# Quantum Computing A Gentle Introduction to Grover's Algorithm

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- 1. Grover's algorithm
  - 1.1. Motivation & Outline
  - 1.2. Steps
- 2. Implementation of Grover's algorithm: 2-Qubit States
  - 2.1. Quantum Circuit
  - 2.2. IBM Implementation
- 3. References

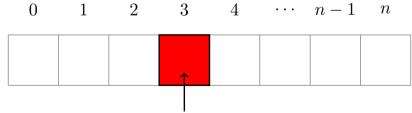
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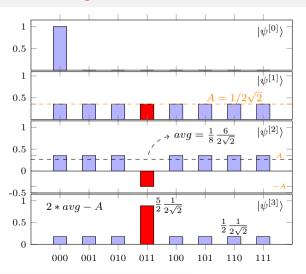
# Grover's algorithm: Motivation

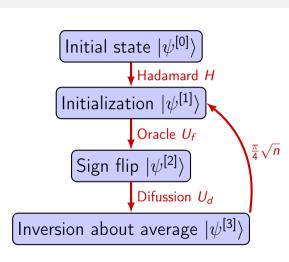
Grover's algorithm performs a search over an unorder set of  $2^n$  items fo find the unique element that satisfies some condition

- Classic approach:  $\sum_{i=1}^{n} \frac{1}{n}i = \frac{n+1}{1} \Rightarrow \mathcal{O}(n)$
- Quantum approach:  $(...) \Rightarrow \mathcal{O}(\sqrt{n})$

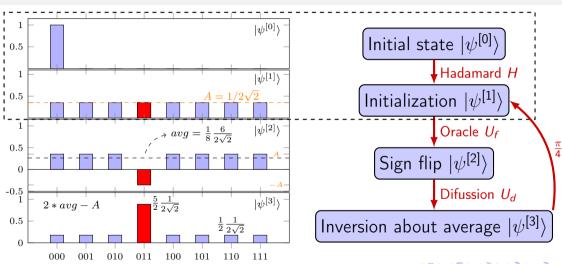


# Grover's algorithm: Outline





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$$(I^{\otimes n} \otimes X) |0\rangle_{n+1} = |0\rangle_n \otimes |1\rangle$$

$$H^{\otimes (n+1)} [(I^{\otimes n} \otimes X) |0\rangle_{n+1}] = H^{\otimes n} |0\rangle_n \otimes H |1\rangle$$

$$= \sum_{j \in \{0,1\}^n} \frac{1}{\sqrt{2^n}} |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$= \sum_{j \in \{0,1\}^n} \alpha_j |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$= |\psi^{[1]}\rangle$$

#### 3-qubit example: Set ancillary qubit to $|1\rangle$

$$(I^{\otimes 3} \otimes X) |0\rangle_{3+1} = I^{\otimes 3} |0\rangle_{3} \otimes X |0\rangle$$

$$= \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$= |0\rangle_{3} \otimes |1\rangle$$

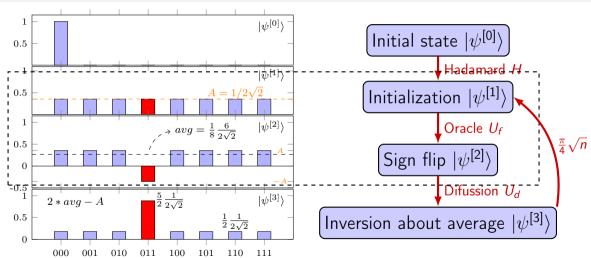
#### 3-qubit example: Apply Hadamard gate

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$$\begin{array}{ll} |\psi^{[1]}\rangle & = & \frac{1}{\sqrt{2^3}} \left(111111111\right)^{\dagger} \otimes \frac{1}{\sqrt{2}} \left(\begin{array}{c} 1 \\ -1 \end{array}\right) \\ & = & \left[\frac{1}{2\sqrt{2}} \left|000\right\rangle + \frac{1}{2\sqrt{2}} \left|001\right\rangle + \frac{1}{2\sqrt{2}} \left|010\right\rangle + \frac{1}{2\sqrt{2}} \left|011\right\rangle \\ & + & \frac{1}{2\sqrt{2}} \left|100\right\rangle + \frac{1}{2\sqrt{2}} \left|101\right\rangle + \frac{1}{2\sqrt{2}} \left|110\right\rangle + \frac{1}{2\sqrt{2}} \left|111\right\rangle\right] \\ & \otimes & \frac{1}{\sqrt{2}} \left(\left|0\right\rangle - \left|1\right\rangle\right) \end{array}$$

### 3-qubit example: Summary

$$\begin{split} |\psi^{[1]}\rangle &= H^{\otimes(3+1)}\left[(I^{\otimes 3}\otimes X)|0\rangle_{n+1}\right] \\ &= \left[\frac{1}{2\sqrt{2}}|000\rangle + \frac{1}{2\sqrt{2}}|001\rangle + \frac{1}{2\sqrt{2}}|010\rangle + \frac{1}{2\sqrt{2}}|011\rangle \\ &+ \frac{1}{2\sqrt{2}}|100\rangle + \frac{1}{2\sqrt{2}}|101\rangle + \frac{1}{2\sqrt{2}}|110\rangle + \frac{1}{2\sqrt{2}}|111\rangle\right] \\ &\otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \\ &= \sum_{j\in\{0,1\}^n} \frac{1}{\sqrt{2^n}}|j\rangle_n \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = |\psi^{[1]}\rangle \end{split}$$



We find a method (unitary operator) which flip the sign of the state of interest.

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#### Quantum Oracle

It is defined the operator  $U_f$ :

$$U_f: |j\rangle_n \otimes |y\rangle_1 \rightarrow |j\rangle_n \otimes |y \oplus f(j)\rangle_1,$$

where  $\oplus$  is the sum operator in mod 2, and  $f(j) = \left\{ \begin{array}{cc} 1 & j = l \\ 0 & j \neq l \end{array} \right\}$ 

Α	В	XOR
0	0	0
0	1	1
1	0	1
1	1	0

We apply  $U_f$  (Quantum Oracle) to the previous state  $\psi^{[1]}$ 

$$|\psi^{[2]}\rangle = U_{f} |\psi^{[1]}\rangle$$

$$= U_{f} \left(\sum_{j \in \{0,1\}^{n}} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)\right)$$

$$= U_{f} \left(\alpha_{I} |I\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)\right)$$

$$= \left(-\alpha_{I} |I\rangle_{n} + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n}\right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

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$$= U_{f} \left( \alpha |I\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right)$$

$$= \left( -\alpha_{I} |I\rangle_{n} + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

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$$= U_{f} \left( \sum_{\text{sign flip!}} j \rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right) \text{Extra qubit!}$$

$$= U_{f} \left( \alpha |I\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right)$$

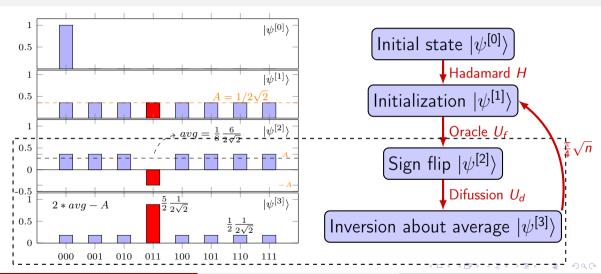
$$= \left( -\alpha_{I} |I\rangle_{n} + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

# Sign flip<sup>1</sup>

### 3-qubit example: Quantum Oracle

$$\begin{split} |\psi^{[2]}\rangle &= U_f |\psi^{[1]}\rangle \\ &= \left[\frac{1}{2\sqrt{2}} |000\rangle + \frac{1}{2\sqrt{2}} |001\rangle + \frac{1}{2\sqrt{2}} |010\rangle - \frac{1}{2\sqrt{2}} |011\rangle \right. \\ &+ \left. \frac{1}{2\sqrt{2}} |100\rangle + \frac{1}{2\sqrt{2}} |101\rangle + \frac{1}{2\sqrt{2}} |110\rangle + \frac{1}{2\sqrt{2}} |111\rangle \right] \\ &\otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \end{split}$$

<sup>&</sup>lt;sup>1</sup>This slide shows the result of applying  $U_f$  but not how is applied. This is because this step of the algorithm depends on the specific problem.



We find a method (unitary operator) to invert the amplitude about the average

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$$\sum_{j \in \{0,1\}^n} \alpha_j |j\rangle_n \xrightarrow{U_d} \sum_{j \in \{0,1\}^n} \left( 2 \left( \sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n} \right) - \alpha_j \right) |j\rangle_n,$$

where  $\sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n}$  is the average.

$$\sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n} = \frac{1}{2^3} \frac{6}{2\sqrt{2}}$$

#### Difussion operator

$$U_{d} = \begin{pmatrix} \frac{2}{2^{n}} & \frac{2}{2^{n}} & \cdots & \frac{2}{2^{n}} \\ \frac{2}{2^{n}} & \frac{2}{2^{n}} & \cdots & \frac{2}{2^{n}} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{2}{2^{n}} & \frac{2}{2^{n}} & \cdots & \frac{2}{2^{n}} \end{pmatrix} - I^{\otimes n} = \cdots = -H^{\otimes n}DH^{\otimes n}, \tag{1}$$

where 
$$D = diag(-1, 1, 1, \cdots, 1)$$

#### 3-qubit example: Difussion operator in python

```
>>> import numpy as np
>>> H1 = 1/np.sqrt(2)*np.array([[1,1],[1,-1]]) # Hadamard operator 1 qubit
>>> H2 = np.kron(H1,H1) # Hadamard operator 2 qubit
>>> H3 = np.kron(H2,H1) # Hadamard operator 3 qubit
>>> D = np.eye(8) # Diagonal operator
>>> D \lceil 0.0 \rceil = -1
>>> Ud = -np.dot(np.dot(H3,D),H3) # Difussion operator
>> psi_2 = 1/(2*np.sqrt(2))*np.array([1,1,1,-1,1,1,1,1]) # psi 2
>>> psi_3 = np.dot(Ud.psi_2)
>>> print(psi_3)
array([0.176, 0.176, 0.176, 0.883, 0.176, 0.176, 0.176, 0.176])
```

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### 3-qubit example: Difussion operator

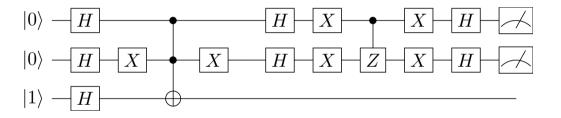
# Measurement & Repetition

For the first iteration we measure

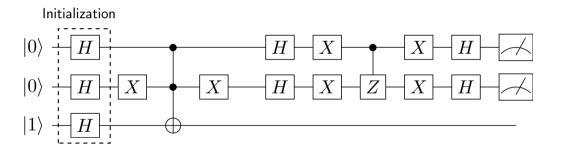
The optimal number of repetitions  $R \simeq \frac{\pi}{4} \sqrt{n} \simeq 2.2$ . For the second iteration (Oracle & Diffusion) we have:

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#### Initialization

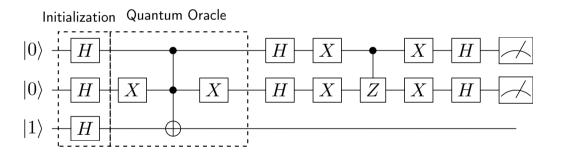


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#### Initialization

$$\begin{array}{lll} \psi^{[0]} & = & |001\rangle \\ \psi^{[1]} & = & H^{\otimes 3}\psi^{[0]} \\ & = & \frac{1}{\sqrt{8}} (|000\rangle + |010\rangle + |100\rangle + |110\rangle - |001\rangle - |011\rangle - |101\rangle - |111\rangle) \end{array}$$

#### Quantum Oracle



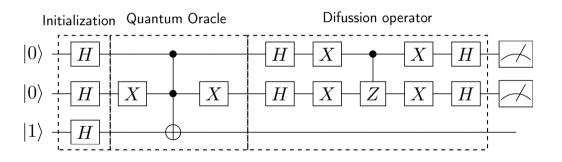
#### Quantum Oracle

$$U_f = (I \otimes X \otimes I) \cdot T \cdot (I \otimes X \otimes I) = egin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \ 0 & 1 & 0 & 0 & 0 & 0 & 0 \ 0 & 0 & 1 & 0 & 0 & 0 & 0 \ 0 & 0 & 0 & 1 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

#### Quantum Oracle

$$\begin{array}{ll} \psi^{[1]} & = & \frac{1}{\sqrt{8}} \left( |000\rangle + |010\rangle + |100\rangle + |110\rangle - |001\rangle - |011\rangle - |101\rangle - |111\rangle \right) \\ \psi^{[2]} & = & U_f |\psi^{[1]}\rangle \\ & = & \frac{1}{\sqrt{8}} \left( |000\rangle + |010\rangle + |101\rangle + |110\rangle - |001\rangle - |011\rangle - |100\rangle - |111\rangle \right) \\ & = & \frac{1}{\sqrt{8}} \left( |00\rangle + |01\rangle - |10\rangle + |11\rangle \right) \otimes \left( |0\rangle - |1\rangle \right) \\ & = & \frac{1}{2} \left( |00\rangle + |01\rangle - |10\rangle + |11\rangle \right) \otimes \frac{1}{\sqrt{2}} \left( |0\rangle - |1\rangle \right) \end{array}$$

#### **Difussion Operator**



#### **Diffusion Operator**

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