# Quantum Computing A Gentle Introduction to Grover's Algorithm

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- 1. Grover's algorithm
  - 1.1. Motivation & Outline
  - 1.2. Steps
- 2. Implementation of Grover's algorithm: 2-Qubit States
  - 2.1. Quantum Circuit
  - 2.2. IBM Implementation
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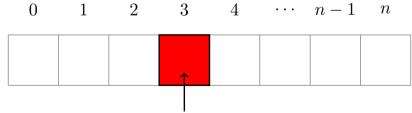
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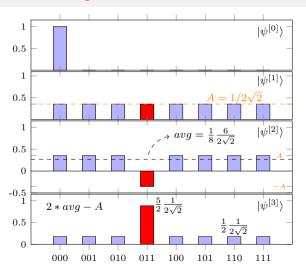
# Grover's algorithm: Motivation

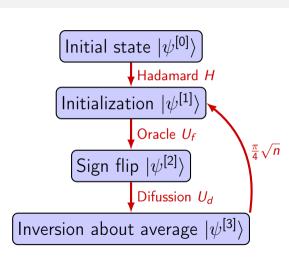
Grover's algorithm performs a search over an unorder set of  $2^n$  items fo find the unique element that satisfies some condition

- Classic approach:  $\sum_{i=1}^{n} \frac{1}{n}i = \frac{n+1}{1} \Rightarrow \mathcal{O}(n)$
- Quantum approach:  $(...) \Rightarrow \mathcal{O}(\sqrt{n})$

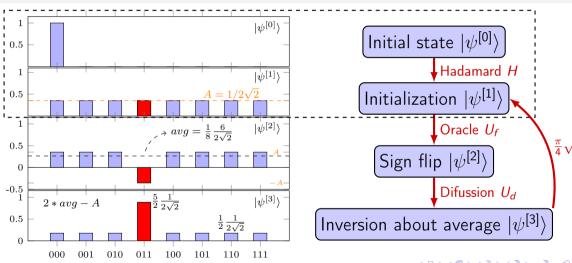


# Grover's algorithm: Outline





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We find a method (unitary operator) to have all the states with the same probability (*principle* of superposition).

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$$(I^{\otimes n} \otimes X) |0\rangle_{n+1} = |0\rangle_n \otimes |1\rangle$$

$$H^{\otimes (n+1)} [(I^{\otimes n} \otimes X) |0\rangle_{n+1}] = H^{\otimes n} |0\rangle_n \otimes H |1\rangle$$

$$= \sum_{j \in \{0,1\}^n} \frac{1}{\sqrt{2^n}} |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$= \sum_{j \in \{0,1\}^n} \alpha_j |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$= |\psi^{[1]}\rangle$$

#### 3-qubit example: Set ancillary qubit to $|1\rangle$

$$(I^{\otimes 3} \otimes X) |0\rangle_{3+1} = I^{\otimes 3} |0\rangle_{3} \otimes X |0\rangle$$

$$= \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$= |0\rangle_{3} \otimes |1\rangle$$

#### 3-qubit example: Apply Hadamard gate

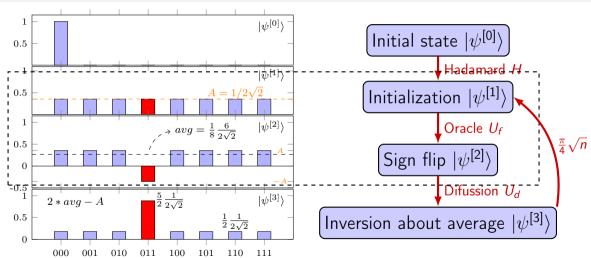
### 3-qubit example: Apply Hadamard gate

$$\begin{split} |\psi^{[1]}\rangle &=& \frac{1}{\sqrt{2^3}} \left(11111111\right)^{\dagger} \otimes \frac{1}{\sqrt{2}} \left(\begin{array}{c} 1 \\ -1 \end{array}\right) \\ &=& \left[\frac{1}{2\sqrt{2}} \left|000\right\rangle + \frac{1}{2\sqrt{2}} \left|001\right\rangle + \frac{1}{2\sqrt{2}} \left|010\right\rangle + \frac{1}{2\sqrt{2}} \left|011\right\rangle \\ &+& \frac{1}{2\sqrt{2}} \left|100\right\rangle + \frac{1}{2\sqrt{2}} \left|101\right\rangle + \frac{1}{2\sqrt{2}} \left|110\right\rangle + \frac{1}{2\sqrt{2}} \left|111\right\rangle \right] \\ &\otimes& \frac{1}{\sqrt{2}} \left(\left|0\right\rangle - \left|1\right\rangle\right) \end{split}$$

#### 3-qubit example: Summary

$$\begin{split} |\psi^{[1]}\rangle &= H^{\otimes(3+1)}\left[(I^{\otimes 3} \otimes X) |0\rangle_{n+1}\right] \\ &= \left[\frac{1}{2\sqrt{2}} |000\rangle + \frac{1}{2\sqrt{2}} |001\rangle + \frac{1}{2\sqrt{2}} |010\rangle + \frac{1}{2\sqrt{2}} |011\rangle \\ &+ \frac{1}{2\sqrt{2}} |100\rangle + \frac{1}{2\sqrt{2}} |101\rangle + \frac{1}{2\sqrt{2}} |110\rangle + \frac{1}{2\sqrt{2}} |111\rangle\right] \\ &\otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \\ &= \sum_{j \in \{0,1\}^n} \frac{1}{\sqrt{2^n}} |j\rangle_n \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) = |\psi^{[1]}\rangle \end{split}$$

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We find a method (unitary operator) which flip the sign of the state of interest.

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#### Quantum Oracle

It is defined the operator  $U_f$ :

$$U_f: |j\rangle_n \otimes |y\rangle_1 \rightarrow |j\rangle_n \otimes |y \oplus f(j)\rangle_1,$$

where  $\oplus$  is the sum operator in mod 2, and  $f(j) = \left\{ \begin{array}{cc} 1 & j = l \\ 0 & j \neq l \end{array} \right\}$ 

Α	В	XOR
0	0	0
0	1	1
1	0	1
1	1	0

We apply  $U_f$  (Quantum Oracle) to the previous state  $\psi^{[1]}$ 

$$|\psi^{[2]}\rangle = U_{f} |\psi^{[1]}\rangle$$

$$= U_{f} \left(\sum_{j \in \{0,1\}^{n}} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)\right)$$

$$= U_{f} \left(\alpha_{I} |I\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)\right)$$

$$= \left(-\alpha_{I} |I\rangle_{n} + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n}\right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

We apply  $U_f$  (Quantum Oracle) to the previous state  $\psi^{[1]}$ 

$$|\psi^{[2]}\rangle = U_{f} |\psi^{[1]}\rangle$$

$$= U_{f} \left( \sum_{\text{sign flip!}} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right)$$

$$= U_{f} \left( \alpha |I\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right)$$

$$= \left( -\alpha_{I} |I\rangle_{n} + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \right) \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

We apply  $U_f$  (Quantum Oracle) to the previous state  $\psi^{[1]}$ 

$$|\psi^{[2]}\rangle = U_{f} |\psi^{[1]}\rangle$$

$$= U_{f} \left( \sum_{\text{sign flip!}} i \rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right) \text{Extra qubit!}$$

$$= U_{f} \left( \alpha |I\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right)$$

$$= \left( -\alpha_{I} |I\rangle_{n} + \sum_{j \in \{0,1\}^{n}; j \neq I} \alpha_{j} |j\rangle_{n} \right) \otimes \left( \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right)$$

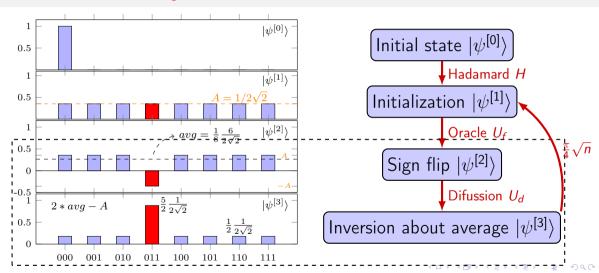
# Sign flip<sup>1</sup>

#### 3-qubit example: Quantum Oracle

$$\begin{split} |\psi^{[2]}\rangle &= U_f |\psi^{[1]}\rangle \\ &= \left[\frac{1}{2\sqrt{2}} |000\rangle + \frac{1}{2\sqrt{2}} |001\rangle + \frac{1}{2\sqrt{2}} |010\rangle - \frac{1}{2\sqrt{2}} |011\rangle \\ &+ \frac{1}{2\sqrt{2}} |100\rangle + \frac{1}{2\sqrt{2}} |101\rangle + \frac{1}{2\sqrt{2}} |110\rangle + \frac{1}{2\sqrt{2}} |111\rangle \right] \\ &\otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \end{split}$$

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<sup>&</sup>lt;sup>1</sup>This slide shows the result of applying  $U_f$  but not how is applied. This is because this step of the algorithm depends on the specific problem.



We find a method (unitary operator) to invert the amplitude about the average

We find a method (unitary operator) to invert the amplitude about the average

$$\sum_{j \in \{0,1\}^n} \alpha_j |j\rangle_n \xrightarrow{U_d} \sum_{j \in \{0,1\}^n} \left( 2 \left( \sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n} \right) - \alpha_j \right) |j\rangle_n,$$

where  $\sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n}$  is the average.

$$\sum_{k \in \{0,1\}^n} \frac{\alpha_k}{2^n} = \frac{1}{2^3} \frac{6}{2\sqrt{2}}$$

#### Difussion operator

$$U_{d} = \begin{pmatrix} \frac{2}{2^{n}} & \frac{2}{2^{n}} & \cdots & \frac{2}{2^{n}} \\ \frac{2}{2^{n}} & \frac{2}{2^{n}} & \cdots & \frac{2}{2^{n}} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{2}{2^{n}} & \frac{2}{2^{n}} & \cdots & \frac{2}{2^{n}} \end{pmatrix} - I^{\otimes n} = \cdots = -H^{\otimes n}DH^{\otimes n}, \tag{1}$$

where  $D = diag(-1, 1, 1, \cdots, 1)$ 

#### 3-qubit example: Difussion operator in python

```
>>> import numpy as np
>>> H1 = 1/np.sqrt(2)*np.array([[1,1],[1,-1]]) # Hadamard operator 1 qubit
>>> H2 = np.kron(H1,H1) # Hadamard operator 2 qubit
>>> H3 = np.kron(H2,H1) # Hadamard operator 3 qubit
>>> D = np.eye(8) # Diagonal operator
>>> D \lceil 0.0 \rceil = -1
>>> Ud = -np.dot(np.dot(H3,D),H3) # Difussion operator
>> psi_2 = 1/(2*np.sqrt(2))*np.array([1,1,1,-1,1,1,1,1]) # psi 2
>>> psi_3 = np.dot(Ud.psi_2)
>>> print(psi_3)
array([0.176, 0.176, 0.176, 0.883, 0.176, 0.176, 0.176, 0.176])
```

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#### 3-qubit example: Difussion operator

# Measurement & Repetition

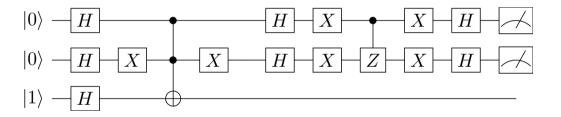
For the first iteration we measure

$$\alpha_{011} = \frac{5}{2} \frac{1}{2\sqrt{2}} \longrightarrow \|\alpha_{011}\|^2 \simeq 78,12\%$$
 $\alpha_j = \frac{1}{2} \frac{1}{2\sqrt{2}} \longrightarrow \|\alpha_j\|^2 \simeq 3,12\% \quad (j \neq |011\rangle)$ 

The optimal number of repetitions  $R \simeq \frac{\pi}{4} \sqrt{n} \simeq 2.2$ . For the second iteration (Oracle & Diffusion) we have:

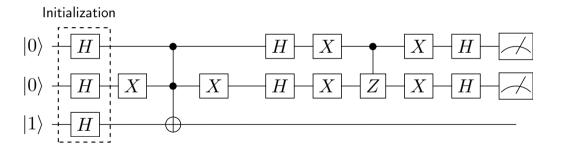
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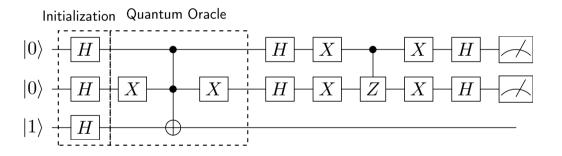
#### Initialization



#### Initialization

$$\begin{array}{lll} \psi^{[0]} & = & |001\rangle \\ \psi^{[1]} & = & H^{\otimes 3}\psi^{[0]} \\ & = & \frac{1}{\sqrt{8}} (|000\rangle + |010\rangle + |100\rangle + |110\rangle - |001\rangle - |011\rangle - |101\rangle - |111\rangle) \end{array}$$

#### Quantum Oracle



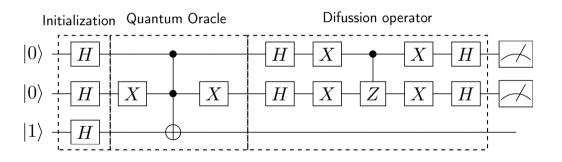
#### Quantum Oracle

$$U_f = (I \otimes X \otimes I) \cdot T \cdot (I \otimes X \otimes I) = egin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \ 0 & 1 & 0 & 0 & 0 & 0 & 0 \ 0 & 0 & 1 & 0 & 0 & 0 & 0 \ 0 & 0 & 0 & 1 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

#### Quantum Oracle

$$\begin{array}{ll} \psi^{[1]} & = & \frac{1}{\sqrt{8}} \left( |000\rangle + |010\rangle + |100\rangle + |110\rangle - |001\rangle - |011\rangle - |101\rangle - |111\rangle \right) \\ \psi^{[2]} & = & U_f |\psi^{[1]}\rangle \\ & = & \frac{1}{\sqrt{8}} \left( |000\rangle + |010\rangle + |101\rangle + |110\rangle - |001\rangle - |011\rangle - |100\rangle - |111\rangle \right) \\ & = & \frac{1}{\sqrt{8}} \left( |00\rangle + |01\rangle - |10\rangle + |11\rangle \right) \otimes \left( |0\rangle - |1\rangle \right) \\ & = & \frac{1}{2} \left( |00\rangle + |01\rangle - |10\rangle + |11\rangle \right) \otimes \frac{1}{\sqrt{2}} \left( |0\rangle - |1\rangle \right) \end{array}$$

#### **Difussion Operator**



#### **Diffusion Operator**

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# **Implementation**

Next slides show the implementation of the algorithm in two different ways:

- Using *qiskit* framework. In this case we have a simulated quantum backend. The reader can copy & run the code in python in order to reproduce the experiment
- 2 Building the gates directly in Quantum IBM computer

In both cases we will observe a histogram with an unique state  $(|10\rangle)$ 

Backend & Quantum Registers

```
from giskit import *
from giskit import OuantumCircuit
from qiskit.circuit.quantumcircuit import QuantumCircuit
from giskit.visualization import plot_histogram
backend = BasicAer.get_backend('gasm_simulator')
#backend = BasicAer.get backend('statevector simulator')
'''Quantum register
https://quantumcomputing.stackexchange.com/questions/4907/
qiskit — is — there — any — way — to — discard — the — results — of — a — measurement '''
q = QuantumRegister(2, 'q')
a = OuantumRegister(1. 'a')
c = ClassicalRegister(2, 'c')
```

#### Quantum Circuit

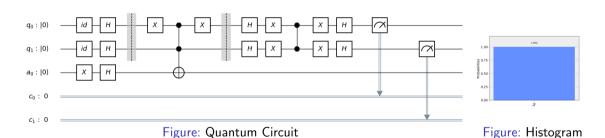
```
#Circuit
circ = QuantumCircuit(q,a,c) # type: qiskit.circuit.quantumcircuit.QuantumCircuit
# ===== prepare the states: psi^{[0]} =====
circ.iden(q[0])
circ.iden(q[1])
circ.x(a[0])
#==== Initialization: psi^{[1]} =====
circ.h(q[0])
circ.h(a[1])
circ.h(a[0])
circ.barrier(q)
#==== Sign flip: psi^{[2]} =====
circ.x(q[0])
circ.ccx(q[0], q[1], a[0])
circ.x(q[0])
circ.barrier(g)
#==== Inversion about average: psi^{[3]} =====
circ.h(a[0])
circ.h(q[1])
circ.x(q[0])
circ.x(q[1])
```

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#### Measurement

```
#==== Meaurement =====
circ.measure(q,c)
#iob
job = execute(circ. backend=backend. shots=1000)
result = job.result()
#circuit
figure = circ.draw(output='mpl')
figure.savefig('../data/images/giskit-circuit.png')
#histogram
counts = result.get_counts(circ)
print("Total counts:")
print(counts)
figure = plot_histogram(counts)
figure.savefig('../data/images/qiskit-histogram.png')
```

#### Measurement



# IBM Implementation: Quantum Experience

https://www.research.ibm.com/ibm-q/technology/experience/

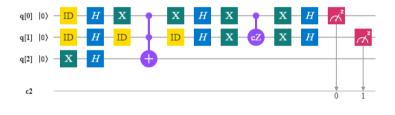


Figure: Quantum Circuit



Figure: Histogram

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- H. Abraham and I. Y. A. et al., "Qiskit: An open-source framework for quantum computing," 2019. [Online]. Available: https://doi.org/10.5281/zenodo.2562111

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