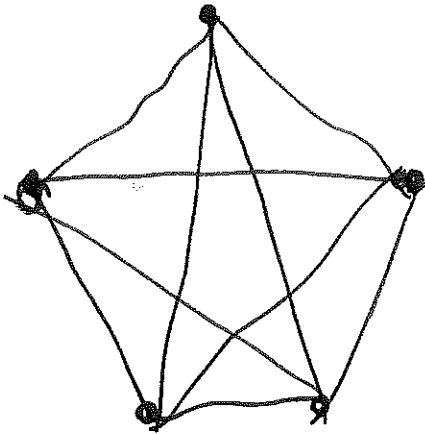


Graph vocab

1. Complete graph on n vertices
 $n=5$



how many edges?

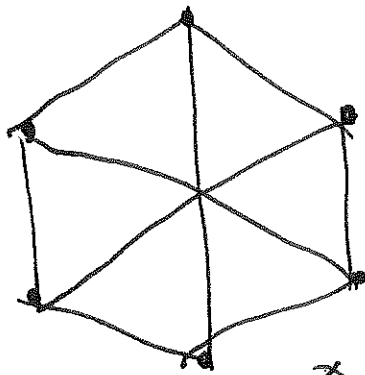
$$\binom{n}{2}$$

2. Regular graph: every vertex has the same degree.

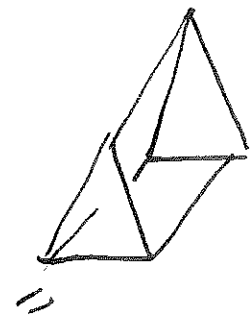
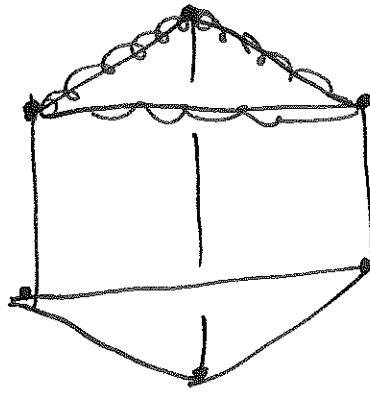
(complete graph is regular)

Puzzle: find 3-regular ^{every vertex has 3 edges} graphs with 6 vertices?
One? two? all?

7 vertices?



$K_{3,3}$

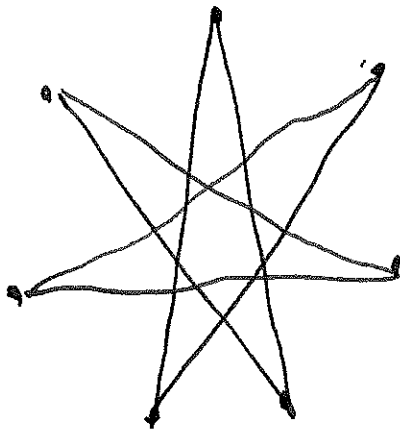


Why different?

contains loops of length 3!
doesn't.

7 vertices

it's ^{not} possible!



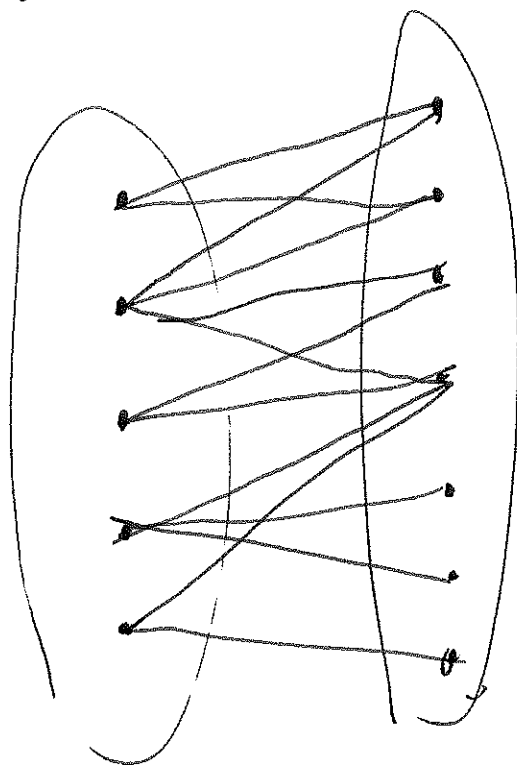
it would need 10.5
edges!

sum of edges

sum of degrees must be
even, but 21 isn't!

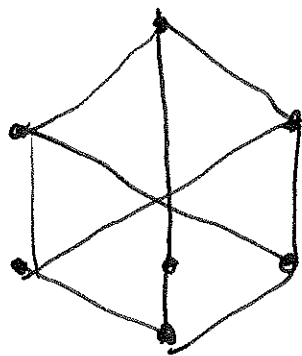
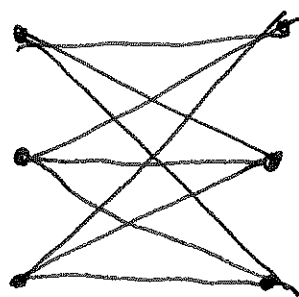
Bipartite graph

G is bipartite: you can split vertices into two sets (think: color some blue, some red) and every edge connected one vertex from each side.

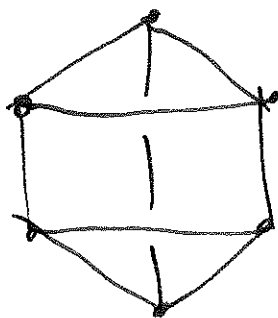


complete bipartite

$K_{m,n}$ = m blue
 n red
make all possible
cross-color
connections



bipartite = $K_{3,3}$



not bipartite
contains a triangle

Thm In a bipartite graph, every cycle has even length.

If a graph has no cycles of odd length, it's automatically bipartite.

→ A graph is Eulerian if there's closed ^{loop.}

trail that touches every edge ~~once~~

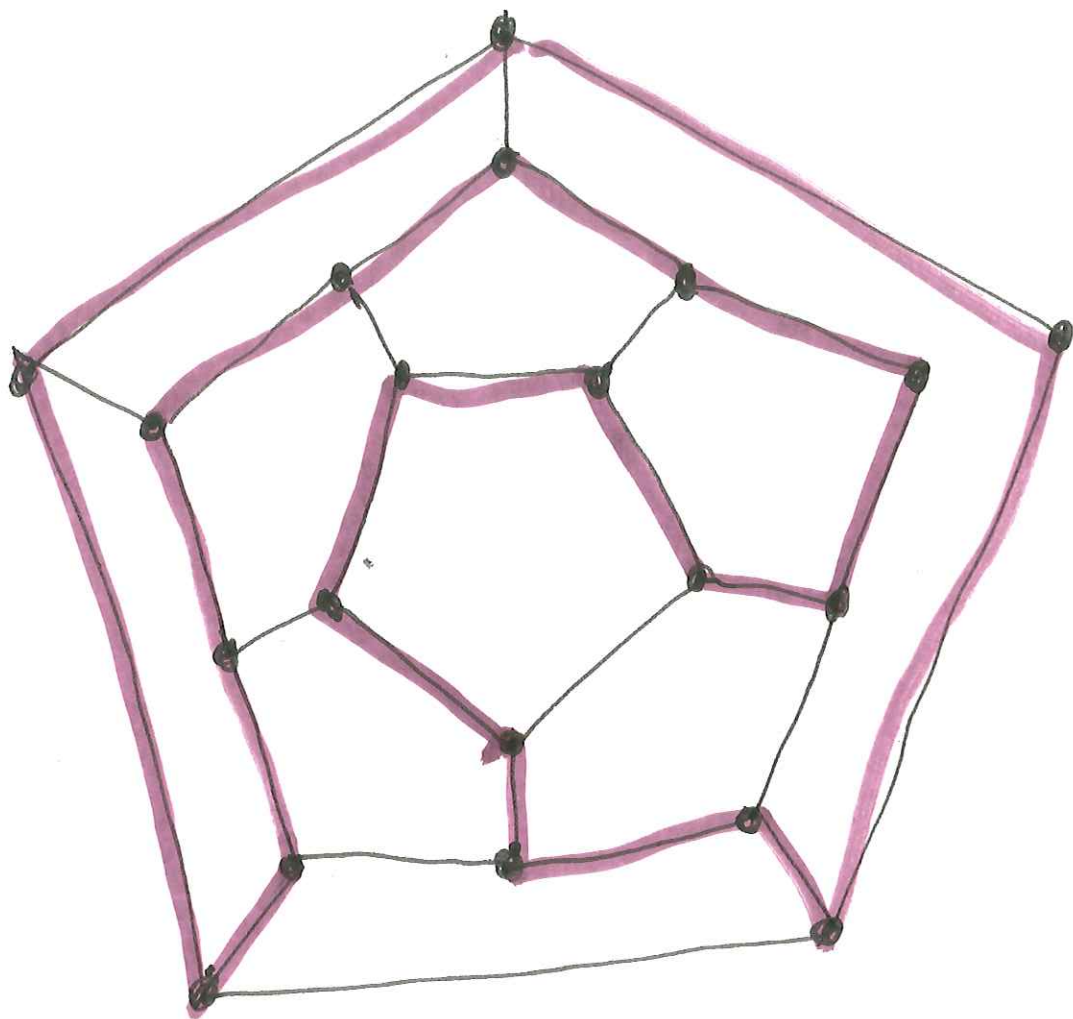
↑
no repeat
edges

(cf Königsberg)

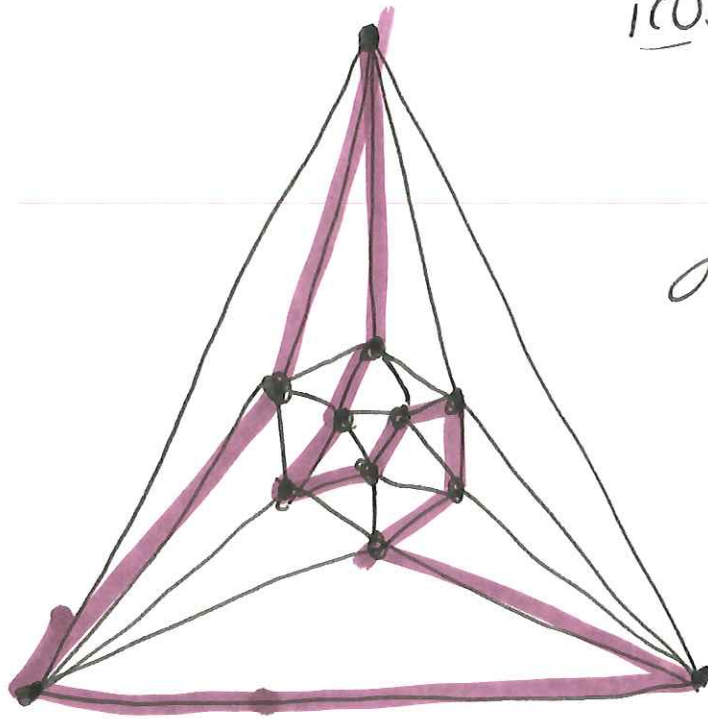
Thm G is Eulerian if and only if every vertex has even degree.

→ A graph is Hamiltonian if there's a closed trail that touches every vertex once.

How to tell?



icosahedral graph!

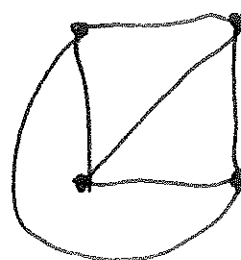
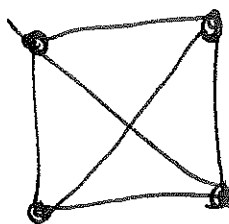


also Hamiltonian

no good algorithm.

Planar graphs:

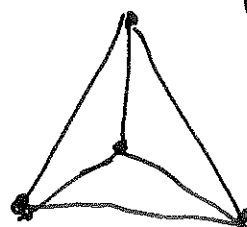
A graph is planar if it can be drawn
where edges don't cross:



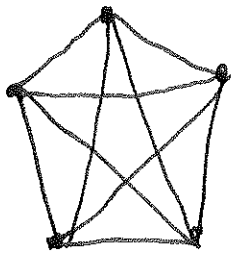
K_4 is planar

Theorem, "Planar" allows ^{curved} ~~bent~~ edges.

In fact any planar graph can be drawn
with straight non-overlapping edges.



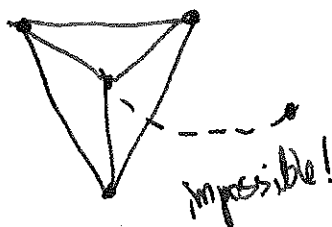
What about K_5 ?



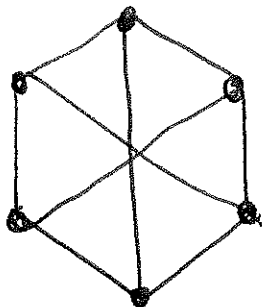
Not planar;

contains K_4 , but need

↙ to add a point!

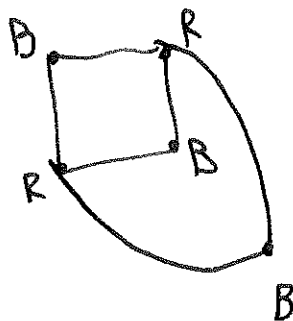


What about $K_{3,3}$?



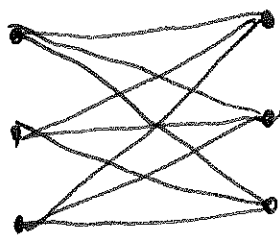
11

Not planar:



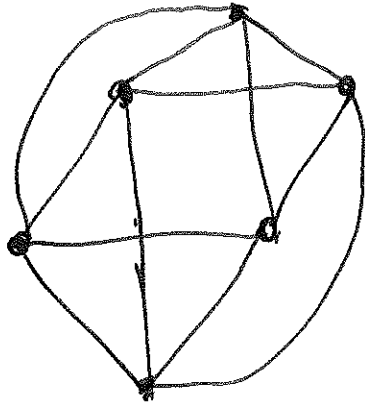
impossible!

I'm convinced.



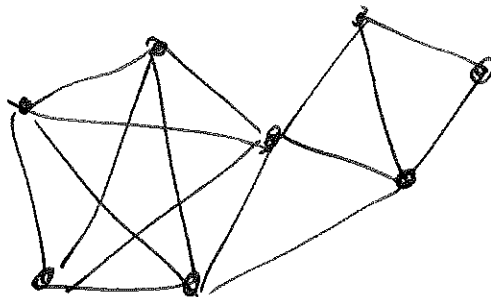
What about O ?

is it planar?



vertices + edges of
octahedron.

Thm \leftarrow "Kuratowski's theorem" G is planar if and only if
it does not
contain a copy of K_5 or $K_{3,3}$.



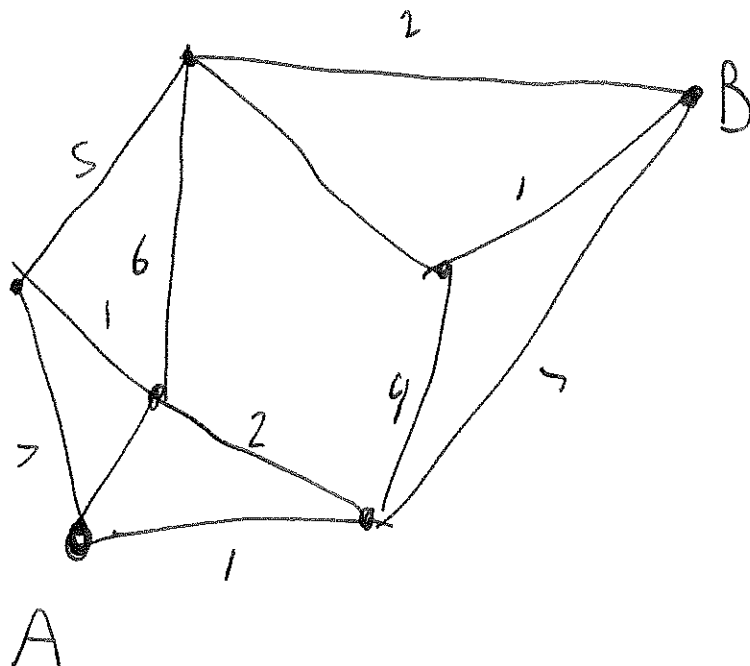
To check planar:

look at every subset of S .
is it K_5 ?

" of G . is it $K_{3,3}$?

Shortest path problem

Suppose G is a graph with a number on each edge ("travel time")

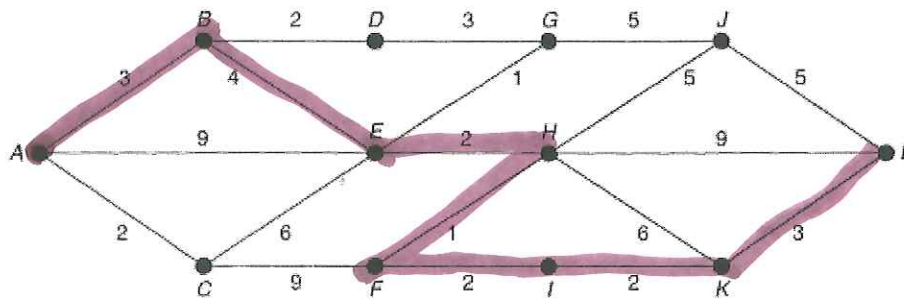


What's the fastest route between A & B?
(path on graph minimizing the sum).

What to do?

Dijkstra's
algorithm.





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