

# CS201 Homework 4

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October 3, 2014

## 1 Book Problems

### 1.1 Exercise 4.4

Using the definition of Big O, show that:

1.  $6n^2 + 3$  is  $O(n^2)$

$$6n^2 + 3 \leq 6n^2 + 3n^2 \text{ for } n \geq 1$$

$$6n^2 + 3 \leq 9n^2 \text{ for } n \geq 1$$

Therefore,

$$6n^2 + 3 \text{ is } O(n^2) \text{ with } c = 9, N=1, g(n) = n^2$$

2.  $n^2 + 17n + 1$  is  $O(n^2)$

$$n^2 + 17n + 1 \leq n^2 + 17n^2 + 1n^2 \text{ for } n \geq 1$$

$$n^2 + 17n + 1 \leq 19n^2 \text{ for } n \geq 1$$

Therefore,

$$n^2 + 17n + 1 \text{ is } O(n^2) \text{ with } c = 19, N=1, g(n) = n^2$$

3.  $5n^3 + 100n^2 - n - 10$  is  $O(n^3)$

$$5n^3 + 100n^2 - n - 10 \leq 5n^3 + 100n^3 - n^3 - 10n^3 \text{ for } n \geq 1$$

$$5n^3 + 100n^2 - n - 10 \leq 94n^3 \text{ for } n \geq 1$$

Therefore,

$$5n^3 + 100n^2 - n - 10 \text{ is } O(n^3) \text{ with } c = 94, N = 1, g(n) = n^3$$

4.  $3n^2 + 2^n$  is  $O(2^n)$

$$3n^2 + 2^n \leq 2^{n+2} \text{ for } n > 4$$

$$2^{n+2} = 4 * 2^n$$

$$3n^2 + 2^n \leq 4 * (2^n)$$

Therefore,

$$3n^2 + 2^n \text{ is } O(2^n) \text{ for } c = 4, N = 4, g(n) = 2^n$$

### 1.2 Exercise 4.9

1. Show that  $7n^2 + 5n$  is not  $O(n)$ .

Proof by contradiction:

Let us assume  $7n^2 + 5n \leq c * n$  - Divide by n

$$7n + 5 \leq c$$

as  $n \Rightarrow \infty, c \geq \infty$  which is impossible.

### 1.3 Exercise 4.12

What is the order of this algorithm?

```
for (int pass = 1; pass <= n; pass++) {  
    for (int index = 0; index < n; index++) {  
        for (int count = 1; count < 10; count++) {  
            . . .  
        }  
    }  
}
```

This algorithm is  $O(n^2)$ , as the first 2 loops require  $n$  operations in the worst case, while the final loop requires a constant of 10 or  $O(1)$

### 1.4 Exercise 4.17

Consider four programs - A, B, C, & D. If each program requires 10 seconds to solve a problem of size 1000, estimate the time required by each program for a problem of size 2000:

$$t = k * n; 10 = k * 1000; k = 1/100;$$

A  $O(\log n)$

.076 seconds.

B  $O(n)$

20 seconds.

C  $O(n^2)$

40,000 seconds.

D  $O(2^n)$

$1.15 * 10^{600}$  seconds.

### 1.5 Exercise 4.18

Suppose that you have a dictionary whose words are not sorted in alphabetical order. As a function of the number,  $n$ , of words, what is the time complexity of searching for a particular word in this dictionary?

The time complexity is  $O(n)$  as in the worst case, the algorithm would have to search the entire array ( $n$  items) to the end to find the word.

### 1.6 Hydra Write-up

1. Using Big O notation, predict the time requirement for this algorithm in terms of the number  $n$  of characters in the initial string.

I estimate that my program will take  $O(n!)$  time. My program creates the child strings (from 'slaying' the larger string) by iterating over each

2. Time the actual execution of the program for various values of  $n$ , and write a chart with your results. (You need not create a plot, simply a series of  $n$  / time pairs will do) For the timing, remove output statements from your program, as simply printing to the screen actually eats up a lot of time.

```
real    293m37.079 s
user    309m28.920 s
sys     1m56.856 s
```