

# Propositional logic

## Solution 1

(a)

false (as  $(\text{true} \Rightarrow \text{false}) \Leftrightarrow \text{false}$ )

(b)

true (as  $(\text{false} \Rightarrow \text{false}) \Leftrightarrow \text{true}$ )

(c)

true (as  $(\text{false} \Rightarrow \text{true}) \Leftrightarrow \text{true}$ )

(d)

true (as  $(\text{false} \Rightarrow \text{false}) \Leftrightarrow \text{true}$ )

(Assuming that pigs can't fly . . . )

## Solution 2

(a)

$p$	$q$	$p \wedge q$	$(p \wedge q) \Rightarrow p$
t	t	t	t
t	f	f	t
f	t	f	t
f	f	f	t

(b)

$p$	$q$	$p \wedge q$	$\neg p$	$\neg p \Rightarrow (p \wedge q)$	$(\neg p \Rightarrow (p \wedge q)) \Leftrightarrow p$
t	t	t	f	t	t
t	f	f	f	t	t
f	t	f	t	f	t
f	f	f	t	f	t

(c)

$p$	$q$	$p \Rightarrow q$	$p \wedge (p \Rightarrow q)$	$(p \wedge (p \Rightarrow q)) \Rightarrow q$
t	t	t	t	t
t	f	f	f	t
f	t	t	f	t
f	f	t	f	t

### Solution 3

(a)

$$\begin{aligned}
 p \Rightarrow \neg p & \\
 & \Leftrightarrow \neg p \vee \neg p & [\Rightarrow] \\
 & \Leftrightarrow \neg p & [\text{idempotence}]
 \end{aligned}$$

(b)

$$\begin{aligned}
 \neg p \Rightarrow p & \\
 & \Leftrightarrow \neg \neg p \vee p & [\Rightarrow] \\
 & \Leftrightarrow p \vee p & [\neg \neg] \\
 & \Leftrightarrow p & [\text{idempotence}]
 \end{aligned}$$

(c)

$$\begin{aligned}
 p \Rightarrow (q \Rightarrow r) & \\
 & \Leftrightarrow \neg p \vee (q \Rightarrow r) & [\Rightarrow] \\
 & \Leftrightarrow \neg p \vee \neg q \vee r & [\Rightarrow] \\
 & \Leftrightarrow \neg p \vee \neg q \vee r & [\text{associativity}] \\
 & \Leftrightarrow \neg (p \wedge q) \vee r & [\text{De Morgan}] \\
 & \Leftrightarrow p \wedge q \Rightarrow r & [\Rightarrow]
 \end{aligned}$$

(d)

$$\begin{aligned}
 q \Rightarrow (p \Rightarrow r) & \\
 & \Leftrightarrow \neg q \vee (p \Rightarrow r) & [\Rightarrow] \\
 & \Leftrightarrow \neg q \vee \neg p \vee r & [\Rightarrow] \\
 & \Leftrightarrow \neg p \vee \neg q \vee r & [\text{associativity} \wedge \text{commutativity}] \\
 & \Leftrightarrow \neg p \vee (q \Rightarrow r) & [\Rightarrow] \\
 & \Leftrightarrow p \Rightarrow (q \Rightarrow r) & [\Rightarrow]
 \end{aligned}$$

(e)

$$\begin{aligned}
p \wedge q &\Leftrightarrow p && [\Leftrightarrow] \\
&\Leftrightarrow (p \wedge q \Rightarrow p) \wedge (p \Rightarrow p \wedge q) && [\Rightarrow] \\
&\Leftrightarrow (\neg(p \wedge q) \vee p) \wedge (\neg p \vee p \wedge q) && [\text{De Morgan}] \\
&\Leftrightarrow (\neg p \vee \neg q \vee p) \wedge (\neg p \vee p \wedge q) && [\text{associativity } \wedge \text{ comm.}] \\
&\Leftrightarrow (\neg q \vee \neg p \vee p) \wedge (\neg p \vee p \wedge q) && [\text{excluded middle}] \\
&\Leftrightarrow (\neg q \vee \text{true}) \wedge (\neg p \vee p \wedge q) && [\vee \wedge \text{true}] \\
&\Leftrightarrow \text{true} \wedge (\neg p \vee p \wedge q) && [\wedge \wedge \text{true}] \\
&\Leftrightarrow \neg p \vee p \wedge q && [\text{distribution}] \\
&\Leftrightarrow (\neg p \vee p) \wedge (\neg p \vee q) && [\text{excluded middle}] \\
&\Leftrightarrow \text{true} \wedge (\neg p \vee q) && [\wedge \wedge \text{true}] \\
&\Leftrightarrow \neg p \vee q && [\Rightarrow] \\
&\Leftrightarrow p \Rightarrow q
\end{aligned}$$

(f)

$$\begin{aligned}
p \vee q &\Leftrightarrow p && [\Leftrightarrow] \\
&\Leftrightarrow (p \vee q \Rightarrow p) \wedge (p \Rightarrow p \vee q) && [\Rightarrow] \\
&\Leftrightarrow (\neg(p \vee q) \vee p) \wedge (\neg p \vee p \vee q) && [\text{De Morgan}] \\
&\Leftrightarrow (\neg p \wedge \neg q \vee p) \wedge (\neg p \vee p \vee q) && [\text{distribution}] \\
&\Leftrightarrow (\neg p \vee p) \wedge (\neg q \vee p) \wedge (\neg p \vee p \vee q) && [\text{excluded middle}] \\
&\Leftrightarrow \text{true} \wedge (\neg q \vee p) \wedge (\neg p \vee p \vee q) && [\wedge \wedge \text{true}] \\
&\Leftrightarrow (\neg q \vee p) \wedge (\neg p \vee p \vee q) && [\text{associativity}] \\
&\Leftrightarrow (\neg q \vee p) \wedge (\text{true} \vee q) && [\text{excluded middle}] \\
&\Leftrightarrow (\neg q \vee p) \wedge \text{true} && [\vee \wedge \text{true}] \\
&\Leftrightarrow \neg q \vee p && [\wedge \wedge \text{true}] \\
&\Leftrightarrow q \Rightarrow p
\end{aligned}$$

#### Solution 4

(a)  $(p \text{ or } q) \Leftrightarrow ((\text{not } p \text{ or not } q) \text{ and } q)$  is not a tautology. You might illustrate this via a truth table or via a chain of equivalences, showing that the proposition is not equivalent to true. Alternatively, you might try and find a combination of values for which the proposition is false. (In this case, the proposition is false when  $p$  and  $q$  are both true.)

(b)  $(p \text{ or } q) \Leftrightarrow ((\text{not } p \text{ and not } q) \text{ or } q)$  is not a tautology. In this case, the proposition is false when  $p$  is true and  $q$  is false.

#### **Solution 5**

#### **Solution 6**

(a)

This is a true proposition: whatever the value of  $x$ , the expression  $x^2 - x + 1$  denotes a natural number. *If we choose to be this natural number, we will find that  $p$  is true.*

(b)

This is a false proposition. We cannot choose a large enough value for  $y$  such that  $p$  will hold for any value of  $x$ .

(c)

This is a false proposition. It is an implication whose antecedent part is true and whose consequent part is false.

(d)

This is a true proposition. It is an implication whose antecedent part is false and whose consequent part is true.

#### **Solution 7**

(a)

We must define a predicate  $p$  that is false for at least one value of  $x$ , and is true for at least one other value. A suitable solution would be  $p \Leftrightarrow x \neq 1$ .

(b)

With the above choice of  $p$ , we require only that  $q$  is sometimes false when  $p$  is true (for else the universal quantification would hold). A suitable solution would be  $q \Leftrightarrow x \neq 3$ .

### Solution 8

(a)

$$\forall x: N \bullet x \geq z$$

## Equality

### Solution 9

(d)

$$\begin{aligned} \exists x: N \bullet x = 1 \wedge x > y \vee x = 2 \wedge x > z & \\ \Leftrightarrow \exists x: N \bullet x = 1 \wedge x > y \vee \exists x: N \bullet x = 2 \wedge x > z & \\ \Leftrightarrow 1 \in N \wedge 1 > y \vee \exists x: N \bullet x = 2 \wedge x > z & \\ \Leftrightarrow 1 \in N \wedge 1 > y \vee 2 \in N \wedge 2 > z & \\ \Leftrightarrow 1 > y \vee 2 > z & \end{aligned}$$

### Solution 10

### Solution 11

### Solution 12

## Deductive proofs

### Solution 13

$$\frac{\frac{\frac{p \wedge (p \Rightarrow q)}{p} [\wedge\text{-elim}^{[1]}] \quad \frac{\frac{p \wedge (p \Rightarrow q)}{p \Rightarrow q} [\wedge\text{-elim}^{[1]}] \quad \frac{p \wedge (p \Rightarrow q)}{p} [\wedge\text{-elim}^{[1]}]}{q} [\Rightarrow \text{elim}]}{p \wedge q} [\wedge \text{intro}]}{p \wedge (p \Rightarrow q) \Rightarrow p \wedge q} [\Rightarrow\text{-intro}^{[1]}]$$

### Solution 14

In one direction:

$$\frac{\frac{\frac{\frac{\frac{}{p \wedge q} [\text{derived}]}{p \wedge q} [\Rightarrow \text{elim from } 1 \wedge 2]}{p \wedge q} [\wedge\text{-elim}^{[3]}]}{\frac{\frac{}{\neg p} [\text{assumption}]}{p} [\neg\text{-intro}^{[2]}]}{p \Rightarrow q} [\Rightarrow\text{-intro}^{[2]}]}{\frac{p \wedge q \Leftrightarrow p \Rightarrow q}{(p \wedge q \Leftrightarrow p) \Rightarrow (p \Rightarrow q)} [\Rightarrow\text{-intro}^{[1]}]}$$

and the other:

$$\frac{\frac{\frac{\frac{}{p \wedge q \Rightarrow p} [\Rightarrow\text{-intro}^{[2]}]}{p \wedge q \Rightarrow p} [\Rightarrow\text{-intro}^{[2]}]}{\frac{\frac{}{p \Rightarrow q} [\text{assumption}]}{p \wedge q \Leftrightarrow p} [\Leftrightarrow \text{intro}]}{\frac{p \Rightarrow q \Rightarrow (p \wedge q \Leftrightarrow p)}{(p \Rightarrow q) \Rightarrow (p \wedge q \Leftrightarrow p)} [\Rightarrow\text{-intro}^{[1]}]}$$

We can then combine these two proofs with  $\Leftrightarrow$  intro.

### Solution 15

$$\frac{\frac{\frac{\frac{\frac{}{p \Rightarrow q} [\text{assumption}]}{p \Rightarrow q} [\Rightarrow \text{elim}]}{q} [\text{false intro}]}{\frac{\frac{}{p} [\text{assumption}]}{false} [\text{false-elim}^{[2]}]}{\frac{(p \Rightarrow q) \wedge \neg q \Rightarrow \neg p}{(p \Rightarrow q) \wedge \neg q \Rightarrow \neg p} [\Rightarrow\text{-intro}^{[1]}]}$$

### Solution 16

In one direction:

[illegible]

In the other:

[illegible]

### Solution 17

In one direction:

$$\frac{\frac{\lceil p \vee q \wedge r \rceil^{[3]} \quad \overline{(p \vee q) \wedge (p \vee r)}}{p \vee q \wedge r \Rightarrow (p \vee q) \wedge (p \vee r)} [\Rightarrow\text{-intro}^{[3]}] \quad [\vee \text{ elim } \wedge \wedge \text{ intro}]$$

and the other:

$$\frac{\frac{\Gamma(p \vee q) \wedge (p \vee r)^{\neg[1]} \quad \Gamma p \vee q \wedge r^{\neg[2]}}{(p \vee q) \wedge (p \vee r) \Rightarrow p \vee q \wedge r} [\Rightarrow\text{-intro}^{[1]}]}{[\Rightarrow\text{-intro}^{[1]}]}$$

### Solution 18

In one direction:

$$\frac{\frac{\Gamma p \Rightarrow q^{\neg[1]} \quad \neg p \vee q}{(p \Rightarrow q) \Rightarrow \neg p \vee q} [\Rightarrow\text{-intro}^{[1]}]}{[\Rightarrow\text{-intro}^{[1]}]}$$

and the other:

$$\frac{\frac{\Gamma \neg p \vee q^{\neg[3]} \quad \frac{\frac{\Gamma p^{\neg[4]} \quad \Gamma q^{\neg[3]}}{p \Rightarrow q} [\Rightarrow\text{-intro}^{[4]}]}{\neg p \vee q \Rightarrow (p \Rightarrow q)} [\Rightarrow\text{-intro}^{[3]}]}{[\Rightarrow\text{-intro}^{[3]}]}$$

## Sets and types

### Solution 19

### Solution 20

### Solution 21

### Solution 22

### Solution 23

(a)

$$\begin{aligned} x \in a \cap a \\ \Leftrightarrow x \in a \wedge x \in a \\ \Leftrightarrow x \in a \end{aligned}$$

(b)



$$\begin{aligned}
x \in a \cup a \\
&\Leftrightarrow x \in a \vee x \in a \\
&\Leftrightarrow x \in a
\end{aligned}$$

**Solution 24**

**Solution 25**

**Solution 26**

## Relations

**Solution 27**

**Solution 28**

**Solution 29**

**Solution 30**

**Solution 31**

**Solution 32**

(a)

$$\begin{aligned}
x \mapsto y \in A \triangleleft B \triangleleft R \\
&\Leftrightarrow x \in A \wedge x \mapsto y \in (B \triangleleft R) \\
&\Leftrightarrow x \in A \wedge x \in B \wedge x \mapsto y \in R \\
&\Leftrightarrow x \in A \cap B \wedge x \mapsto y \in R \\
&\Leftrightarrow x \mapsto y \in A \cap B \triangleleft R
\end{aligned}$$

(b)

$$\begin{aligned}
x \mapsto y \in R \cup S \triangleright C & \\
& \Leftrightarrow x \mapsto y \in R \cup S \wedge y \in C \\
& \Leftrightarrow (x \mapsto y \in R \vee x \mapsto y \in S) \wedge y \in C \\
& \Leftrightarrow x \mapsto y \in R \wedge y \in C \vee x \mapsto y \in S \wedge y \in C \\
& \Leftrightarrow x \mapsto y \in R \triangleright C \vee x \mapsto y \in S \triangleright C \\
& \Leftrightarrow x \mapsto y \in (R \triangleright C) \cup (S \triangleright C)
\end{aligned}$$

## Functions

Solution 33

Solution 34

Solution 35

Solution 36

## Sequences

Solution 37

Solution 38

Solution 39

## Modelling

Solution 40

Solution 41

Solution 42

Solution 43

## Free types and induction

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Solution 45

Solution 46

Solution 47

## Supplementary material : assignment practice

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Solution 51

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