

## Binary Search.

- Binary Search.
- Order agnostic Binary Search.
- 1st and last occurrences of an element.
- Count of element in sorted Array.
- # of times → array is rotated.
- Find an element in rotated sorted Array.
- Searching in nearly sorted Array.
- Floor / ceil of an element.
- Next Letter.
- Index of last 1 in sorted Array.
- Find the position of an element in a sorted Array.
- min diff element in a sorted Array.
- Bitonic Array - max element.
- Search in a Bitonic Array.
- Search in Row wise + Col wise sorted matrix.
- Find Element in sorted Array that appears only once.
- Allocate min # of pages.

→ If Question have sorted key word then there will be more chance to apply Binary search to reduce time complexity.

### Binary Search.

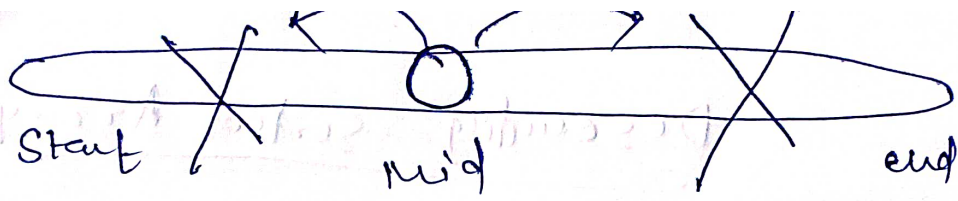
arr[]: 1 2 3 4 5 6 7 8 9 10  
          start                   mid                   end

Key = Search ②

Return index of 2 (i.e. 1)

- One method is of linear search in  $O(n)$
- But here we have sorted array so we can apply Binary search.
- int Binary\_search(arr[], key)  
  {  
    int start = 0;  
    int end = arr.size();  
    // If key element is less than mid then keep end to mid - 1.  
    // If key element is greater than mid so left-hand-side is useless so move start to mid + 1.





Code.

$$\rightarrow \begin{matrix} n \\ n/2 \\ n/4 \\ \dots \\ O(\log_2 n) \end{matrix}$$

while (start < end)

    < int mid = (start + end) / 2 ;

    int mid = start + (end - start) / 2 ;

    ↳ optimized

    if (key >= arr[mid])

        ↳ int overflow

        return mid ;

    else if (key < arr[mid])

        end = mid - 1 ;

    else

        start = mid + 1 ;

    }

    start = (start + end) / 2 ;

    // to avoid integer overflow

## Descending Sorted Array.

$\text{arr}[]: \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 20 & 17 & 15 & 14 & 13 & 12 & 10 & 9 & 8 & 7 \end{matrix}$   
 $\uparrow$   $\text{mid}$   $\uparrow$   $\text{end}$   
 key: 4 start

• Now if key  $<$  less than mid then set  
 $\text{mid} = \text{start} + 1$

• and if key  $>$  greater than mid then set  
 $\text{mid} = \text{end} - 1$

Code

```

    if (key < arr[mid])
        start = mid + 1
    else
        end = mid - 1;
    
```

always keep in mid as

$\text{mid} = \text{start} + (\text{end} - \text{start}) / 2$

for better passes of BST cases =



## Order Agnostic Search

arr[]:



key : 5

Sorted Array, we don't know either if  
it is ascending or descending

$arr[0] < arr[1] \Rightarrow$  ascending  
 $arr[0] > arr[1] \Rightarrow$  descending

• we will have a sorted array and our  
key element, search an element in this  
but we don't know, either if it is ascending  
and descending.

So check

Compare  
any two elements  
 $arr[0] < arr[1] \Rightarrow$  ascending  
 $arr[0] > arr[1] \Rightarrow$  descending

• and check for size as well, if size  
is one then return the same  
key element