

water filling

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1 The CASE $\ell = 1$

For the water-filling method, in the case where $\ell = 1$,

$$value(p^i) = \frac{i - t}{\sum_{j=1}^i \frac{1}{v_j}}$$

$$MaximumValue = \max_{t+1 \leq i \leq n} \left(value(p^i) \right)$$

Where i refers to a specific box

p^i is the probability that box i is selected.

For example, in the puzzle posited by the paper with values 8, 7, 5, 4, and $t = 1$:

Testing $i = 2$: $\frac{2-1}{\frac{1}{8}+\frac{1}{7}} \approx 3.73$

Testing $i = 3$: $\frac{3-1}{\frac{1}{8}+\frac{1}{7}+\frac{1}{5}} \approx 4.27$.

Testing $i = 4$: ≈ 4.18 .

From this, we determine that we should randomize over the first 3 boxes, and ignoring the forth in order to maximize potential value.

2 The Case $\ell > 1$

For each box i , let $p'_i = \mathbb{P}_{S \sim p}(i \in S)$ be the probability that each box i is included in the final set S .

The set of marginal probabilities $p' = (p'_1, \dots, p'_n)$ forms a pseudo-distribution $p' \in \Delta_\ell([n])$.

Each p'_i is between 0 and 1.

The sum of all marginals is exactly ℓ .

3 Pseudocode for water_filling

```

function water_filling(values,  $t$ ,  $\ell$ ):
     $n \leftarrow \text{length}(\text{values})$ 
    if  $t \geq \ell \wedge (\ell = n)$ , return values
    if  $\ell = 1$ :
         $\text{max\_val}, \text{optimal\_p\_prime} \leftarrow \text{solve\_l1}(\text{values}, t, n)$ 
    else:
         $\text{max\_val}, \text{optimal\_p\_prime} \leftarrow \text{solve\_l}(\text{values}, t, n)$ 
    return  $\text{max\_val}, \text{optimal\_p\_prime}$ 

```

4 Pseudocode for solve_l1

```

function solve_l1(values,  $t$ ,  $\ell$ ):
     $\text{max\_value} \leftarrow 0$ 
     $\text{optimal\_i} \leftarrow t$ 
     $\text{inverse\_v\_sum} = 0$ 
    for ( $i = t + 1, i = n + 1, i++$ ):
         $\text{inverse\_v\_sum} \leftarrow \text{inverse\_v\_sum} + \frac{1}{v[i-1]}$ 
         $\text{current\_value} \leftarrow \frac{i-t}{\text{inverse\_v\_sum}}$ 
        if  $\text{current\_value} > \text{max\_value}$ :
             $\text{max\_value} = \text{current\_value}$ 
             $\text{optimal\_i} = i$ 
    if  $\text{max\_value} > 0$ :
         $\text{p\_prime} \leftarrow [0, \dots, 0]$  (of length  $n$ )
        
$$C \leftarrow \sum_{j=0}^{\text{optimal\_i}} \left( \frac{1}{v_j} \right)$$

        
$$\text{final\_p} \leftarrow \sum_{j=0}^n \left( \frac{1}{v_j} \cdot \frac{1}{C} \right) [j < \text{optimal\_i}]$$
 (using anderson notation)
        return  $\text{max\_value}, \text{final\_p}$ 
    else return 0,  $[0, \dots, 0]$ 

```