# Midterm project

The goal of this midterm project is to give you a chance to reflect on what you have been exposed to in this course since the beginning of the semester, including self-reference.

# **Expectations**

You are expected:

- to work in groups, as you did for the weekly handins,
- to share your unit tests, all for one and one for all, and to give credit where credit is due in each of your unit-test functions:

```
let test something or other candidate =
  (* original tests: *)
  (candidate ... = ...)
  & &
  . . .
  & &
  (candidate ... = ...)
  (* my own tests: *)
  (candidate ... = ...)
  & &
  . . .
  & &
  (candidate ... = ...)
  (* Athos's tests: *)
  (candidate ... = ...)
  & &
  . . .
  & &
  (candidate ... = ...)
  (* Porthos's tests: *)
  (candidate ... = ...)
  & &
  . . .
  & &
  (candidate ... = ...)
```

```
&&
    (* Aramis's tests: *)
    (candidate ... = ...)
    &&
    ...
    &&
    (candidate ... = ...)
    (* etc. *);;
```

(If one of Athos's tests helped you debug your code, do mention it somewhere, be it in a comment or in the report.)

• to write an individual report, using your own words.

Throughout, remember to embrace the structure:

- 1. textual description of a computation,
- 2. implementation of a unit-test function,
- 3. inductive specification of the computation,
- 4. implementation of this inductive specification as a structurally recursive function, and
- 5. verification that the implementation passes the unit test.

#### Resources

• The OCaml code for the present midterm project (latest version: 23 Sep 2017).

### Part 1

This part is a warmup.

### Question 1.1

Using (part of or all of) Aristotle's four causes, describe:

- a. a microprocessor
- b. an ordinary computer printer
- c. a binary tree

### Question 1.2

a. Given a processor for x86 programs, an interpreter for Scheme programs written in Python, a self-interpreter for Scheme, a self-interpreter for OCaml, an interpreter for OCaml programs written in Scheme, and a compiler from Python to x86 written in x86, can we execute an OCaml program?

Why?

b. Given a processor for x86 programs, a compiler from OCaml to x86 written in Python, a Scheme interpreter written in x86, a self-interpreter for Python, and a compiler from Python to Scheme written in Scheme, and a chocolate bar, can we execute an OCaml program?

Why?

c. Given a processor for x86 programs, a compiler from OCaml to Scheme written in Scheme, and an interpreter for Scheme programs written in x86, can we execute an OCaml program?

Why?

### Question 1.3

The OCaml library function, Random.int, when applied to a positive integer, returns a random integer that is strictly smaller than this positive integer and greater or equal to 0:

```
# Random.int 3;;
- : int = 0
# Random.int 3;;
- : int = 1
# Random.int 3;;
- : int = 1
# Random.int 3;;
- : int = 2
# Random.int 3;;
- : int = 0
# Random.int 3;;
- : int = 2
```

#### Applying Random.int to a non-positive integer is undefined and so an exception is raised:

```
# Random.int 0;;
Exception: Invalid_argument "Random.int".
# Random.int (-1);;
Exception: Invalid_argument "Random.int".
#
```

Assuming that is\_even denotes a function of type int -> bool that computes whether its argument is even,

- a. what is the result of evaluating is\_even (2 \* (Random.int 5))?
  - 1. true, always
  - 2. false, always
  - 3. sometimes true, sometimes false, but an exception is never raised
  - 4. sometimes true, sometimes false, and sometimes no result because an exception is raised
  - 5. no result because an exception is raised, always
  - 6. we can't say, because applying Random.int might not terminate
- b. what is the result of evaluating is even (Random.int (2 \* 5))?
  - 1. true, always
  - 2. false, always
  - 3. sometimes true, sometimes false, but an exception is never raised
  - 4. sometimes true, sometimes false, and sometimes no result because an exception is raised
  - 5. no result because an exception is raised, always
  - 6. we can't say, because applying Random.int might crash
- c. what is the result of evaluating is\_even (2 \* (Random.int 6))?
  - 1. true, always
  - 2. false, always
  - 3. sometimes true, sometimes false, but an exception is never raised
  - 4. sometimes true, sometimes false, and sometimes no result because an exception is raised
  - 5. no result because an exception is raised, always
  - 6. we can't say, because applying Random.int returns a, you know, random result
- d. what is the result of evaluating is\_even (Random.int (2 \* 6))?
  - 1. true, always
  - 2. false, always
  - 3. sometimes true, sometimes false, but an exception is never raised
  - 4. sometimes true, sometimes false, and sometimes no result because an exception is raised

- 5. no result because an exception is raised, always
- 6. we can't say in general, because every call to random accelerates the half-life decay of our silicon processor, which is why our computers have gotten slower and slower during this course; fortunately, there is Moore's law, but then we would need to buy a new computer, and this new computer would also become slower and slower because of all these calls to random; so all in all, we can't say in general
- e. what is the result of evaluating Random.int (Random.int (6 \* 9))?
  - 1. a really random number, always
  - 2. sometimes a random number, sometimes no result because an exception is raised
  - 3. no result because an exception is raised, always
  - 4. no result because compound randomness is uncomputable
  - 5. a random answer to life, the universe, and everything

Briefly justify each of your answers.

### Part 2

The goal of this part is to study the following alternative data type of binary trees, where the payload (an integer) is not in the leaves but in the nodes:

```
type binary_tree' =
    | Leaf'
    | Node' of binary_tree' * int * binary_tree';;
```

The associated induction principle reads as follows:

```
INDUCTION_BINARY_TREE'

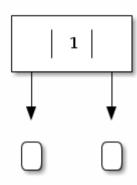
P(Leaf') for all binary trees t1 and t2, for all integers n, P(t1) /\ P(t2) => P(Node' (t1, n, t2))

for all binary trees t, P(t)
```

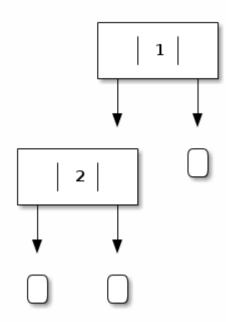
#### For example:

• The binary tree obtained by evaluating Leaf' is depicted as follows:

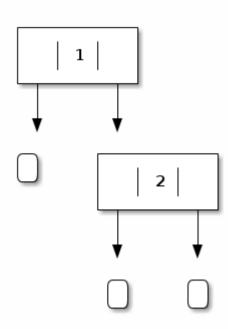
• The binary tree obtained by evaluating Node' (Leaf', 1, Leaf') is depicted as follows:



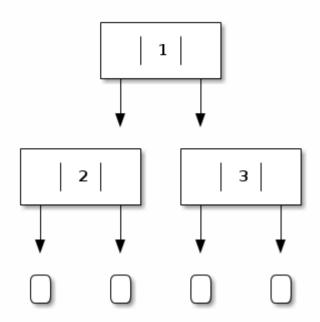
• The binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Leaf') is depicted as follows:



• The binary tree obtained by evaluating Node' (Leaf', 1, Node' (Leaf', 2, Leaf')) is depicted as follows:



• The binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Node' (Leaf', 3, Leaf')) is depicted as follows:



etc.

### Question 2.1

Implement an OCaml function that counts the number of leaves of a given binary tree.

## Question 2.2

Implement an OCaml function that counts the number of nodes of a given binary tree.

# Question 2.3

Is there a relation between the number of leaves and the number of nodes of the same binary tree? If so, state it and prove it.

## Question 2.4

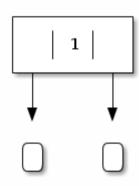
A binary tree is left-balanced if all its right subtrees are leaves.

#### For example:

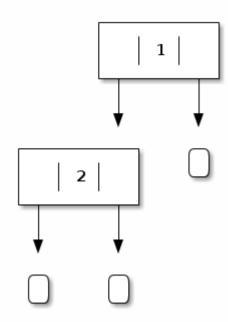
• The binary tree obtained by evaluating Leaf' is vacuously left-balanced:



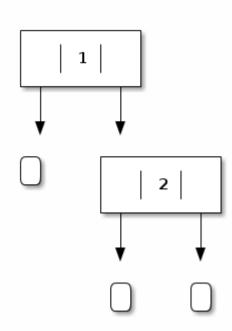
• The binary tree obtained by evaluating Node' (Leaf', 1, Leaf') is left-balanced:



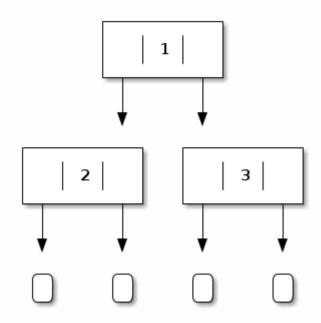
• The binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Leaf') is left-balanced:



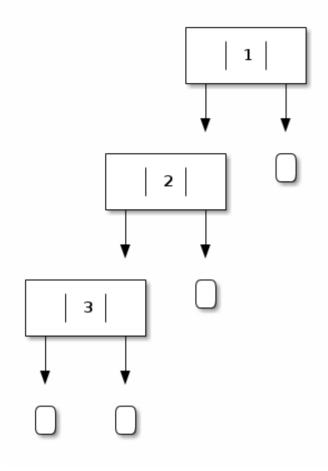
• The binary tree obtained by evaluating Node' (Leaf', 1, Node' (Leaf', 2, Leaf')) is not left-balanced:



• The binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Node' (Leaf', 3, Leaf')) is not left-balanced:



• The binary tree obtained by evaluating Node' (Node' (Node' (Leaf', 3, Leaf'), 2, Leaf'), 1, Leaf') is left-balanced:



etc.

Implement an OCaml function left\_balanced that tests whether a given binary tree is left-balanced.

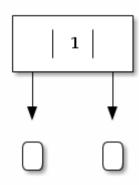
# Question 2.5

A binary tree is right-balanced if all its left subtrees are leaves.

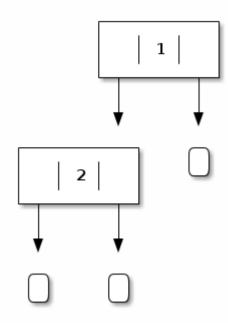
For example:

• The binary tree obtained by evaluating Leaf' is vacuously right-balanced:

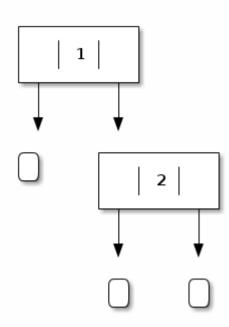
• The binary tree obtained by evaluating Node' (Leaf', 1, Leaf') is right-balanced:



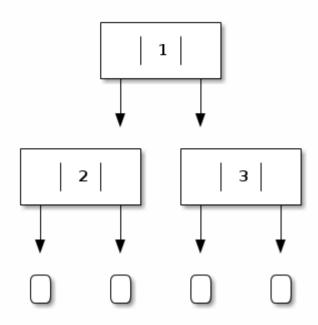
• The binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Leaf') is not right-balanced:



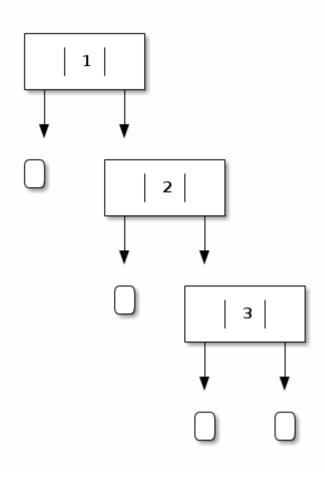
• The binary tree obtained by evaluating Node' (Leaf', 1, Node' (Leaf', 2, Leaf')) is right-balanced:



• The binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Node' (Leaf', 3, Leaf')) is not right-balanced:



• The binary tree obtained by evaluating Node' (Leaf', 1, Node' (Leaf', 2, Node' (Leaf', 3, Leaf'))) is right-balanced:



etc.

Implement an OCaml function right\_balanced that tests whether a given binary tree is right-balanced.

# Question 2.6

How many trees are both left-balanced and right-balanced? If there are finitely many (all with the same integer, e.g., o, in their nodes), enumerate them.

# Part 3

As an alternative to representing left-balanced binary trees as binary trees that satisfy the predicate left\_balanced, let us implement the following dedicated data type, along with the conversion functions from it to binary tree and back:

```
type left binary tree' =
    Left Leaf'
    Left Node' of left binary tree' * int;;
let rec embed left binary tree' into binary tree' t =
     (* embed left binary tree' into binary tree' : left binary tree' -> binary tree' *)
  match t with
  | Left Leaf' ->
     Leaf'
  Left Node' (t1, n) ->
     Node' (embed left binary tree' into binary tree' t1, n, Leaf');;
type option left binary tree' =
    Some left binary tree' of left binary tree'
    None left binary tree';;
let rec project binary tree' into left binary tree' t =
     (* project left binary tree' into binary tree' : binary tree' -> option left binary tree' *)
  match t with
   Leaf' ->
     Some left binary tree' Left Leaf'
   Node' (t1, n, t2) ->
     match t2 with
      Leaf' ->
        (match project binary tree' into left binary tree' t1 with
         | Some left binary tree' t1' ->
            Some left binary tree' (Left Node' (t1', n))
         | None left binary tree' ->
            None left binary tree')
       Node' ->
        None left binary tree';;
```

The associated induction principle reads as follows:

```
INDUCTION_LEFT_BINARY_TREE'

L(Left_Leaf') for all left-balanced binary trees t, for all integers n, L(t) => L(Left_Node' (t, n))

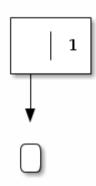
for all left-balanced binary trees t, L(t)
```

For example:

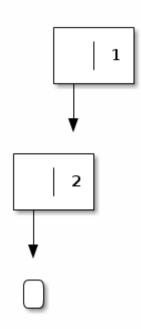
• The binary tree obtained by evaluating <code>Left\_Leaf'</code> is depicted as follows:



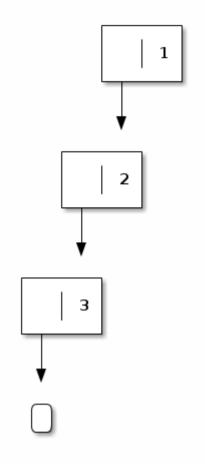
• The binary tree obtained by evaluating Left\_Node' (Left\_Leaf', 1) is depicted as follows:



• The binary tree obtained by evaluating Left\_Node' (Left\_Node' (Left\_Leaf', 2), 1) is depicted as follows:



• The binary tree obtained by evaluating Left\_Node' (Left\_Node' (Left\_Node' (Left\_Leaf', 3), 2), 1) is depicted as follows:



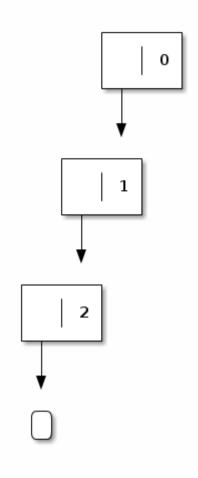
• etc.

## Question 3.1

Why do we need an option type in the co-domain of project\_binary\_tree'\_into\_left\_binary\_tree'?

# Question 3.2 (optional)

Implement an OCaml function left\_nth that indexes a given left-tree at a given depth. For example, consider the following left-tree:



- indexing this tree at depth 0 yields 0,
- indexing this tree at depth 1 yields 1, and
- indexing this tree at depth 2 yields 2.

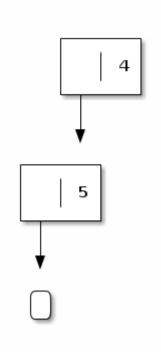
Is your indexing function partial or total?

## Question 3.3

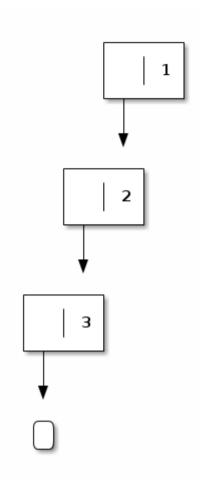
Implement an OCaml function left\_stitch that stitches together two left-balanced binary trees and that satisfies the following unit tests:

```
let test left stitch candidate =
 (* test left stitch : (left binary tree' -> left binary tree' -> left binary tree') -> bool *)
  (candidate Left Leaf'
            Left Leaf'
  = Left Leaf')
  (candidate (Left Node' (Left Leaf', 1))
            Left Leaf'
   = Left Node' (Left Leaf', 1))
  (candidate Left Leaf'
             (Left Node' (Left Leaf', 1))
  = Left Node' (Left Leaf', 1))
  & &
  (candidate (Left Node' (Left Node' (Left Leaf', 2), 1))
            Left Leaf'
   = Left Node' (Left Node' (Left Leaf', 2), 1))
  (candidate Left Leaf'
             (Left Node' (Left Node' (Left Leaf', 2), 1))
   = Left Node' (Left Node' (Left Leaf', 2), 1))
  & &
  (candidate (Left Node' (Left Node' (Left Leaf', 4), 3))
             (Left Node' (Left Node' (Left Leaf', 2), 1), 0))
   = Left Node' (Left Node' (Left Node' (Left Node' (Left Leaf', 4), 3), 2), 1), 0))
  & &
  (candidate (Left Node' (Left Node' (Left Leaf', 4), 3))
             (Left Node' (Left Node' (Left Leaf', 2), 1))
   = Left Node' (Left_Node' (Left_Node' (Left_Leaf', 4), 3), 2), 1))
  & &
  (candidate (Left Node' (Left Node' (Left Leaf', 4), 3))
             (Left Node' (Left Leaf', 2))
   = Left Node' (Left Node' (Left Node' (Left Leaf', 4), 3), 2))
  (candidate (Left Node' (Left Node' (Left Leaf', 4), 3))
             (Left Leaf')
   = Left Node' (Left Node' (Left Leaf', 4), 3))
  (* etc. *);;
```

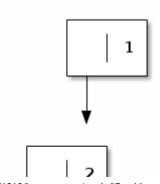
Pictorially, the two left-trees

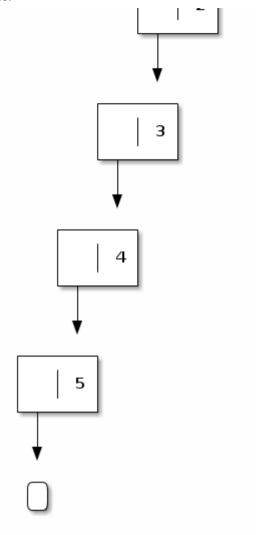


and



are stitched into the following left-tree:





Also, add this pictorial example as a new clause in test\_left\_stitch candidate.

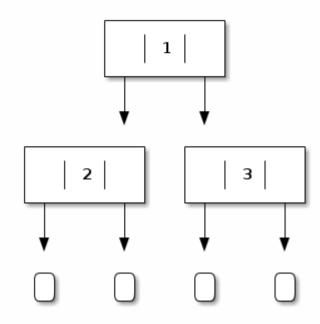
## Question 3.4

Implement an OCaml function <code>left\_rotate</code> that rotates a binary tree to the left and that satisfies the following unit tests:

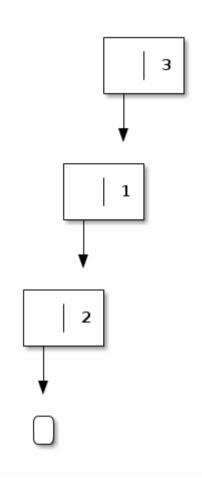
```
let test_left_rotate candidate =
  (* test_left_rotate : (binary_tree' -> left_binary_tree') -> bool *)
```

```
(candidate Leaf'
    = Left Leaf')
& &
(candidate (Node' (Leaf',
                                                                                             1,
                                                                                             Leaf'))
     = Left Node' (Left Leaf',
                                                                          1))
& &
(candidate (Node' (Node' (Leaf',
                                                                                                                                2,
                                                                                                                               Leaf'),
                                                                                            1,
                                                                                             Node'
                                                                                                                           (Leaf',
                                                                                                                                3,
                                                                                                                                Leaf')))
     = Left Node' (Left Node' (Left Leaf',
                                                                                                                                    1),
                                                                          3))
&&
(candidate (Node' (Node' (Leaf',
                                                                                                                                                                   4,
                                                                                                                                                                 Leaf'),
                                                                                                                               2,
                                                                                                                                Node'
                                                                                                                                                             (Leaf',
                                                                                                                                                                   5,
                                                                                                                                                                 Leaf')),
                                                                                             1,
                                                                                             Node' (Node' (Leaf',
                                                                                                                                                                   6,
                                                                                                                                                                 Leaf'),
                                                                                                                                3,
                                                                                                                                Node'
                                                                                                                                                            (Leaf',
                                                                                                                                                                  7,
                                                                                                                                                                 Leaf'))))
     = Left Node' (Left Node' (Left
                                                                                                                                                                                                                                                                                                                                                                                                                                             4),
                                                                                                                                                                                                                                                                                                                                                                                  2),
                                                                                                                                                                                                                                                                                                                     5),
                                                                                                                                                                                                                                                           1),
                                                                                                                                                                                               6),
                                                                                                                                    3),
                                                                          7))
 (* etc. *);;
```

Pictorially, the binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Node' (Leaf', 3, Leaf')), i.e.:



is rotated to the left into the left-tree obtained by evaluating  $Left_Node'$  ( $Left_$ 

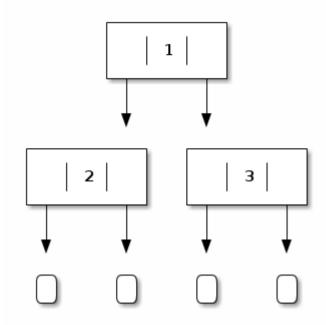


# Question 3.5

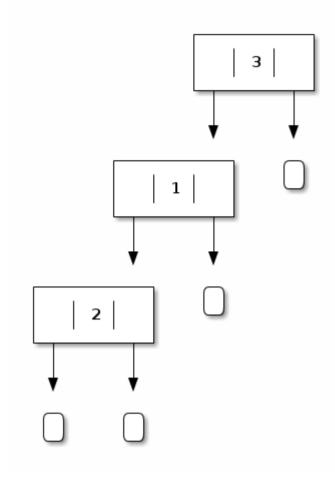
We are now in position to define a function that flattens a binary tree to the left:

```
let left_flatten t =
  (* left_flatten : binary_tree' -> binary_tree' *)
  embed_left_binary_tree'_into_binary_tree' (left_rotate t);;
```

For example, flattening the binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Node' (Leaf', 3, Leaf')), i.e.:



yields the binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Leaf'), 3, Leaf'), i.e.:



Justify why left\_flatten fits the bill and implement a unit-test function for it.

# Part 4

As an alternative to representing right-balanced binary trees as binary trees that satisfy the predicate right\_balanced, let us implement the following dedicated data type, along with the conversion functions from it to binary tree' and back:

```
let rec embed right binary tree' into binary tree' t =
     (* embed right binary tree' into binary tree' : right binary tree' -> binary tree' *)
  match t with
   Right Leaf' ->
     Leaf'
   Right Node' (n, t1) ->
     Node' (Leaf', n, embed right binary tree' into binary tree' t1);;
type option right binary tree' =
    Some right binary tree' of right binary tree'
   None right binary tree';;
let rec project_binary_tree'_into_right binary tree' t =
     (* project binary tree' into right binary tree' : binary tree' -> option right binary tree' *)
  match t with
   Leaf' ->
     Some right binary_tree' Right_Leaf'
   Node' (t1, n, t2) ->
     match t1 with
      Leaf' ->
        (match project binary tree' into right binary tree' t2 with
         | Some right binary tree' t2' ->
            Some right binary tree' (Right Node' (n, t2'))
         | None right binary tree' ->
            None right binary tree')
      Node' ->
        None right binary tree';;
```

The associated induction principle reads as follows:

```
INDUCTION_RIGHT_BINARY_TREE'

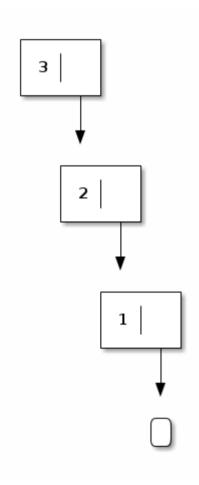
R(Right_Leaf') for all right-balanced binary trees t, for all integers n, R(t) =>
R(Right_Node' (n, t))

for all right-balanced binary trees t, R(t)
```

#### For example:

• The right-tree obtained by evaluating Right Leaf' is depicted as follows:

• The right-tree obtained by evaluating Right\_Node' (3, Right\_Node' (2, Right\_Node' (1, Right\_Leaf'))) is depicted as follows:

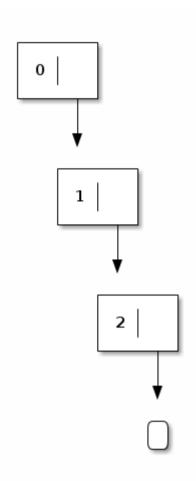


## Question 4.1

Why do we need an option type in the co-domain of project\_binary\_tree'\_into\_right\_binary\_tree'?

# Question 4.2 (optional)

Implement an OCaml function right\_nth that indexes a given right-tree at a given depth. For example, consider the following right-tree:



- indexing this tree at depth 0 yields 0,
- indexing this tree at depth 1 yields 1, and
- indexing this tree at depth 2 yields 2.

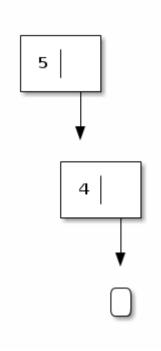
Is your indexing function partial or total?

## Question 4.3

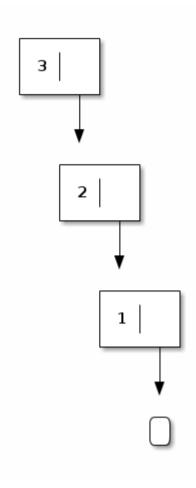
Implement an OCaml function right\_stitch that stitches together two right-balanced binary trees and that satisfies the following unit tests:

```
let test right stitch candidate =
 (* test right stitch : (right binary tree' -> right binary tree' binary tree' -> right binary tree') -> bool *)
  (candidate Right Leaf'
             Right Leaf'
   = Right Leaf')
  &&
  (candidate (Right Node' (1, Right Leaf'))
             Right Leaf'
   = Right Node' (1, Right Leaf'))
  (candidate Right Leaf'
             (Right Node' (1, Right_Leaf'))
   = Right Node' (1, Right Leaf'))
  (candidate (Right Node' (1, Right Node' (2, Right Leaf')))
             Right Leaf'
   = Right Node' (1, Right_Node' (2, Right_Leaf')))
  & &
  (candidate Right Leaf'
             (Right Node' (1, Right Node' (2, Right Leaf')))
   = Right Node' (1, Right Node' (2, Right Leaf')))
  (candidate (Right Node' (1, Right Node' (2, Right Leaf')))
             (Right Node' (3, Right Node' (4, Right Leaf')))
   = Right Node' (1, Right Node' (2, Right Node' (3, Right Node' (4, Right Leaf')))))
  (* etc. *);;
```

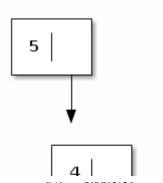
Pictorially, the two right-trees

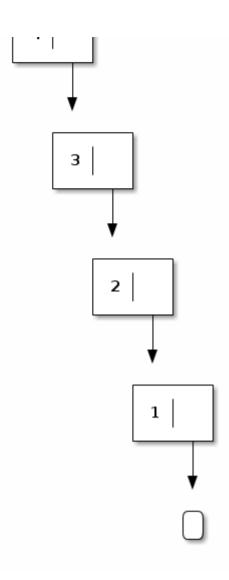


and



are stitched into the following right-tree:





Also, add this pictorial example as a new clause in test\_right\_stitch candidate.

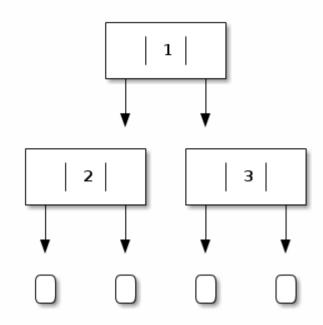
## Question 4.4

Implement an OCaml function right\_rotate that rotates a binary tree to the right and that satisfies the following unit tests:

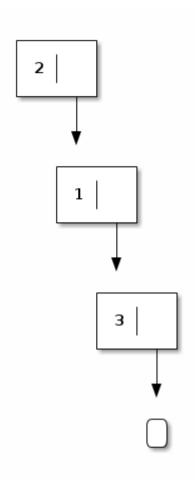
```
let test_right_rotate candidate =
  (* test_right_rotate : (binary_tree' -> right_binary_tree') -> bool *)
```

```
(candidate Leaf'
= Right Leaf')
(candidate (Node' (Leaf',
                   1,
                   Leaf'))
= Right Node' (1,
                Right Leaf'))
&&
(candidate (Node' (Node' (Leaf',
                          2,
                          Leaf'),
                   1,
                   Node'
                         (Leaf',
                          3,
                          Leaf')))
= Right Node' (2,
                Right Node' (1,
                             Right Node' (3,
                                           Right_Leaf'))))
&&
(candidate (Node' (Node' (Leaf',
                                  4,
                                 Leaf'),
                          2,
                          Node'
                                (Leaf',
                                  5,
                                 Leaf')),
                   1,
                   Node'
                         (Node'
                                (Leaf',
                                  6,
                                 Leaf'),
                          3,
                          Node'
                                (Leaf',
                                  7,
                                 Leaf'))))
 = Right Node' (4,
                Right Node' (2,
                             Right_Node' (5,
                                           Right_Node' (1,
                                                        Right Node' (6,
                                                                     Right Node' (3,
                                                                                   Right_Node' (7,
                                                                                                Right_Leaf')))))))
(* etc. *);;
```

Pictorially, the binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Node' (Leaf', 3, Leaf')), i.e.:



is rotated to the right into the right-tree obtained by evaluating Right\_Node' (2, Right\_Node' (1, Right\_Node' (3, Right\_Leaf'))), i.e.:

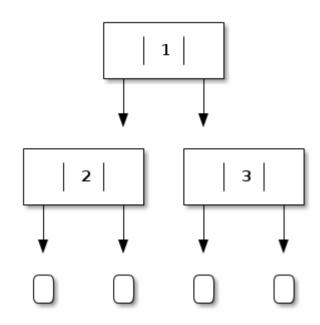


## Question 4.5

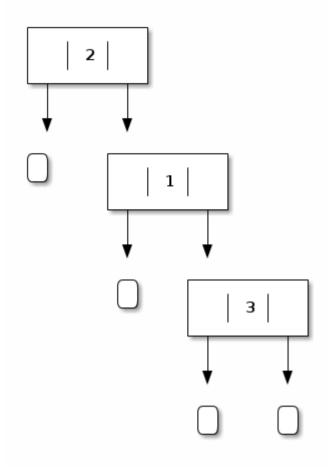
We are now in position to define a function that flattens a binary tree to the right:

```
let right_flatten t =
  (* right_flatten : binary_tree' -> binary_tree' *)
  embed_right_binary_tree'_into_binary_tree' (right_rotate t);;
```

For example, flattening the binary tree obtained by evaluating Node' (Node' (Leaf', 2, Leaf'), 1, Node' (Leaf', 3, Leaf')), i.e.:



yields the binary tree obtained by evaluating Node' (Leaf', 2, Node' (Leaf', 1, Node' (Leaf', 3, Leaf'))), i.e.:



Justify why right\_flatten fits the bill and implement a unit-test function for it.

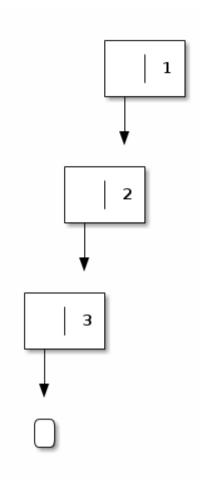
## Part 5

The goal of this part is to study how to rebalance left-trees into right-trees and vice-versa.

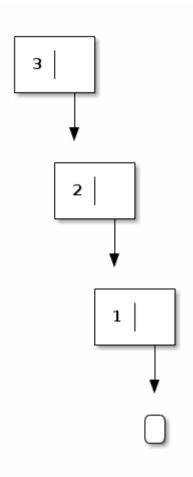
## Question 5.1

Implement an OCaml function that maps a left-balanced tree to a right-balanced tree and that satisfies the following unit tests:

For example, the left-tree obtained by evaluating Left\_Node' (Left\_Node' (Left\_Node' (Left\_Leaf', 3), 2), 1):



is rebalanced into the right-tree obtained by evaluating Right\_Node' (3, Right\_Node' (2, Right\_Node' (1, Right\_Leaf'))):

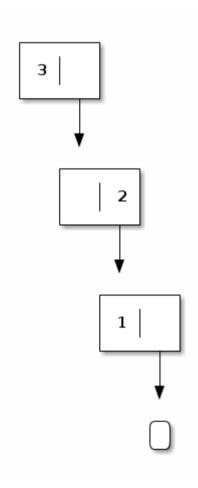


## Question 5.2

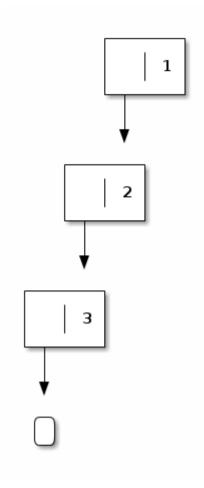
Implement an OCaml function that maps a right-balanced tree to a left-balanced tree and that satisfies the following unit tests:

```
(candidate (Right_Node' (2, Right_Node' (1, Right_Leaf')))
= Left_Node' (Left_Node' (Left_Leaf', 2), 1))
&&
(candidate (Right_Node' (3, Right_Node' (2, Right_Node' (1, Right_Leaf'))))
= Left_Node' (Left_Node' (Left_Leaf', 3), 2), 1))
(* etc. *);;
```

For example, the right-tree obtained by evaluating Right Node' (3, Right Node' (2, Right Node' (1, Right Leaf'))):



is rebalanced into the left-tree obtained by evaluating Left\_Node' (Left\_Node' (Left\_Node' (Left\_Leaf', 3), 2), 1):



# Question 5.3

### Characterize the following two OCaml functions:

#### For example, what are their types?

### Question 5.4

#### Consider the following unit tests:

```
let () = assert (test_left_rotate (fun t -> right_to_left (right_rotate t)));;
let () = assert (test_right_rotate (fun t -> left_to_right (left_rotate t)));;
```

Does your implementation pass these tests? Why?

# The individual report

Your report should include

- a front page with title, name, student number, and date,
- the names of the other members of your group,
- a second page with a table of contents, and
- from the third page and onwards,
  - o an introduction,
  - o a series of sections and subsections reflecting the structure of the project, and
  - o a conclusion where you assess what you did and reflect on how you did it.

Of course you should not paraphrase the code, because what is the point of that.

Pages should be numbered.

An inspiring (and not necessarily humorous, just on topic) quote or three would be welcome.

#### Resources

• The OCaml code for the present midterm project (latest version: 23 Sep 2017).

### Version

Added a handful of clauses in the definition of test\_left\_stitch [29 Sep 2017]

Fixed a typo in Question 2.5, thanks to Shi Tingsheng's eagle eye [28 Sep 2017]

Fixed a typo in Question 4.5, thanks to Shardul Sapkota's eagle eye [27 Sep 2017]

Tightened the narrative, thanks to Kira Kutscher's comments [25 Sep 2017]

Fixed a typo in the notation for Random.int, thanks to Chong Woon Han's eagle eye [24 Sep 2017]

Fine-tuned the narrative [24 Sep 2017]

Created [23 Sep 2017]