Complejidad de Algoritmos

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Potenciación 1.

$$T(n) = 2n + 6$$

T(n)
$$\epsilon$$
 $O(n)$

Busqueda Exhaustiva 2.

$$T(n) = 2n + 6$$

$$T(n) \in O(n)$$

$$T(n) = 2n + 5$$

$$T(n) \in \Omega(n)$$

$$T(n) = n + 3$$

$$T(n) \in \Theta(n)$$

3. Ordenamiento Burbuja

$$T(n) = n^2 + 3n$$

$$T(n) \epsilon O(n^2)$$

$$T(n) = n^2 + 2n$$

$$T(n) \in \Omega(n^2)$$

$$T(n) = \frac{1}{2}n^2 + \frac{3}{2}n$$

$$T(n) \in \Theta(n^2)$$

Ordenamiento de Selección

$$T(n) = n^2 + 3n + 8$$

$$T(n) \epsilon O(n^2)$$

$$T(n) = n^2 + 2n + 4 \qquad \qquad \text{T(n) } \epsilon \ \Omega(n^2)$$

$$T(n) \in \Omega(n^2)$$

$$T(n) = \frac{1}{2}n^2 + \frac{3}{2}n + 4$$

$$T(n) \in \Theta(n^2)$$

Emparejamiento de Cadenas

$$T(n) = n^2 + 5n + 1 \qquad T(n) \ \epsilon \ O(n^2)$$

$$T(n) \in O(n^2)$$

$$T(n) = 4n + 5$$
 $T(n) \epsilon \Omega(n)$

$$T(n) = \frac{1}{2}n^2 + \frac{5}{2}n + \frac{1}{2}$$
 $T(n) \in \Theta(n^2)$