

1 hex digit is 4 bits.

$$16 = 2^4$$

System A block size 16 byte

4 bits

16 elements, each 1 byte.

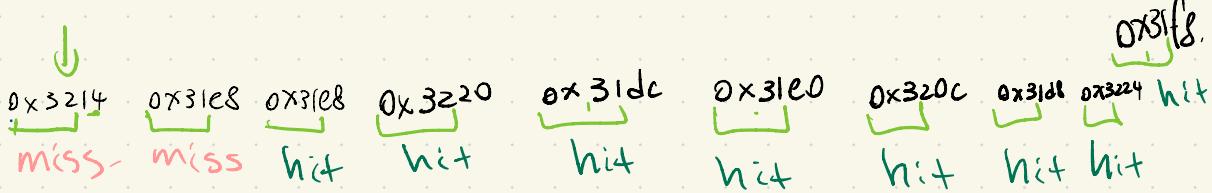
$$\textcircled{D} \quad 0x3214 \rightarrow 0x321f] \\ 0x3210$$

$$\Rightarrow \frac{4}{10} = 0.4.$$

② $0x31e8 \rightarrow 0x31ef$]
③ $0x31e0$
 $0x322f$]
 $0x322f$]

$$B \quad 256 \text{ byte} \Rightarrow 2^8 \cdot 8 \text{ bits}$$

256 elements



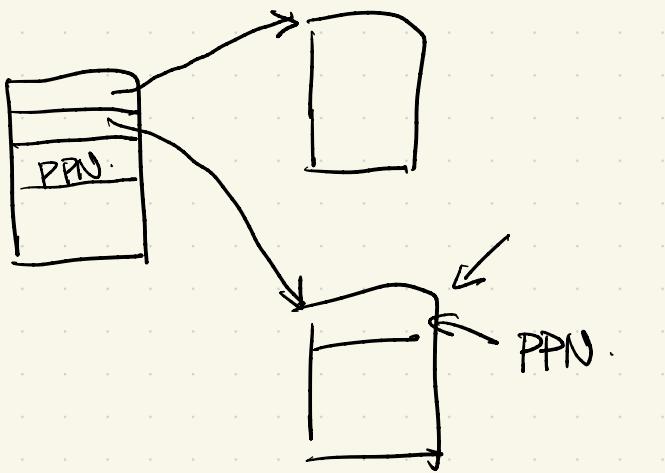
0x32ff] 256 bytes.
0x3200 ΔΔ

$$32 \Rightarrow 2^5.$$

$$\text{hit ratio} = \frac{8}{10} = 0.8.$$

0x3214. $\Rightarrow 0x32$ 0001 0100
 Δ

0x321f] 32 bytes.
0x3200

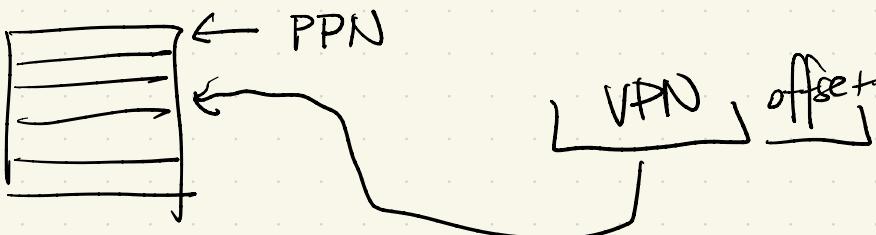


2 level page table.

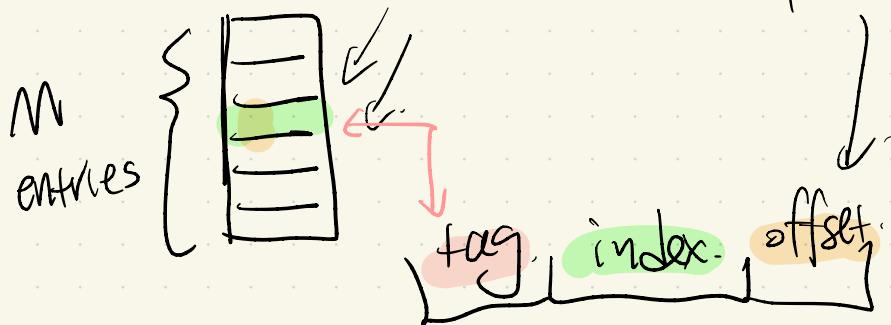
page table size = # entries
PTE

$$\log_2 (\# \text{ entries}) = \text{len of VPN}_i$$

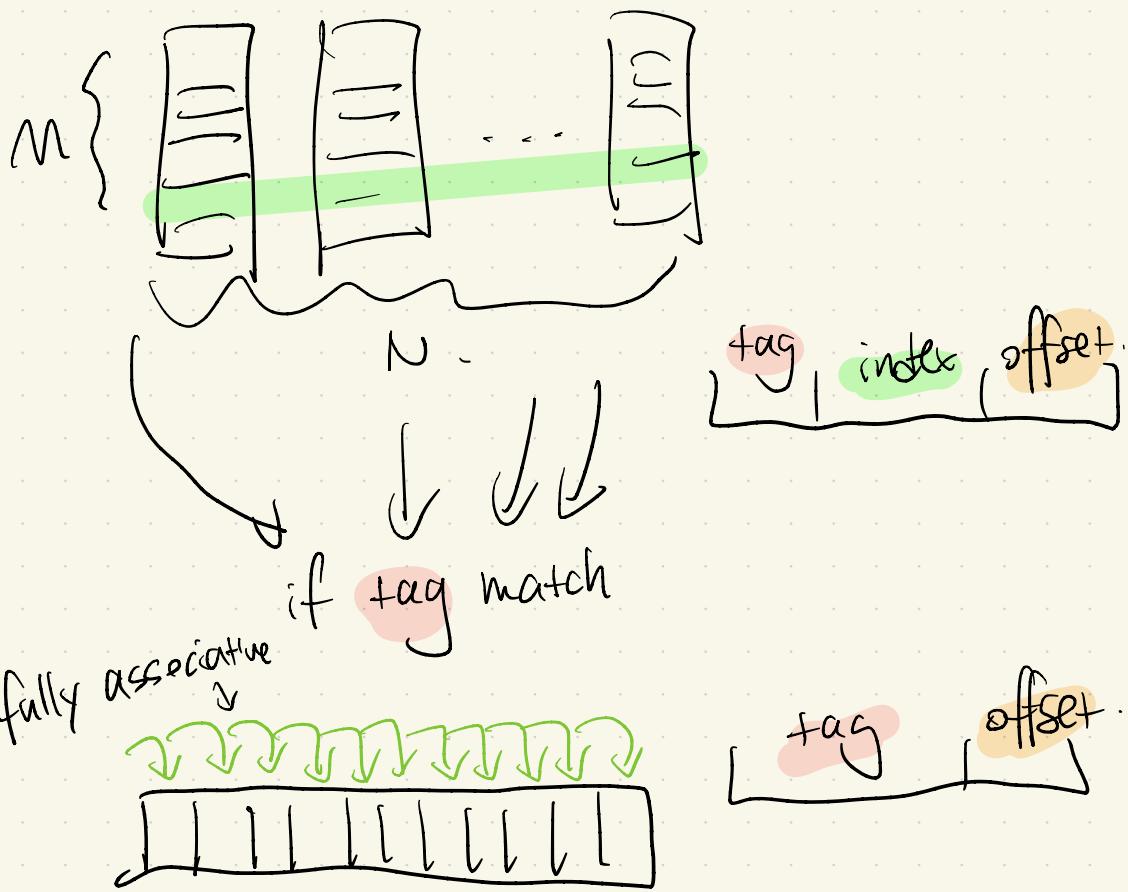
[VPN₀, VPN₁, VPN₂, ..., offset]

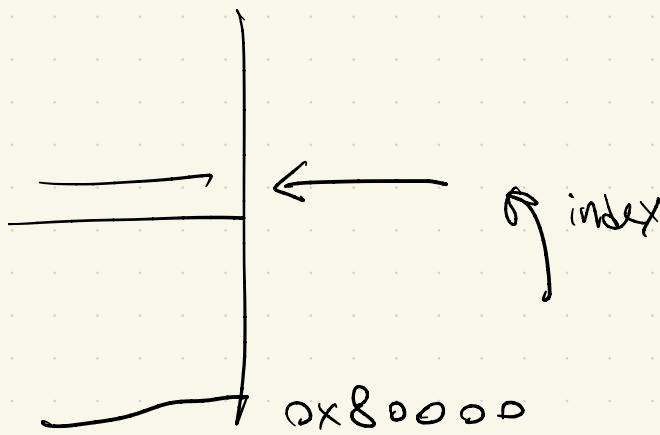
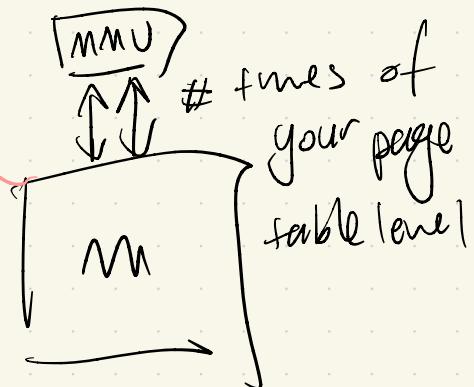
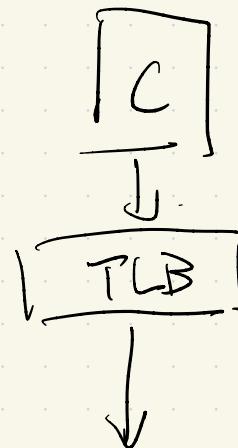
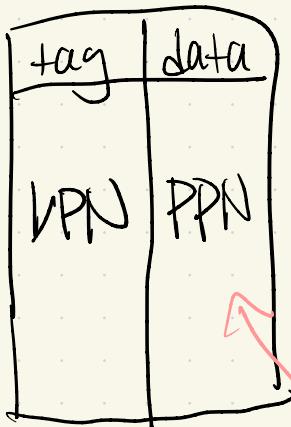


page table size
len(PPN)



N-way association





- The page size is 64KiB
- The physical address size is 32 bits
- The virtual address size is 26 bits
- There is a two level page table (each level is the same width)
- All PTEs are 32 bits.
- The internal PTEs only contains the address for the next level
- The leaf PTEs only contain the PPN, not the full address

SATP = 0x84816c00

Physical address 4 bytes of data

0xfe1f87d4	0x94d7edf8
0xfe1f3c70	0x00002e42
0xa1a6f854	0x0000201e
0xa1a6dd20	0xd8d0c54c
0xa06b9904	0xdc72cb00
0xa06b4868	0x00002558
0x84816c24	0xa1a6f800
③ 0x84816c1c → PTE	0x77cc3c00
0x84816c18	0xfe1f3c00
0x84816c08	0x54728c00
0x84816c04	0xa06b4800
④ 0x77cc3c58 → PPN	0x00000a98
0x77cbf328	0xeb581204
0x64e10590	0xbec813f8
0x5472f62c	0x6da263ec
0x54728c40	0x000039f8
0x44250590	0x805225b4
0x39f80590	0xc7727140
0x2e420590	0xe5af4aa8
0x25580590	0x108e67d4
0x201e0590	0x31094cdc
0x13320590	0xb05fda58
⑤ 0x0a980590 →	0x04950a20
0x01980590	0x81d2a7e4

effective addr.

VPN

0x00f60590 offset

① calculate the offset for page.

$$2^{16} \Rightarrow 16 \text{ bits}$$

② find the VPN

$$\text{len(VAddr)} = 26$$

$$\text{len(Index)} = 26 - 16 = 10 \text{ bits}$$

5 bits 5 bits
VPN0 VPN1

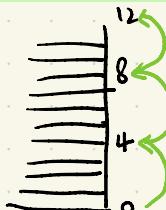
0x00f6
VPN0 VPN1
01111 0110
5 bits 5 bits

↓ ↓
7 22-

28 88
↓ ↓

11100 1011000

VPN₀ = 0x1C VPN₁ = 0x58



* 4 bytes stored Starts that address.

11100 1011000

VPN₀ = 0x1C VPN₁ = 0x58

③ find the first page table entry.

$$SATP = 0x84816c00$$

$$VPN_0 = 0x1c \Rightarrow \underline{0x84816c1c}$$

④ find the physical page number.

$$PTE = 0x77cc3c00$$

$$VPN_1 = 0x58$$

$$\Rightarrow \underline{0x77cc3258}$$

⑤ find Physical address.

$$PPN = 0xa98 .$$

$$\text{offset} = 0x0590$$

$$\begin{array}{l} \text{Padr} \\ \Rightarrow \underline{0xa980590} \end{array}$$

⑥ answer = 0x04950a2d

Q4.

$$t = I \times CPI \times \frac{1}{f}$$

$$t_1 = I \times CPI_1 \times \frac{1}{f_1} \quad 1 \text{ core}$$

$$t_2 = I \times CPI_2 \times \frac{1}{f_2} \quad 2 \text{ cores}$$

$$CPI_2 = CPI_1 \times \left(1 - P + \frac{P}{C}\right)$$

$$\frac{1}{f_2} = \frac{1}{f_1} \times \frac{f_1}{f_2}$$

$$I \times CPI_2 \times \frac{1}{f_2} = I \times CPI_1 \times \left(1 - P + \frac{P}{C}\right) \times \frac{1}{f_1} \times \frac{f_1}{f_2}$$

$$t_2 = t_1 \times \left(1 - P + \frac{P}{C}\right) \times \frac{f_1}{f_2}$$

$$A\bar{M}AT = HT + MP.$$

L1 \rightarrow 2 cycles.

80% hit ratio

Mem \rightarrow 50 cycles

$$80\% \times 2 \text{ cycle.} + 20\%(2 + 50)$$

$$= 80\% \times 2 + 20\% \times 2 + 20\% \times 50.$$

$$= 2 + 20\% \times 50$$