

## Lecture 06: Active Contours & Intelligent Scissors

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### Active Contours

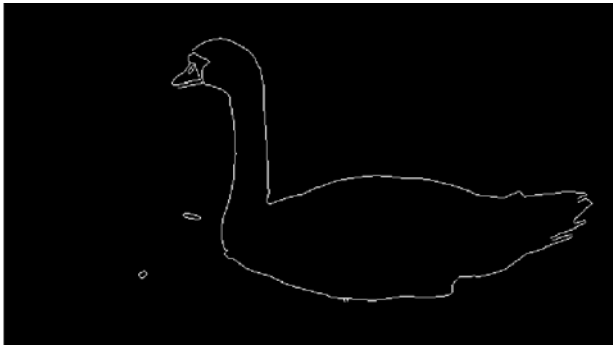
- Raises level of image feature description from edges to boundaries.
- Edge is strong change in pixel intensity.
- Boundary is boundary of an *object*.
  - Smooth (more or less)
  - Closed

Yin,  
Spring

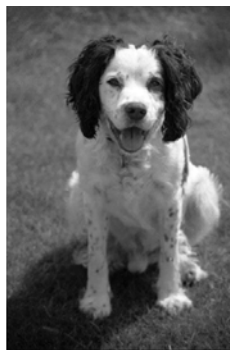


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Sometimes  
edge detectors  
find the  
boundary pretty  
well.



Yin, MST  
Spring, 2014



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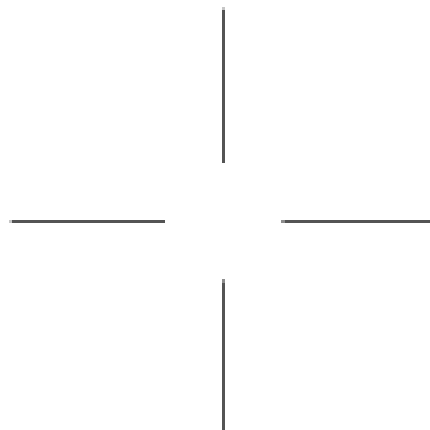
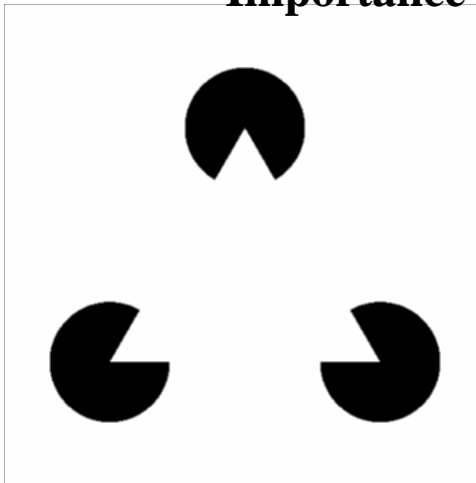
Sometimes  
they don't



## Improved Boundary Detection

- Integrate information over distance.
- Use shape cues
  - Smoothness
  - Closure
- Get User to Help.

## Importance of Continuity

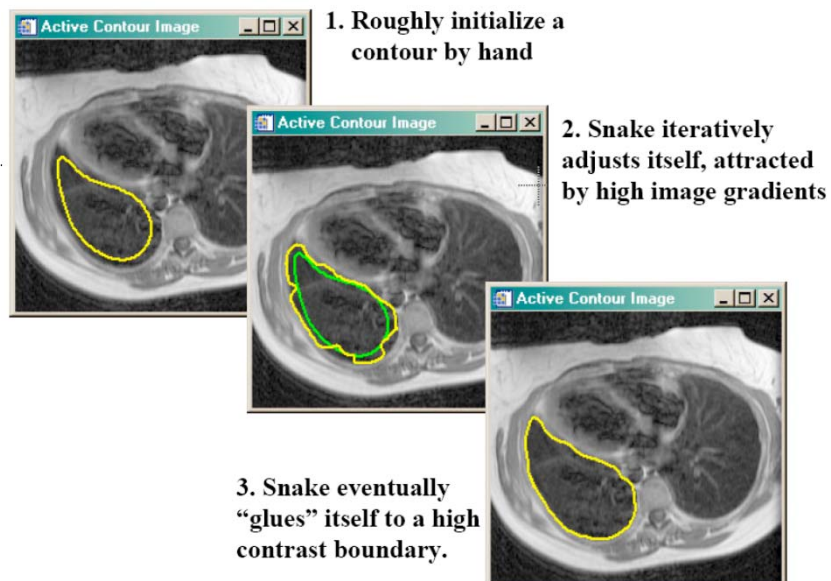


Human perception is good at integrating contour continuity information into boundary detection process

## Active Contours

- They are also called
  - Snakes
  - Deformable Contours
- Think of a snake as an elastic band:
  - of arbitrary shape
  - sensitive to image gradient
  - that can “wiggle” in the image
  - represented as a sequential list of points

## Main Idea:



## The Energy Functional

- Associate to each possible shape and location of the snake a value  $E$ .
  - Values should be s.t. the image contour to be detected has the minimum value.
  - $E$  is called the **energy** of the snake.
- Iteratively adjust points on the snake to achieve a smaller energy  $E$

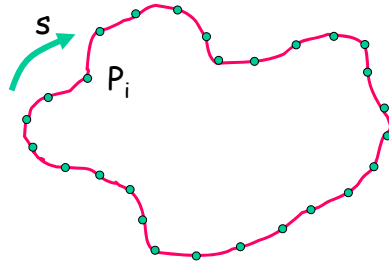
## Energy Functional Design

- We need a function that given a snake state, associates to it an Energy value.
- The function should be designed so that the snake moves towards the contour that we are seeking!

## Snake “state”: Contour Parametrization

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Consider a contour parametrization  $c=c(s)$   
where  $s$  is the “arc length”



Each point  $P_i$  on the contour  
has coordinates  $(x_i(s), y_i(s))$

## What moves the snake?

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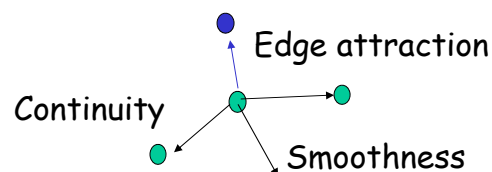
Forces applied to its points

External Forces:

- boundaries in the image (gradients)
- term associated with the DATA

Internal Forces:

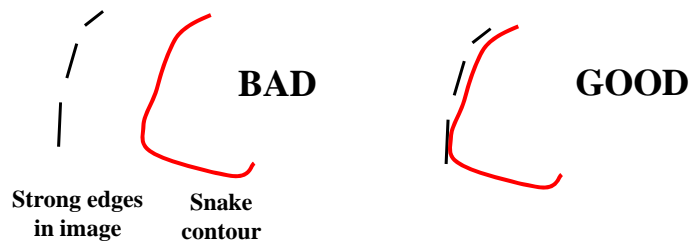
- continuity and curvature
- terms associated with the CURVE



## Forces moving the snake (External)

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- It needs to be attracted to contours:
  - Edge pixels “pull” the snake points.
  - The stronger the edge, the stronger the pull.
  - The force is proportional to  $|\nabla I|$



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## Edgeness Term

Given a snake with  $N$  points  $p_1, p_2, \dots, p_N$

Define the edgeness term of the Energy Functional:

$$E_g(p_i) = -\|\nabla I(p_i)\|$$

Magnitude of the gradient should be **LARGE**  
(which will make this term **SMALL** (very negative))

## Forces preserving the snake (Internal)

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- The snake should not break apart!
  - Points on the snake must stay close to each other
  - The farther the neighbor, the stronger the force to pull them back together
  - The force is proportional to the distance  $|P_i - P_{i-1}|$



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## Continuity Term

Given a snake with  $N$  points  $p_1, p_2, \dots, p_N$

Let  $d$  be the average distance between points

Define the continuity term of the Energy Functional:

$$E_c(p_i) = (d - \underbrace{\|p_i - p_{i-1}\|})^2$$

Distance between points should be kept close to average



## Forces preserving the snake (Internal)

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The snake contour should be “smooth”

- Penalize high curvature.
- Force proportional to snake curvature



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## Smoothness Term

Given a snake with  $N$  points  $p_1, p_2, \dots, p_N$   
Curvature should be kept small

Define the smoothness term of the Energy Functional:

$$E_s(p_i) = \underbrace{\|p_{i-1} - 2p_i + p_{i+1}\|^2}_{\text{Second derivative measures curvature}}$$

## Snake Energy Functional

Given a snake with  $N$  points  $p_1, p_2, \dots, p_N$

Define the following Energy Functional:

$$E = \sum_{i=1}^N a_i E_c(p_i) + b_i E_s(p_i) + c_i E_g(p_i)$$

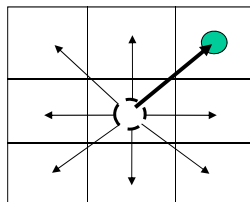
Where:

$E_c$	"Continuity"
$E_s$	"Smoothness"
$E_g$	"Edgeness"

$a_i, b_i, c_i$  are "weights" to control influence

## Greedy Algorithm

Minimize energy one point at a time.  
For each point, consider a finite set of moves  
in a small window around it



Compute the new energy for each candidate location  
Move the point to the one with the minimum value

## Implementation Considerations

- To avoid numerical problems, the terms of the energy function should be normalized.
  - $E_c$  and  $E_s$  are normalized by their maximum in the neighborhood
  - $E_g$  is normalized as  $|\nabla I - m| / (M - m)$ 
    - $M$  and  $m$  are the *max* and *min* value of the gradient magnitude in the neighborhood

**That is, want all terms scaled from 0 to 1 so they are treated equally**

## Implementation Considerations

Keeping high-curvature corners

- Before starting a new iteration:
  - Search for “corners”:
    - max curvature
    - large gradient
  - Corner points should not contribute to the energy (set  $b_i = 0$ )

## Snake Algorithm

- Input:
  - gray scale image  $I$
  - a chain of points  $p_1, p_2, \dots, p_N$
- $f$  is the fraction of points that must move to start a new iteration
- $U(p)$  is a neighborhood around  $p$
- $d$  is the average distance between snake points (computed from the list of points).

## Snake Algorithm

While the fraction of moved points  $> f$

1. For  $i=1, 2, \dots, N$ 
  1. find a point in  $U(p_i)$  s.t. the energy  $E$  is minimum,
  2. move  $p_i$  to this location
2. For  $i=1, 2, \dots, N$ 
  1. Estimate the curvature  $k=|p_{i-1}-2p_i+p_{i+1}|$
  2. Look for local max, and set  $b_{\max} = 0$
3. Update  $d$

Where

$$E = \sum_{i=1}^N a_i E_c(p_i) + b_i E_s(p_i) + c_i E_g(p_i)$$

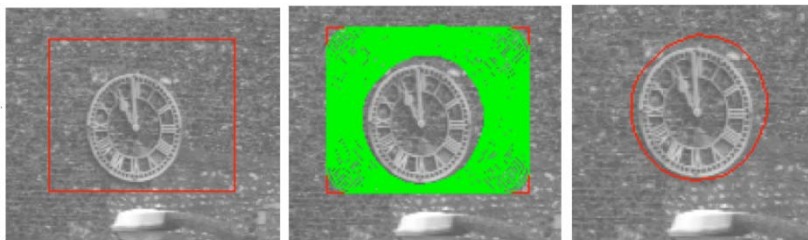
## Issue with this Algorithm

**This algorithm is not guaranteed to find the “best” curve, in the sense of lowest cost.**

**Why? Greedy algorithms do not explore the space of all curves**

## Non-Optimality

**Typical snake behavior is far from optimal**



Snake movement gets “hung up”  
on high contrast stone in wall.

**There have been a \*lot\* of papers written about how to make snakes work robustly: different energy functions, different optimization methods...**

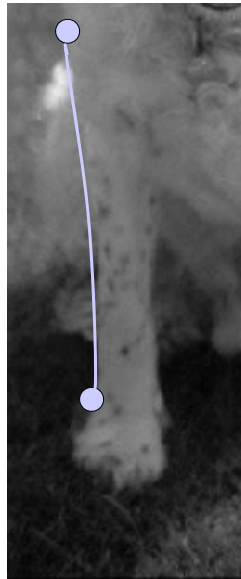
## A More Optimal Strategy

- Given a start and end point, use dynamic programming to determine “best” path from start to end location.
  - Need to determine what is a good path?
  - Need procedure to find best path

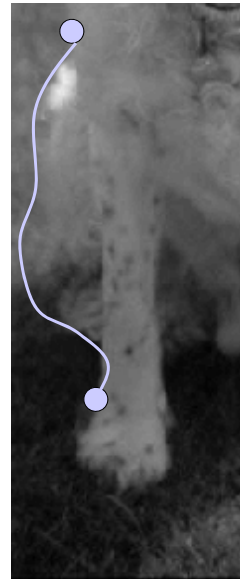
Algorithm we will discuss now is from  
E. N. Mortensen and W. A. Barrett,  
“Intelligent Scissors for Image Composition,”  
in ACM Computer Graphics (SIGGRAPH `95),  
pp. 191-198, 1995



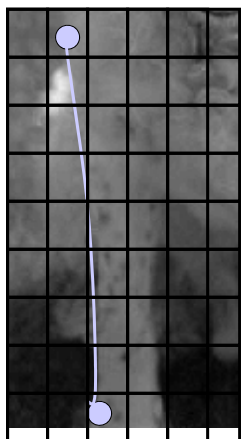
We'll do something easier than finding the whole boundary. Finding the best path between two boundary points.



How do we  
decide how  
good a path is?  
Which of these  
two paths is  
better?

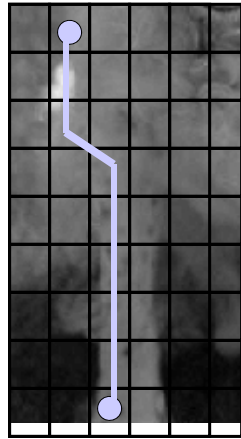


## Discrete Grid



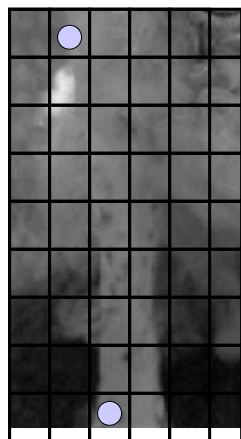
- Contour should be near edge.
  - Strength of gradient.
- Contour should be smooth (good continuation).
  - Low curvature
  - Low change of direction of gradient.

## To start: contour near edge



- For each step from one pixel to another, we measure edge strength (change in intensity across edge).
- Find path with biggest total edge strength.

## So How do we find the best Path?



### Dynamic programming

A Curve is a path through the grid.

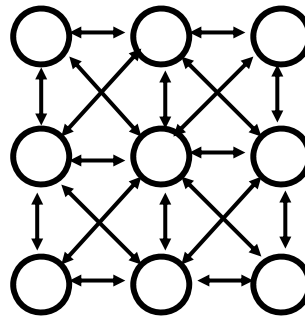
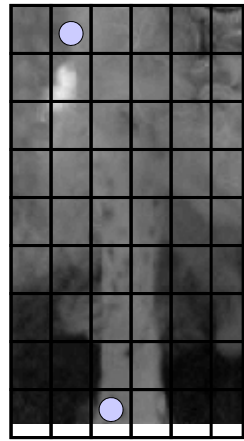
Cost depends on each step of the path.

We want to minimize cost.

Incrementally determine best path, starting from end state



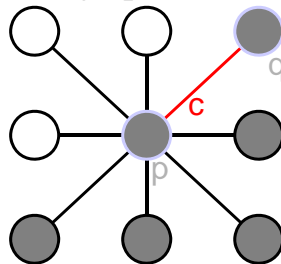
## Map problem to Graph



Weight represents cost of going from one pixel to another. Next term in sum.

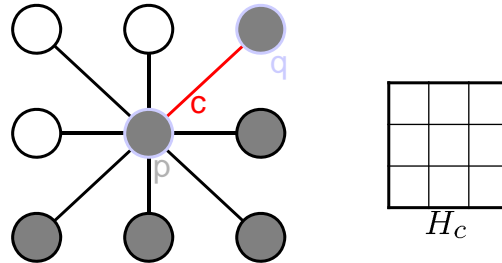
## Defining the costs

- Treat the image as a graph



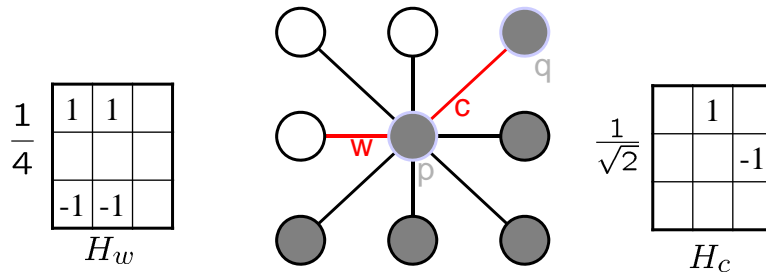
- Want to hug image edges: how to define cost of a link?
  - the link should follow the intensity edge
    - want intensity to change rapidly orthogonal to the link
  - $c \approx -|\text{difference of intensity orthogonal to link}|$

## Defining the costs



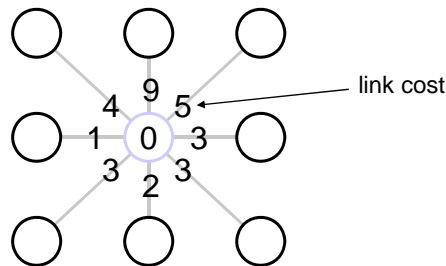
- First, smooth the image to reduce noise.
- $c$  can be computed using a cross-correlation filter
  - assume it is centered at  $p$
- Also typically scale  $c$  by it's length
  - set  $c = (\max\text{-}|\text{filter response}|) * \text{length}(c)$ 
    - where  $\max = \text{maximum } |\text{filter response}| \text{ over all pixels in the image}$

## Defining the costs



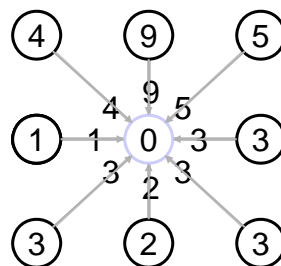
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## Dijkstra's shortest path algorithm



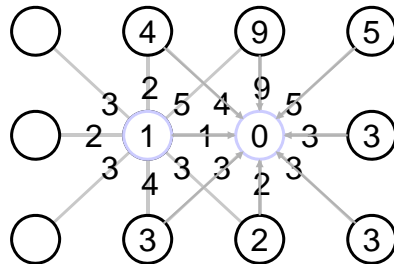
- Algorithm
  1. init node costs to  $\infty$ , set  $p$  = seed point,  $\text{cost}(p) = 0$
  2. expand  $p$  as follows:
    - for each of  $p$ 's neighbors  $q$  that are not expanded
      - set  $\text{cost}(q) = \min(\text{cost}(p) + c_{pq}, \text{cost}(q))$

## Dijkstra's shortest path algorithm



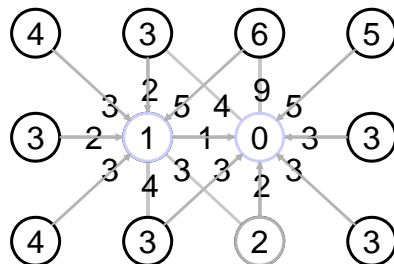
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        - » if  $q$ 's cost changed, make  $q$  point back to  $p$
      - put  $q$  on the ACTIVE list (if not already there)

## Dijkstra's shortest path algorithm



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      - put  $q$  on the ACTIVE list (if not already there)
  3. set  $r$  = node with minimum cost on the ACTIVE list
  4. repeat Step 2 for  $p = r$ .

## Dijkstra's shortest path algorithm



- Algorithm
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5. Stop when next point to expand is goal point. Read off shortest path.

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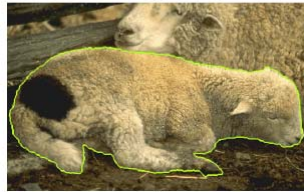
## Intelligent Scissors and Image Composition



(a)



(b)



(c)



## Lessons Learned

- Perceptual organization: contour continuity constraints needed for boundary detection.
- Fully automatic methods for boundary finding (snakes, active contours) are not yet good enough
- Formulate desired solution as a cost function, then optimize it
  - greedy methods, easy but suboptimal
  - dynamic programming --> optimal solution
  - when the method is applicable