

1. Find an upper bound for  $f(n) = n^4 + 10n^2 + 5$ . Write your answer here:  $n^4 + 10n^2 + 5 \leq c(n^4)$  for  $c=4, n=2$  (4 points)

Prove your answer by giving values for  $c$  and  $n$  possible for  $c$ . (4 points)

$$n^4 + 10n^2 + 5 \leq c(n^4)$$

$$\text{When } c = 1, n = 1 \rightarrow 16 \leq 1 \text{ NO}$$

$$\text{When } c = 2, n = 2 \rightarrow 61 \leq 32 \text{ NO}$$

...

$$\text{When } c = 4, n = 2 \rightarrow 16 \leq 32 \text{ YES}$$

The smallest value of  $c$  to work is  $c = 4$ , and  $n = 2$

2. Find an asymptotically tight bound for  $f(n) = 3n^3 - 2n$ . Write answer here:  $C_2(n^2) \leq 2n^2 - n \leq C_1$  (4 points)

Prove your answer by giving values for the constants  $c_1$ ,  $c_2$ , and  $n_0$ . Choose the tightest integral values possible for  $c_1$  and  $c_2$ . (6 points)

$$n_0 \geq 1, C_1 = 2, C_2 = 1$$

3. Is  $3n - 4 \in \Omega(n^2)$ ? Circle your answer: yes / **NO**. (2 points)

If yes, prove your answer by giving values for the constants  $c$  and  $n_0$ . Choose the smallest integral value possible for  $c$ . If no, derive a contradiction. (4 points)

$$\text{When } f(n) = n \dots n \in \Omega(n^2) \rightarrow n \geq n^2$$

4. Write the following asymptotic efficiency classes in **increasing** order of magnitude.  $O(n^2)$ ,  $O(2^n)$ ,  $O(1)$ ,  $O(n \lg n)$ ,  $O(n)$ ,  $O(n!)$ ,  $O(n^3)$ ,  $O(\lg n)$ ,  $O(n^n)$ ,  $O(n^2 \lg n)$  (2 points each)

$$O(1), O(\lg n), O(n), O(n \lg n), O(n^2 \lg n), O(n^2), O(n^3), O(n!), O(n^n), O(2^n)$$

5. Determine the largest size  $n$  of a problem that can be solved in time  $t$ , assuming that the algorithm takes  $f(n)$  milliseconds.  $n$  must be an integer. (2 points each)

a.  $f(n) = n$ ,  $t = 1$  second  $10^3$

b.  $f(n) = n \lg n$ ,  $t = 1$  hour  $2.04 * 10^5$

c.  $f(n) = n^2$ ,  $t = 1$  hour  $1.8974 * 10^3$

d.  $f(n) = n^3$ ,  $t = 1$  day  $4.421 * 10^2$

e.  $f(n) = n!$ ,  $t = 1$  minute  $8!$

6. Suppose we are comparing two sorting algorithms and that for all inputs of size  $n$  the first algorithm runs in  $4n^3$  seconds, while the second algorithm runs in  $64n \lg n$  seconds. For which integral values of  $n$  does the first algorithm beat the second algorithm?  $n = 2, 3, 4$  (4 points)

Explain how you got your answer or paste code that solves the problem (2 point):

$$4n^3 < 64 \lg n \rightarrow n^2/16 < \lg n$$

When...

$n = 1$	$\dots 1/16 < 0$	$\dots$ NO
$n = 2$	$\dots 0.25 < 0$	$\dots$ YES
$n = 3$	$\dots 0.5625 < 1.584$	$\dots$ YES
$n = 4$	$\dots 1 < 2$	$\dots$ YES
$n = 5$	$\dots 1.5625 < 0.698$	$\dots$ NO

7. Give the complexity of the following methods. Choose the most appropriate notation from among  $O$ ,  $\Theta$ , and  $\Omega$ . (8 points each)

```
int function1(int n) {
    int count = 0;
    for (int i = n / 2; i <= n; i++) {
        for (int j = 1; j <= n; j *= 2) {
            count++;
        }
    }
    return count;
}
```

Answer:  $O(n \log n)$

```
int function2(int n) {
    int count = 0;
    for (int i = 1; i * i * i <= n; i++) {
        count++;
    }
    return count;
}
```

Answer:  $\Theta(\sqrt[3]{n})$

```
int function3(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            for (int k = 1; k <= n; k++) {
```

```

        count++;
    }
}
return count;
}

```

Answer:  $\Theta(n^3)$

```

int function4(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            count++;
            break;
        }
    }
    return count;
}

```

Answer:  $\Theta(n)$

```

int function5(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        count++;
    }
    for (int j = 1; j <= n; j++) {
        count++;
    }
    return count;
}

```

Answer:  $O(n)$