

Name: _____

Date: _____

Pledge: _____

Consider the implementation of quickselect below:

```

int lomuto_partition(int array[], int left, int right) {
    int p = array[left], s = left;
    for (int i = left + 1; i <= right; ++i) {
        if (array[i] < p) {
            ++s;
            swap(array[s], array[i]);
        }
    }
    swap(array[left], array[s]);
    return s;
}

int quick_select(int array[], size_t left, size_t right, size_t k) {
    size_t s = lomuto_partition(array, left, right);

    if (s == k - 1) {
        return ____ (a) ____;
    }
    if (s > k - 1) {
        return quick_select(array, left, ____ (b) ____, k);
    }
    return quick_select(array, ____ (c) ____, right, ____ (d) ____);
}

int quick_select(int array[], const size_t length, size_t k) {
    return quick_select(array, 0, length - 1, k);
}

```

1) Fill in the 4 blanks in the quickselect algorithm. (4 points)

- a) _____
 b) _____
 c) _____
 d) _____

2) Show the array [4, 5, 6, 4, 0] after running `lomuto_partition`. (3 points) _____3) Suppose we are sorting an array of eight integers using `quicksort` with `lomuto_partition` and have just finished the first call to `lomuto_partition`. The array now looks as follows:

11 4 20 45 32 60 98 70

Which value or values could have been the pivot? (1 point) _____

4) Suppose mergesort were to cut the array into 3 evenly sized subarrays (instead of 2) and did a 3-way merge after making the recursive calls.

a) Write the recurrence relation for this modified version of mergesort. (1 point)

 $T(n) = \underline{\hspace{2cm}}$

b) Use the Master Theorem to determine its complexity. (1 point, answer depends on correct 4a, indicate base for logarithm)

 $T(n) \in \theta(\underline{\hspace{2cm}})$