

1. Find an upper bound for $f(n) = n^4 + 10n^2 + 5$. Write your answer here: $n^4 + 10n^2 + 5 \leq c(n^4)$ for $c=4, n=2$ (4 points)

Prove your answer by giving values for c and n possible for c . (4 points)

$$n^4 + 10n^2 + 5 \leq c(n^4)$$

$$\text{When } c = 1, n = 1 \rightarrow 16 \leq 1 \text{ NO}$$

$$\text{When } c = 2, n = 2 \rightarrow 61 \leq 32 \text{ NO}$$

...

$$\text{When } c = 4, n = 2 \rightarrow 16 \leq 32 \text{ YES}$$

The smallest value of c to work is $c = 4$, and $n = 2$

2. Find an asymptotically tight bound for $f(n) = 3n^3 - 2n$. Write answer here: $C_2(n^3) \leq 3n^3 - 2n \leq C_1(n^3)$ (4 points)

Prove your answer by giving values for the constants c_1 , c_2 , and n_0 . Choose the tightest integral values possible for c_1 and c_2 . (6 points)

$$n_0 \geq 1, C_1 = 3, C_2 = 1$$

3. Is $3n - 4 \in \Omega(n^2)$? Circle your answer: yes / NO. (2 points)

If yes, prove your answer by giving values for the constants c and n_0 . Choose the smallest integral value possible for c . If no, derive a contradiction. (4 points)

$$\text{When } f(n) = n \dots n \notin \Omega(n^2) \rightarrow n \not\geq n^2$$

4. Write the following asymptotic efficiency classes in **increasing** order of magnitude. $O(n^2)$, $O(2^n)$, $O(1)$, $O(n \lg n)$, $O(n)$, $O(n!)$, $O(n^3)$, $O(\lg n)$, $O(n^n)$, $O(n^2 \lg n)$ (2 points each)

$$O(1), O(\lg n), O(n), O(n \lg n), O(n^2 \lg n), O(n^2), O(n^3), O(n!), O(n^n), O(2^n)$$

5. Determine the largest size n of a problem that can be solved in time t , assuming that the algorithm takes $f(n)$ milliseconds. n must be an integer. (2 points each)

a. $f(n) = n$, $t = 1$ second 10^3

b. $f(n) = n \lg n$, $t = 1$ hour 2.04×10^5

c. $f(n) = n^2, t = 1 \text{ hour}$ $1.8974 * 10^3$

d. $f(n) = n^3, t = 1 \text{ day}$ $4.421 * 10^2$

e. $f(n) = n!, t = 1 \text{ minute}$ $8!$

6. Suppose we are comparing two sorting algorithms and that for all inputs of size n the first algorithm runs in $4n^3$ seconds, while the second algorithm runs in $64n \lg n$ seconds. For which integral values of n does the first algorithm beat the second algorithm? $n = 2, 3, 4$ (4 points)

Explain how you got your answer or paste code that solves the problem (2 point):

$$4n^3 < 64 \lg n \quad \rightarrow \quad n^2/16 < \lg n$$

When...

$n = 1$	$\dots 1/16 < 0$	$\dots \text{NO}$
$n = 2$	$\dots 0.25 < 0$	$\dots \text{YES}$
$n = 3$	$\dots 0.5625 < 1.584$	$\dots \text{YES}$
$n = 4$	$\dots 1 < 2$	$\dots \text{YES}$
$n = 5$	$\dots 1.5625 < 0.698$	$\dots \text{NO}$

7. Give the complexity of the following methods. Choose the most appropriate notation from among O , Θ , and Ω . (8 points each)

```
int function1(int n) {
    int count = 0;
    for (int i = n / 2; i <= n; i++) {
        for (int j = 1; j <= n; j *= 2) {
            count++;
        }
    }
    return count;
}
```

Answer: $O(n \log n)$

```
int function2(int n) {
    int count = 0;
    for (int i = 1; i * i * i <= n; i++) {
        count++;
    }
    return count;
}
```

Answer: $\Theta(\sqrt[3]{n})$

```
int function3(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            for (int k = 1; k <= n; k++) {
```

```

        count++;
    }
}
return count;
}

```

Answer: $\Theta(n^3)$

```

int function4(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            count++;
            break;
        }
    }
    return count;
}

```

Answer: $\Theta(n)$

```

int function5(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        count++;
    }
    for (int j = 1; j <= n; j++) {
        count++;
    }
    return count;
}

```

Answer: $O(n)$