

User - Defined Squash Function

The squash operator appears in the Z notation glossary but is not built into the fuzz type-checker. However, we can define it ourselves as a user-defined function using an axiomatic specification.

Squash takes a partial function from \mathbb{N}_1 (positive naturals) to some type X and compacts it into a sequence by removing gaps in the domain. For example, $\{1 \mapsto a, 3 \mapsto b, 7 \mapsto c\}$ becomes $\langle a, b, c \rangle$.

Example 1 : Axiomatic Definition of Squash

We define squash axiomatically by specifying its key properties using only fuzz-compatible notation.

$$\begin{array}{c} \text{[X]} \\ \hline \text{squashFunc} : (\mathbb{N}_1 \rightarrow X) \rightarrow \text{seq } X \\ \hline \forall f : \mathbb{N}_1 \rightarrow X \bullet \text{ran}(\text{squashFunc}(f)) = \text{ran } f \wedge \#(\text{squashFunc}(f)) = \#(\text{ran } f) \end{array}$$

This axiomatically defines squashFunc (a user-defined version of squash) with two key properties that hold for all partial functions f from \mathbb{N}_1 to X : (1) squashFunc preserves the *range*: $\text{ran}(\text{squashFunc } f) = \text{ran } f$, and (2) the resulting sequence length equals the range cardinality: $\#(\text{squashFunc } f) = \#(\text{ran } f)$. These properties characterize squashFunc as compacting partial functions by removing domain gaps while preserving all elements and their relative order.

Example 2 : Using Squash on a Simple Function

Demonstrate squash on a sparse partial function.

$$\begin{array}{c} f1 : \mathbb{N}_1 \rightarrow \mathbb{N} \\ s1 : \text{seq } \mathbb{N} \\ \hline f1 = \{1 \mapsto 10, 3 \mapsto 30, 5 \mapsto 50\} \wedge s1 = \text{squashFunc}(f1) \end{array}$$

The result $s1$ is $\langle 10, 30, 50 \rangle$, containing exactly the range elements in the order of their original indices.

Example 3 : Squash with Letters

$$\begin{array}{c} \text{[Letter]} \\ \hline a : \text{Letter} \\ b : \text{Letter} \\ c : \text{Letter} \\ d : \text{Letter} \\ f2 : \mathbb{N}_1 \rightarrow \text{Letter} \\ s2 : \text{seq } \text{Letter} \\ \hline f2 = \{2 \mapsto a, 4 \mapsto b, 7 \mapsto c, 9 \mapsto d\} \wedge s2 = \text{squashFunc}(f2) \end{array}$$

The result $s2$ is $\langle a, b, c, d \rangle$. The gaps at positions 1, 3, 5, 6, 8 are removed.

Example 4 : Empty Function

$$\frac{\begin{array}{l} f3 : \mathbb{N}_1 \rightarrow \mathbb{N} \\ s3 : \text{seq } \mathbb{N} \end{array}}{f3 = \{\} \wedge s3 = \text{squashFunc}(f3)}$$

Squashing an empty function yields an empty *sequence* : $s3 = \langle \rangle$.

Example 5 : Continuous Function (No Gaps)

$$\frac{\begin{array}{l} f4 : \mathbb{N}_1 \rightarrow \mathbb{N} \\ s4 : \text{seq } \mathbb{N} \end{array}}{f4 = \{1 \mapsto 100, 2 \mapsto 200, 3 \mapsto 300\} \wedge s4 = \text{squashFunc}(f4)}$$

When the function has no gaps ($\text{dom } f4 = 1..3$), `squashFunc` simply converts it to a *sequence* : $s4 = \langle 100, 200, 300 \rangle$.

Example 6 : Relationship Between Function and Sequence

For any partial function f with a continuous domain starting at 1, `squashFunc` is the identity.

$$\frac{\begin{array}{l} \text{continuous_N} : \mathbb{N}_1 \rightarrow \mathbb{N} \\ s_continuous_N : \text{seq } \mathbb{N} \end{array}}{\text{dom } \text{continuous_N} = 1.. \#(\text{ran } \text{continuous_N}) \wedge s_continuous_N = \text{squashFunc}(\text{continuous_N})}$$

In this case, `continuous_N` is already a sequence, and `squashFunc continuous_N = continuous_N` as a sequence.

Example 7 : Squash Preserves Injectivity

If the input function is injective, the output sequence has no duplicates.

$$\frac{\begin{array}{l} \text{inj_func} : \mathbb{N}_1 \rightarrow \mathbb{N} \\ \text{inj_seq} : \text{seq } \mathbb{N} \end{array}}{(\forall i, j : \text{dom } \text{inj_func} \bullet \text{inj_func}(i) = \text{inj_func}(j) \Rightarrow i = j) \wedge \text{inj_seq} = \text{squashFunc}(\text{inj_func})}$$

Since `inj_func` is injective (one-to-one), `inj_seq` contains no duplicate elements.

Example 8 : Practical Use Case - Video Database

[*VideoID*, *Title*]

$$\frac{\begin{array}{l} \text{sparse_catalog} : \mathbb{N}_1 \rightarrow (\text{VideoID} \times \text{Title}) \\ \text{compact_catalog} : \text{seq } (\text{VideoID} \times \text{Title}) \end{array}}{\text{compact_catalog} = \text{squashFunc}(\text{sparse_catalog})}$$

A sparse catalog with gaps (deleted entries) can be compacted into a continuous sequence while preserving the order of remaining entries.