

Recitation 15 - Homework 6 and 7 Answers and Questions

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December 6, 2006

- ▶ Extra Monday (12/18) Office Hours
 - ▶ 12:30-1:30pm
 - ▶ 5:00-7:00pm
- ▶ Homework 6 Answers
- ▶ Homework 7 Answers
- ▶ Wrap-Up

Problem 1. Show EQ_{CFG} is undecidable.

Problem 2. If $A \leq_M B$ and B is regular, does that imply A is regular? Why or why not?

$$T = \{M \mid M \text{ accepts } w^R \text{ whenever } M \text{ accepts } w\}$$

Problem 3. Show that T is undecidable in two ways.

Problem 4. Consider

$$A = \{(M, w) \mid M \text{ moves left on left most tape pos. on } w\}$$

Show A is undecidable.

Problem 5. Consider

$$A = \{(M, w) \mid M' \text{'s head ever moves left on } w\}$$

Show A is decidable.

Problem 6.

$$J = \{w \mid w = 0x \text{ for some } x \in A_{TM} \text{ or } w = 1y \text{ for some } y \in \overline{A_{TM}}\}$$

Big-O Problems

- ▶ $n^2 = O(n)$
- ▶ $3^n = 2^{O(n)}$
- ▶ $2^{2^n} = O(2^{2^n})$

Little-o Problems

- ▶ $n = o(2n)$
- ▶ $1 = o(1/n)$

Two graphs are isomorphic if the nodes of one graph can be relabelled in such a way that the resulting graph is equal to the second graph.

$$ISO = \{(G, H) \mid G \text{ is isomorphic to } H\}$$

Show $ISO \in NP$.

$$MODEXP = \{(a, b, c, p) \mid (a^b \bmod p) = (c \bmod p)\}$$

Demonstrate $MODEXP \in P$. That is some deterministic TM decides $MODEXP$ in polynomial time.

Problem 5. Part a.

$SPATH = \{(G, a, b, k) \mid G \text{ contains a simple path from } a \text{ to } b \text{ of length at most } k\}$

Show $SPATH \in P$.

Problem 5. Part b.

$LPATH = \{(G, a, b, k) \mid G \text{ contains a simple (loopless) path from } a \text{ to } b \text{ of length at least } k\}$

Show $LPATH$ is NP-Complete.

Problem 6. Problem 7.28.

$$SETSPLITTING = \{(S, C = \{C_i\}) \mid C_i \subseteq S \text{ \& } (S, C) \text{ colorable}\}$$

S can be colored if each element can be chosen to be *red* or *blue* such that no C_i 's elements are all the same color. Show $SETSPLITTING$ is NP-complete.

$$\text{Variables} = \{x_1, x_2, x_3\}$$

$$\phi = \{x_1 \vee x_2 \vee x_3\} \wedge \{\bar{x}_2 \vee x_2 \vee x_3\} \wedge \{\bar{x}_1 \vee x_3 \vee \bar{x}_3\}$$

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$$\phi = \{x_1 \vee x_2 \vee x_3\} \wedge \{\bar{x}_2 \vee x_2 \vee x_3\} \wedge \{\bar{x}_1 \vee x_3 \vee \bar{x}_3\}$$



$$S = \{x_1, \bar{x}_1, x_2, \bar{x}_2, x_3, \bar{x}_3\}$$

$$C = \{\{x_1, x_2, x_3\}, \{\bar{x}_2, \bar{x}_2, \bar{x}_3\}, \{\bar{x}_1, x_3, \bar{x}_3\}\}$$

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$$S = \{\textcolor{red}{x}_1, \textcolor{blue}{\bar{x}}_1, \textcolor{blue}{x}_2, \textcolor{red}{\bar{x}}_2, \textcolor{red}{x}_3, \textcolor{blue}{\bar{x}}_3\}$$

$$C = \{\{\textcolor{red}{x}_1, \textcolor{blue}{x}_2, \textcolor{red}{x}_3\}, \{\textcolor{red}{\bar{x}}_2, \textcolor{red}{\bar{x}}_2, \textcolor{blue}{\bar{x}}_3\}, \{\textcolor{blue}{\bar{x}}_1, \textcolor{red}{x}_3, \textcolor{blue}{\bar{x}}_3\}\}$$

$$\{x_1, x_2\}$$

$$\{x_1 \vee \bar{x}_1 \vee x_2\} \wedge \{x_1 \vee x_1 \vee \bar{x}_2\} \wedge \{\bar{x}_1 \vee \bar{x}_1 \vee x_2\} \wedge \{\bar{x}_1 \vee \bar{x}_1 \vee \bar{x}_2\}$$

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$$S = \{x_1, \bar{x}_1, x_2, \bar{x}_2\}$$

$$\{\{x_1, x_1, x_2\}, \{x_1, x_1, \bar{x}_2\}, \{\bar{x}_1, \bar{x}_1, x_2\}, \{\bar{x}_1, \bar{x}_1, \bar{x}_2\}\}$$

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$$S = \{\textcolor{red}{x}_1, \textcolor{red}{\bar{x}}_1, \textcolor{blue}{x}_2, \textcolor{blue}{\bar{x}}_2\}$$

$$\{\{\textcolor{red}{x}_1, \textcolor{red}{x}_1, \textcolor{blue}{x}_2\}, \{\textcolor{red}{x}_1, \textcolor{red}{x}_1, \textcolor{blue}{\bar{x}}_2\}, \{\textcolor{red}{\bar{x}}_1, \textcolor{red}{\bar{x}}_1, \textcolor{blue}{x}_2\}, \{\textcolor{red}{\bar{x}}_1, \textcolor{red}{\bar{x}}_1, \textcolor{blue}{\bar{x}}_2\}\}$$

$$\{x_1, x_2\}$$

$$\{x_1 \vee x_1 \vee x_2\} \wedge \{x_1 \vee x_1 \vee \bar{x}_2\} \wedge \{\bar{x}_1 \vee \bar{x}_1 \vee x_2\}$$

$$\Downarrow\Downarrow\Downarrow\Downarrow$$

$$S = \{x_1, \bar{x}_1, x_2, \bar{x}_2\}$$

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- ▶ Content of recitations?
- ▶ Pace of recitations?
- ▶ Accessible?
- ▶ Grading?
- ▶ Obtaining help?