

# Teambook Sindicato de Transporte 2880

Universidad Mayor de San Simón



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# Chapter 1

## Introduction

The following document represents the Teambook for the team Sindicato de Transporte 2880. This version was elaborated for the Latin American Regional phase of 2022's ICPC.

The template for the C++ code is presented:

```
#include <bits/stdc++.h>
// #include <ext/pb_ds/assoc_container.hpp>
// #include <ext/pb_ds/assoc_container.hpp>
// #include <ext/pb_ds/tree_policy.hpp>
// #include <ext/rope>
#define int ll
#define mp      make_pair
#define pb      push_back
#define all(a)   (a).begin(), (a).end()
#define sz(a)    (int)a.size()
#define eq(a, b) (fabs(a - b) < EPS)
#define md(a, b) ((a) % b + b) % b
#define mod(a)   md(a, MOD)
#define _max(a, b) ((a) > (b) ? (a) : (b))
#define srt(a)   sort(all(a))
#define mem(a, h) memset(a, (h), sizeof(a))
#define f        first
#define s        second
#define forn(i, n) for(int i = 0; i < n; i++)
#define fore(i, b, e) for(int i = b; i < e; i++)
#define forg(i, b, e, m) for(int i = b; i < e; i+=m)
#define index int mid = (b + e) / 2, l = node * 2 + 1, r = l + 1;
#define DBG(x) cerr<<#x<<"_="<<(x)<<endl
#define RAYA cout<<"===== "<<'\n'
```

```
// int in(){int r=0,c;for(c=getchar();c<=32;c=getchar());if(c=='-')
return -in();for(;c>32;r=(r<<1)+(r<<3)+c-'0',c=getchar());
return r;}
```

```
using namespace std;
// using namespace __gnu_pbds;
// using namespace __gnu_cxx;

// #pragma GCC target ("avx2")
// #pragma GCC optimization ("O3")
// #pragma GCC optimization ("unroll-loops")

typedef long long ll;
typedef long double ld;
typedef unsigned long long ull;
typedef pair<int, int> ii;
typedef pair<pair<int, int>, int> iii;
typedef vector<int> vi;
typedef vector<ii> vii;
typedef vector<ll> vll;
// typedef tree<int,null_type,less<int>,rb_tree_tag,
// tree_order_statistics_node_update> ordered_set;
// find_by_order kth largest order_of_key <
// mt19937 rng(chrono::steady_clock::now().time_since_epoch().count
// ());
// rng
const int tam = 200010;
const int MOD = 1000000007;
```

```
const int MOD1 = 998244353;
const double DINF=1e100;
const double EPS = 1e-9;
const long double PI = acos(-1.0L);

signed main()
{
    ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    // freopen("asd.txt", "r", stdin);
    // freopen("qwe.txt", "w", stdout);

    return 0;
}
```

../Template.cpp

## Chapter 2

# Data Structures

### 2.1 Union Find Disjoint Sets

The Union Find Disjoint Sets data structure is used to keep track of a set of elements partitioned into a number of disjoint (non-overlapping) subsets. It supports two operations:

- **Find:** Determine which subset a particular element is in. This can be used for determining if two elements are in the same subset.
- **Union:** Join two subsets into a single subset.

The implementation is as follows:

```
11 dsu[tam];

int getParent(int x){
    if(dsu[x]<0) return x;
    else return dsu[x] = getParent(dsu[x]);
}

void join(int x, int y){
    x = getParent(x);
    y = getParent(y);
    if(x==y) return;
    if(dsu[x]>dsu[y]) swap(x, y);
    dsu[x] = y;
}
```

```
//With structs
struct unionFind {
    vi p;
    unionFind(int n) : p(n, -1) {}
    int findParent(int v) {
        if (p[v] == -1) return v;
        return p[v] = findParent(p[v]);
    }
    bool join(int a, int b) {
        a = findParent(a);
        b = findParent(b);
        if (a == b) return false;
        p[a] = b;
        return true;
    }
};
```

../Data Structures/UnionFind.cpp

### 2.2 Binary Indexed (Fenwick) Tree

The Binary Indexed Tree (BIT) is a data structure that can efficiently update elements and calculate prefix sums in a table of numbers. It is also called a Fenwick Tree, as it was first described by Peter Fenwick. The implementation is as follows:

```
#define clr(a,h)      memset(a,(h),sizeof(a))
```

```

int BIT[tam];

void update(int pos, int val)
{
    pos++;
    while(pos < 200010)
    {
        BIT[pos] += val;
        pos += (pos & -pos);
    }
}

int query(int pos)
{
    pos++;
    int res = 0;
    while(pos > 0)
    {
        res += BIT[pos];
        pos -= (pos & -pos);
    }
    return res;
}

int main()
{
    clr(BIT,0);
    for(int i = n - 1; i >=0; i--)
    {
        inv +=query(a[i]);
        update(a[i],1);
    }
}

//to update all the values in a range [i,j] the following
//implementation is used

range_update(i,j,val) = update(i,val); update(j+1,-val);

//to find the range result of a range [i,j] for a Range Update

```

### Range Query

```
//rsq(1,j) = rupq.point_query(j)*j-purq.rsq(j)
```

../Data Structures/FenwickTree.cpp

This Data Structure can also be used to find the maximum value in a range of values in an array.

It can also be extended to 2D, 3D, etc. The implementation is as follows:

```

#include <iostream>

using namespace std;

const int tam = 1000;

int BIT[tam][tam];
int n, m;

void update(int row, int col, int val)
{
    row++; col++;
    for (int i = row; i <= n; i += (i & -i))
    {
        for (int j = col; j <= m; j += (j & -j))
        {
            BIT[i][j] += val;
        }
    }
}

int query(int row, int col)
{
    int res = 0;
    row++; col++;
    for (int i = row; i > 0; i -= (i & -i))
    {
        for (int j = col; j > 0; j -= (j & -j))
        {
            res += BIT[i][j];
        }
    }
    return res;
}

```



```
}
```

../Data Structures/FenwickTree2D.cpp

## 2.3 Segment Tree

The Segment Tree is a data structure that allows answering range queries over an array effectively, while still being flexible enough to allow modifying the array. It is, in principle, a static structure. It can answer most queries in  $O(\log n)$ , but its true power is answering range updates. For that, it takes  $O(\log n)$  time per update.

The implementation is as follows:

```
struct node
{
    int val;
};
node join(node a, node b)
{
    a.val += b.val;
    return a;
}
int ar[tam];
node t[4 * tam];
void init(int b, int e, int node)
{
    if(b == e)
    {
        t[node].val = ar[b];
        return;
    }
    int mid = (b + e) / 2, l = node * 2 + 1, r = l + 1;
    init(b, mid, l);
    init(mid + 1, e, r);
    t[node] = join(t[l], t[r]);
}
//b, e always the beginning and end of the segment tree;

node query(int b, int e, int node, int i, int j)
{
    if(b >= i && e <= j)
```

```
    return t[node];
    int mid = (b + e) / 2, l = node * 2 + 1, r = l + 1;
    if(mid < i)
        return query(mid + 1, e, r, i, j);
    if(mid >= j)
        return query(b, mid, l, i, j);
    return join(query(b, mid, l, i, j), query(mid + 1, e, r, i, j));
}
void update(int b, int e, int node, int pos, int val)
{
    if(b == e) {t[node].val = val;return;} // Replaces the value,
    could be updated to add val to the current value
    int mid = (b + e) / 2, l = node * 2 + 1, r = l + 1;
    if(mid < pos)
        update(mid + 1, e, r, pos, val);
    else
        update(b, mid, l, pos, val);
    t[node] = join(t[l], t[r]);
}
```

../Data Structures/SegmentTree.cpp

An iterative implementation is also possible:

Another useful tool while using the Segment Tree is the Lazy Propagation technique. This technique allows us to perform range updates in  $O(\log n)$  time. The implementation is as follows:

In order to do operations on segments without having to worry about the lazy propagation, we can use the following functions:

```
struct segtree {
    int size;
    vi operations;
    void init(int n) {
        size = 1;
        while (size < n) size *= 2;
        operations.assign(2 * size, 0LL);
    }

    void add(int l, int r, int v, int x, int lx, int rx){
        if(lx >= r || l >= rx) return;
        if(lx >= l && rx <= r) {
            operations[x] += v;
            return;
        }
```

```

    }
    int m = (lx + rx) / 2;
    add(l, r, v, 2 * x + 1, lx, m);
    add(l, r, v, 2 * x + 2, m, rx);
}

void add(int l, int r, int v) {
    add(l, r, v, 0, 0, size);
}

int get(int i, int x, int lx, int rx){
    if(rx - lx == 1) return operations[x];
    int m = (lx + rx) / 2;
    int res;
    if(i < m) res = get(i, 2 * x + 1, lx, m);
    else res = get(i, 2 * x + 2, m, rx);
    return res + operations[x];
}

int get(int i) {
    return get(i, 0, 0, size);
}

};

signed main()
{
    ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    // freopen("asd.txt", "r", stdin);
    // freopen("qwe.txt", "w", stdout);

    int n;
    cin >> n;
    int q;
    cin >> q;

    segtree st;

```

```

    st.init(n);

    while(q--){
        int c;
        cin >> c;
        if(c==1){
            int l,r, v;
            cin >> l >> r >> v;
            st.add(l,r,v);
            continue;
        }

        int i;
        cin >> i;
        cout << st.get(i) << '\n';

    }

    return 0;
}

```

../Data Structures/Operations.cpp

This second implementation only works for operations that have commutative and associative properties. Emphasis on the commutativity of the operation.

## Chapter 3

# Mathematics

This chapter is about some useful mathematical tools needed in order to solve problems.

### 3.1 GCD and LCM

In order to find the greatest common divisor (GCD) of two numbers, the Euclidean algorithm can be used. The implementation is as follows:

```
11 gcd(11 a, 11 b){return b==0? a:gcd(b,a%b);}

int x, y, d;
void extendedEuclid(int a, int b)//ecuacion diofantica ax + by = d
{
    if(b==0) {x=1; y=0; d=a; return;}
    extendedEuclid(b,a%b);
    int x1=y;
    y = x-(a/b)*y;
    x=x1;
}
```

../Mathematics/Euclid.cpp

Another (and faster) way to find the GCD is by using the following code:

```
int gcd(int a, int b) {
    if (!a || !b)
        return a | b;
    unsigned shift = __builtin_ctz(a | b);
    a >>= __builtin_ctz(a);
```

```
do {
    b >>= __builtin_ctz(b);
    if (a > b)
        swap(a, b);
    b -= a;
} while (b);
return a << shift;
}
```

../Mathematics/FastGCD.cpp

The way Halim suggests to find the GCD and the LCM is given by the following code:

```
int gcd(int a, int b) { return b == 0 ? a : gcd(b, a%b); }
int lcm(int a, int b) { return a / gcd(a, b) * b; }
```

../Mathematics/HalimGCD.cpp

### 3.2 Prime Numbers

The fastest way to check the primality of a number is by using Erathostenes' sieve. The typical implementation is as follows:

```
bitset<100000> bi;
vi primos; //primos
vector<ll> pric; //primos al cuadrado
void criba()
{
    bi.set();
```

```

for(int i=2;i<100000;i++)
    if(bi[i])
    {
        for(int j=i+i;j<100000;j+=i)
            bi[j]=0;
        primos.push_back(i);
        pric.push_back((ll)i*(ll)i);
    }
}
int euler(int n)
{
    int res=n;
    for(int i=0;pric[i]<=n;i++)
    {
        if(n%primos[i]==0)
        {
            res-= res/primos[i];
            while(n%primos[i]==0) n/=primos[i];
        }
    }
    if(n!=1) res-=res/n;
    return res;
}

```

../Mathematics/Erathostenes.cpp

Nevertheless, the following implementation is faster, since the statement `if (i % prime[j] == 0) break;` terminates the loop when `p` divides `i`. The inner loop is executed only once for each composite. Hence, the code performs in  $O(n)$  complexity, resulting in the 'linear' sieve:

```

// This algorithm allows to find Eratosthenes sieve in  $O(n \log n)$ 
// time.

std::vector<int> prime;
bool is_composite[MAXN];

void sieve (int n) {
    std::fill (is_composite, is_composite + n, false);
    for (int i = 2; i < n; ++i) {
        if (!is_composite[i]) prime.push_back (i);

```

```

        for (int j = 0; j < prime.size () && i * prime[j] < n; ++j) {
            is_composite[i * prime[j]] = true;
            if (i % prime[j] == 0) break;
        }
    }

    // An application of this linear sieve is to find the Euler
    // totient function of a number in  $O(n \log n)$  time.
    std::vector<int> prime;
    bool is_composite[MAXN];
    int phi[MAXN];

    void sieve (int n) {
        std::fill (is_composite, is_composite + n, false);
        phi[1] = 1;
        for (int i = 2; i < n; ++i) {
            if (!is_composite[i]) {
                prime.push_back (i);
                phi[i] = i - 1; //i is prime
            }
            for (int j = 0; j < prime.size () && i * prime[j] < n; ++j) {
                is_composite[i * prime[j]] = true;
                if (i % prime[j] == 0) {
                    phi[i * prime[j]] = phi[i] * prime[j]; //prime[j]
                    divides i
                    break;
                } else {
                    phi[i * prime[j]] = phi[i] * phi[prime[j]]; //prime[j]
                    does not divide i
                }
            }
        }
    }
}

```

../Mathematics/LinearSieve.cpp

### 3.3 Modular Arithmetic

The modular inverse is defined by the following equation:

$$a \cdot a^{-1} \equiv 1 \pmod{m} \quad (3.1)$$

The following code shows how to find the modular inverse of a number:

```
int ModPow(int a, int b, int m) {
    int res = 1;
    while (b > 0) {
        if (b & 1) res = (res * a) % m;
        a = (a * a) % m;
        b >>= 1;
    }
    return res;
}

// Language: java
public static int modPow(int a, int b, int m) {
    int res = 1;
    while (b > 0) {
        if ((b & 1) == 1) res = (res * a) % m;
        a = (a * a) % m;
        b >>= 1;
    }
    return res;
}

int ModInverse(int a, int m) {
    return ModPow(a, m - 2, m);
}
```

../Mathematics/ModularInverse.cpp

Some other useful relationships in modular arithmetic are:

- $(a + b) \pmod{m} = (a \pmod{m} + b \pmod{m}) \pmod{m}$
- $(a - b) \pmod{m} = (a \pmod{m} - b \pmod{m}) \pmod{m}$
- $(a * b) \pmod{m} = (a \pmod{m} * b \pmod{m}) \pmod{m}$
- $(a/b) \pmod{m} = (a \pmod{m} * b^{-1} \pmod{m}) \pmod{m}$
- $(a^b) \pmod{m} = (a \pmod{m})^b \pmod{m}$
- $(a^b) \pmod{m} = (a \pmod{m})^{b \pmod{\phi(m)}} \pmod{m}$

- $(a^b) \pmod{m} = (a \pmod{m})^{b \pmod{m-1}} \pmod{m}$
- $\frac{a}{k} \equiv \frac{a}{k} \pmod{m} \iff a \equiv k \pmod{m}$
- $\frac{a}{k} \equiv \frac{a}{k} \pmod{\frac{n}{\gcd(n,k)}}$

## 3.4 Matrix Exponentiation

The following code shows how to find the nth power of a *mat*, noting that a data structure of type matrix is defined as follows:

```
typedef vector<vector<ll>> mat;
mat ans;
void mult(mat m1, mat m2)
{
    assert(m1[0].size() == m2.size());
    ans.clear();
    ll answer = 0;
    for(i, 0, m1.size())
    {
        vector<ll> fila;
        for(j, 0, m2[0].size())
        {
            answer = 0;
            for(k, 0, m2.size())
                answer = (answer + m1[i][k] * m2[k][j]) % MOD;
            fila.pb(answer);
        }
        ans.pb(fila);
    }
}

void pot(mat base, ll exp)
{
    mat res(base.size(), vector<ll>(base.size(), 0));
    for(i, 0, base.size())
        res[i][i] = 1;
    while(exp)
    {
        if(exp & 1)
        {
            mult(res, base);
        }
    }
}
```

```

    res = ans;
}
mult(base, base);
base = ans;
exp /= 2;
}
ans = res;
}

```

../Mathematics/MatrixPower.cpp

### 3.5 Gauss Jordan Elimination

The following code shows how to solve a system of linear equations using Gauss Jordan elimination:

```

// resuelve Ax = b, dada la matriz a de n * (m + 1), n ecuaciones y
// m variables, siendo la ultima columna el vector b
// The function returns the number of solutions of the system (0,1,
// or INF). if there's at least a solution, it's in ans
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be infinity or a
// big number
int gauss (vector < vector<double> > a, vector<double> & ans) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;

    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {
        int sel = row;
        for (int i=row; i<n; ++i)
            if (abs (a[i][col]) > abs (a[sel][col]))
                sel = i;
        if (abs (a[sel][col]) < EPS)
            continue;
        for (int i=col; i<=m; ++i)
            swap (a[sel][i], a[row][i]);
        where[col] = row;

        for (int i=0; i<n; ++i)
            if (i != row) {

```

```

                double c = a[i][col] / a[row][col];
                for (int j=col; j<=m; ++j)
                    a[i][j] -= a[row][j] * c;
            }
        ++row;
    }

    ans.assign (m, 0);
    for (int i=0; i<m; ++i)
        if (where[i] != -1)
            ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {
        double sum = 0;
        for (int j=0; j<m; ++j)
            sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
            return 0;
    }

    for (int i=0; i<m; ++i)
        if (where[i] == -1)
            return INF;
    return 1;
}

```

../Mathematics/GaussJordan.cpp

### 3.6 Determinant

The following code shows how to find the determinant of a matrix:

```

#include <iostream>

using std::cin;
using std::cout;
using std::endl;

int **submatrix(int **matrix, unsigned int n, unsigned int x,
                unsigned int y) {
    int **submatrix = new int *[n - 1];
    int subi = 0;

```

```

    for (int i = 0; i < n; i++) {
        submatrix[subi] = new int[n - 1];
        int subj = 0;
        if (i == y) {
            continue;
        }
        for (int j = 0; j < n; j++) {
            if (j == x) {
                continue;
            }
            submatrix[subi][subj] = matrix[i][j];
            subj++;
        }
        subi++;
    }
    return submatrix;
}

int determinant(int **matrix, unsigned int n) {
    int det = 0;
    if (n == 2) {
        return matrix[0][0] * matrix[1][1] - matrix[1][0] * matrix[0][1];
    }
    for (int x = 0; x < n; ++x) {
        det += ((x % 2 == 0 ? 1 : -1) * matrix[0][x] * determinant(
            submatrix(matrix, n, x, 0), n - 1));
    }

    return det;
}

int main() {
    int n;
    cin >> n;
    int **matrix = new int *[n];
    for (int i = 0; i < n; ++i) {
        matrix[i] = new int[n];
        for (int j = 0; j < n; ++j) {
            cin >> matrix[i][j];
        }
    }
}

```

```

    }

    cout << determinant(matrix, n);

    return 0;
}

```

../Mathematics/Determinant.cpp

## 3.7 Numerical Integration

The following code shows how to find the integral of a function  $f(x)$  in the interval  $[a,b]$  using Simpson's rule:

```

double simpson(double f(double), double a, double b)
{
    int n = 100000;
    double s = f(a) + f(b);
    double h = (b - a) / n;
    for(i, 1, n)
        s += ((i & 1) ? 4 : 2) * f(a + h * i);
    return s * (h / 3);
}

```

../Mathematics/Simpson.cpp





## Chapter 4

# Graphs

This chapter shows some of the basic algorithms and implementations required to solve problems that include graphs.

### 4.1 Depth First Search (DFS)

The DFS algorithm is a recursive algorithm that visits all the nodes of a graph. It is used to find connected components, topological sorting, and to find bridges and articulation points. The algorithm is as follows:

The implementation can be done as follows:

```
vector<vector<int>> g(tam);
vector<bool> vis(tam);

void dfs(int u){
    vis[u]=true;
    ans++;
    for(int v: g[u]){
        if(!vis[v]){
            dfs(v);
        }
    }
}

signed main()
{
    int n,m;
    cin>>n>>m; // n nodes, m edges
    g.assign(tam,vector<int>());
```

```
    vis.assign(tam, false);
    for(int i=0; i<m;i++){
        int u,v;
        cin>>u>>v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    ll res = 0;
    for(int i=1; i<=n;i++){
        if(!vis[i]){
            ans=0;
            dfs(i);
            res = max(res,ans);
        }
    }
    g.clear();
    vis.clear();
    return 0;
}
```

../Graphs/DFS.cpp

An application of this algorithm in order to find the shortest path between two nodes can be done as follows:

```
// The following code represents the implementation of a DFS
// algorithm
// to find the shortest path between two nodes in a graph.
// The graph is represented as an adjacency list.
// The algorithm is implemented using a stack.
```

**Algorithm 1** Depth First Search (DFS)

---

```

1: procedure DFS( $G$ )
2:    $visited \leftarrow \emptyset$ 
3:    $time \leftarrow 0$ 
4:    $parent \leftarrow \emptyset$ 
5:    $low \leftarrow \emptyset$ 
6:    $disc \leftarrow \emptyset$ 
7:    $AP \leftarrow \emptyset$ 
8:    $bridge \leftarrow \emptyset$ 
9:   for all  $v \in V$  do
10:     $visited[v] \leftarrow false$ 
11:     $parent[v] \leftarrow -1$ 
12:     $low[v] \leftarrow \infty$ 
13:     $disc[v] \leftarrow \infty$ 
14:  end for
15:  for all  $v \in V$  do
16:    if  $visited[v] = false$  then
17:      DFSUtil( $G, v, visited, time, parent, low, disc, AP, bridge$ )
18:    end if
19:  end for
20: end procedure
21: procedure DFSUTIL( $G, v, visited, time, parent, low, disc, AP, bridge$ )
22:    $visited[v] \leftarrow true$ 
23:    $disc[v] \leftarrow time$ 
24:    $low[v] \leftarrow time$ 
25:    $time \leftarrow time + 1$ 
26:    $children \leftarrow 0$ 
27:   for all  $u \in Adj(v)$  do
28:     if  $visited[u] = false$  then
29:        $parent[u] \leftarrow v$ 
30:        $children \leftarrow children + 1$ 
31:       DFSUtil( $G, u, visited, time, parent, low, disc, AP, bridge$ )
32:        $low[v] \leftarrow \min(low[v], low[u])$ 
33:       if  $parent[v] = -1$  and  $children > 1$  then
34:          $AP[v] \leftarrow true$ 
35:       end if
36:       if  $parent[v] \neq -1$  and  $low[u] \geq disc[v]$  then
37:          $AP[v] \leftarrow true$ 
38:       end if
39:       if  $low[u] > disc[v]$  then
40:          $bridge[v][u] \leftarrow true$ 
41:       end if
42:     else
43:        $low[v] \leftarrow \min(low[v], disc[u])$ 
44:     end if
45:   end for

```

```

#include <bits/stdc++.h>

using namespace std;

vector<int> DFS(vector<vector<int>> &adj, int s, int t) {
    stack<vector<int>> path_stack;
    vector<int> path;
    vector<int> visited(adj.size(), 0);
    path_stack.push({s});
    while (!path_stack.empty()) {
        path = path_stack.top();
        path_stack.pop();
        int last = path[path.size() - 1];
        if (last == t) {
            return path;
        }
        if (visited[last] == 0) {
            visited[last] = 1;
            for (int i = 0; i < adj[last].size(); i++) {
                if (visited[adj[last][i]] == 0) {
                    vector<int> new_path(path);
                    new_path.push_back(adj[last][i]);
                    path_stack.push(new_path);
                }
            }
        }
    }
    return {};
}

int main() {
    int n, m;
    cin >> n >> m;
    vector<vector<int>> adj(n, vector<int>());
    for (int i = 0; i < m; i++) {
        int x, y;
        cin >> x >> y;
        adj[x - 1].push_back(y - 1);
        adj[y - 1].push_back(x - 1);
    }
}

```

```

}
int x, y;
cin >> x >> y;
x--, y--;
vector<int> path = DFS(adj, x, y);
for (int i = 0; i < path.size(); i++) {
    cout << path[i] + 1 << " ";
}
}

```

../Graphs/DFS-application.cpp

## 4.2 Breadth First Search (BFS)

The BFS algorithm is a non-recursive algorithm that visits all the nodes of a graph. It is used to find connected components, topological sorting, and to find bridges and articulation points, to better understand it, a propagating fire can be imagined. The algorithm is as follows:

The implementation can be done as follows:

```

#include <bits/stdc++.h>

using namespace std;
signed main()
{
    vector<vector<int>> adj; // adjacency list representation
    int n; // number of nodes
    int s; // source vertex

    queue<int> q;
    vector<bool> used(n);
    vector<int> d(n), p(n);

    q.push(s);
    used[s] = true;
    p[s] = -1;
    while (!q.empty()) {
        int v = q.front();
        q.pop();
        for (int u : adj[v]) {
            if (!used[u]) {

```

---

### Algorithm 2 Breadth First Search (BFS)

---

```

1: procedure BFS(G)
2:   visited ← ∅
3:   time ← 0
4:   parent ← ∅
5:   low ← ∅
6:   disc ← ∅
7:   AP ← ∅
8:   bridge ← ∅
9:   for all v ∈ V do
10:    visited[v] ← false
11:    parent[v] ← -1
12:    low[v] ← ∞
13:    disc[v] ← ∞
14:   end for
15:   for all v ∈ V do
16:    if visited[v] = false then
17:      BFSUtil(G, v, visited, time, parent, low, disc, AP, bridge)
18:    end if
19:   end for
20: end procedure
21: procedure BFSUtil(G, v, visited, time, parent, low, disc, AP, bridge)
22:   visited[v] ← true
23:   disc[v] ← time
24:   low[v] ← time
25:   time ← time + 1
26:   children ← 0
27:   for all u ∈ Adj(v) do
28:    if visited[u] = false then
29:      parent[u] ← v
30:      children ← children + 1
31:      BFSUtil(G, u, visited, time, parent, low, disc, AP, bridge)
32:      low[v] ← min(low[v], low[u])
33:      if parent[v] = -1 and children > 1 then
34:        AP[v] ← true
35:      end if
36:      if parent[v] != -1 and low[u] ≥ disc[v] then
37:        AP[v] ← true
38:      end if
39:      if low[u] > disc[v] then
40:        bridge[v][u] ← true
41:      end if
42:    else
43:      low[v] ← min(low[v], disc[u])
44:    end if
45:   end for

```

```

        used[u] = true;
        q.push(u);
        d[u] = d[v] + 1;
        p[u] = v;
    }
}
return 0;
}

```

../Graphs/BFS.cpp

### 4.3 Finding Bridges and Articulation Points

The following algorithms are used to find bridges and articulation points in a graph. The implementation of these algorithms is done using DFS and BFS. This algorithms are based on Tarjan's algorithm.

Tarjan's algorithm is an algorithm that is used to find bridges and articulation points in a graph. The algorithm is as follows:

```

int n;
vector<vector<int>> adj;

vector<bool> visited;
vector<int> tin, low;
int timer;
vector<vector<int>> comps; // componentes biconexas
stack<int> stk;

void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    stk.push(v);
    int children=0;
    for (int to : adj[v]) {
        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else {
            dfs(to, v);
            low[v] = min(low[v], low[to]);
        }
    }
}

```

---

#### Algorithm 3 Tarjan's Algorithm

---

```

1: procedure TARJAN( $G$ )
2:    $visited \leftarrow \emptyset$ 
3:    $time \leftarrow 0$ 
4:    $parent \leftarrow \emptyset$ 
5:    $low \leftarrow \emptyset$ 
6:    $disc \leftarrow \emptyset$ 
7:    $AP \leftarrow \emptyset$ 
8:    $bridge \leftarrow \emptyset$ 
9:   for all  $v \in V$  do
10:     $visited[v] \leftarrow false$ 
11:     $parent[v] \leftarrow -1$ 
12:     $low[v] \leftarrow \infty$ 
13:     $disc[v] \leftarrow \infty$ 
14:   end for
15:   for all  $v \in V$  do
16:     if  $visited[v] = false$  then
17:       TarjanUtil( $G, v, visited, time, parent, low, disc, AP, bridge$ )
18:     end if
19:   end for
20: end procedure
21: procedure TARJANUTIL( $G, v, visited, time, parent, low, disc, AP, bridge$ )
22:    $visited[v] \leftarrow true$ 
23:    $disc[v] \leftarrow time$ 
24:    $low[v] \leftarrow time$ 
25:    $time \leftarrow time + 1$ 
26:    $children \leftarrow 0$ 
27:   for all  $u \in Adj(v)$  do
28:     if  $visited[u] = false$  then
29:        $parent[u] \leftarrow v$ 
30:        $children \leftarrow children + 1$ 
31:       TarjanUtil( $G, u, visited, time, parent, low, disc, AP, bridge$ )
32:        $low[v] \leftarrow \min(low[v], low[u])$ 
33:       if  $parent[v] = -1$  and  $children > 1$  then
34:          $AP[v] \leftarrow true$ 
35:       end if
36:       if  $parent[v] \neq -1$  and  $low[u] \geq disc[v]$  then
37:          $AP[v] \leftarrow true$ 
38:       end if
39:       if  $low[u] > disc[v]$  then
40:          $bridge[v][u] \leftarrow true$ 
41:       end if
42:     else
43:        $low[v] \leftarrow \min(low[v], disc[u])$ 
44:     end if
45:   end for

```

```

        if (low[to] >= tin[v])
        {
            if (p != -1) IS_CUTPOINT(v);

            comps.push_back({v});
            while (comps.back().back() != to)
            {
                comps.back().push_back(stk.top());
                stk.pop();
            }

        }
        if (low[to] > tin[v])
            IS_BRIDGE(v, to);
        ++children;
    }
}
if(p == -1 && children > 1)
    IS_CUTPOINT(v);
}

void find_cutpoints() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    low.assign(n, -1);
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs (i);
    }
}

vector<int> id;
int curBCTNode;
vector<vector<int>> > tree;
vector<bool> isAP;

void buildTree() {
    curBCTNode = 0;
    id.assign(n, -1);
    tree.clear();

```

```

    isAP.clear();
    fore(v, 0, n) {
        if (cutpoint[v]) {
            id[v] = tree.size();
            tree.pb({});
            isAP.pb(true);
        }
    }
    for (auto comp : comps) {
        int v = tree.size();
        tree.pb({});
        isAP.pb(false);
        for (int x : comp) {
            if (cutpoint[x]) {
                tree[v].pb(id[x]);
                tree[id[x]].pb(v);
            }
            else {
                id[x] = v;
            }
        }
    }
}

```

../Graphs/Tarjan\_y\_BlockCutTree.cpp

### 4.3.1 Bridges

A bridge is an edge that if it is removed, the graph will be divided into two or more components. The implementation is:

```

int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph

vector<bool> visited;
vector<int> tin, low;
int timer;

void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    for (int to : adj[v]) {

```

```

        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else {
            dfs(to, v);
            low[v] = min(low[v], low[to]);
            if (low[to] > tin[v])
                IS_BRIDGE(v, to);
        }
    }
}

void find_bridges() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    low.assign(n, -1);
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs(i);
    }
}

```

../Graphs/FindBridges.cpp

### 4.3.2 Articulation Points

An articulation point is a node that if it is removed, the graph will be divided into two or more components. The implementation is:

```

int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph

vector<bool> visited;
vector<int> tin, low;
int timer;

void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    int children=0;
    for (int to : adj[v]) {

```

```

        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else {
            dfs(to, v);
            low[v] = min(low[v], low[to]);
            if (low[to] >= tin[v] && p != -1)
                IS_CUTPOINT(v);
            ++children;
        }
    }
    if (p == -1 && children > 1)
        IS_CUTPOINT(v);
}

void find_cutpoints() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    low.assign(n, -1);
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs(i);
    }
}

```

../Graphs/FindArticulationPoints.cpp

## 4.4 Flows

The flow is a concept that is used in many algorithms, it is used to find the maximum flow that could go through a system of nodes.

### 4.4.1 Dinic

The Dinic algorithm is a useful algorithm to find the maximum flow that could go through a system of nodes. The implementation of this algorithm is:

```

struct flowEdge
{
    int to, rev, f, cap;

```

```

};

vector<vector<flowEdge> > G;

void addEdge(int st, int en, int cap) {
    // Anade arista (st --> en) con su capacidad
    flowEdge A = {en, (int)G[en].size(), 0, cap};
    flowEdge B = {st, (int)G[st].size(), 0, 0};
    G[st].pb(A);
    G[en].pb(B);
}

int nodes, S, T; // asignar estos valores al armar el grafo G
                // nodes = nodos en red de flujo. Hacer G.clear();
    G.resize(nodes);
vi work, lvl;
int Q[200010];

bool bfs() {
    int qt = 0;
    Q[qt++] = S;
    lvl.assign(nodes, -1);
    lvl[S] = 0;
    for (int qh = 0; qh < qt; qh++) {
        int v = Q[qh];
        for (flowEdge &e : G[v]) {
            int u = e.to;
            if (e.cap <= e.f || lvl[u] != -1) continue;
            lvl[u] = lvl[v] + 1;
            Q[qt++] = u;
        }
    }
    return lvl[T] != -1;
}

int dfs(int v, int f) {
    if (v == T || f == 0) return f;
    for (int &i = work[v]; i < G[v].size(); i++) {
        flowEdge &e = G[v][i];
        int u = e.to;

```

```

        if (e.cap <= e.f || lvl[u] != lvl[v] + 1) continue;
        int df = dfs(u, min(f, e.cap - e.f));
        if (df) {
            e.f += df;
            G[u][e.rev].f -= df;
            return df;
        }
    }
    return 0;
}

int maxFlow() {
    int flow = 0;
    while (bfs()) {
        work.assign(nodes, 0);
        while (true) {
            int df = dfs(S, INF);
            if (df == 0) break;
            flow += df;
        }
    }
    return flow;
}

```

../Graphs/Dinic.cpp

This implementation is done in order to do the Dinic algorithm for a graph with a large number of nodes.

This algorithm is based on the idea of the BFS algorithm, it is used to find the shortest path between two nodes, in this case, the shortest path between the source and the sink. The algorithm is as follows:

#### 4.4.2 Ford Fulkerson

The Ford Fulkerson algorithm is a useful algorithm to find the maximum flow that could go through a system of nodes. The implementation of this algorithm is:

```

// This algorithm solves the max flow problem in a directed graph
// in O(max_flow * E)

// Here the graph is represented by an adjacency matrix, but it can
// be easily changed to an adjacency list

```

**Algorithm 4** Dinic

---

```

1: procedure DINIC( $G$ )
2:    $visited \leftarrow \emptyset$ 
3:    $time \leftarrow 0$ 
4:    $parent \leftarrow \emptyset$ 
5:    $low \leftarrow \emptyset$ 
6:    $disc \leftarrow \emptyset$ 
7:    $AP \leftarrow \emptyset$ 
8:    $bridge \leftarrow \emptyset$ 
9:   for all  $v \in V$  do
10:     $visited[v] \leftarrow false$ 
11:     $parent[v] \leftarrow -1$ 
12:     $low[v] \leftarrow \infty$ 
13:     $disc[v] \leftarrow \infty$ 
14:   end for
15:   for all  $v \in V$  do
16:     if  $visited[v] = false$  then
17:       DinicUtil( $G, v, visited, time, parent, low, disc, AP, bridge$ )
18:     end if
19:   end for
20: end procedure
21: procedure DINICUTIL( $G, v, visited, time, parent, low, disc, AP, bridge$ )
22:    $visited[v] \leftarrow true$ 
23:    $disc[v] \leftarrow time$ 
24:    $low[v] \leftarrow time$ 
25:    $time \leftarrow time + 1$ 
26:    $children \leftarrow 0$ 
27:   for all  $u \in Adj(v)$  do
28:     if  $visited[u] = false$  then
29:        $parent[u] \leftarrow v$ 
30:        $children \leftarrow children + 1$ 
31:       DinicUtil( $G, u, visited, time, parent, low, disc, AP, bridge$ )
32:        $low[v] \leftarrow \min(low[v], low[u])$ 
33:       if  $parent[v] = -1$  and  $children > 1$  then
34:          $AP[v] \leftarrow true$ 
35:       end if
36:       if  $parent[v] \neq -1$  and  $low[u] \geq disc[v]$  then
37:          $AP[v] \leftarrow true$ 
38:       end if
39:       if  $low[u] > disc[v]$  then
40:          $bridge[v][u] \leftarrow true$ 
41:       end if
42:     else
43:        $low[v] \leftarrow \min(low[v], disc[u])$ 
44:     end if
45:   end for

```

---

```

// The algorithm is based on the push-relabel algorithm, which is a
// variant of the relabel-to-front algorithm

// Number of vertices in given graph
#define V 6

/* Returns true if there is a path from source 's' to sink
't' in residual graph. Also fills parent[] to store the
path */
bool bfs(int rGraph[V][V], int s, int t, int parent[])
{
    // Create a visited array and mark all vertices as not visited
    bool visited[V];
    memset(visited, 0, sizeof(visited));

    // Create a queue, enqueue source vertex and mark source vertex
    // as visited
    queue<int> q;
    q.push(s);
    visited[s] = true;
    parent[s] = -1;

    // Standard BFS Loop
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        for (int v = 0; v < V; v++) {
            if (visited[v] == false && rGraph[u][v] > 0) {
                // If we find a connection to the sink node, then
                // there is no point in BFS anymore We just have to set its parent
                // and can return true
                if (v == t) {
                    parent[v] = u;
                    return true;
                }
                q.push(v);
                parent[v] = u;
                visited[v] = true;
            }
        }
    }
}

```



```

}
// We didn't reach sink in BFS starting from source, so return
false
return false;
}

// Returns the maximum flow from s to t in the given graph
int fordFulkerson(int graph[V][V], int s, int t)
{
    int u, v;
    // Create a residual graph and fill the residual graph with
    given capacities in the original graph as residual capacities
    in residual graph
    int rGraph[V][V]; // Residual graph where rGraph[i][j] indicates
    residual capacity of edge from i to j (if there is an edge. If
    rGraph[i][j] is 0, then there is not)
    for (u = 0; u < V; u++)
        for (v = 0; v < V; v++)
            rGraph[u][v] = graph[u][v];

    int parent[V]; // This array is filled by BFS and to store path

    int max_flow = 0; // There is no flow initially
    // Augment the flow while there is path from source to sink
    while (bfs(rGraph, s, t, parent)) {
        // Find minimum residual capacity of the edges along the
        path filled by BFS. Or we can say find the maximum flow through
        the path found.
        int path_flow = INT_MAX;
        for (v = t; v != s; v = parent[v]) {
            u = parent[v];
            path_flow = min(path_flow, rGraph[u][v]);
        }
        // update residual capacities of the edges and reverse
        edges along the path
        for (v = t; v != s; v = parent[v]) {
            u = parent[v];
            rGraph[u][v] -= path_flow;
            rGraph[v][u] += path_flow;
        }
    }
}

```

```

// Add path flow to overall flow
max_flow += path_flow;
}
// Return the overall flow
return max_flow;
}

int main()
{
    // Let us create a graph shown in the above example
    int graph[V][V]
        = { { 0, 16, 13, 0, 0, 0 }, { 0, 0, 10, 12, 0, 0 },
            { 0, 4, 0, 0, 14, 0 }, { 0, 0, 9, 0, 0, 20 },
            { 0, 0, 0, 7, 0, 4 }, { 0, 0, 0, 0, 0, 0 } };

    cout << "The maximum possible flow is \n"
          << fordFulkerson(graph, 0, 5);

    return 0;
}

```

../Graphs/FordFulkerson.cpp

In order to better understand the adjacency matrix in the code Figure 4.1 shows the graph that is used in the code.

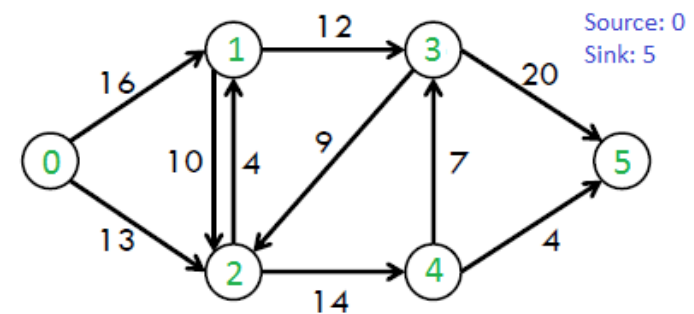


Figure 4.1: Ford Fulkerson

## 4.5 Dijkstra

The Dijkstra algorithm is a useful algorithm to find the shortest path between two nodes. The implementation of this algorithm is:

```
const int INF = 1e9;
vector<vector<pair<int, int>>> adj; //To store the node to which
    the edge flows to and the weight of the edge

void dijkstra(int s, vector<int> & d, vector<int> & p) {
    int n = adj.size();
    d.assign(n, INF);
    p.assign(n, -1);
    vector<bool> u(n, false);

    d[s] = 0;
    for (int i = 0; i < n; i++) {
        int v = -1;
        for (int j = 0; j < n; j++) {
            if (!u[j] && (v == -1 || d[j] < d[v]))
                v = j;
        }

        if (d[v] == INF)
            break;

        u[v] = true;
        for (auto edge : adj[v]) {
            int to = edge.first;
            int len = edge.second;

            if (d[v] + len < d[to]) {
                d[to] = d[v] + len;
                p[to] = v;
            }
        }
    }

    //In order to restore the path, we need to store the parent of each
    node in the shortest path tree
}
```

```
vector<int> restore_path(int s, int t, vector<int> const& p) {
    vector<int> path;

    for (int v = t; v != s; v = p[v]){
        if(v == -1){
            return {};
        }
        path.push_back(v);
    }
    path.push_back(s);

    reverse(path.begin(), path.end());
    return path;
}
```

../Graphs/Dijkstra.cpp

Another implementation of this algorithm is the one that is done using a priority queue, the implementation of this algorithm is:

```
// Dijkstra Modification thanks to pacha2880

vi dijkstra(int n, vector<vii> &g, int s, vi &par)
{
    vi dis(n, MOD), vis(n);
    dis[s] = 0;
    priority_queue<ii> que;
    que.push({0, s});
    while(!que.empty())
    {
        int node = que.top().s;
        que.pop();
        if(vis[node]) continue;
        vis[node] = 1;
        for(ii go : g[node])
            if(dis[go.f] > dis[node] + go.s)
            {
                dis[go.f] = dis[node] + go.s;
                par[go.f] = node;
                que.push({-dis[go.f], go.f});
            }
    }
}
```

```

    return dis;
}

vi path;
vi par(n+1, -1);
int t;
if(dis[t] == MOD){
    do {
        path.pb(t);
        t = par[t];
    } while(t != -1);
}

```

../Graphs/Dijkstra-Mod.cpp

## 4.6 Bellman Ford

The Bellman Ford algorithm is a useful algorithm to find the shortest path between two nodes. Bellman Ford differentiates from Dijkstra in the fact that it can be used in graphs with negative edges. The algorithm is as follows:

The implementation of this algorithm is:

```

struct edge
{
    int a, b, cost;
};

int n, m, v;
vector<edge> e;
const int INF = 1e9;

void solve()
{
    vector<int> d (n, INF);
    d[v] = 0;
    for (int i=0; i<n-1; ++i)
        for (int j=0; j<m; ++j)
            if (d[e[j].a] < INF) // is needed only if the graph
// contains negative weight edges: no such verification would
// result in relaxation from the vertices to which paths have not
// yet found, and incorrect distance

```

---

### Algorithm 5 Bellman Ford

---

```

1: procedure BELLMANFORD( $G$ )
2:    $dist \leftarrow \emptyset$ 
3:    $parent \leftarrow \emptyset$ 
4:   for all  $v \in V$  do
5:      $dist[v] \leftarrow \infty$ 
6:      $parent[v] \leftarrow -1$ 
7:   end for
8:    $dist[s] \leftarrow 0$ 
9:   for  $i = 0$  to  $|V| - 1$  do
10:    for all  $v \in V$  do
11:      for all  $u \in Adj(v)$  do
12:        if  $dist[u] > dist[v] + w(v, u)$  then
13:           $dist[u] \leftarrow dist[v] + w(v, u)$ 
14:           $parent[u] \leftarrow v$ 
15:        end if
16:      end for
17:    end for
18:  end for
19:  for all  $v \in V$  do
20:    for all  $u \in Adj(v)$  do
21:      if  $dist[u] > dist[v] + w(v, u)$  then
22:         $dist[u] \leftarrow -\infty$ 
23:         $parent[u] \leftarrow -1$ 
24:      end if
25:    end for
26:  end for
27: end procedure

```

---

```

        d[e[j].b] = min (d[e[j].b], d[e[j].a] + e[j].cost);
    // display d, for example, on the screen
}

// An improvement in the time of the algorithm could be introduced
// by keeping the flag so that no time is wasted in visiting all
// edges.

void solve()
{
    vector<int> d (n, INF);
    d[v] = 0;
    for (;;) // equivalent to while (true)
    {
        bool any = false;

        for (int j=0; j<m; ++j)
            if (d[e[j].a] < INF)
                if (d[e[j].b] > d[e[j].a] + e[j].cost)
                {
                    d[e[j].b] = d[e[j].a] + e[j].cost;
                    any = true;
                }

        if (!any) break;
    }
    // display d, for example, on the screen
}

//To retrieve the path of thr Bellman Ford algorithm, you need to
//keep the previous vertex in the path, and then go back from the
//end to the beginning.

void solve()
{
    vector<int> d (n, INF);
    d[v] = 0;
    vector<int> p (n, -1);

    for (;;)

```

```

{
    bool any = false;
    for (int j = 0; j < m; ++j)
        if (d[e[j].a] < INF)
            if (d[e[j].b] > d[e[j].a] + e[j].cost)
            {
                d[e[j].b] = d[e[j].a] + e[j].cost;
                p[e[j].b] = e[j].a;
                any = true;
            }
    if (!any) break;
}

if (d[t] == INF)
    cout << "No path from " << v << " to " << t << ".";
else
{
    vector<int> path;
    for (int cur = t; cur != -1; cur = p[cur])
        path.push_back (cur);
    reverse (path.begin(), path.end());

    cout << "Path from " << v << " to " << t << ":";
    for (size_t i=0; i<path.size(); ++i)
        cout << path[i] << ' ';
}

}

//Negative cycle detection

void solve()
{
    vector<int> d (n, INF);
    d[v] = 0;
    vector<int> p (n, -1);
    int x;
    for (int i=0; i<n; ++i)
    {
        x = -1;
        for (int j=0; j<m; ++j)

```

```

        if (d[e[j].a] < INF)
            if (d[e[j].b] > d[e[j].a] + e[j].cost)
            {
                d[e[j].b] = max (-INF, d[e[j].a] + e[j].cost);
                p[e[j].b] = e[j].a;
                x = e[j].b;
            }
    }

    if (x == -1)
        cout << "No negative cycle from " << v;
    else
    {
        int y = x;
        for (int i=0; i<n; ++i)
            y = p[y];

        vector<int> path;
        for (int cur=y; ; cur=p[cur])
        {
            path.push_back (cur);
            if (cur == y && path.size() > 1)
                break;
        }
        reverse (path.begin(), path.end());

        cout << "Negative cycle: ";
        for (size_t i=0; i<path.size(); ++i)
            cout << path[i] << ' ';
    }
}

```

// The SPFA (Shortest Path Faster Algorithm) is an improvement of Bellman Ford which takes advantage of the fact that not all attempts at relaxation will work

```

const int INF = 1e9;
vector<vector<pair<int, int>>> adj;

bool spfa(int s, vector<int>& d) {

```

```

    int n = adj.size();
    d.assign(n, INF);
    vector<int> cnt(n, 0);
    vector<bool> inqueue(n, false);
    queue<int> q;

    d[s] = 0;
    q.push(s);
    inqueue[s] = true;
    while (!q.empty()) {
        int v = q.front();
        q.pop();
        inqueue[v] = false;

        for (auto edge : adj[v]) {
            int to = edge.first;
            int len = edge.second;

            if (d[v] + len < d[to]) {
                d[to] = d[v] + len;
                if (!inqueue[to]) {
                    q.push(to);
                    inqueue[to] = true;
                    cnt[to]++;
                    if (cnt[to] > n)
                        return false; // negative cycle
                }
            }
        }
    }
    return true;
}

```

../Graphs/BellmanFord.cpp



## Chapter 5

# Dynamic Programming

This chapter shows some useful algorithms and implementations required to solve problems that require Dynamic Programming.

Some of the algorithms and implementations are as follows:

```
signed main()
{
    ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    // freopen("asd.txt", "r", stdin);
    // freopen("qwe.txt", "w", stdout);

    int n;
    cin>>n;
    int dp[n+1];
    dp[0]=1;
    //dice combinations
    fore(i,1,n+1){
        dp[i]=0;
        fore(j,1,7){
            if(i-j>=0){
                dp[i]+=dp[i-j];
                dp[i]%=MOD;
            }
        }
    }
    cout << dp[n] << '\n';

    return 0;
}
```

../DP/dp1.cpp

```
signed main()
{
    //La cantidad minimas de monedas para llegar a k o unbounded
    knapsack problem.
    ios::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    // freopen("asd.txt", "r", stdin);
    // freopen("qwe.txt", "w", stdout);

    int n,k;
    cin>>n>>k;
    vector<int> dp(k+1,1e9);
    int a[n];
    for(int i=0;i<n;i++)cin>>a[i];
    dp[0]=0;
    for(int i=1;i<=k;i++){
        for(int j=0;j<n;j++){
            if(i-a[j]>=0){
                dp[i]=min(dp[i],dp[i-a[j]]+1);
            }
        }
    }
    if(dp[k]==1e9)cout<<-1;
    else cout<<dp[k];
}
```

```
    return 0;
}
```

../DP/dp2.cpp

```
11 dp[1000001];

const int MOD = (int) 1e9 + 7;

int main(){
    int n, x; cin >> n >> x;
    vi coins(n);
    for (int i = 0; i < n; i++) {
        cin >> coins[i];
    }
    dp[0] = 1;
    for (int weight = 0; weight <= x; weight++) {
        for (int i = 1; i <= n; i++) {
            if(weight - coins[i - 1] >= 0) {
                dp[weight] += dp[weight - coins[i - 1]];
                dp[weight] %= MOD;
            }
        }
    }
    cout << dp[x] << '\n';
}
```

../DP/dp3.cpp

```
//combination de monedas en orden
11 dp[1000001];

const int MOD = (int) 1e9 + 7;

int main(){
    int n, x; cin >> n >> x;
    vi coins(n);
    for (int i = 0; i < n; i++) {
        cin >> coins[i];
    }
    dp[0] = 1;
    for (int i = 1; i <= n; i++) {
```

```
        for (int weight = 0; weight <= x; weight++) {
            if(weight - coins[i - 1] >= 0) { // prevenir casos bound
                dp[weight] += dp[weight - coins[i - 1]];
                dp[weight] %= MOD;
            }
        }
    }
    cout << dp[x] << '\n';
}
```

../DP/dp4.cpp

```
#include <bits/stdc++.h>
using namespace std;
//You are given an integer n On each step, you may subtract one of
the digits from the number.How many steps are required to make
the number equal to 0
int main() {
    int n;
    cin >> n;
    vector<int> dp(n+1,1e9);
    dp[0] = 0;
    for (int i = 0; i <= n; i++) {
        for (char c : to_string(i)) {
            dp[i] = min(dp[i], dp[i-(c-'0')]+1);
        }
    }
    cout << dp[n] << endl;
}
```

../DP/dp5.cpp

```
#include <bits/stdc++.h>
using namespace std;

int main() {
    // there is one way to reach (0,0), dp[0][0] = 1.
    //Consider an n*n grid whose squares may have traps. It is not
    allowed to move to a square with a trap.Your task is to
    calculate the number of paths from the upper-left square to the
    lower-right square. You can only move right or down
    int mod = 1e9+7;
```



```

int n;
cin >> n;
vector<vector<int>> dp(n, vector<int>(n, 0));
dp[0][0] = 1;
for (int i = 0; i < n; i++) {
    string row;
    cin >> row;
    for (int j = 0; j < n; j++) {
        if (row[j] == '.') {
            if (i > 0) {
                (dp[i][j] += dp[i-1][j]) %= mod;
            }
            if (j > 0) {
                (dp[i][j] += dp[i][j-1]) %= mod;
            }
            else {
                dp[i][j] = 0;
            }
        }
    }
    cout << dp[n-1][n-1] << endl;
}

```

../DP/dp6.cpp

You are in a book shop which sells  $n$  different books. You know the price and number of pages of each book.

You have decided that the total price of your purchases will be at most  $x$ .

What is the maximum number of pages you can buy? You can buy each book at most once. #include <bits/stdc++.h>

```

using namespace std;

int main() {
    int n, x;
    cin >> n >> x;
    vector<int> price(n), pages(n);
    for (int i; i < n; i++) cin >> price[i];
    for (int i; i < n; i++) cin >> pages[i];
    vector<vector<int>> dp(n+1, vector<int>(x+1, 0));

```

```

for (int i = 1; i <= n; i++) {
    for (int j = 0; j <= x; j++) {
        dp[i][j] = dp[i-1][j];
        int left = j - price[i-1];
        if (left >= 0) {
            dp[i][j] = max(dp[i][j], dp[i-1][left] + pages[i-1]);
        }
    }
}
cout << dp[n][x] << endl;
}

```

../DP/dp7.cpp

//You know that an array has  $n$  integers between 1 and  $m$ , and the absolute difference between two adjacent values is at most 1.

//Given a description of the array where some values may be unknown, your task is to count the number of arrays that match the description.

```

#include <bits/stdc++.h>
using namespace std;

int main() {
    int mod = 1e9+7;
    int n, m;
    cin >> n >> m;
    vector<vector<int>> dp(n, vector<int>(m+1, 0));
    int x0;
    cin >> x0;
    if (x0 == 0) {
        fill(dp[0].begin(), dp[0].end(), 1);
    }
    //igual memset pero algo raro
    } else {
        dp[0][x0] = 1;
    }
    for (int i = 1; i < n; i++) {
        int x;
        cin >> x;
        if (x == 0) {
            for (int j = 1; j <= m; j++) {
                for (int k : {j-1, j, j+1}) {
                    if (k >= 1 && k <= m) {

```

```

        (dp[i][j] += dp[i-1][k]) %= mod;
    }
}
}
} else {
    for (int k : {x-1,x,x+1}) {
        if (k >= 1 && k <= m) {
            (dp[i][x] += dp[i-1][k]) %= mod;
        }
    }
}
}
int ans = 0;
for (int j = 1; j <= m; j++) {
    (ans += dp[n-1][j]) %= mod;
}
cout << ans << endl;
}

```

../DP/dp8.cpp

## 5.1 Knapsack Problem

The knapsack problem is a problem that consists of finding the maximum value of a set of items that can be placed in a knapsack of a given weight. The problem can be solved using Dynamic Programming.

The implementation can be done as follows:

```

// A Dynamic Programming based solution for 0-1 Knapsack problem
#include <iostream>

using namespace std;

// A utility function that returns maximum of two integers
int max(int a, int b)
{
    return (a > b) ? a : b;
}

// Returns the maximum value that can be put in a knapsack of
// capacity W

```

```

int knapSack(int W, int wt[], int val[], int n)
{
    int i, w;
    int K[n + 1][W + 1];

    // Build table K[][] in bottom up manner
    for (i = 0; i <= n; i++)
    {
        for (w = 0; w <= W; w++)
        {
            if (i == 0 || w == 0)
                K[i][w] = 0;
            else if (wt[i - 1] <= w)
                K[i][w] = max(val[i - 1] + K[i - 1][w - wt[i - 1]],
                               K[i - 1][w]);
            else
                K[i][w] = K[i - 1][w];
        }
    }

    return K[n][W];
}

int main()
{
    cout << "Enter the number of items in a Knapsack:";
    int n, W;
    cin >> n;
    int val[n], wt[n];
    for (int i = 0; i < n; i++)
    {
        cout << "Enter value and weight for item " << i << ":";
        cin >> val[i];
        cin >> wt[i];
    }

    // int val[] = { 60, 100, 120 };
    // int wt[] = { 10, 20, 30 };
    // int W = 50;
    cout << "Enter the capacity of knapsack";
}

```

```

cin >> W;
cout << knapSack(W, wt, val, n);

return 0;
}

```

../DP/knapsack.cpp

## 5.2 Divide and Conquer

The divide and conquer algorithm is a recursive algorithm that divides the problem into smaller subproblems and solves them recursively. The algorithm is as follows:

---

### Algorithm 6 Divide and Conquer

---

```

1: procedure DIVIDEANDCONQUER( $A$ )
2:   if  $A$  has only one element then
3:     return  $A$ 
4:   end if
5:    $B \leftarrow \text{DivideAndConquer}(A[0..n/2])$ 
6:    $C \leftarrow \text{DivideAndConquer}(A[n/2+1..n])$ 
7:   return  $\text{Merge}(B, C)$ 
8: end procedure

```

---

The implementation can be done as follows:

```

/*
DP[i][j] = min( DP[i-1][k] + C[k][j] )
K[i][j] <= K[i][j+1]
*/

ll lastDP[tam], DP[tam];
int C[tam][tam]; // Cambiar a una funcion de costo si pre-procesar
                  ocupa mucha memoria

void DC(int b, int e, int KL, int KR)
{
    int mid = (b + e) / 2;
    pair<ll, int> best = mp(-1, KL);

    for (int k = KL; k < min(mid, KR+1); k++)

```

```

{
    best = max( best, mp(lastDP[k] + C[k+1][mid], k) );
}

DP[mid] = best.first;
int K = best.second;

if (b <= mid-1)
    DC(b, mid-1, KL, K);
if (mid+1 <= e)
    DC(mid+1, e, K, KR);
}

```

../DP/DivideAndConquer.cpp



# Chapter 6

## Geometry

This chapter is a survey of the main results and algorithms useful to solve geometry problems.

### 6.1 Points and Lines

Some of the most useful functions used while describing points and lines are:

- `atan2(y,x)`: returns the angle between the positive x-axis and the vector (x,y).
- `hypot(x,y)`: returns the Euclidean distance between the origin and the point (x,y).
- `cross(a,b)`: returns the cross product of the vectors a and b.
- `dot(a,b)`: returns the dot product of the vectors a and b.
- `dist(a,b)`: returns the Euclidean distance between the points a and b.
- `dist2(a,b)`: returns the squared Euclidean distance between the points a and b.
- `ccw(a,b,c)`: returns true if the points a, b and c are in counterclockwise order.
- `collinear(a,b,c)`: returns true if the points a, b and c are collinear.
- `angle(a,b)`: returns the angle between the vectors a and b.
- `angle(a,b,c)`: returns the angle between the vectors a-b and c-b.
- `rotate(a,ang)`: returns the vector a rotated by ang radians.

- `rotate(a,ang,center)`: returns the vector a rotated by ang radians around the point center.
- `reflect(a,m)`: returns the reflection of the point a across the line m.
- `project(a,m)`: returns the projection of the point a onto the line m.
- `closest(a,m)`: returns the closest point on the line m to the point a.
- `intersect(a,b,c,d)`: returns true if the lines a-b and c-d intersect.

```
#include <bits/stdc++.h>
#define EPS 1e-9
struct line {double a,b,c;}; // ax + by + c = 0
bool areParallel(line a, line b)
{
    return((fabs(a.a-b.a)<EPS)&&(fabs(a.b-b.b)<EPS));
}
bool areSame(line a, line b)
{
    return areParallel(a,b)&&(fabs(a.c-b.c)<EPS);
}
struct point
{
    double x,y;
    point() {x=y=0;}
    point(double _x, double _y) : x(_x), y(_y) {}
    point operator+(point a) const
    {
```

```

    a.x+=x;
    a.y+=y;
    return a;
}
};
double dist(point a, point b)
{
    return hypot(a.x-b.x,a.y-b.y);
}
void toline(point a, point b, line &l) //dados dos puntos
{
    if(fabs(a.x-b.x)<EPS)
        {l.a = 1, l.b = 0, l.c = -a.x; return;}
    l.a = -(a.y - b.y) / (a.x - b.x);
    l.b = 1;
    l.c = -l.a * a.x - a.y;
}
void tolinegr(point a, double gr, line &l) // a linea dado el
    gradiente
{
    l.a = -gr;
    l.b = 1;
    l.c = a.x * gr - a.y;
}
point tovec(point a, point b)
{
    return point(b.x-a.x,b.y-a.y);
}
point translate (point p, point v)
{
    return point(p.x+v.x,p.y+v.y);
}
point scale(point v, double sc)
{
    return point(v.x*sc,v.y*sc);
}
point rotate(point v, double theta) //rotacion antihorario ccw
{
    theta *= acos(-1)/180.0;
    return point(v.x*cos(theta)-v.y*sin(theta),v.x*sin(theta)+ v.y*
        cos(theta));
}

```

```

}
bool areIntersect(line l1, line l2, point &p) //interseccion de
    lineas
{
    if(areParallel(l1,l2)) return false;
    p.x = (-l1.c*l2.b + l2.c*l1.b) / (l1.a*l2.b-l2.a*l1.b);
    if(fabs(l1.b) > EPS) p.y = -(l1.a*p.x + l1.c);
    else
        p.y = -(l2.a*p.x +l2.c);
    return true;
}
point clos(point a, line l, line &pe) //closest point in a line and
    perpendicular line from a
{
    if(fabs(l.a) < EPS)
    {
        pe.a = 1, pe.b = 0, pe.c = -a.x;
        return point(a.x,-l.c);
    }
    if(fabs(l.b) < EPS)
    {
        pe.a = 0, pe.b = 1, pe.c = -a.y;
        return point(-l.c,a.y);
    }
    tolinegr(a, 1/(l.a),pe);
    areIntersect(l,pe,a);
    return a;
}
point reflexion(point p, point a, point b) // del punto p a linea
    ab
{
    line l,li;
    toline(a,b,li);
    point p1 = clos(p,li,l);
    p1 = p1+(tovec(p,p1));
    return p1;
}
double norm_sq(point a)
{
    return a.x * a.x + a.y * a.y;
}
}

```

```

double dot(point a, point b)
{
    return a.x*b.x + b.y*a.y;
}
double angle(point a, point b, point c) //b el del medio
{
    a = tovec(b,a), b = tovec(b,c);
    double res = dot(a,b);
    res = acos(res / (sqrt(norm_sq(a))*sqrt(norm_sq(b))));
    res*= 180.0/acos(-1);
    return res;
}
double cross(point a, point b) //producto cruz
{
    return a.x * b.y - a.y * b.x;
}
bool left(point a, point b, point c) //ccw
{
    c = tovec(b,c);
    a = tovec(b,a);
    return (cross(c,a)>0.0);
}
bool colinear(point a, point b, point c)
{
    c = tovec(b,c);
    a = tovec(b,a);
    return (fabs(cross(c,a))<EPS);
}
double distToLine(point p, point a, point b, point &c) //con
    producto punto halla el punto
{
    //c = a + u*ab
    point ab = tovec(a,b), ap = tovec(a,p);
    double u = dot(ab,ap) / norm_sq(ab);
    c = translate(a, scale(ab,u));
    return dist(p,c);
}
double distToline1(point p, point a, point b) //con producto cruz
    solo distancia
{

```

```

    point ap = tovec(a,p), ab = tovec(a,b);
    return fabs(cross(ab,ap)/hypot(ab.x,ab.y));
}

```

../Geometry/PointsAndLines.cpp

### 6.1.1 Lines

In order to represent lines and find their intersection, we can use the following struct:

```

struct line
{
    double a, b, c;
    line(point p, point q)
    {
        a = p.y - q.y;
        b = q.x - p.x;
        c = -a * p.x - b * p.y;
    };
    void setOrigin(point p) { c += a * p.x + b * p.y; } //trasladar
        linea como si p fuera el origen
};

double det(double a, double b, double c, double d)
{
    return a * d - b * c;
}

point intersec(line a, line b) //primero estar seguro si no son
    paralelas
{
    double d = -det(a.a, a.b, b.a, b.b);
    return point(det(a.c, a.b, b.c, b.b) / d, det(a.a, a.c, b.a, b.c)
        / d);
}

```

../Geometry/Line.cpp