Lab 0 — SystemVerilog

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Abstract

This document helps you familiarize yourself with SystemVerilog HDL (Hardware Description Language), which you will use to design digital logic blocks, and then you will learn the basic of verifying your RTL design.

In case you meet an error or have any improvement in this document, please email the TA: cxhai.sdh221@hcmut.edu.vn with the subject

"[COMPARCH203: FEEDBACK]"

1 Objectives

- Familiarize yourself with basic design flow.
- Review some basics of SystemVerilog HDL with existing codes.
- Implement the design on KIT.

1.1 Prerequisites

You should prepare two essential softwares:

- Ubuntu
- Quartus

Additionally, you need to learn some common commands to work in UNIX/Linux environment. You may find this cheatsheet useful.

2 Design Flow

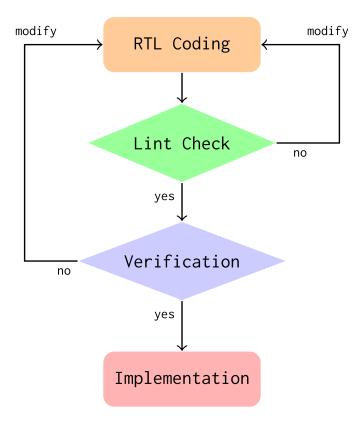


Figure 1: RTL design flow

Figure 1 shows the basic design flow of design a digital logic block. In this diagram, designers have to follow four steps:

- 1. **RTL** Coding is consisted of designing algorithms, FSMs, etc. to handle the requirements of the problem or expected designs. In this document, you will use SystemVerilog HDL, a versatile HDL for both design and verification.
- 2. **Lint Check** is a process of static code analysis to check the quality of the code. Lint Check will check your code with a list of rules, to find any violations syntax errors or potential code lines or blocks that may cause errors or bugs when running. This step is needed to perform right after you finish RTL Coding. If there are any flags or errors, you need to read them carefully to know exactly the rules you violated to fix your code.
- 3. **Verification** is a process of running your code to check correctness your design. Based on the requirements, each input will have only one expected output, so you have to interpret the log or the waveform generated from this step to fix your design. After fixing your design, don't forget linting again.
- 4. **Implementation** is to implement the design into real hardwares. Your simulation and the reality are sometimes not the same.

3 Combinational Logic Modeling

3.1 Problem

Design the combinational logic in Figure 2.

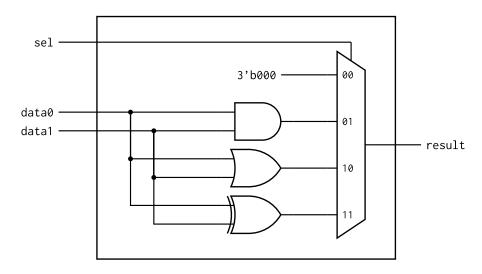


Figure 2: Sample combinational logic

 $! \rightarrow$ The data width of data0, data1, and result is 3, and that of sel is 2.

3.2 Design

3.2.1 Creating Project

A project should be organized systematically to manage and maintain easily later. First, in the terminal, type the below command to download the example projects of **lab 0** from GitHub to your local computer.

```
git clone https://github.com/joachimcao/lab0
```

Inside the new-downloaded directory, labo, there are two directories, exo1 and exo2. The former is of this section, go to that directory by typing cd labo/exo1. Type ls to see all files and directories in the current directory; an alternative is tree which displays the tree directory.

```
~/gi/c/lab0/ex01 $ ls
drwxr-xr-x@ - ioachimus 21 Feb 15:53 quartus
drwxr-xr-x@ - ioachimus 21 Feb 15:53 src
drwxr-xr-x@ - ioachimus 21 Feb 15:55 tb
```

In this sample project, you may notice:

src contains source files — SystemVerilog files.

tb contains testbench files

- makefile is a bash script to run several command lines conveniently.
- filelist contains all design files' names.
- top.sv is to instantiate the design to be tested.
- driver.cpp is a C++ file to drive inputs.
- tb_top.cpp is a C++ testbench file.

quartus contains DE2/DE10 related files.

- de2_pin.qsf is a pin assignments file for DE2 with its pins named according to this file.
- de10_pin.qsf is a pin assignments file for DE10 with its pins named according to this file.
- wrapper.sv is to connect the design's pins to those of DE2/DE10.
- $! \rightarrow$ All design files have to be placed in src.
- ! → In tb, only filelist, top.sv, and driver.cpp will be configured. In quartus, only wrapper.sv will be edited.

```
Question

Why do you need wrapper.sv?
```

3.2.2 RTL Coding

As all design files are placed into <u>src</u>, type <u>gedit src/design_1.sv</u> to open that file while you're in directory <u>ex01</u>. If you want to use <u>vim</u>, you could check this website for some basic commands in vim.

```
module design_1 (
      // input
      input logic [2:0] data0_i,
      input logic [2:0] data1_i,
      input logic [1:0] sel_i,
      // output
      output logic [2:0] result_o
    );
10
      // local declaration
11
      logic [2:0] and_tmp; // temporary for and result
      logic [2:0] or_tmp; // temporary for or result
13
      logic [2:0] xor_tmp; // temporary for xor result
14
      assign and_tnp = data0_i & data1_i;
      assign or_tmp = data0_i | data1_i;
17
      assign xor_tmp = data0_i ~^ data1_i;
18
19
      always_comb begin : proc_mux
20
        case (sel_i)
          2'b00: result_o =
22
          2'b01: result_o = and_tmp;
23
          2'b10: result_o = or_tmp;
24
          2'b10: result_o = xor_tmp;
25
        endcase
26
      end
27
28
    endmodule : design_1
```

3.2.3 Lint Check

To perform lint check, go to directory to by using cd tb. In this directory, all design files are listed in filelist. This example only has design_1.sv in src directory, the content of filelist is:

```
../src/design_1.sv
```

Then, in tb directory, run make lint, powered by Verilator here.

```
~/gi/c/l/ex01/tb $ make lint
   -----> LINT CHECK <-----
   %Warning-IMPLICIT: ../src/design_1.sv:16:10: Signal definition not found,
    : ... Suggested alternative:
                                              → 'and_tmp'
      16 |
             assign and_tnp = data0_i & data1_i;
                    ^~~~~~
                      ... For warning description see
                      → https://verilator.org/warn/IMPLICIT?v=5.003
                      ... Use "/* verilator lint_off IMPLICIT */" and lint_on
                      → around source to disable this message.
   %Warning-WIDTH: ../src/design_1.sv:16:18: Operator ASSIGNW expects 1 bits
    \hookrightarrow on the Assign RHS, but Assign RHS's AND generates 3 bits.
                                           : ... In instance design_1
10
      16
             assign and_tnp = data0_i & data1_i;
11
   %Warning-UNDRIVEN: ../src/design_1.sv:12:15: Signal is not driven:
13
    → 'and_tmp'
                                              : ... In instance design_1
14
      12 |
             logic [2:0] and_tmp;
15
16
   %Warning-UNUSEDSIGNAL: ../src/design_1.sv:16:10: Signal is not used:
17
    → 'and_tnp'
                                                  : ... In instance design_1
18
             assign and_tnp = data0_i & data1_i;
      16 |
19
20
   %Warning-CASEOVERLAP: ../src/design_1.sv:25:7: Case values overlap (example
    → pattern 0x2)
      25 |
                 2'b10: result_o = xor_tmp;
22
         23
   %Warning-CASEINCOMPLETE: ../src/design_1.sv:21:5: Case values incompletely

→ covered (example pattern 0x3)

      21 |
               case (sel_i)
25
               ^~~~
         26
   %Error: Exiting due to 6 warning(s)
27
   make: *** [makefile:47: lint] Error 1
   ~/gi/c/l/ex01/tb $
```

The example code has 6 warnings. For the first one — <code>%Warning-IMPLICIT</code>, the log specifically states the violation, a suggestion to tackle it, and also a link for more information on that particular warning. The original design, indeed, has a typo in line 16, fix it and lint again.

```
~/gi/c/l/ex01/tb $ make lint
   -----> LINT CHECK <-----
   %Warning-CASEOVERLAP: ../src/design_1.sv:25:7: Case values overlap (example
    → pattern 0x2)
                 2'b10: result_o = xor_tmp;
      25 |
         5
                         ... For warning description see
                         → https://verilator.org/warn/CASEOVERLAP?v=5.003
                         ... Use "/* verilator lint_off CASEOVERLAP */" and
                         → lint_on around source to disable this message.
   %Warning-CASEINCOMPLETE: ../src/design_1.sv:21:5: Case values incompletely

→ covered (example pattern 0x3)

      21 |
               case (sel_i)
9
10
   %Error: Exiting due to 2 warning(s)
   make: *** [makefile:47: lint] Error 1
   ^{\prime}gi/c/l/ex01/tb $
```

This time, linting alerts the violation of CASEOVERLAP, which is, according to the log, "case values overlap" because of the value 2'b10 is repeated twice. Fix this error at line 25 and lint again, and there are no warnings or errors.

```
~/gi/c/l/ex01/tb $ make lint
2 -----> LINT CHECK <-----
3 ~/gi/c/l/ex01/tb $
```

Now, the design is free of syntax errors and potential bugs, but, still, the next step will prove its correctness.

```
Change the value [2:0] to [1:0] in the line 13 of design_1.sv?
What is the error that lint reports?
Click on the link that Verilator provides and read the information.
```

3.3 Verification

3.3.1 **Setup**

This process is to verify the correctness of the design — for each set of inputs, only one set of outputs is right. Although designers could drive a list of inputs and compare the corresponding outputs they've already computed with the design's outputs — directed tests, verifying a design with a vast number of inputs will be improbable and tedious. Hence, randomly driving the inputs and automatically monitoring the outputs will provides a more accurate and reliable verification methodology as in Figure 3.

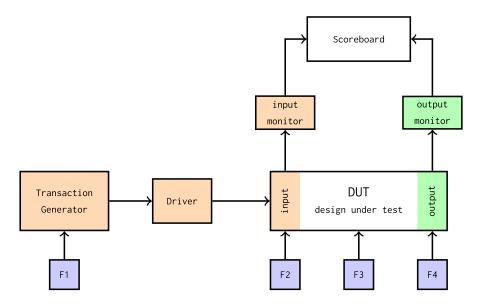


Figure 3: Verification methodology

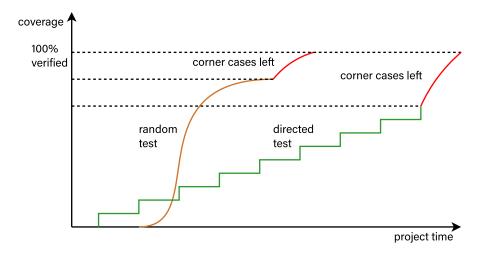


Figure 4: Compare time to verify a design

 $! \rightarrow Do NOT modify tb_top.cpp$. You may do so later.

First, preparing top.sv, which is alreay written and placed into tb directory.

```
module top (
      // inputs
2
      input logic
                          clk_i,
      input logic [2:0]
                          data0_i,
      input logic [2:0]
                          data1_i,
      input logic [1:0]
                          sel_i ,
      // outputs
9
      output logic [2:0] result_o
    );
10
11
      design_1 dut (
12
        .data0_i (data0_i ),
13
        .data1_i (data1_i ),
14
        .sel_i
                (sel_i
15
        .result_o(result_o)
16
      );
17
18
      always @(posedge clk_i) begin : proc_assertions
19
        if (sel_i == 2'b00)
20
          assert (result_o == '0);
21
        if (sel_i == 2'b01)
          assert (result_o == (data0_i & data1_i));
23
        if (sel_i == 2'b10)
24
          assert (result_o == (data0_i | data1_i));
25
        if (sel_i == 2'b11)
26
          assert (result_o == (data0_i ^ data1_i));
27
28
      end
29
    endmodule : top
30
```

Second, according to Figure 3, the design under test, design_1, is instantiated in top.sv, line 12 — 17. Then, line 19 — 28 set some assertions or required contraints (F3), which can easily be inferred from Figure 2. assert keyword with the expression inside will signal the EDA to stop when the expression failed and report it on the monitor or screen. As this design is simple and has four cases in total, this example uses if to set up the "check point." You may notice, even though the design is de facto a combinational logic, you have to set a clock because assert needs it as a reference time to check the expression. F2 and F4 are already configured in tb_top.cpp, and driver.cpp will set the inputs of the DUT.

```
#define MAX_SIM 20

void set_random(Vtop *dut, vluint64_t sim_unit) {
  dut->data0_i = rand()%8;
  dut->data1_i = rand()%8;
  dut->sel_i = rand()%4;
}
```

MAX_SIM is the number of inputs generated. Basically, illustrated in the code above, the randomized values are set by following a syntax of dut-> and the inputs with rand() function. This function in C++ generates a 32-bit number, but because an input varies in width, the modulo operator will restrict the value, which rand() assigns to the input.

```
Question

If data0_i varies from 3 to 7, how to set it in driver.cpp?
```

3.3.2 Simulation

In directory tb, run make sim to simulate.

The error states that the assertion at line 27 failed. In top.sv, that line indicates the design failed to handle the case 2'b11, the xor operator. To investigate the waveform for debugging, run make wave. Figure 5 shows the result 3'b101 instead of 3'b010.

A little bit further, design_1.sv in src directory containing the source code has a bug in line 18. The operator is intentionally altered, which is xnor instead. Fix it and run simulation again.

```
~/gi/c/l/ex01/tb $ make sim

> BUILD -----

...

> SIMULATING -----

~/gi/c/l/ex01/tb $
```

FABULOUS!! The design is successfully debugged!

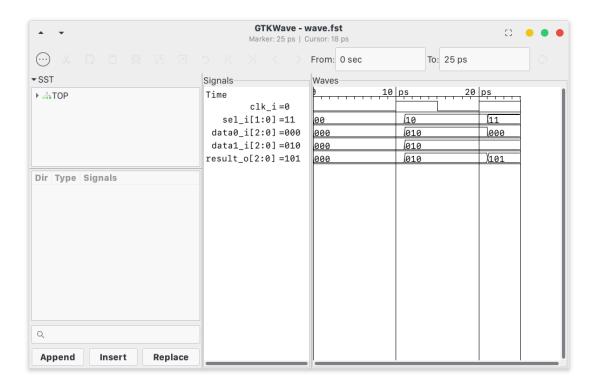


Figure 5: Waveform of the first simulation

3.4 Implementation

3.4.1 Setup

To implement the design in FPGA, wrapper.sv is to instantiate the design and connect FPGA pins to it. In this example, DE2 is used, so you could look up the pin names in this link. You may check DE2 User Manual in quartus directory. design_1.sv uses switches for inputs and LEDs for outputs. The example wrapper.sv is already written and placed into quartus directory.

```
module wrapper (
      // inputs
      input logic [7:0] SW,
      // outputs
      output logic [2:0] LEDR
5
    );
6
      design_1 dut (
        .data0_i ( SW[2:0]),
        .data1_i ( SW[5:3]),
10
        .sel_i
                 (SW[7:6]),
11
        .result_o(LEDR[2:0])
12
      );
13
14
    endmodule : wrapper
```

Question

Explain how FPGA's pins connect to the design.

It's time to use Quartus.

- Create Quartus Project: File → New Project Wizard
 The working directory is quartus in the example project.
 Name the new project wrapper to match the wrapper.sv module.
- 2. Select all source codes, design_1.sv and wrapper.sv.

 Make sure in the Type column, all files are SystemVerilog HDL.
- 3. For DE2, the device is EP2C35F672C6.
- 4. Import pin assignment of DE2. The file is de2_pin.qsf placed in quartus directory already.
- 5. Ctrl L or Processing \rightarrow Start Compilation. Mkae sure there are no errors or critical warnings.
- 6. Connect DE2 with your computer via **Blaster**.
- 7. Check the connection: Tools \rightarrow Programmers
- 8. If the device is not connected, select Hardware Setup
- 9. Click Start to start the implementation.

4 Sequential Logic/FSM Modeling

Problem

Design a counter to count each time a button is pressed then display the number using 7-segment LED.

4.1 Analysis

The design requires a button and a 7-segment LED. It, absolutely, needs a clock and an active low reset. Based on the requirements:

- A module to receive the signal from a button and its output will be high in only one cycle if the button is pressed.
- A counter to count, let's say, from 0 to 9 (maximum).
- A module to decode a binary number to data that drives a 7-segment LED.

The design is shown in Figure 6

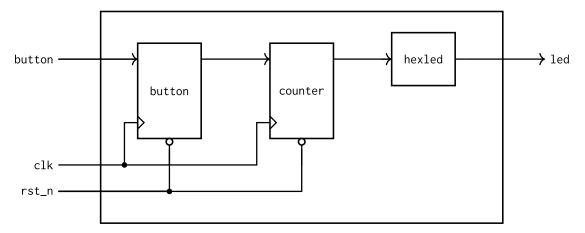


Figure 6: Block diagram of the counter

The button module is designed as an FSM with three states: IDLE, PRESS, and HOLD.

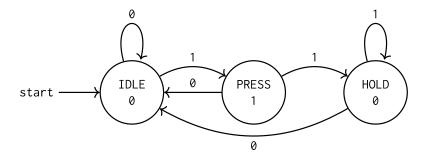


Figure 7: Button FSM

4.2 Design

This project is ex02.

4.3 Verification

Before diving into linting and simulating, here are some notices regarding writing assertions:

- 1. This "more complex" example needs to assert the data between module button (output) and counter (input), so the assertions have to be placed in design_2.sv. The "clock requirement" of using assertions is satisfied.
- 2. \$past keyword is to sample the data in the previous cycle, and thus line 2–6 and line 14 are needed to sample the write data. Also, When using \$past, remember to AND the pastvld.
- 3. Line 1 and 15 are to make sure the assertions are checked when running simulation with Verilator but ignored when running Quartus.

```
`ifdef VERILATOR
      /*verilator lint_off UNUSED*/
      logic pastvld;
      always @(posedge clk_i) begin : proc_setup_past
        pastvld <= 1'b1;</pre>
      end
      always @(posedge clk_i) begin : proc_assertions
        if (pastvld && $past(inc))
          assert(!inc);
10
11
        assert(counter <= 4'h9);</pre>
12
      end
13
      /*verilator lint_on UNUSED*/
    `endif
```

```
Question

Explain the two assertions in design_2.sv.
```

Here are some notices regarding driver.cpp:

1. This design has the maximum number of 9, so MAX_SIM must be set to a large number.

2. Because driver.cpp is a C++ file, assigned values could be manipulated usefully. Such as, the reset should be low for several cycles from the beginning, so using sim_unit, illustrated in Figure 8, as in line 4, it will certainly be low for 4 cycles from the beginning. Yet, it is just a part of the story.

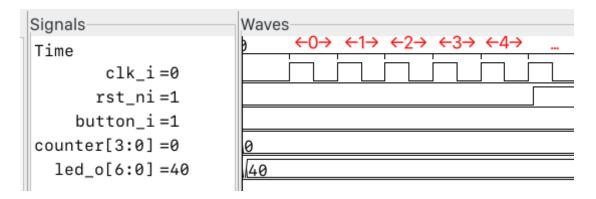


Figure 8: sim_unit

3. The reset should be 0 at some points, to prove that the counter will reset to 0 and count again, but its probability should be low to show the case that the counter only counts to 9 and stop. The expression [rand()%30] = 0 has 1/30 = 3.3% chance of being FALSE or 0. It is sufficient to create a reset to verify this design.

```
#define MAX_SIM 200

void set_random(Vtop *dut, vluint64_t sim_unit) {
  dut->rst_ni = (sim_unit > 4) && (rand()%30 != 0);
  dut->button_i = (rand()%8 >= 2);
}
```

Question

Compute the probability $Pr(button_i = 1)$. Explain the reason for choosing that value.

Question

Lint and debug this design.

4.4 Implementation

Question

Implement this design into DE2.