#### Scheduling and Analys of Limited-Preemptive Modable Gang Tasks

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- Parallel threads executed together as a "gang"
- Execution does not start until there are enough free cores



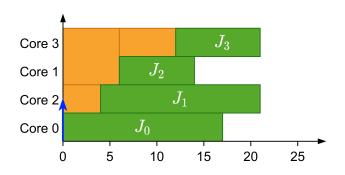
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First of all, let's explain what is gang scheduling.

It's the execution of multiple parallel threads together as a "gang". In these threads the execution does not start until there are enough cores to execute them together.

- 1.We have these jobs assigned to different cores.
- 2.If we release them as a gang job then we should pack them like a single task. Obtaining the following group. <click>
- 3.Then,<click> we have obtained the gang task result of merging the jobs

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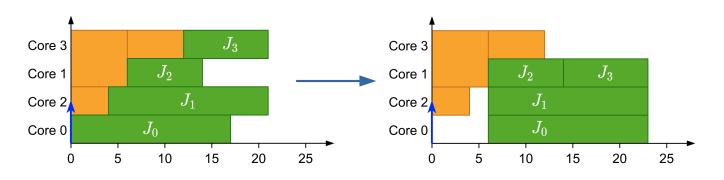
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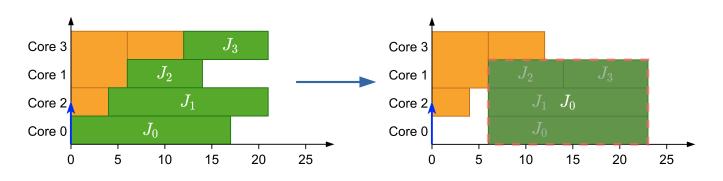
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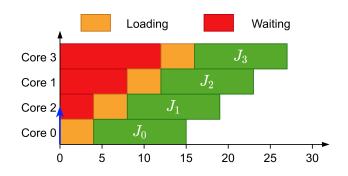
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- Avoids overhead when loading initial data



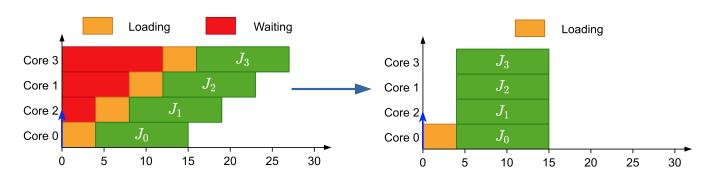
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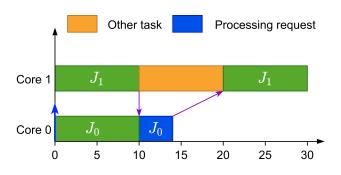


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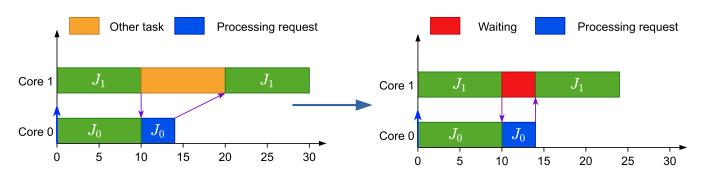
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#### Types of gang

- Rigid: number of cores set by programmer
- Moldable: number of cores assigned during scheduling
- Malleable: number of cores can change during runtime



12

We can have three types of gang scheduling:

- Rigid gang, where the number of cores is set by the programmer and fixed during execution
- Moldable gang, where the number of cores can be in a range of values and the actual chosen value is decided when scheduling the task
- Malleable gang, where the number of cores can vary during the execution of the job

#### Previous work



# Rigid gang solutions



# Moldable gang solutions



## Malleable gang solutions



# Schedulability analysis



#### Our work



1:

#### **LPMRGS**

