Operating Systems Synchronization

An Introduction

Mindset

Think: Listen, read, explore, try, think

Motivation for Synchronization

- Cooperating Computers/Systems/Processes/Threads/Tasks
 - Communication
 - Information/Resourse Sharing
 - Modularity/Convenience
 - Economy/Performance/Responsiveness/Scalability

Synchronization Overview

- Characterizing Cooperating Processes/Synchronization
- Sections (e.g. Critical Section)
- Some Mechanisms for dealing with Critical Section
 - Mutex
 - Semaphore
 - Monitor
- Other topics
 - Peterson's Solution
 - Classical Problems in Synchronization: Dining Philosopher, Bounded Buffer Problem, Producer-Consumer Problems

Issues of Cooperating Process (Part 1)

- Race Condition
 - When several processes manipulate same data concurrently and outcome depends on order of execution.
- Can you give examples of race conditions?

Issues of Cooperating Process Example

- Race Condition Example
 - When several processes manipulate same data concurrently and outcome depends on order of execution.

Banker Makes a Deposit

Step 1: Banker reads balance from an account

Step 2: Banker calculates the new balance (e.g. adding deposit/subtracting withdrawal)

Step 3: Banker stores or writes amount

• Entry Section, Critical Sections, Exit Section and Remainder Section do {

Entry Section Critical Section Exit Section Remainder Section } while (Some Condition)

K G Osafo-Maafo, Learning Operating Systems

Characterizing Cooperating Process Example

• Entry Section, Critical Sections, Exit Section and Remainder Section

do { //Banker Makes a Deposit

Entry Section	some work before starting the transaction
Critical Section	Step 1: Banker reads balance from an account
	Step 2: Banker calculates the new balance (e.g. adding deposit/subtracting withdrawal)
	Step 3: Banker stores or writes amount
Exit Section	some work after starting the transaction
Remainder Section	

} while (SomeCondition)

Issues of Cooperating Process (Part 2)

- Race Condition
 - When several processes manipulate same data concurrently and outcome depends on order of execution.
- Characterizing parts of code
 - Entry Section, Critical Section, Exit Section, Remainder Section
- Critical Section Requirements
 - Mutual Exclusion
 - Progress
 - Bounded Waiting

Approaches to handling Critical Sections by OS

- Preemptive Kernel vs Non-preemptive Kernels
 - Preemptive Kernel allows process to be preempted while in kernel mode
 - Non-preemptive Kernel allows process to run unless it blocks or yield control of CPU
- Cooperating Process/Tasks and Race Conditions in Kernels

- Behavior of Preemptive and Non-Preemptive Kernels
 - Preemptive Kernels can provide responsiveness
 - Non-preemptive Kernel avoids race conditions

Concepts in Synchronization

- Locking
- Atomicity
- Busy Waiting
- Spin Lock

Shared Data: shared_item; Boolean item_available

```
do {
```

```
Entry Section while (litem_available) {
}
item_available = false;

Critical Section // access and or modify shared item

Exit Section ltem_available = true;

Remainder Section
```

} while (SomeCondition)

Shared Data: currency account_balance; Boolean account_available = true

do {	//Banker Makes a Deposit	
Entry Section	while (!account_available) { } account_available = false;	some work before starting the transaction
Critical Section	var balance = account_balance	Step 1: Banker reads balance from an account
	balance = balance + deposit_amount	Step 2: Banker calculates the new balance
	account_balance = balance	Step 3: Banker stores or writes amount
Exit Section	account_available = true;	some work after starting the transaction
Remainder Section		some work before starting the transaction
3 1 11		

} while
(SomeCondition)

Shared Data: currency account_balance; Boolean account_available

do {	//Banker Makes a Deposit
Entry Section	lock_account();
Critical Section	var balance = account_balance
	balance = balance + deposit_amount account_balance = balance
Exit Section	release_account();
Remainder Section	

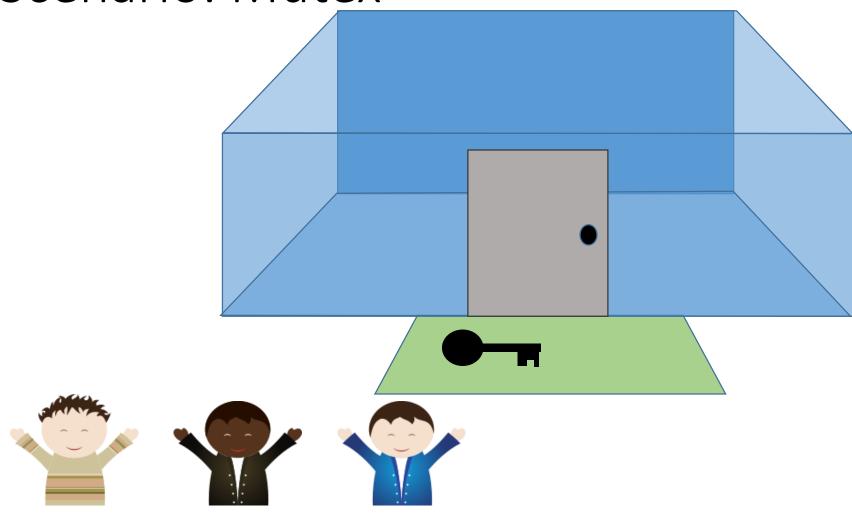
```
function lock_account()
{
    while (!account_available) {
    }
    account_available = false;
}

function release_account()
{
    account_available = true;
}
```

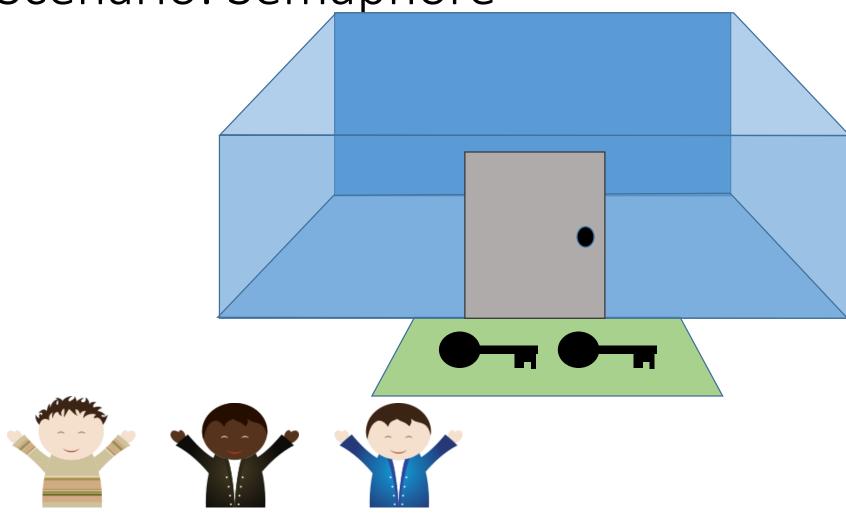
Synchronization Techniques

- Mutex Locks
 - A mutually exclusive lock with methods to
 - lock()/acquire()/hold()
 - unlock()/release()
- Semaphores
 - Integer variable accessible through
 - wait()/decrement()/acquire()
 - signal()/increment()/release()
- Monitors

Scenario: Mutex



Scenario: Semaphore



Scenario: Monitor

