

JONATHAN PARLETT

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- A simplified encoding scheme from polynomials with complex coefficients to vectors with complex coefficients
- The actual encoding scheme from polynomials with integer coefficients to vectors of complex coefficients.

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- Another important property of cyclotomics is that there roots are complex conjugates of each other. To see this lets look at the 8-th cyclotomic $X^4 + 1$

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- Another important property of cyclotomics is that there roots are complex conjugates of each other. To see this lets look at the 8-th cyclotomic $X^4 + 1$
- $\Phi_8(x) = (x - e^{2i\pi \frac{1}{8}})(x - e^{2i\pi \frac{3}{8}})(x - e^{2i\pi \frac{5}{8}})(x - e^{2i\pi \frac{7}{8}})$

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- Grouping real and imaginary parts we can see that the 1st root is the complex of the 4th and the 2nd root is the complex conjugate of the 3rd

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- This fact will become important when we are discussing the mapping of polynomials with integer coefficients to complex vectors

SIMPLIFIED ENCODING SCHEME

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- So the first big thing to understand in order to fully understand CKKS is this encoding scheme. How do we get from our complex vectors to our plaintext polynomials
- The ultimate goal is to be able to fully understand the map that defines CKKS encoding algorithm

$$\sigma^{-1} : \mathbb{C}^N \rightarrow \frac{\mathbb{Z}[X]}{X^N + 1}$$

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- The forward map may be easier to consider

$$\sigma : \frac{\mathbb{C}[X]}{X^N + 1} \rightarrow \mathbb{C}^N$$

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- In CKKS they usually consider N to be a power of 2, $N = 2^k$. So we need cyclotomic polynomial of degree $N = 2^k$.
- Since as we noted earlier the degree of the n -th cyclotomic is equal to $\phi(n)$ it is easy to find a cyclotomic of the appropriate degree as $\phi(2^{k+1}) = 2^k$, since the only numbers that have a common factor with 2^k will be only the even numbers less than 2^k .

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- Here we can see we have the coefficient vector **a** that uniquely determines the polynomial and the output vector **b**.

SIMPLIFIED ENCODING SCHEME



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- Its clear from this equation that given one we can compute the other and since this is a square matrix there is one and only one solution

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- Its clear from this equation that given one we can compute the other and since this is a square matrix there is one and only one solution
- so it should be intuitive that this transformation defines an isomorphism between \mathbb{C}^4 and $\frac{\mathbb{C}[X]}{X^4+1}$

SIMPLIFIED ENCODING SCHEME

- At this point we have defined our simplified map and its inverse as essentially the equation we showed in the previous slide

$$\sigma^{-1} : \mathbb{C}^N \rightarrow \frac{\mathbb{C}[X]}{X^N + 1}$$

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- The CKKS map/algorithm $\sigma^{-1} : \mathbb{C}^N \rightarrow \frac{\mathbb{Z}[X]}{X^N + 1}$ adds further structure in order to place restrictions on this map to ensure that we encode our complex vectors as polynomials with integer coefficients only