Concurrency & Parallelism

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1

1.1 A

```
x = x + 1;
                    x = y + 1;
                                         z = x + y;
y = y + 1;
                    y = 2;
                                            Listing 3: Thread 3
  Listing 1: Thread 1
                    Listing 2: Thread 2
x: 1, 2, 3, 4
y: 1, 2, 3
z: 0, 1, 2, 3, 4, 5, 6, 7
1.2 B
x = x + 1;
                    y = y + 1;
                                         y = x + 1;
  Listing 4: Thread 1
                      Listing 5: Thread 2
                                            Listing 6: Thread 3
x: 1, 5, 6
y: 1, 2, 6, 7
(x, y): (1, 1), (1, 2),
         (5, 1), (5, 6),
         (6, 1), (6, 7)
```

2

False The volatile keyword does not stop multiple threads from modifying the variable. The expression x++ can be broken down into a read-modify-write sequence, which can be interleaved between threads.

3

- 1. Safety
- 2. Liveness
- 3. Liveness
- 4. Safety
- 5. Safety

4

- 1. G
- 2. A
- 3. D
- 4. B
- 5. C
- 6. F
- 7. H
- 8. E

5

Because the barrier will ensure that all threads wait for each other.

6

6.1 a

The number of combinations is 3! = 6 and the possible combinations are:

- 1 3 5
- 1 5 3
- 3 1 5
- 3 5 1
- 5 1 3
- 5 3 1

```
6.2 b
1 3 5 2 4 6
6.3
    \mathbf{c}
3 6 2 5 4 1
7
public class BarrierN {
    private int a, b;
    public BarrierN (int howmany) {
        b = howmany;
        a = 0;
    public synchronized void arrive() {
        a++;
        if (a < b) {
             this.wait();
             return;
        this.notifyAll();
    }
}
8
```

8.1 a

In the line 11, the program read the values [20, 25, 30], before M_2 runs the credit statement, M_1 transferTo is able to execute completely and then M_2 credit statement runs, adding an old interest value.

8.2 b

The way it is written the program cannot deadlock, unless the client can have accounts that are references to other accounts.