

# Mathematical Foundation of DNN : HW 8

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For  $1 \leq i \leq m/2, 1 \leq j \leq n/2$

$$[\mathcal{T}(X)]_{ij} = \frac{1}{4} (X_{2i-1,2j-1} + X_{2i-1,2j} + X_{2i,2j-1} + X_{2i,2j}) \equiv \frac{1}{4} \text{sum} (X_{2i-1:2i,2j-1:2j})$$

From the equation above, one can observe that  $\mathcal{T}$  is linear operator.

$$\begin{aligned} [\mathcal{T}(X + Y)]_{ij} &= \frac{1}{4} \text{sum} \left( (X + Y)_{2i-1:2i,2j-1:2j} \right) \\ &= \frac{1}{4} \text{sum} (X_{2i-1:2i,2j-1:2j} + Y_{2i-1:2i,2j-1:2j}) \\ &= \frac{1}{4} (\text{sum}(X_{2i-1:2i,2j-1:2j}) + \text{sum}(Y_{2i-1:2i,2j-1:2j})) \\ &= [\mathcal{T}(X)]_{ij} + [\mathcal{T}(Y)]_{ij} \end{aligned}$$

$$\begin{aligned} \mathcal{T}(aX) &= \frac{1}{4} \text{sum} ((aX)_{2i-1:2i,2j-1:2j}) \\ &= a \times \frac{1}{4} \text{sum} ((X)_{2i-1:2i,2j-1:2j}) \\ &= a\mathcal{T}(X) \end{aligned}$$

$$\begin{aligned} &\sum_{i=1}^{m/2} \sum_{j=1}^{n/2} Y_{ij} (\mathcal{T}(\mathcal{X}))_{ij} \\ &= \sum_{i=1}^{m/2} \sum_{j=1}^{n/2} Y_{ij} \frac{1}{4} \text{sum} (X_{2i-1:2i,2j-1:2j}) \\ &= \frac{1}{4} \left[ \sum_{i=1}^{m/2} \sum_{j=1}^{n/2} Y_{ij} X_{2i-1,2j-1} + \sum_{i=1}^{m/2} \sum_{j=1}^{n/2} Y_{ij} X_{2i-1,2j} + \sum_{i=1}^{m/2} \sum_{j=1}^{n/2} Y_{ij} X_{2i,2j-1} + \sum_{i=1}^{m/2} \sum_{j=1}^{n/2} Y_{ij} X_{2i,2j} \right] \\ &= \frac{1}{4} \left[ \sum_{\substack{i=1,3,\dots,m-1 \\ j=1,3,\dots,n-1}} Y_{\frac{i+1}{2}, \frac{j+1}{2}} X_{ij} + \sum_{\substack{i=1,3,\dots,m-1 \\ j=2,4,\dots,n}} Y_{\frac{i+1}{2}, \frac{j}{2}} X_{ij} + \sum_{\substack{i=2,4,\dots,m \\ j=1,3,\dots,n-1}} Y_{\frac{i}{2}, \frac{j+1}{2}} X_{ij} + \sum_{\substack{i=2,4,\dots,m \\ j=2,4,\dots,n}} Y_{\frac{i}{2}, \frac{j}{2}} X_{ij} \right] \\ &= \frac{1}{4} \sum_{i=1}^m \sum_{j=1}^n Y_{\lfloor \frac{i+1}{2} \rfloor, \lfloor \frac{j+1}{2} \rfloor} X_{ij} \end{aligned}$$

Consider the upsampling operator  $\mathcal{U} : \mathbb{R}^{m/2 \times n/2} \rightarrow \mathbb{R}^{m \times n}$ . For  $1 \leq i \leq m, 1 \leq j \leq n$ ,

$$[\mathcal{U}(Y)]_{ij} = \begin{cases} Y_{\frac{i+1}{2}, \frac{j+1}{2}} & \text{if } i = \text{odd}, j = \text{odd} \\ Y_{\frac{i+1}{2}, \frac{j}{2}} & \text{if } i = \text{odd}, j = \text{even} \\ Y_{\frac{i}{2}, \frac{j+1}{2}} & \text{if } i = \text{even}, j = \text{odd} \\ Y_{\frac{i}{2}, \frac{j}{2}} & \text{if } i = \text{even}, j = \text{even} \end{cases}$$

The mapping above can be written in short as follows:  $[\mathcal{U}(Y)]_{ij} = Y_{\lfloor \frac{i+1}{2} \rfloor, \lfloor \frac{j+1}{2} \rfloor}$ . Thus,

$$\sum_{i=1}^{m/2} \sum_{j=1}^{n/2} Y_{ij} (\mathcal{T}(\mathcal{X}))_{ij} = \sum_{i=1}^m \sum_{j=1}^n (\mathcal{T}^T(Y))_{ij} X_{ij} = \sum_{i=1}^m \sum_{j=1}^n \frac{1}{4} (\mathcal{U}(Y))_{ij} X_{ij}$$

## 2

Suppose input  $X \in \mathbb{R}^{C_{in} \times H_{in} \times W_{in}}$  and the output  $Y \in \mathbb{R}^{C_{out} \times H_{out} \times W_{out}}$ . Noting the **ConvTranspose2d()** operation,  $H_{out} = (H_{in} - 1) \times \text{striding} + k$  without any consideration of padding. To make **ConvTranspose2d()** operation and **Upsample()** operation are equivalent, the filter size  $k$  and striding should satisfy the following equation.

$$H_{out} = (H_{in} - 1) \times \text{striding} + k = rH_{in}$$

Selecting  $k = r$ ,  $\text{striding} = r$  is the simplest way to hold the condition above. Also, to do nearest upsampling, the weight of the filter should be initialized to be 1 with dimension  $r \times r$ . Thus,

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layer = nn.ConvTranspose2d(C_in, C_out, kernel_size = r, stride = r)
layer.weight.data = torch.ones(r, r)
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## 3

a

$$\begin{aligned} D_f(X||Y) &= \int f\left(\frac{P_X(x)}{P_Y(x)}\right) P_Y(x) dx \\ &= \mathbb{E}\left[f\left(\frac{P_X}{P_Y}\right)\right] \\ &\geq f\left(\mathbb{E}\left[\frac{P_X}{P_Y}\right]\right) \\ &= f\left(\int \frac{P_X(x)}{P_Y(x)} P_Y(x) dx\right) \\ &= f(1) = 0 \end{aligned}$$

b

To show that  $y = -\log x$  and  $y = x \log x$  are convex, I used the following theorem : The differentiable function  $f$  is strictly convex *if and only if* its derivate  $f'$  is strictly increasing (Refer to the appendix to see the proof of the following theorem). Noting that logarithm function is defined on  $\{x \in \mathbb{R} | x > 0\}$ ,

$$\begin{aligned} \frac{d^2}{dx^2}(-\log x) &= \frac{d}{dx}\left(-\frac{1}{x}\right) = \frac{1}{x^2} \geq 0 \\ \frac{d^2}{dx^2}(x \log x) &= \frac{d}{dx} \log x + 1 = \frac{1}{x} \geq 0 \end{aligned}$$

Also, it is obvious that  $\log 1 = 0, 1 \times \log 1 = 0$ . Thus, the  $f$ -divergence of both  $-\log x, x \log x$  are well defined. Then, one can find that  $f$ -divergence of both functions are KL-divergence.

$$\begin{aligned} D_f(X||Y) &= \int -\log\left(\frac{P_X(x)}{P_Y(x)}\right) P_Y(x) dx = \int P_Y(x) \log\left(\frac{P_Y(x)}{P_X(x)}\right) dx = D_{KL}(P_Y||P_X) \\ D_f(X||Y) &= \int P_Y(x) \frac{P_X(x)}{P_Y(x)} \log\left(\frac{P_X(x)}{P_Y(x)}\right) dx = \int P_X(x) \log\left(\frac{P_X(x)}{P_Y(x)}\right) dx = D_{KL}(P_X||P_Y) \end{aligned}$$

## 4

$$\mathbb{P}(G(U) \leq t) = \mathbb{P}(\inf \{x \in \mathbb{R} | U \leq F(x)\} \leq t)$$

Define two event  $A, B$  as follows :

$$\begin{aligned} A &= \{\inf \{x \in \mathbb{R} | U \leq F(x)\} \leq t\} \\ B &= \{U \leq F(t)\} \end{aligned}$$

Since the CDF of  $U$ ,  $F_U(u)$ , is  $F_U(u) = u$  for  $u \in [0, 1]$ ,  $\mathbb{P}(B) = \mathbb{P}(U \leq F(t)) = F_U(F(t)) = F(t)$ . If  $A = B$ , then  $\mathbb{P}(A) = \mathbb{P}(B) = F(x)$  and the proof is done. To show that  $A = B$ , consider the both cases :  $B \subset A, A \subset B$ .

(i) :  $A \subset B$

Suppose  $A$  is true. Then,  $\inf \{x \in \mathbb{R} | U \leq F(x)\} \leq t$ . This implies  $\exists x_0 \in \mathbb{R}, F(x_0) \geq U$  for  $x_0 \leq t$ . Since the CDF  $F$  is nondecreasing,  $F(t) \geq F(x_0)$ . Thus,  $U \leq F(x_0) \leq F(t)$ , which concludes that  $B$  is true.

(ii) :  $B \subset A$

Suppose  $B$  is true. Then,  $U \leq F(t)$ . Consider a set  $S = \{x | F(x) \geq U\}$ . Then,  $t \in S$ . This implies  $\inf S = \inf \{x | F(x) \geq U\} \leq t$ . Thus, the event  $A$  is true. As a result,  $A = B$  and this ends up with  $\mathbb{P}(G(U) \leq t) = F(t)$

## 5

From  $X = AY + b$ ,  $\varphi(X) = Y = A^{-1}(X - b)$ . According to the property of  $p_X(x)$  produced by the condition of the problem,

$$\begin{aligned} P_X(x) &= P_Y(\varphi(x)) \left| \det \frac{\partial \varphi}{\partial x}(x) \right| \\ &= P_Y(\varphi(x)) |\det A^{-1}| \\ &= \frac{1}{\det A} \frac{1}{(2\pi)^{n/2}} e^{-\frac{1}{2} \|A^{-1}(x-b)\|^2} \\ &= \frac{1}{\sqrt{(2\pi)^n \det \Sigma}} e^{-\frac{1}{2} (x-b)^T (A^{-1})^T A^{-1} (x-b)} (\because \det \Sigma = \det A \det A^T = (\det A)^2) \\ &= \frac{1}{\sqrt{(2\pi)^n \det \Sigma}} e^{-\frac{1}{2} (x-b)^T \Sigma^{-1} (x-b)} (\because \Sigma^{-1} = (AA^T)^{-1} = (A^{-1})^T A^{-1}) \end{aligned}$$

## 6

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### Algorithm 1 Inverse Permutation Pseudo Code

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1: Let  $n$  be the length of the list  $\sigma$ 
2: Initialize an empty list called result
3: for  $i = 1$  to  $n$  do
4:   for  $j = 1$  to  $n$  do
5:     if  $\sigma[i]$  equals  $j + 1$  then
6:       Append  $j + 1$  to the result list
7:     end if
8:   end for
9: end for

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## 7

**a**

$$[P_\sigma x]_i = e_{\sigma(i)}^T x = x_{\sigma(i)}$$

**b**

$$\begin{aligned} (P_\sigma P_\sigma^T)_{ij} &= \sum_k P_\sigma^{(ik)} P_\sigma^{(jk)} \\ &= \sum_k e_{\sigma(i)}^{(k)} e_{\sigma(j)}^{(k)} \\ &= \sum_k \delta_{\sigma(i),k} \delta_{\sigma(j),k} \\ &= \delta_{\sigma(i),\sigma(j)} \\ &= \delta_{ij} (\because \sigma \text{ is bijection}) \end{aligned}$$

$$\begin{aligned}
(P_{\sigma^{-1}}P_{\sigma})_{ij} &= \sum_k P_{\sigma^{-1}}^{(ik)} P_{\sigma}^{(kj)} \\
&= \sum_k e_{\sigma^{-1}(i)}^{(k)} e_{\sigma(k)}^{(j)} \\
&= \sum_k \delta_{\sigma^{-1}(i),k} \delta_{\sigma(k),j} \\
&= \delta_{ij} \quad (\because \delta_{\sigma^{-1}(i),k} \delta_{\sigma(k),j} = 1 \text{ if } k = \sigma^{-1}(i), \sigma(k) = j)
\end{aligned}$$

**c**

$$\begin{aligned}
\det(P_{\sigma}P_{\sigma}^T) &= (\det P_{\sigma})^2 = 1 \\
\therefore |\det P_{\sigma}| &= 1
\end{aligned}$$

## A Proof of Theorem in problem 3-b

**proof**

$\implies$  :  $a, b \in \mathbb{R} (a < b)$ . Let  $x_1, x_2, x_3 \in \mathbb{R}$  be chosen s.t.  $a < x_1 < x_2 < x_3 < b$ .

From Chordal Slope Lemma, which is famous lemma applicable to convex function, the following inequalities are held.

$$\frac{f(x_1) - f(a)}{x_1 - a} < \frac{f(x_2) - f(x_1)}{x_2 - x_1} < \frac{f(x_3) - f(x_2)}{x_3 - x_2} < \frac{f(b) - f(x_3)}{b - x_3} \quad (1)$$

$$\frac{f(x_2) - f(a)}{x_2 - a} < \frac{f(b) - f(x_2)}{b - x_2} \quad (2)$$

Taking limit  $x_1 \rightarrow a, x_3 \rightarrow b$  on equation 1,

$$f'(a) \leq \frac{f(x_2) - f(a)}{x_2 - a} < \frac{f(b) - f(x_2)}{b - x_2} \leq f'(b) \quad (3)$$

and, therefore,  $f'(a) < f'(b)$  for arbitrary  $a, b \in \mathbb{R}$ .

$\Leftarrow$  : Let  $x_1, x_2, x_3 \in \mathbb{R}$  s.t.  $x_1 < x_2 < x_3$ . From Mean Value Theorem,

$$\exists a \text{ s.t. } \frac{f(x_2) - f(x_1)}{x_2 - x_1} = f'(a) \quad (4)$$

$$\exists b \text{ s.t. } \frac{f(x_3) - f(x_2)}{x_3 - x_2} = f'(b) \quad (5)$$

Note that  $x_1 < a < x_2 < b < x_3$ . Since  $f'$  is strictly increasing,  $f'(a) < f'(b)$ , which implies

$$\frac{f(x_2) - f(x_1)}{x_2 - x_1} < \frac{f(x_3) - f(x_2)}{x_3 - x_2} \quad (6)$$

The result of equation 6 is equivalent that  $f$  is strictly convex according to Chordal Slope Lemma. ■