# beamer named overlay specifications with beanoves

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### Abstract

This package allows the management of multiple named slide number sets in beamer documents. Named slide number sets are very handy both during edition and to manage complex and variable beamer overlay specifications. In particular, they allow to replace raw numbers in beamer <...> overlay specifications by logical identifiers. Demonstration files are available for download as part of the development repository.

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# 1 Minimal example

The document below is a contrived example to show how the  $\verb"beamer"$  overlay specifications have been extended.

```
\documentclass {beamer}
  \RequirePackage {beanoves}
  \begin{document}
   \Beanoves {
       A = 1:3,
      B = A.next::3,
       C = B.next,
  \begin{frame}
  {\Large Frame \insertframenumber}
11 {\Large Slide \insertslidenumber}
12 \visible<?(A.1)> {Only on slide 1}\\
  \visible < ?(B.1) - ?(B.last) > {Only on slide 3 to 5} 
  \visible < (C.1) > {Only on slide 6} \
15 \visible<?(A.2)> \{0nly on slide 2\}\\
_{16} \ \ensuremath{\mbox{visible}\mbox{?(B.2:B.last)> {Only on slide 4 to 5}}\ensuremath{\mbox{\lambda}}
17 \visible<?(C.2)> \{0nly on slide 7\}\\
18 \visible<?(A.next)-> {From slide 3}\\
19 \visible<?(B.3:B.last)> \{0nly on slide 5\}\\
21 \end{frame}
  \end{document}
```

On line 4, we use the \Beanoves command to declare named overlay sets. On line 5, we declare an overlay set named 'A', which is a range starting at slide 1 and ending at slide 3. On line 12, the extended named overlay specification ?(A.1) stands for 1 because 1 is the first index of the overlay set named A. On line 15, ?(A.2) stands for 2 whereas on line 18, ?(A.next) stands for 3. On line 6, we declare a second overlay set named 'B', starting after the 2 slides of 'A' namely 3. Its length is 3 meaning that its last slide number is 5, thus each ?(B.last) is replaced by 5. The next slide number after slide range 'B' is 6 which is also the start of the third slide range due to line 7.

# 2 Named overlay sets

### 2.1 Presentation

Within a beamer frame, there are different slides that appear in turn according to overlay specifications. The main overlay sets is a range of integers covering all the slide numbers, from one to the total amount of slides. In general, an overlay set is a range of positive integers identified by a unique name. The main practical interest is that such sets may be defined relative to one another, we can even have lists of overlay sets. Finally, we can use these lists to build and organize beamer overlay specifications logically.

### 2.2 Named overlay reference

A.1, C.2 are named overlay references, as well as A and Y!C.2. More precisely, they are string identifiers, each one representing a well defined static integer to be used in beamer overlay specifications. They can take one of the next forms.

```
(short name) : like A and C,
```

(frame id)!(short name): denoted by qualified names, like X!A and Y!C.

(short name) (dotted path) : denoted by full names like A.1 and C.2,

 $\langle frame id \rangle! \langle short name \rangle \langle dotted path \rangle$ : denoted by qualified full names like X!A.1 and Y!C.2.

The *short names* and *frame ids* are alphanumerical case sensitive identifiers, with possible underscores but no space nor leading digit. Unicode symbols above U+00AO are allowed if the underlying TeX engine supports it. Identifiers consisting only of lowercase letters and underscores are reserved by the package.

The dotted path is a string  $.\langle component_1 \rangle . \langle component_2 \rangle . . . . \langle component_n \rangle$ , where each  $\langle component_i \rangle$  denotes either an integer, eventually signed, or a  $\langle short\ name \rangle$ . The dotted path can be empty for which n is 0.

The mapping from named overlay references to integers is defined at the global T<sub>E</sub>X level to allow its use in \begin{frame}<...> and to share the same overlay sets between different frames. Hence the frame id due to the need to possibly target a particular frame.

### 2.3 Defining named overlay sets

In order to define *named overlay sets*, we can either execute the next \Beanoves command before a beamer frame environment, or use the beanoves option of this environment. The value of the beanoves option is similar to the argument of the \Beanoves commands, but the latter takes precedence on the former. This behaviour may be useful to input the very same source code into different frames and have different combinations of slides.

beanoves

```
\texttt{beanoves} \ = \ \{\langle \textit{ref}_1 \rangle = \langle \textit{spec}_1 \rangle \text{,} \ \langle \textit{ref}_2 \rangle = \langle \textit{spec}_2 \rangle \text{,...,} \ \langle \textit{ref}_n \rangle = \langle \textit{spec}_n \rangle \}
```

\Beanoves

```
\Beanoves\{\langle ref_1\rangle = \langle spec_1\rangle, \langle ref_2\rangle = \langle spec_2\rangle, \ldots, \langle ref_n\rangle = \langle spec_n\rangle\}
```

Each  $\langle ref \rangle$  key is a named overlay reference whereas each  $\langle spec \rangle$  value is an overlay set specifier. When the same  $\langle ref \rangle$  key is used multiple times, only the last one is taken into account.

### 2.3.1 Basic case

In the possible values for  $\langle spec \rangle$  below,  $\langle value \rangle$ ,  $\langle first \rangle$ ,  $\langle length \rangle$  and  $\langle last \rangle$  are algebraic expression possibly involving any named overlay reference defined above.

 $\langle value \rangle$ , the simple value specifiers for the whole signed integers set. If only the  $\langle key \rangle$  is provided, the  $\langle value \rangle$  defaults to 1.

 $\langle first \rangle$ : and  $\langle first \rangle$ ::, for the infinite range of signed integers starting at and including  $\langle first \rangle$ .

 $:\langle last \rangle$ , for the infinite range of signed integers ending at and including  $\langle last \rangle$ .

 $\langle first \rangle : \langle last \rangle$ ,  $\langle first \rangle : : \langle length \rangle$ ,  $: \langle length \rangle$ ,  $: : \langle length \rangle : \langle last \rangle$ , are variants for the finite range of signed integers starting at and including  $\langle first \rangle$ , ending at and including  $\langle last \rangle$ . At least one of  $\langle first \rangle$  or  $\langle last \rangle$  must be provided. We always have  $\langle first \rangle + \langle length \rangle = \langle last \rangle + 1$ .

When performed at the document level, the \Beanoves command starts by cleaning what was set by previous calls. When performed inside LATEX environments, each call cumulates with the previous. Notice that the argument of this function can contain macros: they will be exhaustively expanded at resolution time.

### 2.3.2 List specifiers

Also possible values are *list specifiers* which are comma separated lists of  $\langle ref \rangle = \langle spec \rangle$  definitions. The definition

```
\begin{split} \langle \mathit{key} \rangle = & \{\langle \mathit{ref}_1 \rangle = \langle \mathit{spec}_1 \rangle \,, \  \, \langle \mathit{ref}_2 \rangle = \langle \mathit{spec}_2 \rangle \,, \dots, \  \, \langle \mathit{ref}_n \rangle = \langle \mathit{spec}_n \rangle \} \\ \text{is a convenient shortcut for} \\ & \langle \mathit{key} \rangle \,. \, \langle \mathit{ref}_1 \rangle = \langle \mathit{spec}_1 \rangle \,, \\ & \langle \mathit{key} \rangle \,. \, \langle \mathit{ref}_2 \rangle = \langle \mathit{spec}_2 \rangle \,, \\ & \dots, \\ & \langle \mathit{key} \rangle \,. \, \langle \mathit{ref}_n \rangle = \langle \mathit{spec}_n \rangle \,. \end{split}
```

The rules above can apply individually to each line.

To support an array like syntax, we can omit the  $\langle ref \rangle$  key. The first missing key is replaced by 1, the second by 2, and so on.

#### 2.3.3 .n specifiers

 $\langle key \rangle$ .n= $\langle value \rangle$  is used to set the value of the index counter defined below.

# 3 Named overlay resolution

Turning a named overlay reference into the static integer it represents, as when above  $\langle ?(A.1) \rangle$  was replaced by 1, is denoted by named overlay resolution or simply resolution. This section is devoted to resolution rules depending on the definition of the named overlay set. Here  $\langle i \rangle$  denotes an integer whereas  $\langle first \rangle$ ,  $\langle last \rangle$  and  $\langle length \rangle$  stand for integers, or integer valued expressions.

### 3.1 Simple definitions

 $\langle key \rangle = \langle value \rangle$  For an unlimited range

reference	resolution
$\langle \texttt{key} \rangle$ .1	$\langle \textit{value}  angle$
$\langle  exttt{key}  angle$ . 2	$\langle {\it value}  angle + 1$
$ackslash \langle \mathtt{key}  angle . \langle \mathtt{i}  angle$	$\langle  extit{value}  angle + \langle  extit{i}  angle - 1$

 $\langle key \rangle = \langle first \rangle$ : as well as  $\langle first \rangle$ :.. For a range limited from below:

reference	resolution
$\langle \texttt{key} \rangle$ .1	$\langle first  angle$
$\langle \mathtt{key}  angle$ . 2	$\langle  extit{first}  angle + 1$
$\langle  exttt{key}  angle$ . $\langle  exttt{i}  angle$	$ig  \langle  extit{first}  angle + \langle  extit{i}  angle - 1$
$\langle  extit{ extit{key}}  angle$ . $ exttt{ exttt{previous}}$	$\langle first  angle - 1$

 $\langle key \rangle = : \langle last \rangle$  For a range limited from above:

reference	resolution
$\langle key \rangle$ .1	$\langle last \rangle$
$\langle  extit{ extit{key}}  angle$ . 0	$\langle last \rangle - 1$
$\langle  exttt{ extit{key}}  angle$ . $\langle  exttt{ extit{i}}  angle$	$\langle last \rangle + \langle i \rangle - 1$
$\langle  extit{key}  angle$ . next	$\langle last \rangle + 1$

 $\langle key \rangle = \langle first \rangle : \langle last \rangle$  as well as variants  $\langle first \rangle : : \langle length \rangle, :: \langle length \rangle : \langle last \rangle$  or  $: \langle last \rangle :: \langle length \rangle$ , which are equivalent provided  $\langle first \rangle + \langle length \rangle = \langle last \rangle + 1$ .

For a range limited from both above and below:

reference	resolution
$\langle \texttt{key} \rangle$ .1	$\langle first  angle$
$\langle \mathtt{key}  angle$ . 2	$\langle first  angle + 1$
$\langle \mathtt{key}  angle$ . $\langle \mathtt{i}  angle$	$ig \langle  extit{first}  angle + \langle  extit{i}  angle - 1 ig )$
$\langle  extit{key} angle$ . <code>previous</code>	$\langle first  angle - 1$
$\langle  extit{key}  angle$ . last	$\langle  \mathit{last}  angle$
$\langle  exttt{ extit{key}}  angle$ . $ exttt{ exttt{ exttt{next}}}$	$\langle last  angle + 1$
$\langle  exttt{ extit{key}}  angle$ . length	$\langle \mathit{length}  angle$
$\langle  exttt{ extit{key}}  angle$ . range	$\max(0,\langle first\rangle)$ ''-'' $\max(0,\langle last\rangle)$

Notice that the resolution of  $\langle key \rangle$ .range is not an algebraic difference, and negative integers do not make sense there while in beamer context.

For example

```
\Beanoves {
    A = 3:8, % or equivalently A = 3::6, A = ::6:8 and A = :8::6
3 }
4 \begin{frame} {Frame \insertframenumber} {Slide \insertslidenumber}
5 \ttfamily
6 \BeanovesEval[see](A.1)
  \BeanovesEval[see](A.-1)
8 \BeanovesEval[see](A.previous) == 2,
9 \BeanovesEval[see](A.last)
                                  == 8,
10 \BeanovesEval[see](A.next)
                                  == 9,
11 \BeanovesEval[see](A.length)
                                  == 6,
12 \BeanovesEval[see](A.range)
                                  == 3-8,
13 \end{frame}
```

For example both ?(A.next), ?(A.last+1), ?(A.1+A.length) give the same result as soon as the slide range named 'A' has been properly defined with a starting value and a length.

### 3.2 Counters

Each named overlay set defined has a dedicated value counter which is some kind of variable that can be used and incremented. A simple  $\langle key \rangle$  named value reference is resolved into the position of this value counter. For each frame, this variable is initialized to the first available amongst  $\langle value \rangle$ ,  $\langle key \rangle$ .first or  $\langle key \rangle$ .last. If none is available, an error is raised.

For each named overlay set defined, we also have an implicit index counter always starting at 1, its actual value is an integer denoted  $\langle n \rangle$ . The  $\langle \textit{key} \rangle$ .n named index reference is resolved into  $\langle \textit{key} \rangle . \langle n \rangle$ , which in turn is resolved according to the preceding rules.

Additionally, resolution rules are provided for the named value references:

- $\langle \text{key} \rangle + = \langle \text{integer expression} \rangle$ , resolve  $\langle \text{integer expression} \rangle$  into  $\langle \text{integer} \rangle$ , advance the value counter by  $\langle \text{integer} \rangle$  and use the new position. Here  $\langle \text{integer expression} \rangle$  is the longest character sequence with no space<sup>1</sup>.
- ++ $\langle key \rangle$ , advance the value counter for  $\langle key \rangle$  by 1 and use the new position.
- $\langle \mathbf{key} \rangle + +$ , use the actual position and advance the value counter for  $\langle key \rangle$  by 1.

We have resolution rules as well for the named index references:

- $\langle \mathtt{key} \rangle$ .n+= $\langle \mathtt{integer\ expression} \rangle$ , resolve  $\langle \mathtt{integer\ expression} \rangle$  into  $\langle \mathtt{integer} \rangle$ , advance the implicit index counter associate to  $\langle \mathtt{key} \rangle$  by  $\langle \mathtt{integer} \rangle$  and use the resolution of  $\langle \mathtt{key} \rangle$ .n.
  - Here again,  $\langle integer\ expression \rangle$  denotes the longest character sequence with no space.
- $\langle key \rangle$ .++n, ++ $\langle key \rangle$ .n, advance the implicit index counter associate to  $\langle key \rangle$  by 1 and use the resolution of  $\langle key \rangle$ .n,
- $\langle \texttt{key} \rangle$ .n++, use the resolution of  $\langle \texttt{key} \rangle$ .n and increment the implicit index counter associate to  $\langle \texttt{key} \rangle$  by 1.

In order to decrement a counter, one can increment with a negative value, no dedicated syntax is provided yet.

These counters are reset to their default value for each new frame, which is 1 for the  $\langle key \rangle$ .n counter, and whichever  $\langle key \rangle$ .first or  $\langle key \rangle$ .last is defined for the  $\langle key \rangle$  counter.

### 3.3 Dotted paths

 $\langle key \rangle . \langle i \rangle = \langle spec \rangle$ , All the preceding rules are overriden by this particular one and  $\langle key \rangle . \langle i \rangle$  resolves to the resolution of  $\langle spec \rangle$ .

In the frame example below, we use the \BeanovesEval command for the demonstration. It is mainly used for debugging and testing purposes.

```
1 \Beanoves {
2   A = 3,
3   A.3 = 0,
4 }
5 \begin{frame} {Frame \insertframenumber} {Slide \insertslidenumber}
6 \ttfamily
7 \BeanovesEval[see] (A.1) == 3,
8 \BeanovesEval[see] (A.2) == 4,
9 \BeanovesEval[see] (A.-1)== 1,
10 \BeanovesEval[see] (A.3) == 0,
11 \end{frame}
```

 $<sup>^{1}</sup>$ The parser for algebraic expression is very rudimentary.

Without line 3, A.3 would be evaluated to 5.

 $\langle \textbf{key} \rangle. \langle \textbf{c}_1 \rangle. \langle \textbf{c}_2 \rangle \dots \langle \textbf{c}_k \rangle = \langle \textbf{range spec} \rangle \text{ When a dotted path has more than one component, a named overlay reference like A.1.2 needs some well defined resolution rule to avoid ambiguity. To resolve one level of such a reference <math display="block">\langle \textbf{key} \rangle. \langle \textbf{c}_1 \rangle. \langle \textbf{c}_2 \rangle \dots \langle \textbf{c}_n \rangle,$  we replace the longest  $\langle \textbf{key} \rangle. \langle \textbf{c}_1 \rangle. \langle \textbf{c}_2 \rangle \dots \langle \textbf{c}_k \rangle \text{ where } 0 \leq \textbf{k} \leq \textbf{n} \text{ by its definition } \langle \textbf{name'} \rangle. \langle \textbf{c'}_1 \rangle \dots \langle \textbf{c'}_p \rangle \text{ if any (the path can be empty). beanoves uses this one level resolution as many times as possible, but no more than a predefined limit to catch circular reference that would lead to an infinite TeX loop. One final resolution occurs with rules above if possible or an error is raised.$ 

For a named indexed reference like  $\langle key \rangle . \langle c_1 \rangle . \langle c_2 \rangle ... \langle c_n \rangle$ .n, we must first resolve  $\langle key \rangle . \langle c_1 \rangle . \langle c_2 \rangle ... \langle c_n \rangle$  into  $\langle name' \rangle$  with an empty dotted path, then retrieve the value of  $\langle name' \rangle$ .n denoted as  $\langle n' \rangle$  and finally use the resolved  $\langle key \rangle . \langle c_1 \rangle . \langle c_2 \rangle ... \langle c_n \rangle . \langle n' \rangle$ .

### 3.4 Frame id

Except for very special situations, the frame ids can be left unspecified. When no frame id was explicitly provided, beanoves uses the last frame id. At the beginning of each frame, the last frame id is set to the frame id of the current frame, which is denoted current frame id and defaults to ?. Then it gets updated after each named reference resolution. For example, the first time A.1 reference is resolved within a given frame, it is first translated to  $\langle current \ frame \ id \rangle$ !A.1, but when used just after Y!C.2, it becomes a shortcut to Y!A.1 because the last frame id was then Y.

In order to set the *frame id* of the current frame to  $\langle frame\ id \rangle$ , use the new beanoves id option of the beamer frame environment.

beanoves id

beanoves  $id=\langle frame\ id \rangle$ ,

We can use the same frame id for different frames to share named overlay sets.

# 4 ?(...) query expressions

This is the key feature of the beanoves package, extending beamer overlay specifications included between pointed brackets. Before the overlay specifications are processed by the beamer class, the beanoves package scans them for any occurrence of '?( $\langle queries \rangle$ )'. Each one is then evaluated and replaced by its resolved static counterpart. The overall result is finally forwarded to the beamer class.

The  $\langle queries \rangle$  argument is a comma separated list of individual  $\langle query \rangle$ 's of next table. Sometimes, using  $\langle key \rangle$  range is not allowed because the resolution would be interpreted as an algebraic difference instead of a beamer range. If it is not possible, an error is raised.

query	resolution	limitation
$\langle  extsf{start expr}  angle$	$\langle start  angle$	
$\langle  extit{start expr} \rangle$ :	$\langle start  angle$ -	$\operatorname{no}\langle {\it key} \rangle$ .range
$\langle  exttt{start expr}  angle : \langle  exttt{end expr}  angle$	$\langle start  angle$ - $\langle end  angle$	$\operatorname{no}\langle {\it key}  angle$ .range
$:: \langle \mathtt{length} \ \mathtt{expr} \rangle : \langle \mathtt{end} \ \mathtt{expr} \rangle$	$\langle start  angle$ - $\langle end  angle$	$\operatorname{no}\langle  extit{key} angle$ .range
: $\langle  exttt{end expr}  angle$	$-\langle end \rangle$	$\operatorname{no}\langle {\it key} \rangle$ .range
:	_	
$\langle  extit{start expr}  angle ::$	$\langle start  angle$ -	$\operatorname{no}\langle {\it key} \rangle$ .range
$\langle  exttt{start expr}  angle :: \langle  exttt{length expr}  angle$	$\langle start  angle$ - $\langle end  angle$	$\operatorname{no}\langle {\it key} \rangle$ .range
$:\langle  ext{end expr} \rangle :: \langle  ext{length expr} \rangle$	$\langle start  angle$ - $\langle end  angle$	$\operatorname{no}\langle  extit{key} angle$ .range
::	_	

Here  $\langle start\ expr \rangle$ ,  $\langle end\ expr \rangle$  and  $\langle length\ expr \rangle$  both denote algebraic expressions possibly involving named overlay references and counters. As integers, they are respectively resolved into  $\langle start \rangle$ ,  $\langle end \rangle$  and  $\langle length \rangle$ .

Notice that nesting ?(...) query expressions is not supported.

# 5 Support

See https://github.com/jlaurens/beanoves. One can report issues.

# 6 Implementation

Identify the internal prefix (IATEX3 DocStrip convention, unused).

1 (@@=bnvs)

Reserved namespace: identifiers containing the case insensitive string beanoves or the case insensitive string bnvs delimited by two non characters.

### 6.1 Package declarations

- 2 \NeedsTeXFormat{LaTeX2e}[2020/01/01]
- 3 \ProvidesExplPackage
- {beanoves}
- 5 {2023/01/07}
- 6 {1.0}
- {Named overlay specifications for beamer}

## 6.2 Facility layer: definitions and naming

In order to make the code shorter and easier to read, we add a layer over LATEX3. The c and v argument specifiers take a different meaning when used in a function which name contains with bnvs or BNVS. Where LATEX3 would transform 1\_bnvs\_key\_tl into \1\_bnvs\_key\_tl, bnvs will directly transform name into \1\_bnvs\_key\_tl. The type of the local variable used depends on the context and may be seq or int for example. There are however a pair of exceptions mentionned below. For a better reading experience, 'key' will generally stand for \1\_bnvs\_key\_tl, whereas 'path sequence' will generally stand for \1 bnvs path seq. Other similar shortcuts are used as well.

Functions with BNVS in their names are management functions. They belong to a deeper layer and do not contain any beanoves specific logic.

```
\BNVS:c
                            \BNVS:c \{\langle cs \ core \ name \rangle\}
         \BNVS_1:cn
                            \BNVS_1:cn \{\langle local\ variable\ core\ name \rangle\}\ \{\langle\ type\ \rangle\}
         \BNVS_g:cn
                            \BNVS_g:cn \{\langle global\ variable\ core\ name\rangle\}\ \{\langle\ type\ \rangle\}
                            \BNVS_g_prop:c \{\langle global\ prop\ core\ name\rangle\}
                            These are naming functions.
                               8 \cs_new:Npn \BNVS:c
                                                                    #1
                                                                              { __bnvs_#1
                               9 \cs_new:Npn \BNVS_1:cn #1 #2 { 1__bnvs_#1_#2 }
                               10 \cs_new:Npn \BNVS_g:cn #1 #2 { g__bnvs_#1_#2 }
\BNVS_use_raw:c
                            \BNVS_use_raw:c \{\langle cs name \rangle\}
                            \verb|\BNVS_use_raw:Nc| \langle \mathit{function} \rangle \ \{\langle \mathit{cs} \ \mathit{name} \rangle\}|
\BNVS_use_raw:Nc
\BNVS_use_raw:nc
                            \BNVS_use_raw:nc \{\langle tokens \rangle\} \{\langle cs name \rangle\}
\BNVS_use:c
                            \BNVS_use:c \{\langle cs \ core \rangle\}
\BNVS_use:Nc
                            \BNVS_use:Nc \langle function \rangle \ \{\langle cs \ core \rangle\}
\BNVS_use:nc
                            \BNVS_use:nc \{\langle tokens \rangle\}\ \{\langle cs\ core \rangle\}
```

\BNVS\_use\_raw:c is a wrapper over \use:c. possibly prepended with some code. It needs 3 expansion steps just like \BNVS\_use:c. The other are used to expand \use:c twice before usage by  $\langle function \rangle$  or  $\langle tokens \rangle$ . The first argument of  $\langle function \rangle$  has type N. The next token after  $\langle tokens \rangle$  will have type N too.  $\langle cs\ name \rangle$  is a full cs name whereas  $\langle cs\ core \rangle$  will be prepended with the appropriate prefix.

```
11 \cs new:Npn \BNVS use raw:N #1 { #1 }
12 \cs_new:Npn \BNVS_use_raw:c #1 {
    \exp_args:NNo
13
    \exp_last_unbraced:No
    \BNVS_use_raw:N { \use:c { #1 } }
16 }
17 \cs_new:Npn \BNVS_use:c #1 {
    \BNVS_use_raw:c { \BNVS:c { #1 } }
19 }
20 \cs_new:Npn \BNVS_use_raw:NN #1 #2 {
21
22 }
23 \cs_new:Npn \BNVS_use_raw:nN #1 #2 {
24
25 }
26 \cs_new:Npn \BNVS_use_raw:Nc #1 #2 {
27
    \exp_args:NNNo
    \exp_last_unbraced:NNo
    \BNVS_use_raw:NN #1 { \use:c { #2 } }
29
30 }
31 \exp_args_generate:n { NNno }
32 \cs_new:Npn \BNVS_use_raw:nc #1 #2 {
    \exp_args:NNno
33
    \exp_last_unbraced:Nno
    \BNVS_use_raw:nN { #1 } { \use:c { #2 } }
37 \cs_new:Npn \BNVS_use:Nc #1 #2 {
    \BNVS_use_raw:Nc #1 { \BNVS:c { #2 } }
38
40 \cs_new:Npn \BNVS_use:nc #1 #2 {
    \BNVS_use_raw:nc { #1 } { \BNVS:c { #2 } }
42 }
```

\BNVS\_new:cpn \BNVS\_set:cpn \BNVS\_use:c

\BNVS\_new:cpn is like \cs\_new:cpn except that the name argument is tagged for beanoves package. Similarly for \BNVS\_set:cpn and \BNVS\_use:c.

```
43 \cs_new:Npn \BNVS_log:n #1 { }
44 \cs_generate_variant:Nn \BNVS_log:n { x }
45 \cs_new:Npn \BNVS_DEBUG_on: {
    \cs_set:Npn \BNVS_DEBUG_log:n { \BNVS_log:n }
47 }
48 \cs_new:Npn \BNVS_DEBUG_off: {
    \cs_set:Npn \BNVS_DEBUG_log:n { \use_none:n }
49
50 }
51 \BNVS_DEBUG_off:
52 \cs_new:Npn \BNVS_new:cpn #1 {
    \cs_new:cpn { \BNVS:c { #1 } }
53
<sub>54</sub> }
55 \cs_new:Npn \BNVS_set:cpn #1 {
    \cs_set:cpn { \BNVS:c { #1 } }
56
57 }
58 \cs_generate_variant:Nn \cs_generate_variant:Nn { c }
59 \cs_new:Npn \BNVS_generate_variant:cn #1 {
    \cs_generate_variant:cn { \BNVS:c { #1 } }
61 }
```

## 6.3 logging

Utility message.

```
62 \msg_new:nnn { beanoves } { :n } { #1 }
63 \msg_new:nnn { beanoves } { :nn } { #1~(#2) }
64 \BNVS_new:cpn { warning:n } {
    \msg_warning:nnn { beanoves } { :n }
65
67 \BNVS_generate_variant:cn { warning:n } { x }
68 \BNVS_new:cpn { error:n } {
    \msg_error:nnn { beanoves } { :n }
70 }
71 \BNVS_new:cpn { error:x } {
    \msg_error:nnx { beanoves } { :n }
72
73 }
74 \BNVS_new:cpn { fatal:n } {
    \msg_fatal:nnn { beanoves } { :n }
75
76 }
77 \BNVS_new:cpn { fatal:x } {
    \msg_fatal:nnx { beanoves } { :n }
78
79 }
```

# 6.4 Facility layer: Variables

\BNVS\_N\_new:c \BNVS\_v\_new:c

```
\BNVS_N_new:n \{\langle type \rangle\}
```

Creates typed utility functions, see usage below. Undefined when no longer used.  $\langle type \rangle$  is one of t1, seq...

```
\cs_new:Npn \BNVS_N_new:c #1 {
     \cs_new:cpn { BNVS_#1:c } ##1 {
81
       1 \BNVS:c{ ##1 } \tl_if_empty:nF { ##1 } { _ } #1
82
83
     \cs_new:cpn { BNVS_#1_new:c } ##1 {
84
       \use:c { #1_new:c } { \use:c { BNVS_#1:c } { ##1 } }
85
86
     \cs_new:cpn { BNVS_#1_use:c } ##1 {
87
       \use:c { \use:c { BNVS_#1:c } { ##1 } }
88
89
     \cs_new:cpn { BNVS_#1_use:Nc } ##1 ##2 {
90
       \BNVS_use_raw:Nc
91
         ##1 { \use:c { BNVS_#1:c } { ##2 } }
92
93
     \cs_new:cpn { BNVS_#1_use:nc } ##1 ##2 {
94
       \BNVS_use_raw:nc
95
         { ##1 } { \use:c { BNVS_#1:c } { ##2 } }
96
     }
97
98 }
   \cs_new:Npn \BNVS_v_new:c #1 {
99
     \cs_new:cpn { BNVS_#1_use:Nv } ##1 ##2 {
100
       \BNVS_use_raw:nc
101
         { \exp_args:NV ##1 }
         { \BNVS_use_raw:c { BNVS_#1:c } { ##2 } }
103
104
     \cs_new:cpn { BNVS_#1_use:nv } ##1 ##2 {
105
       \BNVS_use_raw:nc
106
         { \exp_args:NnV \use:n { ##1 } }
107
         { \BNVS_use_raw:c { BNVS_#1:c } { ##2 } }
108
     }
109
110 }
111 \BNVS_N_new:c { bool }
112 \BNVS_N_new:c { int }
113 \BNVS_v_new:c { int }
114 \BNVS_N_new:c { tl }
115 \BNVS_v_new:c { tl }
116 \BNVS_N_new:c { str }
117 \BNVS_v_new:c { str }
118 \BNVS_N_new:c { seq }
119 \BNVS_v_new:c { seq }
  \cs_undefine:N \BNVS_N_new:c
  \cs_new:Npn \BNVS_use:Ncn #1 #2 #3 {
     \BNVS_use_raw:c { BNVS_#3_use:Nc }
                                                 { #2 }
123 }
  \cs_new:Npn \BNVS_use:ncn #1 #2 #3 {
124
     \BNVS_use_raw:c { BNVS_#3_use:nc } { #1 } { #2 }
125
126 }
   \cs_new:Npn \BNVS_use:Nvn #1 #2 #3 {
127
128
     \BNVS_use_raw:c { BNVS_#3_use:Nv }
                                                 { #2 }
129 }
  \cs_new:Npn \BNVS_use:nvn #1 #2 #3 {
     \BNVS_use_raw:c { BNVS_#3_use:nv } { #1 } { #2 }
132 }
```

```
\cs_new:Npn \BNVS_use:Ncncn #1 #2 #3 {
     \BNVS_use:ncn {
134
       \BNVS_use:Ncn
                              { #2 } { #3 }
                         #1
135
136
137 }
   \cs_new:Npn \BNVS_use:ncncn #1 #2 #3 {
138
     \BNVS_use:ncn {
139
       \BNVS_use:ncn { #1 } { #2 } { #3 }
140
141
142 }
   \cs_new:Npn \BNVS_use:Nvncn #1 #2 #3 {
143
     \BNVS_use:ncn {
144
       \BNVS_use:Nvn
                              { #2 } { #3 }
                         #1
145
146
147 }
   \cs_new:Npn \BNVS_use:nvncn #1 #2 #3 {
148
     \BNVS_use:ncn {
149
       \BNVS_use:nvn { #1 } { #2 } { #3 }
150
     }
151
152 }
   \cs_new:Npn \BNVS_use:Ncncncn #1 #2 #3 #4 #5 {
153
     \BNVS_use:ncn {
154
       \BNVS_use:Ncncn
                           #1
                                 { #2 } { #3 } { #4 } { #5 }
155
     }
156
157 }
   \cs_new:Npn \BNVS_use:ncncncn #1 #2 #3 #4 #5 {
158
     \BNVS_use:ncn {
159
       \BNVS_use:ncncn { #1 } { #2 } { #3 } { #4 } { #5 }
160
161
162 }
163 \cs_new:Npn \BNVS_new_c:nc #1 #2 {
     \BNVS_new:cpn { #1_#2:c } {
       \label{local_bnvs_use_raw:c} $$BNVS_use_raw:c { $\mu_1$_use:nc } { BNVS_use_raw:c { $\mu_1$_2:N } }
165
     }
166
167 }
   \cs_new:Npn \BNVS_new_cn:nc #1 #2 {
168
     \BNVS_new:cpn { #1_#2:cn } ##1 {
169
170
       \BNVS_use:ncn { \BNVS_use_raw:c { #1_#2:Nn } } { ##1 } { #1 }
171
172 }
   \cs_new:Npn \BNVS_new_cnn:ncN #1 #2 #3 {
173
174
     \BNVS_new:cpn { #2:cnn } ##1 {
       \BNVS_use:Ncn { #3 } { ##1 } { #1 }
175
176
177 }
   \cs_new:Npn \BNVS_new_cnn:nc #1 #2 {
178
     \BNVS_use_raw:nc {
179
       \BNVS_new_cnn:ncN { #1 } { #1_#2 }
180
     } { #1_#2:Nnn }
181
182 }
   \cs_new:Npn \BNVS_new_cnv:ncN #1 #2 #3 {
     \BNVS_new:cpn { #2:cnv } ##1 ##2 {
184
       \BNVS_tl_use:nv {
185
         \BNVS_use:Ncn #3 { ##1 } { #1 } { ##2 }
186
```

```
}
187
188
189 }
   \cs_new:Npn \BNVS_new_cnv:nc #1 #2 {
190
     \BNVS_use_raw:nc {
191
       \BNVS_new_cnv:ncN { #1 } { #1_#2 }
192
     } { #1_#2:Nnn }
193
194
   \cs_new:Npn \BNVS_new_cnx:ncN #1 #2 #3 {
     \BNVS_new:cpn { #2:cnx } ##1 ##2 {
       \exp_args:Nnx \use:n {
197
         \BNVS_use:Ncn #3 { ##1 } { #1 } { ##2 }
198
199
     }
200
201 }
   \cs_new:Npn \BNVS_new_cnx:nc #1 #2 {
202
     \BNVS_use_raw:nc {
203
       \BNVS_new_cnx:ncN { #1 } { #1_#2 }
204
     } { #1_#2:Nnn }
205
206 }
   \cs_new:Npn \BNVS_new_cc:ncNn #1 #2 #3 #4 {
207
     \BNVS_new:cpn { #2:cc } ##1 ##2 {
208
       \BNVS_use:Ncncn #3 { ##1 } { ##2 } { #4 }
209
     }
210
211 }
   \cs_new:Npn \BNVS_new_cc:ncn #1 #2 {
212
     \BNVS_use_raw:nc {
213
       \BNVS_new_cc:ncNn { #1 } { #1_#2 }
214
     } { #1_#2:NN }
215
217 \cs_new:Npn \BNVS_new_cc:nc #1 #2 {
     \BNVS_new_cc:ncn { #1 } { #2 } { #1 }
218
219 }
  \cs_new:Npn \BNVS_new_cn:ncNn #1 #2 #3 #4 {
220
     \BNVS_new:cpn { #2:cn } ##1 {
221
       \BNVS_use:Ncn #3 { ##1 } { #1 }
223
224 }
225
   \cs_new:Npn \BNVS_new_cn:ncn #1 #2 {
     \BNVS_use_raw:nc {
       \BNVS_new_cn:ncNn { #1 } { #1_#2 }
228
     } { #1_#2:Nn }
229 }
   \cs_new:Npn \BNVS_new_cv:ncNn #1 #2 #3 #4 {
230
     \BNVS_new:cpn { #2:cv } ##1 ##2 {
       \BNVS_use:nvn {
232
         \BNVS_use:Ncn #3 { ##1 } { #1 }
       } { ##2 } { #4 }
234
     }
235
236 }
   \cs_new:Npn \BNVS_new_cv:ncn #1 #2 {
238
     \BNVS_use_raw:nc {
       \BNVS_new_cv:ncNn { #1 } { #1_#2 }
239
     } { #1_#2:Nn }
```

```
242 \cs_new:Npn \BNVS_new_cv:nc #1 #2 {
             \BNVS_new_cv:ncn { #1 } { #2 } { #1 }
  243
  244 }
tl clear:NSUCCESS
  ^{245} \sl new:\sl n
             \BNVS_use_raw:Nc #1
                                                               { \BNVS_1:cn { #2 } { #3 } }
  247 }
        \cs_new:Npn \BNVS_1_use:ncn #1 #2 #3 {
  248
             \BNVS_use_raw:nc { #1 } { \BNVS_1:cn { #2 } { #3 } }
  249
  250 }
        \cs_new:Npn \BNVS_g_use:Ncn #1 #2 #3 {
             \BNVS_use_raw:Nc #1 { \BNVS_g:cn { #2 } { #3 } }
  253 }
  254
        \cs_new:Npn \BNVS_g_use:ncn #1 #2 #3 {
             \BNVS_use_raw:nc { #1 } { \BNVS_g:cn { #2 } { #3 } }
  255
  256 }
        \cs_new:Npn \BNVS_g_prop_use:Nc #1 #2 {
  257
             \BNVS_use_raw:Nc #1 { \BNVS_g:cn { #2 } { prop } }
  258
  259 }
        \cs_new:Npn \BNVS_g_prop_use:nc #1 #2 {
  260
  261
             \BNVS_use_raw:nc { #1 } { \BNVS_g:cn { #2 } { prop } }
  262 }
        \cs_new:Npn \BNVS_exp_args:Nvvv #1 #2 #3 #4 {
             \BNVS_use:ncncncn { \exp_args:NVVV #1 }
                  { #2 } { t1 } { #3 } { t1 } { #4 } { t1 }
  265
  266 }
  267 \cs_generate_variant:Nn \prg_new_conditional:Npnn { c }
  268 \cs_new:Npn \BNVS_new_conditional:cpnn #1 {
             \prg_new_conditional:cpnn { \BNVS:c { #1 } }
  269
  270 }
  271 \cs_generate_variant:Nn \prg_generate_conditional_variant:Nnn { c }
  272 \cs_new:Npn \BNVS_generate_conditional_variant:cnn #1 {
             \prg_generate_conditional_variant:cnn { \BNVS:c { #1 } }
  274 }
  275 \cs_new:Npn \BNVS_new_conditional_vn:cNnn #1 #2 #3 #4 {
             \BNVS_new_conditional:cpnn { #1:vn } ##1 ##2 { #4 } {
  276
                  \BNVS_use:Nvn #2 { ##1 } { #3 } { ##2 } {
  277
                      \prg_return_true:
  278
  279
                      \prg_return_false:
  280
  281
             }
  282
  283 }
        \cs_new:Npn \BNVS_new_conditional_vn:cnn #1 #2 {
  284
             \BNVS_use:nc {
  285
                 \BNVS_new_conditional_vn:cNnn { #1 }
  286
             } { #1:nn TF } { #2 }
  287
  288 }
         \cs_new:Npn \BNVS_new_conditional_vc:cNnn #1 #2 #3 #4 {
  289
             \BNVS_new_conditional:cpnn { #1:vc } ##1 ##2 { #4 } {
  290
                  \BNVS_use:Nvn #2 { ##1 } { #3 } { ##2 } {
  291
  292
                      \prg_return_true:
                 } {
```

```
\prg_return_false:
                     }
   295
               }
   296
   297 }
           \cs_new:Npn \BNVS_new_conditional_vc:cnn #1 {
   298
                \BNVS_use:nc {
   299
                      \BNVS_new_conditional_vc:cNnn { #1 }
   300
                } { #1:ncTF }
   301
   302 }
          \cs_new:Npn \BNVS_new_conditional_vc:cNn #1 #2 #3 {
   303
                \BNVS_new_conditional:cpnn { #1:vc } ##1 ##2 { #3 } {
   304
                      \BNVS_tl_use:Nv #2 { ##1 } { ##2 } {
   305
                           \prg_return_true:
   306
                     } {
   307
                           \prg_return_false:
   308
   309
                }
   310
   311 }
          \cs_new:Npn \BNVS_new_conditional_vc:cn #1 {
   312
                \BNVS_use:nc {
   313
                     \BNVS_new_conditional_vc:cNn { #1 }
   314
                } { #1:ncTF }
   315
   316 }
6.4.1 Regex
   317 \cs_new:Npn \BNVS_regex_use:Nc #1 #2 {
                \BNVS_use_raw:Nc #1 { c \BNVS:c { #2 } _regex }
   319 }
\label{local_noise} $$\sum_{\substack{n=0\\ n \in \mathbb{N}}} {\operatorname{conce}:\operatorname{NnTF} \ \langle regex\ variable} \ \{\langle expression \rangle\}$$
\{\langle yes\ code \rangle\}\ \{\langle\ \rangle\}no code
\verb|\cline | $$ \subseteq \operatorname{InTF} \{\langle \operatorname{regex} \rangle\} \ \{\langle \operatorname{expression} \rangle\} 
\{\langle yes\ code \rangle\}\ \{\langle\ \rangle\}no code
\verb|\cline| Lorentz | Lore
}}no code
These are shortcuts to
        • \regex_match_once:NnNTF with the match sequence as N argument
        • \regex_match_once:nnNTF with the match sequence as N argument
        • \regex_split:NnNTF with the split sequence as last N argument
   320 \BNVS_new_conditional:cpnn { match_once:Nn } #1 #2 { T, F, TF } {
                \BNVS_use:ncn {
                      \regex_extract_once:NnNTF #1 { #2 }
   322
                } { match } { seq } {
   324
                     \prg_return_true:
               } {
   325
```

\\_\_bnvs\_match\_once:Nn*TF* 

\\_\_bnvs\_match\_once:Nv<u>TF</u> \\_\_bnvs\_match\_once:nn<u>TF</u>

\\_\_bnvs\_regex\_split:cnTF

\prg\_return\_false:

326

327 } 328 }

```
329 \BNVS_new_conditional:cpnn { match_once:Nv } #1 #2 { T, F, TF } {
     \BNVS_seq_use:nc {
330
       \BNVS_tl_use:nv {
331
         \regex_extract_once:NnNTF #1
332
       } { #2 }
333
     } { match } {
334
       \prg_return_true:
335
     } {
336
337
       \prg_return_false:
     }
338
339 }
   \BNVS_new_conditional:cpnn { match_once:nn } #1 #2 { T, F, TF } {
340
     \BNVS_seq_use:nc {
341
       \regex_extract_once:nnNTF { #1 } { #2 }
342
     } { match } {
343
       \prg_return_true:
344
345
       \prg_return_false:
346
     }
347
348 }
   \BNVS_new_conditional:cpnn { regex_split:cnc } #1 #2 #3 { T, F, TF } {
349
     \BNVS_seq_use:nc {
350
       \BNVS_regex_use:Nc \regex_split:NnNTF { #1 } { #2 }
351
     } { #3 } {
352
       \prg_return_true:
353
     } {
354
       \prg_return_false:
355
     }
356
357 }
_{\mbox{\scriptsize 358}} \BNVS_new_conditional:cpnn { regex_split:cn } #1 #2 { T, F, TF } {
     \BNVS_seq_use:nc {
359
       \BNVS_regex_use:Nc \regex_split:NnNTF { #1 } { #2 }
360
     } { split } {
361
       \prg_return_true:
362
    } {
363
       \prg_return_false:
364
365
366 }
```

#### 6.4.2 Token lists

391 }

```
\__bnvs_tl_clear:c
                                       \verb|\__bnvs_tl_clear:c {| \langle core \ key_t | l \rangle }|
                                       \ \ \_bnvs\_tl\_use:c \ \{\langle core \rangle\}
\__bnvs_tl_use:c
                                       \ \ \_bnvs\_tl\_count:c \ \{\langle core \rangle\}
\__bnvs_tl_set_eq:cc
                                       \verb|\ \_bnvs_tl_set_eq:cc | \{\langle \mathit{lhs} \ \mathit{core} \ \mathit{name} \rangle\} | \{\langle \mathit{rhs} \ \mathit{core} \ \mathit{name} \rangle\}|
\__bnvs_tl_set:cn
                                       \verb|\__bnvs_tl_set:cn | \{\langle core \rangle\} | \{\langle tl \rangle\}|
\__bnvs_tl_set:(cv|cx)
                                        \ \ \_bnvs\_tl\_set:cv \ \{\langle core \rangle\} \ \{\langle value\ core\ name \rangle\}
\__bnvs_tl_put_left:cn
                                       \verb|\__bnvs_tl_put_left:cn | \{\langle \mathit{core}\rangle\} | \{\langle \mathit{tl}\rangle\}|
\__bnvs_tl_put_right:cn
\__bnvs_tl_put_right:(cx|cv)
                                       \ __bnvs_tl_put_right:cn \{\langle core \rangle\} \{\langle tl \rangle\}
                                       \verb|\__bnvs_tl_put_right:cv = \{\langle \mathit{core}\rangle\} \ \{\langle \mathit{value core name}\rangle\}|
                              These are shortcuts to
                                  • \tl_clear:c {l_bnvs_\(core\)_tl}
                                  • \tl_use:c {l__bnvs_\langle core \_tl}
                                  • \tl_set_eq:cc {l__bnvs_\langle lhs core\_tl}{l__bnvs_\langle rhs core\_tl}
                                  • \tl_set:cv {l__bnvs_\( core \)_tl}{l__bnvs_\( value core \)_tl}
                                  • tl_set:cx \{l_bnvs_\langle core \rangle_tl\}\{\langle tl \rangle\}
                                  • tl_put_left:cn \{l_bnvs_{core}_tl\}\{\langle tl \rangle\}
                                  • tl\_put\_right:cn \{l\_bnvs\_\langle core \rangle\_tl\}\{\langle tl \rangle\}
                                  • \tl_put_right:cv {l__bnvs_\(core\)_tl}{l__bnvs_\(value core\)_tl}
                                367 \cs_new:Npn \BNVS_new_conditional_vnc:cNn #1 #2 #3 {
                                      \BNVS_new_conditional:cpnn { #1:vnc } ##1 ##2 ##3 { #3 } {
                                         \BNVS_tl_use:Nv #2 { ##1 } { ##2 } { ##3 } {
                               369
                                           \prg_return_true:
                               370
                                        } {
                               371
                                            \prg_return_false:
                               372
                                        }
                               373
                                      }
                               374
                               375 }
                                   \cs_new:Npn \BNVS_new_conditional_vnc:cn #1 {
                                      \BNVS_use:nc {
                               377
                                        \BNVS_new_conditional_vnc:cNn { #1 }
                               378
                                      } { #1:nncTF }
                               379
                               380 }
                                   \cs_new:Npn \BNVS_new_conditional_vvnc:cNn #1 #2 #3 {
                               381
                                      \BNVS_new_conditional:cpnn { #1:vvnc } ##1 ##2 ##3 ##4 { #3 } {
                               382
                                         \BNVS_tl_use:nv {
                               383
                                           \BNVS_tl_use:Nv #2 { ##1 }
                               384
                                         } { ##2 } { ##3 } { ##4 } {
                               385
                                           \prg_return_true:
                                         } {
                                            \prg_return_false:
                               389
                                     }
                               390
```

```
\cs_new:Npn \BNVS_new_conditional_vvnc:cn #1 {
     \BNVS_use:nc {
393
       \BNVS_new_conditional_vvnc:cNn { #1 }
394
     } { #1:nnncTF }
395
396 }
   \cs_new:Npn \BNVS_new_conditional_vvvc:cNn #1 #2 #3 {
397
     \BNVS_new_conditional:cpnn { #1:vvvc } ##1 ##2 ##3 ##4 { #3 } {
398
       \BNVS_tl_use:nv {
399
         \BNVS_tl_use:nv {
           \BNVS_tl_use:Nv #2 { ##1 }
401
         } { ##2 }
402
       } { ##3 } { ##4 } {
403
         \prg_return_true:
404
       } {
405
         \prg_return_false:
406
407
     }
408
409 }
   \cs_new:Npn \BNVS_new_conditional_vvvc:cn #1 {
411
     \BNVS_use:nc {
       \BNVS_new_conditional_vvvc:cNn { #1 }
412
     } { #1:nnncTF }
413
414 }
   \cs_new:Npn \BNVS_new_conditional_vvc:cNn #1 #2 #3 {
415
     \BNVS_new_conditional:cpnn { #1:vvc } ##1 ##2 ##3 { #3 } {
416
       \BNVS_tl_use:nv {
417
         \BNVS_tl_use:Nv #2 { ##1 }
418
       } { ##2 } { ##3 } {
419
         \prg_return_true:
420
       } {
422
         \prg_return_false:
423
       }
     }
424
425 }
  \cs_new:Npn \BNVS_new_conditional_vvc:cn #1 {
426
     \BNVS_use:nc {
427
       \BNVS_new_conditional_vvc:cNn { #1 }
428
429
     } { #1:nncTF }
430 }
  \cs_new:Npn \BNVS_new_tl_c:c {
     \BNVS_new_c:nc { tl }
433 }
434 \BNVS_new_tl_c:c { clear }
  \BNVS_new_tl_c:c { use }
   \BNVS_new_tl_c:c { count }
437
   \BNVS_new:cpn { tl_set_eq:cc } #1 #2 {
438
     \BNVS_use:ncncn { \tl_set_eq:NN } { #1 } { t1 } { #2 } { t1 }
439
440 }
   \cs_new:Npn \BNVS_new_tl_cn:c {
     \BNVS_new_cn:nc { tl }
443 }
444 \cs_new:Npn \BNVS_new_tl_cv:c #1 {
     \BNVS_new_cv:ncn { tl } { #1 } { tl }
```

```
446 }
                               447 \BNVS_new_tl_cn:c { set }
                               448 \BNVS_new_tl_cv:c { set }
                               449 \BNVS_new:cpn { tl_set:cx } {
                                    \exp_args:Nnx \__bnvs_tl_set:cn
                               451 }
                               452 \BNVS_new_tl_cn:c { put_right }
                               453 \BNVS_new_tl_cv:c { put_right }
                               454 % \BNVS_generate_variant:cn { tl_put_right:cn } { cx }
                               455 \BNVS_new:cpn { tl_put_right:cx } {
                                    \exp_args:Nnnx \BNVS_use:c { tl_put_right:cn }
                               457
                               458 \BNVS_new_tl_cn:c { put_left }
                               459 \BNVS_new_tl_cv:c { put_left }
                               460 % \BNVS_generate_variant:cn { tl_put_left:cn } { cx }
                               461 \BNVS_new:cpn { tl_put_left:cx } {
                                    \exp_args:Nnnx \BNVS_use:c { tl_put_left:cn }
                               462
                               463 }
                             \verb|\__bnvs_tl_if_empty:cTF  | \{\langle core \rangle\} | \{\langle yes | code \rangle\} | \{\langle no | code \rangle\}|
\__bnvs_tl_if_empty:cTF
                             \verb|\ \_bnvs_tl_if_blank:vTF  | \{\langle core \rangle\} | \{\langle yes | code \rangle\} | \{\langle no | code \rangle\}|
\__bnvs_tl_if_blank:vTF
                             \ \ \_bnvs\_tl\_if\_eq:cnTF \ \{\langle core \rangle\} \ \{\langle tl \rangle\} \ \{\langle yes\ code \rangle\} \ \{\langle no\ code \rangle\}
\__bnvs_tl_if_eq:cn\overline{\mathit{TF}}
                              These are shortcuts to
                                 • \t = \frac{1_{\text{bnvs}}(core)_{\text{tl}}}{(yes\ code)} {(no\ code)}
                                 • tl_if_eq:cnTF \{l_bnvs_(core)_tl\}\{(tl)\} \{(yes\ code)\} \{(no\ code)\}\}
                               464 \cs_new:Npn \BNVS_new_conditional_c:ncNn #1 #2 #3 #4 {
                                    \BNVS_new_conditional:cpnn { #2 } ##1 { #4 } {
                                       \BNVS_use:Ncn #3 { ##1 } { #1 } {
                                         \prg_return_true:
                                       } {
                               469
                                         \prg_return_false:
                                       }
                               470
                                    }
                               471
                               472 }
                                  \cs_new:Npn \BNVS_new_conditional_c:ncn #1 #2 {
                               473
                                    \BNVS_use_raw:nc {
                               474
                                       \BNVS_new_conditional_c:ncNn { #1 } { #1_#2:c }
                               475
                                    } { #1_#2:NTF }
                               476
                               478 \BNVS_new_conditional_c:ncn { tl } { if_empty } { p, T, F, TF }
                                  \BNVS_new_conditional:cpnn { tl_if_blank:v } #1 { T, F, TF } {
                                    \BNVS_tl_use:Nv \tl_if_blank:nTF { #1 } {
                               480
                                       \prg_return_true:
                               481
                                    } {
                               482
                                       \prg_return_false:
                               483
                               484
                               485 }
                                  \cs_new:Npn \BNVS_new_conditional_cn:ncNn #1 #2 #3 #4 {
                                    \BNVS_new_conditional:cpnn { #2:cn } ##1 ##2 { #4 } {
                                       \BNVS_use:Ncn #3 { ##1 } { ##2 } {
```

```
491
                                  \prg_return_false:
                               }
                        492
                        493
                        494 }
                           \cs_new:Npn \BNVS_new_conditional_cn:ncn #1 #2 {
                        495
                             \BNVS_use_raw:nc {
                               \BNVS_new_conditional_cn:ncNn { #1 } { #1_#2 }
                             } { #1_#2:NnTF }
                        498
                        499 }
                           \BNVS_new_conditional_cn:ncn { tl } { if_eq } { T, F, TF }
                        500
                           \cs_new:Npn \BNVS_new_conditional_cv:ncNn #1 #2 #3 #4 {
                        501
                             \BNVS_new_conditional:cpnn { #2:cv } ##1 ##2 { #4 } {
                        502
                               \BNVS_use:nvn {
                        503
                                  \BNVS_use:Ncn #3 { ##1 } { #1 }
                        504
                               } { ##2 } { #1 } {
                        505
                                  \prg_return_true:
                        506
                               } {
                                  \prg_return_false:
                               }
                        509
                             }
                        510
                        511 }
                           \cs_new:Npn \BNVS_new_conditional_cv:ncn #1 #2 {
                        512
                             \BNVS_use_raw:nc {
                        513
                        514
                               \BNVS_new_conditional_cv:ncNn { #1 } { #1_#2 }
                             } { #1_#2:NnTF }
                        515
                        516 }
                        517 \BNVS_new_conditional_cv:ncn { tl } { if_eq } { T, F, TF }
                       6.4.3 Strings
_bnvs_str_if_eq:vnTF
                       These are shortcuts to
                          • tl_if_empty:cTF \{l_bnvs_\langle core\rangle_tl\}\{\langle yes\ code\rangle\} \{\langle no\ code\rangle\}
                           \cs_new:Npn \BNVS_new_conditional_vn:ncNn #1 #2 #3 #4 {
                        518
                        519
                             \BNVS_new_conditional:cpnn { #2:vn } ##1 ##2 { #4 } {
                               \BNVS_use:Nvn #3 { ##1 } { #1 } { ##2 } {
                        520
                                  \prg_return_true:
                        521
                               } {
                                  \prg_return_false:
                               }
                        524
                             }
                        525
                        526 }
                           \cs_new:Npn \BNVS_new_conditional_vn:ncn #1 #2 {
                        527
                             \BNVS_use_raw:nc {
                        528
                               \BNVS_new_conditional_vn:ncNn { #1 } { #1_#2 }
                        529
                             } { #1_#2:nnTF }
                        530
                        531 }
                        _{\rm 532} \BNVS_new_conditional_vn:ncn { str } { if_eq } { T, F, TF }
```

489

490

} {

\prg\_return\_true:

```
\cs_new:Npn \BNVS_new_conditional_vv:ncNn #1 #2 #3 #4 {
     \BNVS_new_conditional:cpnn { #2:vv } ##1 ##2 { #4 } {
534
       \BNVS_use:nvn {
535
         \BNVS_use:Nvn #3 { ##1 } { #1 }
536
         { ##2 } { #1 } {
537
         \prg_return_true:
538
       } {
539
         \prg_return_false:
       }
541
     }
542
543 }
   \cs_new:Npn \BNVS_new_conditional_vv:ncn #1 #2 {
     \BNVS_use_raw:nc {
545
       \BNVS_new_conditional_vv:ncNn { #1 } { #1_#2 }
546
     } { #1_#2:nnTF }
547
548 }
549 \BNVS_new_conditional_vv:ncn { str } { if_eq } { T, F, TF }
```

### 6.4.4 Sequences

```
\__bnvs_seq_count:c
                                           \ \ \_bnvs\_seq\_new:c \ \{\langle core \rangle\}
\__bnvs_seq_clear:c
                                           \ \ \_bnvs\_seq\_count:c \ \{\langle core \rangle\}
\__bnvs_seq_set_eq:cc
                                           \ \ \_bnvs\_seq\_clear:c \ \{\langle core \rangle\}
\__bnvs_seq_use:cn
                                           \ __bnvs_seq_set_eq:cc \{\langle core_1 \rangle\} \{\langle core_2 \rangle\}
                                           \verb|\| \_bnvs\_seq\_use:cn | \{\langle core \rangle\} | \{\langle separator \rangle\}|
\__bnvs_seq_item:cn
                                           \__bnvs_seq_remove_all:cn
                                           \verb|\ \_bnvs_seq_remove_all:cn  | \{\langle \mathit{core}\rangle\} | \{\langle \mathit{tl}\rangle\}|
\__bnvs_seq_put_left:cv
                                           \verb|\ \_bnvs_seq_put_right:cn {$\langle seq\ core\rangle$} {\langle tl\rangle}$
\__bnvs_seq_put_right:cn
                                           \__bnvs_seq_put_right:cv
                                           \verb|\ \_bnvs\_seq\_set\_split:cnn | \{\langle seq\ core \rangle\} | \{\langle tl \rangle\} | \{\langle separator \rangle\}|
\__bnvs_seq_set_split:cnn
\__bnvs_seq_set_split:(cnv|cnx)
                                          \ __bnvs_seq_pop_left:cc \{\langle core_1 \rangle\} \{\langle core_2 \rangle\}
\__bnvs_seq_pop_left:cc
```

These are shortcuts to

- \seq\_set\_eq:cc  $\{1\_bnvs\_\langle core_1\rangle\_seq\}$   $\{1\_bnvs\_\langle core_2\rangle\_seq\}$  \seq\_count:c  $\{1\_bnvs\_\langle core\rangle\_seq\}$  \seq\_use:cn  $\{1\_bnvs\_\langle core\rangle\_seq\}$  {\langle separator\rangle} \seq\_item:cn  $\{1\_bnvs\_\langle core\rangle\_seq\}$  {\langle integer expression\rangle} \seq\_remove\_all:cn  $\{1\_bnvs\_\langle core\rangle\_seq\}$  {\langle tl\rangle}
- \\_bnvs\_seq\_clear:c {l\_\_bnvs\_\(core\)\_seq}
- \seq\_put\_right:cv {l\_\_bnvs\_ $\langle seq\ core \rangle$ \_seq} {l\_\_bnvs\_ $\langle tl\ core \rangle$ \_tl}
- \seq\_set\_split:cnn{l\_\_bnvs\_ $\langle seq\ core \rangle$ \_seq}{l\_\_bnvs\_ $\langle tl\ core \rangle$ \_tl}{ $\langle tl \rangle$ }

```
550 \BNVS_new_c:nc { seq } { count }
551 \BNVS_new_c:nc { seq } { clear }
552 \BNVS_new_cn:nc { seq } { use }
```

```
554 \BNVS_new_cn:nc { seq } { remove_all }
                              555 \BNVS_new_cn:nc { seq } { map_inline }
                              556 \BNVS_new_cc:nc { seq } { set_eq }
                              557 \BNVS_new_cv:ncn { seq } { put_left } { tl }
                              558 \BNVS_new_cn:ncn { seq } { put_right } { tl }
                              559 \BNVS_new_cv:ncn { seq } { put_right } { tl }
                              560 \BNVS_new_cnn:nc { seq } { set_split }
                              561 \BNVS_new_cnv:nc { seq } { set_split }
                              562 \BNVS_new_cnx:nc { seq } { set_split }
                              563 \BNVS_new_cc:ncn { seq } { pop_left } { tl }
                               564 \BNVS_new_cc:ncn { seq } { pop_right } { tl }
                              \verb|\ \_bnvs\_seq\_if\_empty:cTF \ \{\langle seq\ core\ name\rangle\}\ \{\langle yes\ code\rangle\}\ \{\langle no\ code\rangle\}
  _bnvs_seq_if_empty:cTF
\__bnvs_seq_get_right:cc<u>TF</u>
                             code\rangle}
\__bnvs_seq_pop_left:cc<u>TF</u>
\_\_bnvs_seq_pop_right:cc{\it TF}
                               565 \cs_new:Npn \BNVS_new_conditional_cc:ncnn #1 #2 #3 #4 {
                                    \BNVS_new_conditional:cpnn { #1_#2:cc } ##1 ##2 { #4 } {
                                      \BNVS_use:ncncn {
                               567
                                        \verb|\BNVS_use_raw:c { #1_#2:NNTF }|
                               568
                                      } { ##1 } { #1 } { ##2 } { #3 } {
                               569
                                        \prg_return_true:
                               570
                               571
                                      } {
                               572
                                        \prg_return_false:
                                      }
                               573
                                    }
                               574
                              575 }
                              576 \BNVS_new_conditional_c:ncn { seq } { if_empty } { T, F, TF }
                              {\tt 577} \verb|\BNVS_new_conditional_cc:ncnn|\\
                                    { seq } { get_right } { tl } { T, F, TF }
                              579 \BNVS_new_conditional_cc:ncnn
                                    { seq } { pop_left } { tl } { T, F, TF }
                              581 \BNVS_new_conditional_cc:ncnn
                                   { seq } { pop_right } { tl } { T, F, TF }
```

553 \BNVS\_new\_cn:nc { seq } { item }

### 6.4.5 Integers

```
\__bnvs_int_new:c
                         \ __bnvs_int_new:c
                                                   \{\langle core \rangle\}
                         \ __bnvs_int_use:c
                                                   \{\langle core \rangle\}
\_\_bnvs_int_use:c
                         \ __bnvs_int_incr:c \{\langle core \rangle\}
\_\_bnvs\_int\_inc:c
                         \ \ \_bnvs\_int\_decr:c \ \{\langle core \rangle\}
\__bnvs_int_decr:c
                         \ __bnvs_int_set:cn \{\langle core \rangle\}\ \{\langle value \rangle\}
\__bnvs_int_set:cn
 __bnvs_int_set:cv
                        These are shortcuts to
                            • \int_new:c
                                                 \{l\_bnvs\_\langle core \rangle\_int\}
                                                 \{l\_bnvs\_\langle core \rangle\_int\}
                             • \int_use:c
                             • \int_incr:c {l_bnvs_\langle core \rangle_int}
                             • \int_idecr:c {l__bnvs_\( core \)_int}
                             • \int_set:cn \{l\_bnvs\_\langle core \rangle\_int\} \langle value \rangle
                          583 \BNVS_new_c:nc
                                                  { int } { new }
                          584 \BNVS_new_c:nc
                                                 { int } { use
                                                 { int } { zero }
                          585 \BNVS_new_c:nc
                          586 \BNVS_new_c:nc
                                                 { int } { incr }
                          587 \BNVS_new_c:nc { int } { decr }
                          588 \BNVS_new_cn:nc { int } { set }
                          589 \BNVS_new_cv:ncn { int } { set } { int }
```

### 6.4.6 Prop

```
\__bnvs_prop_get:NncTF
```

```
590 \BNVS_new_conditional:cpnn { prop_get:Nnc } #1 #2 #3 { T, F, TF } {
591   \BNVS_use:ncn {
592   \prop_get:NnNTF #1 { #2 }
593      } { #3 } { tl } {
594      \prg_return_true:
595      } {
596      \prg_return_false:
597      }
598 }
```

### 6.5 Debug facilities

Typesetting file beanoves.dtx creates both beanoves and beanoves-debug style files. The former is intended for everyday use whereas the latter contains supplemental debugging and testing facilities which are intentionally left undocumented. In particular, we have aliases for \group\_begin: and \group\_end: to allow the display of supplemental informations while debugging.

- 6.6 Debug messages
- 6.7 Variable facilities
- 6.8 Testing facilities

### 6.9 Local variables

We make heavy use of local variables and function scopes. Many functions are executed within a TeX group, which ensures no name collision with the caller stack. The number of variables used has not been optimized, nor the TeX groups used. Optimization often goes against readability.

```
599 \tl_new:N \l__bnvs_id_last_tl
600 \tl_set:Nn \l__bnvs_id_last_tl { ?! }
601 \tl_new:N \l__bnvs_a_tl
602 \tl_new:N \l__bnvs_b_tl
603 \tl_new:N \l__bnvs_c_tl
604 \tl_new:N \l__bnvs_V_tl
605 \tl_new:N \l__bnvs_A_tl
606 \tl_new:N \l__bnvs_L_tl
607 \tl_new:N \l__bnvs_Z_tl
608 \tl_new:N \l__bnvs_ans_tl
609 \tl_new:N \l__bnvs_key_tl
610 \tl_new:N \l__bnvs_key_base_tl
611 \tl_new:N \l__bnvs_id_tl
612 \tl_new:N \l__bnvs_n_tl
613 \tl_new:N \l__bnvs_path_tl
614 \tl_new:N \l__bnvs_group_tl
615 \tl_new:N \l__bnvs_scan_tl
616 \tl_new:N \l__bnvs_query_tl
_{617} \tl_new:N \l__bnvs_token_tl
618 \tl_new:N \l__bnvs_root_tl
619 \tl_new:N \l__bnvs_n_incr_tl
620 \tl_new:N \l__bnvs_incr_tl
621 \tl_new:N \l__bnvs_post_tl
622 \tl_new:N \l__bnvs_suffix_tl
623 \int_new:N \g__bnvs_call_int
624 \int_new:N \l__bnvs_int
625 \seq_new:N \g__bnvs_def_seq
626 \seq_new:N \l__bnvs_a_seq
  \seq_new:N \l__bnvs_b_seq
  \seq_new:N \l__bnvs_ans_seq
  \seq_new:N \l__bnvs_match_seq
630 \seq_new:N \l__bnvs_split_seq
631 \seq_new:N \l__bnvs_path_seq
632 \seq_new:N \l__bnvs_path_base_seq
633 \seq_new:N \l__bnvs_query_seq
634 \seq_new:N \l__bnvs_token_seq
635 \bool_new:N \l__bnvs_in_frame_bool
636 \bool_set_false:N \l__bnvs_in_frame_bool
637 \bool_new:N \l__bnvs_parse_bool
```

In order to implement the provide feature, we add getters and setters

```
638 \bool_new:N \l__bnvs_provide_bool
                                                                 \BNVS_new:cpn { provide_on: } {
                                                                      \bool_set_true:N \l__bnvs_provide_bool
                                                         640
                                                         641 }
                                                                 \BNVS_new:cpn { provide_off: } {
                                                                      \bool_set_false:N \l__bnvs_provide_bool
                                                          645 \__bnvs_provide_off:
                                                       \_ bnvs_if_provide:TF \{\langle yes\ code \rangle\} \{\langle no\ code \rangle\}
_bnvs_if_provide:TF
                                                       Execute \langle yes\ code \rangle when in provide mode, \langle no\ code \rangle otherwise.
                                                                 \BNVS_new_conditional:cpnn { if_provide: } { p, T, F, TF } {
                                                                       \bool_if:NTF \l__bnvs_provide_bool {
                                                         647
                                                                            \prg_return_true:
                                                          648
                                                                      } {
                                                          649
                                                          650
                                                                            \prg_return_false:
                                                                      }
                                                          651
                                                         652 }
                                                                         Infinite loop management
                                                       6.10
                                                       Unending recursivity is managed here.
                                                      Some functions calls, as well as some loop bodies, decrement this counter. When this
     \g__bnvs_call_int
                                                       counter reaches 0, an error is raised or a computation is aborted.
                                                       (End\ definition\ for\ \g_bnvs_call_int.)
                                                         653 \int_const:Nn \c__bnvs_max_call_int { 2048 }
  _bnvs_call_greset:
                                                       \__bnvs_call_greset:
                                                       Reset globally the call stack counter to its maximum value.
                                                         654 \cs_set:Npn \__bnvs_call_greset: {
                                                                      \int_gset:Nn \g__bnvs_call_int { \c__bnvs_max_call_int }
                                                         656 }
          \ \ \sum_{a} bnvs_call: TF
                                                       \mbox{\colored} \mbox{\color
                                                       Decrement the \g_{\text{pnvs_call_int}} counter globally and execute \langle yes\ code\ \rangle if we have
                                                       not reached 0, \langle no\ code \rangle otherwise.
                                                                 \BNVS_new_conditional:cpnn { call: } { T, F, TF } {
                                                                      \int_gdecr:N \g__bnvs_call_int
                                                          658
                                                                      \int_compare:nNnTF \g__bnvs_call_int > 0 {
                                                          659
                                                                            \prg_return_true:
                                                          660
                                                          661
                                                                            \prg_return_false:
                                                          662
```

663 }

## 6.11 Overlay specification

# 6.12 Basic functions

 $\langle key \rangle - \langle value \rangle$  property list to store the named overlay sets. The basic keys are, assuming \g\_\_bnvs\_prop  $\langle id \rangle! \langle key \rangle$  is a fully qualified overlay set name,  $\langle id \rangle! \langle key \rangle / V$  for the value  $\langle id \rangle! \langle key \rangle / A$  for the first index  $\langle id \rangle! \langle key \rangle / L$  for the length when provided  $\langle id \rangle! \langle key \rangle/Z$  for the last index when provided The implementation is private, in particular, keys may change in future versions. 665 \prop\_new:N \g\_\_bnvs\_prop (End definition for  $\g_bnus_prop.$ ) \_bnvs\_gput:nnn  $\_\bnvs_gput:nnn {\langle subkey \rangle} {\langle key \rangle} {\langle value \rangle}$  $\label{locality} $$\sum_{\substack{b \in Subkey}} {\langle key \rangle} $$ \_bnvs\_gput:nnn \\_\_bnvs\_item:nn  $\label{local_subkey} $$\sum_{\text{bnvs\_gremove:nn }} {\langle subkey \rangle} {\langle key \rangle}$$  $\_$ \_bnvs\_gclear:n  $\{\langle key \rangle\}$ \\_\_bnvs\_gremove:nn  $\_\_$ bnvs $\_$ gclear:n \\_\_bnvs\_gclear: \_\_bnvs\_gclear:v Convenient shortcuts to manage the storage, it makes the code more concise and readable.  $\_\_$ bnvs $\_$ gclear: This is a wrapper over IATEX3 eponym functions. The key argument is  $\langle key \rangle / \langle subkey \rangle$ . 666 \BNVS\_new:cpn { gput:nnn } #1 #2 { \prop\_gput:Nnn \g\_\_bnvs\_prop { #2 / #1 } 667 668 } \BNVS\_new:cpn { gput:nnv } #1 #2 { 669 \BNVS\_tl\_use:nv { 670 \\_\_bnvs\_gput:nnn { #1 } { #2 } 671 672 673 }

```
674 \BNVS_new:cpn { item:nn } #1 #2 {
     \prop_item: Nn \g__bnvs_prop { #2 / #1 }
676 }
677 \BNVS_new:cpn { gremove:nn } #1 #2 {
     \prop_gremove:Nn \g__bnvs_prop { #2 / #1 }
678
679 }
680 \BNVS_new:cpn { gclear:n } #1 {
     \clist_map_inline:nn { V, A, Z, L } {
681
       \__bnvs_gremove:nn { ##1 } { #1 }
682
683
     \__bnvs_cache_gclear:n { #1 }
684
685 }
686 \BNVS_new:cpn { gclear: } {
     \prop_gclear:N \g__bnvs_prop
687
688 }
689 \BNVS_generate_variant:cn { gclear:n } { V }
  \BNVS_new:cpn { gclear:v } {
      \BNVS_tl_use:Nc \__bnvs_gclear:V
691
692 }
```

Convenient shortcuts to test for the existence of  $\langle key \rangle / \langle subkey \rangle$ , it makes the code more concise and readable. The version with no  $\langle subkey \rangle$  is the or combination for keys V, A and Z.

```
\BNVS_new_conditional:cpnn { if_in:nn } #1 #2 { p, T, F, TF } {
     \prop_if_in:NnTF \g__bnvs_prop { #2 / #1 } {
       \prg_return_true:
     } {
696
697
       \prg_return_false:
698
699 }
   \BNVS_new_conditional:cpnn { if_in:n } #1 { p, T, F, TF } {
700
     \bool_if:nTF {
701
           \__bnvs_if_in_p:nn V { #1 }
702
       || \__bnvs_if_in_p:nn A { #1 }
703
       || \__bnvs_if_in_p:nn Z { #1 }
     } {
705
706
       \prg_return_true:
     } {
707
708
       \prg_return_false:
709
710 }
   \BNVS_new_conditional:cpnn { if_in:v } #1 { p, T, F, TF } {
711
     \BNVS_tl_use:Nv \__bnvs_if_in:nTF { #1 }
712
       { \prg_return_true: } { \prg_return_false: }
713
714 }
```

\_\_bnvs\_gprovide:nnnT

 $\verb|\__bnvs_gprovide:nnnT| \{\langle subkey\rangle\} \ \{\langle key\rangle\} \ \{\langle value\rangle\} \ \{\langle true\ precode\rangle\}$ 

Execute  $\langle true\ precode \rangle$  before providing, or  $\langle false\ precode \rangle$  before not providing.

```
715 \BNVS_new:cpn { gprovide:nnnT } #1 #2 #3 #4 {
716 \prop_if_in:NnF \g__bnvs_prop { #2 / #1 } {
717 #4
718 \prop_gput:Nnn \g__bnvs_prop { #2 / #1 } { #3 }
719 }
720 }
```

\_\_bnvs\_get:nnc*TF* 

```
\label{localization} $$\sum_{\substack{b \in \mathbb{N} \\ \text{ode}}} {\langle ubkey \rangle} {\langle tl \ core \ name \rangle} {\langle yes \ code \rangle} {\langle no \ code \rangle} $$
```

Convenient shortcuts to retrieve the value with branching, it makes the code more concise and readable. Execute  $\langle yes\ code \rangle$  when the item is found,  $\langle no\ code \rangle$  otherwise. In the latter case, the content of the  $\langle tl\ variable \rangle$  is undefined, on resolution only. NB: the predicate won't work because  $\prop_get:NnNTF$  is not expandable.

```
721 \BNVS_new_conditional:cpnn { get:nnc } #1 #2 #3 { T, F, TF } {
722 \BNVS_tl_use:nc {
723 \prop_get:NnNTF \g_bnvs_prop { #2 / #1 }
724 } { #3 } {
```

```
\prg_return_true:
     } {
726
       \prg_return_false:
728
729 }
   \BNVS_new_conditional:cpnn { get:nvc } #1 #2 #3 { T, F, TF } {
730
     \BNVS_tl_use:nv {
731
       \__bnvs_get:nncTF { #1 }
732
733
     } { #2 } { #3 } {
       \prg_return_true:
734
     } {
735
       \prg_return_false:
736
     }
737
738 }
```

### 6.13 Functions with cache

\g\_\_bnvs\_prop

 $\langle key \rangle - \langle value \rangle$  property list to store the named overlay sets. Other keys are eventually used to cache results when some attributes are defined from other slide ranges.

- (id)!(key)/V for the cached static value of the value
- $\langle id \rangle! \langle key \rangle / A$  for the cached static value of the first index
- $\langle id \rangle! \langle key \rangle / L$  for the cached static value of the length
- $\langle id \rangle! \langle key \rangle / Z$  for the cached static value of the last index
- $\langle id \rangle! \langle key \rangle/P$  for the cached static value of the previous index
- $\langle id \rangle! \langle key \rangle/N$  for the cached static value of the next index

The implementation is private, in particular, keys may change in future versions.

```
739 \prop_new:N \g__bnvs_cache_prop (End definition for \g__bnvs_prop.)
```

Wrapper over the functions above for  $\langle key \rangle / \langle subkey \rangle$ .

```
740 \BNVS_new:cpn { cache_gput:nnn } #1 #2 {
741 \prop_gput:Nnn \g__bnvs_cache_prop { #2 / #1 }
742 }
```

```
\cs_generate_variant:Nn \__bnvs_cache_gput:nnn { nV, nnV }
    \BNVS_new:cpn { cache_gput:nvn } #1 {
      \BNVS_tl_use:nc {
 745
        \__bnvs_cache_gput:nVn { #1 }
 746
 747
 748 }
    \BNVS_new:cpn { cache_gput:nnv } #1 #2 {
 749
      \BNVS_tl_use:nc {
 750
        \__bnvs_cache_gput:nnV { #1 } { #2 }
 752
 753 }
    \BNVS_new:cpn { cache_item:nn } #1 #2 {
 754
      \prop_item:Nn \g_bnvs_cache_prop { #2 / #1 }
 755
 756 }
    \BNVS_new:cpn { cache_gremove:nn } #1 #2 {
 757
      \prop_gremove: Nn \g_bnvs_cache_prop { #2 / #1 }
 758
 759 }
    \BNVS_new:cpn { cache_gclear:n } #1 {
 760
      \clist_map_inline:nn { V, A, Z, L, P, N } { }
        \prop_gremove:Nn \g__bnvs_cache_prop { #1 / ##1 }
 763
 764 }
    \BNVS_new:cpn { cache_gclear: } {
 765
      \prop_gclear:N \g__bnvs_cache_prop
 766
 767 }
\verb|\__bnvs_cache_if_in:nTF {|\langle subkey\rangle| } {|\langle key\rangle| } {|\langle yes\ code\rangle| } {|\langle no\ code\rangle|}
```

\\_\_bnvs\_cache\_if\_in\_p:nn  $\star$ \\_\_bnvs\_cache\_if\_in:nnTF  $\star$ 

Convenient shortcuts to test for the existence of  $\langle subkey \rangle / \langle key \rangle$ , it makes the code more concise and readable.

```
768 \prg_new_conditional:Npnn \__bnvs_cache_if_in:nn #1 #2 { p, T, F, TF } {
769    \prop_if_in:NnTF \g__bnvs_cache_prop { #2 / #1 } {
770    \prg_return_true:
771    } {
772    \prg_return_false:
773    }
774 }
```

 $\_{ ext{\_bnvs\_cache\_get:nnc}}$ 

```
\__bnvs_cache_get:nncTF {\langle subkey \rangle} {\langle tl\ core\ name \rangle} {\langle yes\ code \rangle} {\langle no\ code \rangle}
```

Convenient shortcuts to retrieve the value with branching, it makes the code more concise and readable. Execute  $\langle yes\ code \rangle$  when the item is found,  $\langle no\ code \rangle$  otherwise. In the latter case, the content of the  $\langle tl\ variable \rangle$  is undefined. NB: the predicate won't work because  $\prop_get:NnNTF$  is not expandable.

```
775 \BNVS_new_conditional:cpnn { cache_get:nnc } #1 #2 #3 { p, T, F, TF } {
776 \BNVS_tl_use:nc {
777 \prop_get:NnNTF \g_bnvs_cache_prop { #2 / #1 }
778 } { #3 } {
779 \prg_return_true:
780 } {
```

#### 6.13.1 Implicit value counter

The implicit value counter is local to the current frame. It is defined at the global level because changes made at any depth must be made at the frame depth. If the frame were a closure, this counter would belong to that closure. When used for the first time, it either defaults to the first index or last index.

\g\_\_bnvs\_v\_prop

 $\langle key \rangle - \langle value \rangle$  property list to store the contents or the named value counters. The keys are  $\langle id \rangle! \langle key \rangle$ .

```
784 \prop_new:N \g__bnvs_v_prop
(End definition for \g__bnvs_v_prop.)
```

```
\__bnvs_v_gput:nn
\__bnvs_v_gput:(nV|Vn)
\__bnvs_v_item:n
\__bnvs_v_gremove:n
\__bnvs_v_gclear:
```

```
\label{lem:continuous} $$\sum_{\substack{\ \subseteq \\ \ DNVS_v_{item:n \ \{\langle key \rangle\} \ \\ \ DNVS_v_{gremove:n \ \{\langle key \rangle\} \ \\ \ DNVS_v_{gclear:}}} $$
```

Convenient shortcuts to manage the storage, it makes the code more concise and readable. This is a wrapper over LATEX3 eponym functions.

```
\BNVS_new:cpn { v_gput:nn } {
786
     \prop_gput:Nnn \g_bnvs_v_prop
787 }
788
  \BNVS_new:cpn { v_gput:nv } #1 {
     \BNVS_tl_use:nv {
789
       \__bnvs_v_gput:nn { #1 }
790
791
792 }
   \BNVS_new:cpn { v_item:n } #1 {
793
     \prop_item:Nn \g__bnvs_v_prop { #1 }
794
795 }
   \BNVS_new:cpn { v_gremove:n } {
796
     \prop_gremove:Nn \g_bnvs_v_prop
797
798 }
799 \BNVS_new:cpn { v_gclear: } {
     \prop_gclear:N \g__bnvs_v_prop
801 }
```

```
\__bnvs_v_if_in_p:n *
\__bnvs_v_if_in:n<u>TF</u> *
```

```
\label{local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_local_loc
```

Convenient shortcuts to test for the existence of the  $\langle key \rangle$  value counter.

```
802 \BNVS_new_conditional:cpnn { v_if_in:n } #1 { p, T, F, TF } {
803    \prop_if_in:NnTF \g__bnvs_v_prop { #1 } {
804    \prg_return_true:
805    } {
806    \prg_return_false:
807    }
808 }
```

\\_\_bnvs\_v\_get:nc*TF* 

```
\cline{1.5} \cli
```

Convenient shortcuts to retrieve the value with branching, it makes the code more concise and readable. Execute  $\langle yes\ code \rangle$  when the item is found,  $\langle no\ code \rangle$  otherwise. In the latter case, the content of the  $\langle tl\ variable \rangle$  is undefined. NB: the predicate won't work because  $\prop_get:NnNTF$  is not expandable.

```
809 \BNVS_new_conditional:cpnn { v_get:nc } #1 #2 { T, F, TF } {
810    \BNVS_tl_use:nc {
811    \prop_get:\Nn\NTF \g_bnvs_v_prop { #1 }
812    } { #2 } {
813    \prg_return_true:
814    } {
815    \prg_return_false:
816    }
817 }
```

The key must include the frame id. Reset the value counter to the given  $\langle initial\ value \rangle$ . The \_all version also cleans the cached values. If the  $\langle key \rangle$  is known,  $\langle true\ code \rangle$  is executed, otherwise  $\langle false\ code \rangle$  is executed.

```
\BNVS_new_conditional:cpnn { v_greset:nn } #1 #2 { T, F, TF } {
     \__bnvs_v_if_in:nTF { #1 } {
       \_\_bnvs_v_gremove:n { #1 }
820
821
       \tl_if_empty:nF { #2 } {
         \__bnvs_v_gput:nn { #1 } { #2 }
822
823
824
       \prg_return_true:
     } {
825
       \prg_return_false:
826
827
828 }
   \BNVS_new_conditional:cpnn { v_greset:vn } #1 #2 { T, F, TF } {
829
     \BNVS_tl_use:Nv \__bnvs_v_greset:nnTF { #1 } { #2 }
       { \prg_return_true: } { \prg_return_false: }
831
832 }
   \BNVS_new_conditional:cpnn { greset_all:nn } #1 #2 { T, F, TF } {
833
     \__bnvs_if_in:nTF { #1 } {
834
       \BNVS begin:
835
       \clist_map_inline:nn { V, A, Z, L } {
836
         \_bnvs_get:nncT { ##1 } { #1 } { a } {
837
            \__bnvs_quark_if_nil:cT { a } {
838
             \_bnvs_cache_get:nncTF { ##1 } { #1 } { a } {
839
                \__bnvs_gput:nnv { ##1 } { #1 } { a }
             } {
                \_bnvs_gput:nnn { ##1 } { #1 } { 1 }
842
             }
843
           }
844
         }
845
       }
846
```

```
\BNVS_end:
847
        \__bnvs_cache_gclear:n { #1 }
848
        \__bnvs_v_greset:nnT { #1 } { #2 } {}
849
        \prg_return_true:
850
851
        \prg_return_false:
852
853
854 }
   \BNVS_new_conditional:cpnn { greset_all:vn } #1 #2 { T, F, TF } {
     \label{lem:bnvs_tl_use:Nv } $$\BNVS_tl_use:Nv \__bnvs_greset_all:nnTF { #1 } { #2 } $$
        { \prg_return_true: } { \prg_return_false: }
857
858
```

\\_\_bnvs\_gclear\_all:n
\\_\_bnvs\_gclear\_all:

```
\__bnvs_gclear_all:n \{\langle key \rangle\} \__bnvs_gclear_all:
```

Convenient shortcuts to clear all the storage, for the given key in the first case.

```
\BNVS_new:cpn { gclear_all: } {
     \__bnvs_gclear:
     \__bnvs_cache_gclear:
861
     \__bnvs_n_gclear:
862
     \__bnvs_v_gclear:
863
864 }
  \BNVS_new:cpn { gclear_all:n } #1 {
865
     \__bnvs_gclear:n { #1 }
866
     \_bnvs_cache_gclear:n { #1 }
867
     \__bnvs_n_gremove:n { #1 }
     \__bnvs_v_gremove:n { #1 }
870 }
```

### 6.13.2 Implicit index counter

The implicit index counter is also local to the current frame. It is defined at the global level because changes made at any depth must be made at the frame depth. When used for the first time, it defaults to 1.

\g\_\_bnvs\_n\_prop

 $\langle key \rangle - \langle value \rangle$  property list to store the contents of the named index counters. The keys are  $\langle id \rangle! \langle key \rangle$ .

```
871 \prop_new:N \g__bnvs_n_prop
(End definition for \g__bnvs_n_prop.)
```

```
\__bnvs_n_gput:nn
\__bnvs_n_gput:(nv|vn)
\__bnvs_n_gprovide:nn
\__bnvs_n_item:n
\__bnvs_n_gremove:n
\__bnvs_n_gremove:v
\__bnvs_n_gclear:
```

```
\label{eq:continuous_n_gput:nn} $$ \sum_{\substack{b \in \mathcal{Y} \\ -bnvs_n_{sem} \in \{\langle key \rangle\}}} $$ \sum_{\substack{b \in \mathcal{Y} \\ -bnvs_n_{sem}}} $$ $$ \sum_{\substack{b \in \mathcal{Y} \\ -bnvs_n_{sem}}} $$
```

Convenient shortcuts to manage the storage, it makes the code more concise and readable. This is a wrapper over LATEX3 eponym functions.

```
872 \BNVS_new:cpn { n_gput:nn } {
873   \prop_gput:Nnn \g__bnvs_n_prop
874 }
```

```
\cs_generate_variant:Nn \__bnvs_n_gput:nn { nV }
                          \BNVS_new:cpn { n_gput:nv } #1 {
                            \BNVS_tl_use:nc {
                       877
                              \__bnvs_n_gput:nV { #1 }
                       878
                       879
                       880 }
                           \BNVS_new:cpn { n_gprovide:nn } #1 #2 {
                       881
                            \prop_if_in:NnF \g__bnvs_n_prop { #1 } {
                       882
                              \prop_gput:\nn \g__bnvs_n_prop { #1 } { #2 }
                       884
                       885 }
                          \BNVS_new:cpn { n_item:n } #1 {
                       886
                            \prop_item:Nn \g__bnvs_n_prop { #1 }
                       887
                       888 }
                          \BNVS_new:cpn { n_gremove:n } {
                       889
                            \prop_gremove:Nn \g_bnvs_n_prop
                       890
                       891 }
                          \BNVS_generate_variant:cn { n_gremove:n } { V }
                       892
                          \BNVS_new:cpn { n_gremove:v } {
                            \BNVS_tl_use:nc {
                              \__bnvs_n_gremove:V
                       896
                       897 }
                          \BNVS_new:cpn { n_gclear: } {
                            \prop_gclear:N \g__bnvs_n_prop
                       899
                       900 }
                       901 \cs_generate_variant:\n \__bnvs_n_gremove:n { V }
                      \__bnvs_n_if_in_p:nn \{\langle key \rangle\}
_bnvs_n_if_in_p:n \star
_bnvs_n_if_in:nTF *
                      Convenient shortcuts to test for the existence of the \langle key \rangle value counter.
                          \prg_new_conditional:Npnn \__bnvs_n_if_in:n #1 { p, T, F, TF } {
                            \prop_if_in:NnTF \g__bnvs_n_prop { #1 } {
                       903
                              \prg_return_true:
                       904
                            } {
                       905
                       906
                              \prg_return_false:
                            }
                       907
                       908 }
```

 $\_{ t mget:nc}$ 

 $\verb|\__bnvs_n_get:ncTF| \{\langle key \rangle\} \ \langle \textit{tl variable} \rangle \ \{\langle \textit{yes code} \rangle\} \ \{\langle \textit{no code} \rangle\}$ 

Convenient shortcuts to retrieve the value with branching, it makes the code more concise and readable. Execute  $\langle yes\ code \rangle$  when the item is found,  $\langle no\ code \rangle$  otherwise. In the latter case, the content of the  $\langle tl\ variable \rangle$  is undefined. NB: the predicate won't work because  $\prop_get:NnNTF$  is not expandable.

```
909 \prg_new_conditional:Npnn \__bnvs_n_get:nc #1 #2 { T, F, TF } {
910    \__bnvs_prop_get:NncTF \g__bnvs_n_prop { #1 } { #2 } {
911    \prg_return_true:
912    } {
913    \prg_return_false:
914    }
915 }
```

### 6.13.3 Regular expressions

The name of a slide range consists of a non void list of alphanumerical characters and \c\_\_bnvs\_name\_regex underscore, but with no leading digit. 916 \regex\_const:Nn \c\_\_bnvs\_name\_regex { 917 [[:alpha:]\_][[:alnum:]\_]\* 918 } (End definition for \c\_\_bnvs\_name\_regex.) The name of a slide range consists of a non void list of alphanumerical characters and \c\_\_bnvs\_id\_regex underscore, but with no leading digit. 919 \regex\_const:Nn \c\_\_bnvs\_id\_regex { 920 (?: \ur{c\_bnvs\_name\_regex} | [?] )? ! 921 } (End definition for \c bnvs id regex.) \c\_\_bnvs\_path\_regex A sequence of . (positive integer) items representing a path. 922 \regex\_const:Nn \c\_\_bnvs\_path\_regex { (?: \. \ur{c\_bnvs\_name\_regex} | \. [-+]? \d+ )\*  $(End\ definition\ for\ \c_\_bnvs\_path\_regex.)$ A key is the name of an overlay set possibly followed by a dotted path. Matches the \c\_\_bnvs\_A\_key\_Z\_regex whole string.  $(End\ definition\ for\ \c_bnvs_A_key_Z_regex.)$ 925 \regex\_const:Nn \c\_\_bnvs\_A\_key\_Z\_regex { 1: The range name including the slide  $\langle id \rangle$  and question mark if any 2: slide  $\langle id \rangle$  including the question mark \A ( ( \ur{c\_bnvs\_id\_regex} ? ) \ur{c\_bnvs\_name\_regex} ) 3: the path, if any. (  $\ur\{c\_bnvs\_path\_regex\}$  )  $\Z$ A key is the name of an overlay set possibly followed by a dotted path. Matches the \c\_bnvs\_TEST\_A\_key\_n\_Z\_regex whole string. Catch the ending .n. (End definition for \c\_bnvs\_TEST\_A\_key\_n\_Z\_regex.) 929 \regex\_const:Nn \c\_\_bnvs\_TEST\_A\_key\_n\_Z\_regex { 1: The full match 2: The overlay set name including the slide  $\langle id \rangle$  and question mark if any, the dotted

path but excluding the trailing .n

3: slide  $\langle id \rangle$  including the question mark

```
\A ( ( \ur{c_bnvs_id_regex} ? )
                           930
                                      \ur{c__bnvs_name_regex}
                           931
                                      (?: \. \ur{c_bnvs_name_regex} | \. [-+]? \d+ )*? )
                           932
                            4: the last .n component if any.
                                      (\. n)?\Z
\c_bnvs_colons_regex For ranges defined by a colon syntax.
                           935 \regex_const:Nn \c__bnvs_colons_regex { :(:+)? }
                          (End definition for \c__bnvs_colons_regex.)
                         Used to parse slide list overlay specifications in queries. Next are the 9 capture groups.
\c__bnvs_split_regex
                          Group numbers are 1 based because the regex is used in splitting contexts where only
                         capture groups are considered and not the whole match.
                           936 \regex_const:Nn \c__bnvs_split_regex {
                                \s* ( ? :
                          We start with ++ instrussions<sup>2</sup>.
                                     \+\+
                             • 1: \langle key \rangle of a slide range
                              • 2: \langle id \rangle of a slide range including the exclamation mark
                                   ( ( \ur{c_bnvs_id_regex}? ) \ur{c_bnvs_name_regex} )
                             • 3: optionally followed by a dotted path
                                   ( \ur{c_bnvs_path_regex} )
                             • 4: \langle key \rangle of a slide range
                              • 5: \langle id \rangle of a slide range including the exclamation mark
                                   | ( ( \ur{c_bnvs_id_regex}? ) \ur{c_bnvs_name_regex} )
                              • 6: optionally followed by a dotted path
                                     ( \ur{c_bnvs_path_regex} )
                          We continue with other expressions
                              • 7: the \langle ++n \rangle attribute
```

(?: \.(\+)\+n

943

• 8: the poor man integer expression after '+=', which is the longest sequence of black characters, which ends just before a space or at the very last character. This tricky definition allows quite any algebraic expression, even those involving parenthesis.

<sup>&</sup>lt;sup>2</sup>At the same time an instruction and an expression... this is a synonym of exprection

```
944 | \s* \+= \s* ( \S+ )

• 9: the post increment

945 | (\+)\+

946 | )?

947 | ) \s*
948 }

(End definition for \c_bnvs_split_regex.)
```

### 6.13.4 beamer.cls interface

Work in progress.

```
949 \RequirePackage{keyval}
950 \define@key{beamerframe}{beanoves~id}[]{
     \tl_set:Nx \l__bnvs_id_last_tl { #1 ! }
951
952 }
953 \AddToHook{env/beamer@frameslide/before}{
     \__bnvs_n_gclear:
954
     \__bnvs_v_gclear:
955
     \bool_set_true:N \l__bnvs_in_frame_bool
956
957 }
958 \AddToHook{env/beamer@frameslide/after}{
     \verb|\bool_set_false:N \l|_bnvs_in_frame_bool|
960 }
```

# 6.13.5 Defining named slide ranges

\\_\_bnvs\_range\_set:cccn*TF* 

```
\__bnvs_range_set:cccnTF \langle core\ first \rangle \langle core\ end \rangle \langle core\ length \rangle \{\langle tl \rangle\} \{\langle yes\ code \rangle\} \{\langle no\ code \rangle\}
```

Parse  $\langle tl \rangle$  as a range according to \c\_\_bnvs\_colons\_regex and set the variables accordingly.  $\langle tl \rangle$  is expected to only contain colons and integers.

```
961 \BNVS_new_conditional:cpnn { split_pop_left:c } #1 { T, F, TF } {
     \__bnvs_seq_pop_left:ccTF { split } { #1 } {
962
       \prg_return_true:
963
964
       \prg_return_false:
     }
966
967 }
  \exp_args_generate:n { VVV }
  \BNVS_new_conditional:cpnn { range_set:cccn } #1 #2 #3 #4 { T, F, TF } {
969
     \BNVS_begin:
970
     \__bnvs_tl_clear:c { a }
971
     \__bnvs_tl_clear:c { b }
972
     \__bnvs_tl_clear:c { c }
973
     \__bnvs_regex_split:cnTF { colons } { #4 } {
974
       \__bnvs_seq_pop_left:ccT { split } { a } {
```

```
a may contain the \langle start \rangle.
           \__bnvs_seq_pop_left:ccT { split } { b } {
              \__bnvs_tl_if_empty:cTF { b } {
 977
This is a one colon range.
 978
                \__bnvs_split_pop_left:cTF { b } {
b may contain the \langle end \rangle.
                  \__bnvs_seq_pop_left:ccT { split } { c } {
                    \__bnvs_tl_if_empty:cTF { c } {
 980
A :: was expected:
                       \__bnvs_error:n { Invalid~range~expression(1):~#4 }
 981
                    } {
 982
 983
                       \int_compare:nNnT { \__bnvs_tl_count:c { c } } > { 1 } {
                         \__bnvs_error:n { Invalid~range~expression(2):~#4 }
 984
                       }
 985
                       \__bnvs_split_pop_left:cTF { c } {
 986
\label{local_local_local} $$ l_bnvs_c_tl may contain the $$ (length).
                         \__bnvs_seq_if_empty:cF { split } {
 987
                            \__bnvs_error:n { Invalid~range~expression(3):~#4 }
 988
 989
                       } {
 990
                         \__bnvs_error:n { Internal~error }
 991
                    }
                  }
               } {
 995
                }
 996
             } {
 997
This is a two colon range component.
                \int_compare:nNnT { \__bnvs_tl_count:c { b } } > { 1 } {
 998
                  \__bnvs_error:n { Invalid~range~expression(4):~#4 }
 999
1000
                \__bnvs_seq_pop_left:ccT { split } { c } {
1001
c contains the \langle length \rangle.
                  \__bnvs_split_pop_left:cTF { b } {
1002
                     \__bnvs_tl_if_empty:cTF { b } {
1003
                       \__bnvs_seq_pop_left:cc { split } { b }
1004
b may contain the \langle end \rangle.
                       \__bnvs_seq_if_empty:cF { split } {
1005
                         \__bnvs_error:n { Invalid~range~expression(5):~#4 }
1006
1007
1008
                         _bnvs_error:n { Invalid~range~expression(6):~#4 }
1009
                    }
1010
1011
                    \__bnvs_tl_clear:c { b }
                  }
1013
               }
1014
             }
1015
           }
1016
         }
1017
```

Providing both the  $\langle start \rangle$ ,  $\langle length \rangle$  and  $\langle end \rangle$  of a range is not allowed, even if they happen to be consistent.

```
\cs_set:Npn \BNVS_next: { }
1018
        \_bnvs_tl_if_empty:cT { a } {
1019
          \__bnvs_tl_if_empty:cT { b } {
1020
            \__bnvs_tl_if_empty:cT { c } {
              \cs_set:Npn \BNVS_next: {
                 \__bnvs_error:n { Invalid~range~expression(7):~#3 }
1023
              }
1024
            }
1025
          }
1026
        }
1027
        \BNVS_next:
1028
        \cs_set:Npn \BNVS:nnn ##1 ##2 ##3 {
1029
          \BNVS_end:
1030
          \__bnvs_tl_set:cn { #1 } { ##1 }
          \_bnvs_tl_set:cn { #2 } { ##2 }
          \__bnvs_tl_set:cn { #3 } { ##3 }
1033
1034
        \BNVS_exp_args:Nvvv \BNVS:nnn { a } { b } { c }
1035
        \prg_return_true:
1036
     } {
1037
        \BNVS end:
1038
        \prg_return_false:
1039
     }
1040
1041 }
```

\\_\_bnvs\_range:nnnn \\_\_bnvs\_range:nvvv  $\verb|\__bnvs_range:nnnn| \{\langle key \rangle\} \ \{\langle start \rangle\} \ \{\langle \mathit{end} \rangle\} \ \{\langle \mathit{length} \rangle\}$ 

Auxiliary function called within a group. Setup the model to define a range.

```
\BNVS_new:cpn { range:nnnn } #1 {
1042
      \__bnvs_if_provide:TF {
1043
        \__bnvs_if_in:nnTF A { #1 } {
1044
          \use_none:nnn
1045
1046
           \__bnvs_if_in:nnTF Z { #1 } {
1047
            \use_none:nnn
1049
          } {
             \__bnvs_if_in:nnTF L { #1 } {
1050
               \use_none:nnn
1051
1052
               \__bnvs_do_range:nnnn { #1 }
1053
1054
          }
1055
1056
        {
1057
        \__bnvs_do_range:nnnn { #1 }
   \BNVS_new:cpn { range:nvvv } #1 #2 #3 #4 {
1061
      \BNVS_tl_use:nv {
1062
        \BNVS_tl_use:nv {
1063
          \BNVS_tl_use:nv {
1064
```

```
\BNVS_use:c { range:nnnn } { #1 }
                              1065
                                          } { #2 }
                              1066
                                        } { #3 }
                              1067
                                     } { #4 }
                              1068
                              1069 }
                                         \__bnvs_parse_record:n {\langle full name \rangle}
   _bnvs_parse_record:n
                                          \_\_bnvs\_parse\_record:nn \{\langle full\ name \rangle\}\ \{\langle value \rangle\}
\__bnvs_parse_record:v
                                         \__bnvs_n_parse_record:n {\langle full name \rangle}
\__bnvs_parse_record:nn
\__bnvs_parse_record:(xn|vn)
                                         \label{local_norm} $$\sum_{n\_parse\_record:nn } {\langle full name \rangle} {\langle value \rangle}$
\__bnvs_n_parse_record:n
\__bnvs_n_parse_record:v
\__bnvs_n_parse_record:nn
\__bnvs_n_parse_record:(xn|vn)
```

Auxiliary function for  $\_$  bnvs\_parse:n and  $\_$  bnvs\_parse:nn below. If  $\langle value \rangle$  does not correspond to a range, the V key is used. The \_n variant concerns the index counter. This is a bottleneck.

```
\BNVS_new:cpn { parse_record:n } #1 {
      \__bnvs_if_provide:TF {
        \__bnvs_gprovide:nnnT V { #1 } { 1 } {
1072
1073
          \__bnvs_gclear:n { #1 }
1074
      } {
1075
          _bnvs_gclear:n { #1 }
1076
        \__bnvs_gput:nnn V { #1 } { 1 }
1077
1078
1079 }
    \cs_generate_variant:Nn \__bnvs_parse_record:n { V }
1080
    \BNVS_new:cpn { parse_record:v } {
1081
      \BNVS_tl_use:nc {
        \__bnvs_parse_record:V
      }
1084
1085 }
    \BNVS_new:cpn { parse_record:nn } #1 #2 {
1086
      \__bnvs_range_set:cccnTF { a } { b } { c } { #2 } {
1087
        \__bnvs_range:nvvv { #1 } { a } { b } { c }
1088
1089
        \__bnvs_if_provide:TF {
1090
          \_bnvs_gprovide:nnnT V { #1 } { #2 } {
1091
            \__bnvs_gclear_all:n { #1 }
1092
          }
1093
        } {
1094
          \__bnvs_gclear_all:n { #1 }
1095
          \__bnvs_gput:nnn V { #1 } { #2 }
1096
        }
1097
      }
1098
1099 }
    \cs_generate_variant:Nn \__bnvs_parse_record:nn { x, V }
1100
    \BNVS_new:cpn { parse_record:vn } {
1101
      \BNVS_tl_use:nc {
        \__bnvs_parse_record:Vn
1103
1104
```

```
1105 }
    \BNVS_new:cpn { n_parse_record:n } #1 {
1106
      \bool_if:NTF \l__bnvs_n_provide_bool {
        \__bnvs_n_gprovide:nn
1108
1109
          _bnvs_n_gput:nn
      { #1 } { 1 }
1112
1113 }
    \cs_generate_variant:Nn \__bnvs_n_parse_record:n { V }
1114
    \BNVS_new:cpn { n_parse_record:v } {
1115
      \BNVS_tl_use:nc {
1116
        \__bnvs_n_parse_record:V
1118
1119
    \BNVS_new:cpn { n_parse_record:nn } #1 #2 {
1120
      \__bnvs_range_set:cccnTF { a } { b } { c } { #2 } {
        \__bnvs_error:n { Unexpected~range:~#2 }
1122
1123
          _bnvs_if_provide:TF {
          \__bnvs_n_gprovide:nn { #1 } { #2 }
1125
1126
           __bnvs_n_gput:nn { #1 } { #2 }
1127
        }
1128
     }
1129
1130 }
    \cs_generate_variant:Nn \__bnvs_n_parse_record:nn { x, V }
1131
    \BNVS_new:cpn { n_parse_record:vn } {
      \BNVS_tl_use:Nc \__bnvs_n_parse_record:Vn
1134 }
```

\\_\_bnvs\_name\_id\_n\_get:n<u>TF</u> \\_\_bnvs\_name\_id\_n\_get:v<u>TF</u>  $\verb|\__bnvs_name_id_n_set:nTF {|\langle key \rangle} {| {\langle yes \ code \rangle}} {| {\langle no \ code \rangle}}$ 

If the  $\langle key \rangle$  is a key, put the name it defines into the key tl variable, the frame id in the id tl variable, then execute  $\langle yes\ code \rangle$ . The n tl variable is empty except when  $\langle key \rangle$  ends with .n. Otherwise execute  $\langle no\ code \rangle$ . If  $\langle key \rangle$  does not contain a frame id, then key is prepended with then id\_last and id is set to this value as well.

```
\BNVS_new:cpn { name_id_n_end_export: } {
     \cs_set:Npn \BNVS:nnn ##1 ##2 ##3 {
1136
       \BNVS_end:
        \__bnvs_tl_set:cn { key } { ##1 }
1138
        \__bnvs_tl_set:cn { id } { ##2 }
1139
        \__bnvs_tl_set:cn { n } { ##3 }
1140
        _bnvs_tl_if_empty:cTF { id } {
1142
        \BNVS_exp_args:Nvvv
1143
        \BNVS:nnn { key } { id_last } { n }
1144
        \__bnvs_tl_put_left:cv { key } { id_last }
1145
     } {
1146
        \BNVS_exp_args:Nvvv
        \BNVS:nnn { key } { id } { n }
1148
        \__bnvs_tl_set:cv { id_last } { id }
1149
     }
1150
```

```
1151 }
    \BNVS_new_conditional:cpnn { name_id_n_get:n } #1 { T, F, TF } {
      \BNVS_begin:
        _bnvs_match_once:NnTF \c__bnvs_TEST_A_key_n_Z_regex { #1 } {
1154
        \__bnvs_match_pop_left:cTF { key } {
           \__bnvs_match_pop_left:cTF { key } {
1156
             \__bnvs_match_pop_left:cTF { id } {
               \__bnvs_match_pop_left:cTF { n } {
1158
                 \__bnvs_name_id_n_end_export:
1159
                 \prg_return_true:
1160
              } {
1161
                 \BNVS_end:
1162
                 \__bnvs_error:n { LOGICALLY_UNREACHABLE_A_key_n_Z/n }
1163
                 \prg_return_false:
1164
               }
1165
            } {
1166
               \BNVS_end:
1167
               \__bnvs_error:n { LOGICALLY_UNREACHABLE_A_key_n_Z/id }
1168
               \prg_return_false:
            }
          } {
            \BNVS_end:
1172
            \__bnvs_error:n { LOGICALLY_UNREACHABLE_A_key_n_Z/name }
            \prg_return_false:
1174
1175
        } {
1176
          \BNVS_end:
1177
          \__bnvs_error:n { LOGICALLY_UNREACHABLE_A_key_n_Z/n }
1178
          \prg_return_false:
1179
1180
        }
      } {
1181
        \BNVS_end:
1182
1183
        \prg_return_false:
      }
1184
1185
    \BNVS_new_conditional:cpnn { name_id_n_get:v } #1 { T, F, TF } {
1186
      \BNVS_tl_use:nv { \BNVS_use:c { name_id_n_get:nTF } } { #1 } {
1187
1188
        \prg_return_true:
1189
        \prg_return_false:
1191
      }
1192 }
```

\\_\_bnvs\_parse:n
\\_\_bnvs\_parse:nn

```
\_bnvs_parse:n \{\langle key \rangle\}
\_bnvs_parse:nn \{\langle key \rangle\} \{\langle definition \rangle\}
```

Auxiliary functions called within a group by  $\ensuremath{\mbox{keyval\_parse:nnn.}} \langle key \rangle$  is the overlay reference key, including eventually a dotted path and a frame identifier,  $\langle definition \rangle$  is the corresponding definition.

```
\l__bnvs_match_seq Local storage for the match result.
```

```
(End definition for \l_bnvs_match_seq.)

1193 \BNVS_new:cpn { parse:n } #1 {
```

```
\peek_remove_spaces:n {
1194
        \peek_catcode:NTF \c_group_begin_token {
1195
          \__bnvs_tl_if_empty:cTF { root } {
1196
            \__bnvs_error:n { Unexpected~list~at~top~level. }
1197
1198
          \BNVS_begin:
1199
          \__bnvs_int_incr:c { }
1200
          \__bnvs_tl_set:cx { root } { \__bnvs_int_use:c { } . }
1201
          \cs_set:Npn \bnvs:nw ####1 ####2 \s_stop {
            \regex_match:nnT { \S* } { ####2 } {
               \__bnvs_error:n { Unexpected~####2 }
            }
1205
            \keyval_parse:nnn {
1206
              \__bnvs_parse:n
1207
            } {
1208
              \__bnvs_parse:nn
1209
            } { ####1 }
            \BNVS_end:
1211
         }
          \bnvs:nw
       } {
1214
          \__bnvs_tl_if_empty:cTF { root } {
1215
            \_\bnue_id_n_get:nTF { #1 } {
1216
              \__bnvs_tl_if_empty:cTF { n } {
                \__bnvs_parse_record:v
1218
              } {
1219
                   _bnvs_n_parse_record:v
              }
              { key }
1222
            } {
              \__bnvs_error:n { Unexpected~key:~#1 }
1224
1225
            }
          } {
1226
            \__bnvs_int_incr:c { }
            \__bnvs_tl_if_empty:cTF { n } {
1228
               \__bnvs_parse_record:xn
1229
            } {
1230
              \__bnvs_n_parse_record:xn
1232
            } {
              \__bnvs_tl_use:c { root } . \__bnvs_int_use:c { }
            } { #1 }
          }
1236
          \use_none_delimit_by_s_stop:w
       }
     }
1238
     #1 \s_stop
1239
1240 }
    \BNVS_new:cpn { do_range:nnnn } #1 #2 #3 #4 {
1241
        \__bnvs_gclear_all:n { #1 }
1242
1243
     \tl_if_empty:nTF { #4 } {
        \tl_if_empty:nTF { #2 } {
1245
          \tl_if_empty:nTF { #3 } {
            \__bnvs_error:n { Not~a~range:~:~#1 }
1246
          } {
1247
```

```
\__bnvs_gput:nnn Z { #1 } { #3 }
            \_bnvs_gput:nnn V { #1 } { \q_nil }
1249
1250
       } {
          \_bnvs_gput:nnn A { #1 } { #2 }
1252
          \__bnvs_gput:nnn V { #1 } { \q_nil }
1253
          \tl_if_empty:nF { #3 } {
1254
            1255
            \_\_bnvs\_gput:nnn L { #1 } { q\_nil }
          }
1257
       }
1258
     } {
1259
        \tl_if_empty:nTF { #2 } {
1260
          \__bnvs_gput:nnn L { #1 } { #4 }
1261
          \tl_if_empty:nF { #3 } {
1262
            \__bnvs_gput:nnn Z { #1 } { #3 }
1263
            \__bnvs_gput:nnn A { #1 } { \q_nil }
1264
            \_bnvs_gput:nnn V { #1 } { \q_nil }
1265
          }
       } {
          \__bnvs_gput:nnn A { #1 } { #2 }
          \__bnvs_gput:nnn L { #1 } { #4 }
1269
          \__bnvs_gput:nnn Z { #1 } { \q_nil }
          \__bnvs_gput:nnn V { #1 } { \q_nil }
1272
     }
1273
1274 }
    \cs_new:Npn \BNVS_exp_args:NNcv #1 #2 #3 #4 {
1275
     \BNVS_tl_use:nc { \exp_args:NNnV #1 #2 { #3 } }
1276
1277
        { #4 }
1278 }
   \cs_new:Npn \BNVS_end_tl_set:cv #1 #2 {
     \BNVS_tl_use:nv {
1280
        \BNVS_end: \__bnvs_tl_set:cn { #1 }
1281
     } { #2 }
1282
1283 }
   \BNVS_new:cpn { parse:nn } #1 #2 {
1284
      \BNVS_begin:
1285
1286
      \__bnvs_tl_set:cn { a } { #1 }
      \__bnvs_tl_put_left:cv { a } { root }
      \__bnvs_name_id_n_get:vTF { a } {
        \mbox{regex_match:nnTF { \S } { #2 } {}
1289
1290
          \peek_remove_spaces:n {
            \peek_catcode:NTF \c_group_begin_token {
1291
```

The value is a comma separated list, go recursive. But before we warn about an unexpected .n suffix, if any.

```
\regex_match:nnT { \S } { ##2 } {
1299
                   \__bnvs_error:n { Unexpected~value~#2 }
1300
1301
                 \keyval_parse:nnn {
1302
                    \__bnvs_parse:n
1303
                   {
1304
                   \__bnvs_parse:nn
1305
                 } { ##1 }
1306
                 \BNVS_end:
               }
               \BNVS:nn
             } {
               \__bnvs_tl_if_empty:cTF { n } {
1311
                  \__bnvs_parse_record:vn
1312
                    _bnvs_n_parse_record:vn
               { key } { #2 }
1316
               \use_none_delimit_by_s_stop:w
             }
          }
1319
          #2 \s_stop
        } {
1321
Empty value given: remove the reference.
           \__bnvs_tl_if_empty:cTF { n } {
1322
             \__bnvs_gclear:v
1323
             \__bnvs_n_gremove:v
1325
          }
1326
1327
          { key }
      } {
1329
         \__bnvs_error:n { Invalid~key:~#2 }
1330
1331
We export \l__bnvs_id_last_tl:
      \BNVS_end_tl_set:cv { id_last } { id_last }
1332
1333 }
1334 \BNVS_new:cpn { parse_prepare:N } #1 {
      \tl_set:Nx #1 #1
1335
      \bool_set_false:N \l__bnvs_parse_bool
1336
      \bool_do_until:Nn \l__bnvs_parse_bool {
         \tl_if_in:NnTF #1 {%---[
1338
        ]} {
1339
           \rcell{regex_replace_all:nnNF} { ([^\]%---)}
1340
          ]*%---[(
1341
          ) \] } { { \1 } } #1 {
             \bool_set_true:N \l__bnvs_parse_bool
          }
1344
        } {
1345
           \bool_set_true:N \l__bnvs_parse_bool
1346
        }
1347
      }
1348
```

```
\tl_if_in:NnTF #1 {%---[
1349
      ]} {
1350
         \__bnvs_error:n { Unbalanced~%---[
1351
        ]}
1352
      } {
1353
         \tl_if_in:NnT #1 { [%---]
1354
        } {
1355
           \__bnvs_error:n {    Unbalanced~[ %---]
1356
        }
1358
      }
1359
1360 }
```

\Beanoves

```
\Beanoves \{\langle key--value\ list \rangle\}
```

The keys are the slide overlay references. When no value is provided, it defaults to 1. On the contrary,  $\langle key-value \rangle$  items are parsed by \\_\_bnvs\_parse:nn.

```
\cs_new:Npn \BNVS_end_tl_put_right:cv #1 #2 {
      \BNVS_tl_use:nv {
1362
         \BNVS_end:
1363
         \__bnvs_tl_put_right:cn { #1 }
1364
      } { #2 }
1365
1366 }
    \cs_new:Npn \BNVS_end_v_gput:nc #1 #2 {
1367
      \BNVS_tl_use:nv {
1368
         \BNVS_end:
         \__bnvs_v_gput:nn { #1 }
1370
      } { #2 }
1371
1372
    \NewDocumentCommand \Beanoves { sm } {
1373
      \tl_if_empty:NTF \@currenvir {
1374
We are most certainly in the preamble, record the definitions globally for later use.
         \seq_gput_right: Nn \g_bnvs_def_seq { #2 }
      } {
1376
        \tl_if_eq:NnT \@currenvir { document } {
1377
At the top level, clear everything.
1378
           \__bnvs_gclear:
        }
1379
        \BNVS_begin:
1380
         \__bnvs_tl_clear:c { root }
1381
         \int_zero:N \l__bnvs_int
1382
         \__bnvs_tl_set:cn { a } { #2 }
1383
         \tl_if_eq:NnT \@currenvir { document } {
1384
At the top level, use the global definitions.
           \seq_if_empty:NF \g__bnvs_def_seq {
1385
             \__bnvs_tl_put_left:cx { a } {
               \seq_use:Nn \g__bnvs_def_seq , ,
1387
             }
1388
          }
1389
        }
1390
         \__bnvs_parse_prepare:N \l__bnvs_a_tl
1391
```

```
\IfBooleanTF {#1} {
1392
          \__bnvs_provide_on:
1393
        } {
1394
             _bnvs_provide_off:
1395
1396
        \BNVS_tl_use:nv {
1397
          \keyval_parse:nnn { \__bnvs_parse:n } { \__bnvs_parse:nn }
1398
1399
        \BNVS_end_tl_set:cv { id_last } { id_last }
        \ignorespaces
1403 }
```

If we use the frame beanoves option, we can provide default values to the various name ranges.

1404 \define@key{beamerframe}{beanoves}{\Beanoves\*{#1}}

# 6.13.6 Scanning named overlay specifications

Patch some beamer commands to support ?(...) instructions in overlay specifications.

\beamer@frame \beamer@masterdecode

Preprocess (overlay specification) before beamer reads it.

\l\_\_bnvs\_ans\_tl

Storage for the translated overlay specification, where ?(...) instructions are replaced by their static counterparts.

```
(End\ definition\ for\ \verb|\l_bnvs_ans_tl|)
```

Save the original macro \beamer@masterdecode and then override it to properly preprocess the argument.

```
\cs_set_eq:NN \__bnvs_beamer@frame \beamer@frame
   \cs_set:Npn \beamer@frame < #1 > {
1406
     \BNVS_begin:
      \__bnvs_tl_clear:c { ans }
      \__bnvs_scan:nNc { #1 } \__bnvs_eval:nc { ans }
1410
      \BNVS_tl_use:nv {
        \BNVS_end:
1411
        \__bnvs_beamer@frame <
1412
     } { ans } >
1413
1414 }
   \cs_set_eq:NN \__bnvs_beamer@masterdecode \beamer@masterdecode
1415
   \cs_set:Npn \beamer@masterdecode #1 {
1416
     \BNVS_begin:
1417
      \__bnvs_tl_clear:c { ans }
1418
      \__bnvs_scan:nNc { #1 } \__bnvs_eval:nc { ans }
1419
      \BNVS_tl_use:nv {
1420
        \BNVS_end:
1421
        \__bnvs_beamer@masterdecode
1422
     } { ans }
1423
1424 }
```

#### \\_\_bnvs\_scan:nNc

```
\label{local_norm} $$\sum_{\mathrm{nnc}} {\langle named\ overlay\ expression \rangle} \ \langle eval \rangle \ \langle tl\ core \rangle $$
```

Scan the  $\langle named\ overlay\ expression \rangle$  argument and feed the  $\langle tl\ variable \rangle$  replacing ?(...) instructions by their static counterpart with help from the  $\langle eval \rangle$  function, which is \\_\_bnvs\_eval:nN. A group is created to use local variables:

\l\_\_bnvs\_ans\_tl

The token list that will be appended to  $\langle tl \ variable \rangle$  on return.

 $(End\ definition\ for\ \l_bnvs_ans_tl.)$ 

\l\_\_bnvs\_in

Store the depth level in parenthesis grouping used when finding the proper closing parenthesis balancing the opening parenthesis that follows immediately a question mark in a ?(...) instruction.

(End definition for \l\_\_bnvs\_int.)

\l\_\_bnvs\_query\_tl

Storage for the overlay query expression to be evaluated.

 $(End\ definition\ for\ \l_bnvs_query_tl.)$ 

\l\_\_bnvs\_token\_seq

The  $\langle overlay \ expression \rangle$  is split into the sequence of its tokens.

\l\_\_bnvs\_token\_tl

Storage for just one token.

(End definition for  $\l_bnvs_token_tl$ .) Next are helpers.

### \_\_bnvs\_scan\_question:T

```
\_\begin{tabular}{ll} \ \cline{-100} \ \cline{-10
```

At top level state, scan the tokens of the  $\langle named\ overlay\ expression \rangle$  looking for a '?' character. If a '?(...)' is found, then the  $\langle code \rangle$  is executed.

```
\BNVS_new:cpn { scan_question:T } #1 {
       __bnvs_seq_pop_left:ccT { token } { token } {
        \__bnvs_tl_if_eq:cnTF { token } { ? } {
          \__bnvs_scan_require_open:
          #1
1429
       } {
1430
           \__bnvs_tl_put_right:cv { ans } { token }
1431
1432
1433
          _bnvs_scan_question:T { #1 }
1434
1435
1436 }
```

### \_\_bnvs\_scan\_require\_open:

### \require\_open:

We just found a '?', we first gobble tokens until the next '(', whatever they may be. In general, no tokens should be silently ignored.

```
1437 \BNVS_new:cpn { scan_require_open: } {
```

Get next token.

```
\__bnvs_seq_pop_left:ccTF { token } { token } {
\tau_{1439} \tau_{tl_if_eq:NnTF \l__bnvs_token_tl } { ( %)
\tau_{1440} \} {
\text{We found the '(' after the '?'. Set the parenthesis depth to 1 (on first passage).}
\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\t
```

Record the forthcomming content in the \l\_\_bnvs\_query\_tl variable, up to the nex balancing ')'.

Ignore this token and loop.

```
1445 \__bnvs_scan_require_open:
1446 }
1447 } {
```

End reached but no opening parenthesis found, raise.

\\_\_bnvs\_scan\_require\_close: \require\_close:

We found a '?(', we record the forthcomming content in the query variable, up to the next balancing ')'.

```
1452 \BNVS_new:cpn { scan_require_close: } {
```

Get next token.

We found a '(', increment the depth and append the token to query, then scan again for a).

This is not a '('.

```
1460 \__bnvs_tl_if_eq:cnTF { token } { %(---
1461 )
1462 } {
```

We found a balancing ')', we decrement and test the depth.

```
\_bnvs_int_decr:c {}

int_compare:nNnTF { \_bnvs_int_use:c {} } = 0 {
```

The depth level has reached 0: we found our balancing parenthesis of the ?(...) instruction. We can append the evaluated slide ranges token list to ans and look for the next ?.

```
1465 } {
```

The depth has not yet reached level 0. We append the ')' to query because it is not yet the end of sequence marker.

The scanned token is not a '(' nor a ')', we append it as is to query and look for a balancing).

Above ends the code for Not a '('. We reached the end of the sequence and the token list with no closing ')'. We raise and terminate. As recovery we feed query with the missing ')'.

```
\__bnvs_error:x { Missing~%(---
         `)'}
1477
       \__bnvs_tl_put_right:cx { query } {
         \prg_replicate:nn { \l_bnvs_int } {\%(---)}
1478
1479
         )}
1480
     }
1481
1482 }
   \BNVS_new:cpn { scan:nNc } #1 #2 #3 {
     \BNVS_begin:
     \BNVS_set:cpn { fatal:x } ##1 {
       \msg_fatal:nnx { beanoves } { :n }
1486
         { \tl_to_str:n { #1 }:~##1}
1487
1488
     \BNVS_set:cpn { error:x } ##1 {
1489
       \msg_error:nnx { beanoves } { :n }
1490
         { \tl_to_str:n { #1 }:~##1}
1491
1492
1493
        \__bnvs_tl_clear:c { ans }
     \__bnvs_seq_clear:c { token }
```

Explode the  $\langle named\ overlay\ expression \rangle$  into a list of individual tokens:

```
\regex_split:nnN { } { #1 } \l__bnvs_token_seq
```

Run the top level loop to scan for a '?' character:

#### 6.13.7 Resolution

Given a name, a frame id and an integer path, we resolve any intermediate standalone reference. For example, with A=B and B=C, A is resolved in C. But with A=B+1 and B=C, A is not resolved in C+1. With A=B:D and B=C, A is not resolved in C:D as well.

\\_\_bnvs\_kip:cccTF

```
\c \sum_{kip:cccTF} \{\langle key \rangle\} \{\langle id \rangle\} \{\langle path \rangle\} \{\langle yes\ code \rangle\} \{\langle no\ code \rangle\}
```

Auxiliary function. On input, the  $\langle key \rangle$  tl variable contains a set name whereas the  $\langle id \rangle$  tl variable contains a frame id. If  $\langle key \rangle$  tl variable contents is a recorded key, on return,  $\langle key \rangle$  tl variable contains the resolved name,  $\langle id \rangle$  tl variable contains the used frame id,  $\langle path \rangle$  seq variable is prepended with new dotted path components,  $\langle yes\ code \rangle$  is executed, otherwise  $\langle no\ code \rangle$  is executed.

```
\exp_args_generate:n { VVx }
    \quark_new:N \q__bnvs
   \BNVS_new:cpn { end_kip_export_seq:nnnccc } #1 #2 #3 #4 #5 #6 {
      \BNVS_end:
      \tl_if_empty:nTF { #2 } {
1509
        \__bnvs_tl_set:cn { #4 } { #1 }
1510
        \__bnvs_tl_put_left:cv { #4 } { #5 }
1511
1512
          _bnvs_tl_set:cn { #4 } { #1 }
1513
        \__bnvs_tl_set:cn { #5 } { #2 }
1514
1515
      \_bnvs_seq_set_split:cnn { #6 } { \q_bnvs } { #3 }
1516
      \__bnvs_seq_remove_all:cn { #6 } { }
1517
1518 }
    \BNVS_new:cpn { end_kip_export:ccc } {
1519
1520
      \exp_args:Nnnx \BNVS_tl_use:nv {
        \BNVS_tl_use:Nv \__bnvs_end_kip_export_seq:nnnccc { key }
1521
     } { id } {
1522
        \_bnvs_seq_use:cn { path } { \q_bnvs }
1523
1524
1525 }
    \BNVS_new_conditional:cpnn { match_pop_kip: } { T, F, TF } {
1526
      \_bnvs_match_pop_left:cTF { key } {
1527
        \__bnvs_match_pop_left:cTF { key } {
1528
          \__bnvs_match_pop_left:cTF { id } {
            \_bnvs_match_pop_left:cTF { path } {
1530
               \__bnvs_seq_set_split:cnv { path } { . } { path }
1531
              \__bnvs_seq_remove_all:cn { path } { }
1532
              \prg_return_true:
1533
            }
1534
               \prg_return_false:
1535
1536
          }
1537
             \prg_return_false:
1538
          }
        } {
1541
          \prg_return_false:
        }
1542
     } {
1543
        \prg_return_false:
1544
1545
```

```
1546 }
                                 \BNVS_new_conditional:cpnn { kip:ccc } #1 #2 #3 { T, F, TF } {
                            1547
                                   \BNVS begin:
                            1548
                                   \__bnvs_match_once:NvTF \c__bnvs_A_key_Z_regex { #1 } {
                            1549
                            This is a correct key, update the path sequence accordingly.
                                         _bnvs_match_pop_kip:TF {
                            1550
                                         \__bnvs_end_kip_export:ccc { #1 } { #2 } { #3 }
                            1551
                                         \prg_return_true:
                             1552
                                      }
                                         \BNVS_end:
                             1554
                                         \prs_return_false:
                             1556
                                      {
                             1557
                                      \BNVS_end:
                            1558
                                      \prg_return_false:
                            1559
                            1560
                            1561 }
  _bnvs_kip_n_path_resolve:TF
                                       \mbox{\line bnvs\_kip\_n\_path\_resolve:TF } {\langle yes\ code \rangle} {\langle no\ code \rangle}
\__bnvs_kip_x_path_resolve:{\it TF}
                                       \mbox{\line bnvs\_kip\_x\_path\_resolve:TF } {\langle yes\ code \rangle} {\langle no\ code \rangle}
```

{\langle yes code}} will be executed once resolution has occurred, {\langle no code}} otherwise. The key and id variables as well as the path sequence are meant to contain proper information on input and on output as well. On input, \l\_\_bnvs\_key\_tl contains a slide range name, \l\_\_bnvs\_id\_tl contains a frame id and \l\_\_bnvs\_path\_seq contains the components of an integer path, possibly empty. On return, the variable \l\_\_bnvs\_key\_tl contains the resolved range name, \l\_\_bnvs\_id\_tl contains the frame id used and \l\_\_bnvs\_path\_seq contains the sequence of integer path components that could not be resolved.

To resolve one level of a named one slide specification like  $\langle qualified\ name \rangle.\langle i_1 \rangle...\langle i_n \rangle$ , we replace the shortest  $\langle qualified\ name \rangle.\langle i_1 \rangle...\langle i_k \rangle$  where  $0 \le k \le n$  by its definition  $\langle qualified\ name' \rangle.\langle j_1 \rangle...\langle j_p \rangle$  if any. The \\_\_bnvs\_resolve\_?:NNNTF function uses this one level resolution as many times as possible, but no more than a predefined limit to catch circular reference that would lead to an infinite loop.

- 1. If \l\_\_bnvs\_key\_tl content is the name of an unlimited range, and the first item of this range is exactly another name range with eventually a heading frame identifier or a trailing integer path, then \l\_\_bnvs\_key\_tl is replaced by this name, the \l\_\_bnvs\_id\_tl and \l\_\_bnvs\_id\_tl are updates accordingly and the \langle path seq var\rangle is prepended with the integer path.
- 2. If  $\langle path \ seq \ var \rangle$  is not empty, append to the right of \l\_\_bnvs\_key\_tl after a separating dot, all its left elements but the last one and loop. Otherwise return.

In the \_n variant, the resolution is driven only when there is a non empty dotted path.

In the  $\underline{\ x \ }$  variant, the resolution is driven one step further: if  $\langle path \ seq \ var \rangle$  is empty,  $\langle name \ tl \ var \rangle$  can contain anything, including an integer for example.

```
\c = bnvs_{ip} = accepted by the content of the c
```

```
1562 \BNVS_new:cpn { kip_x_path_resolve:TFF } #1 #2 {
1563    \__bnvs_kip_x_path_resolve:TF {
1564    \__bnvs_seq_if_empty:cTF { path } { #1 } { #2 }
1565    }
1566 }
```

# Local variables:

- \l\_bnvs\_a\_tl contains the name with a partial index path currently resolved.
- \l\_\_bnvs\_a\_seq contains the index path components currently resolved.
- \l\_bnvs\_b\_tl contains the resolution.
- \l\_bnvs\_b\_seq contains the index path components to be resolved.

```
\BNVS_new:cpn { end_kip_export: } {
1567
      \exp_args:Nnnx
      \BNVS_tl_use:nv {
        \BNVS_tl_use:Nv \__bnvs_end_kip_export_seq:nnnccc { key }
1570
     } { id } {
1571
          _bnvs_seq_use:cn { path } { \q_bnvs }
1572
     } { key } { id } { path }
1573
1574
    \BNVS_new:cpn { seq_merge:cc } #1 #2 {
1575
      \ \ \sum_{\text{bnvs\_seq\_if\_empty:cF}} \{ \#2 \} 
1576
        \__bnvs_seq_set_split:cnx { #1 } { \q__bnvs } {
1577
          \__bnvs_seq_use:cn { #1 } { \q__bnvs }
1578
          \exp_not:n { \q_bnvs }
          \_bnvs_seq_use:cn { #2 } { \q_bnvs }
1581
         __bnvs_seq_remove_all:cn { #1 } { }
1582
     }
1583
1584
   }
    \BNVS_new:cpn { kip_x_path_resolve:nFF } #1 #2 #3 {
1585
      \_bnvs_get:nvcTF #1 { a } { b } {
1586
        \__bnvs_kip:cccTF { b } { id } { path } {
1587
          \__bnvs_tl_set_eq:cc { key } { b }
1588
          \__bnvs_seq_merge:cc { path } { b }
          \__bnvs_seq_clear:c { b }
          \__bnvs_seq_set_eq:cc { a } { path }
1591
            __bnvs_kip_x_path_resolve_loop_or_end_return:
1592
1593
          \__bnvs_seq_if_empty:cTF { b } {
1594
            \_bnvs_tl_set_eq:cc { key } { b }
1595
            \__bnvs_seq_clear:c { path }
1596
            \__bnvs_seq_clear:c { a }
1597
            \__bnvs_kip_x_path_resolve_loop_or_end_return:
1598
          } {
            #2
1600
          }
        }
1602
     } {
1603
        #3
1604
     }
1605
1606 }
```

```
\BNVS_new:cpn { kip_x_path_resolve_VAL_loop_or_end_return:F } #1 {
     \__bnvs_kip_x_path_resolve:nFF V { #1 } {
1608
          _bnvs_kip_x_path_resolve:nFF A { #1 } {
1609
          1610
1611
     }
1612
1613 }
    \BNVS_new:cpn {    kip_x_path_resolve_end_return_true:    } {
1614
     \__bnvs_seq_pop_left:ccTF { path } { a } {
        \_bnvs_seq_if_empty:cTF { path } {
          \__bnvs_tl_clear:c { b }
1617
          \__bnvs_index_can:vTF { key } {
1618
            \__bnvs_index_append:vvcTF { key } { a } { b } {
1619
              \__bnvs_tl_set:cv { key } { b }
1620
            }
1621
              \__bnvs_tl_set:cv { key } { a }
1622
1623
         }
1624
            \_\_bnvs\_tl\_set:cv { key } { a }
         3
       } {
          \__bnvs_error:x { Path~too~long~.\BNVS_tl_use:c { a }
1628
            .\_bnvs_seq_use:cn { path } . }
1629
       }
1630
     }
       {
1631
          _bnvs_value_resolve:vcT { key } { key } {}
1632
1633
1634
     \__bnvs_end_kip_export:
     \prg_return_true:
1635
1636 }
   \BNVS_new_conditional:cpnn { kip_x_path_resolve: } { T, F, TF } {
1637
1638
     \BNVS_begin:
1639
     \__bnvs_seq_set_eq:cc { a } { path }
     \__bnvs_seq_clear:c { b }
1640
     \__bnvs_kip_x_path_resolve_loop_or_end_return:
1641
1642 }
   \BNVS_new:cpn { kip_x_path_resolve_loop_or_end_return: } {
1643
      \__bnvs_call:TF {
1644
1645
        \_\_bnvs_tl_set_eq:cc { a } { key }
          _bnvs_seq_if_empty:cTF { a } {
            _bnvs_kip_x_path_resolve_VAL_loop_or_end_return:F {
            \__bnvs_kip_x_path_resolve_end_return_true:
         }
1649
       } {
1650
          \__bnvs_tl_put_right:cx { a } { . \__bnvs_seq_use:cn { a } . }
1651
          \__bnvs_kip_x_path_resolve_VAL_loop_or_end_return:F {
1652
            \__bnvs_seq_pop_right:ccT { a } { c } {
1653
              \__bnvs_seq_put_left:cv { b } { c }
1654
1655
              _bnvs_kip_x_path_resolve_loop_or_end_return:
1656
         }
1658
       }
     } {
1659
        \BNVS_end:
1660
```

```
1661
                              \prg_return_false:
                      1662
                      1663 }
                          \BNVS_new:cpn { kip_n_path_resolve_or_end_return:nF } #1 #2 {
                      1664
                            \__bnvs_get:nvcTF { #1 } { a } { b } {
                      1665
                              \__bnvs_kip:cccTF { b } { id } { path } {
                      1666
                                 \__bnvs_tl_set_eq:cc { key } { b }
                      1667
                                 \__bnvs_seq_merge:cc { path } { b }
                      1668
                                 \__bnvs_seq_set_eq:cc { a } { path }
                       1669
                                 \__bnvs_seq_clear:c { b }
                                 \__bnvs_kip_n_path_resolve_loop_or_end_return:
                      1671
                              } {
                      1672
                                 \__bnvs_seq_pop_right:ccTF { a } { c } {
                      1673
                                   \__bnvs_seq_put_left:cv { b } { c }
                      1674
                                   \__bnvs_kip_n_path_resolve_loop_or_end_return:
                      1675
                                  {
                      1676
                                   \__bnvs_kip_n_path_resolve_end_return_true:
                      1677
                      1678
                              }
                      1680
                            } {
                              #2
                            }
                      1682
                      1683 }
                          \BNVS_new:cpn { kip_n_path_resolve_VAL_loop_or_end_return: } {
                      1684
                            \__bnvs_kip_n_path_resolve_or_end_return:nF V {
                      1685
                                 _bnvs_kip_n_path_resolve_or_end_return:nF A {
                      1686
                                 \__bnvs_kip_n_path_resolve_or_end_return:nF L {
                      1687
                                   \__bnvs_seq_pop_right:ccTF { a } { c } {
                      1688
                                     \__bnvs_seq_put_left:cv { b } { c }
                      1689
                                     \__bnvs_kip_n_path_resolve_loop_or_end_return:
                                  }
                      1692
                                      \__bnvs_kip_n_path_resolve_end_return_true:
                      1693
                                }
                      1694
                              }
                      1695
                      1696
                      1697
                          \BNVS_new:cpn { kip_n_path_resolve_end_return_false: } {
                      1698
                            \BNVS_end:
                      1699
                            \prg_return_false:
                      1700
                      1701
                          \BNVS_new:cpn { kip_n_path_resolve_end_return_true: } {
                            \__bnvs_end_kip_export:
                      1704
                            \prg_return_true:
                      1705 }
_bnvs_kip_n_path_resolve_loop_or_end_return:
                      Loop to resolve the path.
```

```
1706 \BNVS_new:cpn { kip_n_path_resolve_loop_or_end_return: } {
     \__bnvs_call:TF {
```

```
\__bnvs_tl_set_eq:cc { a } { key }
          _bnvs_seq_if_empty:cTF { a } {
1709
          \_bnvs_seq_if_empty:cTF { b } {
            \__bnvs_kip_n_path_resolve_end_return_true:
1711
               _bnvs_kip_n_path_resolve_VAL_loop_or_end_return:
1714
       }
         {
1715
            _bnvs_tl_put_right:cx { a } { . \__bnvs_seq_use:cn { a } . }
1716
1717
          \__bnvs_kip_n_path_resolve_VAL_loop_or_end_return:
1718
     } {
1719
        \BNVS_end:
1720
        \prg_return_false:
1721
1723 }
```

\_bnvs\_kip\_n\_path\_resolve:

This is the entry point to resolve the path. Local variables:

- \...key\_tl, \...id\_tl, \...path\_seq contain the resolution.
- $\bullet$  ...a\_tl contains the name with a partial index path currently resolved.
- \...a\_seq contains the dotted path components to be resolved. It equals \...path\_seq at the beginning
- \...b\_seq is used as well. Initially empty.

```
1724 \BNVS_new_conditional:cpnn { kip_n_path_resolve: } { T, F, TF } {
1725 \BNVS_begin:
1726 \__bnvs_seq_set_eq:cc { a } { path }
1727 \__bnvs_seq_clear:c { b }
1728 \__bnvs_kip_n_path_resolve_loop_or_end_return:
1729 }
```

## 6.13.8 Evaluation bricks

We start by helpers.

```
\__bnvs_round_ans:n
\__bnvs_round:c
\__bnvs_round_ans:
```

```
\_bnvs_round:c \langle tl \ core \ name \rangle \_bnvs_round_ans: \_bnvs_round_ans:n \{\langle expression \rangle\}
```

The first function replaces the variable content with its rounded floating point evaluation. The second function replaces ans tl variable content with its rounded floating point evaluation. The last function appends to the ans tl variable the rounded floating point evaluation of the argument.

```
1730 \BNVS_new:cpn { round_ans:n } #1 {

1731 \tl_if_empty:nTF { #1 } {

1732 \__bnvs_tl_put_right:cn { ans } { 0 }
```

```
1733
                                                                                                                                                                                      _bnvs_tl_put_right:cx { ans } { \fp_eval:n { round(#1) } }
                                                                                                                                1734
                                                                                                                                1735
                                                                                                                                1736 }
                                                                                                                                                     \BNVS_new:cpn { round:N } #1 {
                                                                                                                                1737
                                                                                                                                                               \tl_if_empty:NTF #1 {
                                                                                                                                1738
                                                                                                                                                                          \tl_set:Nn #1 { 0 }
                                                                                                                                1739
                                                                                                                                1740
                                                                                                                                                                          \tl_set:Nx #1 { \fp_eval:n { round(#1) } }
                                                                                                                                1741
                                                                                                                                1742
                                                                                                                                1743 }
                                                                                                                                                    \BNVS_new:cpn { round:c } {
                                                                                                                                1744
                                                                                                                                                              \BNVS_tl_use:Nc \__bnvs_round:N
                                                                                                                                1745
                                                                                                                                1746 }
\BNVS_end_return_false:
                                                                                                                                                  \BNVS_end_return_false:x
                                                                                                                                                                                                                                                                                                      \__bnvs_end_return_false:
                                                                                                                                                                                                                                                                                                     \__bnvs_end_return_false:x {\message\}
                                                                                                                             End a group and calls \prg_return_false:. The message is for debugging only.
                                                                                                                                                  \cs_new:Npn \BNVS_end_return_false: {
                                                                                                                                                              \BNVS_end:
                                                                                                                                1748
                                                                                                                                                               \prg_return_false:
                                                                                                                                1749
                                                                                                                                1750
                                                                                                                                                  \cs_new:Npn \BNVS_end_return_false:x #1 {
                                                                                                                                1751
                                                                                                                                                               \__bnvs_error:x { #1 }
                                                                                                                                1752
                                                                                                                                                              \BNVS_end_return_false:
                                                                                                                                1753
                                                                                                                               1754 }
\__bnvs_value_resolve:ncTF
                                                                                                                                                                                 \cline{1.8} \cli
\__bnvs_value_resolve:vcTF
                                                                                                                                                                                \cline{1.5cm} 
\__bnvs_value_append:ncTF
       __bnvs_value_append:(xc|vc){\it TF}
```

Resolve the content of the  $\langle key \rangle$  value counter into the  $\langle tl\ variable \rangle$  or append this value to the right of the variable. Execute  $\langle yes\ code \rangle$  when there is a  $\langle value \rangle$ ,  $\langle no\ code \rangle$  otherwise. Inside the  $\langle no\ code \rangle$  branch, the content of the  $\langle tl\ variable \rangle$  is undefined. Implementation detail: we return the first in the cache for subkey V and in the general prop for subkey V. Once we have found a value, we feed the previous items such that the next search stops at the first item. The cache contains an integer which is the computed value from the general prop. A group is created while appending but not while resolving.

```
1755 \BNVS_new:cpn { value_resolve_return:nnnT } #1 #2 #3 #4 {
1756    \__bnvs_tl_if_empty:cTF { #3 } {
1757    \prg_return_false:
1758    } {
1759    \__bnvs_cache_gput:nnv V { #2 } { #3 }
1760    #4
1761    \prg_return_true:
1762    }
1763 }
```

```
\BNVS_new_conditional:cpnn { quark_if_nil:c } #1 { T, F, TF } {
      \BNVS_tl_use:Nc \quark_if_nil:NTF { #1 } {
1765
         \prg_return_true:
1766
      }
        {
1767
         \prg_return_false:
1768
1769
1770 }
    \BNVS_new_conditional:cpnn {    quark_if_no_value:c } #1 { T, F, TF } {
1771
      \BNVS_tl_use:Nc \quark_if_no_value:NTF { #1 } {
1772
         \prg_return_true:
1773
      } {
1774
         \prg_return_false:
1775
1776
1777 }
    \BNVS_new_conditional:cpnn { value_resolve:nc } #1 #2 { T, F, TF } {
1778
      \_bnvs_cache_get:nncTF V { #1 } { #2 } {
1779
         \prg_return_true:
1780
1781
         \_bnvs_get:nncTF V { #1 } { #2 } {
1782
           \__bnvs_quark_if_nil:cTF { #2 } {
We can retrieve the value from either the first or last index.
             \__bnvs_gput:nnn V { #1 } { \q_no_value }
1784
             \__bnvs_first_resolve:ncTF { #1 } { #2 } {
1785
               \label{lem:lem:nnnT} A \ \{ \ \#1 \ \} \ \{ \ \#2 \ \} \ \{
1786
                 \__bnvs_gput:nnn V { #1 } { \q_nil }
1787
               }
1788
             } {
1789
               \__bnvs_last_resolve:ncTF { #1 } { #2 } {
1790
                  \__bnvs_value_resolve_return:nnnT Z { #1 } { #2 } {
                    \__bnvs_gput:nnn V { #1 } { \q_nil }
                 }
               } {
1794
                  \__bnvs_gput:nnn V { #1 } { \q_nil }
1795
                 \prg_return_false:
1796
               }
1797
1798
          }
1799
1800
                _bnvs_quark_if_no_value:cTF { #2 } {
               \__bnvs_fatal:n {Circular~definition:~#1}
             } {
Possible recursive call.
               \_bnvs_if_resolve:vcTF { #2 } { #2 } {
                  \__bnvs_value_resolve_return:nnnT V { #1 } { #2 } {
1804
                    \_bnvs_gput:nnn V { #1 } { \q_nil }
1805
1806
               } {
1807
                    _bnvs_gput:nnn V { #1 } { \q_nil }
1808
                 \prg_return_false:
1809
1810
1811
1812
        } {
1813
```

```
\prg_return_false:
1814
        }
1815
      }
1816
1817 }
    \BNVS_new_conditional:cpnn { value_resolve:vc } #1 #2 { T, F, TF } {
1818
      \BNVS_tl_use:Nv \__bnvs_value_resolve:ncTF { #1 } { #2 } {
1819
        \prg_return_true:
1820
        {
1821
        \prg_return_false:
1822
      }
1823
1824 }
    1825
      \BNVS_tl_use:nv {
1826
        \BNVS_end:
1827
        \__bnvs_tl_put_right:cn { #2 }
1828
      } { #1 }
1829
1830 }
    \BNVS_new_conditional:cpnn { value_append:nc } #1 #2 { T, F, TF } {
1831
      \BNVS_begin:
1832
      \__bnvs_value_resolve:ncTF { #1 } { #2 } {
1834
        \BNVS_end_tl_put_right:cv { #2 } { #2 }
1835
        \prg_return_true:
      } {
1836
        \BNVS_end:
1837
        \prg_return_true:
1838
1839
1840 }
    \BNVS_new_conditional_vc:cn { value_append } { T, F, TF }
cTF:nnnnvalueFIRST2222
```

```
\__bnvs_first_resolve:nc\overline{TF} \__bnvs_first_resolve:ncTF {\langle key\rangle} \langle tl variable \ {\langle yes code\rangle} \ {\langle no code\rangle} \ \__bnvs_first_append:ncTF {\langle key\rangle} \ \tauriable \ \ {\langle yes code\rangle} \ {\langle no code\rangle} \ \__bnvs_first_append:(xc|vc)\overline{TF}
```

Resolve the first index of the  $\langle key \rangle$  slide range into the  $\langle tl \ variable \rangle$  or append the first index of the  $\langle tey \rangle$  slide range to the  $\langle tl \ variable \rangle$ . If no resolution occurs the content of the  $\langle tl \ variable \rangle$  is undefined in the first case and unmodified in the second. Cache the result. Execute  $\langle yes \ code \rangle$  when there is a  $\langle first \rangle$ ,  $\langle no \ code \rangle$  otherwise.

The first index must be computed separately from the length and the last index.

```
\__bnvs_gput:nnn A { #1 } { \q_nil }
                  \prg_return_false:
1856
               }
                 {
1857
                    _bnvs_gput:nnn A { #1 } { \q_nil }
1858
                  \__bnvs_cache_gput:nnv A { #1 } { #2 }
1859
                  \prg_return_true:
1860
               }
1861
             } {
1862
                \__bnvs_error:n {
   \__bnvs_gput:nnn A { #1 } { \q_nil }
                \prg_return_false:
1866
             }
1867
           } {
1868
              \__bnvs_error:n {
1869
   Unavailable~last~for~#1~(\token_to_str:N\__bnvs_first_resolve:ncTF/1) }
1870
              \__bnvs_gput:nnn A { #1 } { \q_nil }
1871
              \prg_return_false:
1872
           }
         } {
              _bnvs_quark_if_no_value:cTF { a } {
              \__bnvs_fatal:n {Circular~definition:~#1}
1876
           } {
1877
             \__bnvs_if_resolve:vcTF { #2 } { #2 } {
1878
                \__bnvs_cache_gput:nnv A { #1 } { #2 }
1879
                \prg_return_true:
1880
             } {
1881
1882
                \prg_return_false:
             }
1883
           }
         }
       } {
1887
          \prs_return_false:
1888
     }
1889
1890
   \BNVS_new_conditional_vc:cn { first_resolve } { T, F, TF }
1891
   \BNVS_new_conditional:cpnn { first_append:nc } #1 #2 { T, F, TF } {
1892
     \BNVS_begin:
1893
       _bnvs_first_resolve:ncTF { #1 } { #2 } {
       \BNVS_end_tl_put_right:cv { #2 } { #2 }
       \prg_return_true:
     } {
1897
       \prg_return_false:
1898
     }
1899
1900 }
```

\\_\_bnvs\_last\_resolve:nc<u>TF</u> \\_\_bnvs\_last\_append:nc<u>TF</u>

```
\__bnvs_last_resolve:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\ \__bnvs_last_append:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\
```

Resolve the last index of the fully qualified  $\langle key \rangle$  range into or to the right of the right of the  $\langle tl \ variable \rangle$ , when possible. Execute  $\langle yes \ code \rangle$  when a last index was given,  $\langle no \ code \rangle$  otherwise.

```
\BNVS_new_conditional:cpnn { last_resolve:nc } #1 #2 { T, F, TF } {
      \__bnvs_cache_get:nncTF Z { #1 } { #2 } {
1902
        \prg_return_true:
1903
1904
        \__bnvs_get:nncTF Z { #1 } { #2 } {
1905
          \__bnvs_quark_if_nil:cTF { #2 } {
1906
             \_bnvs_gput:nnn Z { #1 } { \q_no_value }
1907
The last index must be computed separately from the start and the length.
             \__bnvs_first_resolve:ncTF { #1 } { #2 } {
               \__bnvs_tl_put_right:cn { #2 } { + }
1909
               \__bnvs_length_append:ncTF { #1 } { #2 } {
1910
                 \__bnvs_tl_put_right:cn { #2 } { - 1 }
1911
                 \__bnvs_round:c { #2 }
1912
                 \__bnvs_cache_gput:nnv Z { #1 } { #2 }
1913
                 \__bnvs_gput:nnn Z { #1 } { \q_nil }
1914
                 \prg_return_true:
1915
              } {
                 \__bnvs_error:x {
     Unavailable~length~for~#1~(\token_to_str:N \__bnvs_last_resolve:ncTF/1) }
1918
                 \__bnvs_gput:nnn Z { #1 } { \q_nil }
1919
1920
                 \prg_return_false:
              }
1921
            } {
1922
               \__bnvs_error:x {
1923
    Unavailable~first~for~#1~(\token_to_str:N \__bnvs_last_resolve:ncTF/1) }
1924
               \__bnvs_gput:nnn Z { #1 } { \q_nil }
1925
               \prg_return_false:
1926
            }
1927
          } {
1928
             \_bnvs_quark_if_no_value:cTF { #2 } {
               \__bnvs_fatal:n {Circular~definition:~#1}
1930
            } {
1931
               \__bnvs_if_resolve:vcTF { #2 } { #2 } {
1932
                 \_bnvs_cache_gput:nnv Z { #1 } { #2 }
1933
                 \prg_return_true:
1934
1935
               }
1936
                 \prg_return_false:
               }
            }
          }
        } {
1940
           \prg_return_false:
1941
        }
1942
      }
1943
1944
    \BNVS_new_conditional_vc:cn { last_resolve } { T, F, TF }
1945
    \prg_new_conditional:Npnn \__bnvs_last_append:nc #1 #2 { T, F, TF } {
1946
1947
      \BNVS_begin:
      \__bnvs_last_resolve:ncTF { #1 } { #2 } {
1949
        \BNVS_end_tl_put_right:cv { #2 } { #2 }
        \prg_return_true:
1950
      } {
1951
```

```
1952 \BNVS_end:
1953 \prg_return_false:
1954 }
1955 }
1956 \BNVS_new_conditional_vc:cn { last_append } { T, F, TF }
```

```
\__bnvs_length_resolve:nc\overline{\mathit{TF}} \__bnvs_length_resolve:ncTF \{\langle key \rangle\}\ \langle \mathit{tl\ variable} \rangle\ \{\langle \mathit{yes\ code} \rangle\}\ \{\langle \mathit{no\ code} \rangle\}\ \__bnvs_length_append:ncTF \{\langle \mathit{key} \rangle\}\ \langle \mathit{tl\ variable} \rangle\ \{\langle \mathit{yes\ code} \rangle\}\ \{\langle \mathit{no\ code} \rangle\}\
```

Resolve the length of the  $\langle key \rangle$  slide range into  $\langle tl \ variable \rangle$ , or append the length of the  $\langle key \rangle$  slide range to  $\langle tl \ variable \rangle$ . Execute  $\langle yes \ code \rangle$  when there is a  $\langle length \rangle$ ,  $\langle no \ code \rangle$  otherwise.

The length must be computed separately from the start and the last index.

```
\__bnvs_last_resolve:ncTF { #1 } { #2 } {
1964
              \__bnvs_tl_put_right:cn { #2 } { - }
1965
              \__bnvs_first_append:ncTF { #1 } { #2 } {
1966
                 \__bnvs_tl_put_right:cn { #2 } { + 1 }
1967
                 \__bnvs_round:c { #2 }
1968
                 \__bnvs_gput:nnn L { #1 } { \q_nil }
1969
                 \__bnvs_cache_gput:nnv L { #1 } { #2 }
                 \prg_return_true:
              } {
1972
                 \__bnvs_error:n {
1973
   Unavailable~first~for~#1~(\__bnvs_length_resolve:ncTF/2) }
1974
                 \return_false:
1975
              }
1976
            } {
1977
              \__bnvs_error:n {
1978
   Unavailable~last~for~#1~(\__bnvs_length_resolve:ncTF/1) }
1979
              \return_false:
            }
1981
          } {
            \__bnvs_quark_if_no_value:cTF { #2 } {
1983
              \__bnvs_fatal:n {Circular~definition:~#1}
1984
            } {
1985
              \_bnvs_if_resolve:vcTF { #2 } { #2 } {
1986
                 \_bnvs_cache_gput:nnv L { #1 } { #2 }
1987
                 \prg_return_true:
1988
              } {
1989
                 \prg_return_false:
              }
            }
          }
1993
       } {
1994
```

```
\prg_return_false:
1995
       }
1996
     }
1997
   }
1998
    \BNVS_new_conditional_vc:cn { length_resolve } { T, F, TF }
1999
    \BNVS_new_conditional:cpnn { length_append:nc } #1 #2 { T, F, TF } {
2000
      \BNVS_begin:
2001
      \__bnvs_length_resolve:ncTF { #1 } { #2 } {
2002
        \BNVS_end_tl_put_right:cv { #2 } { #2 }
        \prg_return_true:
2004
     } {
2005
        \prg_return_false:
2006
2007
2008
   \BNVS_new_conditional_vc:cn { length_append } { T, F, TF }
```

\_\_bnvs\_range\_resolve:nc<u>TF</u> \_\_bnvs\_range\_append:nc<u>TF</u>

```
\__bnvs_range_resolve:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\ \__bnvs_range_append:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\
```

Resolve the range of the  $\langle key \rangle$  slide range into the  $\langle tl \ variable \rangle$  or append this range to the  $\langle tl \ variable \rangle$ . Execute  $\langle yes \ code \rangle$  when there is a  $\langle range \rangle$ ,  $\langle no \ code \rangle$  otherwise, in that latter case the content the  $\langle tl \ variable \rangle$  is undefined on resolution only.

```
\BNVS_new_conditional:cpnn { range_append:nc } #1 #2 { T, F, TF } {
      \BNVS_begin:
2011
        _bnvs_first_resolve:ncTF { #1 } { a } {
2012
        \BNVS_tl_use:Nv \int_compare:nNnT { a } < 0 {
2013
           __bnvs_tl_set:cn { a } { 0 }
2014
2015
        \__bnvs_last_resolve:ncTF { #1 } { b } {
2016
Limited from above and below.
          \BNVS_tl_use:Nv \int_compare:nNnT { b } < 0 {
2018
             2019
          \__bnvs_tl_put_right:cn { a } { - }
          \__bnvs_tl_put_right:cv { a } { b }
2021
          \BNVS_end_tl_put_right:cv { #2 } { a }
2022
          \prg_return_true:
2023
        } {
2024
Limited from below.
          \BNVS_end_tl_put_right:cv { #2 } { a }
2025
          \__bnvs_tl_put_right:cn { #2 } { - }
2026
           \prg_return_true:
2027
2028
      }
        \__bnvs_last_resolve:ncTF { #1 } { b } {
Limited from above.
          \BNVS_tl_use:Nv \int_compare:nNnT { b } < 0 {
2031
            \__bnvs_tl_set:cn { b } { 0 }
2032
          }
2033
          \__bnvs_tl_put_left:cn { b } { - }
2034
          \BNVS_end_tl_put_right:cv { #2 } { b }
2035
```

```
2036
          \prg_return_true:
        } {
2037
             _bnvs_value_resolve:ncTF { #1 } { b } {
2038
          \BNVS_tl_use:Nv \int_compare:nNnT { b } < 0 {
2039
             \__bnvs_tl_set:cn { b } { 0 }
2040
2041
Unlimited range.
             \BNVS_end_tl_put_right:cv { #2 } { b }
2042
             \__bnvs_tl_put_right:cn { #2 } { - }
2043
             \prg_return_true:
2044
          } {
             \BNVS_end:
             \prg_return_false:
          }
2048
        }
2049
      }
2050
    }
2051
    \BNVS_new_conditional_vc:cn { range_append } { T, F, TF }
2052
    \BNVS_new_conditional:cpnn { range_resolve:nc } #1 #2 { T, F, TF } {
2053
      \_bnvs_tl_clear:c { #2 }
2054
      \__bnvs_range_append:ncTF { #1 } { #2 } {
2055
        \prg_return_true:
      } {
2057
        \prg_return_false:
2058
      }
2059
2060 }
2061 \BNVS_new_conditional_vc:cn { range_resolve } { T, F, TF }
```

 $\label{lem:code} $$\sum_{\substack{-b \text{nvs\_previous\_append:ncTF} \\ code}} $$ \left(\frac{TF}{\langle key \rangle} \right) \left(\frac{TF}{\langle key \rangle}$ 

Resolve the index after the  $\langle key \rangle$  slide range into the  $\langle tl \ variable \rangle$ , or append this index to the variable. Execute  $\langle yes \ code \rangle$  when there is a  $\langle next \rangle$  index,  $\langle no \ code \rangle$  otherwise. In the latter case, the  $\langle tl \ variable \rangle$  is undefined on resolution only.

```
\BNVS_new_conditional:cpnn { previous_resolve:nc } #1 #2 { T, F, TF } {
2062
      \__bnvs_cache_get:nncTF P { #1 } { #2 } {
2063
        \prg_return_true:
2064
2065
        \__bnvs_first_resolve:ncTF { #1 } { #2 } {
2066
          \__bnvs_tl_put_right:cn { #2 } { -1 }
2067
          \__bnvs_round:c { #2 }
          \__bnvs_cache_gput:nnv P { #1 } { #2 }
          \prg_return_true:
2070
       } {
2071
          \prs_return_false:
2072
       }
2073
     }
2074
2075 }
   \BNVS_new_conditional_vc:cn { previous_resolve } { T, F, TF }
   \BNVS_new_conditional:cpnn { previous_append:nc } #1 #2 { T, F, TF } {
```

\\_\_bnvs\_next\_resolve:nc<u>TF</u> \\_\_bnvs\_next\_append:nc<u>TF</u>

```
\__bnvs_next_resolve:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\ \__bnvs_next_append:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\
```

Resolve the index after the  $\langle key \rangle$  slide range into the  $\langle tl \ variable \rangle$ , or append this index to this variable. Execute  $\langle yes \ code \rangle$  when there is a  $\langle next \rangle$  index,  $\langle no \ code \rangle$  otherwise. In the latter case, the content of the  $\langle tl \ variable \rangle$  is undefined, on resolution only.

```
\BNVS_new_conditional:cpnn { next_resolve:nc } #1 #2 { T, F, TF } {
2088
      \__bnvs_cache_get:nncTF N { #1 } { #2 } {
2089
        \prg_return_true:
2090
2091
          _bnvs_last_resolve:ncTF { #1 } { #2 } {
          \__bnvs_tl_put_right:cn { #2 } { +1 }
          \__bnvs_round:c { #2 }
          \__bnvs_cache_gput:nnv N { #1 } { #2 }
          \prg_return_true:
2096
       } {
2097
          \prg_return_false:
2098
       }
2099
     }
2100
2101
   \BNVS_new_conditional_vc:cn { next_resolve } { T, F, TF }
2102
   \BNVS_new_conditional:cpnn { next_append:nc } #1 #2 { T, F, TF } {
      \BNVS_begin:
2104
      \__bnvs_next_resolve:ncTF { #1 } { #2 } {
2105
        \BNVS_end_tl_put_right:cv { #2 } { #2 }
2106
        \prg_return_true:
2107
     } {
2108
        \BNVS_end:
2109
        \prg_return_true:
2110
2111
2112 }
   \BNVS_new_conditional_vc:cn { next_append } { T, F, TF }
```

\\_\_bnvs\_v\_resolve:nc<u>TF</u> \\_\_bnvs\_v\_append:nc<u>TF</u>

```
\__bnvs_v_resolve:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\ \__bnvs_v_append:ncTF \{\langle key \rangle\}\ \langle tl\ variable \rangle\ \{\langle yes\ code \rangle\}\ \{\langle no\ code \rangle\}\
```

Resolve the value of the  $\langle key \rangle$  overlay set into the  $\langle tl \ variable \rangle$  or append this value to the right of this variable. Execute  $\langle yes \ code \rangle$  when there is a  $\langle value \rangle$ ,  $\langle no \ code \rangle$  otherwise. In the latter case, the content of the  $\langle tl \ variable \rangle$  is undefined, on resolution only.

```
2114 \BNVS_new_conditional:cpnn { v_resolve:nc } #1 #2 { T, F, TF } {
```

```
\__bnvs_v_get:ncTF { #1 } { #2 } {
2115
        \__bnvs_quark_if_no_value:cTF { #2 } {
2116
          \__bnvs_fatal:n {Circular~definition:~#1}
2117
          \prg_return_false:
2118
        } {
2119
          \prg_return_true:
2120
        }
2121
      } {
2122
        \__bnvs_v_gput:nn { #1 } { \q_no_value }
2123
        \__bnvs_value_resolve:ncTF { #1 } { #2 } {
2124
          \__bnvs_v_gput:nv { #1 } { #2 }
2125
          \prg_return_true:
2126
        } {
2127
          \__bnvs_first_resolve:ncTF { #1 } { #2 } {
2128
            \__bnvs_v_gput:nv { #1 } { #2 }
2129
            \prg_return_true:
2130
          } {
2131
            \__bnvs_last_resolve:ncTF { #1 } { #2 } {
2132
            \__bnvs_v_gput:nv { #1 } { #2 }
               \prg_return_true:
            } {
               \__bnvs_v_gremove:n { #1 }
2136
               \prg_return_false:
2137
2138
          }
2139
        }
2140
      }
2141
2142 }
    \BNVS_new_conditional_vc:cn { v_resolve } { T, F, TF }
2143
    \BNVS_new_conditional:cpnn { v_append:nc } #1 #2 { T, F, TF } {
      \BNVS_begin:
2145
      \__bnvs_v_resolve:ncTF { #1 } { #2 } {
2146
        \BNVS_end_tl_put_right:cv { #2 } { #2 }
2147
        \prg_return_true:
2148
2149
        \BNVS_end:
2150
        \prg_return_false:
2152
2153 }
2154 \BNVS_new_conditional_vc:cn { v_append } { T, F, TF }
```

Resolve the index associated to the  $\langle key \rangle$  and  $\langle integer \rangle$  slide range into the  $\langle tl \ variable \rangle$  or append this index to the right of this variable. When  $\langle integer \rangle$  is 1, this is the first index, when  $\langle integer \rangle$  is 2, this is the second index, and so on. When  $\langle integer \rangle$  is 0, this is the index, before the first one, and so on. If the computation is possible,  $\langle yes \ code \rangle$  is executed, otherwise  $\langle no \ code \rangle$  is executed. In the latter case, the content of the  $\langle tl \ variable \rangle$  is undefined, on resolution only. The computation may fail when too many recursion calls are made.

```
\BNVS_new_conditional:cpnn { index_can:n } #1 { p, T, F, TF } {
2155
      \bool_if:nTF {
2156
            \__bnvs_if_in_p:nn V { #1 }
2157
        || \__bnvs_if_in_p:nn A { #1 }
2158
        || \__bnvs_if_in_p:nn Z { #1 }
2159
      } {
2160
2161
        \prg_return_true:
      } {
2162
        \prg_return_false:
2163
2164
2165 }
    \BNVS_new_conditional:cpnn { index_can:v } #1 { p, T, F, TF } {
2166
      \BNVS_tl_use:Nv \__bnvs_index_can:nTF { #1 } {
2167
        \prg_return_true:
2168
2169
      ን ፈ
2170
        \prg_return_false:
      }
2171
2172 }
    \BNVS_new_conditional:cpnn { index_resolve:nnc } #1 #2 #3 { T, F, TF } {
2173
      \exp_args:Nx \__bnvs_value_resolve:ncTF { #1.#2 } { #3 } {
2174
           \prg_return_true:
2175
2176
2177
           _bnvs_first_resolve:ncTF { #1 } { #3 } {
           \__bnvs_tl_put_right:cn { #3 } { + #2 - 1 }
        \__bnvs_round:c { #3 }
2179
           \prg_return_true:
Limited overlay set.
        } {
           \__bnvs_last_resolve:ncTF { #1 } { #3 } {
2182
             \_bnvs_tl_put_right:cn { #3 } { + #2 - 1 }
           \__bnvs_round:c { #3 }
2184
             \prg_return_true:
2185
          } {
2186
             \__bnvs_value_resolve:ncTF { #1 } { #3 } {
               \__bnvs_tl_put_right:cn { #3 } { + #2 - 1 }
             \__bnvs_round:c { #3 }
2189
               \prg_return_true:
2190
             } {
2191
```

```
2192
            \prg_return_false:
            }
2193
          }
2194
     }
2195
      }
2196
2197
    \BNVS_new_conditional:cpnn { index_resolve:nvc } #1 #2 #3 { T, F, TF } {
2198
      \BNVS_tl_use:nv {
2199
        \__bnvs_index_resolve:nncTF { #1 }
      } { #2 } { #3 } {
2201
2202
        \prg_return_true:
      } {
2203
        \prg_return_false:
2204
      }
2205
2206 }
    \BNVS_new_conditional:cpnn { index_resolve:vvc } #1 #2 #3 { T, F, TF } {
2207
      \BNVS_tl_use:nv {
2208
        \BNVS_tl_use:Nv \__bnvs_index_resolve:nncTF { #1 }
2209
      } { #2 } { #3 } {
        \prg_return_true:
      } {
2212
2213
        \prg_return_false:
      }
2215 }
    \BNVS_new_conditional:cpnn { index_append:nnc } #1 #2 #3 { T, F, TF } {
2216
      \BNVS_begin:
      \__bnvs_index_resolve:nncTF { #1 } { #2 } { #3 } {
2218
        \BNVS_end_tl_put_right:cv { #3 } { #3 }
2219
        \prg_return_true:
2220
      } {
        \BNVS_end:
2222
2223
        \prg_return_false:
      }
2224
2225 }
   \BNVS_new_conditional:cpnn { index_append:vvc } #1 #2 #3 { T, F, TF } {
2226
      \BNVS_tl_use:nv {
        \BNVS_tl_use:Nv \__bnvs_index_append:nncTF { #1 }
2228
      } { #2 } { #3 } {
2229
2230
        \prg_return_true:
      } {
        \prg_return_false:
2233
      }
2234 }
```

# 6.13.9 Index counter

```
\__bnvs_n_resolve:ncTF
\__bnvs_n_append:ncTF
\__bnvs_n_append:VcTF
```

```
\verb|\__bnvs_n_resolve:ncTF| \{\langle key \rangle\} \ \langle \textit{tl variable} \rangle \ \{\langle \textit{yes code} \rangle\} \ \{\langle \textit{no code} \rangle\}
```

Evaluate the n counter associated to the  $\{\langle key \rangle\}$  overlay set into  $\langle tl \ variable \rangle$ . Initialize this counter to 1 on the first use.  $\langle no \ code \rangle$  is never executed.

```
\mbox{\sc bnvs_new\_conditional:cpnn} { n_resolve:nc } #1 #2 { T, F, TF } {
```

```
_bnvs_n_get:ncF { #1 } { #2 } {
        \__bnvs_tl_set:cn { #2 } { 1 }
        \__bnvs_n_gput:nn { #1 } { 1 }
2238
2239
      \prg_return_true:
2240
2241 }
    \BNVS_new_conditional:cpnn { n_append:nc } #1 #2 { T, F, TF } {
2242
      \BNVS_begin:
2243
      \__bnvs_n_resolve:ncTF { #1 } { #2 } {
        \BNVS_end_tl_put_right:cv { #2 } { #2 }
2245
2246
        \prg_return_true:
     } {
2247
        \BNVS_end:
2248
        \prg_return_false:
2249
2250
2251 }
   \BNVS_new_conditional_vc:cn { n_append } { T, F, TF }
```

```
\__bnvs_n_index_resolve:nc<u>TF</u>
\__bnvs_n_index_append:nc<u>TF</u>
\__bnvs_n_index_resolve:nnc<u>TF</u>
\__bnvs_n_index_append:nnc<u>TF</u>
```

Resolve the index for the value of the n counter associated to the  $\{\langle key \rangle\}$  overlay set into the  $\langle tl \ variable \rangle$  or append this value the right of this variable. Initialize this counter to 1 on the first use. If the computation is possible,  $\langle yes \ code \rangle$  is executed, otherwise  $\langle no \ code \rangle$  is executed. In the latter case, the content of the  $\langle tl \ variable \rangle$  is undefined on resolution only.

```
\BNVS_new_conditional:cpnn { n_index_resolve:nc } #1 #2 { T, F, TF } {
2253
     \__bnvs_n_resolve:ncTF { #1 } { #2 } {
2254
       \__bnvs_index_resolve:nvcTF { #1 } { #2 } { #2 } {
        \prg_return_true:
2256
      } {
2257
         \prg_return_false:
2258
      }
2259
     } {
2260
       \prg_return_false:
2261
     }
2262
2263 }
   2264
     \__bnvs_n_resolve:ncTF { #1 } { #3 } {
2265
       \__bnvs_tl_put_left:cn { #3 } { #2. }
2266
       \__bnvs_if_resolve:vcTF { #3 } { #3 } {
2267
        \prg_return_true:
2268
      } {
2269
         \prg_return_false:
2270
      }
     } {
```

```
\prg_return_false:
     }
2274
2275 }
    \BNVS_new_conditional:cpnn { n_index_append:nc } #1 #2 { T, F, TF } {
2276
      \BNVS_begin:
      \__bnvs_n_index_resolve:ncTF { #1 } { #2 } {
2278
        \BNVS_end_tl_put_right:cv { #2 } { #2 }
2279
        \prg_return_true:
2280
     } {
        \BNVS_end:
2282
2283
        \prg_return_false:
     }
2284
2285 }
    \BNVS_new_conditional:cpnn { n_index_append:nnc } #1 #2 #3 { T, F, TF } {
2286
      \BNVS_begin:
2287
      \__bnvs_n_index_resolve:nncTF { #1 } { #2 } { #3 } {
2288
        \BNVS_end_tl_put_right:cv { #3 } { #3 }
2289
        \prg_return_true:
2290
     } {
        \BNVS_end:
        \prg_return_false:
     }
2294
2295 }
   \BNVS_new_conditional_vc:cn { n_index_append } { T, F, TF }
   \BNVS_new_conditional_vvc:cn { n_index_append } { T, F, TF }
```

## 6.13.10 Value counter

Increment the value counter position accordingly. When requested, put the result in the  $\langle tl \ variable \rangle$ . In the second version, the result will lay within the declared range.

```
\BNVS_new_conditional:cpnn { v_incr_resolve:nnc } #1 #2 #3 { T, F, TF } {
2299
      \__bnvs_if_resolve:ncTF { #2 } { #3 } {
        \BNVS_tl_use:Nv \int_compare:nNnTF { #3 } = 0 {
          \__bnvs_v_resolve:ncTF { #1 } { #3 } {
2302
            \prg_return_true:
          } {
2303
            \prg_return_false:
2304
          }
2305
       } {
2306
          \__bnvs_tl_put_right:cn { #3 } { + }
2307
          \__bnvs_v_append:ncTF { #1 } { #3 } {
2308
2309
            \__bnvs_round:c { #3 }
            \__bnvs_v_gput:nv { #1 } { #3 }
            \prg_return_true:
2312
          } {
```

```
2313
            \prg_return_false:
2314
       }
     } {
2316
        \prg_return_false:
2317
2318
2319 }
    \BNVS_new_conditional_vnc:cn {    v_incr_resolve } { T, F, TF }
2320
   \BNVS_new_conditional:cpnn { v_incr_append:nnc } #1 #2 #3 { T, F, TF } {
     \BNVS_begin:
2322
      \_bnvs_v_incr_resolve:nncTF { #1 } { #2 } { #3 } {
2323
        \BNVS_end_tl_put_right:cv { #3 } { #3 }
2324
        \prg_return_true:
     } {
2326
        \prg_return_false:
2327
2328
2329
   \BNVS_new_conditional_vnc:cn { v_incr_append } { T, F, TF }
2330
   \BNVS_new_conditional_vvc:cn { v_incr_append } { T, F, TF }
   \BNVS_new_conditional:cpnn { v_post_resolve:nnc } #1 #2 #3 { T, F, TF } {
      \__bnvs_v_resolve:ncTF { #1 } { #3 } {
        \BNVS_begin:
2334
        \__bnvs_if_resolve:ncTF { #2 } { a } {
2335
          \BNVS_tl_use:Nv \int_compare:nNnTF { a } = 0 {
2336
            \BNVS_end:
            \prg_return_true:
2338
          } {
2339
            \_bnvs_tl_put_right:cn { a } { + }
2340
            \__bnvs_tl_put_right:cv { a } { #3 }
2341
            \__bnvs_round:c { a }
            \BNVS_end_v_gput:nc { #1 } { a }
2343
2344
            \prg_return_true:
         }
2345
       } {
2346
          \BNVS_end:
2347
          \prg_return_false:
2348
2349
     } {
2350
2351
          \prg_return_false:
     }
2352
   \BNVS_new_conditional_vvc:cn { v_post_resolve } { T, F, TF }
   \BNVS_new_conditional:cpnn { v_post_append:nnc } #1 #2 #3 { T, F, TF } {
2355
      \BNVS_begin:
2356
      \_bnvs_v_post_resolve:nncTF { #1 } { #2 } { #3 } {
2357
        \BNVS_end_tl_put_right:cv { #3 } { #3 }
2358
        \prg_return_true:
2359
     } {
2360
2361
        \prg_return_true:
2362
   \BNVS_new_conditional_vnc:cn { v_post_append } { T, F, TF }
   \BNVS_new_conditional_vvc:cn { v_post_append } { T, F, TF }
```

```
\__bnvs_n_incr_resolve:nnncTF
                                                                                                                                                                                                                                                                                                                                                                                                           \label{locality} $$\sum_{n=1}^{\infty} \frac{\langle key \rangle}{\langle base\ key \rangle} \ {\langle offset \rangle} \ \langle tl\ core\ busyed
   \__bnvs_n_incr_resolve:vvncTF
                                                                                                                                                                                                                                                                                                                                                                                                         name\ {\langle yes code\} {\langle no code\}
                                                                                                                                                                                                                                                                                                                                                                                                           \verb|\__bnvs_n_incr_resolve:nncTF| \{\langle key \rangle\} \ \{\langle offset \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \
\__bnvs_n_incr_resolve:nnc<u>TF</u>
                         _bnvs_n_incr_resolve:vvc<u>TF</u>
                                                                                                                                                                                                                                                                                                                                                                                                           code\rangle} {\langle no \ code \rangle}
                                                                                                                                                                                                                                                                                                                                                                                                           \__bnvs_n_incr_append:nnncTF
\__bnvs_n_incr_append:nnc<u>TF</u>
                                                                                                                                                                                                                                                                                                                                                                                                         name \rangle \{\langle yes code \rangle\} \{\langle no code \rangle\}
                                                                                                                                                                                                                                                                                                                                                                                                         \verb|\__bnvs_n_incr_append:nncTF| \{\langle key \rangle\} \ \{\langle offset \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \{\langle yes \rangle\} \ \langle tl \ core \ name \rangle \ \langle tl \ core 
 \__bnvs_n_incr_append:(vnc|vvc)\underline{\mathit{TF}}
                                                                                                                                                                                                                                                                                                                                                                                                         code} {\langle no \ code \rangle}
\__bnvs_n_post_resolve:nnc_TF
\__bnvs_n_post_append:nncTF
```

Increment the implicit n counter accordingly. When requested, put the resulting index in the variable with  $\langle tl \ core \ name \rangle$ .

```
\BNVS_new_conditional:cpnn { n incr resolve:nnnc } #1 #2 #3 #4 { T, F, TF } {
      \_bnvs_if_resolve:ncTF { #3 } { #4 } {
2367
    \BNVS_tl_use:Nv \int_compare:nNnTF { #4 } = 0 {
2368
          \_bnvs_n_resolve:ncTF { #1 } { #4 } {
            \_bnvs_index_resolve:nvcTF { #1 } { #4 } { #4 } {
              \prg_return_true:
            } {
2372
2373
              \prg_return_false:
2374
          } {
            \prg_return_false:
2376
          }
2377
       } {
2378
          \__bnvs_tl_put_right:cn { #4 } { + }
2379
          \__bnvs_n_append:ncTF { #1 } { #4 } {
2380
            \__bnvs_round:c { #4 }
            \__bnvs_n_gput:nv { #1 } { #4 }
2382
            \__bnvs_index_resolve:nvcTF { #2 } { #4 } { #4 } {
2383
2384
              \prg_return_true:
            } {
2385
              \prg_return_false:
2386
2387
          }
2388
            \prg_return_false:
2389
        }
     } {
2393
        \prg_return_false:
2394
2395 }
   \BNVS new conditional:cpnn { n incr resolve:nnc } #1 #2 #3 { T, F, TF } {
2396
      bnvs if resolve:ncTF { #2 } { #3 } {
2397
   \BNVS_tl_use:Nv \int_compare:nNnTF { #3 } = 0 {
2398
          \_bnvs_n_resolve:ncTF { #1 } { #3 } {
2399
            \_bnvs_index_resolve:nvcTF { #1 } { #3 } { #3 } {
2400
              \prg_return_true:
            } {
              \prg_return_false:
2403
            }
2404
          } {
2405
```

```
}
                         2407
                                 } {
                         2408
                                      _bnvs_tl_put_right:cn { #3 } { + }
                         2409
                                    \__bnvs_n_append:ncTF { #1 } { #3 } {
                         2410
                                      \__bnvs_round:c { #3 }
                         2411
                                      \__bnvs_n_gput:nv { #1 } { #3 }
                         2412
                                      \__bnvs_index_resolve:nvcTF { #1 } { #3 } { #3 } {
                         2413
                                        \prg_return_true:
                                      } {
                                        \prg_return_false:
                         2417
                                   } {
                         2418
                                      \prg_return_false:
                         2419
                         2420
                                  }
                         2421
                         2422
                                  \prg_return_false:
                         2423
                         2424
                            }
                         2425
                             \BNVS_new_conditional_vnc:cn { n_incr_resolve } { T, F, TF }
                             \BNVS_new_conditional_vvc:cn { n_incr_resolve } { T, F, TF }
                             \BNVS_new_conditional:cpnn { n_incr_append:nnnc } #1 #2 #3 #4 { T, F, TF } {
                         2428
                               \BNVS begin:
                         2429
                               \_bnvs_n_incr_resolve:nnncTF { #1 } { #2 } { #3 } { #4 }{
                         2430
                                  \BNVS_end_tl_put_right:cv { #4 } { #4 }
                         2431
                         2432
                                  \prg_return_true:
                               } {
                         2433
                                  \BNVS_end:
                         2434
                                  \prg_return_false:
                               }
                         2436
                         2437 }
                             \BNVS_new_conditional_vvnc:cn { n_incr_append } { T, F, TF }
                         2438
                             \BNVS_new_conditional_vvvc:cn { n_incr_append } { T, F, TF }
                             \BNVS_new_conditional:cpnn { n_incr_append:nnc } #1 #2 #3 { T, F, TF } {
                         2440
                               \BNVS_begin:
                         2441
                                \_bnvs_n_incr_resolve:nncTF { #1 } { #2 } { #3 } {
                         2442
                         2443
                                  \BNVS_end_tl_put_right:cv { #3 } { #3 }
                                  \prg_return_true:
                               } {
                                  \BNVS_end:
                         2447
                                  \prg_return_false:
                               }
                         2448
                         2449 }
                             \BNVS_new_conditional_vnc:cn { n_incr_append } { T, F, TF }
                         2450
                             \BNVS_new_conditional_vvc:cn { n_incr_append } { T, F, TF }
\__bnvs_v_post_resolve:nnc_<u>TF</u>
                                      \__bnvs_v_post_resolve:vvc<u>TF</u>
                                      code} {\langle no \ code \rangle}
\_{\rm bnvs\_v\_post\_append:nnc}
                                      \cline{1.8} L_bnvs_v_post_append:nncTF {\langle key \rangle} {\langle offset \rangle} {\langle tl \ variable \rangle} {\langle yes \rangle}
\__bnvs_v_post_append:(vnN|vvN)TF
                                      code\rangle} {\langle no \ code \rangle}
```

2406

\prg\_return\_false:

Resolve the value of the free counter for the given  $\langle key \rangle$  into the  $\langle tl \ variable \rangle$  then increment this free counter position accordingly. The append version, appends the value to the right of the  $\langle tl \ variable \rangle$ . The content of the  $\langle tl \ variable \rangle$  is undefined while in the  $\{\langle no \ code \rangle\}$  branch and on resolution only.

```
\BNVS_new_conditional:cpnn { n_post_resolve:nnc } #1 #2 #3 { T, F, TF } {
     \__bnvs_n_resolve:ncTF { #1 } { #3 } {
2453
        \BNVS_begin:
2454
        \__bnvs_if_resolve:ncTF { #2 } { #3 } {
2455
          \BNVS_tl_use:Nv \int_compare:nNnTF { #3 } = 0 {
2456
            \BNVS_end:
2457
            \_bnvs_index_resolve:nvcTF { #1 } { #3 } { #3 } {
2458
              \prg_return_true:
2459
            } {
              \prg_return_false:
2462
            7
          } {
2463
            \__bnvs_tl_put_right:cn { #3 } { + }
2464
            \__bnvs_n_append:ncTF { #1 } { #3 } {
2465
              \__bnvs_round:c { #3 }
2466
              \__bnvs_n_gput:nv { #1 } { #3 }
2467
              \BNVS_end:
2468
              \__bnvs_index_resolve:nvcTF { #1 } { #3 } { #3 } {
2469
                 \prg_return_true:
              } {
2472
                 \prg_return_false:
              }
            } {
2474
              \BNVS_end:
2475
              \prg_return_false:
2476
            }
2477
          }
2478
       } {
2479
          \BNVS_end:
          \prg_return_false:
        }
2483
     } {
2484
        \prg_return_false:
     }
2485
2486
   \BNVS_new_conditional:cpnn { n_post_append:nnc } #1 #2 #3 { T, F, TF } {
2487
      \BNVS_begin:
2488
      \__bnvs_n_post_resolve:nncTF { #1 } { #2 } { #3 } {
2489
        \BNVS_end_tl_put_right:cv { #3 } { #3 }
2490
        \prg_return_true:
2491
     } {
        \BNVS_end:
        \prg_return_false:
     }
2495
2496 }
   \BNVS_new_conditional_vnc:cn { n_post_append } { T, F, TF }
   \BNVS_new_conditional_vvc:cn { n_post_append } { T, F, TF }
```

## 6.13.11 Evaluation

}

2522 2523 }

```
_bnvs_round_ans:
                           \__bnvs_rslv_round:
                           Helper function to round the \l__bnvs_ans_tl variable. For ranges only, this will be set
                           to \prg_do_nothing because we do not want to interpret the - sign as a minus operator.
                           2499 \BNVS_set:cpn { round_ans: } {
                                 \__bnvs_round:c { ans }
                           2501 }
                           6.13.12
                                      Functions for the resolution
                           They manily start with \__bnvs_if_resolve_
     _bnvs_if_resolve_end_return_false:n
                                             \verb|\_bnvs_if_resolve_end_return_false:n {|\langle message \rangle|}
                           Close one TEX group, display a message and return false.
bnvs_path_resolve_n:TFF
                           _{2502} \ \BNVS_new:cpn \ \{ path_resolve_n:TFF \ \} \ \#1 \ \#2 \ \{
                                 \__bnvs_kip_n_path_resolve:TF  {
                                   \__bnvs_seq_if_empty:cTF { path } { #1 } { #2 }
                           2505
                           2506 }
 _bnvs_path_resolve_n:T
                           \_ bnvs_path_resolve_n:T {\langle yes\ code \rangle}
                           Resolve the path and execute \langle yes \ code \rangle on success.
                               \BNVS_new:cpn { if_resolve_end_return_false:n } #1 {
                                 \BNVS_end:
                           2508
                                 \prg_return_false:
                           2509
                           2510 }
                           2511
                               \BNVS_new:cpn { path_resolve_n:T } #1 {
                           2512
                                 \__bnvs_path_resolve_n:TFF {
                                   #1
                                 } {
                           2514
                                   \__bnvs_if_resolve_end_return_false:n {
                           2516
                                     Too~many~dotted~components
                                   }
                           2517
                                 } {
                           2518
                                      _bnvs_if_resolve_end_return_false:n {
                           2519
                                     Unknown~dotted~path
                           2520
                           2521
```

```
\BNVS_set:cpn { resolve_x:T } #1 {
      \__bnvs_kip_x_path_resolve:TFF {
2525
        #1
2526
      }
        {
2527
           _bnvs_if_resolve_end_return_false:n {
2528
          Too~many~dotted~components
2529
2530
      } {
2531
          _bnvs_if_resolve_end_return_false:n {    Unknown~dotted~path }
2532
      }
2533
2534 }
```

\_\_bnvs\_path\_suffix:n*TF* 

If the last item of  $\l_bnvs_path_seq$  is  $\langle suffix \rangle$ , then execute  $\langle yes\ code \rangle$  otherwise execute  $\langle no\ code \rangle$ . The suffix is n in the second case.

For  $\_$ \_bnvs\_if\_resolve\_pop\_kip:TTF. If the split sequence is empty, execute  $\langle end code \rangle$ . Otherwise pops the 3 heading items of the split sequence into the three t1 variables key, id, path. If key is blank then execute  $\langle blank \ code \rangle$ , otherwise execute  $\langle black \ code \rangle$ .

For  $\_$ \_bnvs\_if\_resolve\_pop\_complete\_white:T: pops the three heading items of the split sequence into the three variables n\_incr, incr, post. Then execute  $\langle blank \ code \rangle$ .

For  $\_$ \_bnvs\_if\_resolve\_pop\_complete\_black:T: pops the six heading items of the split sequence then execute  $\langle blank \ code \rangle$ .

```
\BNVS_new:cpn { if_resolve_pop_kip_complete: } {
2539
      \_bnvs_tl_if_blank:vT { id } {
2540
        \__bnvs_tl_put_left:cv { key } { id_last }
2541
        \__bnvs_tl_set:cv { id } { id_last }
2542
        _bnvs_tl_if_blank:vTF { path } {
        \__bnvs_seq_clear:c { path }
2545
2546
          _bnvs_seq_set_split:cnv { path } { . } { path }
2547
        \__bnvs_seq_remove_all:cn { path } { }
2548
2549
        _bnvs_tl_set_eq:cc { key_base } { key }
2550
      \__bnvs_seq_set_eq:cc { path_base } { path }
2551
2552 }
2553 \BNVS_new:cpn { if_resolve_pop_kip:TTF } #1 #2 #3 {
```

The first 3 capture groups are empty, and the 3 next ones are expected to contain the expected information.

```
2558
             }
2559
               \__bnvs_if_resolve_pop_kip_complete:
               #2
             }
          } {
       _bnvs_end_unreachable_return_false:n { if_resolve_pop_kip:TTF/2 }
2565
          }
        } {
2566
       _bnvs_end_unreachable_return_false:n {    if_resolve_pop_kip:TTF/1 }
2567
        }
2568
      } { #3 }
2569
2570 }
```

\_\_bnvs\_if\_resolve\_pop\_complete:nNT

```
\__bnvs_if_resolve_pop_kip:FFTF \{\langle empty\ key\ code\rangle\}\ \{\langle no\ id\ code\rangle\}\ \{\langle true\ code\rangle\}\ \{\langle no\ capture\ code\rangle\}\ \__bnvs_if_resolve_pop_complete:nNT \{\langle tl\rangle\}\ \langle tl\ var\rangle\ \{\langle true\ code\rangle\}\
```

 $\langle tl \rangle$  and  $\langle tl \ var \rangle$  are the arguments of the \\_\_bnvs\_if\_resolve:nc conditionals. conditional variants.

\\_\_bnvs\_if\_resolve\_pop\_kip:FFTF locally sets the key, id and path t1 variables to the 3 heading items of the split sequence, which correspond to the 3 eponym capture groups. If no capture group is available,  $\langle no\ capture\ code \rangle$  is executed. If the capture group for the key is empty, then  $\langle empty\ key\ code \rangle$  is executed. If there is no capture group for the id, then  $\langle no\ id\ code \rangle$  is executed. Otherwise  $\langle true\ code \rangle$  is executed.

\\_\_bnvs\_rslv\_pop\_end: T locally sets the three tl variables n\_incr, incr and post to the three heading items of the split sequence, which correspond to the last 3 eponym capture groups.

```
\BNVS_new:cpn { if_resolve_pop_complete_white:T } #1 {
      \__bnvs_split_pop_left:cTF { n_incr } {
        \__bnvs_split_pop_left:cTF { incr } {
2573
          \__bnvs_split_pop_left:cTF { post } {
2574
            #1
2575
         } {
2576
      _bnvs_end_unreachable_return_false:n {    if_resolve_pop_complete_white:T/3 }
2577
         }
2578
       } {
2579
    \__bnvs_end_unreachable_return_false:n {    if_resolve_pop_complete_white:T/2 }
2580
   \__bnvs_end_unreachable_return_false:n { if_resolve_pop_complete_white:T/1 }
2584
2585 }
   \BNVS_new:cpn { if_resolve_pop_complete_black:T } #1 {
2586
      \__bnvs_split_pop_left:cTF { a } {
2587
        \_bnvs_split_pop_left:cTF { a } {
2588
```

```
\__bnvs_split_pop_left:cTF { a } {
2589
           \__bnvs_split_pop_left:cTF { a } {
2590
             \__bnvs_split_pop_left:cTF { a } {
2591
               \__bnvs_split_pop_left:cTF { a } {
2592
2593
              } {
2594
   \__bnvs_end_unreachable_return_false:n {    if_resolve_pop_complete_black:T/6 }
              }
2596
            } {
   \__bnvs_end_unreachable_return_false:n {    if_resolve_pop_complete_black:T/5 }
            }
          } {
2600
   \__bnvs_end_unreachable_return_false:n {    if_resolve_pop_complete_black:T/4 }
2601
          }
2602
        } {
2603
   \__bnvs_end_unreachable_return_false:n {    if_resolve_pop_complete_black:T/3 }
2604
2605
       } {
2606
   \__bnvs_end_unreachable_return_false:n {    if_resolve_pop_complete_black:T/2 }
      }
     } {
}
2611
2612 }
```

```
\__bnvs_if_resolve:ncTF
                                  \verb|\__bnvs_if_append:ncTF| \{\langle expression \rangle\} \ \langle tl \ variable \rangle \ \{\langle yes \ code \rangle\} \ \{\langle no \ code \rangle\} 
  _bnvs_if_resolve:vc<u>TF</u>
                                  Evaluates the \langle expression \rangle, replacing all the named overlay specifications by their static
\__bnvs_if_append:ncTF
                                  counterpart then put the rounded result in \( \lambda t \) variable \( \rangle \) when resolving or to the right of
\_{
m bnvs\_if\_append:(vc|xc)
                                  the (tl variable) when appending. Executed within a group. Heavily used by \__bnvs_-
                                  query_eval:nc, where \( \langle integer expression \rangle \) was initially enclosed inside '?(\( \ldots \))'. Local
                                  variables:
                                  To feed \langle tl \ variable \rangle with.
             \l__bnvs_ans_tl
                                  (End definition for \l_bnvs_ans_tl.)
                                  The sequence of catched query groups and non queries.
          \l__bnvs_split_seq
                                  (End\ definition\ for\ \l_bnvs_split_seq.)
                                  Is the index of the non queries, before all the catched groups.
          \l__bnvs_split_int
                                  (End\ definition\ for\ \l_\_bnvs\_split\_int.)
                                   2613 \BNVS_int_new:c { split }
                                  Storage for split sequence items that represent names.
             \l__bnvs_key_tl
                                  (End\ definition\ for\ \l_bnvs_key_tl.)
                                  Storage for split sequence items that represent integer paths.
            \l__bnvs_path_tl
                                  (End definition for \l_bnvs_path_tl.)
                                  Catch circular definitions. Open a main T<sub>F</sub>X group to define local functions and variables,
                                  sometimes another grouping level is used. The main T<sub>F</sub>X group is closed in the various
                                  \...end_return... functions.
                                       \BNVS_new:cpn { kip_x_path_resolve_or_end_return_false:nT } #1 #2 {
                                   2614
                                          \__bnvs_kip_x_path_resolve:TFF {
                                   2615
                                            #2
                                   2616
                                         } {
                                   2617
                                            \BNVS_end_return_false:x { Too~many~dotted~components:~#1 }
                                   2618
                                   2619
                                            \BNVS_end_return_false:x { Unknown~dotted~path:~#1 }
                                   2620
                                          }
                                   2621
                                   2622 }
                                       \BNVS_new_conditional:cpnn { if_append:nc } #1 #2 { T, F, TF } {
                                   2623
                                          \BNVS begin:
                                   2624
                                          \_bnvs_if_resolve:ncTF { #1 } { #2 } {
                                   2625
                                            \BNVS_end_tl_put_right:cv { #2 } { #2 }
                                   2626
                                   2627
                                       \BNVS_DEBUG_log_if_append_ncTF:nn { ... } { ...TRUE }
                                   2628
                                       \langle /! final \rangle
                                   2629
                                            \prg_return_true:
                                   2630
                                          } {
                                   2631
                                            \BNVS_end:
                                       \langle *!final \rangle
                                       \BNVS_DEBUG_log_if_append_ncTF:nn { ... } { ...FALSE }
                                   2634
                                   <sub>2635</sub> ⟨/!final⟩
```

\prg\_return\_false:

2636 2637

```
2638
    \BNVS_new:cpn { end_unreachable_return_false:n } #1 {
2639
      \_bnvs_error:x { UNREACHABLE/#1 }
2640
      \BNVS_end:
2641
      \prg_return_false:
2642
2643
    \BNVS_new_conditional:cpnn { if_resolve:nc } #1 #2 { T, F, TF } {
2644
      \__bnvs_call:TF {
2645
        \BNVS_begin:
        \BNVS_set:cpn { if_resolve_warning:n } ##1 {
2647
           \__bnvs_warning:n { #1:~##1 }
           \BNVS_set:cpn { if_resolve_warning:n } {
2649
             \use_none:n
2650
2651
2652
This T<sub>F</sub>X group will be closed just before returning. Implementation:
      \__bnvs_regex_split:cnTF { split } { #1 } {
The leftmost item is not a special item: we start feeding \l_bnvs_ans_tl with it.
           \BNVS_set:cpn { if_resolve_end_return_true: } {
Normal and unique end of the loop.
             \__bnvs_if_resolve_round_ans:
             \BNVS_tl_use:nv {
               \BNVS_end:
2657
               \__bnvs_tl_set:cn { #2 }
             } { ans }
             \prg_return_true:
2660
2661
           \BNVS_set:cpn { if_resolve_round_ans: } { \__bnvs_round_ans: }
2662
           \__bnvs_tl_clear:c { ans }
2663
           \__bnvs_if_resolve_loop_or_end_return:
2664
2665
           \_\_bnvs_tl_clear:c { ans }
2666
           \_\_bnvs\_round\_ans:n { #1 }
           \BNVS_end_tl_set:cv { #2 } { ans }
2669
           \prg_return_true:
        }
2670
      } {
2671
           _bnvs_error:n { TOO_MANY_NESTED_CALLS/Resolution }
2672
        \prg_return_false:
2673
2674
2675
    \BNVS_new_conditional:cpnn { if_append:vc } #1 #2 { T, F, TF } {
2676
      \BNVS_tl_use:Nv \__bnvs_if_append:ncTF { #1 } { #2 } {
        \prg_return_true:
2678
      } {
2679
2680
        \prg_return_false:
      }
2681
2682 }
    \BNVS_new_conditional:cpnn { if_resolve:vc } #1 #2 { T, F, TF } {
2683
      \BNVS_tl_use:Nv \__bnvs_if_resolve:ncTF { #1 } { #2 } {
2684
        \prg_return_true:
2685
```

```
2686 } {
2687 \prg_return_false:
2688 }
2689 }
```

Next functions are helpers for the  $\_$ \_bnvs\_if\_resolve:nc conditional variants. When present, their two first arguments  $\langle tl \rangle$  and  $\langle tl \ var \rangle$  are exactly the ones given to the variants.

\\_\_bnvs\_if\_resolve\_loop\_or\_end\_return: \\_\_bnvs\_if\_resolve\_loop\_or\_end\_return:

May call itself at the end.

```
\BNVS_new:cpn { if_resolve_loop_or_end_return: } {
     \__bnvs_split_pop_left:cTF { a } {
        \__bnvs_tl_put_right:cv { ans } { a }
2692
2693
        \__bnvs_if_resolve_pop_kip:TTF {
          \__bnvs_if_resolve_pop_kip:TTF {
2694
      _bnvs_end_unreachable_return_false:n {    if_resolve_loop_or_end_return:/3 }
2695
         } {
2696
            \__bnvs_if_resolve_pop_complete_white:T {
2697
              \__bnvs_tl_if_blank:vTF { n_incr } {
2698
                \__bnvs_tl_if_blank:vTF { incr } {
2699
                  \__bnvs_tl_if_blank:vTF { post } {
2700
                    \__bnvs_if_resolve_value_loop_or_end_return_true:F {
```

Only the dotted path, branch according to the last component.

```
\__bnvs_seq_pop_right:ccTF { path } { a } {
2702
                        \BNVS_tl_use:Nv \str_case:nnF { a } {
2703
2704 { n
               } { \BNVS_use:c { if_resolve_loop_or_end_return[.n]: } }
   { length
                } { \BNVS_use:c { if_resolve_loop_or_end_return[.length]: } }
   { last
                } { \BNVS_use:c { if_resolve_loop_or_end_return[.last]: } }
   { range
                } { \BNVS_use:c { if_resolve_loop_or_end_return[.range]: } }
               } { \BNVS_use:c { if_resolve_loop_or_end_return[.previous]: } }
   { previous
   { next
                } { \BNVS_use:c { if_resolve_loop_or_end_return[.next]: } }
                } { \BNVS_use:c { if_resolve_loop_or_end_return[.reset]: } }
     reset_all } { \BNVS_use:c { if_resolve_loop_or_end_return[.reset_all]: } }
                        } {
    \BNVS_use:c { if_resolve_loop_or_end_return[...<integer>]: }
2713
2714
                      } {
    \BNVS_use:c { if_resolve_loop_or_end_return[...]: }
2716
                    }
2718
                  } {
2719
    \BNVS_use:c { if_resolve_loop_or_end_return[...++]: }
                  }
               } {
                    _bnvs_path_suffix:nTF { n } {
   \BNVS_use:c { if_resolve_loop_or_end_return[...n+=...]: }
2724
                  } {
2725
    \BNVS_use:c { if_resolve_loop_or_end_return[...+=...]: }
2726
                  }
2727
                }
2728
2729
             } {
```

```
\BNVS_use:c { if_resolve_loop_or_end_return[...++n]: }
                                   }
                              }
2732
                        } {
         % split sequence empty
2734
          \__bnvs_end_unreachable_return_false:n {    if_resolve_loop_or_end_return:/2 }
2735
2736
                   } {
2737
                          \__bnvs_if_resolve_pop_complete_black:T {
2738
                              \_\begin{tabular}{ll} \_\begin{tabular}{ll} \_\begin{tabular}{ll} \\\ \begin{tabular}{ll} 
2739
          \BNVS_use:c { if_resolve_loop_or_end_return[++...n]: }
                             } {
2741
          \BNVS_use:c { if_resolve_loop_or_end_return[++...]: }
2742
2743
                              }
2744
                   } {
2745
                           \__bnvs_if_resolve_end_return_true:
2746
2747
              } {
2748
           \__bnvs_end_unreachable_return_false:n { if_resolve_loop_or_end_return:/1 }
2750
2751 }
         \BNVS_set:cpn { if_resolve_value_loop_or_end_return_true:F } #1 {
2752
               \__bnvs_tl_set:cx { a } {
2753
                    \BNVS_tl_use:c { key } \BNVS_tl_use:c { path }
2754
               \__bnvs_v_resolve:vcTF { a } { a } {
2756
                    \__bnvs_tl_put_right:cv { ans } { a }
2757
                    \__bnvs_if_resolve_loop_or_end_return:
2758
              } {
2759
                    \__bnvs_value_resolve:vcTF { a } { a } {
2760
                         \__bnvs_tl_put_right:cv { ans } { a }
2761
                         \__bnvs_if_resolve_loop_or_end_return:
2762
                   } {
2763
                         #1
2764
                   }
2765
2766
2767
2768
         \BNVS_new:cpn { end_return_error:n } #1 {
                          \__bnvs_error:n { #1 }
                         \BNVS_end:
2771
                         \prg_return_false:
2772 }
2773 \BNVS_new:cpn { if_resolve_loop_or_end_return[.n]: } {
         • Case ...n.
               \__bnvs_path_resolve_n:T {
2774
                    \__bnvs_base_resolve_n:
2775
                    \__bnvs_n_index_append:vvcTF { key } { key_base } { ans } {
2776
                          \__bnvs_if_resolve_loop_or_end_return:
                   } {
                                _bnvs_end_return_error:n {
2779
                              Undefined~dotted~path
2780
```

```
}
        }
2782
      }
2783
2784 }
    \BNVS_new_conditional:cpnn { path_suffix:n } #1 { T, F, TF } {
      \__bnvs_seq_get_right:ccTF { path } { a } {
        \__bnvs_tl_if_eq:cnTF { a } { #1 } {
2787
           \__bnvs_seq_pop_right:ccT { path } { a } { }
2788
          \prg_return_true:
2789
        } {
2790
           \prs_return_false:
2792
      } {
2793
        \prg_return_false:
2794
2795
2796 }
   \BNVS_new:cpn { if_resolve_loop_or_end_return[.length]: } {
   • Case ...length.
      \__bnvs_path_resolve_n:T {
        \_bnvs_length_append:vcTF { key } { ans } {
2800 % \begin{bnvs.gobble}
    \langle *!final \rangle
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../length }
   \langle /! final \rangle
   % \end{bnvs.gobble}
          \__bnvs_if_resolve_loop_or_end_return:
2805
        } {
2806
           \__bnvs_if_resolve_end_return_false:n { NO~length }
2807
2808
      }
2809
    \BNVS_new:cpn { if_resolve_loop_or_end_return[.last]: } {
   • Case ...last.
      \__bnvs_path_resolve_n:T {
        \__bnvs_last_append:vcTF { key } { ans } {
2814 % \begin{bnvs.gobble}
   \langle *!final \rangle
<code>2816 \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../last }</code>
   \langle /! final \rangle
2817
   % \end{bnvs.gobble}
2818
          \__bnvs_if_resolve_loop_or_end_return:
2819
2820
          \BNVS_end_return_false:x { NO~last }
        }
      }
2824 }
2825 \BNVS_new:cpn { if_resolve_loop_or_end_return[.range]: } {
   • Case ...range.
```

```
\__bnvs_path_resolve_n:T {
        \__bnvs_range_append:vcTF { key } { ans } {
2827
          \BNVS_set:cpn { if_resolve_round_ans: } { \prg_do_nothing: }
2828
2829 % \begin{bnvs.gobble}
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../range }
   \langle /! final \rangle
   % \end{bnvs.gobble}
          \__bnvs_if_resolve_loop_or_end_return:
           __bnvs_if_resolve_end_return_false:n {    NO~range }
2836
2837
2838
2839 }
    \BNVS_new:cpn { if_resolve_loop_or_end_return[.previous]: } {
   • Case ...previous.
      \__bnvs_path_resolve_n:T {
        \__bnvs_previous_append:vcTF { key } { ans } {
2843 % \begin{bnvs.gobble}
   ⟨*!final⟩
2845 \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../previous }
   ⟨/!final⟩
   % \end{bnvs.gobble}
2847
          \__bnvs_if_resolve_loop_or_end_return:
          \__bnvs_if_resolve_end_return_false:n { NO~previous }
2852
     }
2853 }
   \BNVS_new:cpn { if_resolve_loop_or_end_return[.next]: } {
   • Case ...next.
      \__bnvs_path_resolve_n:T {
2855
        \_bnvs_next_append:vcTF { key } { ans } {
   % \begin{bnvs.gobble}
    \langle *!final \rangle
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../next }
   ⟨/!final⟩
2861
   % \end{bnvs.gobble}
          \__bnvs_if_resolve_loop_or_end_return:
2862
2863
          \__bnvs_if_resolve_end_return_false:n { NO~next }
2864
2865
     }
2866
    \BNVS_new:cpn { if_resolve_loop_or_end_return[.reset]: } {
   • Case ...reset.
      \__bnvs_path_resolve_n:T {
2869
        \__bnvs_v_greset:vnT { key } { } { }
2870
        \__bnvs_value_append:vcTF { key } { ans } {
2871
```

```
2872 % \begin{bnvs.gobble}
2873 (*!final)
2874 \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../reset }
   \langle /! final \rangle
   % \end{bnvs.gobble}
2876
          \__bnvs_if_resolve_loop_or_end_return:
2877
2878
           \__bnvs_if_resolve_end_return_false:n { NO~reset }
     }
2881
2882 }
   \BNVS_new:cpn { if_resolve_loop_or_end_return[.reset_all]: } {
   • Case ...reset_all.
      \__bnvs_path_resolve_n:T {
2884
        \_bnvs_greset_all:vnT { key } { } { }
        \_bnvs_value_append:vcTF { key } { ans } {
   % \begin{bnvs.gobble}
   \langle *!final \rangle
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../reset }
   ⟨/!final⟩
2891 % \end{bnvs.gobble}
          \__bnvs_if_resolve_loop_or_end_return:
2892
        } {
2893
             _bnvs_if_resolve_end_return_false:n { NO~reset }
2894
2895
     }
2896
   \BNVS_set:cpn { if_resolve_loop_or_end_return[...<integer>]: } {
   • Case ...\langle integer \rangle.
      \__bnvs_path_resolve_n:T {
2899
        \__bnvs_index_append:vvcTF { key } { a } { ans } {
2900
2901 % \begin{bnvs.gobble}
   \langle *!final \rangle
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../<integer> }
   \langle /! final \rangle
   % \end{bnvs.gobble}
          \__bnvs_if_resolve_loop_or_end_return:
           \__bnvs_if_resolve_end_return_false:n { NO~integer }
2908
2909
     }
2910
2911 }
   \BNVS_set:cpn { if_resolve_loop_or_end_return[...]: } {
2912
   • Case . . . .
      \__bnvs_path_resolve_n:T {
        \_bnvs_value_append:vcTF { key } { ans } {
2915 % \begin{bnvs.gobble}
2916 (*!final)
2917 \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../... }
2918 (/!final)
```

```
2919 % \end{bnvs.gobble}
          \__bnvs_if_resolve_loop_or_end_return:
2920
2921
             _bnvs_if_resolve_end_return_false:n { NO~value }
2922
2923
     }
2924
2925 }
    \BNVS_set:cpn { if_resolve_loop_or_end_return[...++]: } {
   • Case ...++.
      \_bnvs_path_suffix:nTF { reset } {
2927
        \__bnvs_path_resolve_n:T {
2928
          \__bnvs_v_greset:vnT { key } { } { }
2929
          \__bnvs_v_post_append:vncTF { key } { 1 } { ans } {
2930
            \__bnvs_if_resolve_loop_or_end_return:
2931
          } {
2932
            \__bnvs_if_resolve_end_return_false:n { NO~post }
2933
2934
2935
       }
2936
     } {
          _bnvs_path_suffix:nTF { reset_all } {
          \__bnvs_path_resolve_n:T {
            \__bnvs_greset_all:vnT { key } { } { }
2939
            \_bnvs_v_post_append:vncTF { key } { 1 } { ans } {
2940
               \__bnvs_if_resolve_loop_or_end_return:
2941
            } {
2942
                 _bnvs_if_resolve_end_return_false:n { NO~post }
2943
2944
          }
2945
        } {
          \__bnvs_path_resolve_n:T {
            \__bnvs_v_post_append:vncTF { key } { 1 } { ans } {
               \__bnvs_if_resolve_loop_or_end_return:
            }
2950
               \__bnvs_if_resolve_end_return_false:n { NO~post }
2951
2952
          }
2953
       }
2954
2955
   \BNVS_set:cpn { if_resolve_loop_or_end_return[...n+=...]: } {
   • Case ....n+=\langle integer \rangle.
      \__bnvs_path_resolve_n:T {
        \__bnvs_base_resolve_n:
2959
        \__bnvs_n_incr_append:vvvcTF { key } { key_base } { incr } { ans } {
2961 % \begin{bnvs.gobble}
   ⟨*!final⟩
   \label{log_log_if_resolve_ncTF:nn} $$ \dots $$ { \dots / \dots n+= \dots }$
   \langle /! final \rangle
   % \end{bnvs.gobble}
          \__bnvs_if_resolve_loop_or_end_return:
```

```
_bnvs_if_resolve_end_return_false:n {
            NO~n~incrementation
2969
2970
        }
2971
      }
2972
2973 }
    \BNVS_set:cpn { if_resolve_loop_or_end_return[...+=...]: } {
2974
   • Case A + = \langle integer \rangle.
      \__bnvs_path_resolve_n:T {
        \__bnvs_v_incr_append:vvcTF { key } { incr } { ans } {
2977 % \begin{bnvs.gobble}
   \langle *!final \rangle
2978
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../...+=... }
2979
    \langle /! final \rangle
   % \end{bnvs.gobble}
2981
          \__bnvs_if_resolve_loop_or_end_return:
2982
            __bnvs_if_resolve_end_return_false:n {
2985
            {\tt NO\text{-}incremented-} {\tt value}
2986
2987
        }
      }
2988
2989
    \BNVS_new:cpn { base_resolve_n: } {
2990
      \_bnvs_seq_if_empty:cF { path_base } {
2991
        \__bnvs_seq_pop_right:cc { path_base } { a }
2992
        \__bnvs_seq_if_empty:cF { path_base } {
          \__bnvs_tl_put_right:cx { key_base } {
              \__bnvs_seq_use:cn { path_base } { . }
          }
2996
        }
2997
      }
2998
2999 }
    \BNVS_new:cpn { base_resolve: } {
3000
      \_bnvs_seq_if_empty:cF { path_base } {
3001
        \__bnvs_tl_put_right:cx { key_base } {
3002
            \__bnvs_seq_use:cn { path_base } { . }
3003
        }
      }
3005
3006 }
   \BNVS_new:cpn { if_resolve_loop_or_end_return[...++n]: } {

    Case ...++n.

      \__bnvs_path_resolve_n:T {
3008
        \__bnvs_base_resolve:
3009
        \_bnvs_n_incr_append:vvncTF { key } { key_base } { 1 } { ans } {
3010
          \__bnvs_if_resolve_loop_or_end_return:
3011
        } {
           \__bnvs_if_resolve_end_return_false:n { N0~...++n }
3013
        }
3014
      }
3015
3016 }
3017 \BNVS_set:cpn { if_resolve_loop_or_end_return[++...n]: } {
```

```
• Case ++...n.
      \__bnvs_path_resolve_n:T {
3018
        \__bnvs_base_resolve_n:
3019
        \_bnvs_n_incr_append:vvncTF { key } { key_base } { 1 } { ans } {
3020
           \__bnvs_if_resolve_loop_or_end_return:
3021
           \__bnvs_if_resolve_end_return_false:n { NO~++...n }
        }
      7
3025
3026 }
   \BNVS_new:cpn { if_resolve_loop_or_end_return[++...]: } {
3027
   • Case ++....
      \__bnvs_path_suffix:nTF { reset } {
3028
        \__bnvs_path_resolve_n:T {
3029
           \_bnvs_v_incr_append:vncTF { key } { 1 } { ans } {
3030
3031 % \begin{bnvs.gobble}
    \langle *!final \rangle
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../++...reset }
   \langle /! final \rangle
3035 % \end{bnvs.gobble}
         \begin{macrocode}
3036 %
             \__bnvs_v_greset:vnT { key } { } { }
3037
             \__bnvs_if_resolve_loop_or_end_return:
3038
            {
3039
             \__bnvs_v_greset:vnT { key } { } { }
3040
             \__bnvs_if_resolve_end_return_false:n { No~increment }
3041
          }
3042
        }
      } {
3044
        \__bnvs_path_suffix:nTF { reset_all } {
3045
           \__bnvs_path_resolve_n:T {
3046
             \__bnvs_v_incr_append:vncTF { key } { 1 } { ans } {
3047
3048 % \begin{bnvs.gobble}
   \langle *!final \rangle
3049
   \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../++...reset_all }
3050
   \langle /! final \rangle
3051
   % \end{bnvs.gobble}
3052
         \begin{macrocode}
3053
               \__bnvs_greset_all:vnT { key } { } { }
               \__bnvs_if_resolve_loop_or_end_return:
3055
            } {
3056
               \__bnvs_greset_all:vnT { key } { } { }
3057
               \__bnvs_if_resolve_end_return_false:n { No~increment }
3058
            }
3059
          }
3060
        }
          {
3061
           \__bnvs_path_resolve_n:T \{
3062
             \__bnvs_v_incr_append:vncTF { key } { 1 } { ans } {
3064 % \begin{bnvs.gobble}
   \langle *! final \rangle
3066 \BNVS_DEBUG_log_if_resolve_ncTF:nn { ... } { .../++... }
3067 (/!final)
```

```
_bnvs_query_eval:ncTF
                        Evaluates the single (overlay query), which is expected to contain no comma. Extract a
                       range specification from the argument, replaces all the named overlay specifications by
                        their static counterparts, make the computation then append the result to the right of
                        \l__bnvs_ans_tl. Ranges are supported with the colon syntax. This is executed within
                       a local TEX group managed by the caller. Below are local variables and constants.
                       Storage for a single value out of a range.
        \l__bnvs_V_tl
                        (End definition for \l__bnvs_V_tl.)
   \l__bnvs_TEST_A_tl Storage for the first component of a range.
                        (End definition for \l__bnvs_TEST_A_tl.)
        \l__bnvs_Z_tl Storage for the last component of a range.
                       \l_bnvs_L_tl Storage for the length component of a range.
                       (End definition for \l bnvs L tl.)
                       Used to parse slide range overlay specifications. A, A:Z, A::L on one side, :Z, :Z::L and
\c__bnvs_A_cln_Z_regex
                       ::L:Z on the other sides. Next are the capture groups.
                        (End definition for \c__bnvs_A_cln_Z_regex.)
                        3078 \regex_const:Nn \c__bnvs_A_cln_Z_regex {
                             \A \s* (?:
                           • 2: V
                               ( [^:]+? )
                           • 3, 4, 5: A : Z? or A :: L?
                                | ( [^:]+? ) \s* : (?: ( \s* [^:]*? ) | : ( \s* [^:]*? ) )
                           • 6, 7: ::(L:Z)?
                              | :: \s* (?: ( [^:]+? ) \s* : \s* ( [^:]+? ) )?
                           • 8, 9: :(Z::L)?
                              | : \s* (?: ( [^:]+? ) \s* :: \s* ( [^:]*? ) )?
                             \s* \Z
                        3085
                        3086 }
                           \BNVS_set:cpn { query_eval_end_return_true: } {
                              \group_end:
                              \prg_return_true:
                        3089
                        3090 }
                        3091 \BNVS_new:cpn { query_eval_end_return_false: } {
                             \BNVS end:
                        3092
                              \prg_return_false:
                        3093
```

```
3094 }
   \BNVS_new:cpn { query_eval_end_return_false:n } #1 {
3095
      \BNVS end:
3096
      \prg_return_false:
3097
3098 }
    \BNVS_new:cpn {    query_eval_error_end_return_false:n } #1 {
3099
      \__bnvs_error:x { #1 }
3100
      \__bnvs_query_eval_end_return_false:
3101
3102 }
    \BNVS_new:cpn { query_eval_unreachable: } {
3103
      \__bnvs_query_eval_error_end_return_false:n { UNREACHABLE }
3104
3105
    \BNVS_new:cpn { if_blank:cTF } #1 {
3106
      \BNVS_tl_use:Nc \tl_if_blank:VTF { #1 }
3107
3108 }
    \BNVS_new_conditional:cpnn { match_pop_left:c } #1 { T, F, TF } {
3109
      \BNVS_tl_use:nc {
3110
        \BNVS_seq_use:Nc \seq_pop_left:NNTF { match }
3111
      } { #1 } {
3112
        \prg_return_true:
3113
      } {
3114
3115
        \prg_return_false:
      }
3116
3117 }
```

\\_\_bnvs\_query\_eval\_match\_branch:  $\overline{\mathit{TF}}$  \\_\_bnvs\_query\_eval\_match\_branch: TF  $\{\langle \mathit{true}\ \mathit{code} \rangle\}$   $\{\langle \mathit{false}\ \mathit{code} \rangle\}$ 

Puts the proper items of \l\_\_bnvs\_match\_seq in \l\_\_bnvs\_V\_tl, \l\_\_bnvs\_TEST\_A\_tl, \l\_\_bnvs\_Z\_tl, \l\_\_bnvs\_L\_tl then branches accordingly on one of the returning \\_\_bnvs\_query\_eval\_return[\langle description \rangle]: functions. All these functions properly set the ...ans\_tl variable and they end with either \prg\_return\_true: or \prg\_return\_false:. This is not inlined for readability.

```
\BNVS_new_conditional:cpnn { query_eval_match_branch: } { T, F, TF } {
3118
      \__bnvs_match_pop_left:cT V {
3119
        \__bnvs_match_pop_left:cT V {
3120
          \_\_bnvs_if_blank:cTF V {
3121
             \__bnvs_match_pop_left:cT A {
               \__bnvs_match_pop_left:cT Z {
3123
                 \__bnvs_match_pop_left:cT L {
3124
                   \__bnvs_if_blank:cTF A {
3125
                     \__bnvs_match_pop_left:cT L {
3126
                       \__bnvs_match_pop_left:cT Z {
3127
                         \_bnvs_if_blank:cTF Z {
3128
                            \__bnvs_if_blank:cTF L {
3129
                              \__bnvs_match_pop_left:cT Z {
3130
                                \__bnvs_match_pop_left:cT L {
3131
                                  \__bnvs_if_blank:cTF L {
3132
                                    \__bnvs_if_blank:cTF Z {
3133
                                       \BNVS_use:c { query_eval_return[:]: }
3134
                                    } {
3135
                                       \BNVS_use:c { query_eval_return[:Z]: }
3136
3137
                                  } {
3138
```

```
\__bnvs_if_blank:cTF Z {
3139
    3140
                                } {
3141
                                   \BNVS_use:c { query_eval_return[:Z::L]: }
3142
3143
3144
3145
                           }
3146
                         } {
    }
3149
                       } {
3150
                           _bnvs_if_blank:cTF L {
3151
                           \__bnvs_query_eval_unreachable:
3152
3153
                           \BNVS_use:c { query_eval_return[:Z::L]: }
3154
3155
3156
                    }
3157
                   }
3158
                 } {
                   \__bnvs_if_blank:cTF Z {
3160
                     \__bnvs_if_blank:cTF L {
3161
                       \BNVS_use:c { query_eval_return[A:]: }
3162
                    } {
3163
                       \BNVS_use:c { query_eval_return[A::L]: }
3164
                     }
3165
                   } {
3166
                     \__bnvs_if_blank:cTF L {
3167
                       \BNVS_use:c { query_eval_return[A:Z]: }
3168
                     } {
                       \__bnvs_query_eval_error_end_return_false:n {
3171
                         Only~two~of~first,~last~or~length
3172
3173
3174
3175
3176
3177
           }
3178
         } {
3179
           \BNVS_use:c { query_eval_return[V]: }
         }
3181
       }
3182
     }
3183
3184 }
   \BNVS_new:cpn { query_eval_return[V]: } {
3185
Single value
     \__bnvs_if_resolve:vcTF { V } { ans } {
3186
       \prg_return_true:
3187
     } {
3188
       \prg_return_false:
3189
```

```
3191 }
3192 \BNVS_new:cpn { query_eval_return[A:Z]: } {
\P \langle first \rangle : \langle last \rangle range
       \__bnvs_if_resolve:vcTF { A } { ans } {
         \__bnvs_tl_put_right:cn { ans } { - }
3194
         \__bnvs_if_append:vcTF { Z } { ans } {
3195
           \prg_return_true:
3196
        } {
3197
           \prg_return_false:
3198
3199
3200
         \prg_return_false:
3201
3202
3203 }
    \BNVS_new:cpn { query_eval_return[A::L]: } {
\P \langle first \rangle :: \langle length \rangle range
       \_bnvs_if_resolve:vcTF { A } { A } {
3205
         \_bnvs_if_resolve:vcTF { L } { ans } {
3206
           \__bnvs_tl_put_right:cn { ans } { + }
3207
           \__bnvs_tl_put_right:cv { ans } { A }
3208
           \__bnvs_tl_put_right:cn { ans } { -1 }
           \__bnvs_round_ans:
           \__bnvs_tl_put_left:cn { ans } { - }
3211
           \__bnvs_tl_put_left:cv { ans } { A }
3212
           \prg_return_true:
3213
3214
           \prg_return_false:
3215
3216
3217
      }
3218
         \prg_return_false:
3219
      }
3220 }
    \BNVS_new:cpn { query_eval_return[A:]: } {
3221
\P \langle first \rangle: and \langle first \rangle:: range
       \__bnvs_if_resolve:vcTF { A } { ans } {
         \__bnvs_tl_put_right:cn { ans } { - }
3223
         \prg_return_true:
3224
3225
      } {
         \prg_return_false:
3226
      }
3227
3228 }
    \BNVS_new:cpn { query_eval_return[:Z::L]: } {
■ :Z::L or ::L:Z range
       \__bnvs_if_resolve:vcTF { Z } { Z } {
3230
         \__bnvs_if_resolve:vcTF { L } { ans } {
3231
           \_bnvs_tl_put_left:cn { ans } { 1-}
3232
           \__bnvs_tl_put_right:cn { ans } { + }
3233
           \__bnvs_tl_put_right:cv { ans } { Z }
3234
           \__bnvs_round_ans:
3235
           \__bnvs_tl_put_right:cn { ans } { - }
3236
```

```
\__bnvs_tl_put_right:cv { ans } { Z }
3237
           \prg_return_true:
3238
         } {
3239
           \prg_return_false:
3240
3241
      } {
3242
         \prg_return_false:
3243
3244
    \BNVS_new:cpn { query_eval_return[:]: } {
\blacksquare: or :: range
       \__bnvs_tl_set:cn { ans } { - }
3247
       \prg_return_true:
3248
3249 }
3250 \BNVS_new:cpn { query_eval_return[:Z]: } {
\blacksquare ::\langle last \rangle range
       \__bnvs_tl_set:cn { ans } { - }
3251
       \__bnvs_if_append:vcTF { Z } { ans } {
3252
         \prg_return_true:
3253
      } {
3254
3255
         \prg_return_false:
      }
3256
3257 }
    \BNVS_new_conditional:cpnn { query_eval:nc } #1 #2 { T, F, TF } {
3258
       \__bnvs_call_greset:
3259
       \__bnvs_match_once:NnTF \c__bnvs_A_cln_Z_regex { #1 } {
3260
         \BNVS_begin:
3261
         \__bnvs_query_eval_match_branch:TF {
3262
3263
           \BNVS_end_tl_set:cv { #2 } { ans }
3264
           \prg_return_true:
         } {
3265
           \BNVS_end:
           \prg_return_false:
3267
         }
3268
      } {
3269
Error
         \__bnvs_error:n { Syntax~error:~#1 }
3270
         \prg_return_false:
3271
      }
3272
3273 }
```

```
\verb|\clustriangle| = $$ \sum_{\substack{b \in \mathbb{N} \\ \text{overlay query } list}} \langle tl \ variable \rangle
```

This is called by the *named overlay specifications* scanner. Evaluates the comma separated list of  $\langle overlay \ query \rangle$ 's, replacing all the named overlay specifications and integer expressions by their static counterparts by calling  $\\_$ \_bnvs\_query\_eval:nc, then append the result to the right of the  $\langle tl \ variable \rangle$ . This is executed within a local group. Below are local variables and constants used throughout the body of this function.

\l\_\_bnvs\_query\_seq

Storage for a sequence of  $\langle query \rangle$ 's obtained by splitting a comma separated list.

\l\_\_bnvs\_ans\_seq

Storage of the evaluated result.

```
(End\ definition\ for\ \l_bnvs_ans_seq.)
```

\c\_\_bnvs\_comma\_regex

Used to parse slide range overlay specifications.

In this main evaluation step, we evaluate the integer expression and put the result in a variable which content will be copied after the group is closed. We authorize comma separated expressions and  $\langle first \rangle :: \langle last \rangle$  range expressions as well. We first split the expression around commas, into \l\_query\_seq.

```
regex_split:NnN \c__bnvs_comma_regex { #1 } \l__bnvs_query_seq
```

Then each component is evaluated and the result is stored in \l\_\_bnvs\_ans\_seq that we have clear before use.

We have managed all the comma separated components, we collect them back and append them to the tl variable.

\BeanovesEval

```
\BeanovesEval \ [\langle tl \ variable \rangle] \ \{\langle overlay \ queries \rangle\}
```

 $\langle overlay \; queries \rangle$  is the argument of ?(...) instructions. This is a comma separated list of single  $\langle overlay \; query \rangle$ 's.

This function evaluates the  $\langle overlay\ queries \rangle$  and store the result in the  $\langle tl\ variable \rangle$  when provided or leave the result in the input stream. Forwards to  $\_$ \_bnvs\_eval:nN within a group.  $\.$ ..ans\_tl is used locally to store the result.

```
\NewDocumentCommand \BeanovesEval { O{} m } {
      \BNVS_begin:
3295
      \keys_define:nn { BeanovesEval } {
        in:N .tl_set:N = \l__bnvs_BeanovesEval_tl,
        in:N .initial:n = { },
3298
        see .bool_set:N = \l__bnvs_BeanovesEval_bool,
3299
       see .default:n = true,
3300
        see .initial:n = false,
3301
3302
      \keys_set:nn { BeanovesEval } { #1 }
3303
      \__bnvs_tl_clear:c { ans }
3304
      \__bnvs_eval:nc { #2 } { ans }
3305
      \__bnvs_tl_if_empty:cTF { BeanovesEval } {
        \bool_if:nTF { \l__bnvs_BeanovesEval_bool } {
          \BNVS_tl_use:Nv \BNVS_end: { ans }
       } {
          \BNVS_end:
       }
3311
     } {
3312
        \bool_if:nTF { \l__bnvs_BeanovesEval_bool } {
3313
          \cs_set:Npn \BNVS_end:Nn ##1 ##2 {
3314
            \BNVS_end:
3315
            \tl_set:Nn ##1 { ##2 }
3316
            ##2
          }
3319
          \BNVS_tl_use:nv {
            \exp_last_unbraced:NV \BNVS_end:Nn \l__bnvs_BeanovesEval_tl
          } { ans }
3321
        } {
3322
          \cs_set:Npn \BNVS_end:Nn ##1 ##2 {
3323
            \BNVS_end:
3324
            \tl_set:Nn ##1 { ##2 }
3325
          \BNVS_tl_use:nv {
            \exp_last_unbraced:NV \BNVS_end:Nn \l__bnvs_BeanovesEval_tl
          } { ans }
3329
3330
     }
3331
3332 }
```

## 6.13.13 Reseting counters

\BeanovesReset \BeanovesReset\*

```
\beanovesReset [\langle first\ value \rangle] \{\langle key \rangle\} \beanovesReset* [\langle first\ value \rangle] \{\langle key \rangle\}
```

Forwards to \\_bnvs\_v\_greset:nnF or \\_bnvs\_greset\_all:nnF when starred.

```
\_{\rm bnvs\_name\_id\_n\_get:nTF} { #3 } {
3334
       \BNVS_tl_use:nv {
3335
         \IfBooleanTF { #1 } {
3336
           \__bnvs_greset_all:nnF
3337
3338
           \__bnvs_v_greset:nnF
       } { key } { #2 } {
          \__bnvs_warning:n { Unknown~name:~#3 }
3342 %
3343
     } {
3344
       \__bnvs_warning:n { Bad~name:~#3 }
3345
3346
     \ignorespaces
3347
3348 }
   \makeatother
3350 \ExplSyntaxOff
   <*internal> </internal>
```