

Quantum Computing

A Gentle Introduction to Grover's Algorithm

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v1.0

Outline

1. Grover's algorithm
 - 1.1. Motivation & Outline
 - 1.2. Steps
2. Implementation of Grover's algorithm: 2-Qubit States
 - 2.1. Quantum Circuit
 - 2.2. IBM Implementation

Outline

1. Grover's algorithm

1.1. Motivation & Outline

1.2. Steps

2. Implementation of Grover's algorithm: 2-Qubit States

2.1. Quantum Circuit

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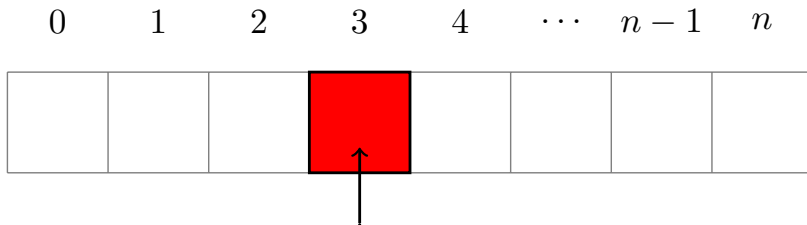
2.1. Quantum Circuit

2.2. IBM Implementation

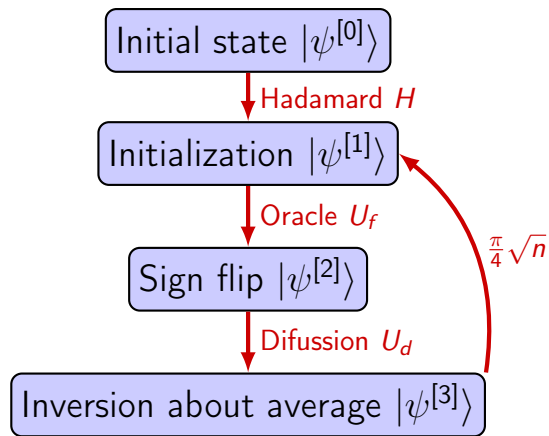
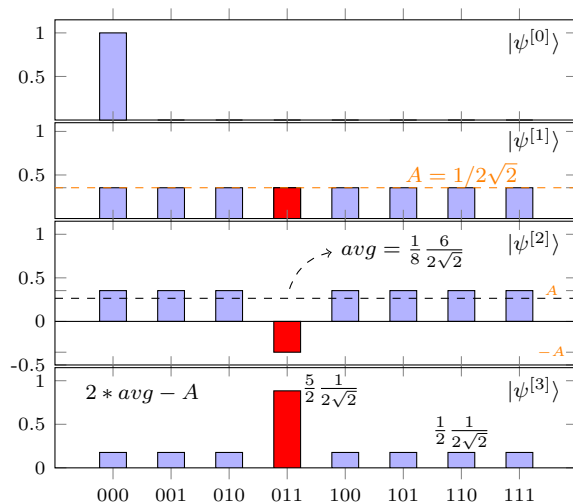
Grover's algorithm: Motivation

Grover's algorithm performs a search over an unordered set of 2^n items to find the unique element that satisfies some condition

- Classic approach: $\sum_{i=1}^n \frac{1}{n} i = \frac{n+1}{2} \Rightarrow \mathcal{O}(n)$
- Quantum approach: (...) $\Rightarrow \mathcal{O}(\sqrt{n})$



Grover's algorithm: Outline



Outline

1. Grover's algorithm

1.1. Motivation & Outline

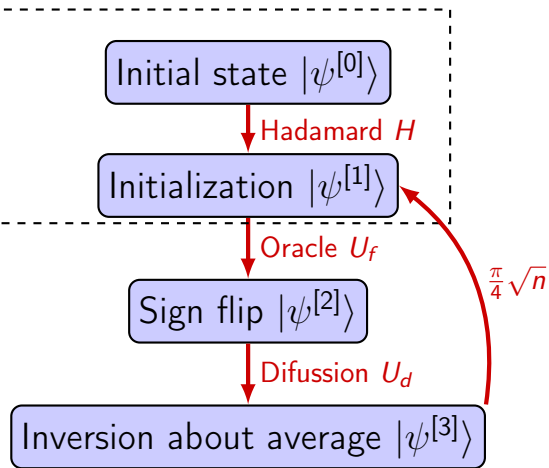
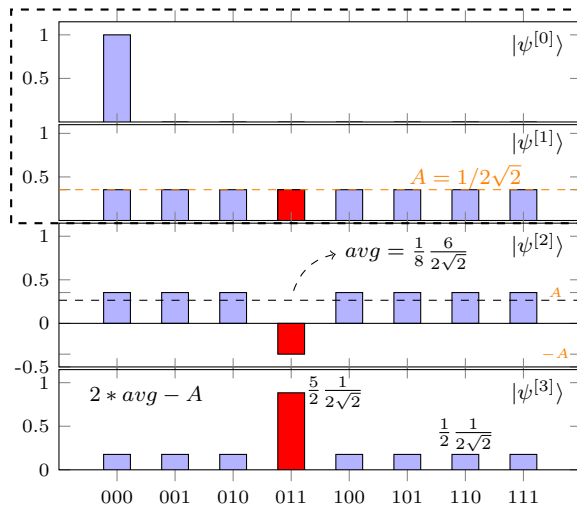
1.2. Steps

2. Implementation of Grover's algorithm: 2-Qubit States

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Initialization



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We find a method (unitary operator) to have all the states with the same probability (*principle of superposition*).

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We find a method (unitary operator) to have all the states with the same probability (*principle of superposition*).

$$\begin{aligned}
 (I^{\otimes n} \otimes X) |0\rangle_{n+1} &= |0\rangle_n \otimes |1\rangle \\
 H^{\otimes(n+1)} [(I^{\otimes n} \otimes X) |0\rangle_{n+1}] &= H^{\otimes n} |0\rangle_n \otimes H |1\rangle \\
 &= \sum_{j \in \{0,1\}^n} \frac{1}{\sqrt{2^n}} |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \\
 &= \sum_{j \in \{0,1\}^n} \alpha_j |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \\
 &= |\psi^{[1]}\rangle
 \end{aligned}$$

Initialization

3-qubit example: Set ancillary qubit to $|1\rangle$

$$\begin{aligned}
 (I^{\otimes 3} \otimes X) |0\rangle_{3+1} &= I^{\otimes 3} |0\rangle_3 \otimes X |0\rangle \\
 &= \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\
 &= |0\rangle_3 \otimes |1\rangle
 \end{aligned}$$

Initialization

3-qubit example: Apply Hadamard gate

$$\begin{aligned}
 |\psi^{[1]}\rangle &= H^{\otimes(3+1)} [(I^{\otimes 3} \otimes X) |0\rangle_{n+1}] \\
 &= H^{\otimes 3} |0\rangle_3 \otimes H |1\rangle \\
 &= \frac{1}{\sqrt{2^3}} \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ 1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 \\ 1 & 1 & -1 & -1 & -1 & -1 & 1 & 1 \\ 1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \otimes \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix}
 \end{aligned}$$

Initialization

3-qubit example: Apply Hadamard gate

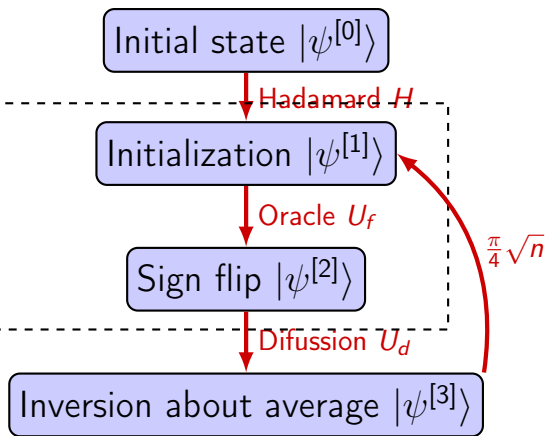
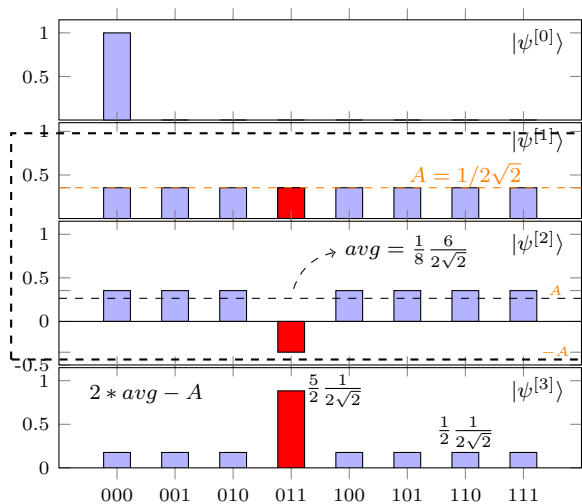
$$\begin{aligned}
 |\psi^{[1]}\rangle &= \frac{1}{\sqrt{2^3}} (11111111)^\dagger \otimes \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix} \\
 &= \left[\frac{1}{2\sqrt{2}} |000\rangle + \frac{1}{2\sqrt{2}} |001\rangle + \frac{1}{2\sqrt{2}} |010\rangle + \frac{1}{2\sqrt{2}} |011\rangle \right. \\
 &\quad \left. + \frac{1}{2\sqrt{2}} |100\rangle + \frac{1}{2\sqrt{2}} |101\rangle + \frac{1}{2\sqrt{2}} |110\rangle + \frac{1}{2\sqrt{2}} |111\rangle \right] \\
 &\quad \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)
 \end{aligned}$$

Initialization

3-qubit example: Summary

$$\begin{aligned}
 |\psi^{[1]}\rangle &= H^{\otimes(3+1)} [(I^{\otimes 3} \otimes X) |0\rangle_{n+1}] \\
 &= \left[\frac{1}{2\sqrt{2}} |000\rangle + \frac{1}{2\sqrt{2}} |001\rangle + \frac{1}{2\sqrt{2}} |010\rangle + \frac{1}{2\sqrt{2}} |011\rangle \right. \\
 &\quad \left. + \frac{1}{2\sqrt{2}} |100\rangle + \frac{1}{2\sqrt{2}} |101\rangle + \frac{1}{2\sqrt{2}} |110\rangle + \frac{1}{2\sqrt{2}} |111\rangle \right] \\
 &\quad \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \\
 &= \sum_{j \in \{0,1\}^n} \frac{1}{\sqrt{2^n}} |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) = |\psi^{[1]}\rangle
 \end{aligned}$$

Sign flip



Sign flip

We find a method (unitary operator) which flip the sign of the state of interest.

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Quantum Oracle

It is defined the operator U_f :

$$U_f : |j\rangle_n \otimes |y\rangle_1 \rightarrow |j\rangle_n \otimes |y \oplus f(j)\rangle_1,$$

where \oplus is the sum operator in mod 2, and $f(j) = \begin{cases} 1 & j = I \\ 0 & j \neq I \end{cases}$

A	B	XOR
0	0	0
0	1	1
1	0	1
1	1	0

Sign flip

We apply U_f (*Quantum Oracle*) to the previous state $\psi^{[1]}$

$$\begin{aligned}
 |\psi^{[2]}\rangle &= U_f |\psi^{[1]}\rangle \\
 &= U_f \left(\sum_{j \in \{0,1\}^n} \alpha_j |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right) \\
 &= U_f \left(\alpha_l |l\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^n; j \neq l} \alpha_j |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right) \\
 &= \left(-\alpha_l |l\rangle_n + \sum_{j \in \{0,1\}^n; j \neq l} \alpha_j |j\rangle_n \right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)
 \end{aligned}$$

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 &= \left(-\alpha_l |l\rangle_n + \sum_{j \in \{0,1\}^n; j \neq l} \alpha_j |j\rangle_n \right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)
 \end{aligned}$$

Note: A red arrow points from the boxed term $-\alpha_l |l\rangle_n$ in the final equation to the boxed text "sign flip!" in the second equation, indicating the sign change.

Sign flip

We apply U_f (Quantum Oracle) to the previous state $\psi^{[1]}$

$$\begin{aligned}
 |\psi^{[2]}\rangle &= U_f |\psi^{[1]}\rangle \\
 &= U_f \left(\sum_{j \neq l} \alpha_j |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right) \quad \text{sign flip!} \quad \text{Extra qubit!} \\
 &= U_f \left(\alpha_l |l\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^n; j \neq l} \alpha_j |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right) \\
 &= \left(-\alpha_l |l\rangle_n + \sum_{j \in \{0,1\}^n; j \neq l} \alpha_j |j\rangle_n \right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)
 \end{aligned}$$

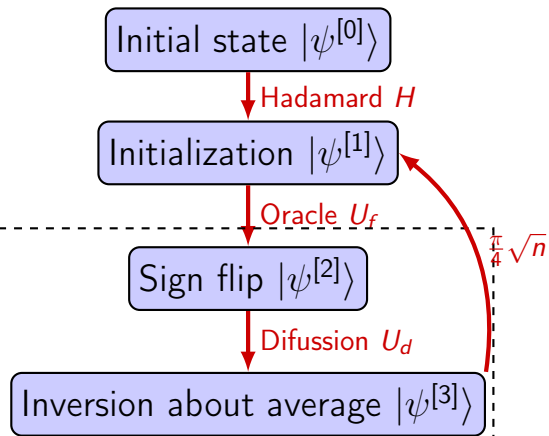
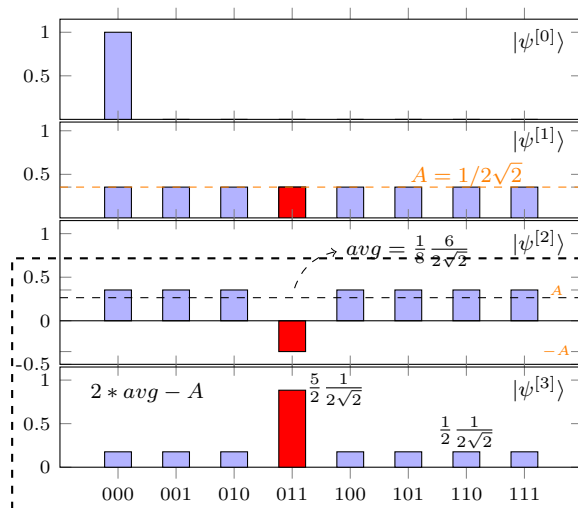
Sign flip¹

3-qubit example: Quantum Oracle

$$\begin{aligned}
 |\psi^{[2]}\rangle &= U_f |\psi^{[1]}\rangle \\
 &= \left[\frac{1}{2\sqrt{2}} |000\rangle + \frac{1}{2\sqrt{2}} |001\rangle + \frac{1}{2\sqrt{2}} |010\rangle - \frac{1}{2\sqrt{2}} |011\rangle \right. \\
 &\quad \left. + \frac{1}{2\sqrt{2}} |100\rangle + \frac{1}{2\sqrt{2}} |101\rangle + \frac{1}{2\sqrt{2}} |110\rangle + \frac{1}{2\sqrt{2}} |111\rangle \right] \\
 &\quad \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)
 \end{aligned}$$

¹This slide shows the result of applying U_f but not how is applied. This is because this step of the algorithm depends on the specific problem.

Inversion about average



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We find a method (unitary operator) to invert the amplitude about the average

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$$\sum_{j \in \{0,1\}^n} \alpha_j |j\rangle_n \xrightarrow{U_d} \sum_{j \in \{0,1\}^n} \left(2 \left(\sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n} \right) - \alpha_j \right) |j\rangle_n,$$

where $\sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n}$ is the average.

$$\sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n} = \frac{1}{2^3} \frac{6}{2\sqrt{2}}$$

Inversion about the average

Difussion operator

$$U_d = \begin{pmatrix} \frac{2}{2^n} & \frac{2}{2^n} & \cdots & \frac{2}{2^n} \\ \frac{2}{2^n} & \frac{2}{2^n} & \cdots & \frac{2}{2^n} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{2}{2^n} & \frac{2}{2^n} & \cdots & \frac{2}{2^n} \end{pmatrix} - I^{\otimes n} = \cdots = -H^{\otimes n} D H^{\otimes n}, \quad (1)$$

where $D = \text{diag}(-1, 1, 1, \dots, 1)$

Inversion about the average

3-qubit example: Difussion operator in python

```
>>> import numpy as np
>>> H1 = 1/np.sqrt(2)*np.array([[1,1],[1, -1]]) # Hadamard operator 1 qubit
>>> H2 = np.kron(H1,H1) # Hadamard operator 2 qubit
>>> H3 = np.kron(H2,H1) # Hadamard operator 3 qubit

>>> D = np.eye(8) # Diagonal operator
>>> D[0,0] = -1

>>> Ud = -np.dot(np.dot(H3,D),H3) # Difussion operator
>>> psi_2 = 1/(2*np.sqrt(2))*np.array([1,1,1,-1,1,1,1,1]) # psi_2

>>> psi_3 = np.dot(Ud,psi_2)
>>> print(psi_3)
array([0.176 , 0.176 , 0.176 , 0.883, 0.176, 0.176, 0.176, 0.176])
```

Inversion about the average

3-qubit example: Difussion operator

$$|\psi^{[3]}\rangle = U_d |\psi^{[2]}\rangle$$

$$|\psi^{[3]}\rangle = -\frac{1}{8} \frac{1}{2\sqrt{2}} \begin{pmatrix} 6 & -2 & -2 & -2 & -2 & -2 & -2 & -2 \\ -2 & 6 & -2 & -2 & -2 & -2 & -2 & -2 \\ -2 & -2 & 6 & -2 & -2 & -2 & -2 & -2 \\ -2 & -2 & -2 & 6 & -2 & -2 & -2 & -2 \\ -2 & -2 & -2 & -2 & 6 & -2 & -2 & -2 \\ -2 & -2 & -2 & -2 & -2 & 6 & -2 & -2 \\ -2 & -2 & -2 & -2 & -2 & -2 & 6 & -2 \\ -2 & -2 & -2 & -2 & -2 & -2 & -2 & 6 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \\ 1 \\ -1 \\ 1 \\ 1 \\ 1 \\ 1 \end{pmatrix} = \frac{1}{4\sqrt{2}} \begin{pmatrix} 1 \\ 1 \\ 1 \\ 5 \\ 1 \\ 1 \\ 1 \\ 1 \end{pmatrix}$$

Measurement & Repetition

For the first iteration we measure

$$\begin{aligned}\alpha_{011} &= \frac{5}{2} \frac{1}{2\sqrt{2}} \longrightarrow \|\alpha_{011}\|^2 \simeq 78,12\% \\ \alpha_j &= \frac{1}{2} \frac{1}{2\sqrt{2}} \longrightarrow \|\alpha_j\|^2 \simeq 3,12\% \quad (j \neq |011\rangle)\end{aligned}$$

The optimal number of repetitions $R \simeq \frac{\pi}{4} \sqrt{n} \simeq 2.2$. For the second iteration (Oracle & Diffusion) we have:

$$\begin{aligned}\alpha_{011} &= \frac{11}{4} \frac{1}{2\sqrt{2}} \longrightarrow \|\alpha_{011}\|^2 \simeq 94,5\% \\ \alpha_j &= \frac{-1}{4} \frac{1}{2\sqrt{2}} \longrightarrow \|\alpha_j\|^2 \simeq 0,78\% \quad (j \neq |011\rangle)\end{aligned}$$

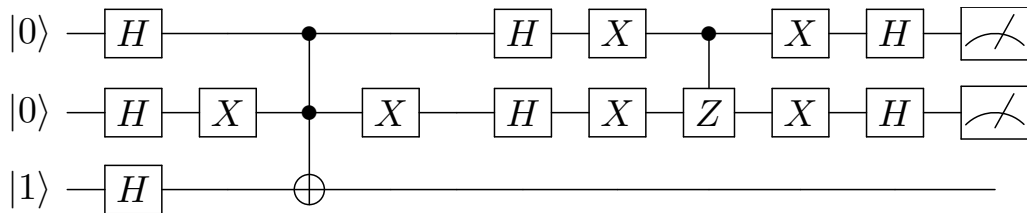
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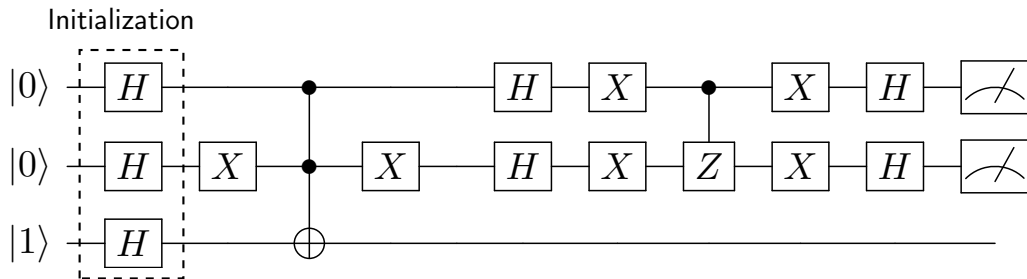
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2-Qubit Quantum Circuit



2-Qubit Quantum Circuit

Initialization



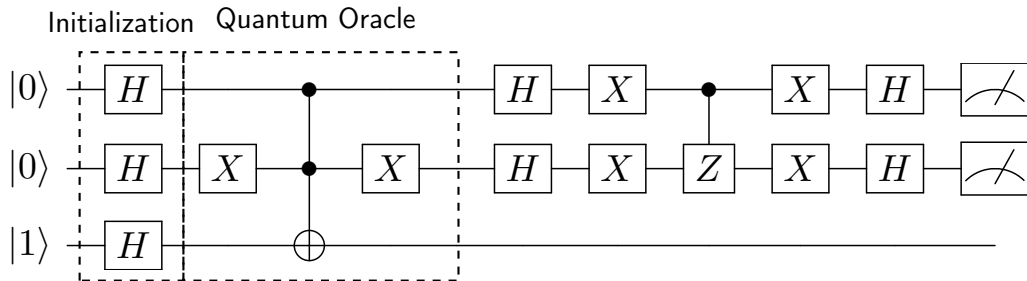
2-Qubit Quantum Circuit

Initialization

$$\begin{aligned}\psi^{[0]} &= |001\rangle \\ \psi^{[1]} &= H^{\otimes 3} \psi^{[0]} \\ &= \frac{1}{\sqrt{8}} (|000\rangle + |010\rangle + |100\rangle + |110\rangle - |001\rangle - |011\rangle - |101\rangle - |111\rangle)\end{aligned}$$

2-Qubit Quantum Circuit

Quantum Oracle



2-Qubit Quantum Circuit

Quantum Oracle

$$U_f = (I \otimes X \otimes I) \cdot T \cdot (I \otimes X \otimes I) = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

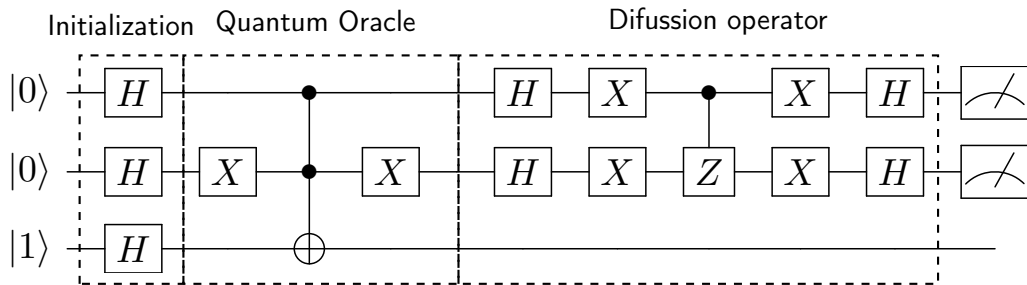
2-Qubit Quantum Circuit

Quantum Oracle

$$\begin{aligned}
 \psi^{[1]} &= \frac{1}{\sqrt{8}} (|000\rangle + |010\rangle + |100\rangle + |110\rangle - |001\rangle - |011\rangle - |101\rangle - |111\rangle) \\
 \psi^{[2]} &= U_f |\psi^{[1]}\rangle \\
 &= \frac{1}{\sqrt{8}} (|000\rangle + |010\rangle + |101\rangle + |110\rangle - |001\rangle - |011\rangle - |100\rangle - |111\rangle) \\
 &= \frac{1}{\sqrt{8}} (|00\rangle + |01\rangle - |10\rangle + |11\rangle) \otimes (|0\rangle - |1\rangle) \\
 &= \frac{1}{2} (|00\rangle + |01\rangle - |10\rangle + |11\rangle) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)
 \end{aligned}$$

2-Qubit Quantum Circuit

Difussion Operator



2-Qubit Quantum Circuit

Diffusion Operator

$$U_d = (H \otimes H) \cdot (X \otimes X) \cdot (CZ) (X \otimes X) \cdot (H \otimes H) = \frac{1}{4} \begin{pmatrix} 2 & -2 & -2 & -2 \\ -2 & 2 & -2 & -2 \\ -2 & -2 & 2 & -2 \\ -2 & -2 & -2 & 2 \end{pmatrix}$$

$$\psi^{[3]} = U_d \psi^{[2]}$$

$$\frac{1}{8} \begin{pmatrix} 2 & -2 & -2 & -2 \\ -2 & 2 & -2 & -2 \\ -2 & -2 & 2 & -2 \\ -2 & -2 & -2 & 2 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} = \frac{1}{8} \begin{pmatrix} 2 & -2 & 2 & -2 \\ -2 & 2 & 2 & -2 \\ -2 & -2 & -2 & -2 \\ -2 & -2 & 2 & 2 \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ 0 \\ -1 \\ 0 \end{pmatrix} = -|10\rangle$$

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