Modern Optimization

Jean-Luc Bouchot

School of Mathematics and Statistics Beijing Institute of Technology jlbouchot@bit.edu.cn

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Geometric programming

Outline

Geometric programming

Definition 2.1

Let $f_0: \mathbb{R}^d_{>0} o \mathbb{R}$ be defined as

$$f_0(x) := \sum_{j=1}^m c_j \left(\prod_{i=1}^d x_i^{a_{ij}} \right)$$

for some coefficients $c_j>0$ and some real exponents a_{ij} . f_0 is called a posynomial.

Definition 2.2

A geometric programming problem is an minimization problem in which the objective function is a posynomial and constrained to positive variables.

For the minimization problem to be meaningful, we require some of the a_{ij} to be negative.

Definition 2.3 (DGP)

The dual geometric program to a GP problem is the following optimization problem

$$\text{Maximize}_{\delta \in \mathbb{R}^m} g_0(\delta) \tag{2.1}$$

subject to
$$\sum_{j=1}^{m} \delta_j = 1$$
 (2.2)

and
$$\sum_{j=1}^{m} a_{ij} \delta_j = 0 \tag{2.3}$$

where we define

$$g_0(x) := \prod_{j=1}^m \left(\frac{c_j}{\delta_j}\right)^{\delta_j}.$$

If we cannot find such a δ we say that the DGP is **inconsistent**.

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Definition 2.4 (Duality gap)

Assume both the primal problem (GP) and the dual problem (DGP) have solutions and let p^* and d^* denote these solutions respectively. The **duality** gap is the difference between the two optimal values: $p^* - d^*$.

Theorem 2.1

If $x^* > 0$ minimizes $f_0(x)$, then DGP is consistent. Moreover

$$\delta_j^* := \frac{u_j(x^*)}{f_0(x^*)}$$

is feasible and solves the DGP. Finally, the duality gap is 0, i.e.

$$f_0(x^*) = g_0(\delta^*).$$

The solution is implemented in two steps:

- Find a solution to the dual problem (MART: Multiplicative algebraic reconstruction technique)
- ② Remembering that $\delta_j^* = \frac{u_j(x^*)}{g_0(\delta^*)}$, we can solve for x^* (probably use some logs to help here!)

Remark 2.3 (MART Algo – setup)

The MART algorithm is used to find nonnegative solutions to linear systems

$$Ax = b$$
,

where $A \in \mathbb{R}^{I \times J}$ has nonnegative entries and b has positive entries. Moreover, we request that

$$s_j = \sum_{i=1}^{I} A_{i,j} > 0.$$

Definition 2.5 (The MART Algo)

The MART algorithm is an iterative algo with the following steps:

- Start with a positive vector x^0
- **2** Iterate starting with k = 0, setting $i = k \pmod{I} + 1$.
- Increment $k: k \leftarrow k+1$

Definition 2.6 (KL Divergence)

The (non-symmetric) Kullback-Leibler divergence between two positive numbers a and b is defined as

$$KL(a,b) = a\log(a/b) + b - a.$$

It is extended to 0 numbers as $KL(a,0)=+\infty$ and KL(0,b)=b. For J dimensional vectors, it is defined as

$$KL(a,b) = \sum_{j=1}^{J} KL(a_j, b_j).$$

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Proposition 2.1 (admitted)

Let x^0 be a given non negative vector and A be a consistent matrix for the MART algorithm. Then the sequence of iterates $\{x^k\}_k$ converges to the nonnegative solution of b=Ax which minimizes $KL(x,x^0)$.

Since the a_{ij} may be negative, we cannot apply MART directly:

- ① Lift all the equations in the system as much as needed adding a constant times the last row of the matrix. This generates a system $B\delta=\widetilde{b}$.
- ② Solve $B\delta = \widetilde{b}$ which is equivalent to finding a solution to $A\delta = b$. (why?)
- **3** Set the starting point $\delta^0 = c = (c_1, \dots, c_m)^T$.

 δ^k converges to a solution which minimizes

$$KL(\delta, c) = \sum_{i=1}^{m} \delta_i \log \left(\frac{\delta_i}{c_i}\right) + c_i - \delta_i$$
$$= -\sum_{i=1}^{m} \delta_i + \sum_i c_i - \sum_{i=1}^{m} \delta_i \log \left(\frac{c_i}{\delta_i}\right)$$
$$= -1 + \sum_i c_i - \log \left(g_0(\delta)\right).$$

Example 2.1 (in-class)

Find the minimum to the following function in 3 dimensions

 $f_0(x) = \frac{40}{x_1x_2x_3} + 20x_1x_3 + 10x_1x_2 + 40x_2x_3$ for strictly positive entries.

- Compute the first 2-3 iterations of MART
- ② Assuming MART converges to $\delta^* = (0.4, 0.2, 0.2, 0.2)^T$, what are the optimal values for x^* and the minimal cost?

Exercise 2.1

Minimize the function $f_0(x,y) = \frac{1}{xy} + xy + x + y$ for x > 0 and y > 0. (You might need some help of a computer.)