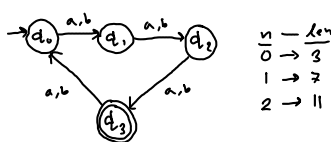
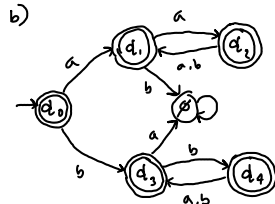


Set A

1. a)  $4n+3$



0 0 1 1  
1 1 2 2  
2 2 3 3  
3 3 4 4



Problem 1 (CO1): DFA and Regular Languages (15 points)

Let  $\Sigma = \{a, b\}$ . Consider the following languages over  $\Sigma$ .

- $L_1 = \{w : \text{length of } w \text{ is three more than multiple of four}\}$
- $L_2 = \{w : \text{every even position letter in } w \text{ is the same as the first letter of } w\}$
- $L_3 = \{w : \text{every } 2k+1 \text{ position in } w \text{ is } a, \text{ where } k \geq 0\}$
- $L_4 = \{w : \text{every } 2k+1 \text{ position in } w \text{ is } b, \text{ where } k \geq 0\}$

Now solve the following problems.

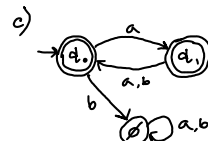
- Give the state diagram for a DFA that recognizes  $L_1$ . (3 points)
- Give the state diagram for a DFA that recognizes  $L_2$ . (3 points)
- Give the state diagram for a DFA that recognizes  $L_3$ . (3 points)
- If you were to use the "cross product" construction to obtain a DFA for the language  $L_2 \cap (L_3 \cup L_4)$ , how many states would it have? (1 point)
- Find all four-letter strings in  $L_2 \cap (L_3 \cup L_4)$ . (1 point)
- Give the state diagram for a DFA that recognizes  $L_2 \cap (L_3 \cup L_4)$  using only four states. (2 points)
- Find a four-letter string in  $L_3 \cup L_4$ . [Recall:  $\bar{L}$  denotes the complement of the language  $L$  i.e.,  $\bar{L} = \Sigma^* - L$ ] (1 point)
- Is  $L_3 = L_4 = \bar{L}_3$ ? Give justification for your answer. (1 point)

g) babb

h) No. Suppose,

$$\bar{L}_3 = b, L_4 = ba$$

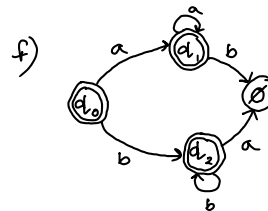
$\therefore \bar{L}_3 \cdot L_4 = bba$ , this is not an accepted str in  $\bar{L}_3$



d)  $L_4$  states  $\rightarrow 3$

$$\text{total State} = 6 \times 3 \times 3 = 54$$

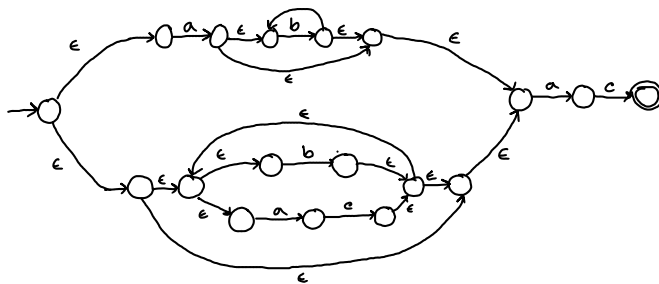
e) aaaaa, bbbbbb,



Problem 2 (CO2): Converting Regular Expressions to NFA (10 points)

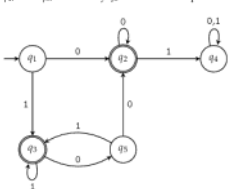
Convert the following regular expression over  $\Sigma = \{a, b, c\}$  into an equivalent NFA. Note that  $R_1 + R_2$  is the same as  $R_1 \cup R_2$ .

$$(ab^* + (b+ac)^*)bc$$

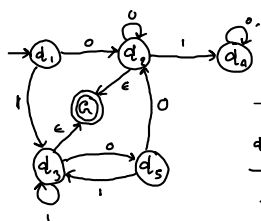


Problem 3 (CO2): Converting Finite Automata to Regular Expressions (10 points)

Convert the following finite automata into an equivalent regular expression using the state elimination method. You must eliminate  $q_2$  first, then  $q_4$ , and finally  $q_3$ . Show each step of the process.

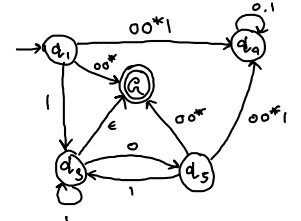


$$q_2 \rightarrow q_1 \rightarrow q_3 \rightarrow q_4$$

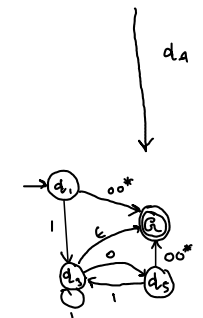


$$d_2$$

$$\begin{array}{ll} d_1 \text{ to } q & d_3 \text{ to } q \\ \rightarrow d_1 d_2 q & \rightarrow d_3 d_2 q \\ - 00^* & - 00^* \\ d_1 \text{ to } d_4 & d_5 \text{ to } d_4 \\ \rightarrow d_1 d_4 d_4 & \rightarrow d_5 d_4 d_4 \\ - 00^*1 & - 00^*1 \end{array}$$

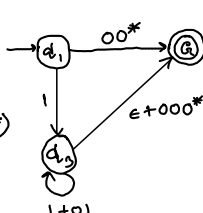


$d_4$  trap state



$$d_5$$

$$\begin{array}{ll} d_3 \text{ to } q & \\ \rightarrow d_3 d_5 q & - 000^* \\ d_5 \text{ to } d_3 & \\ \rightarrow d_3 d_5 d_3 & - 01 \end{array}$$



$$\begin{array}{ll} d_3 & \\ \rightarrow d_1 & 00^* + 1(1+0)^*(\epsilon + 000^*) \\ d_1 \text{ to } q & \\ \rightarrow d_1 d_3 q & \\ - 1(1+0)^*(\epsilon + 000^*) & \end{array}$$

ANS:  $00^* + 1(1+0)^*(\epsilon + 000^*)$

# Problem 4 (CO1): Regular Expressions (10 points)

Let  $\Sigma = \{0, 1\}$ . Consider the following languages over  $\Sigma$ .

$$\begin{aligned} L_1 &= \{w : \text{length of } w \text{ is exactly } 4\} \\ L_2 &= \{w : \text{the third last digit of } w \text{ is } 0\} \\ L_3 &= \{w : w \text{ contains at most two } 1\text{'s}\} \\ L_4 &= L_1^c \cap L_2 \end{aligned}$$

Now solve the following problems.

- Give a regular expression for the language  $L_1$ . (1 point)
- Give a regular expression for the language  $L_1^c$ . (1 point)
- Give a regular expression for the language  $L_3^c$ . [Recall:  $\bar{L}$  denotes the complement of the language  $L$  i.e.,  $\bar{L} = \Sigma^* - L$ ] (2 points)
- Give a regular expression for the language  $L_2$ . (2 points)
- Give a regular expression for the language  $L_3$ . (2 points)
- Give a regular expression for the language  $L_4$ . (2 points)

- $\Sigma\Sigma\Sigma\Sigma$
- $(\Sigma\Sigma\Sigma\Sigma)^*$
- $(\Sigma\Sigma\Sigma\Sigma)^*\Sigma\Sigma\Sigma\Sigma$
- $\Sigma^*0\Sigma\Sigma$
- $\Sigma^*11\Sigma^*11\Sigma^*11\Sigma^*$
- $(\Sigma\Sigma\Sigma\Sigma)^*1\Sigma\Sigma\Sigma$

3rd last  $\rightarrow 1$   
not a mul. of 4

# Problem 5 (CO2): Subset Construction Method (5 points)

Consider the following NFA:

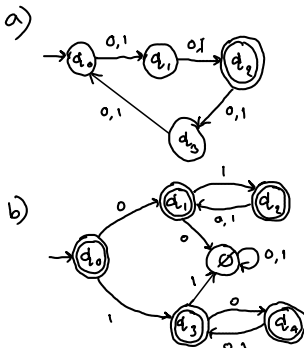


Now answer the following questions. [Note: You do not need to convert the given NFA into its equivalent DFA to answer the questions.]

- If you convert the given NFA into an equivalent DFA using the subset construction method, what is the maximum number of states that the DFA can have? (1 point)
- What is the maximum number of accepting states that the equivalent DFA can have? (1 point)
- Write the  $\epsilon$ -closure of state  $q_1$  in the given NFA. (1 point)
- Write the subset of states of the given NFA that will be the starting state in its equivalent DFA. (1 point)
- What is  $\delta(\{q_1, q_2\}, b)$  in the given NFA? [Recall:  $\delta(\{q\}, a)$  is the set of states in which the NFA transitions when it is in state  $q$  and receives input  $a$ .] (1 point)

- $2^5 = 32$
- reject states  $\rightarrow 4$   
 $\therefore$  accept states  $\rightarrow 32 - 4 = 28$
- $q_1 \xrightarrow{\epsilon} \{q_1, q_2\}$
- $q_1, q_2$
- $q_1, q_2 \xrightarrow{b} q_3, q_4$

## Set B



$$\begin{array}{r} 4n+2 \\ n \quad 4n \\ 0 \quad 2 \\ 1 \quad 6 \\ 2 \quad 8 \end{array}$$

- e) ✓ d) ✓ e) ✓ f) ✓ g) ✓ h) ✓

# Problem 1 (CO1): DFA and Regular Languages (15 points)

Let  $\Sigma = \{a, b\}$ . Consider the following languages over  $\Sigma$ .

$$\begin{aligned} L_1 &= \{w : \text{length of } w \text{ is two more than multiple of four}\} \\ L_2 &= \{w : \text{every even position letter in } w \text{ is different from the first letter of } w\} \\ L_3 &= \{w : \text{every } 2k+1 \text{ position in } w \text{ is } b, \text{ where } k \geq 0\} \\ L_4 &= \{w : \text{every } 2k+1 \text{ position in } w \text{ is } a, \text{ where } k \geq 0\} \end{aligned}$$

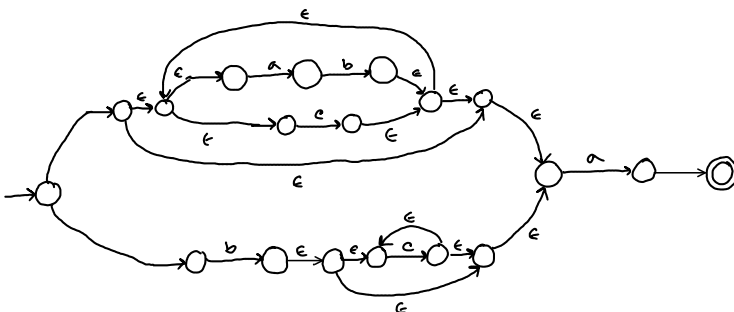
Now solve the following problems.

- Give the state diagram for a DFA that recognizes  $L_1$ . (3 points)
- Give the state diagram for a DFA that recognizes  $L_2$ . (3 points)
- Give the state diagram for a DFA that recognizes  $L_3$ . (3 points)
- If you were to use the "cross product" construction to obtain a DFA for the language  $L_2 \cap (L_3 \cup L_4)$ , how many states would it have? (1 point)
- Find all four-letter strings in  $L_2 \cap (L_3 \cup L_4)$ . (1 point)
- Give the state diagram for a DFA that recognizes  $L_2 \cap (L_3 \cup L_4)$  using only four states. (2 points)
- Find a four-letter string in  $\bar{L}_3 \cap L_4$ . [Recall:  $\bar{L}$  denotes the complement of the language  $L$  i.e.,  $\bar{L} = \Sigma^* - L$ ] (1 point)
- Is  $\bar{L}_3 \cap L_4 = \bar{L}_3$ ? Give justification for your answer. (1 point)

# Problem 2 (CO2): Converting Regular Expressions to NFA (10 points)

Convert the following regular expression over  $\Sigma = \{a, b, c\}$  into an equivalent NFA. Note that  $R_1 + R_2$  is the same as  $R_1 \cup R_2$ .

$$((ab + c)^+ + bc^*)ac$$



3. ✓ 4. ✓ 5. ✓