

# Writing an Interpreter in Go

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# Paul Graham - YCombinator

Your mind is like a compiled program you've lost the source of.  
It works, but you don't know why.

# What is a Compiler?

- program that transforms a formal language into another formal language
- immutable - same input language always yields same output language
- output language - called the *target* - executes "fastest"

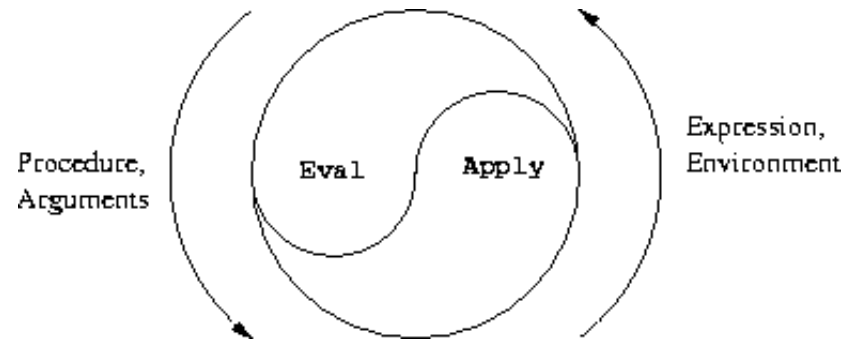
Does a compiler define semantics of either source or target language?

# Examples of Typically Compiled Languages

- go
- c
- java
- perl5 is **NOT** compiled
- **perl6** will eventually be compiled
- YACC compiles Backus-Naur into Go (and many other languages)
- Ken Thompson (gofather) compiled regular expressions into pdp11 assembler
- compiling SQL into native code (Oracle, PostgreSQL)

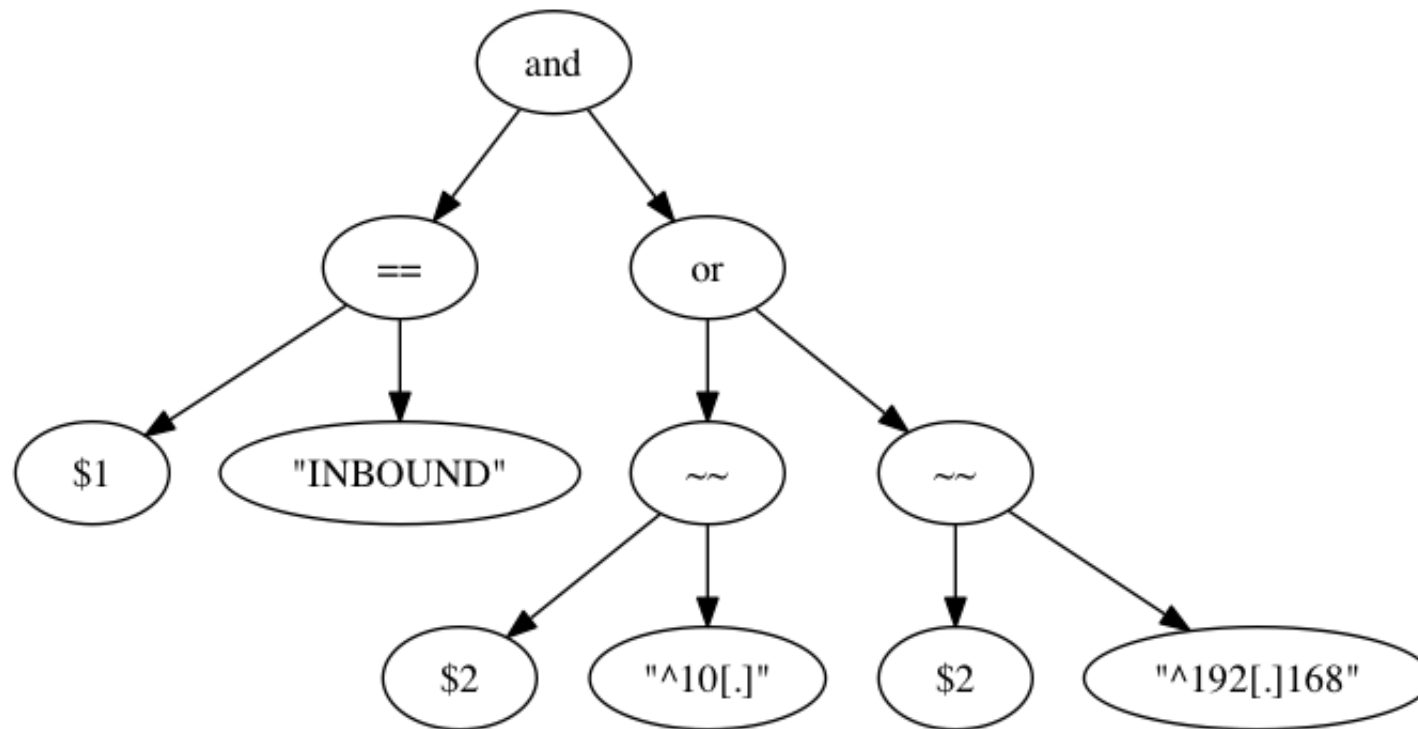
# What is an Interpreter?

- program that directly executes the source program (often text)



- bourne shell, awk, javascript, ruby, SQL
- lisp was the first interpreter
- the java command is an interpreter (strictly) of compiled java class files

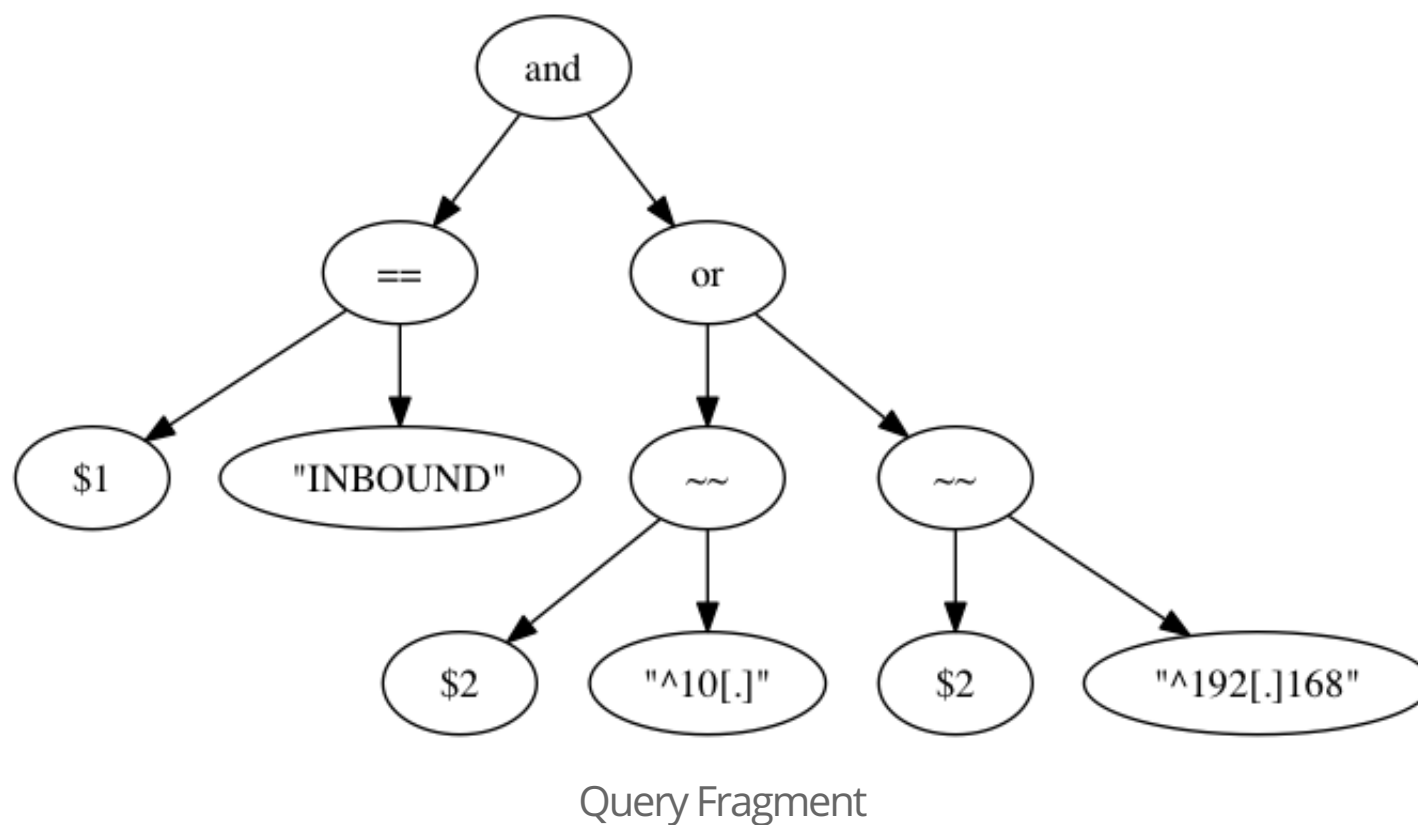
# Interpreter Translates Source into Internal Format



Abstract Syntax Tree (perl5, awk)

- line by line (many unix shells)
- virtual machine instruction (python, ruby, vbasic, , perl6)
- java translates jvm opcodes into native hardware instructions

# Abstract Syntax Tree of a Query Fragment in HOQ



```
$1 == "INBOUND" and ( $2 ~~ "^10[.]" or $2 ~~ "^192[.]168" )
```

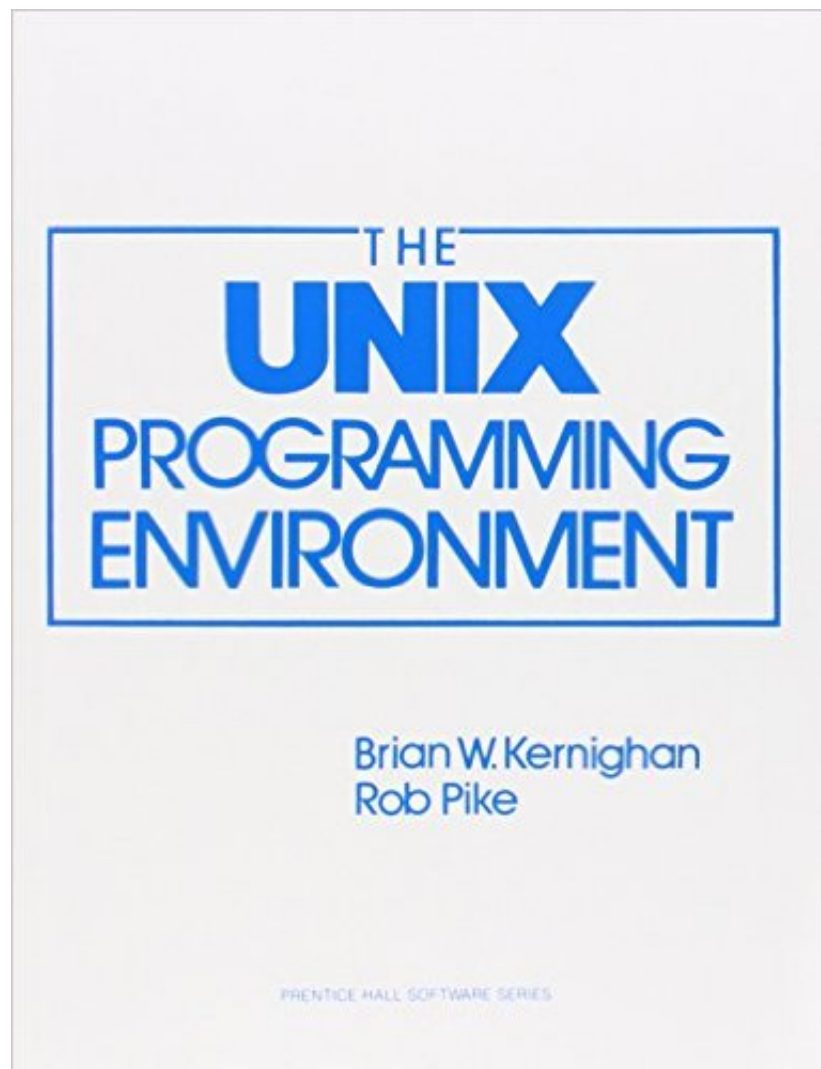
# What is HOQ?

- Higher Order Query
- toy interpreter that demonstrates the famous YACC compiler for Go language
- also, hopefully, demonstrates "Communicating Sequential Processes"

Did I bite off more than I can chew?



# Inspired by "HOC" Calculator in Unix Programming Environment



Higher Order Calculator - Chapter 8

Mighty Fine Book

## How Do We Invoke HOQ?

- command line program invoked with script as first argument

```
hoq file.hoq < data.in
```

- hoq script *may* invoke other unix programs and wait for their exit status

# HOQ Execution is Driven by Standard Input

- consuming text drives the execution of the whole hoq script

```
tr ', ' '\t' quarterly.csv | hoq quarterly.hoq
```

Terminal

- hoq terminates after reading the the final line and waiting for all processes to exit.

# HOQ Splits Input Into Fields

- input lines `strings.Split()` on tab separated boundaries

```
line = strings.TrimRight(line, "\n")
...
fields:  strings.SplitN(line, "\t", 255),
```

- \$1 is first tab separated field, \$2 is second field, ...
- \$0 is whole, current line, including tabs, minus terminating new line

# HOQ is Mostly Declarative

- qualify on patterns in tab separated field (\$1, \$2)
- qualify on process exit status codes (uint8)
- boolean combinations (logical and, or, not) on any qualifications
- subprocess are executed when boolean qualifications are true

Execution order of subprocess is directed acyclic graph (DAG)

# Hello, World

```
command say
{
    path = "echo";
}

exec say("hello, world");
```

say.hoq

# Invoke Say

```
command say
{
    path = "echo";
}

exec say("hello, world");
```

say.hoq

```
$ echo | hoq say.hoq
hello, world
```

Terminal

# Good Bye, Cruel World

```
command say1 {  
    path = "echo";  
}  
command say2 {  
    path = "echo";  
}  
  
exec say2("good bye, cruel world");  
exec say1("hello, world");
```

say-bye.hoq

```
$ echo | hoq say-bye.hoq  
hello, world  
good bye, cruel world  
  
$ echo | hoq say-bye.hoq  
good bye, cruel world  
hello, world
```



# Trinity

```
command say1 {  
    path = "echo";  
}  
command say2 {  
    path = "echo";  
}  
command say3 {  
    path = "echo";  
}  
  
exec say1("hello, world");           # always executes  
  
exec say2("to be or not to be")    when say1.exit_status == 0;  
exec say3("good bye, cruel world") when say2.exit_status == 0;
```

say-trinity.hoq

```
$ echo | hoq say-trinity.hoq  
hello, world  
to be or not to be  
good bye, cruel world
```

# A Complex Qualification

```
command blob_on_network {
    path = "true";                # path to /bin/true
}
command merge_blob {
    path = "merge-blob";         # path to executable program
}

exec blob_on_network()
    when ((                        # $3 is a blob request action
        $3 == "put" or $3 == "get" or
        $3 == "eat" or $3 == "wrap" or
        $3 == "roll"
    )
        and $5 == "ok"           # $5 is request status
    ) or (
        $3 == "give" and $5 == "ok,ok"
    );

exec merge_blob($1, $2)          # merge blob request into database
    when
        blob_on_network.exit_status == 0;
```

Edited from <https://github.com/jmscott/setspace/blob/master/schema/setspace/setspace.flow.example>

# Is Wu Wei Nothing?

# Idiom #1 - Turn Any Function into a Channel

```
package main
import "fmt"

func fib(i uint64) uint64 {
    if i < 2 {
        return 1
    }

    return fib(i - 1) + fib(i - 2)
}

func main() {
    fmt.Println(fib(24))
}
```

[Run](#)

Fibonacci is a Sequential Function

# Fibonacci Becomes a Channel

```
func fib_chan(in chan int) (out chan int) {  
  
    out = make(chan int)  
  
    go func() {  
        defer close(out)  
  
        for i := range in {  
            out <- fib(i)  
        }  
    }()  
  
    return out  
}
```

Idiom #1 - Turn Any Function into a Channel

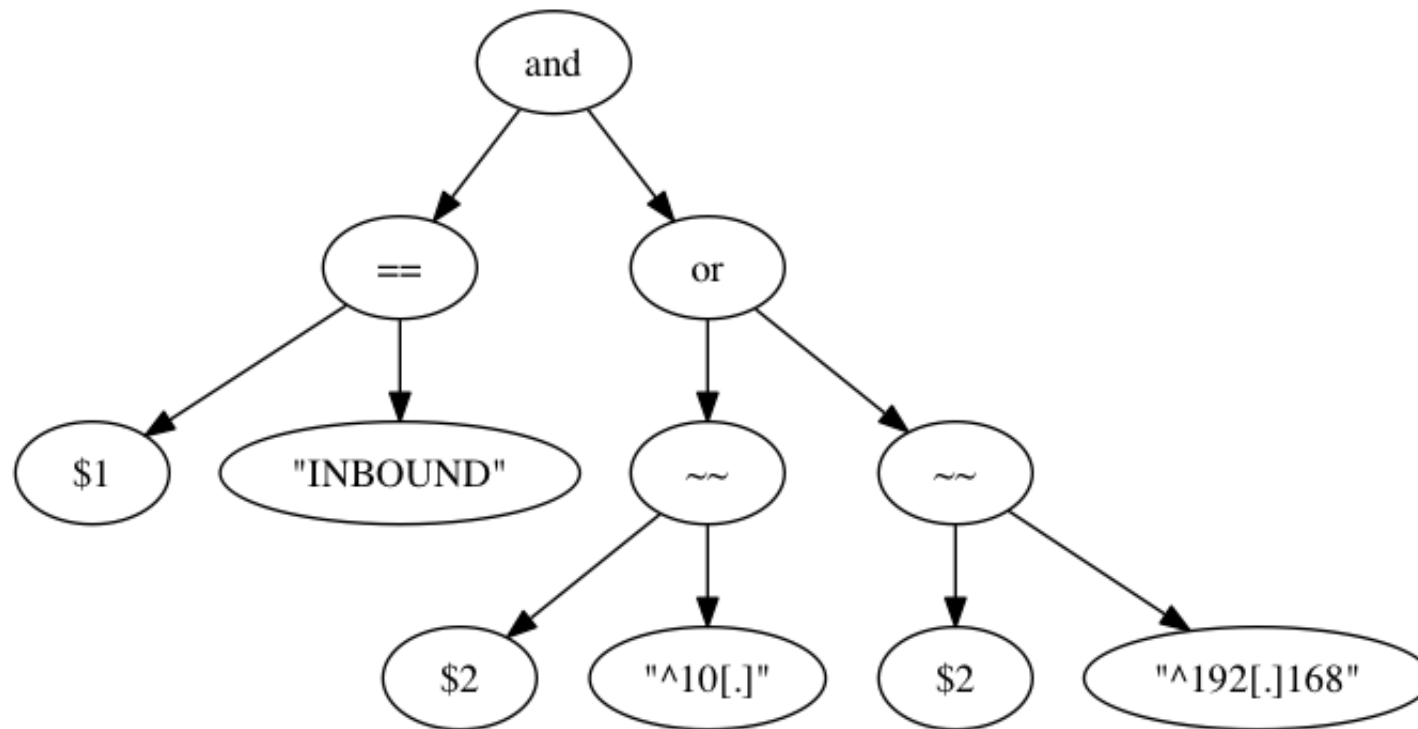
# Idiom #1 - Run the Fibonacci Series

```
func main() {  
    in := make(chan int)  
  
    out := fib_chan(in)  
  
    i := 1;  
    for {  
        in <- i  
        fmt.Println(i, "=", <-out)  
        i++  
    }  
}
```

[Run](#)

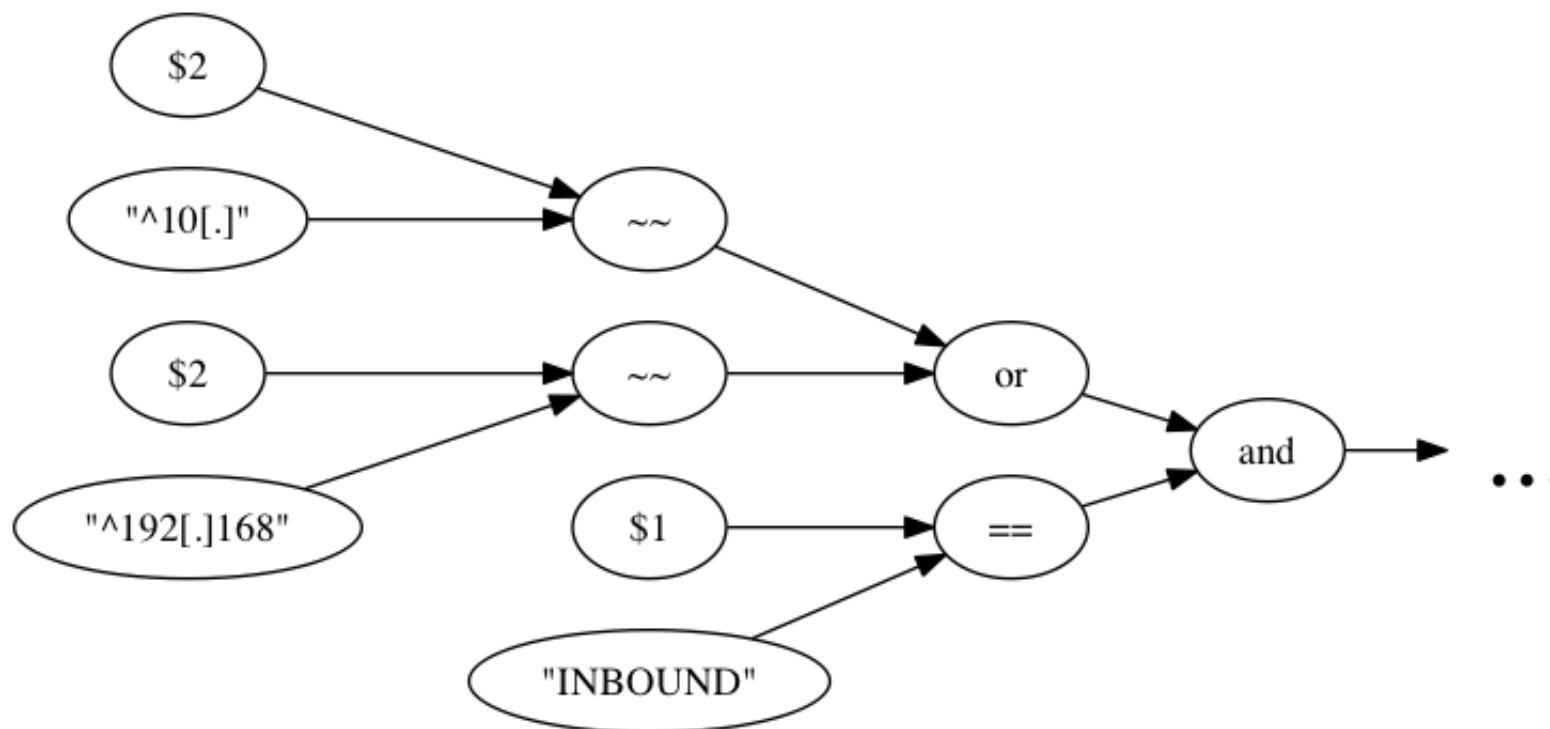
Idiom #1 - Turn Any Function into a Channel

# Abstract Syntax Tree is a Flow Graph



```
$1 == "INBOUND" and ( $2 ~~ "^10[.]" or $2 ~~ "^192[.]168[.]" )
```

# Data Flows From Leaves to Root

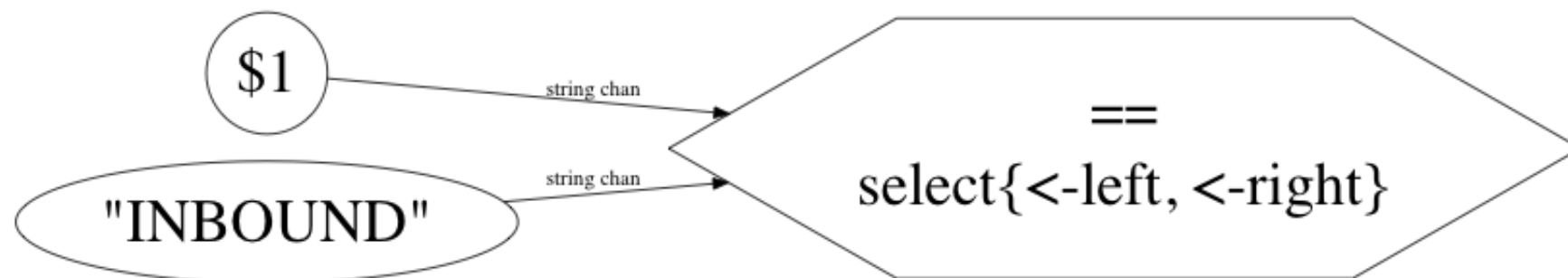


Each Edge is a Go Channel

```
$1 == "INBOUND" and ( $2 ~~ "^10[.]" or $2 ~~ "^192[.]168[.]" )
```

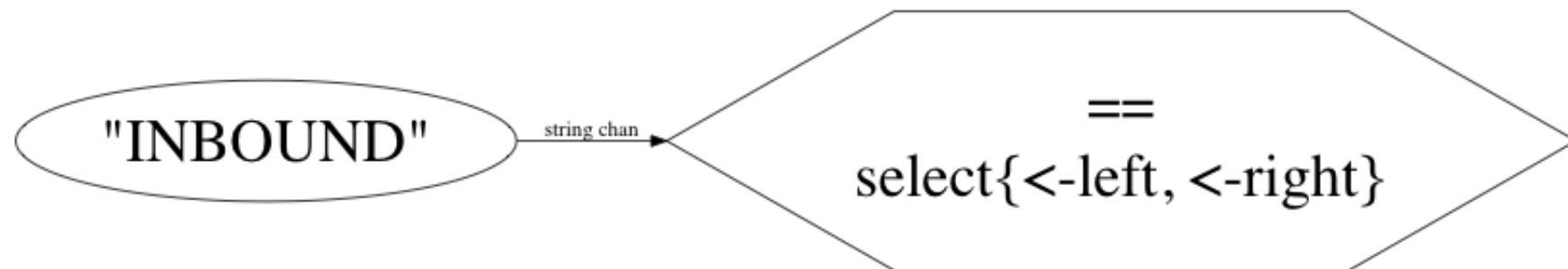


# AST Query Fragment: Data Flows Towards String Equality Operand



```
$1 == "INBOUND"
```

# AST Query Fragment: Send String Constant to Equality Operand



Similar to Push in Sequential Machine

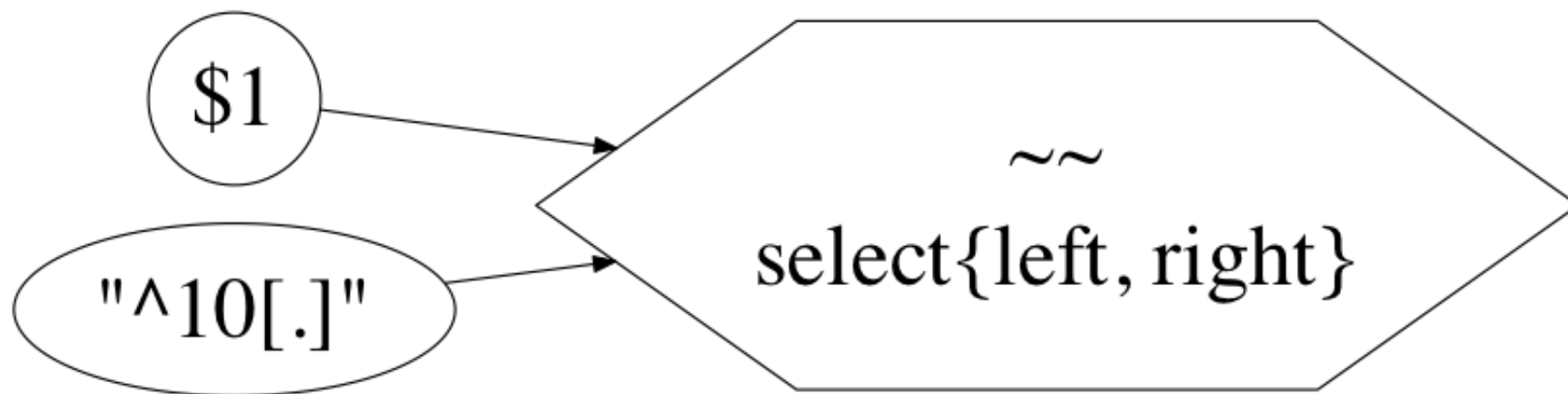
# Opcode: const\_string()

```
type string_value struct {  
    string  
    is_null bool  
}  
type string_chan chan *string_value
```

```
func (flo *flow) const_string(s string) (out string_chan) {  
    out = make(string_chan)  
  
    go func() {  
        defer close(out)  
  
        for flo = flo.get(); flo != nil; flo = flo.get() {  
            out <- &string_value{ // send string upstream to next opcode  
                string: s,  
            }  
        }  
    }()  
  
    return out  
}
```

'Push' a String Constant

## Idiom #2 - `select{}` on Two Channels Allows Concurrent Operands



```
$1 ~~ "^10[.]"
```

# The regex.Match() is a Binary Relation on Two Strings

```
import "regexp"

func re_match(sample, re string) bool {

    matched, err := regexp.MatchString(re, sample)
    if err != nil {
        panic(err)
    }

    return matched
}
```

import "regexp"

```
type bool_value struct {
    bool
    is_null bool
}
type bool_chan chan *bool_value
```

Idiom #2 - select{} on Two Channels Allows Concurrent Operands

## - OpCode: Relational Binary Operator on Two Strings

```
func (flo *flow) string_rel2(  
    rel2 func(left, right string) bool,  
    in_left,  
    in_right string_chan,  
) (out bool_chan) {  
  
    out = make(bool_chan)  
    go func() {  
        defer close(out)
```

Upper Half of string\_rel2()

*... Main Loop in Next Slide ...*



```
}()  
return out  
}
```

Lower Half of string\_rel2()

Idiom #2 - select{} on Two Channels Allows Concurrent Operands

## Relational Binary Operator on Two Strings - Main Loop (Idiom #2)

```
for flo = flo.get(); flo != nil; flo = flo.get() {
    var left, right *string_value

    for left == nil || right == nil {
        select {
            case lv := <-in_left:
                if lv == nil {
                    return
                }
                left = lv
            case rv := <-in_right:
                if rv == nil {
                    return
                }
                right = rv
        }
    }
    bv := &bool_value {is_null: left.is_null || right.is_null}
    if bv.is_null == false {
        bv.bool = rel2(left.string, right.string)
    }
    out <- bv          // send channel to next opcode
}
```

# 16 Opcodes of Flow Machine

- `const_string`, `const_bool`, `const_uint8`
- `to_string_uint8`, `to_string_bool`
- `dollar`, `dollar0`
- `string_rel2`, `uint8_rel2`, `bool_rel2`
- `argv0`, `argv1`, `argv`
- `exec`
- `fanout_uint8`, `fanin_uint8`



# Exec() of Command Sends Exit Status When Qualification is True

```
command xtrue {  
    path = "true";  
}  
command say1 ("hello from say1:") {  
    path = "echo";  
}  
command say2 ("hello from say2:") {  
    path = "echo";  
}
```

**exec xtrue();      # always called**

```
exec say1("xtrue exited 0")  
    when  
        xtrue.exit_status == 0  
;  
exec say2("xtrue exited 1")  
    when  
        xtrue.exit_status == 1;
```

```
$ echo | hoq exec-xtrue.hoq  
hello from say1: xtrue exited 0
```

# Exec() of Command Sends Null When Qualification is False or Null

```
command xtrue {  
    path = "true";  
}  
command say1 ("hello from say1:") {  
    path = "echo";  
}  
command say2 ("hello from say2:") {  
    path = "echo";  
}
```

**# xtrue never called, so exit\_status sends null bool upstream to all qualifications  
exec xtrue() when false;**

```
exec say1("xtrue exited 0")  
    when  
        xtrue.exit_status == 0  
;  
exec say2("xtrue exited 1")  
    when  
        xtrue.exit_status == 1  
;
```

```
$ echo | hoq exec-false.hoq
```

# Logical Boolean Operators Follow Strict SQL Semantics

## Logical AND

```
false and * => false
* and false => false
null and *  => null
* and null  => null
*           => true
```

## Logical OR

```
true or *   => true
* or true   => true
null or *   => null
* or null   => null
*           => false
```

All other binary operators are null if either operand is null

## Idiom #3 - Channels Over Channels

```
// each opcode requests another flow

func (flo *flow) get() *flow {

    // wait for entire query to resolve

    <-flo.resolved

    // next active flow arrives on this channel

    reply := make(flow_chan)

    // request another flow by sending reply channel to main()

    flo.next <- reply

    // return next flow to the this opcode

    return <-reply
}
```

# Flow Structure Synchronizes Operators

```
// a flow tracks the firing of rules over a single line of input text.

type flow struct {

    // request a new flow from this channel, reading reply on sent side-channel
    next chan flow_chan

    // channel is closed when all exec()s make no further progress
    resolved chan struct{}

    // the whole line of input with trailing new line removed
    line string

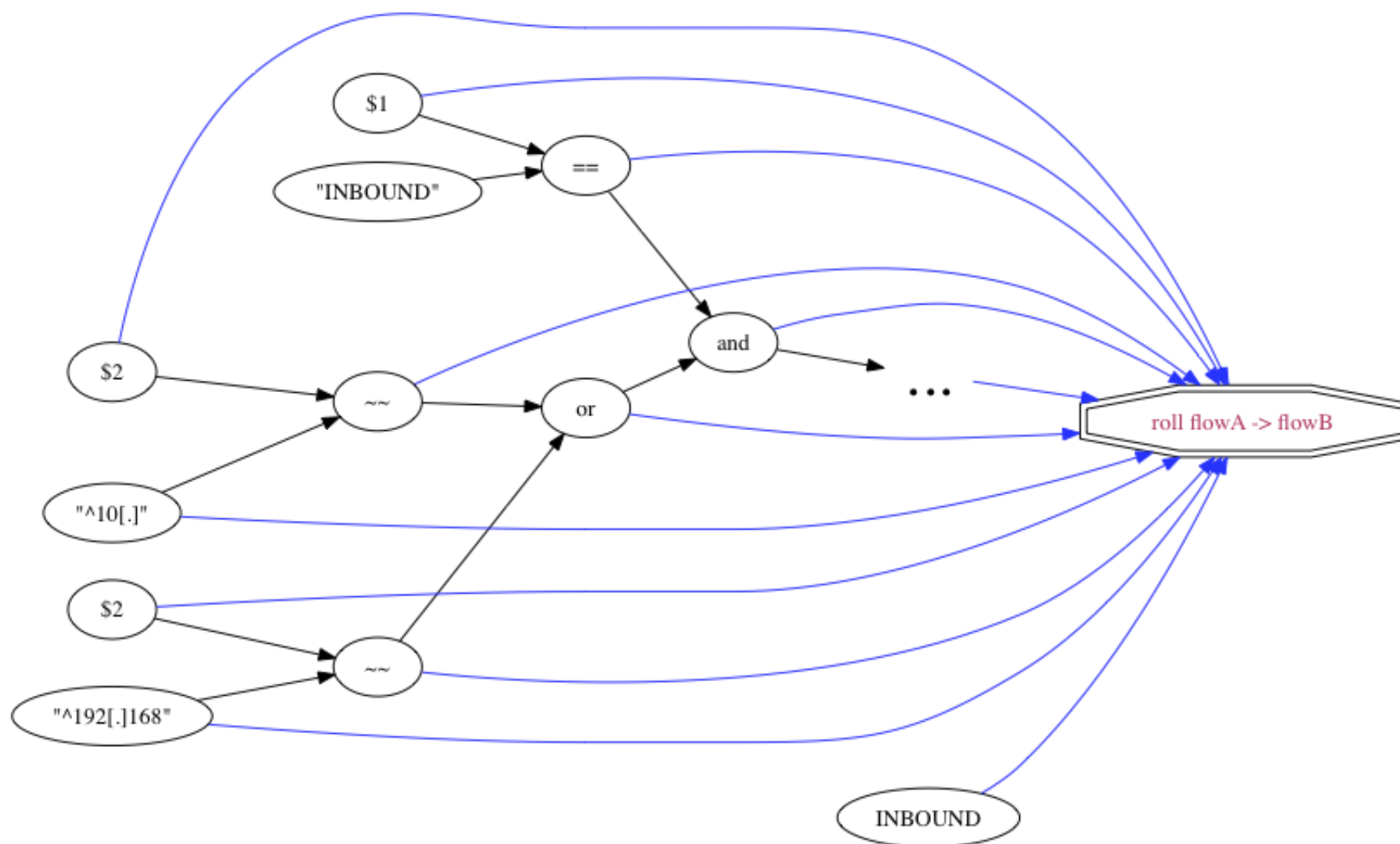
    // tab separated fields split out from the line read from standard input
    fields []string

    // count of go routines/operators still resolving qualifications
    confluent_count int
}

type flow_chan chan *flow
```

Idiom #3 - Channels Over Channels

# Each Opcode Syncs with main() on Each "Tick"



`$1 == "INBOUND" and ( $2 ~~ "^10[.]" or $2 ~~ "^192[.]168[.]" )`

## Idiom #4 - Close(channel) is a Cheap Broadcast

```
// set up first flow to compile all ast nodes into a single flow graph.
// each flow terminates by sending a count of the fired unix processes.

flowA := &flow{
    next:      make(chan flow_chan),
    resolved: make(chan struct{}),
}
uc := flowA.compile(ast, depend_order)           // compile using gnu tsort
close(flowA.resolved)                             // flowA never runs

// start pumping standard input to the flow graph of nodes

in := bufio.NewReader(os.Stdin)
for {
    line, err := in.ReadString('\n')
    ...
}
```

*... Main Loop in Next Slide ...*



```
}
```

# Flow A and Flow B Run Concurrently

```
line = strings.TrimRight(line, "\n")           // trim and split the input line of text
flowB := &flow{
    line:    line,
    fields:  strings.SplitN(line, "\t", 255),
    next:    make(chan flow_chan),
    resolved: make(chan struct{}),
}
for flowA.confluent_count > 0 {                 // push flowA to flowB
    reply := <-flowA.next
    flowA.confluent_count--
    reply <- flowB
    flowB.confluent_count++
}
if <-uc == nil {                               // wait for flowB to finish
    break
}
close(flowB.resolved)                          // broadcast to all nodes in flowB the flow is done
flowA = flowB                                  // and so the wheel turns
```

Idiom #4 - Close(channel) is a Cheap Broadcast



# What is Go YACC?

- Compiles a Flavor of Backus-Naur Form (BNF) into Go Code
- Original Golang Grammar written in YACC
- Domain Specific Languages (chemical reactions, SAT solvers)
- Mutation Coverage

# Mechanics of YACC

- YACC generated Go code reads a stream of integers called <TOKEN>s
- Patterns are recognized in stream of <TOKEN>s
- Each pattern has an associated block of manually written Go code
- When the pattern is recognized in the stream then associated block of Go is invoked
- Patterns are recursive

# Backus-Naur Grammar of Backus-Naur Form

```
alpha:      'a' | 'b' | ... 'Z'
            ;
alphanum:   alpha | '0' | '1' ... '9'
            ;
name:       alpha
            |
            name alphanum
            ;
term:       <TOKEN>
            |
            name
            |
            term term
            ;
expression: term
            |
            expression '|' term
            ;
production: name ':' expression ';'
            ;
grammar:    production
            |
            grammar production
            ;
```

Whitespace/Comment not specified. <TOKEN> is scanned integers, 'a'=97

# YACC Grammar for HOQ Language (Backus-Naur Form)

```

exp:
    TRUE      |    FALSE      |    STRING      |    UINT8      |    '$' UINT8      |
    XCOMMAND '.' EXIT_STATUS
    exp RE_MATCH exp      |    exp RE_NMATCH exp
    exp AND exp |    exp OR exp |    exp EQ exp |    exp NEQ exp
    NOT exp      |    '(' exp ')'
;

exp_list:
    exp      |    exp_list ',' exp
;

argv:
    /*empty*/      |    exp_list
;

qualification:
    /*empty*/      |    WHEN exp
;

string_list:
    STRING      |    string_list ',' STRING
;

command_argv:
    /*empty*/      |    '(' ')'      |    '(' string_list ')'
;

statement:
    COMMAND NAME command_argv '{' PATH '=' STRING ';' '}'
    EXEC XCOMMAND '(' argv ')' qualification ';'
;

statement_list:
    statement      |    statement_list statement
;

```

Stripped from <https://github.com/jmscott/play/blob/master/hoq/src/parser.y>

# YACC/Go Code Snippet of Qualification Expression in HOQ Grammar

```
exp AND exp          // logical and of two qualifications
{
    l := yylex.(*yyLexState)
    $$ = l.bool_node(AND, $1, $3)
    if $$ == nil {
        return 0
    }

    if $1.go_type != reflect.Bool {
        l.error("logical 'and' requires boolean operands")
        return 0
    }
}
|
 '(' exp ')'          // change precedence of qualification
{
    $$ = $2
}
```

In YACC \$1 is value of first term; \$2 is value of second term ...  
\$\$ is value of new, reduced term

# YACC Needs a Lexer

- lexer easy to write by hand !!!
- just a big switch{}
- busts text into stream of integer tokens
- scanned token has a simple value related to scan yylval.(string, float, uint8)
- token const defined by YACC (see %token in grammar)
- co-routines invented to write lexers (Rob Pike)

# Go Files on [github.com/jmsscott](https://github.com/jmsscott)

- `hoq.go`
- `ast.go`
- `command.go`
- `compile.go`
- `opcode.go`
- `parser.y`
- `rummy.go`
- `tsort.go`

[github.com/jmsscott/play/tree/master/hoq/src](https://github.com/jmsscott/play/tree/master/hoq/src) (<https://github.com/jmsscott/play/tree/master/hoq/src>)



# YACC Resources

Early GoLang in YACC (<https://docs.google.com/document/d/1P3BLR31VA8cvLJLfMibSuTdwTuF7WWLux71CYD0eeD8/edit>)

Book 'Unix Programming Environment' by Kernighan and Pike (<http://www.amazon.com/Unix-Programming-Environment-Prentice-Hall-Software/dp/013937681X>)

C Written in Yacc (<http://heim.ifi.uio.no/inf2270/programmer/historien-om-C.pdf>)

Simple Go Calculator (<https://github.com/golang-samples/yacc>)

Go AST Library (<https://golang.org/pkg/go/ast/>)

EBNF in GoLang (<https://godoc.org/golang.org/x/exp/ebnf>)

# Lex Resources

Lexer Talk by Rob Pike ([https://www.youtube.com/watch?v=HxaD\\_trXwRE](https://www.youtube.com/watch?v=HxaD_trXwRE))

Ragle is a Lexical Compiler (<http://www.colm.net/open-source/ragel/>)

Nex (New Lex) (<http://www-cs-students.stanford.edu/~blynn/nex/>)

# CSP Resources

[research.swtch.com/power](http://research.swtch.com/power) (<http://research.swtch.com/power>)

# Thank you

## John Scott

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