

! Try again once you are ready

TO PASS 70% or higher

Try again

GRADE

62.5%

Module 1 Quiz

LATEST SUBMISSION GRADE

62.5%

1. Which of the following primitive operations on threads were introduced in Lecture 1.1?

0.75 / 1 point

Please choose all options that are correct.

! Incorrect

2. Which of the following statements are true, based on what you learned about threads in Lecture 1.1?

0.5 / 1 point

Please choose all options that are correct.

! Incorrect

Please choose all options that are correct.

! Incorrect

4. Consider a parallel program with two structured locks, L1 and L2, (as introduced in Lecture 1.2), and 100 threads that are started in parallel such that each of them executes the following snippet of code:

1 / 1 point

```
1 synchronized(L1) {  
2   synchronized(L2) {  
3     S1  
4   }  
5 }
```

The statement "This program is guaranteed to be free from deadlocks" is:

✓ Correct

5. Which of the following are true regarding unstructured read-write locks introduced in Lecture 1.3?

0.667 / 1 point

Please choose all options that are correct.

-

6. Which of the following best describe the differences between structured and unstructured locks introduced in Lectures 1.2 and 1.3?

0.667 / 1 point

Please choose all options that are correct.

! Incorrect

7. Please classify the following five scenarios, each employing two threads (T1 and T2) running concurrently, as possibly leading to deadlock, livelock or neither. Recall that deadlock and livelock were discussed in Lecture 1.4.

0 / 1 point

Scenario 1:

```
1 // Thread T1  
2 {  
3   . . . // sequential code  
4   T2.join();  
5 }  
6  
7 // Thread T2  
8 {  
9   . . . // sequential code  
10  T1.join();  
11 }
```

Scenario 2: A and B are declared and initialized as distinct objects

```
1 // Thread T1  
2 {  
3   synchronized(A) {  
4     synchronized(B) {  
5       . . . // sequential code  
6     }  
7   }  
8 }  
9
```

```
12 synchronized(A) {  
13   synchronized(B) {  
14     . . . // sequential code  
15   }  
16 }  
17 }
```

Scenario 3: A and B are declared and initialized as distinct objects

```
1 // Thread T1  
2 {  
3   synchronized(A) {  
4     synchronized(B) {  
5       . . . // sequential code  
6     }  
7   }  
8 }  
9  
10 // Thread T2  
11 {  
12   synchronized(B) {  
13     synchronized(A) {  
14       . . . // sequential code  
15     }  
16   }  
17 }
```

Scenario 4: x is declared and initialized as a shared Integer object

```
1 // Thread T1  
2 {  
3   while (x > 4) {  
4     synchronized(x) {  
5       x--;  
6     }  
7   }  
8 }  
9  
10 // Thread T2  
11 {  
12   while (x < 4) {  
13     synchronized(x) {  
14       x++;  
15     }  
16   }  
17 }
```

Scenario 5: x is declared and initialized as a shared Integer object

```
1 // Thread T1  
2 {  
3   while (x < 6) {  
4     synchronized(x) {  
5       x++;  
6     }  
7   }  
8 }  
9  
10 // Thread T2  
11 {  
12   while (x < 8) {  
13     synchronized(x) {  
14       x++;  
15     }  
16   }  
17 }
```

! Incorrect

8. What is the difference between deadlock and livelock? (Recall that deadlock and livelock were discussed in Lecture 1.4.)

1 / 1 point

9. Many algorithms have been proposed to address the Dining Philosophers problem introduced in Lecture 1.5. In this problem, you will evaluate one such algorithm.

0 / 1 point

The following simple algorithm can deadlock with all philosophers holding their left chopsticks.

```
1 // All philosophers:  
2 while (True) {  
3   acquire philosopher's left chopstick  
4   acquire philosopher's right chopstick  
5   eat  
6   release right chopstick  
7   release left chopstick  
8 }
```

So, we develop a slight variant that attempts to avoid this problem — adjacent philosophers pick up the chopsticks in the opposite order. More precisely, we propose separate algorithms for the even-numbered and odd-numbered philosophers. Assume the philosophers are numbered 1, ..., n clockwise around the table.

```
1 // All even-numbered philosophers:  
2 while (True) {  
3   acquire philosopher's left chopstick  
4   acquire philosopher's right chopstick  
5   eat  
6   release right chopstick  
7   release left chopstick  
8 }  
9  
10 // All odd-numbered philosophers:  
11 while (True) {  
12   acquire philosopher's right chopstick  
13   acquire philosopher's left chopstick  
14   eat  
15   release left chopstick  
16   release right chopstick  
17 }
```

What liveness issues could arise if we have an even number of philosophers? What about if we had an odd number of philosophers?

! Incorrect

10. Towards the end of Lecture 1.5, Dr. Sarkar mentioned that deadlock and livelock can be avoided in the Dining Philosophers problem by modifying the "all acquire left fork first" algorithm such that n-1 philosophers attempt to acquire their left fork first, and 1 philosopher attempts to acquire its right fork first. However, nothing was mentioned about the impact of this modification on another important liveness issue: starvation. How could starvation occur, or not occur, with this modification presented at the end of Lecture 1.5?

1 / 1 point