

# Concurrency & Parallelism

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## 1

### 1.1 A

<code>x = x + 1;</code>	<code>x = y + 1;</code>	<code>z = x + y;</code>
<code>y = y + 1;</code>	<code>y = 2;</code>	Listing 3: Thread 3
Listing 1: Thread 1	Listing 2: Thread 2	

x: 1, 2, 3, 4  
y: 1, 2, 3  
z: 0, 1, 2, 3, 4, 5, 6, 7

### 1.2 B

<code>x = x + 1;</code>	<code>y = y + 1;</code>	<code>x = 5;</code>
Listing 4: Thread 1	Listing 5: Thread 2	<code>y = x + 1;</code>
		Listing 6: Thread 3

x: 1, 5, 6  
y: 1, 2, 6, 7  
(x, y): (1, 1), (1, 2),  
          (5, 1), (5, 6),  
          (6, 1), (6, 7)

## 2

**False** The `volatile` keyword does not stop multiple threads from modifying the variable. The expression `x++` can be broken down into a read-modify-write sequence, which can be interleaved between threads.

### 3

1. Safety
2. Liveness
3. Liveness
4. Safety
5. Safety

### 4

1. G
2. A
3. D
4. B
5. C
6. F
7. H
8. C

### 5

Because the barrier will ensure that all threads wait for each other.

### 6

#### 6.1 a

The number of combinations is  $3! = 6$  and the possible combinations are:

- 1 3 5
- 1 5 3
- 3 1 5
- 3 5 1
- 5 1 3
- 5 3 1

## 6.2 b

1 3 5 2 4 6

## 6.3 c

3 6 2 5 4 1

## 7

```
public class BarrierN {
    private int a, b;

    public BarrierN (int howmany) {
        b = howmany;
        a = 0;
    }

    public synchronized void arrive() {
        a++;
        if (a < b) {
            this.wait();
            return;
        }
        this.notifyAll();
    }
}
```

## 8

### 8.1 a

In the line 11, the program read the values [20,25,30], before  $M_2$  runs the **credit** statement,  $M_1$  **transferTo** is able to execute completely and then  $M_2$  **credit** statement runs, adding an old interest value.

### 8.2 b

The way it is written the program cannot deadlock, unless the client can have accounts that are references to other accounts.