Conditional Entropy

H(Y | X) is the expected value of H(Y | X=x)

$$\begin{aligned} &H(Y\mid X) = \ \Sigma_{x\in X} \ p(X=x) \ H(Y\mid X=x) \\ &= \ \Sigma_{x\in X} \ p(X=x) \ [\ -\Sigma_{y\in Y} \ p(Y=y\mid X=x) \ \log \ p(Y=y\mid X=x)] \\ &= - \ \Sigma_{x\in X} \ \Sigma_{y\in Y} \ p(X=x, \ Y=y) \ \log \ p(Y=y\mid X=x) \\ &= \ E_{p(X,Y)} \ (- \ \log \ p(Y=y\mid X=x)) \end{aligned}$$

 \Box H(Y | X) is the expected value of - log p(Y=y | X=x)

In the example

- Find H(Y|X)
- H(Y|X) = p(X=0) H(Y|X=0) + p(X=1) H(Y|X=1)
- p(X = 1) = p(X = 0) = 1/2

P(XY)	Y=0	Y=1
X=0	1/3	1/6
X=1	1/6	1/3

- $H(Y|X=0) = -\sum_{y \in \{0,1\}} p(Y=y|X=0) \log p(Y=y|X=0)$
- P(Y=0|X=0)= p(Y=0, X=0)/ p(X=0) = (2/6)/(1/2)=2/3
- P(Y=1|X=0)= p(Y=1, X=0)/ p(X=0) = (1/6)/(1/2)=1/3
- $H(Y|X=1) = -\sum_{y \in \{0,1\}} p(Y=y|X=1) \log p(Y=y|X=1)$
- P(Y=0|X=1)= p(Y=0, X=1)/ p(X=1) = (1/6)/(1/2)=1/3
- P(Y=1|X=1)= p(Y=1, X=1)/ p(X=1) = (2/6)/(1/2)=2/3

H(Y|X)

$$H(Y \mid X = 0) = -[p(Y = 1 \mid X = 0)\log p(Y = 1 \mid X = 0) + p(Y = 0 \mid X = 0)\log p(Y = 0 \mid X = 0)] = -(\frac{1}{3}\log\frac{1}{3} + \frac{2}{3}\log\frac{2}{3}) = 0.918$$

$$H(Y \mid X = 1) = -[p(Y = 1 \mid X = 1)\log p(Y = 1 \mid X = 1) + p(Y = 0 \mid X = 1)\log p(Y = 0 \mid X = 1)] = -(\frac{1}{3}\log\frac{1}{3} + \frac{2}{3}\log\frac{2}{3}) = 0.918$$

$$H(Y \mid X) = p(X = 0)H(Y \mid X = 0) + p(X = 1)H(Y \mid X = 1)$$

$$= (\frac{1}{2} \times 0.918) + (\frac{1}{2} \times 0.918) = 0.918 \quad bit$$

- For this example: H(X,Y)= H(X) + H(Y|X)
- 1.918 = 1 + 0.918

Theorem

• H(X, Y) = H(X) + H(Y|X)= H(Y) + H(X|Y)

• (Theorem 2.2.1 - CT)

• (Example 2.2.1 - CT)

H(X|Y)
H(X|Y)
H(Y|X)

Note that H(Y|X) ≠ H(X|Y)

Mutual Information

- X and Y are two random variables with joint probability distribution p(x,y).
- The information that observing Y gives about X is
 - -I(X;Y) = H(X) H(X|Y)
- This is the expected reduction in uncertainty about X when Y is known.
 - Similarly the information that X gives about Y is
 - -I(Y;X) = H(Y) H(Y|X)

Mutual Information

• I(X;Y) = H(X) - H(X|Y) $= -\sum_{x \in X} \sum_{y \in Y} p(X=x,Y=y) \log p(X=x)$ $-(-\Sigma_{x \in X} \Sigma_{y \in Y} p(X=x,Y=y) \log p(X=x|Y=y)$ = $\sum_{x \in X} \sum_{y \in Y} p(X=x,Y=y) \log [p(X=x|Y=y) / p(X=x)]$ $= \sum_{x \in X} \sum_{y \in Y} p(X=x,Y=y) \log [p(X=x,Y=y) / p(X=x)p(Y=y)]$ $= \sum_{x \in X} \sum_{y \in Y} p(X=x,Y=y) \log [p(Y=y|X=x) / p(Y=y)] = I(Y;X)$ That is: I(X;Y)=I(Y;X)

X says as much about Y as Y says about X.

Example: Erasure Channel

An information source:

$$p(X=0)= 2/3$$
, $p(X=1)= 1/3$
H(X) =.918 bits



Example: Erasure Channel

An information source:

$$p(X=0)= 2/3$$
, $p(X=1)= 1/3$
H(X) =.918 bits

How much is H(Y)?

Find p(y):

•
$$p(Y=0) = p(X=0)p(Y=0|X=0)$$

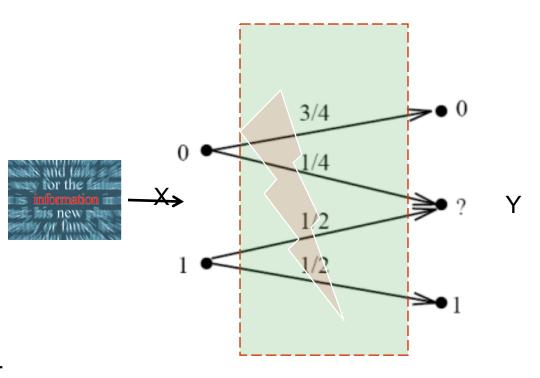
=2/3 x 3/4 =1/2

•
$$p(Y=1) = p(X=1)p(Y=1|X=1)$$

=1/3 x 1/2 =1/6

•
$$p(Y=?) = p(X=1)p(Y=?|X=1) + p(X=0)p(Y=?|X=0)$$

$$=1/3 \times 1/2 + 2/3 \times 1/4 = 1/3$$



Entropy of Y

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• Find H(Y) = -\Sigma_{y \in Y} p(y)log p(y)
= 1/2 \times 1 + 1/6 \times 2.58 + 1/3 \times 1.58
= 1.95 bits
```

H(Y) > H(X). Has information increased?

 I(X;Y) =H(X) - H(X|Y) is the information about X that Bob receives.

$H(X|Y) = \sum_{y \in Y} p(Y=y)H(X|Y=y)$

3/4

- H(X|Y=0) = 0
- H(X|Y=1) = 0
- $H(X|Y=?) = -\sum_{x \in X} p(x|Y=?) \log p(x|Y=?)$

•
$$p(X = 0|Y =?) = p(X = 0, Y =?)/p(Y =?)$$

= $p(X=0)p(Y=?|X=0) / p(Y=?)$
= $(2/3 \times 1/4)/(1/3) = 1/2$

•
$$p(X = 1|Y =?) = p(X = 1, Y =?)/p(Y =?)$$

= $p(X=1)p(Y=? | X=1) / p(Y=?)$
= $(1/3 \times 1/2)/(1/3) = 1/2$

Entropies

- $H(X|Y) = \sum_{y \in Y = \{0,?,1\}} p(y)H(X|Y=y)$ = $1/3 \times 1 = 0.33$ bits
- I(X; Y) = H(X) H(X|Y)= 0.9183 - 0.3333 = 0.5850 bits

- Note that H(X|Y =?) = 1 bit
- H(X|Y =?) > H(X) (= 0.9183 bits)

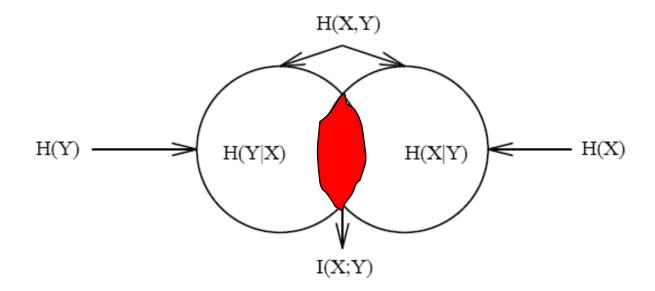
I(X;Y)

- Symmetry: The symmetry of this function implies:
 - information theory is not concerned with the cause/effect type of relation. It merely looks at statistical dependence.

- If X and Y are independent, we have
 - -H(X|Y) = H(X) and I(X;Y) = 0
 - Y does not give any information about X.

I(X;Y)

• I(X;Y) = H(X) + H(Y) - H(X,Y)



H(X,Y)			
H(X)			
	Н	(Y)	
$H(X \mid Y)$	I(X;Y)	H(Y X)	

DM Info Theory..