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Department of Computer Science

COS212: Practical 6

Release: Thursday 11 April 2019, 18:00
Deadline: Friday 12 April 2019, 18:00

Objectives

This practical has the following objectives:

- To observe the similarities between trees and graphs.
- To implement a Trie using a graph structure.
- To utilize the Trie data structure for approximate string matching.

Instructions

Complete the task below. Certain classes have been provided for you in the *files* zip archive of the practical. You have also been given a main file which will test some code functionality, but it is by no means intended to provide extensive test coverage. You are encouraged to edit this file and test your code more thoroughly. Remember to test boundary cases. Upload **only** the given source files with your changes in a zip archive before the deadline. Please comment your name **and** student number in at the top of each file.

Introduction

You will be required to implement a *trie* using a graph structure. Generally a tree only has nodes, where nodes store reference to other nodes. In this practical a trie will be implemented that as nodes as well as edges. A node will store reference to edges instead of nodes. An edge will contain a label and reference to the target node.

The Trie will have a root node that never contains a key. If the Trie is empty, then the root node will have no edges going from it. Edges and their labels will be used to distinguish keys from each other. Full keys will not be stored in the nodes, only the suffix. The key can be retrieved by combining the labels of the edges followed with the suffix in the node. If the *hasKey* flag is set and the suffix is an empty string, then the key is contained entirely in the series of labels followed to the node.

Figure 1: Legend for Figures in this document

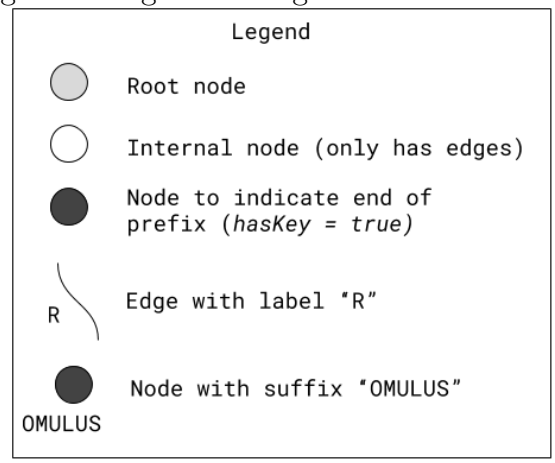

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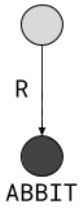
Figure 2: Trie with the key "RABBIT"



By following edges from the root node, common prefixes can be found. Figure. 2 shows the Trie with one key: "RABBIT". The root node has one edge labeled "R". The edge has a target node of the node containing the suffix "ABBIT". Since there is not a common prefix the only edge in the Trie is the edge from root with the label "R".

Figure 3: After adding the key "CAT" to the trie

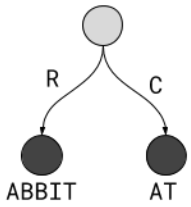


Figure. 3 follows from Figure. 2 with the key "CAT" added. There are no common prefixes so only a single edge from the root is added.

Figure 4: After adding the keys "RABIES" and "RAB" to the trie

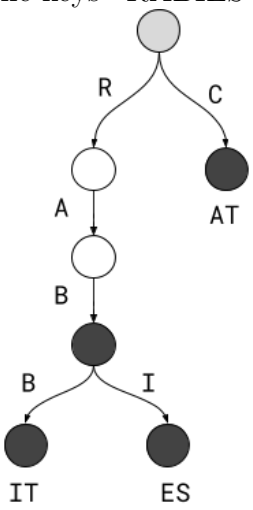


Figure. 4 follows from Figure. 3 with the keys "RABIES" and "RAB" added. This time "RAB" becomes a common prefix so the nodes and edges for "A" and "B" are created in

series. Since "RABBIT" and "RABIES" have the same prefix of "RAB", the edges labeled "B" and "I" are created to distinguish the two keys from each other. Since the key "RAB" is fully contained in the path of edges, the node contains an empty suffix.

Nodes that indicate that a key can be retrieved have the*hasKey* flag set. The key will be composed from the series of edge labels from the root to the node, combined with the suffix stored in the node.

The *edges* array should be sorted alphabetically according to the edge label.

For information on tries, see section 7.2 in the textbook. For information on graphs see section 8.1 and 8.3 in the textbook. The trie in this practical will function similar to the trie described in the textbook, however it will be represented using nodes and edges and only the suffix will be stored in nodes.

With the graph representation we lose the random access provided by the implicit relationship between characters in the alphabet and their indices in the array of pointers to sub trees. For the purpose of this practical that compromise is acceptable.

Since tries provide efficient searching and retrieval of strings they are a good data structure for applications such as autocorrect. You are also required to implement a function to return a word from the trie that is closest to a given word in terms in Edit Distance. The edit distance indicates the minimum number of insert/substitute/delete operations to edit a string to be the same as another string.

Task 1: [22]

Implement the following methods in the *Trie* Class according to the given specification:

```
void insert(String key)
```

Insert the given key into the trie. See the Introduction for information.

```
Boolean contains(String key)
```

Return true if the key exists in the trie, otherwise false.

```
String closestMatch(String str, Integer maxDistance)
```

Return the closest string in the trie using Levenshtein distance. Only consider distances that are in the range [0, maxDistance]. If no match could be found in the given range, return null. If multiple strings have the same Levenshtein distance, return the string that comes first alphabetically.

The Levenshtein distance function [1] is provided to you and has the following signature:

```
int Distance.LD(String s, String t)
```

You may use your own helper functions to assist in implementing the specification. However you may not modify any of the given method signatures.

Submission

You need to submit your source files on the Assignment website (assignments.cs.up.ac.za). All methods need to be implemented (or at least stubbed) before submission. Place all the source files including a makefile in a zip or tar/gzip archive named uXXXXXXXX.zip or uXXXXXXXX.tar.gz where XXXXXXXX is your student number. There should be no folders in your archive. You have 24 hours to finish this practical, regardless of which practical session you attend. Upload your archive to the *2019 Prac 6 - Friday* slot on the Assignment website. Submit your work before the deadline. **No late submissions will be accepted!**

References

[1] Michael Gilleland. Levenshtein distance. http://people.cs.pitt.edu/~kirk/cs1501/Pruhs/Fall2006/Assignments/editdistance/Levenshtein%20Distance.htm.