nN` `\__bnvs_fp_round:nN {⟨expression⟩} ⟨tl variable⟩`
`\__bnvs_fp_round:N` `\__bnvs_fp_round:N ⟨tl variable⟩`

Shortcut for `\fp_eval:n{round(⟨expression⟩)}` appended to ⟨*tl variable*⟩. The second variant replaces the variable content with its rounded floating point evaluation.

```
\cs_new:Npn \__bnvs_fp_round:nN #1 #2 {
  \__bnvs_DEBUG:x { ROUND:\tl_to_str:n{#1}/\string#2=\tl_to_str:V #2}
  \tl_if_empty:nTF { #1 } {
    \__bnvs_DEBUG:x { ...ROUND:~EMPTY }
  } {
    \tl_put_right:Nx #2 {
      \fp_eval:n { round(#1) }
    }
    \__bnvs_DEBUG:x { ...ROUND:~\tl_to_str:V #2 => \string#2}
  }
}
\cs_generate_variant:Nn \__bnvs_fp_round:nN { VN, xN }
\cs_new:Npn \__bnvs_fp_round:N #1 {
  \tl_if_empty:VTF #1 {
    \__bnvs_DEBUG:x { ROUND:~EMPTY }
  } {
    \__bnvs_DEBUG:x { ROUND~IN:~\tl_to_str:V #1 }
    \tl_set:Nx #1 {
      \fp_eval:n { round(#1) }
    }
    \__bnvs_DEBUG:x { ROUND~OUT:~\tl_to_str:V #1 }
  }
}
```

`\__bnvs_raw_first:nNTF` `\__bnvs_raw_first:nNTF {⟨name⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}`
`\__bnvs_raw_first:(xN|VN)TF`

Append the first index of the ⟨*name*⟩ slide range to the ⟨*tl variable*⟩. Cache the result. Execute ⟨*true code*⟩ when there is a ⟨*first*⟩, ⟨*false code*⟩ otherwise.

```
\cs_set:Npn \__bnvs_return_true:nnN #1 #2 #3 {
  \tl_if_empty:NTF \l_ans_tl {
    \__bnvs_group_end:
```

```
\__bnvs_DEBUG:n { RETURN_FALSE/key=#1/type=#2/EMPTY }
    \__bnvs_gremove:n { #1//#2 }
    \prg_return_false:
  } {
    \__bnvs_fp_round:N \l_ans_tl
    \__bnvs_gput:nV { #1//#2 } \l_ans_tl
    \exp_args:NNNV
    \__bnvs_group_end:
    \tl_put_right:Nn #3 \l_ans_tl
\__bnvs_DEBUG:x { RETURN_TRUE/key=#1/type=#2/ans=\l_ans_tl/ }
    \prg_return_true:
  }
}
\cs_set:Npn \__bnvs_return_false:nn #1 #2 {
\__bnvs_DEBUG:n { RETURN_FALSE/key=#1/type=#2/ }
  \__bnvs_group_end:
  \__bnvs_gremove:n { #1//#2 }
  \prg_return_false:
}
\prg_new_conditional:Npnn \__bnvs_raw_first:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { RAW_FIRST/
    key=\tl_to_str:n{#1}/\string #2=/\tl_to_str:V #2/}
  \__bnvs_if_in:nTF { #1//A } {
\__bnvs_DEBUG:n { RAW_FIRST/#1/CACHED }
    \tl_put_right:Nx #2 { \__bnvs_item:n { #1//A } }
    \prg_return_true:
  } {
\__bnvs_DEBUG:n { RAW_FIRST/key=#1/NOT_CACHED }
    \__bnvs_group_begin:
    \tl_clear:N \l_ans_tl
    \__bnvs_get:nNTF { #1/A } \l_a_tl {
\__bnvs_DEBUG:x { RAW_FIRST/key=#1/A=\l_a_tl }
      \__bnvs_if_append:VNTF \l_a_tl \l_ans_tl {
        \__bnvs_return_true:nnN { #1 } A #2
      } {
        \__bnvs_return_false:nn { #1 } A
      }
    } {
\__bnvs_DEBUG:n { RAW_FIRST/key=#1/A/F }
      \__bnvs_get:nNTF { #1/L } \l_a_tl {
\__bnvs_DEBUG:n { RAW_FIRST/key=#1/L=\l_a_tl }
        \__bnvs_get:nNTF { #1/Z } \l_b_tl {
\__bnvs_DEBUG:n { RAW_FIRST/key=#1/Z=\l_b_tl }
          \__bnvs_if_append:xNTF {
            \l_b_tl - ( \l_a_tl ) + 1
          } \l_ans_tl {
            \__bnvs_return_true:nnN { #1 } A #2
          } {
            \__bnvs_return_false:nn { #1 } A
          }
        } {
\__bnvs_DEBUG:n { RAW_FIRST/key=#1/Z/F/ }
          \__bnvs_return_false:nn { #1 } A
        }
```

```
      } {
\__bnvs_DEBUG:n { RAW_FIRST/key=#1/L/F/ }
        \__bnvs_return_false:nn { #1 } A
      }
    }
  }
}
\prg_generate_conditional_variant:Nnn
    \__bnvs_raw_first:nN { VN, xN } { T, F, TF }
```

__bnvs_if_first:nN*TF* `\__bnvs_if_first:nNTF {⟨name⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}`

Append the first index of the ⟨*name*⟩ slide range to the ⟨*tl variable*⟩. If no first index was explicitely given, use the counter when available and 1 hen not. Cache the result. Execute ⟨*true code*⟩ when there is a ⟨*first*⟩, ⟨*false code*⟩ otherwise.

```
\prg_new_conditional:Npnn \__bnvs_if_first:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { IF_FIRST/\tl_to_str:n{#1}/\string #2=\tl_to_str:V #2}
  \__bnvs_raw_first:nNTF { #1 } #2 {
    \prg_return_true:
  } {
    \__bnvs_get:nNTF { #1/C } \l_a_tl {
\__bnvs_DEBUG:n { IF_FIRST/#1/C/T/\l_a_tl }
      \bool_set_true:N \l_no_counter_bool
      \__bnvs_if_append:xNTF \l_a_tl \l_ans_tl {
        \__bnvs_return_true:nnN { #1 } A #2
      } {
        \__bnvs_return_false:nn { #1 } A
      }
    } {
      \regex_match:NnTF \c__bnvs_A_key_Z_regex { #1 } {
        \__bnvs_gput:nn { #1/A } { 1 }
        \tl_set:Nn #2 { 1 }
\__bnvs_DEBUG:x{IF_FIRST_MATCH:
  key=\tl_to_str:n{#1}/\string #2=\tl_to_str:V #2 /}
        \__bnvs_return_true:nnN { #1 } A #2
      } {
\__bnvs_DEBUG:x{IF_FIRST_NO_MATCH:
  key=\tl_to_str:n{#1}/\string #2=\tl_to_str:V #2 /}
        \__bnvs_return_false:nn { #1 } A
      }
    }
  }
}
```

__bnvs_first:nN
__bnvs_first:VN

`\__bnvs_first:nN {⟨name⟩} ⟨tl variable⟩`

Append the start of the ⟨*name*⟩ slide range to the ⟨*tl variable*⟩. Cache the result.

```
\cs_new:Npn \__bnvs_first:nN #1 #2 {
  \__bnvs_if_first:nNF { #1 } #2 {
    \msg_error:nnn { beanoves } { :n } { Range~with~no~first:~#1 }
  }
}
\cs_generate_variant:Nn \__bnvs_first:nN { VN }
```

__bnvs_raw_length:nN*TF* `\__bnvs_raw_length:nNTF {⟨name⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}`

Append the length of the ⟨*name*⟩ slide range to ⟨*tl variable*⟩ Execute ⟨*true code*⟩ when there is a ⟨*length*⟩, ⟨*false code*⟩ otherwise.

```
\prg_new_conditional:Npnn \__bnvs_raw_length:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { \string\__bnvs_raw_length:nN/#1/\string#2/ }
  \__bnvs_if_in:nTF { #1//L } {
    \tl_put_right:Nx #2 { \__bnvs_item:n { #1//L } }
\__bnvs_DEBUG:x { RAW_LENGTH/CACHED/key:#1/\__bnvs_item:n { #1//L } }
    \prg_return_true:
  } {
\__bnvs_DEBUG:x { RAW_LENGTH/NOT_CACHED/key:#1/ }
    \__bnvs_gput:nn { #1//L } { 0 }
    \__bnvs_group_begin:
    \tl_clear:N \l_ans_tl
    \__bnvs_if_in:nTF { #1/L } {
      \__bnvs_if_append:xNTF {
        \__bnvs_item:n { #1/L }
      } \l_ans_tl {
        \__bnvs_return_true:nnN { #1 } L #2
      } {
        \__bnvs_return_false:nn { #1 } L
      }
    } {
      \__bnvs_get:nNTF { #1/A } \l_a_tl {
        \__bnvs_get:nNTF { #1/Z } \l_b_tl {
          \__bnvs_if_append:xNTF {
            \l_b_tl - (\l_a_tl) + 1
          } \l_ans_tl {
            \__bnvs_return_true:nnN { #1 } L #2
          } {
            \__bnvs_return_false:nn { #1 } L
          }
        } {
          \__bnvs_return_false:nn { #1 } L
        }
      } {
        \__bnvs_return_false:nn { #1 } L
      }
    }
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_raw_length:nN { VN } { T, F, TF }
```

__bnvs_raw_last:nN*TF* `\__bnvs_raw_last:nNTF {⟨name⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}`

Put the last index of the fully qualified ⟨*name*⟩ range to the right of the ⟨*tl variable*⟩, when possible. Execute ⟨*true code*⟩ when a last index was given, ⟨*false code*⟩ otherwise.

```
\prg_new_conditional:Npnn \__bnvs_raw_last:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:n { RAW_LAST/#1 }
  \__bnvs_if_in:nTF { #1//Z } {
    \tl_put_right:Nx #2 { \__bnvs_item:n { #1//Z } }
```

```
    \prg_return_true:
  }  {
    \__bnvs_gput:nn { #1//Z } { 0 }
    \__bnvs_group_begin:
    \tl_clear:N \l_ans_tl
    \__bnvs_if_in:nTF { #1/Z } {
\__bnvs_DEBUG:x { NORMAL_RAW_LAST:~\__bnvs_item:n { #1/Z } }
      \__bnvs_if_append:xNTF {
        \__bnvs_item:n { #1/Z }
      } \l_ans_tl {
        \__bnvs_return_true:nnN { #1 } Z #2
      } {
        \__bnvs_return_false:nn { #1 } Z
      }
    } {
      \__bnvs_get:nNTF { #1/A } \l_a_tl {
        \__bnvs_get:nNTF { #1/L } \l_b_tl {
          \__bnvs_if_append:xNTF {
            \l_a_tl + (\l_b_tl) - 1
          } \l_ans_tl {
            \__bnvs_return_true:nnN { #1 } Z #2
          } {
            \__bnvs_return_false:nn { #1 } Z
          }
        } {
          \__bnvs_return_false:nn { #1 } Z
        }
      } {
        \__bnvs_return_false:nn { #1 } Z
      }
    }
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_raw_last:nN { VN } { T, F, TF }
```

__bnvs_last:nN
__bnvs_last:VN

`\__bnvs_last:nN {⟨name⟩} ⟨tl variable⟩`

Append the last index of the fully qualified ⟨*name*⟩ slide range to ⟨*tl variable*⟩

```
\cs_new:Npn \__bnvs_last:nN #1 #2 {
  \__bnvs_raw_last:nNF { #1 } #2 {
    \msg_error:nnn { beanoves } { :n } { Range~with~no~last:~#1 }
  }
}
\cs_generate_variant:Nn \__bnvs_last:nN { VN }
```

__bnvs_if_next:nN*TF*

`\__bnvs_if_next:nNTF {⟨name⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}`

Append the index after the ⟨*name*⟩ slide range to the ⟨*tl variable*⟩. Execute ⟨*true code*⟩ when there is a ⟨*next*⟩ index, ⟨*false code*⟩ otherwise.

```
\prg_new_conditional:Npnn \__bnvs_if_next:nN #1 #2 { T, F, TF } {
  \__bnvs_if_in:nTF { #1//N } {
    \tl_put_right:Nx #2 { \__bnvs_item:n { #1//N } }
```

```
    \prg_return_true:
  } {
    \__bnvs_group_begin:
    \cs_set:Npn \__bnvs_return_true: {
      \tl_if_empty:NTF \l_ans_tl {
        \__bnvs_group_end:
        \prg_return_false:
      } {
        \__bnvs_fp_round:N \l_ans_tl
        \__bnvs_gput:nV { #1//N } \l_ans_tl
        \exp_args:NNNV
        \__bnvs_group_end:
        \tl_put_right:Nn #2 \l_ans_tl
        \prg_return_true:
      }
    }
    \cs_set:Npn \return_false: {
      \__bnvs_group_end:
      \prg_return_false:
    }
    \tl_clear:N \l_a_tl
    \__bnvs_raw_last:nNTF { #1 } \l_a_tl {
      \__bnvs_if_append:xNTF {
        \l_a_tl + 1
      } \l_ans_tl {
        \__bnvs_return_true:
      } {
        \return_false:
      }
    } {
      \return_false:
    }
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_if_next:nN { VN } { T, F, TF }
```

__bnvs_next:nN
__bnvs_next:VN

`\__bnvs_next:nN {⟨name⟩} ⟨tl variable⟩`

Append the index after the ⟨*name*⟩ slide range to the ⟨*tl variable*⟩.

```
\cs_new:Npn \__bnvs_next:nN #1 #2 {
  \__bnvs_if_next:nNF { #1 } #2 {
    \msg_error:nnn { beanoves } { :n } { Range~with~no~next:~#1 }
  }
}
\cs_generate_variant:Nn \__bnvs_next:nN { VN }
```

__bnvs_if_index:nnNTF
__bnvs_if_index:VVNTF
__bnvs_if_index:nnnNTF

__bnvs_if_index:nnNTF {⟨*name*⟩} {⟨*integer*⟩} ⟨*tl variable*⟩ {⟨*true code*⟩} {⟨*false code*⟩}

Append the index associated to the {⟨*name*⟩} and {⟨*integer*⟩} slide range to the right of ⟨*tl variable*⟩. When ⟨*integer shift*⟩ is 1, this is the first index, when ⟨*integer shift*⟩ is 2, this is the second index, and so on. When ⟨*integer shift*⟩ is 0, this is the index, before the first one, and so on. If the computation is possible, ⟨*true code*⟩ is executed, otherwise ⟨*false code*⟩ is executed. The computation may fail when too many recursion calls are made.

```
\prg_new_conditional:Npnn \__bnvs_if_index:nnN #1 #2 #3 { T, F, TF } {
\__bnvs_DEBUG:x { IF_INDEX:key=#1/index=#2/\string#3/ }
  \__bnvs_group_begin:
  \tl_clear:N \l_ans_tl
  \__bnvs_raw_first:nNTF { #1 } \l_ans_tl {
    \tl_put_right:Nn \l_ans_tl { + (#2) - 1}
    \exp_args:NNV
    \__bnvs_group_end:
    \__bnvs_fp_round:nN \l_ans_tl #3
\__bnvs_DEBUG:x { IF_INDEX_TRUE:key=#1/index=#2/
  \string#3=\tl_to_str:N #3 }
    \prg_return_true:
  } {
\__bnvs_DEBUG:x { IF_INDEX_FALSE:key=#1/index=#2/ }
    \prg_return_false:
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_if_index:nnN { VVN } { T, F, TF }
```

__bnvs_if_range:nNTF

__bnvs_if_range:nNTF {⟨*name*⟩} ⟨*tl variable*⟩ {⟨*true code*⟩} {⟨*false code*⟩}

Append the range of the ⟨*name*⟩ slide range to the ⟨*tl variable*⟩. Execute ⟨*true code*⟩ when there is a ⟨*range*⟩, ⟨*false code*⟩ otherwise.

```
\prg_new_conditional:Npnn \__bnvs_if_range:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x{ RANGE:key=#1/\string#2/}
  \bool_if:NTF \l__bnvs_no_range_bool {
    \prg_return_false:
  } {
    \__bnvs_if_in:nTF { #1/ } {
      \tl_put_right:Nn { 0-0 }
    } {
      \__bnvs_group_begin:
      \tl_clear:N \l_a_tl
      \tl_clear:N \l_b_tl
      \tl_clear:N \l_ans_tl
      \__bnvs_raw_first:nNTF { #1 } \l_a_tl {
        \__bnvs_raw_last:nNTF { #1 } \l_b_tl {
          \exp_args:NNNx
          \__bnvs_group_end:
          \tl_put_right:Nn #2 { \l_a_tl - \l_b_tl }
\__bnvs_DEBUG:x{ RANGE_TRUE_A_Z:key=#1/\string#2=#2/}
          \prg_return_true:
        } {
```

```
          \exp_args:NNNx
          \__bnvs_group_end:
          \tl_put_right:Nn #2 { \l_a_tl - }
\__bnvs_DEBUG:x{ RANGE_TRUE_A:key=#1/\string#2=#2/}
          \prg_return_true:
        }
      } {
        \__bnvs_raw_last:nNTF { #1 } \l_b_tl {
\__bnvs_DEBUG:x{ RANGE_TRUE_Z:key=#1/\string#2=#2/}
          \exp_args:NNNx
          \__bnvs_group_end:
          \tl_put_right:Nn #2 { - \l_b_tl }
          \prg_return_true:
        } {
\__bnvs_DEBUG:x{ RANGE_FALSE:key=#1/}
          \__bnvs_group_end:
          \prg_return_false:
        }
      }
    }
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_if_range:nN { VN } { T, F, TF }
```

`\__bnvs_range:nN` `\__bnvs_range:VN`

`\__bnvs_range:nN {`⟨*name*⟩`}` ⟨*tl variable*⟩

Append the range of the ⟨*name*⟩ slide range to the ⟨*tl variable*⟩.

```
\cs_new:Npn \__bnvs_range:nN #1 #2 {
  \__bnvs_if_range:nNF { #1 } #2 {
    \msg_error:nnn { beanoves } { :n } { No~range~available:~#1 }
  }
}
\cs_generate_variant:Nn \__bnvs_range:nN { VN }
```

`\__bnvs_if_free_counter:nN`*TF* `\__bnvs_if_free_counter:VN`*TF*

`\__bnvs_if_free_counter:nNTF {`⟨*name*⟩`}` ⟨*tl variable*⟩ `{`⟨*true code*⟩`}` `{`⟨*false code*⟩`}`

Set the ⟨*tl variable*⟩ to the value of the counter associated to the `{`⟨*name*⟩`}` slide range.

```
\prg_new_conditional:Npnn \__bnvs_if_free_counter:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { IF_FREE:  key=\tl_to_str:n{#1}/
    value=\__bnvs_item:n {#1/C}/cs=\string #2/ }
  \__bnvs_group_begin:
  \tl_clear:N \l_ans_tl
  \__bnvs_get:nNF { #1/C } \l_ans_tl {
    \__bnvs_raw_first:nNF { #1 } \l_ans_tl {
      \__bnvs_raw_last:nNF { #1 } \l_ans_tl { }
    }
  }
\__bnvs_DEBUG:x { IF_FREE_2:\string \l_ans_tl=\tl_to_str:V \l_ans_tl/}
  \tl_if_empty:NTF \l_ans_tl {
    \__bnvs_group_end:
```

```
    \regex_match:NnTF \c__bnvs_A_key_Z_regex { #1 } {
      \__bnvs_gput:nn { #1/C } { 1 }
      \tl_set:Nn #2 { 1 }
\__bnvs_DEBUG:x { IF_FREE_MATCH_TRUE:
  key=\tl_to_str:n{#1}\string #2=\tl_to_str:V #2 / }
      \prg_return_true:
    } {
\__bnvs_DEBUG:x { IF_FREE_NO_MATCH_FALSE:
  key=\tl_to_str:n{#1}\string #2=\tl_to_str:V #2/ }
      \prg_return_false:
    }
  } {
    \__bnvs_gput:nV { #1/C } \l_ans_tl
    \exp_args:NNNV
    \__bnvs_group_end:
    \tl_set:Nn #2 \l_ans_tl
\__bnvs_DEBUG:x { IF_FREE_TRUE(2):  /
  key=\tl_to_str:n{#1}/\string #2=\tl_to_str:V #2}
    \prg_return_true:
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_if_free_counter:nN { VN } { T, F, TF }
```

__bnvs_if_counter:nN*TF*
__bnvs_if_counter:VN*TF*

`\__bnvs_if_counter:nNTF {`⟨*name*⟩`} ⟨`*tl variable*⟩ {⟨*true code*⟩} {⟨*false code*⟩}`

Append the value of the counter associated to the {⟨*name*⟩} slide range to the right of ⟨*tl variable*⟩. The value always lays in between the range, whenever possible.

```
\prg_new_conditional:Npnn \__bnvs_if_counter:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { IF_COUNTER:key=
    \tl_to_str:n{#1}/\string #2=\tl_to_str:V #2 }
  \__bnvs_group_begin:
  \__bnvs_if_free_counter:nNTF { #1 } \l_ans_tl {
```

If there is a ⟨*first*⟩, use it to bound the result from below.

```
    \tl_clear:N \l_a_tl
    \__bnvs_raw_first:nNT { #1 } \l_a_tl {
      \fp_compare:nNnT { \l_ans_tl } < { \l_a_tl } {
        \tl_set:NV \l_ans_tl \l_a_tl
      }
    }
```

If there is a ⟨*last*⟩, use it to bound the result from above.

```
    \tl_clear:N \l_a_tl
    \__bnvs_raw_last:nNT { #1 } \l_a_tl {
      \fp_compare:nNnT { \l_ans_tl } > { \l_a_tl } {
        \tl_set:NV \l_ans_tl \l_a_tl
      }
    }
    \exp_args:NNV
    \__bnvs_group_end:
    \__bnvs_fp_round:nN \l_ans_tl #2
\__bnvs_DEBUG:x {IF_COUNTER_TRUE:key=\tl_to_str:n{#1}/
  \string #2=\tl_to_str:V #2 }
```

```
    \prg_return_true:
  } {
\__bnvs_DEBUG:x {IF_COUNTER_FALSE:key=\tl_to_str:n{#1}/
  \string #2=\tl_to_str:V #2 }
    \prg_return_false:
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_if_counter:nN { VN } { T, F, TF }
```

__bnvs_if_incr:nn*TF*
__bnvs_if_incr:nnN*TF*
__bnvs_if_incr:(VnN|VVN)*TF*

__bnvs_if_incr:nnTF {⟨*name*⟩} {⟨*offset*⟩} {⟨*true code*⟩} {⟨*false code*⟩}
__bnvs_if_incr:nnNTF {⟨*name*⟩} {⟨*offset*⟩} ⟨*tl variable*⟩ {⟨*true code*⟩} {⟨*false code*⟩}

Increment the free counter position accordingly. When requested, put the result in the ⟨*tl variable*⟩. In the second version, the result will lay within the declared range.

```
\prg_new_conditional:Npnn \__bnvs_if_incr:nn #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { IF_INCR:\tl_to_str:n{#1}/\tl_to_str:n{#2} }
  \__bnvs_group_begin:
  \tl_clear:N \l_a_tl
  \__bnvs_if_free_counter:nNTF { #1 } \l_a_tl {
    \tl_clear:N \l_b_tl
    \__bnvs_if_append:xNTF { \l_a_tl + (#2) } \l_b_tl {
      \__bnvs_fp_round:N \l_b_tl
      \__bnvs_gput:nV { #1/C } \l_b_tl
      \__bnvs_group_end:
\__bnvs_DEBUG:x { IF_INCR_TRUE:#1/#2 }
      \prg_return_true:
    } {
      \__bnvs_group_end:
\__bnvs_DEBUG:x { IF_INCR_FALSE(1):#1/#2 }
      \prg_return_false:
    }
  } {
    \__bnvs_group_end:
\__bnvs_DEBUG:x { IF_INCR_FALSE(2):#1/#2 }
    \prg_return_false:
  }
}
\prg_new_conditional:Npnn \__bnvs_if_incr:nnN #1 #2 #3 { T, F, TF } {
  \__bnvs_if_incr:nnTF { #1 } { #2 } {
    \__bnvs_if_counter:nNTF { #1 } #3 {
      \prg_return_true:
    } {
      \prg_return_false:
    }
  } {
    \prg_return_false:
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_if_incr:nnN { VnN, VVN } { T, F, TF }
```

5.5.8 Evaluation

__bnvs_if_append:nNTF
__bnvs_if_append:(VN|xN)TF

__bnvs_if_append:nNTF {⟨*integer expression*⟩} ⟨*tl variable*⟩ {⟨*true code*⟩} {⟨*false code*⟩}

Evaluates the ⟨*integer expression*⟩, replacing all the named specifications by their static counterpart then put the result to the right of the ⟨*tl variable*⟩. Executed within a group. Heavily used by __bnvs_eval_query:nN, where ⟨*integer expression*⟩ was initially enclosed in '?(...)'. Local variables:

\l_ans_tl To feed ⟨*tl variable*⟩ with.

(*End definition for* \l_ans_tl. *This variable is documented on page* ??.)

\l_split_seq The sequence of catched query groups and non queries.

(*End definition for* \l_split_seq. *This variable is documented on page* ??.)

\l_split_int Is the index of the non queries, before all the catched groups.

(*End definition for* \l_split_int. *This variable is documented on page* ??.)

```
\int_if_exist:NF \l_split_int {
  \int_new:N  \l_split_int
}
```

\l_name_tl Storage for \l_split_seq items that represent names.

(*End definition for* \l_name_tl. *This variable is documented on page* ??.)

\l_path_tl Storage for \l_split_seq items that represent integer paths.

(*End definition for* \l_path_tl. *This variable is documented on page* ??.)

Catch circular definitions.

```
\prg_new_conditional:Npnn \__bnvs_if_append:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { \string\__bnvs_if_append:nNTF/
    \tl_to_str:n { #1 } / \string #2/
}
  \__bnvs_call:TF {
\__bnvs_DEBUG:x { IF_APPEND...}
    \__bnvs_group_begin:
```

Local variables:

```
    \int_zero:N  \l_split_int
    \seq_clear:N \l_split_seq
    \tl_clear:N  \l_id_tl
    \tl_clear:N  \l_name_tl
    \tl_clear:N  \l_path_tl
    \tl_clear:N  \l_group_tl
    \tl_clear:N  \l_ans_tl
    \tl_clear:N  \l_a_tl
```

Implementation:

```
    \regex_split:NnN \c__bnvs_split_regex { #1 } \l_split_seq
\__bnvs_DEBUG:x { IF_APPEND_SPLIT_SEQ: /
  \#=\seq_count:N \l_split_seq /
```

```
  \seq_use:Nn \l_split_seq / /
}
    \int_set:Nn \l_split_int { 1 }
    \tl_set:Nx \l_ans_tl {
      \seq_item:Nn \l_split_seq { \l_split_int }
    }
\__bnvs_DEBUG:x { ANS: \l_ans_tl }
```

\switch:nTF `\switch:nTF {⟨capture group number⟩} {⟨black code⟩} {⟨white code⟩}`

Helper function to locally set the `\l_group_tl` variable to the captured group ⟨*capture group number*⟩ and branch.

```
    \cs_set:Npn \switch:nNTF ##1 ##2 ##3 ##4 {
      \tl_set:Nx ##2 {
        \seq_item:Nn \l_split_seq { \l_split_int + ##1 }
      }
\__bnvs_DEBUG:x { IF_APPEND_SWITCH/##1/
  \int_eval:n { \l_split_int + ##1 } /
  \string##2=\tl_to_str:N##2/
}
      \tl_if_empty:NTF ##2 {
\__bnvs_DEBUG:x { IF_APPEND_SWITCH_WHITE/##1/
  \int_eval:n { \l_split_int + ##1 }
}
        ##4 } {
\__bnvs_DEBUG:x { IF_APPEND_SWITCH_BLACK/##1/
  \int_eval:n { \l_split_int + ##1 }
}
        ##3
      }
    }
```

`\prg_return_true:` and `\prg_return_false:` are wrapped locally to close the group and return the proper value.

```
    \cs_set:Npn \return_true: {
      \fp_round:
      \exp_args:NNNV
      \__bnvs_group_end:
      \tl_put_right:Nn #2 \l_ans_tl
\__bnvs_DEBUG:x { IF_APPEND_TRUE:\tl_to_str:n { #1 } /
  \string #2=\tl_to_str:V #2 /}
\log_g_prop:
      \prg_return_true:
    }
    \cs_set:Npn \fp_round: {
      \__bnvs_fp_round:N \l_ans_tl
    }
    \cs_set:Npn \return_false: {
      \__bnvs_group_end:
\__bnvs_DEBUG:x { IF_APPEND_FALSE:\tl_to_str:n { #1 } /
  \string #2=\tl_to_str:V #2 / }
      \prg_return_false:
    }
    \cs_set:Npn \:NnnT ##1 ##2 ##3 ##4 {
```

```
      \switch:nNTF { ##2 } \l_id_tl { } {
        \tl_set_eq:NN \l_id_tl \l__bnvs_id_tl
        \tl_put_left:NV \l_name_tl \l_id_tl
      }
      \switch:nNTF { ##3 } \l_path_tl {
        \seq_set_split:NnV \l_path_seq { . } \l_path_tl
        \seq_remove_all:Nn \l_path_seq { }
\__bnvs_DEBUG:x { PATH_SEQ:\l_path_tl==\seq_use:Nn\l_path_seq .}
      } {
        \seq_clear:N \l_path_seq
      }
\__bnvs_DEBUG:x { PATH_SEQ:\l_path_tl==\seq_use:Nn\l_path_seq .}
\__bnvs_DEBUG:x { \string ##1 }
      ##1 \l_id_tl \l_name_tl \l_path_seq {
        \cs_set:Npn \: {
          ##4
\__bnvs_DEBUG:x {AFTER_ANS~=~\l_ans_tl}
        }
      } {
        \cs_set:Npn \: { \cs_set_eq:NN \loop: \return_false: }
      }
      \:
\__bnvs_DEBUG:x {AFTER_AFTER_ANS~=~\l_ans_tl}
    }
    \cs_set:Npn \:T ##1 {
      \seq_if_empty:NTF \l_path_seq { ##1 } {
        \cs_set_eq:NN \loop: \return_false:
      }
    }
```

Main loop.

```
    \cs_set:Npn \loop: {
\__bnvs_DEBUG:x { IF_APPEND_LOOP:\int_use:N\l_split_int /
  \seq_count:N \l_split_seq / }
      \int_compare:nNnTF {
        \l_split_int } < { \seq_count:N \l_split_seq
      } {
        \switch:nNTF 1 \l_name_tl {
```

- Case ++⟨*name*⟩⟨*integer path*⟩.n.

```
          \:NnnT \__bnvs_resolve_n:NNNTF 2 3 {
            \__bnvs_if_incr:VnNF \l_name_tl 1 \l_ans_tl {
              \cs_set_eq:NN \loop: \return_false:
            }
          }
        } {
          \switch:nNTF 4 \l_name_tl {
```

- Cases ⟨*name*⟩⟨*integer path*⟩....

```
            \switch:nNTF 7 \l_a_tl {
              \:NnnT \__bnvs_resolve:NNNTF 5 6 {
                \:T {
                  \__bnvs_raw_length:VNF \l_name_tl \l_ans_tl {
```

```
                \cs_set_eq:NN \loop: \return_false:
              }
            }
          }
```

- Case `...length`.

```
        } {
          \switch:nNTF 8 \l_a_tl {
```

- Case `...last`.

```
            \:NnnT \__bnvs_resolve:NNNTF 5 6 {
              \:T {
                \__bnvs_raw_last:VNF \l_name_tl \l_ans_tl {
                  \cs_set_eq:NN \loop: \return_false:
                }
              }
            }


          } {
            \switch:nNTF 9 \l_a_tl {
```

- Case `...next`.

```
              \:NnnT \__bnvs_resolve:NNNTF 5 6 {
                \:T {
                  \__bnvs_if_next:VNF \l_name_tl \l_ans_tl {
                    \cs_set_eq:NN \loop: \return_false:
                  }
                }
              }
            } {
              \switch:nNTF { 10 } \l_a_tl {
```

- Case `...range`.

```
                \:NnnT \__bnvs_resolve:NNNTF 5 6 {
                  \:T {
                    \__bnvs_if_range:VNTF \l_name_tl \l_ans_tl {
                      \cs_set_eq:NN \fp_round: \prg_do_nothing:
                    } {
                      \cs_set_eq:NN \loop: \return_false:
                    }
                  }
                }
              } {
                \switch:nNTF { 11 } \l_a_tl {
```

- Case `...n`.

```
                  \switch:nNTF { 12 } \l_a_tl {
```

- Case `...+=`⟨*integer*⟩.

```
                        \:NnnT \__bnvs_resolve_n:NNNTF 5 6 {
                          \:T {
\__bnvs_DEBUG:x {NAME=\l_name_tl}
\__bnvs_DEBUG:x {INCR=\l_a_tl}
\__bnvs_if_incr:VVNF \l_name_tl \l_a_tl \l_ans_tl {
  \cs_set_eq:NN \loop: \return_false:
}
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
                          }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
                        }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
                      } {
                        \:NnnT \__bnvs_resolve_n:NNNTF 5 6 {
                          \seq_if_empty:NTF \l_path_seq {
\__bnvs_if_counter:VNF \l_name_tl \l_ans_tl {
  \cs_set_eq:NN \loop: \return_false:
}
                          } {
                            \seq_pop_left:NN \l_path_seq \l_a_tl
                            \seq_if_empty:NTF \l_path_seq {
\__bnvs_if_incr:VVNF \l_name_tl \l_a_tl \l_ans_tl {
\__bnvs_DEBUG:x { INCR~FALSE }
  \cs_set_eq:NN \loop: \return_false:
}
\__bnvs_DEBUG:x { INCR~ANS=\l_ans_tl }
                            } {
\msg_error:nnx { beanoves } { :n } { Too~many~.<integer>~components:~#1 }
\cs_set_eq:NN \loop: \return_false:
                            }
                          }
                        }
                      }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}


                    } {
                      \:NnnT \__bnvs_resolve_n:NNNTF 5 6 {
                        \seq_if_empty:NTF \l_path_seq {
\__bnvs_if_counter:VNF \l_name_tl \l_ans_tl {
  \cs_set_eq:NN \loop: \return_false:
}
                        } {
                          \seq_pop_left:NN \l_path_seq \l_a_tl
                          \seq_if_empty:NTF \l_path_seq {
\__bnvs_if_index:VVNF \l_name_tl \l_a_tl \l_ans_tl {
  \cs_set_eq:NN \loop: \return_false:
}
                          } {
\msg_error:nnx { beanoves } { :n } { Too~many~.<integer>~components:~#1 }
\cs_set_eq:NN \loop: \return_false:
                          }
                        }
```

```
                      }
                    }
                  }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
                }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
              }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
            }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
          } {
```

No name.

```
          }
\__bnvs_DEBUG:x {ANS~=~\l_ans_tl}
        }
\__bnvs_DEBUG:x {ITERATE~ANS=\l_ans_tl }
        \int_add:Nn \l_split_int { 13 }
        \tl_put_right:Nx \l_ans_tl {
          \seq_item:Nn \l_split_seq { \l_split_int }
        }
\__bnvs_DEBUG:x {ITERATE~ANS=\l_ans_tl }
        \loop:
      } {
\__bnvs_DEBUG:x {END_OF_LOOP~ANS=\l_ans_tl }
        \return_true:
      }
    }
    \loop:
  } {
    \msg_error:nnx { beanoves } { :n } { Too~many~calls:~ #1 }
    \prg_return_false:
  }
}
\prg_generate_conditional_variant:Nnn
  \__bnvs_if_append:nN { VN, xN } { T, F, TF }
```

__bnvs_if_eval_query:nN*TF* `\__bnvs_if_eval_query:nNTF {⟨overlay query⟩} ⟨tl variable⟩ {⟨true code⟩} {⟨false code⟩}`

Evaluates the single ⟨*overlay query*⟩, which is expected to contain no comma. Extract a range specification from the argument, replaces all the *named overlay specifications* by their static counterparts, make the computation then append the result to the right of the ⟨*seq variable*⟩. Ranges are supported with the colon syntax. This is executed within a local group. Below are local variables and constants.

`\l_a_tl` Storage for the first index of a range.

(*End definition for* `\l_a_tl`. *This variable is documented on page* ??.)

`\l_b_tl` Storage for the last index of a range, or its length.

(*End definition for* `\l_b_tl`. *This variable is documented on page* ??.)

`\c__bnvs_A_cln_Z_regex` Used to parse slide range overlay specifications. Next are the capture groups.

(*End definition for* `\c__bnvs_A_cln_Z_regex`.)

```
\regex_const:Nn \c__bnvs_A_cln_Z_regex {
  \A \s* (?:
```

- 2: ⟨*first*⟩

```
      ( [^:]* ) \s* :
```

- 3: second optional colon

```
      (:)? \s*
```

- 4: ⟨*length*⟩

```
      ( [^:]* )
```

- 5: standalone ⟨*first*⟩

```
    | ( [^:]+ )
  ) \s* \Z
}

\prg_new_conditional:Npnn \__bnvs_if_eval_query:nN #1 #2 { T, F, TF } {
\__bnvs_DEBUG:x { EVAL_QUERY:#1/
    \tl_to_str:n{#1}/\string#2=\tl_to_str:N #2}
  \__bnvs_call_reset:
  \regex_extract_once:NnNTF \c__bnvs_A_cln_Z_regex {
    #1
  } \l_match_seq {
\__bnvs_DEBUG:x { EVAL_QUERY:#1/
  \string\l_match_seq/\seq_use:Nn \l_match_seq //}
    \bool_set_false:N \l__bnvs_no_counter_bool
    \bool_set_false:N \l__bnvs_no_range_bool
```

`\switch:nNTF` `\switch:nNTF {⟨capture group number⟩} ⟨tl variable⟩ {⟨black code⟩} {⟨white code⟩}`

Helper function to locally set the ⟨*tl variable*⟩ to the captured group ⟨*capture group number*⟩ and branch depending on the emptyness of this variable.

```
    \cs_set:Npn \switch:nNTF ##1 ##2 ##3 ##4 {
\__bnvs_DEBUG:x { EQ_SWITCH:##1/ }
      \tl_set:Nx ##2 {
        \seq_item:Nn \l_match_seq { ##1 }
      }
\__bnvs_DEBUG:x { \string ##2/ \tl_to_str:N ##2/}
      \tl_if_empty:NTF ##2 { ##4 } { ##3 }
    }
    \switch:nNTF 5 \l_a_tl {
```

Single expression

```
      \bool_set_false:N \l__bnvs_no_range_bool
      \__bnvs_if_append:VNTF \l_a_tl #2 {
        \prg_return_true:
      } {
        \prg_return_false:
      }
    } {
      \switch:nNTF 2 \l_a_tl {
        \switch:nNTF 4 \l_b_tl {
          \switch:nNTF 3 \l_c_tl {
```

⟨*first*⟩::⟨*last*⟩ range

```
            \__bnvs_if_append:VNTF \l_a_tl #2 {
              \tl_put_right:Nn #2 { - }
              \__bnvs_if_append:VNTF \l_b_tl #2 {
                \prg_return_true:
              } {
                \prg_return_false:
              }
            } {
              \prg_return_false:
            }
          } {
```

⟨*first*⟩:⟨*length*⟩ range

```
            \__bnvs_if_append:VNTF \l_a_tl #2 {
              \tl_put_right:Nx #2 { - }
              \tl_put_right:Nx \l_a_tl { + ( \l_b_tl ) - 1}
              \__bnvs_if_append:VNTF \l_a_tl #2 {
                \prg_return_true:
              } {
                \prg_return_false:
              }
            } {
              \prg_return_false:
            }
          }
        } {
```

⟨*first*⟩: and ⟨*first*⟩:: range

```
          \__bnvs_if_append:VNTF \l_a_tl #2 {
            \tl_put_right:Nn #2 { - }
            \prg_return_true:
          } {
```

```
            \prg_return_false:
          }
        }
      } {
        \switch:nNTF 4 \l_b_tl {
          \switch:nNTF 3 \l_c_tl {
```

::⟨*last*⟩ range

```
            \tl_put_right:Nn #2 { - }
            \__bnvs_if_append:VNTF \l_a_tl #2 {
              \prg_return_true:
            } {
              \prg_return_false:
            }
          } {
\msg_error:nnx { beanoves } { :n } { Syntax~error(Missing~first):~#1 }
          }
        } {
```

: or :: range

```
          \seq_put_right:Nn #2 { - }
        }
      }
    }
  } {
```

Error

```
    \msg_error:nnn { beanoves } { :n } { Syntax~error:~#1 }
  }
}
```

`\__bnvs_eval:nN` `\__bnvs_eval:nN {`⟨*overlay query list*⟩`}` ⟨*tl variable*⟩

This is called by the *named overlay specifications* scanner. Evaluates the comma separated list of ⟨*overlay query*⟩'s, replacing all the named overlay specifications and integer expressions by their static counterparts by calling `\__bnvs_eval_query:nN`, then append the result to the right of the ⟨*tl variable*⟩. This is executed within a local group. Below are local variables and constants used throughout the body of this function.

`\l_query_seq` Storage for a sequence of ⟨*query*⟩'s obtained by splitting a comma separated list.

(*End definition for* `\l_query_seq`. *This variable is documented on page* ??.)

`\l_ans_seq` Storage of the evaluated result.

(*End definition for* `\l_ans_seq`. *This variable is documented on page* ??.)

`\c__bnvs_comma_regex` Used to parse slide range overlay specifications.

```
\regex_const:Nn \c__bnvs_comma_regex { \s* , \s* }
```

(*End definition for* `\c__bnvs_comma_regex`.)

No other variable is used.

```
\cs_new:Npn \__bnvs_eval:nN #1 #2 {
\__bnvs_DEBUG:x {\string\__bnvs_eval:nN:\tl_to_str:n{#1}/\string#2=\tl_to_str:V #2}
  \__bnvs_group_begin:
```

Local variables declaration

```
\__bnvs_DEBUG:x { WILL_CLEAN }
  \seq_if_exist:NTF \l_query_seq {
    \seq_clear:N \l_query_seq
  } {
    \seq_new:N \l_query_seq
  }
  \seq_if_exist:NTF \l_ans_seq {
    \seq_clear:N \l_ans_seq
  } {
    \seq_new:N \l_ans_seq
  }
```

In this main evaluation step, we evaluate the integer expression and put the result in a variable which content will be copied after the group is closed. We authorize comma separated expressions and ⟨*first*⟩`::`⟨*last*⟩ range expressions as well. We first split the expression around commas, into `\l_query_seq`.

```
  \regex_split:NnN \c__bnvs_comma_regex { #1 } \l_query_seq
```

Then each component is evaluated and the result is stored in `\l_ans_seq` that we have clear before use.

```
  \seq_map_inline:Nn \l_query_seq {
    \tl_clear:N \l_ans_tl
    \__bnvs_if_eval_query:nNTF { ##1 } \l_ans_tl {
      \seq_put_right:NV \l_ans_seq \l_ans_tl
    } {
      \seq_map_break:n {
        \msg_fatal:nnn { beanoves } { :n } { Circular~dependency~in~#1}
      }
    }
  }
```

We have managed all the comma separated components, we collect them back and append them to ⟨*tl variable*⟩.

```
  \exp_args:NNNx
  \__bnvs_group_end:
  \tl_put_right:Nn #2 { \seq_use:Nn \l_ans_seq , }
}
\cs_generate_variant:Nn \__bnvs_eval:nN { VN, xN }
```

`\BeanovesEval` `\BeanovesEval [`⟨*tl variable*⟩`] {`⟨*overlay queries*⟩`}`

⟨*overlay queries*⟩ is the argument of `?(...)` instructions. This is a comma separated list of single ⟨*overlay query*⟩'s.

This function evaluates the ⟨*overlay queries*⟩ and store the result in the ⟨*tl variable*⟩ when provided or leave the result in the input stream. Forwards to `\__bnvs_eval:nN` within a group. `\l_ans_tl` is used locally to store the result.

```
\NewDocumentCommand \BeanovesEval { s o m } {
  \__bnvs_group_begin:
  \tl_clear:N \l_ans_tl
  \IfBooleanTF { #1 } {
    \bool_set_true:N  \l__bnvs_no_counter_bool
  } {
    \bool_set_false:N \l__bnvs_no_counter_bool
```

```
  }
  \__bnvs_eval:nN { #3 } \l_ans_tl
  \IfValueTF { #2 } {
    \exp_args:NNNV
    \__bnvs_group_end:
    \tl_set:Nn #2 \l_ans_tl
  } {
    \exp_args:NV
    \__bnvs_group_end: \l_ans_tl
  }
}
```

5.5.9 Reseting slide ranges

\BeanovesReset `\beanovesReset [⟨first value⟩] {⟨Slide range name⟩}`

```
\NewDocumentCommand \BeanovesReset { O{1} m } {
  \__bnvs_reset:nn { #1 } { #2 }
  \ignorespaces
}
```

Forwards to `\__bnvs_reset:nn`.

__bnvs_reset:nn `\__bnvs_reset:nn {⟨first value⟩} {⟨slide range name⟩}`

Reset the counter to the given ⟨*first value*⟩. Clean the cached values also.

```
\cs_new:Npn \__bnvs_reset:nn #1 #2 {
  \bool_if:nTF {
    \__bnvs_if_in_p:n { #2/A } || \__bnvs_if_in_p:n { #2/Z }
  } {
    \__bnvs_gremove:n { #2/C }
    \__bnvs_gremove:n { #2//A }
    \__bnvs_gremove:n { #2//L }
    \__bnvs_gremove:n { #2//Z }
    \__bnvs_gremove:n { #2//N }
    \__bnvs_gput:nn { #2/C0 } { #1 }
  } {
    \msg_warning:nnn { beanoves } { :n } { Unknown~name:~#2 }
  }
}
\makeatother
\ExplSyntaxOff
```

⟨/package⟩