Propositional logic

Solution 1

(a)

$$false(as(true \Rightarrow false) \Leftrightarrow false)$$

(b)

$$true(as(false \Rightarrow false) \Leftrightarrow true)$$

(c)

$$true(as(false \Rightarrow true) \Leftrightarrow true)$$

(d)

$$true(as(false \Rightarrow false) \Leftrightarrow true)$$

(Assuming that pigs can't fly . . .)

Solution 2

(a)

p	q	$p \wedge q$	$(p \wedge q) \Rightarrow p$
t	t	t	t
t	f	f	t
f	t	f	t
f	f	f	t

(b)

p	q	$p \wedge q$	$\neg\, p$	$\neg\, p \Rightarrow (p \wedge q)$	$(\neg\, p \Rightarrow (p \wedge q)) \Leftrightarrow p$
t	t	t	f	t	t
t	f	f	f	t	t
f	t	f	t	f	t
f	f	f	t	f	t

(c)

p	q	$p \Rightarrow q$	$p \wedge (p \Rightarrow q)$	$(p \wedge (p \Rightarrow q)) \Rightarrow q$
t	t	t	t	t
t	f	f	f	t
f	t	t	f	t
f	f	t	f	t

Solution 3

(a)

$$\begin{array}{lll} p \Rightarrow \neg\, p & & \\ & \Leftrightarrow \neg\, p \vee \neg\, p & [\Rightarrow] \\ & \Leftrightarrow \neg\, p & [\text{idempotence}] \end{array}$$

(b)

$$\begin{array}{lll} \neg\, p \Rightarrow p & & \\ & \Leftrightarrow \neg\,\neg\, p \vee p & [\Rightarrow] \\ & \Leftrightarrow p \vee p & [\neg\ \neg\,] \\ & \Leftrightarrow p & [\text{idempotence}] \end{array}$$

(c)

$$\begin{array}{lll} p \Rightarrow (q \Rightarrow r) & & \\ & \Leftrightarrow \neg\, p \vee (q \Rightarrow r) & [\Rightarrow] \\ & \Leftrightarrow \neg\, p \vee \neg\, q \vee r & [\Rightarrow] \\ & \Leftrightarrow \neg\, p \vee \neg\, q \vee r & [\text{associativity}] \\ & \Leftrightarrow \neg\, (p \wedge q) \vee r & [\text{De Morgan}] \\ & \Leftrightarrow p \wedge q \Rightarrow r & [\Rightarrow] \end{array}$$

(d)

$$\begin{array}{lll} q \Rightarrow (p \Rightarrow r) & & \\ & \Leftrightarrow \neg\, q \vee (p \Rightarrow r) & [\Rightarrow] \\ & \Leftrightarrow \neg\, q \vee \neg\, p \vee r & [\Rightarrow] \\ & \Leftrightarrow \neg\, p \vee \neg\, q \vee r & [\text{associativity} \wedge \text{commutativity}] \\ & \Leftrightarrow \neg\, p \vee (q \Rightarrow r) & [\Rightarrow] \\ & \Leftrightarrow p \Rightarrow (q \Rightarrow r) & [\Rightarrow] \end{array}$$

(e)

$$
\begin{aligned}
&p \wedge q \Leftrightarrow p \\
&\qquad \Leftrightarrow (p \wedge q \Rightarrow p) \wedge (p \Rightarrow p \wedge q) && [\Leftrightarrow] \\
&\qquad \Leftrightarrow (\neg\, (p \wedge q) \vee p) \wedge (\neg\, p \vee p \wedge q) && [\Rightarrow] \\
&\qquad \Leftrightarrow (\neg\, p \vee \neg\, q \vee p) \wedge (\neg\, p \vee p \wedge q) && [\text{De Morgan}] \\
&\qquad \Leftrightarrow (\neg\, q \vee \neg\, p \vee p) \wedge (\neg\, p \vee p \wedge q) && [\text{associativity} \wedge \text{comm .}] \\
&\qquad \Leftrightarrow (\neg\, q \vee \mathit{true}) \wedge (\neg\, p \vee p \wedge q) && [\text{excluded middle}] \\
&\qquad \Leftrightarrow \mathit{true} \wedge (\neg\, p \vee p \wedge q) && [\vee \wedge \text{true}] \\
&\qquad \Leftrightarrow \neg\, p \vee p \wedge q && [\wedge \wedge \text{true}] \\
&\qquad \Leftrightarrow (\neg\, p \vee p) \wedge (\neg\, p \vee q) && [\text{distribution}] \\
&\qquad \Leftrightarrow \mathit{true} \wedge (\neg\, p \vee q) && [\text{excluded middle}] \\
&\qquad \Leftrightarrow \neg\, p \vee q && [\wedge \wedge \text{true}] \\
&\qquad \Leftrightarrow p \Rightarrow q && [\Rightarrow]
\end{aligned}
$$

(f)

$$
\begin{aligned}
&p \vee q \Leftrightarrow p \\
&\qquad \Leftrightarrow (p \vee q \Rightarrow p) \wedge (p \Rightarrow p \vee q) && [\Leftrightarrow] \\
&\qquad \Leftrightarrow (\neg\, (p \vee q) \vee p) \wedge (\neg\, p \vee p \vee q) && [\Rightarrow] \\
&\qquad \Leftrightarrow (\neg\, p \wedge \neg\, q \vee p) \wedge (\neg\, p \vee p \vee q) && [\text{De Morgan}] \\
&\qquad \Leftrightarrow (\neg\, p \vee p) \wedge (\neg\, q \vee p) \wedge (\neg\, p \vee p \vee q) && [\text{distribution}] \\
&\qquad \Leftrightarrow \mathit{true} \wedge (\neg\, q \vee p) \wedge (\neg\, p \vee p \vee q) && [\text{excluded middle}] \\
&\qquad \Leftrightarrow (\neg\, q \vee p) \wedge (\neg\, p \vee p \vee q) && [\wedge \wedge \text{true}] \\
&\qquad \Leftrightarrow (\neg\, q \vee p) \wedge (\neg\, p \vee p \vee q) && [\text{associativity}] \\
&\qquad \Leftrightarrow (\neg\, q \vee p) \wedge (\mathit{true} \vee q) && [\text{excluded middle}] \\
&\qquad \Leftrightarrow (\neg\, q \vee p) \wedge \mathit{true} && [\vee \wedge \text{true}] \\
&\qquad \Leftrightarrow \neg\, q \vee p && [\wedge \wedge \text{true}] \\
&\qquad \Leftrightarrow q \Rightarrow p && [\Rightarrow]
\end{aligned}
$$

Solution 4

(a) (p or q) $\Leftrightarrow$ ((not p or not q) and q) is not a tautology. You might illustrate this via a truth table or via a chain of equivalences, showing that the proposition is not equivalent to true. Alternatively, you might try and find a combination of values for which the proposition is false. (In this case, the proposition is false when p and q are both true.)

(b) (p or q) ⇔ ((not p and not q) or q) is not a tautology. In this case, the proposition is false when p is true and q is false.

Solution 5

(a)

$$\exists\, d : \mathit{Dog} \bullet \mathit{gentle}(d) \land \mathit{well\_trained}(d)$$

(b)

$$\forall\, d : \mathit{Dog} \bullet \mathit{neat}(d) \land \mathit{well\_trained}(d) \Rightarrow \mathit{attractive}(d)$$

(c)

(Requires nested quantifier in implication - parser limitation)

Solution 6

(a)

This is a true $proposition : whatever the value of x, the expression x^2$ - x + 1 denotes a natural number. If we choose y to be this natural number, we will find that p is true.

(b)

This is a false proposition. We cannot choose a large enough value for y such that p will hold for any value of x.

(c)

This is a false proposition. It is an implication whose antecedent part is true and whose consequent part is false.

(d)

This is a true proposition. It is an implication whose antecedent part is false and whose consequent part is true.

Solution 7

(a)

We must define a predicate p that is false for at least one value of x, and is true for at least one other value. A suitable solution would be $p \Leftrightarrow x > 1$.

(b)

With the above choice of p, we require only that q is sometimes false when p is true (for else the universal quantification would hold). A suitable solution would be $q \Leftrightarrow x > 3$.

Solution 8

(a)

$$\forall\, x : N \bullet x \geq z$$

Equality

Solution 9

(d)

$$\begin{aligned}
&\exists\, x : N \bullet x = 1 \land x > y \lor x = 2 \land x > z \\
&\qquad \Leftrightarrow \exists\, x : N \bullet x = 1 \land x > y \lor \exists\, x : N \bullet x = 2 \land x > z \\
&\qquad \Leftrightarrow 1 \in N \land 1 > y \lor \exists\, x : N \bullet x = 2 \land x > z \\
&\qquad \Leftrightarrow 1 \in N \land 1 > y \lor 2 \in N \land 2 > z \\
&\qquad \Leftrightarrow 1 > y \lor 2 > z
\end{aligned}$$

Solution 10

As discussed, the quantifier $\exists_1$ can help give rise to a 'test' or 'precondition' to ensure that an application of mu will work.

So, as a simple example, as the proposition

$$\exists_1 n\colon N \bullet \forall m\colon N \bullet n \leq m$$

is equivalent to true, we can be certain that the statement

$$\mu n\colon N \bullet \forall m\colon N \bullet n \leq m$$

will return a result (which happens to be 0).

Solution 11

(a)

$(\mu a\colon N \bullet a = a) = 0$ is a provable statement, since 0 is the only natural number with the specified property.

(b)

$(\mu b\colon N \bullet b = b) = 1$ is not provable. The specified property is true of both 0 and 1, and thus the value of the mu-expression is undefined.

(c)

$(\mu c\colon N \bullet c > c) = (\mu c\colon N \bullet c > c)$ is a provable statement. Neither expression is properly defined, but we may conclude that they are equal; there is little else that we can prove about them.

(d)

$(\mu d\colon N \bullet d = d) = 1$ is not a provable statement. We cannot confirm that 1 is the only natural number with the specified property; we do not know what value is taken by undefined operations.

Solution 12

(Requires mu-operator with expression part - not yet implemented)

(a)

$$\mu\, m\colon \mathit{Mountain} \mid \forall\, n\colon \mathit{Mountain} \bullet \mathit{height}(n) \leq \mathit{height}(m) \bullet \mathit{height}(m)$$

(b)

$$\mu\, c\colon \mathit{Chapter} \mid \exists_1\, d\colon \mathit{Chapter} \bullet \mathit{length}(d) > \mathit{length}(c) \bullet \mathit{length}(c)$$

(c)

Assuming the existence of a suitable function, max: $(\mu\, n\colon N \bullet n = \mathit{max}(\{m\colon N \mid 8 * m < 100.8 * m\})$. 100 - n)

Deductive proofs

Solution 13

$$\dfrac{\ulcorner p \wedge (p \Rightarrow q)\urcorner^{[1]} \qquad \dfrac{\dfrac{p \wedge (p \Rightarrow q)}{p}\;[\wedge\text{-elim}^{[1]}] \qquad \dfrac{\dfrac{p \wedge (p \Rightarrow q)}{p \Rightarrow q}\;[\wedge\text{-elim}^{[1]}] \qquad \dfrac{p \wedge (p \Rightarrow q)}{p}\;[\wedge\text{-elim}^{[1]}]}{q}\;[\Rightarrow \text{elim}]}{p \wedge q}\;[\wedge \text{ intro}]}{p \wedge (p \Rightarrow q) \Rightarrow p \wedge q}\;[\Rightarrow\text{-intro}^{[1]}]$$

Solution 14

In one direction:

$$\dfrac{\ulcorner p \wedge q \Leftrightarrow p\urcorner^{[1]} \qquad \dfrac{\ulcorner p\urcorner^{[2]} \qquad \dfrac{\dfrac{\dfrac{}{p \wedge q}\;[\text{derived}]}{p \wedge q}\;[\Rightarrow \text{elim from } 1 \wedge 2]}{q}\;[\wedge\text{-elim}^{[3]}]}{p \Rightarrow q}\;[\Rightarrow\text{-intro}^{[2]}]}{(p \wedge q \Leftrightarrow p) \Rightarrow (p \Rightarrow q)}\;[\Rightarrow\text{-intro}^{[1]}]$$

and the other:

$$\dfrac{\ulcorner p \Rightarrow q \urcorner^{[1]} \quad \dfrac{\dfrac{\ulcorner p \wedge q \urcorner^{[2]} \quad \ulcorner p \urcorner^{[2]}}{p \wedge q \Rightarrow p}\;[\Rightarrow\text{-intro}^{[2]}] \quad \dfrac{\ulcorner p \urcorner^{[3]} \quad \ulcorner p \wedge q \urcorner^{[1]}}{p \Rightarrow p \wedge q}\;[\Rightarrow\text{-intro}^{[3]}]}{p \wedge q \Leftrightarrow p}\;[\Leftrightarrow \text{ intro}]}{(p \Rightarrow q) \Rightarrow (p \wedge q \Leftrightarrow p)}\;[\Rightarrow\text{-intro}^{[1]}]$$

We can then combine these two proofs *with* $\Leftrightarrow$ *intro*.

Solution 15

$$\dfrac{\ulcorner (p \Rightarrow q) \wedge \neg\, q \urcorner^{[1]} \quad \dfrac{\ulcorner p \urcorner^{[2]} \quad \dfrac{\dfrac{\ulcorner p \Rightarrow q \urcorner^{[1]} \quad \ulcorner p \urcorner^{[2]}}{q}\;[\Rightarrow \text{ elim}] \quad \ulcorner \neg\, q \urcorner^{[1]}}{\mathit{false}}\;[\text{false intro}]}{\neg\, p}\;[\text{false-elim}^{[2]}]}{(p \Rightarrow q) \wedge \neg\, q \Rightarrow \neg\, p}\;[\Rightarrow\text{-intro}^{[1]}]$$

Solution 16

In one direction:

$$\dfrac{\ulcorner p \wedge (q \vee r) \urcorner^{[1]} \quad \dfrac{\ulcorner q \vee r \urcorner^{[1]} \quad \dfrac{\dfrac{\ulcorner p \urcorner^{[1]} \quad \dfrac{}{q}\;[\text{case assumption}]}{p \wedge q}\;[\wedge \text{ intro}]}{p \wedge q \vee p \wedge r}\;[\vee \text{ intro}] \quad \dfrac{\dfrac{\ulcorner p \urcorner^{[1]} \quad \dfrac{}{r}\;[\text{case assumption}]}{p \wedge r}\;[\wedge \text{ intro}]}{p \wedge q \vee p \wedge r}\;[\vee \text{ intro}]}{p \wedge q \vee p \wedge r}\;[\vee\text{-elim}^{[2]}]}{p \wedge (q \vee r) \Rightarrow p \wedge q \vee p \wedge r}\;[\Rightarrow\text{-intro}^{[1]}]$$

In the other:

$$
\dfrac{\ulcorner p \wedge q \vee p \wedge r \urcorner^{[3]} \qquad \dfrac{\ulcorner case1 \vee case2 \urcorner^{[3]} \qquad \dfrac{\dfrac{\dfrac{}{p}\;[\wedge \text{ elim}]}{q \vee r}\;[\vee \text{ intro}]}{p \wedge (q \vee r)}\;[\wedge \text{ intro}] \qquad \dfrac{\dfrac{\dfrac{}{p}\;[\wedge \text{ elim}]}{q \vee r}\;[\vee \text{ intro}]}{p \wedge (q \vee r)}\;[\wedge \text{ intro}]}{p \wedge (q \vee r)}\;[\vee\text{-elim}^{[4]}]}{p \wedge q \vee p \wedge r \Rightarrow p \wedge (q \vee r)}\;[\Rightarrow\text{-intro}^{[3]}]
$$

Solution 17

In one direction:

$$
\dfrac{\ulcorner p \vee q \wedge r \urcorner^{[3]} \qquad \dfrac{}{(p \vee q) \wedge (p \vee r)}\;[\vee \text{ elim} \wedge \wedge \text{ intro}]}{p \vee q \wedge r \Rightarrow (p \vee q) \wedge (p \vee r)}\;[\Rightarrow\text{-intro}^{[3]}]
$$

and the other:

$$
\dfrac{\ulcorner (p \vee q) \wedge (p \vee r) \urcorner^{[1]} \qquad \ulcorner p \vee q \wedge r \urcorner^{[2]}}{(p \vee q) \wedge (p \vee r) \Rightarrow p \vee q \wedge r}\;[\Rightarrow\text{-intro}^{[1]}]
$$

Solution 18

In one direction:

$$
\dfrac{\ulcorner p \Rightarrow q \urcorner^{[1]} \qquad \neg\, p \vee q}{(p \Rightarrow q) \Rightarrow \neg\, p \vee q}\;[\Rightarrow\text{-intro}^{[1]}]
$$

and the other:

$$
\dfrac{\ulcorner \neg\, p \vee q \urcorner^{[3]} \qquad \dfrac{\ulcorner p \urcorner^{[4]} \qquad \ulcorner q \urcorner^{[3]}}{p \Rightarrow q}\;[\Rightarrow\text{-intro}^{[4]}]}{\neg\, p \vee q \Rightarrow (p \Rightarrow q)}\;[\Rightarrow\text{-intro}^{[3]}]
$$

Sets and types

Solution 19

(a)

1 in $\{4, 3, 2, 1\}$ is true.

(b)

$\{1\}$ in $\{1, 2, 3, 4\}$ is undefined.

(c)

$\{1\}$ in $\{\{1\}, \{2\}, \{3\}, \{4\}\}$ is true.

(d)

The empty set in $\{1, 2, 3, 4\}$ is undefined.

Solution 20

(a)

$$\{1\} \times \{2, 3\}$$

is the set $\{(1, 2), (1, 3)\}$

(b)

The empty set cross $\{2, 3\}$ is the empty set

(c)

$$\mathbb{P}\,\emptyset \times \{1\}$$

is the set $\{(\emptyset, 1)\}$

(d)

$\{(1,2)\}$ cross $\{3,4\}$ is the set $\{((1,2),3),((1,2),4)\}$

Solution 21

There are various ways of describing these sets via set comprehensions. Examples are given below.

(a)

$$\{z \colon Z \mid 0 \leq z \wedge z \leq 100\}$$

(b)

$$\{z \colon Z \mid z = 10\}$$

(c)

$$\{z \colon Z \mid z \bmod 2 = 0 \vee z \bmod 3 = 0 \vee z \bmod 5 = 0\}$$

Solution 22

(a)

$$\{n \colon N \mid n \leq 4 \bullet n^2\}$$

(b)

$$\{n \colon N \mid n \leq 4 \bullet (n, n^2)\}$$

(c)

$n : P0,\ 1$

(d)

$\{n\colon \mathbb{P}\{0,1\} \mid \mathit{true} \bullet (n, \#n)\}$

Solution 23

(a)

$$\begin{aligned} & x \in a \cap a \\ & \quad \Leftrightarrow x \in a \land x \in a \\ & \quad \Leftrightarrow x \in a \end{aligned}$$

(b)

$$\begin{aligned} & x \in a \cup a \\ & \quad \Leftrightarrow x \in a \lor x \in a \\ & \quad \Leftrightarrow x \in a \end{aligned}$$

Solution 24

(a)

The set of all pairs of integers is Z cross Z. To give it a name, we could write:

Pairs == Z $\times Z$

(b)

The set of all integer pairs in which each element is strictly greater than zero could be defined by:

StrictlyPositivePairs == { m, n : $Z \mid m > 0 \land n > 0 \bullet (m, n)\}$

(c)

It is intuitive to use a singular noun for the name of a basic type; we define the set of all people by writing:

$[Person]$

(d)

The set of all couples could be defined by:

Couples == { s : $\mathbb{P}$ $Person \mid \#s = 2\}$

Solution 25

(Requires generic set notation and Cartesian product)

Solution 26

(Requires generic parameters and relation type notation)

Relations

Solution 27

(a)

The power set of $\{(0,0),(0,1),(1,0),(1,1)\}$ is:

$\{\emptyset, \{(0,0)\}, \{(0,1)\}, \{(1,0)\}, \{(1,1)\}, \{(1,0),(1,1)\}, \{(0,0),(0,1)\}, \{(0,1),(1,1)\}, \{(0,1),(1,0)\}, \{(0,0),(1,1)\}, \{$

(b)

$$\{\emptyset, \{(0,0)\}, \{(0,1)\}, \{(0,0),(0,1)\}\}$$

(c)

$$\{\emptyset\}$$

(d)

$$\{\emptyset\}$$

Solution 28

(a)

$$\operatorname{dom} R = \{0, 1, 2\}$$

(b)

$$\operatorname{ran} R = \{1, 2, 3\}$$

(c)

$$\{1, 2\} \lhd R = \{1 \mapsto 2, 1 \mapsto 3, 2 \mapsto 3\}$$

Solution 29

(a)

$$\{2 \mapsto 4, 3 \mapsto 3, 3 \mapsto 4, 4 \mapsto 2\}$$

(b)

$$\{1 \mapsto 3, 2 \mapsto 2, 2 \mapsto 3, 3 \mapsto 1\}$$

(c)

$$\{1 \mapsto 1, 2 \mapsto 2, 2 \mapsto 3, 3 \mapsto 2, 3 \mapsto 3, 4 \mapsto 4\}$$

(d)

$$\{1 \mapsto 4, 2 \mapsto 2, 2 \mapsto 3, 3 \mapsto 2, 3 \mapsto 3, 4 \mapsto 1\}$$

Solution 30

| $childOf : Person \leftrightarrow Person$

(a)

parentOf == childOf^{-1}

This is a good example of how there are many different ways of writing the same thing. An alternative abbreviation is:

parentOf == { x, y : $Person \mid x \mapsto y \in childOf \bullet y \mapsto x\}$

Or, via an axiomatic definition:

$$\begin{array}{|l} parentOf : Person \leftrightarrow Person \\ \hline parentOf = childOf^{-1} \end{array}$$

(b)

siblingOf == (childOf $\circ parentOf) \setminus id$

(c)

cousinOf == childOf $\circ siblingOf \circ parentOf$

(d)

ancestorOf == parentOf$^+$

Solution 31

(Requires compound identifiers with operators - R+, R*)

(a)

R == { a, b : $N \mid b = a \vee b = a\}$

(b)

S == { a, b : $N \mid b = a \vee b = a\}$

(c)

R+ == $\{a, b \colon N \mid b > a\}$

(d)

R* == $\{a, b \colon N \mid b \geq a\}$

Solution 32

(a)

$$
\begin{aligned}
&x \mapsto y \in A \lhd B \lhd R \\
&\qquad \Leftrightarrow x \in A \wedge x \mapsto y \in (B \lhd R) \\
&\qquad \Leftrightarrow x \in A \wedge x \in B \wedge x \mapsto y \in R \\
&\qquad \Leftrightarrow x \in A \cap B \wedge x \mapsto y \in R \\
&\qquad \Leftrightarrow x \mapsto y \in A \cap B \lhd R
\end{aligned}
$$

(b)

$$
\begin{aligned}
&x \mapsto y \in R \cup S \rhd C \\
&\qquad \Leftrightarrow x \mapsto y \in R \cup S \wedge y \in C \\
&\qquad \Leftrightarrow (x \mapsto y \in R \vee x \mapsto y \in S) \wedge y \in C \\
&\qquad \Leftrightarrow x \mapsto y \in R \wedge y \in C \vee x \mapsto y \in S \wedge y \in C \\
&\qquad \Leftrightarrow x \mapsto y \in R \rhd C \vee x \mapsto y \in S \rhd C \\
&\qquad \Leftrightarrow x \mapsto y \in (R \rhd C) \cup (S \rhd C)
\end{aligned}
$$

Functions

Solution 33

The set of 9 functions:

$$\{\emptyset, \{(0,0)\}, \{(0,1)\}, \{(1,1)\}, \{(1,0)\}, \{(0,0),(1,1)\}, \{(0,1),(1,1)\}, \{(1,0),(0,0)\}, \{(0,1),(1,0)\}\}$$

(a)

The set of total functions:

$$\{\{(0,0),(1,1)\}, \{(0,1),(1,1)\}, \{(1,0),(0,0)\}, \{(0,1),(1,0)\}\}$$

(b)

The set of functions which are neither injective nor surjective:

$$\{\{(0,1),(1,1)\},\{(0,0),(1,0)\}\}$$

(c)

The set of functions which are injective but not surjective:

$$\{\emptyset,\{(0,0)\},\{(0,1)\},\{(1,0)\},\{(1,1)\}\}$$

(d)

There are no functions (of this type) which are surjective but not injective.

(e)

The set of bijective functions:

$$\{\{(0,0),(1,1)\},\{(0,1),(1,0)\}\}$$

Solution 34

(a)

$$\{1 \mapsto a, 2 \mapsto b, 3 \mapsto c, 4 \mapsto b\}$$

(b)

$$\{1 \mapsto c, 2 \mapsto b, 3 \mapsto c, 4 \mapsto d\}$$

(c)

$$\{1 \mapsto c, 2 \mapsto b, 3 \mapsto c, 4 \mapsto b\}$$

(d)

$$\{1 \mapsto c, 2 \mapsto b, 3 \mapsto c, 4 \mapsto b\}$$

Solution 35

(Requires power set notation P and relational image)

(a)

$$\begin{array}{|l} children : Person \to \mathbb{P}\ Person \\ \hline children = \{p \colon Person \bullet p \mapsto parentOf(\!|\ \{p\}\ |\!)\} \end{array}$$

(b)

$$\begin{array}{|l} number\_of\_grandchildren : Person \to N \\ \hline number\_of\_grandchildren = \{p \colon Person \bullet p \mapsto \#parentOf \circ parentOf(\!|\ \{p\}\ |\!)\} \end{array}$$

Solution 36

$(Note : Thissolutiondemonstratesrelationtypesinquantifierdomains)$

$$\begin{array}{|l} number\_of\_drivers : Drivers \leftrightarrow Cars \to (Cars \to N) \\ \hline number\_of\_drivers = \lambda\, r : Drivers \leftrightarrow Cars \bullet \{c \colon \operatorname{ran} r \bullet c \mapsto \#\{d \colon Drivers \mid d \mapsto c \in r\}\} \end{array}$$

Sequences

Solution 37

(a)

$$\langle a \rangle$$

(b)

$$\{1 \mapsto a, 2 \mapsto b, 2 \mapsto a, 3 \mapsto c, 3 \mapsto b, 4 \mapsto d\}$$

(c)

$$\{2 \mapsto b, 3 \mapsto c, 4 \mapsto d\}$$

(d)

$$\{1, 2, 3, 4\}$$

(e)

$$\{a, b\}$$

(f)

$$\{a \mapsto 1, b \mapsto 2, c \mapsto 3, d \mapsto 4\}$$

(g)

$$\langle a, b \rangle$$

(h)

$$\{3 \mapsto b\}$$

(i)

$$\{a\}$$

(j)

$$c$$

Solution 38

(a)

$$\begin{array}{|l} f : \mathit{Place} \to \mathbb{P}\ \mathit{Place} \\ \hline \forall\, p : \mathit{Place} \bullet f(p) = \{q : \mathit{Place} \mid p \mapsto q \in \operatorname{ran} \mathit{trains}\} \end{array}$$

(b)

$$\{p\colon Place \mid \exists_1 x\colon \operatorname{dom} trains \bullet trains(x).2 = p\}$$

(c)

$$\mu\, p\colon Place \bullet \forall\, q\colon Place \bullet p \neq q \land \#\{x\colon \operatorname{dom} trains \mid trains(x).2 = p\} > \#\{x\colon \operatorname{dom} trains \mid trains(x).2 = q\}$$

Solution 39

(a)

$large\_coins : Collection \to N$

$\forall\, c\colon Collection \bullet large\_coins(c) = c(large)$

(Blocked $by : underscoreinidentifierforfuzzcompatibility$)

(b)

$add\_coin : Collection * Coin \to Collection$

$\forall\, c\colon Collection \bullet \forall\, d\colon Coin \bullet add\_coin(c, d) = c \cup [\![d]\!]$

(Blocked $by : underscoreinidentifierandbagunion$)

Modelling

Solutions 40-52 are work in progress - many require features not yet implemented

Solution 40

(Work in progress - requires semicolon-separated bindings in set comprehensions)

(a)

$hd : seq(Title * Length * Viewed)$

$cumulative\_total(hd) \leq 12000$

$\forall p\colon \mathrm{ran}\, hd \bullet p.2 \leq 360$

Note that cumulative$_t otalisdefinedinpart(d)$.

(b)

$$\{p\colon \mathrm{ran}\, hd \mid p.2 > 120 \bullet p.1\}$$

(c)

These can be defined recursively:

$viewed : \mathrm{seq}\ Programme \rightarrow \mathrm{seq}\ Programme$

$viewed(\langle\rangle) = \langle\rangle \wedge \forall x\colon Programme \bullet \forall s\colon\ \mathrm{seq}\ Programme \bullet viewed(\langle x\rangle \frown s) = (\text{if } x.3 = yes \text{ then } \langle x\rangle \frown$

or otherwise (omitted - requires semicolon-separated bindings in set comprehension)

(d)

$cumulative\_total : \mathrm{seq}\ Title * Length * Viewed \rightarrow N$

$cumulative\_total(\langle\rangle) = 0\, \forall x\colon\ Title * Length * Viewed \bullet \forall s\colon\ \mathrm{seq}\ Title * Length * Viewed \bullet cumulative\_$

(e)

(mu p : $ranhd$ $|\forall q\colon\ \mathrm{ran}\, hd \bullet p \neq q \wedge p.2 > q.2$— p.1)

(This, of course, assumes that there is a unique element with this property.)

(f)

(f) Omitted - requires semicolon-separated bindings in nested set comprehension

(g)

axdef

$g : seq(Title * Length * Viewed) \rightarrow seq(Title * Length * Viewed)$

where

$\forall s \colon \mathrm{seq}\ Title{*}Length{*}Viewed \bullet g(s) = s$—¿ $\{x \colon \mathrm{ran}\, s \mid x \neq longest\_viewed(s)\}$

end

Where longest$_v$*iewedisdefinedas*

axdef

$longest\_viewed : seq(Title * Length * Viewed) + \rightarrow Title * Length * Viewed$

where

$\forall s \colon \mathrm{seq}\ Title * Length * Viewed \bullet longest_viewed(s) = (\mu\, p \colon \mathrm{ran}\, s \bullet p.3 = yes$and $\forall q \colon \mathrm{ran}\, s \bullet p \neq q \wedge q.3 = yes \wedge p.2 > q.2)$

end

(Blocked $by : nestedquantifiersinmuexpressions - parserlimitation$)

This, of course, assumes that there is at least one viewed programme (and one of a unique maximum length).

(h)

$$\begin{array}{|l} s : \mathrm{seq}\ Title * Length * Viewed \rightarrow \mathrm{seq}\ Title * Length * Viewed \\ \hline \forall x : \mathrm{seq}\ Title * Length * Viewed \bullet items(s(x)) = items(x) \land \forall i, j : \mathrm{dom}\, s(x) \bullet i < j \Rightarrow s(x)(i).2 \geq s(x \end{array}$$

Solution 41

(a)

axdef

$records : Year- |> Table$

where

dom(records) = 1993..current

$\forall y : \mathrm{dom}\, records \bullet \#records(y) \leq 50$

$\forall\ y : dom(records)\ |\forall e : \mathrm{ran}\, records(y) \bullet year(e.1) = y$

$\forall\ r : ran(records)\ |\forall i1, i2 : \mathrm{dom}\, r \bullet i1 \neq i2 \land r(i1).1 = r(i2).1 \Rightarrow r(i1).3 \neq r(i2).3$

end

(Blocked $by : nestedquantifiersinpredicates - parserlimitation$)

(b)

(i)

$$\{e\colon Entry \mid \exists\, r\colon \operatorname{ran} records \bullet e \in \operatorname{ran} r \land e.3 = 479\}$$

ii

$$\{e\colon Entry \mid \exists\, r\colon \operatorname{ran} records \bullet e \in \operatorname{ran} r \land e.6 > e.5\}$$

iii

$$\{e\colon Entry \mid \exists\, r\colon \operatorname{ran} records \bullet e \in \operatorname{ran} r \land e.7 \geq 70\}$$

iv

$$\{c\colon Course \mid \forall\, r\colon \operatorname{ran} records \bullet \forall\, e\colon \operatorname{ran} r \bullet e.2 = c \Rightarrow e.7 \geq 70\}$$

v

$y : Year \mid yindomrecords.$ y —-¿ $l : Lecturer \mid c : ran(recordsy) \mid c.4 =$ l ¿ 6

(c)

axdef

where

$\forall\, x\colon Entry \bullet \forall\, s\colon \operatorname{seq} Entry \bullet 479_c ourses(\langle\rangle) = \langle\rangle$ and $479_c ourses(\langle x \rangle\ ^s) =$ $if x.3 = 479$then ¡x>^{4} $79_c ourses(s) else 479_c ourses(s)$

end

(Blocked $by : underscoreinidentifier - usecamelCaseforfuzzcompatibility$)

(d)

$\forall x\colon Entry \bullet \forall s\colon \mathrm{seq}\ Entry \bullet total(\langle\rangle) = 0 \land total(\langle x\rangle \frown s) = x.5 + total(s)$

Solution 42

$[Person]$

axdef

$State : P(seq(iseq(Person)))$

where

$\forall\ s : State\ |\forall\, i, j\colon\ \mathrm{dom}\, s \bullet i \neq j \land \mathrm{ran}\, s(i) \cap \mathrm{ran}\, s(j) = \{\}$

end

(Blocked $by : nestedquantifierswithsemicolonbindings - parserlimitation$)

(b)

axdef

$add : N * Person * State - |> State$

where

$\forall n\colon N \bullet \forall p\colon Person \bullet \forall s\colon State \bullet n \in \mathrm{dom}\, s \land p \notin \bigcup \mathrm{ran}\,\mathrm{ran}\, s$—

add(n, p, s) = s ++ $\{n \mapsto s(n)^{\langle p\rangle}\}$

end

(Blocked by: -—¿ operator not implemented)

Solution 43

(a)

(i) $\forall\ i : dombookings\ |\forall\, x, y\colon\ bookings(i) \bullet x \neq y \wedge x.2 \mathinner{..} x.3 \cap y.2 \mathinner{..} y.3 = \{\}$

(ii) $\forall\, i\colon$ dom $bookings \bullet forall$l $x : bookings(i)\ |\{x.2, x.3\}$ subseteq 1..max(i.1)

(iii) $\forall\ i : dombookings\ |\forall\, b\colon\ bookings(i) \bullet b.2 \leq b.3$

(iv) This is enforced by the constraint for part (i).

(Blocked $by : nestedquantifiers - parserlimitation)$

(b)

(i) $\{i\colon$ dom $bookings \mid i.1 = Banbury \bullet i.2\}$

(ii) $\{i\colon$ dom $bookings \mid i.1 = Banbury \wedge \exists\, b\colon\ bookings(i) \bullet 50 \in b.2 \mathinner{..} b.3\}$

(iii) $r : Room;\ s : N\ |\exists\, i\colon$ dom $bookings \bullet i.1 = r \wedge i.2 = s.$ (r, s)

(iv) $r : Room\ |\exists\, i\colon$ dom $bookings \bullet i.1 = r \wedge \#bookings(i) \geq 10$

(Blocked $by : semicolonbindingsinsetcomprehensionsandnestedquantifiers)$

Free types and induction

$[N]$

$Tree ::= stalk \mid leaf\langle\langle N\rangle\rangle \mid branch\langle\langle Tree \times Tree\rangle\rangle$

Solution 44

The two cases of the proof are established by equational $reasoning : thefirstby$

reverse $(\langle\rangle\ ^{t}) = reverset[cat.1a] = (reverset)\langle\rangle$ [cat.1b]

where cat.1a is $\langle\rangle$ s = s$andcat.1biss\langle\rangle$ = s

and the second by

reverse $((\langle x\rangle\ ^{u})^{t}) = reverse(\langle x\rangle\ ^{(}u^{t}))[cat.2]$

= reverse (u $^{t})\langle x\rangle$ [reverse.2]

= (reverse t $^{r}everseu)\langle x\rangle$ [anti-distributive]

= reverse t $^{(}reverseu\langle x\rangle)$ [cat.2]

= reverse t $^{r}everse(\langle x\rangle\ ^{u})[reverse.2]$

Solution 45

The base case:

reverse (reverse $\langle\rangle$) = reverse $\langle\rangle$ [reverse.1] = $\langle\rangle$ [reverse.1]

The inductive step:

reverse (reverse ($\langle x \rangle$ ^{t}))

= reverse ((reverse t) $\langle x \rangle$) [reverse.2]

= reverse ($\langle x \rangle$) $^r everse(reverset)[anti - distributive]$

= reverse ($\langle x \rangle$ $\langle\rangle$) $^r everse(reverset)[cat.1]$

= ((reverse $\langle\rangle$) $\langle x \rangle$) $^r everse(reverset)[reverse.2]$

= ($\langle\rangle$ $\langle x \rangle$) $^r everse(reverset)[reverse.1]$

= $\langle x \rangle$ $^r everse(reverset)[cat.1]$

= $\langle x \rangle$ $^t[reverse(reverset) = t]$

Solution 46

(a)

$count : Tree \to N$

count $stalk = 0$

$\forall\, n : N \bullet count(\text{leaf n}) = 1$

$\forall\, t1, t2 : Tree \bullet count(branch(t1, t2)) = count\text{t1} + \text{count t2}$

(Blocked $by : recursivefreetypesandpatternmatching$)

(b)

$flatten : Tree \to seqN$

flatten stalk = $\langle\rangle$

$\forall\, n : N \bullet flatten(\text{leaf n}) = \langle n \rangle$

$\forall\, t1, t2 : Tree \bullet flatten(branch(t1, t2)) = flatten\text{t1}\ ^f lattent2$

(Blocked $by : recursivefreetypesandpatternmatching$)

Solution 47

First, exhibit the induction principle for the free type:

P stalk and ($\forall\, n : N \bullet P(\text{leaf n})$) and ($\forall\, t1, t2 : Tree \bullet \mathbb{P}\ t1 \land \mathbb{P}\ t2 \Rightarrow \mathbb{P}\ branch(t1, t2)$)

implies $\forall\, t : Tree \bullet \mathbb{P}\ t$

This gives three cases for the proof:

(flatten stalk) = $\langle\rangle$ [flatten] = 0 [] = count stalk [count]

(Remaining cases omitted - require equational reasoning with recursive functions)

Supplementary material : assignment practice

Solution 48

$[\mathit{SongId}, \mathit{UserId}, \mathit{PlaylistId}, \mathit{Playlist}]$

$\mathit{songs} : \mathbb{F}\ \mathit{SongIdusers} : \mathbb{F}\ \mathit{UserIdplaylists} : \mathit{PlaylistId} \pfun \mathit{Playlistplaylist Owner} : \mathit{PlaylistId} \pfun \mathit{UserIdpla}$

$\forall\, i : \operatorname{dom} \mathit{playlists} \bullet \operatorname{ran} \mathit{playlists}(i) \mathit{subseteqsongs} \operatorname{dom} \mathit{playlistOwnersubseteq} \operatorname{dom} \mathit{playlists} \operatorname{ran} \mathit{playlistOu}$

Solution 49

$\mathit{hated} : \mathit{UserId} \pfun \mathbb{F}\ \mathit{SongIdloved} : \mathit{UserId} \pfun \mathbb{F}\ \mathit{SongId}$

$\operatorname{dom} \mathit{hatedsubsetequsers} \,\forall\, i : \operatorname{dom} \mathit{hated} \bullet \mathit{hated}(i) \mathit{subseteqsongs} \operatorname{dom} \mathit{lovedsubsetequsers} \,\forall\, i : \operatorname{dom} \mathit{loved} \bullet$

Solution 50

(a)

abbrev

A == users $\setminus \bigcup \operatorname{ran} \mathit{playlistSubscribers}$

(b)

abbrev

B == { p : $\operatorname{dom} \mathit{playlistSubscribers} \mid \#\mathit{playlistSubscribers}(p) \geq 100\}$

(c)

C == $\mu\, u : \operatorname{dom} \mathit{loved} \bullet \forall\, v : \operatorname{dom} \mathit{loved} \bullet u \neq v \wedge \#\mathit{loved}(u) > \#\mathit{loved}(v)$

(d)

D == $\mu\, s : \mathit{songs} \bullet \forall\, t : \mathit{songs} \bullet s \neq t \wedge \#\{u : \mathit{UserId} \mid s \in \mathit{loved}(u)\} > \#\{u : \mathit{UserId} \mid t \in \mathit{loved}(u)\}$

Solution 51

(a)

Let's first define two helper functions:

$loveHateScore : SongId+ \rightarrow N$

$\forall\ i : songs \mid \{u\colon UserId \mid i \in loved(u)\}$ ¿= $\{u\colon UserId \mid i \in hated(u)\} \Rightarrow$

loveHateScore(i) = $\{u\colon UserId \mid i \in loved(u)\}$ - $\{u\colon UserId \mid i \in hated(u)\}$

and

$\forall\ i : songs \mid \{u\colon UserId \mid i \in loved(u)\}$ ¡ $\{u\colon UserId \mid i \in hated(u)\} \Rightarrow$

loveHateScore(i) = 0

$$\begin{array}{|l} playlistCount : SongId \twoheadrightarrow N \\ \hline \forall\, i\colon songs \bullet playlistCount(i) = \#\{p\colon \operatorname{dom} playlist \mid i \in \operatorname{ran} playlist(p)\} \end{array}$$

We then have:

$$\begin{array}{|l} length : SongId \twoheadrightarrow Npopularity : SongId \twoheadrightarrow N \\ \hline \operatorname{dom} lengthsubseteqsongs \operatorname{dom} popularitysubseteqsongs \forall\, i\colon songs \bullet popularity(i) = loveHateScore(i) + p \end{array}$$

(b)

$mostPopular : SongId$

$(\exists_1\ i : songs\ |\forall\, j\colon songs \bullet i \neq j \wedge popularity(i) > popularity(j)) \Rightarrow$

mostPopular = (mu $i : songs\ |\forall\, j\colon songs \bullet i \neq j \wedge popularity(i) > popularity(j))$

and

$\neg\ \exists_1\, i\colon songs \bullet \forall\, j\colon songs \bullet i \neq j \wedge popularity(i) > popularity(j) \Rightarrow mostPopular = nullSong$

(c)

playlistsContainingMostPopularSong == $\{i\colon \operatorname{dom} \mathit{playlists} \mid \mathit{mostPopular} \in \operatorname{ran} \mathit{playlists}(i)\}$

Solution 52

(a)

$\mathit{premiumPlays} : \mathit{seq}(\mathit{Play}) \rightarrow \mathit{seq}(\mathit{Play})$

premiumPlays($\langle\rangle$) = $\langle\rangle$

$\forall\ x : \mathit{Play};\ s : \mathit{seq}(\mathit{Play}) \mid$

premiumPlays($\langle x \rangle\ ^s$) =$\langle x \rangle\ ^p \mathit{remiumPlays}(s) \mathit{ifuserStatus}(x.2) = \mathit{premium}$

premiumPlays(s) if userStatus(x.2) = standard

$(\mathit{Note} : \mathit{UsescamelCaseforfuzzcompatibility})$

(b)

$\mathit{standardPlays} : \mathit{seq}(\mathit{Play}) \rightarrow \mathit{seq}(\mathit{Play})$

standardPlays($\langle\rangle$) = $\langle\rangle$

$\forall\ x : \mathit{Play};\ s : \mathit{seq}(\mathit{Play}) \mid$

standardPlays($\langle x \rangle\ ^s$) =$\langle x \rangle\ ^s \mathit{tandardPlays}(s) \mathit{ifuserStatus}(x.2) = \mathit{standard}$

standardPlays(s) if userStatus(x.2) = premium

$(Note: UsescamelCaseforfuzzcompatibility)$

(c)

$cumulativeLength : seq(Play) \rightarrow N$

cumulativeLength$(\langle\rangle) = 0$

$\forall\ x : Play;\ s : seq(Play) \mid$

cumulativeLength$(\langle x \rangle\ ^{s}) = length(x.1) + cumulativeLength(s)$

$(Note: UsescamelCaseforfuzzcompatibility)$