# Program 4 Design and Analysis

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12/9/2014

## 1 Part 1 - Shared Memory

## 1.1 Program Description

## 1.2 Design

Following Foster's design methodology, we have a brief overview:

## 1.2.1 Partitioning

- Reading values to be hashed
- Hashing
- Sending
- Recieving
- Inserting

### 1.2.2 Communication

- Sharing data between computations
- Determine amount and pattern of communication

### 1.2.3 Aglomeration

• Combine or group tasks to improve performance

## 1.2.4 Mapping

• Assign agglomerated tasks to threads/physical processors.

## 1.3 Performance Testing

# 2 Part 2 - Distributed memory

## 2.1 Program Description

This distributed memory program will allow for the typical hash table operations, but on a set of programs on potentially different computers.

Each process on each machine will be responsible for its own hash table, and respond to messages from a managing producer thread. The managing thread will give values to store and request table lookups.

The data is currently random strings read in from a file. A few sample inputs are provided.

Figure 1: Timing Evaluation of increasing problem size

Test 1. Increasing Number of Values Hashed

Thread Count = 2, Chunk Size = 5

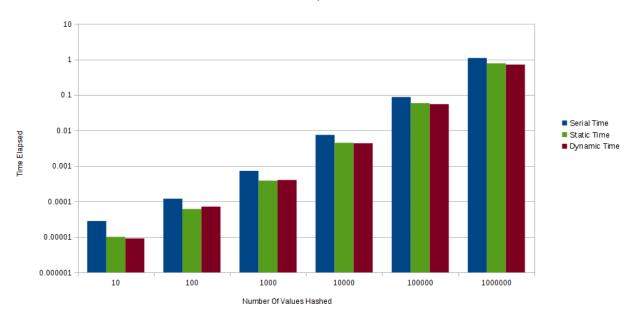


Figure 2: Speedup / Efficiency Evaluation of increasing problem size

Test 1. Increasing Number of Values Hashed

Speedup / Efficiency Analysis

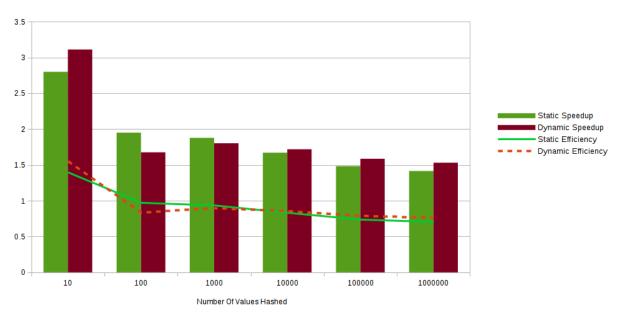


Figure 3: Timing Evaluation of increasing thread count

Test 2. Increasing Number of Threads Used

Values Hashed = 1000, Chunk Size = 1

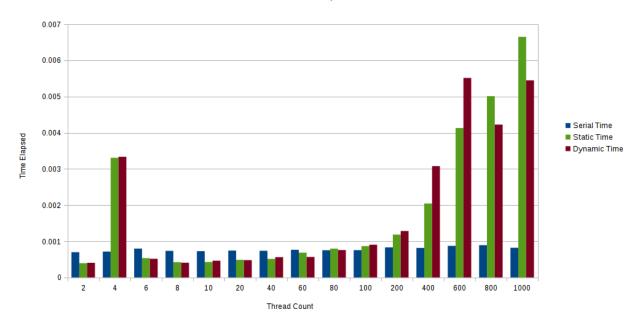


Figure 4: Speedup / Efficiency Evaluation of increasing thread count

Test 2. Increasing Number of Threads Used

Speedup / Efficiency Analysis

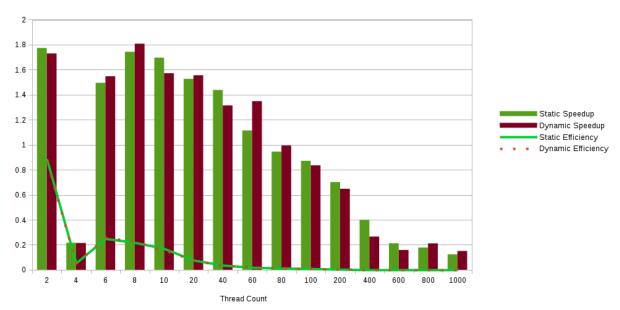


Figure 5: Timing Evaluation of increasing chunk size

Test 3. Increasing Chunk Size

Values Hashed = 8000, Thread Count = 2

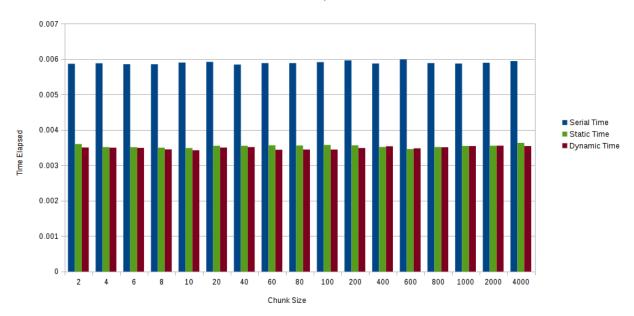
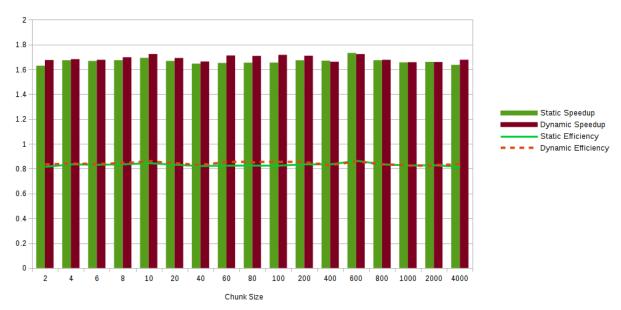


Figure 6: Speedup / Efficiency Evaluation of increasing chunk size

Test 3. Increasing Chunk Size

Speedup / Efficiency Analysis



## 2.2 Design

Following Foster's design methodology, we have a brief overview:

#### 2.2.1 Partitioning

For our distributed memory solution, as in our shared, we will be doing data-centric partitioning. We will assign to each piece of data (in this case the strings) the tasks of hashing that data and storing it into a table.

This will create a large number of tasks, much larger than our number of processors. There are no redundant computations, and each tak will be roughly the same size.

#### 2.2.2 Communication

For each piece of data, in order to place it in a table, it must first be hashed according to the hash function. This is an example of local communication. Therfore, since the tasks are dependent, we can consolidate them into one single task, hashing and storing a value.

In addition, the tasks associated with a piece of data will work very independently. Very little global communication will be needed.

Communication is balanced, minimized, and all tasks can be performed concurrently.

## 2.2.3 Aglomeration

As discussed in the last section, each piece of data can be consolidated into one combined operation.

Other than that, there is not much agglomeration that can be done. Because the work associated with each piece of data is independent of the solution as a whole, and the locality of the table insert is minimal, we cannot group together tasks in a meaningful or advantageous way.

#### 2.2.4 Mapping

For the distributed memory solution, we will employ a producer-consumer structure. We will have one root thread that produces values from an input file, handing out the values to the consumer threads. Each consumer thread will take the piece of data, hash it, and place it in the table.

It is very important that we balance the load that each consumer thread experiences, making sure that they are not overwhelemed with data that will make computation take longer or potentially fill up their table. At the same time, we want to make sure that the storing of data is determinstic, that way we can find the data in a later loopup operation. In order to do this, the data is hashed twice. Once by the root producer thread in order to determine which worker thread will get the data, and again by the worker to determine where in its table to place the data.

## 2.3 Performance Testing