

1 Exclusive Buyer Mechanism

Suppose there are $i = 1, \dots, N$ bidders in an auction for a good with $j = 1, 2$ quality levels. Suppose each bidder's valuation for the good is given by $x^i = (x_1^i, x_2^i) \in [\underline{x}_1^i, \bar{x}_1^i] \times [\underline{x}_2^i, \bar{x}_2^i]$. We assume bidders' valuations are identical and independently distributed but we allow an individual bidder's valuations for quality grades to be arbitrarily correlated.

We now construct an *exclusive-buyer mechanism*. This mechanism is a little different¹ than that of (Belloni, Lopomo, and Wang 2010). Since bidder valuations are symmetric, we omit the superscript i (i.e., $X = X^i$). For any given *reserve price* $p = (p_1, p_2)$, let

$$\beta_1 = x_1 - p_1 \quad \text{and} \quad \beta_2 = x_2 - p_2 \quad (1)$$

and define the interim allocations² as:

$$Q_1(x; p) = \mathbb{1}\{\beta_1 > \beta_2 \wedge \beta_1 \geq 0\} \left(\int_{\underline{x}}^{(x_1, \min\{\bar{x}_2, p_2 + \beta_1\})} \mathbb{1}\{t_1 - p_1 > t_2 - p_2 \wedge t_1 - p_1 \geq 0\} dF(t) \right)^{N-1} \quad (2)$$

$$Q_2(x; p) = \mathbb{1}\{\beta_2 > \beta_1 \wedge \beta_2 \geq 0\} \left(\int_{\underline{x}}^{(\min\{\bar{x}_1, p_1 + \beta_2\}, x_2)} \mathbb{1}\{t_2 - p_2 > t_1 - p_1 \wedge t_2 - p_2 \geq 0\} dF(t) \right)^{N-1} \quad (3)$$

Then the expected revenue from price p is given by:

$$R_{EBM}(p) = N \int_X Q(x; p) \cdot (p - c) dF(x) \quad (4)$$

where costs are given by $c = (c_1, c_2)$.

2 Calculations

For ease of calculation, we divide up the type space according to which quality level of the good is preferred:

$$A := \{x \in X \mid x_1 - p_1 > x_2 - p_2 \text{ and } x_1 - p_1 \geq 0\} \quad (5)$$

$$B := \{x \in X \mid x_2 - p_2 > x_1 - p_1 \text{ and } x_2 - p_2 \geq 0\} \quad (6)$$

¹It does not rely on a "add-on price" in addition to a "reserve price". Instead, we assume a price is given for each quality grade.

²In the event of ties where $\beta_1 = \beta_2$ then both allocations Q_1, Q_2 are equal.

2.1 $N = 1$

Notice that for $N = 1$,

$$R(p) = \int_{x \in A} (p_1 - c_1) dF(x) + \int_{x \in B} (p_2 - c_2) dF(x) \quad (7)$$

2.1.1 $X = U[0, 1]^2$ (Pavlov 2011)

$p^* = (\frac{1}{\sqrt{3}}, \frac{1}{\sqrt{3}})$ DONE

2.1.2 $X = U[6, 8] \times U[9, 11], c = [.9, 5]$ (Belloni, Lopomo, and Wang 2010)

$p^* = (9, 10.5)$ DONE

$revenue^* \approx 5.12$ DONE

2.2 $N = 2$

2.2.1 $X = U[6, 8] \times U[9, 11], c = [.9, 5]$ (Belloni, Lopomo, and Wang 2010)

$p^* = (6, 10.1)$ TODO

revenue* ≈ 5.89 TODO

$$R_{EBM}(p) = N \left(\int_{x \in A} Q_1(x; p)(p_1 - c_1) dF(x) + \int_{x \in B} Q_2(x; p)(p_2 - c_2) dF(x) \right) \quad (8)$$

$$= N \int_{x \in A} \left(\int_{\underline{x}}^{(x_1, \min\{\bar{x}_2, p_2 + x_1 - p_1\})} \mathbb{1}\{t_1 - p_1 > t_2 - p_2 \wedge t_1 - p_1 \geq 0\} dF(t) \right)^{N-1} (p_1 - c_1) dF(x) \quad (9)$$

$$+ N \int_{x \in B} \left(\int_{\underline{x}}^{(\min\{\bar{x}_1, p_1 + x_2 - p_2\}, x_2)} \mathbb{1}\{t_2 - p_2 > t_1 - p_1 \wedge t_2 - p_2 \geq 0\} dF(t) \right)^{N-1} (p_2 - c_2) dF(x) \quad (10)$$

$$= 2 \int_6^8 \int_9^{\min\{11, 10.1 + x_1 - 6\}} \left(\int_6^{x_1} \int_9^{\min\{11, 10.1 + x_1 - 6\}} \mathbb{1}\{t_1 - 6 > t_2 - 10.1\} f(t_1, t_2) dt_2 dt_1 \right) (6 - .9) f(x_1, x_2) dx_2 dx_1 \quad (11)$$

$$+ 2 \int_6^{6.9} \int_{10.1 + x_1 - 6}^{11} \left(\int_6^{6 + x_2 - 10.1} \int_{10.1}^{x_2} \mathbb{1}\{t_2 - 10.1 > t_1 - 6\} f(t_1, t_2) dt_2 dt_1 \right) (10.1 - 5) f(x_1, x_2) dx_2 dx_1 \quad (12)$$

$$= \frac{(6 - .9)}{8} \int_6^{6.9} \int_9^{10.1 + x_1 - 6} \left(\int_6^{x_1} \int_9^{10.1 + t_1 - 6} dt_2 dt_1 \right) dx_2 dx_1 \quad (13)$$

$$+ \frac{(6 - .9)}{8} \int_{6.9}^8 \int_9^{11} \left(\int_6^{x_1} \int_9^{\min\{11, 10.1 + t_1 - 6\}} dt_2 dt_1 \right) dx_2 dx_1 \quad (14)$$

$$+ \frac{(10.1 - 5)}{8} \int_6^{6.9} \int_{10.1 + x_1 - 6}^{11} \left(\int_{10.1}^{x_2} \int_6^{6 + t_2 - 10.1} dt_1 dt_2 \right) dx_2 dx_1 \quad (15)$$

$$= 4.17 \dots \quad (16)$$

3 References

- Belloni, Alexandre, Giuseppe Lopomo, and Shouqiang Wang (2010). “Multidimensional mechanism design: Finite-dimensional approximations and efficient computation”. In: *Operations Research* 58.4-part-2, pp. 1079–1089.
- Pavlov, Gregory (2011). “Optimal mechanism for selling two goods”. In: *The BE Journal of Theoretical Economics* 11.1.