

water filling

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November 2025

1 The CASE $\ell = 1$

For the water-filling method, in the case where $\ell = 1$,

$$\begin{aligned} \text{value}(p^i) &= \frac{i - t}{\sum_{j=1}^i \frac{1}{v_j}} \\ \text{MaximumValue} &= \max_{t+1 \leq i \leq n} \left(\text{value}(p^i) \right) \end{aligned}$$

Where i refers to a specific box

p^i is the probability that box i is selected.

For example, in the puzzle posited by the paper with values 8, 7, 5, 4, and $t = 1$:

Testing $i = 2$: $\frac{2-1}{\frac{1}{8} + \frac{1}{7}} \approx 3.73$

Testing $i = 3$: $\frac{3-1}{\frac{1}{8} + \frac{1}{7} + \frac{1}{5}} \approx 4.27$.

Testing $i = 4$: ≈ 4.18 .

From this, we determine that we should randomize over the first 3 boxes, and ignoring the forth in order to maximize potential value.

2 The Case $\ell > 1$

For each box i , let $p'_i = \mathbb{P}_{S \sim p}(i \in S)$ be the probability that each box i is included in the final set S .

The set of marginal probabilities $p' = (p'_1, \dots, p'_n)$ forms a pseudo-distribution $p' \in \Delta_\ell([n])$.

Each p'_i is between 0 and 1.

The sum of all marginals is exactly ℓ .

3 Pseudocode for water_filling

```

function water_filling(values,  $t$ ,  $\ell$ ):
     $n \leftarrow \text{length}(\text{values})$ 
    if  $t \geq \ell \wedge (l = n)$ , return values
    if  $l = 1$ :
         $\max\_val, \text{optimal\_p\_prime} \leftarrow \text{solve\_l1}(\text{values}, t, n)$ 
    else:
         $\max\_val, \text{optimal\_p\_prime} \leftarrow \text{solve\_l}(\text{values}, t, n)$ 
    return  $\max\_val, \text{optimal\_p\_prime}$ 

```

4 Pseudocode for solve_l1

```

function solve_l1(values,  $t$ ,  $\ell$ ):
     $\max\_value \leftarrow 0$ 
     $\text{optimal\_i} \leftarrow t$ 
     $\text{inverse\_v\_sum} = 0$ 
    for ( $i = t + 1, i = n + 1, i += 1$ ):
         $\text{inverse\_v\_sum} \leftarrow \text{inverse\_v\_sum} + \frac{1}{v[i-1]}$ 
         $\text{current\_value} \leftarrow \frac{i-t}{\text{inverse\_v\_sum}}$ 
        if  $\text{current\_value} > \max\_value$ :
             $\max\_value = \text{current\_value}$ 
             $\text{optimal\_i} = i$ 
    if  $\max\_value > 0$ :
         $p\_prime \leftarrow [0, \dots, 0]$  (of length  $n$ )
         $C \leftarrow \sum_{j=0}^{\text{optimal\_i}} \left( \frac{1}{v_j} \right)$ 
         $final\_p \leftarrow \sum_{j=0}^n \left( \frac{1}{v_j} \cdot \frac{1}{C} \right) [j < \text{optimal\_i}]$  (using anderson notation)
        return  $\max\_value, final\_p$ 
    else return  $0, [0, \dots, 0]$ 

```