Discrete Optimization

Constraint Programming: Part II

Goals of the Lecture

Illustrating more complex constraint propagation

Showing that constraints have dedicated algorithms

- Branch and prune
 - pruning
 - reduce the search space as much as possible
 - branching
 - decompose the problem into subproblems and explore the subproblems

Branch and prune

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 - reduce the search space as much as possible
- branching
 - decompose the problem into subproblems and explore the subproblems

► Pruning

• use constraints to remove, from the variable domains, values that cannot belong to any solution

Branch and prune

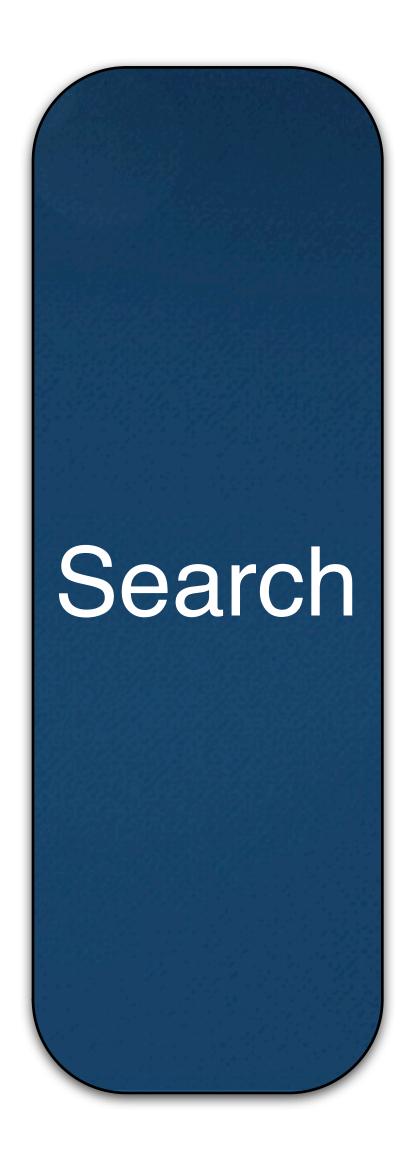
- pruning
 - reduce the search space as much as possible
- branching
 - decompose the problem into subproblems and explore the subproblems

► Pruning

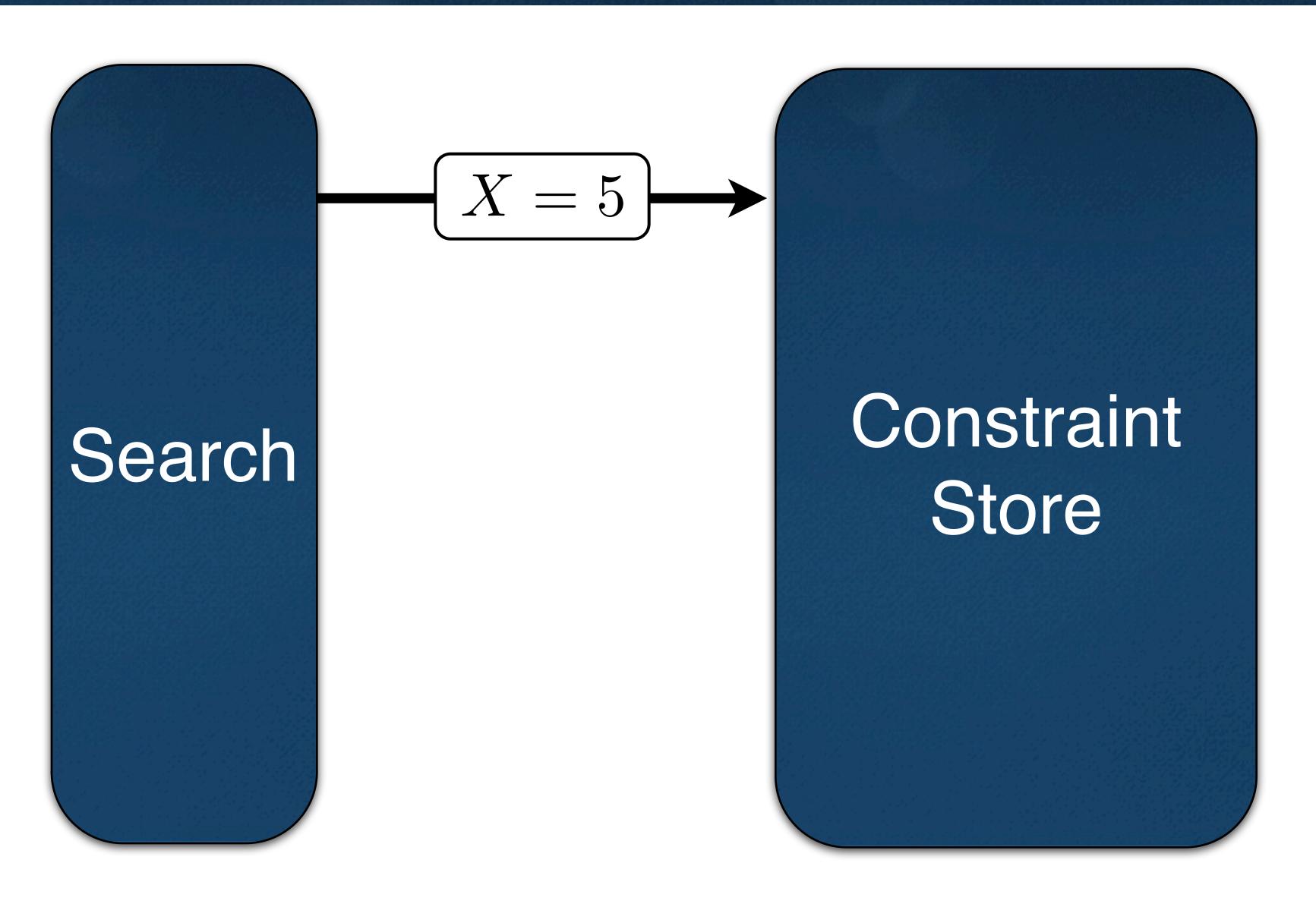
• use constraints to remove, from the variable domains, values that cannot belong to any solution

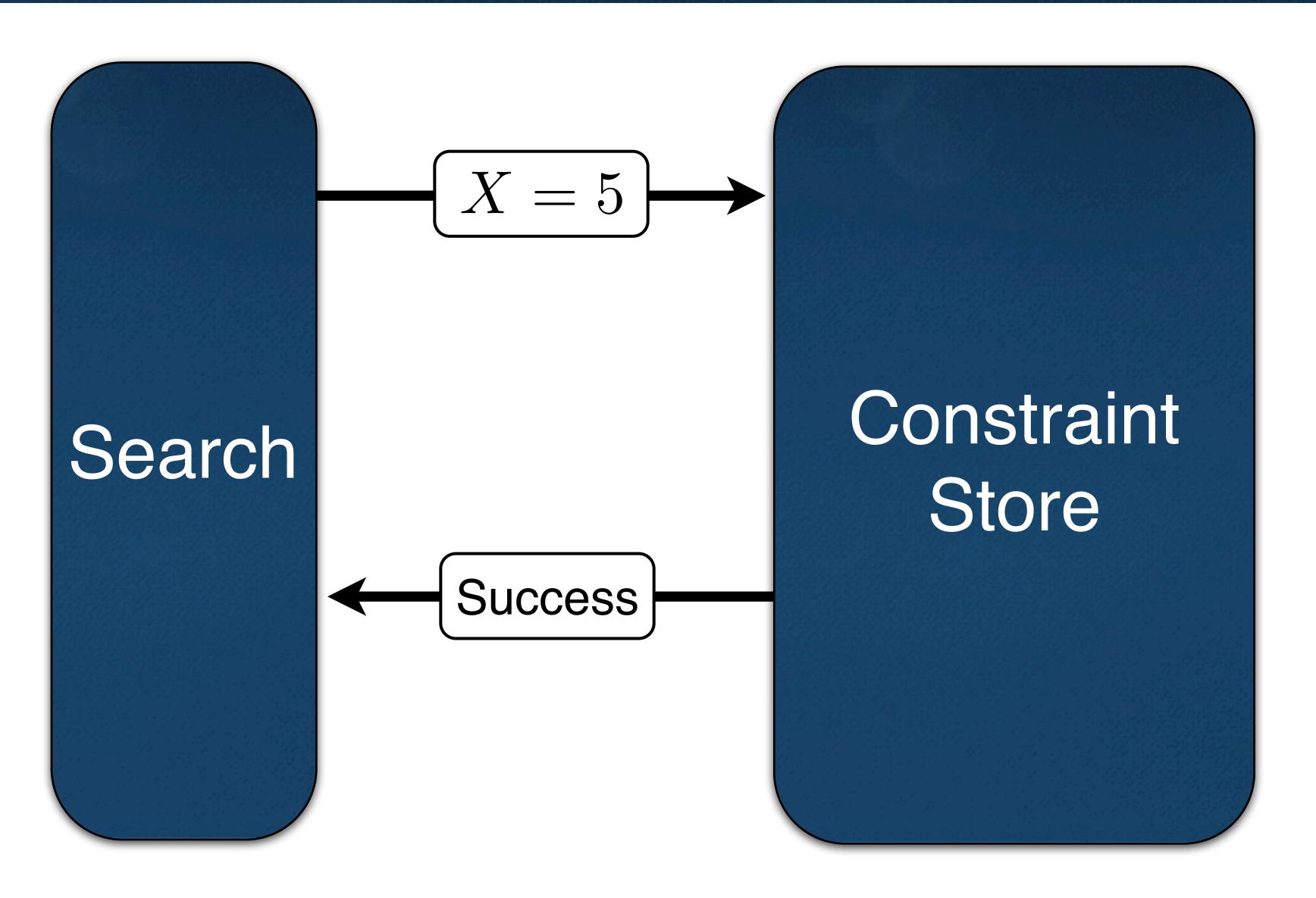
Branching

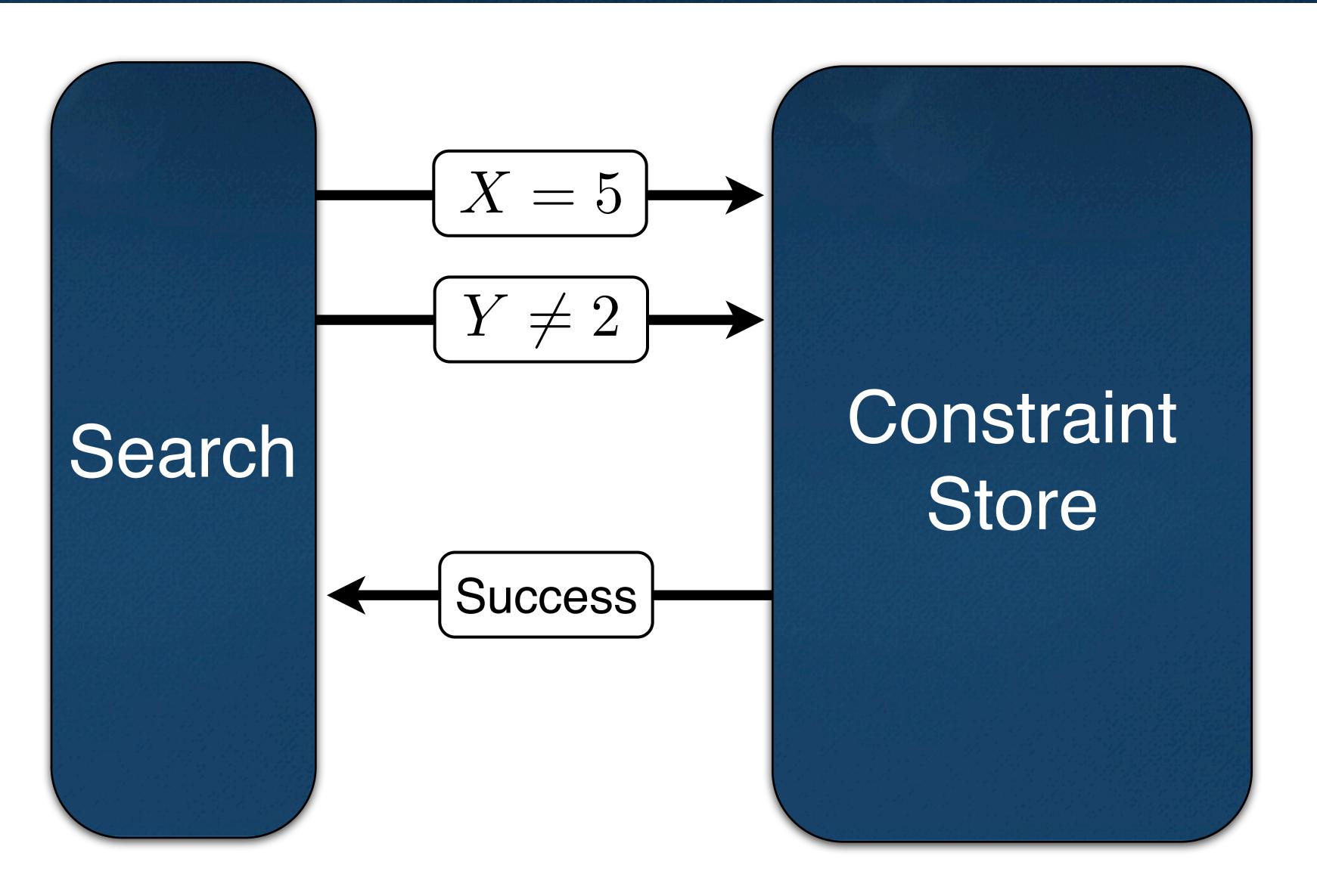
• e.g., try all the possible values of a variable until a solution is found or it can be proven that no solution exists

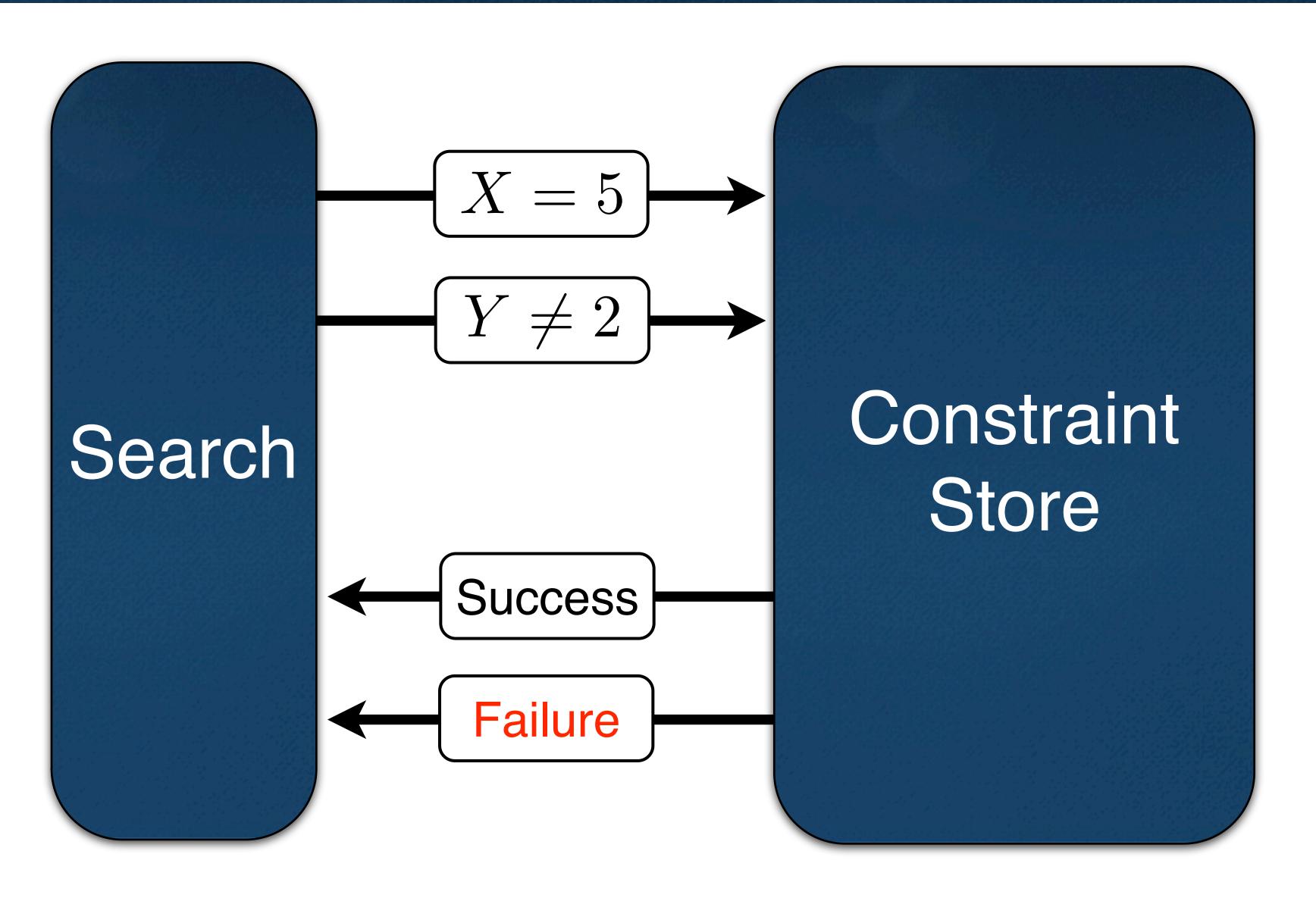


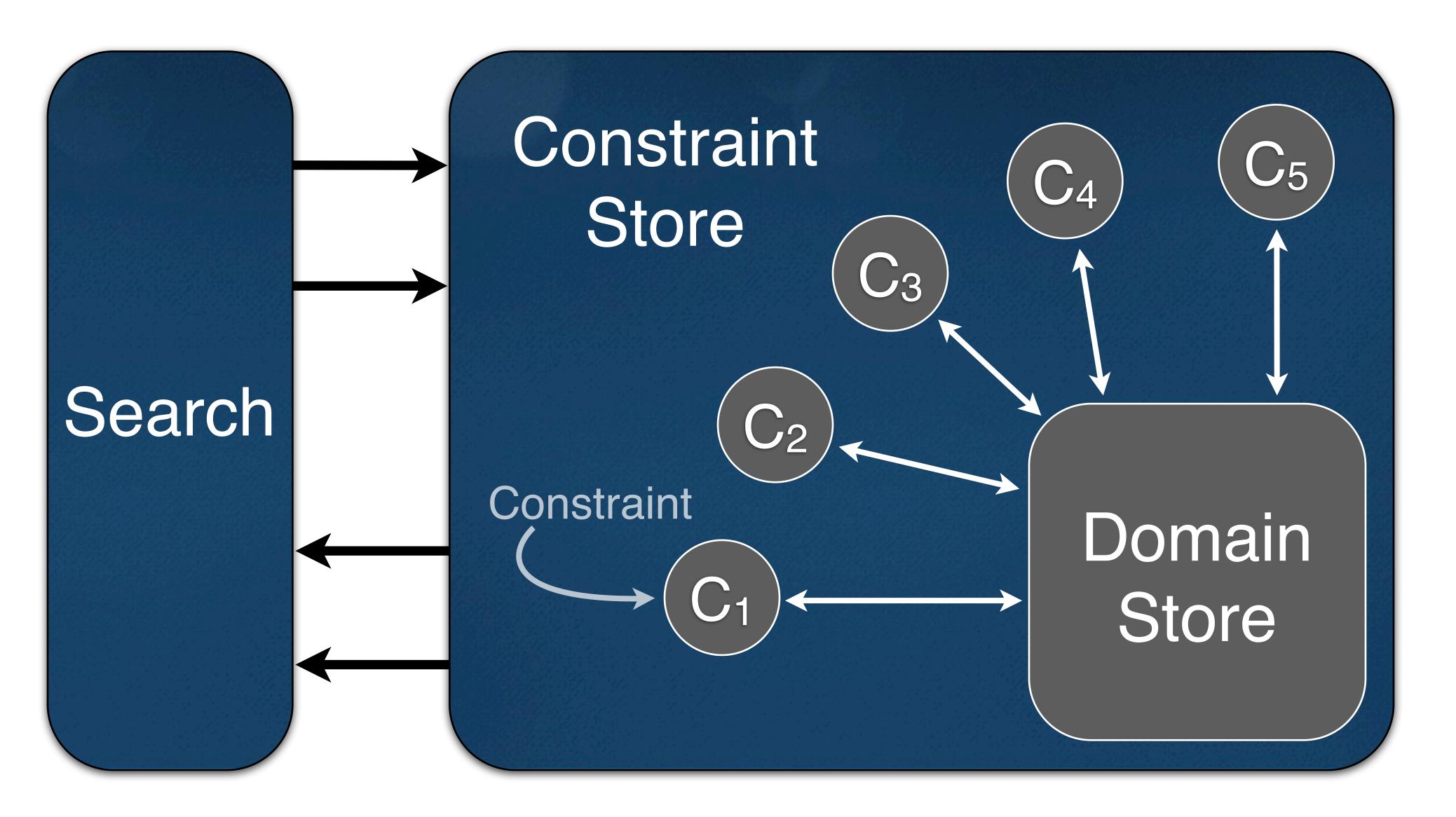












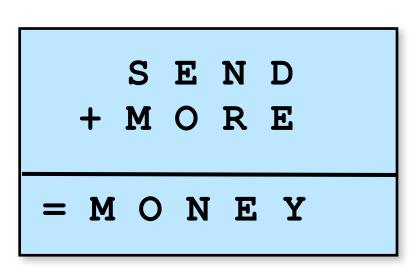
- What does a constraint do?
 - feasibility checking
 - pruning

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- Feasibility checking
 - can a constraint be satisfied given the values in the domains of its variables

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- ► Pruning
 - if satisfiable, determine which values in the domains cannot be part of any solution

- What does a constraint do?
 - feasibility checking
 - pruning
- Feasibility checking
 - can a constraint be satisfied given the values in the domains of its variables
- ► Pruning
 - if satisfiable, determine which values in the domains cannot be part of any solution
- The algorithms use dedicated algorithms for each constraint
 - they exploit the structure and properties of the constraint

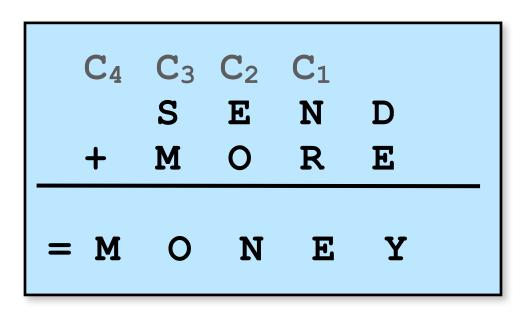
- Specification
 - assign different digits to letters to satisfy the addition



- Specification
 - assign different digits to letters to satisfy the addition
- Model to illustrate constraint propagation
 - -no claim that this is a good model: it is not

```
SEND
+ MORE
= MONEY
```

- Basic modeling
 - add carries explicitly like in kindergarten



- Basic modeling
 - add carries explicitly like in kindergarten
- What are the decision variables
 - there is a variable for each letter to denote the value of the letters
 - there is a variable for each carry

```
C<sub>4</sub> C<sub>3</sub> C<sub>2</sub> C<sub>1</sub>
S E N D
+ M O R E

= M O N E Y
```

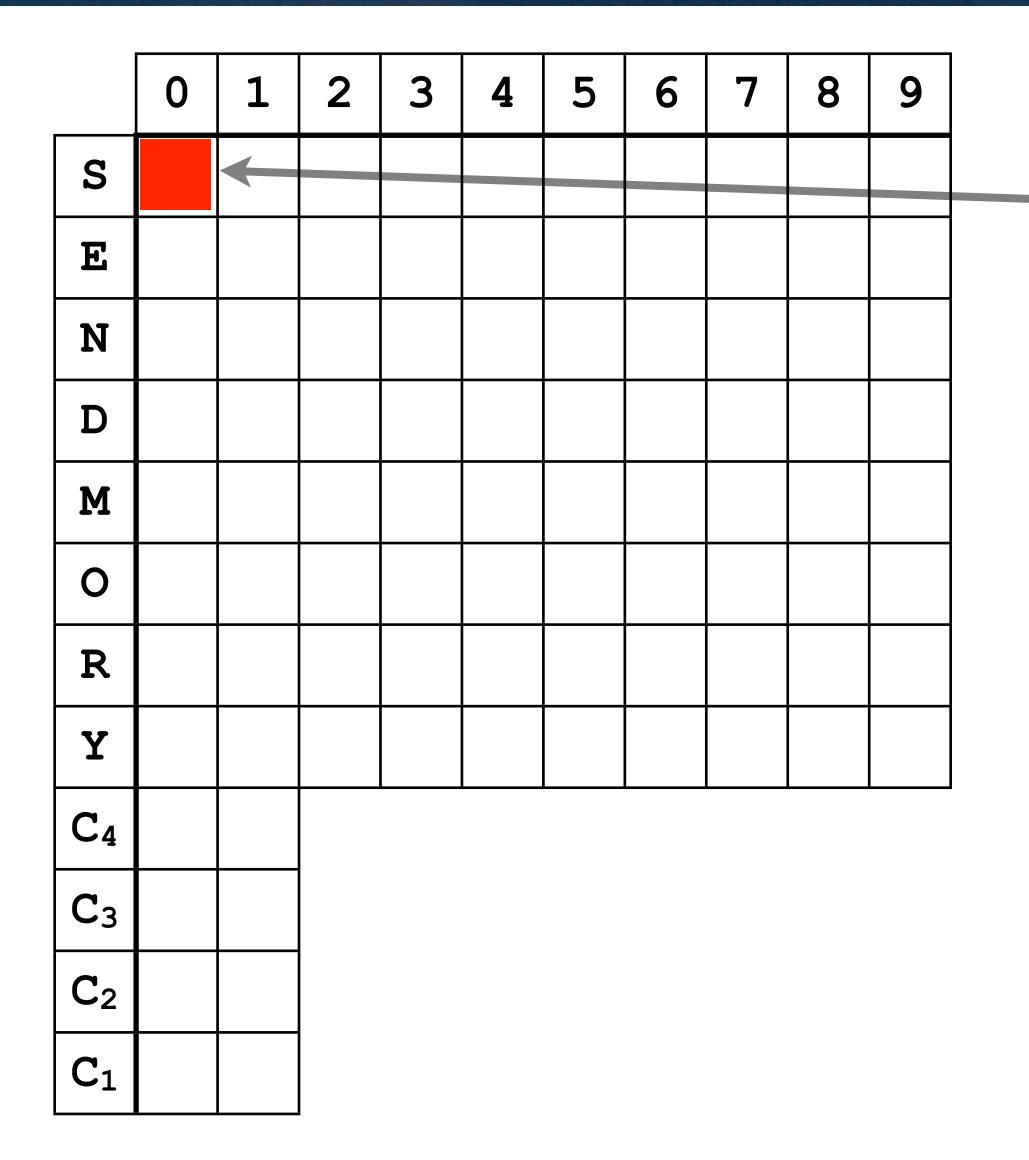
```
enum Letters = { S, E, N, D, M, O, R, Y};
range Digits = 0..9;
var{int} value[Letters] in Digits;
var{int} carry[1..4] in 0..1;
solve {
 forall(i in Letters, j in Letters: i < j)</pre>
   value[i] ≠ value[j];
 value[S] \neq 0;
 value[M] \neq 0;
                                 = value[M];
 carry[4]
 carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];
 carry[2] + value[E] + value[O] = value[N] + 10 * carry[3];
 carry[1] + value[N] + value[R] = value[E] + 10 * carry[2];
            value[D] + value[E] = value[Y] + 10 * carry[1];
```

```
C_4 C_3 C_2 C_1
S E N D
+ M O R E
= M O N E Y
```

_										
	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
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R										
Y										
C ₄										
C ₃										
C ₂										
C_1										

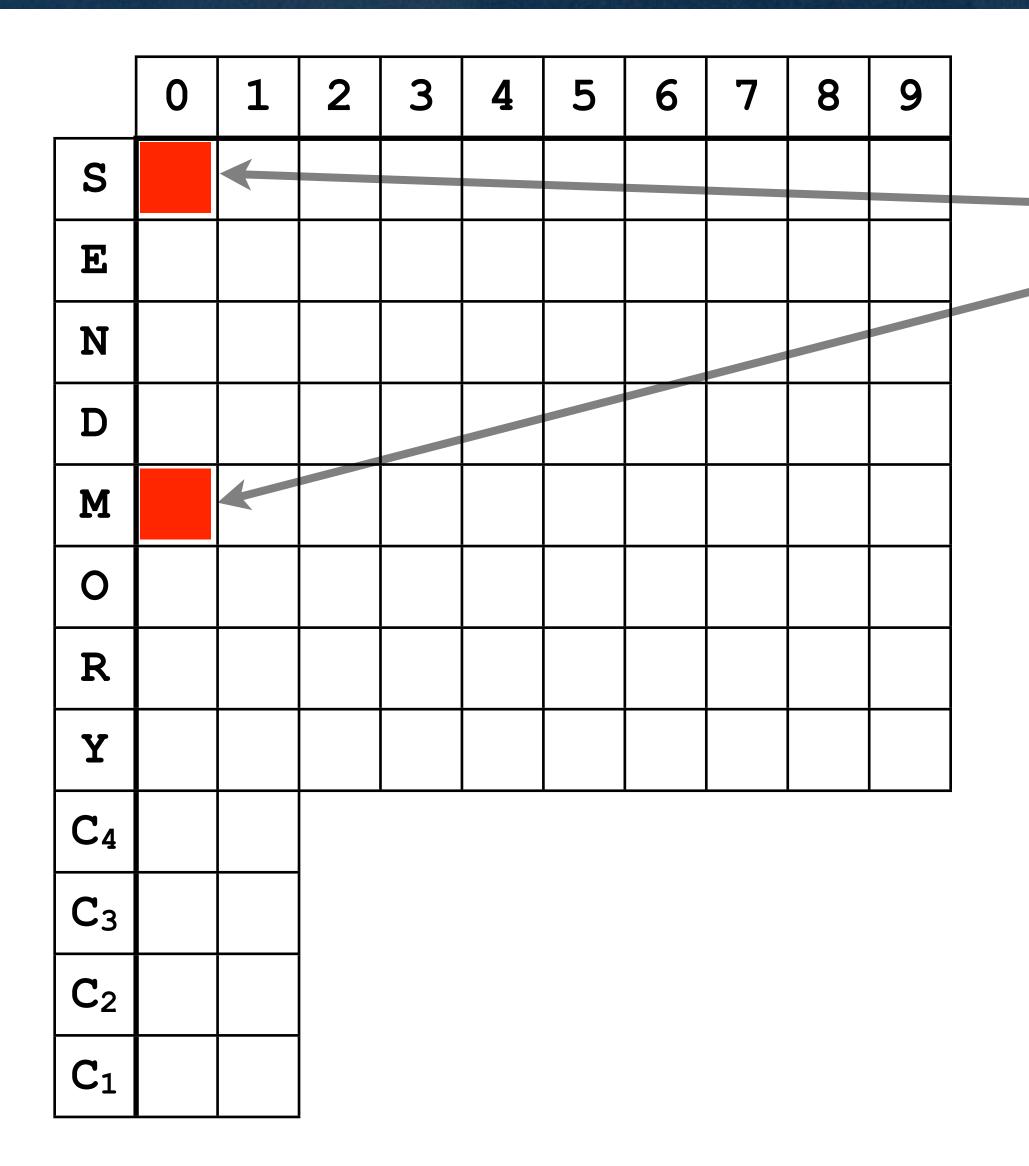
	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
0										
R										
Y										
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C ₂										
C ₁										

```
forall(i in Letters, j in Letters: i < j)
   value[i] ≠ value[j];
value[S] ≠ 0;
value[M] ≠ 0;
carry[4] = value[M];</pre>
```



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forall(i in Letters, j in Letters: i < j)
    value[i] ≠ value[j];

— value[S] ≠ 0;
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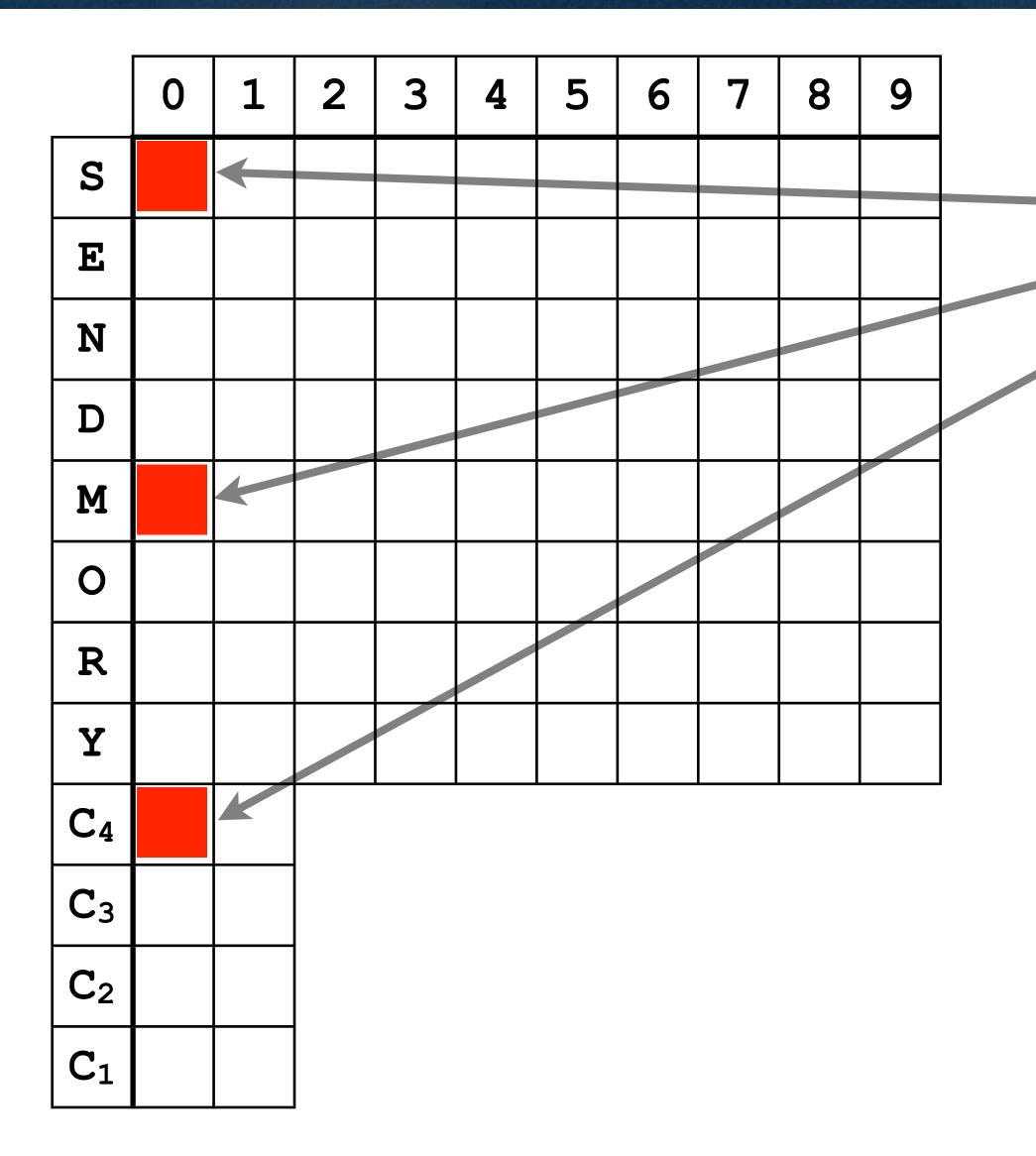


```
forall(i in Letters, j in Letters: i < j)
    value[i] ≠ value[j];

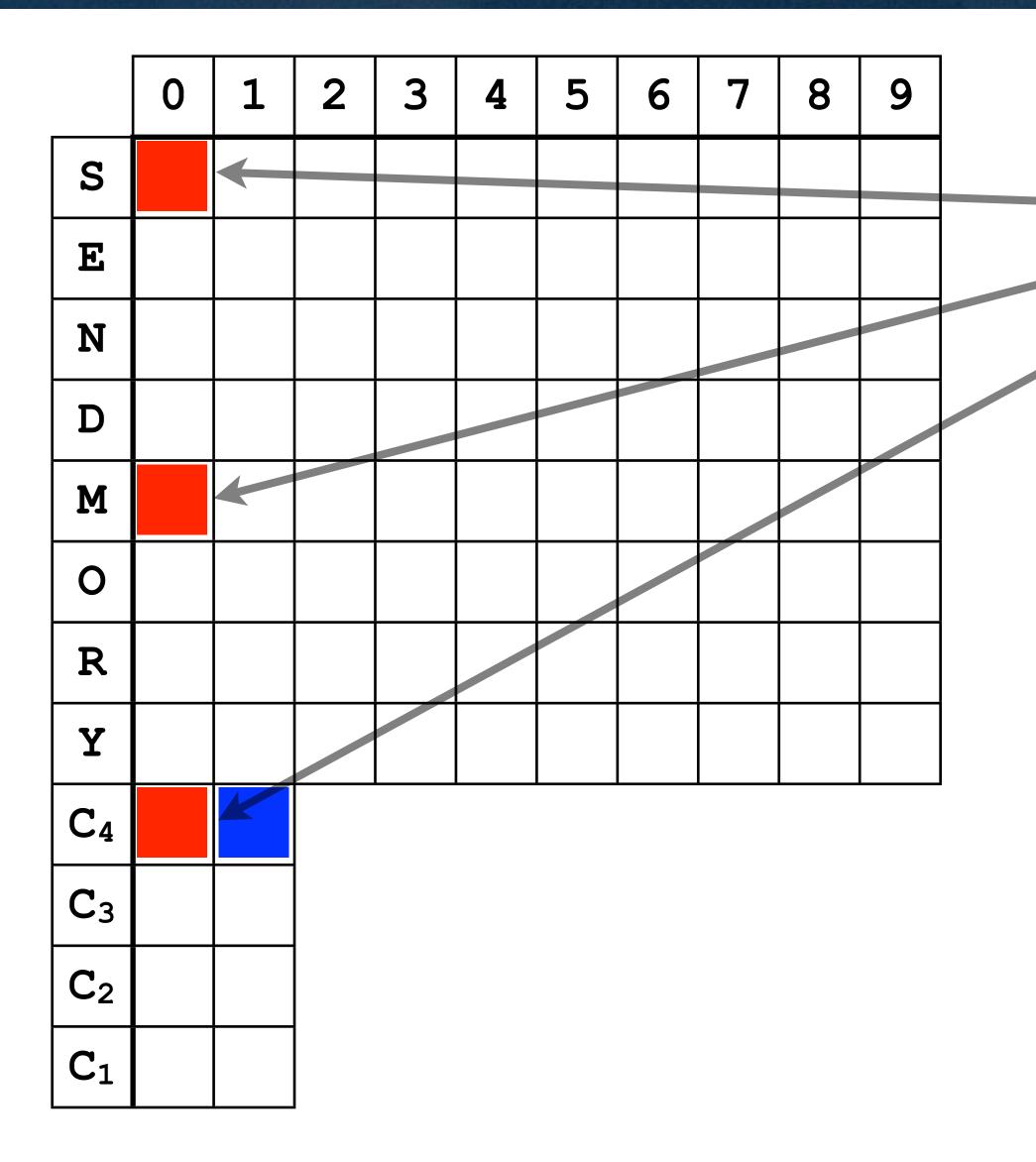
value[S] ≠ 0;

value[M] ≠ 0;

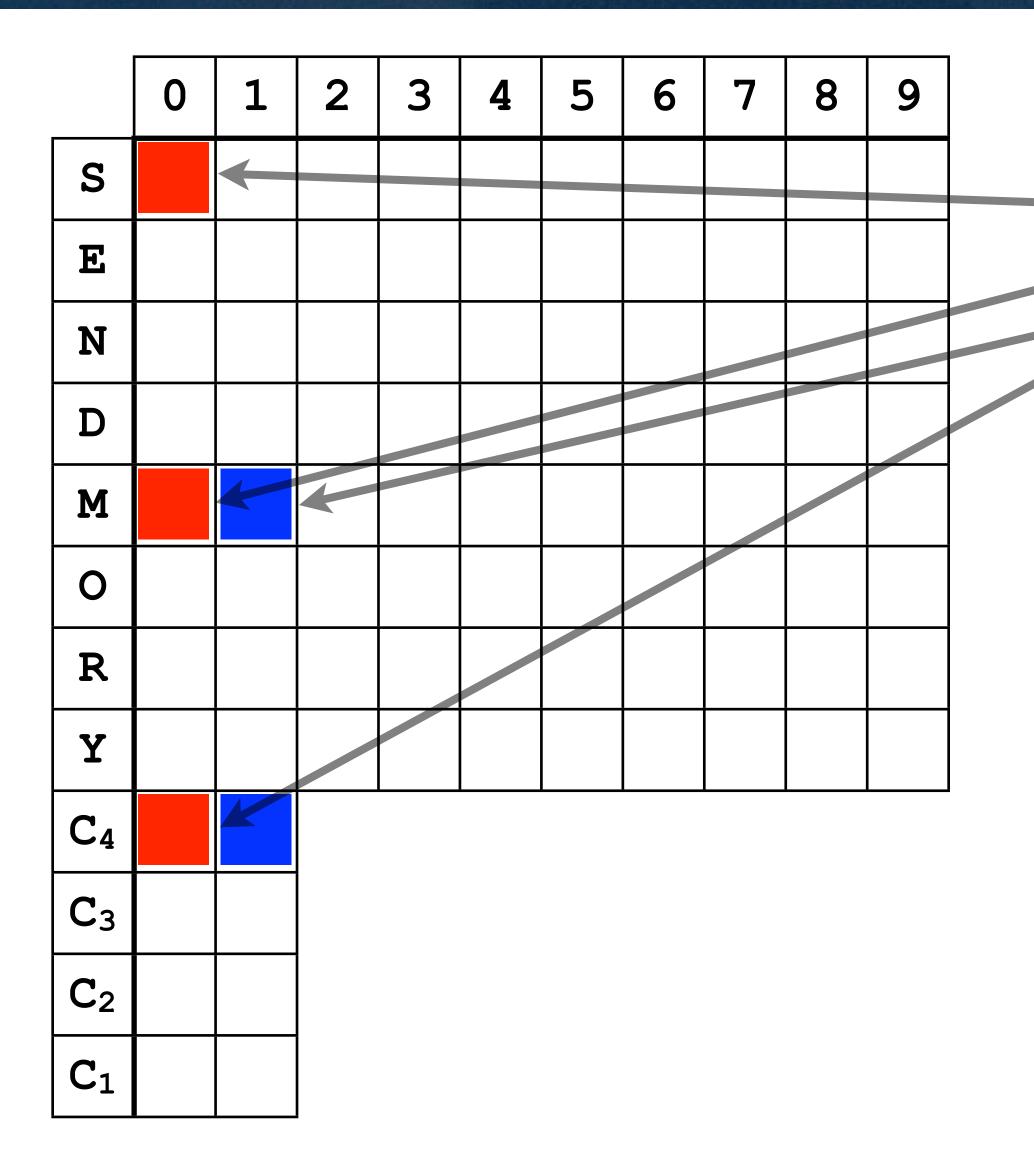
carry[4] = value[M];</pre>
```



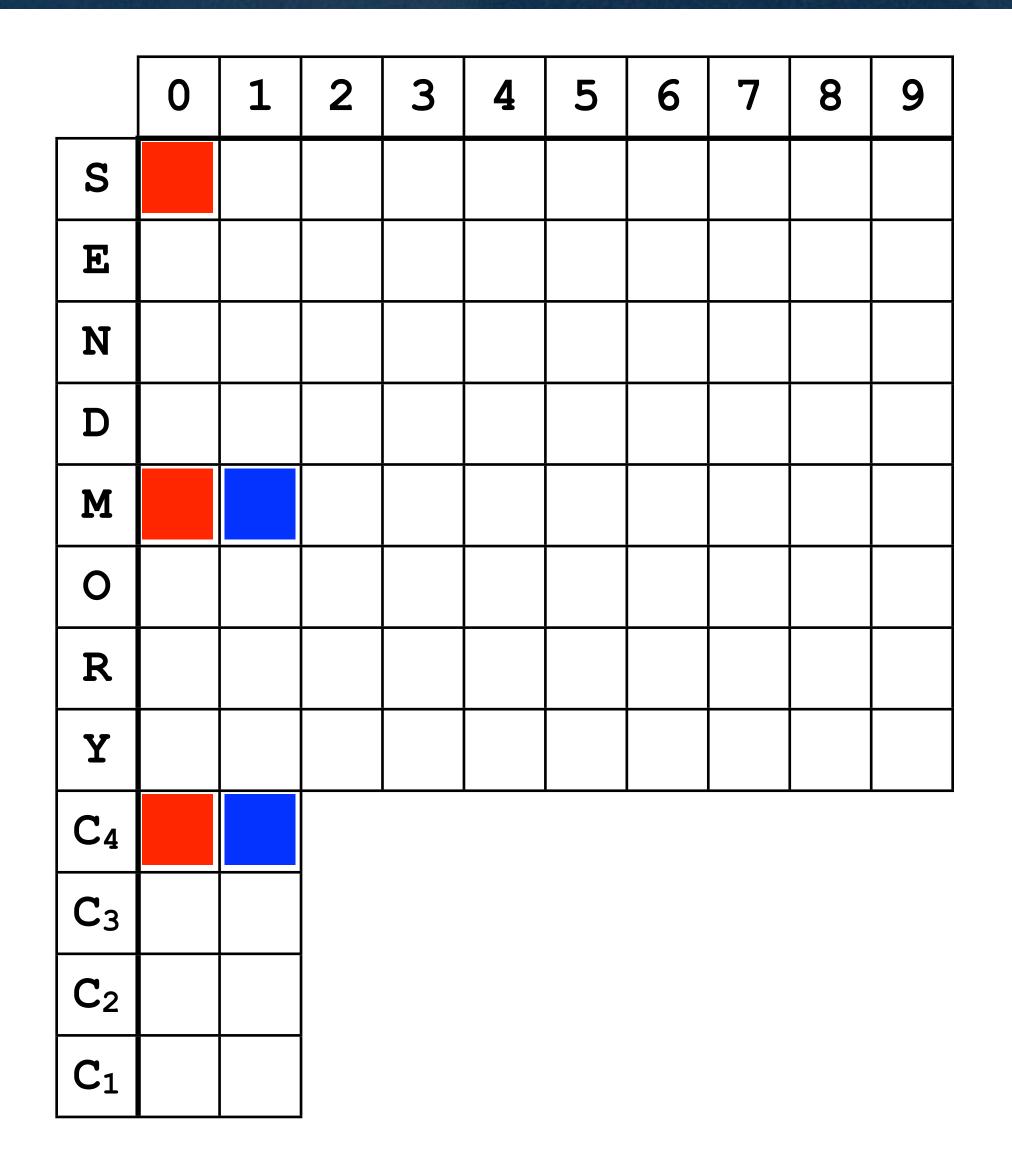
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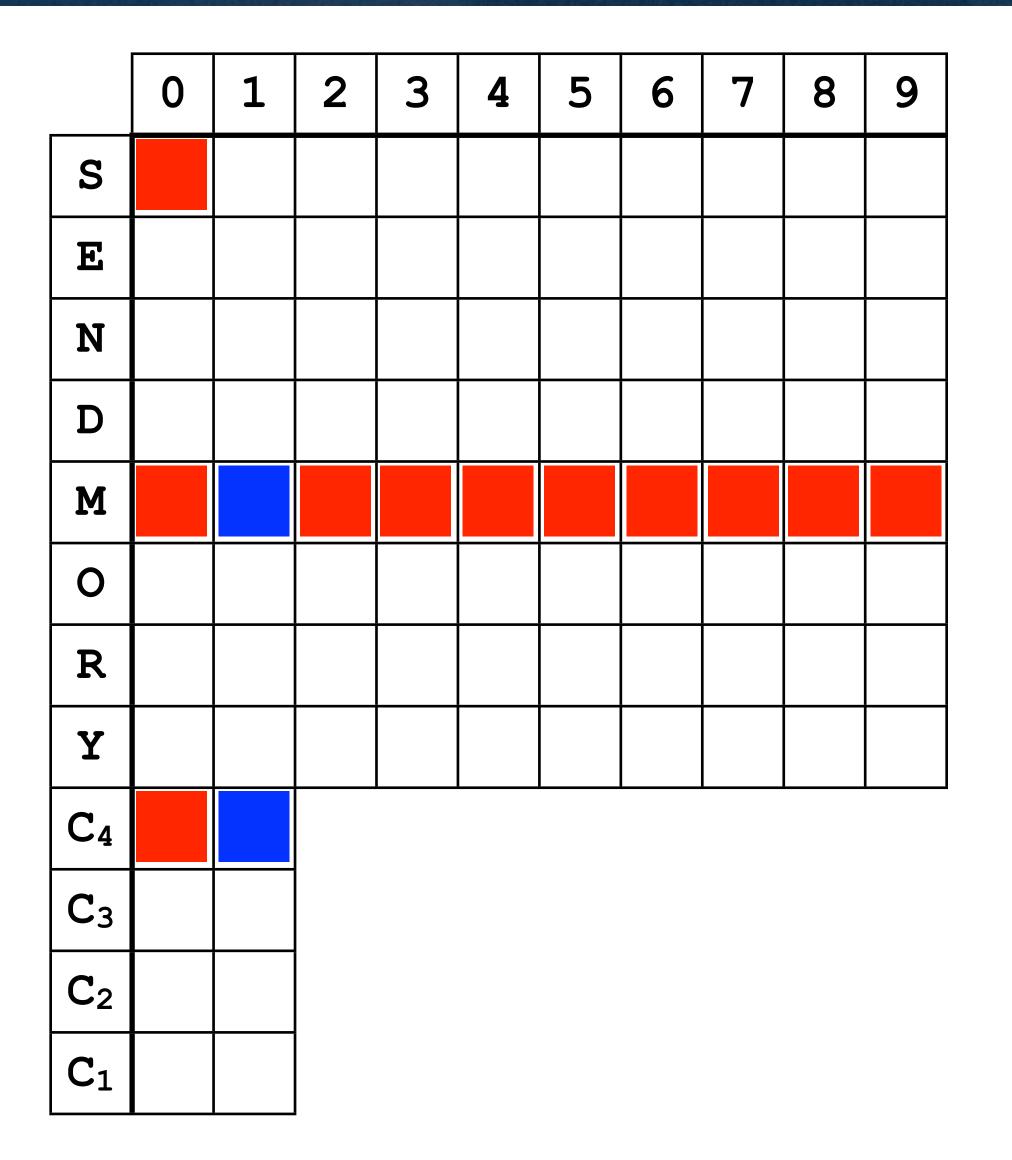
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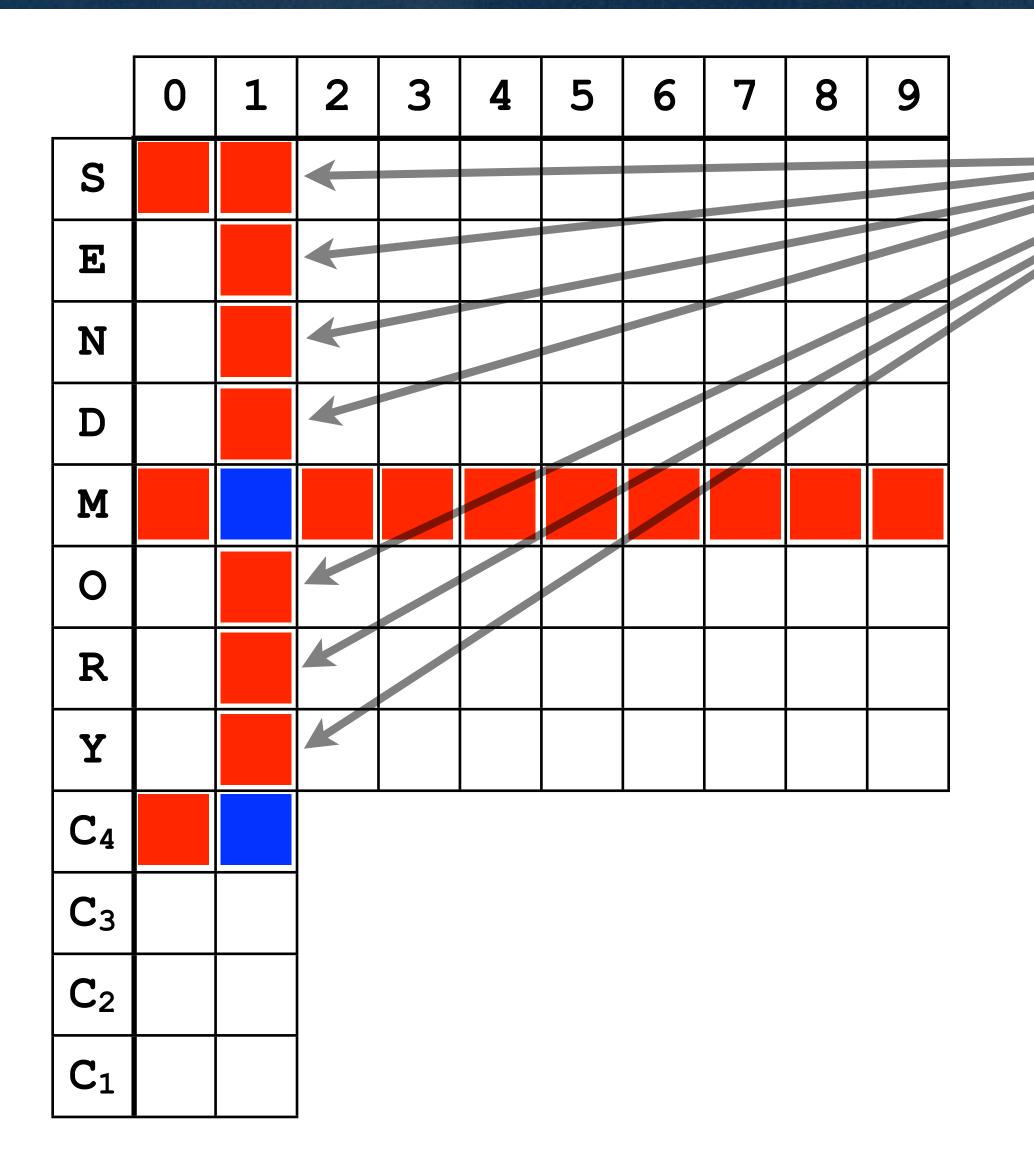
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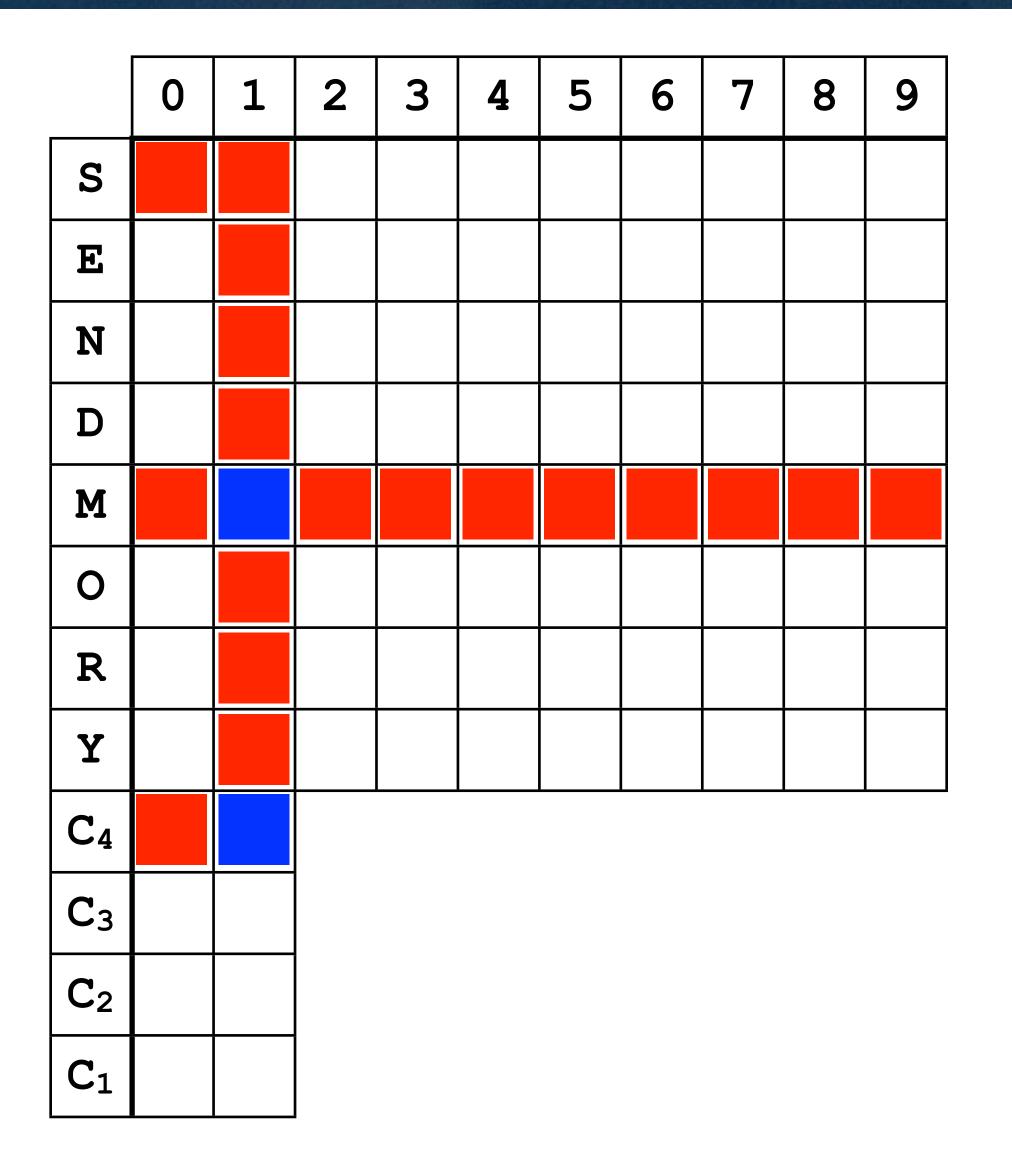
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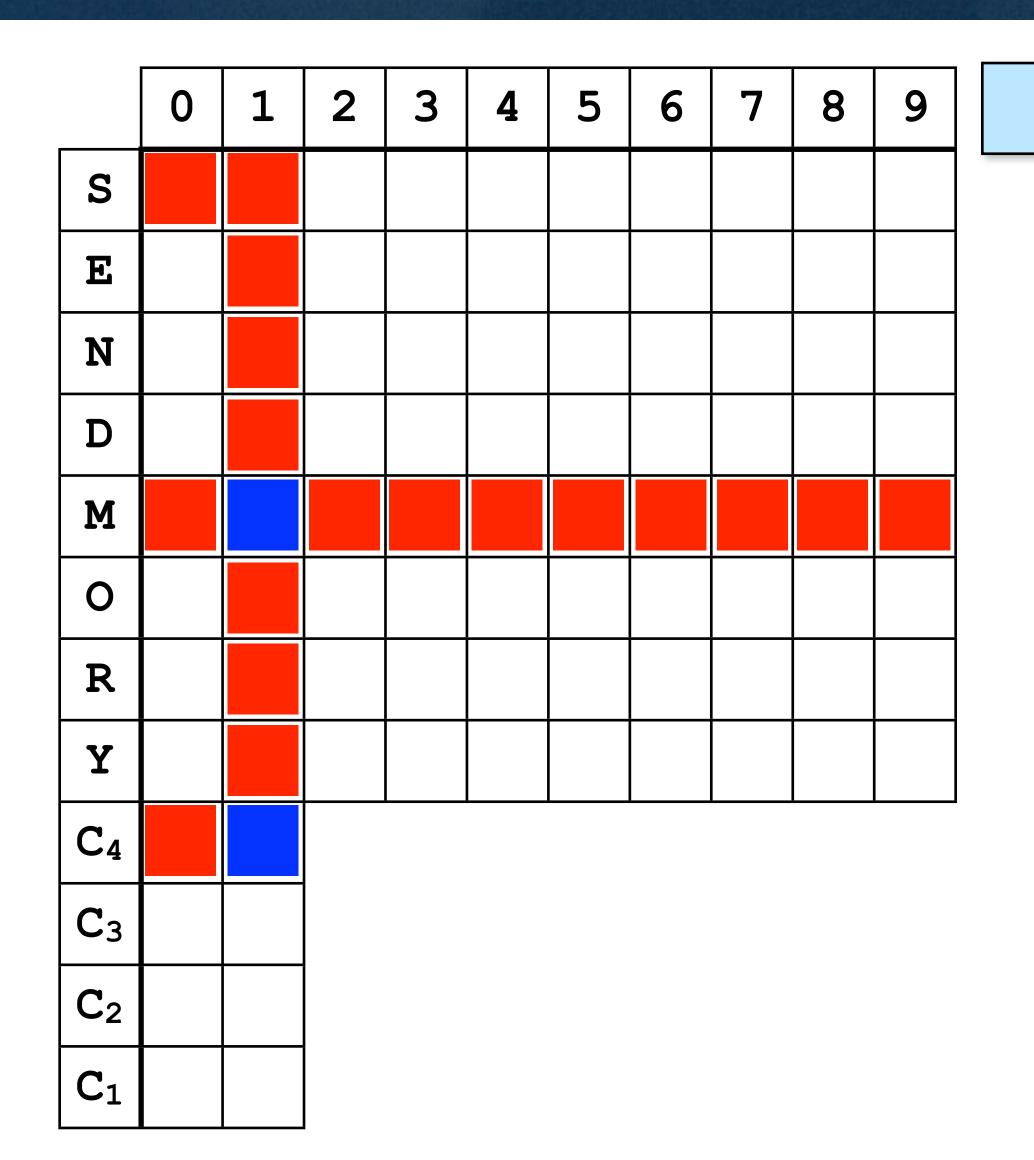
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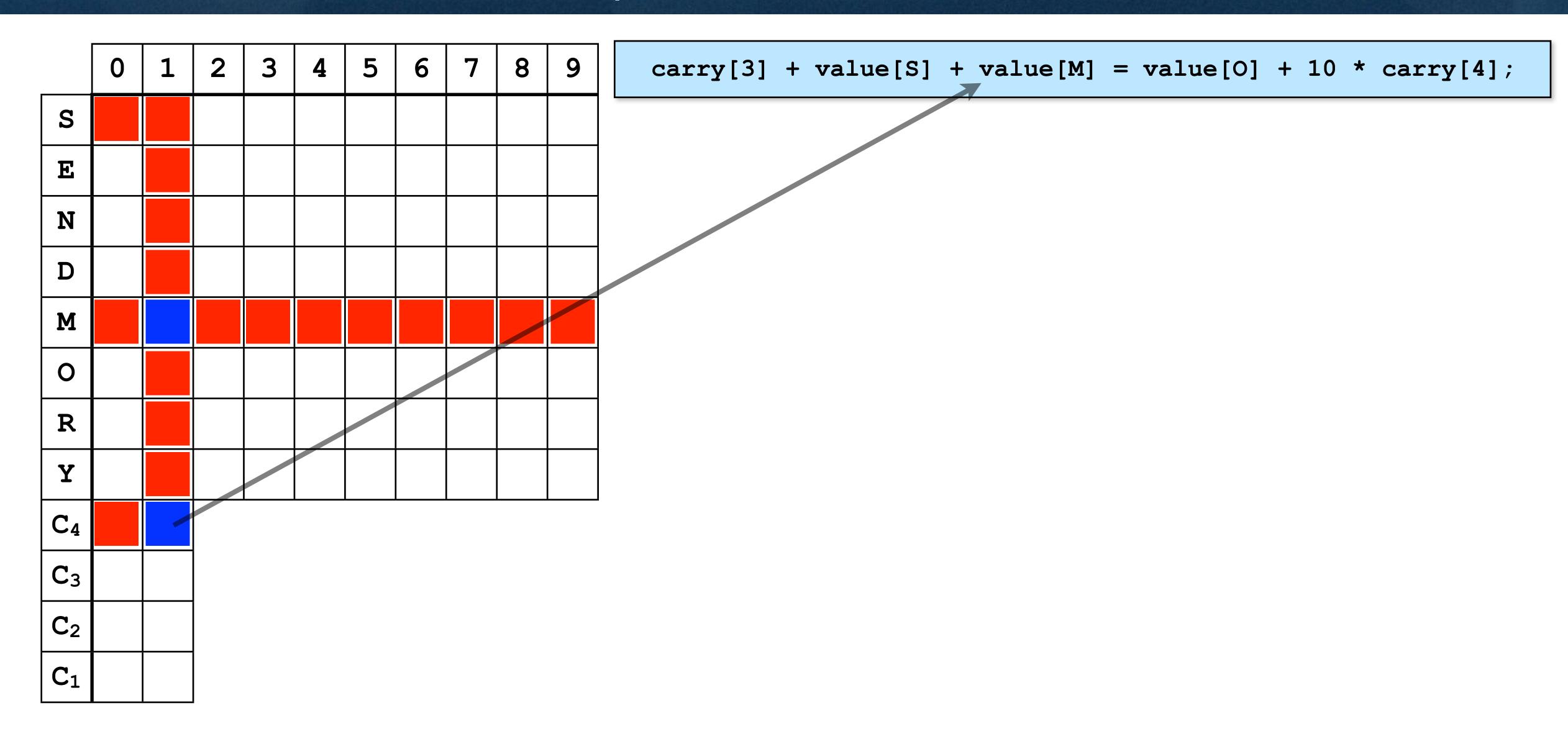
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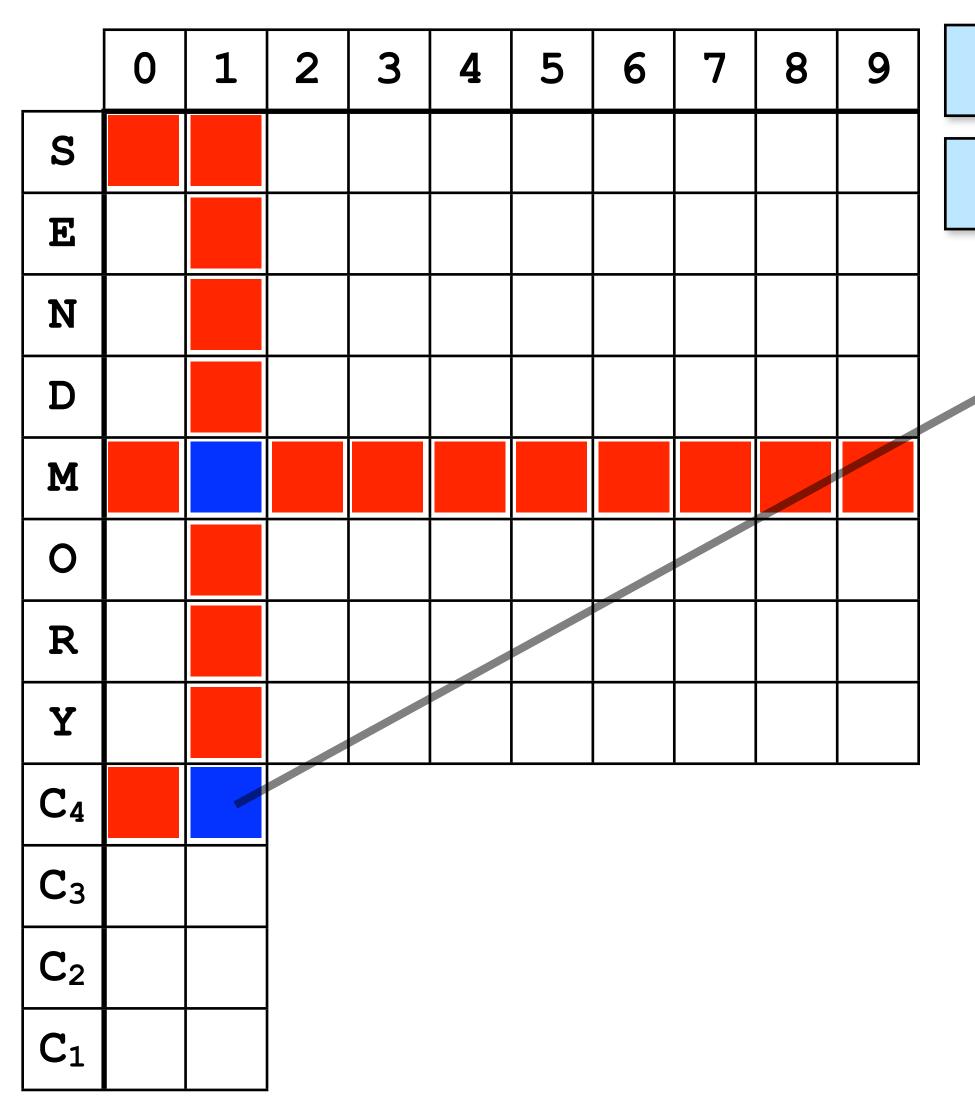
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	0	1	2	3	4	5	6	7	8	9
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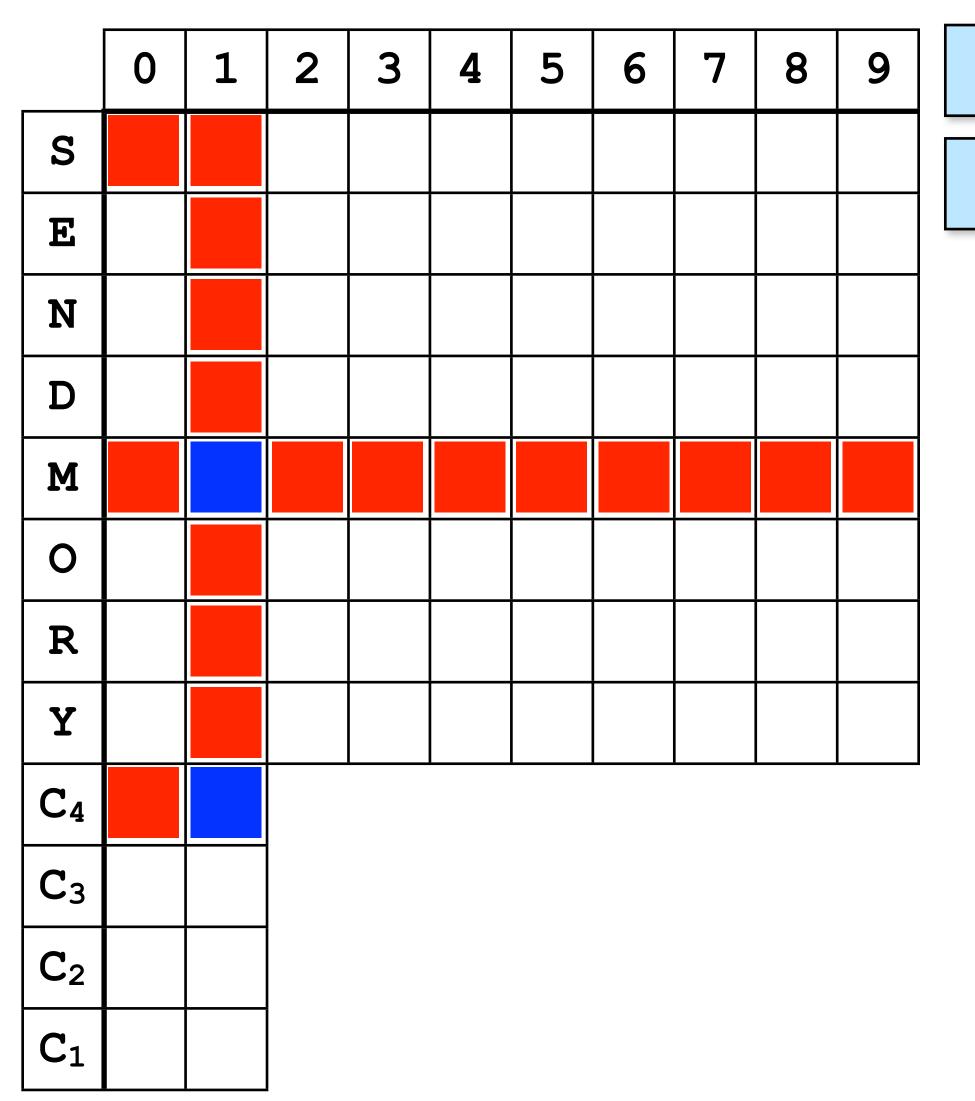


carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];

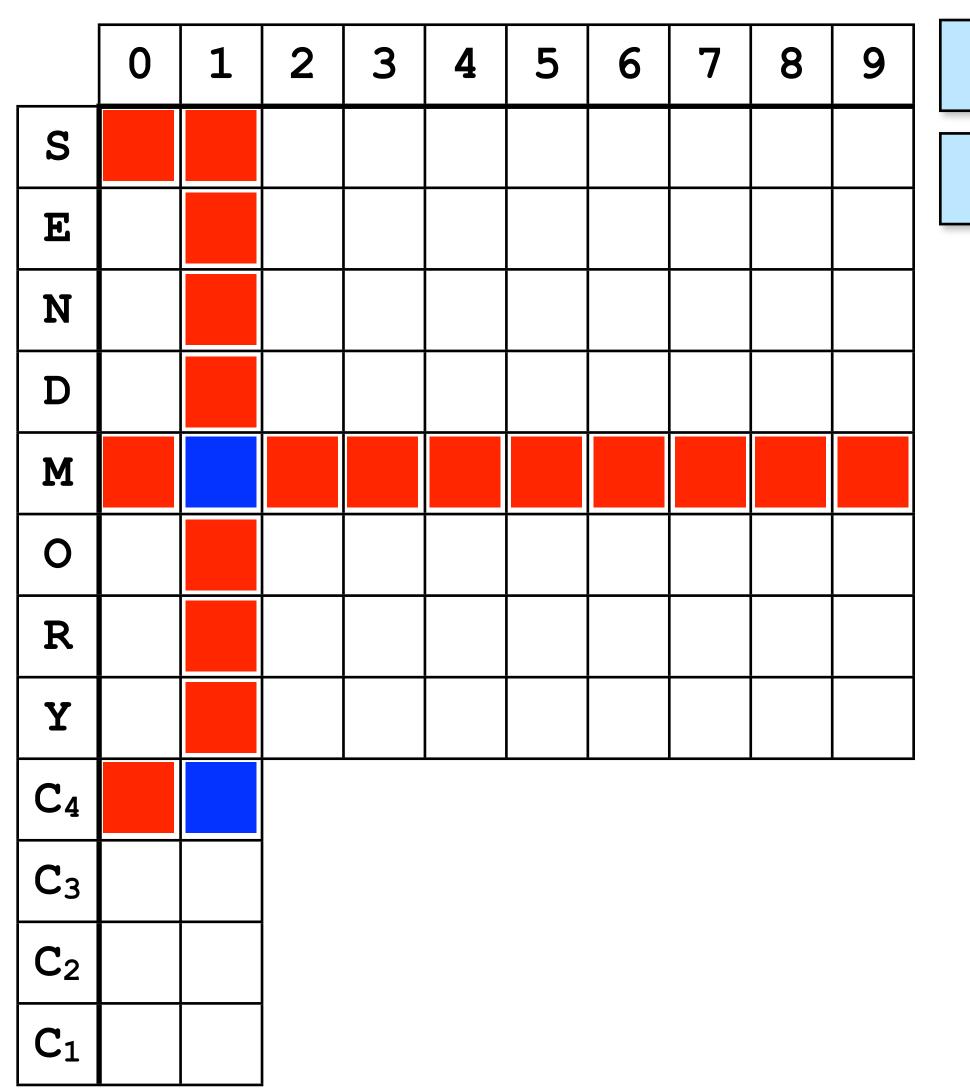


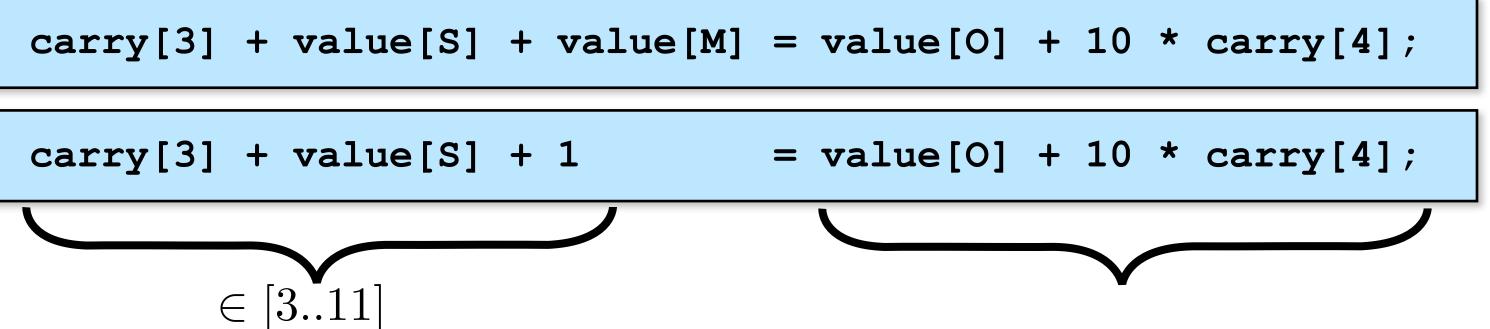


```
carry[3] + value[S] + value[M] = value[0] + 10 * carry[4];
carry[3] + value[S] + 1 = value[0] + 10 * carry[4];
```

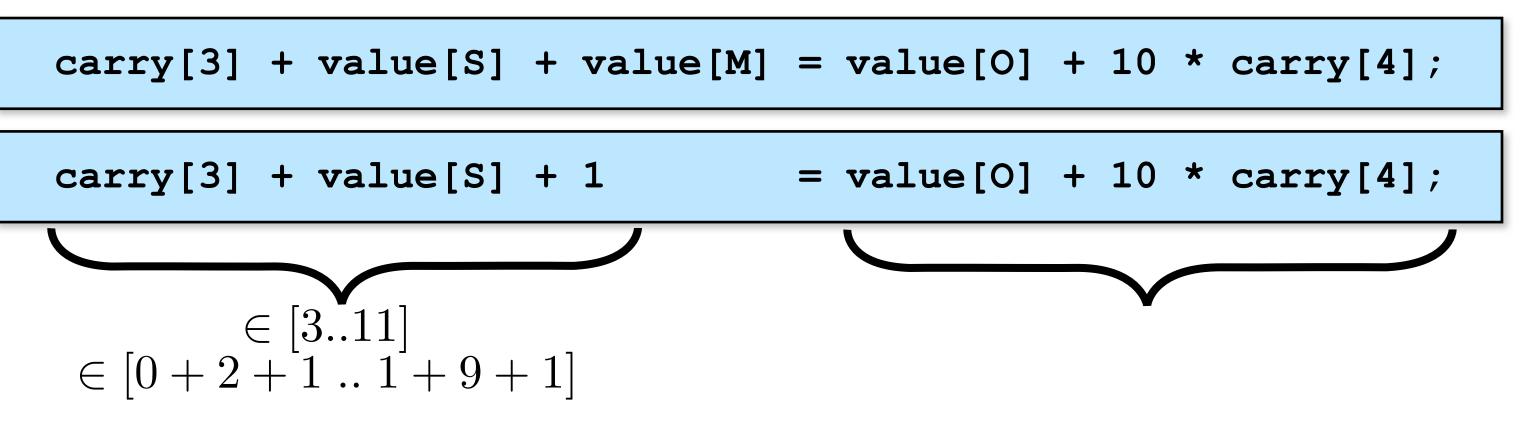


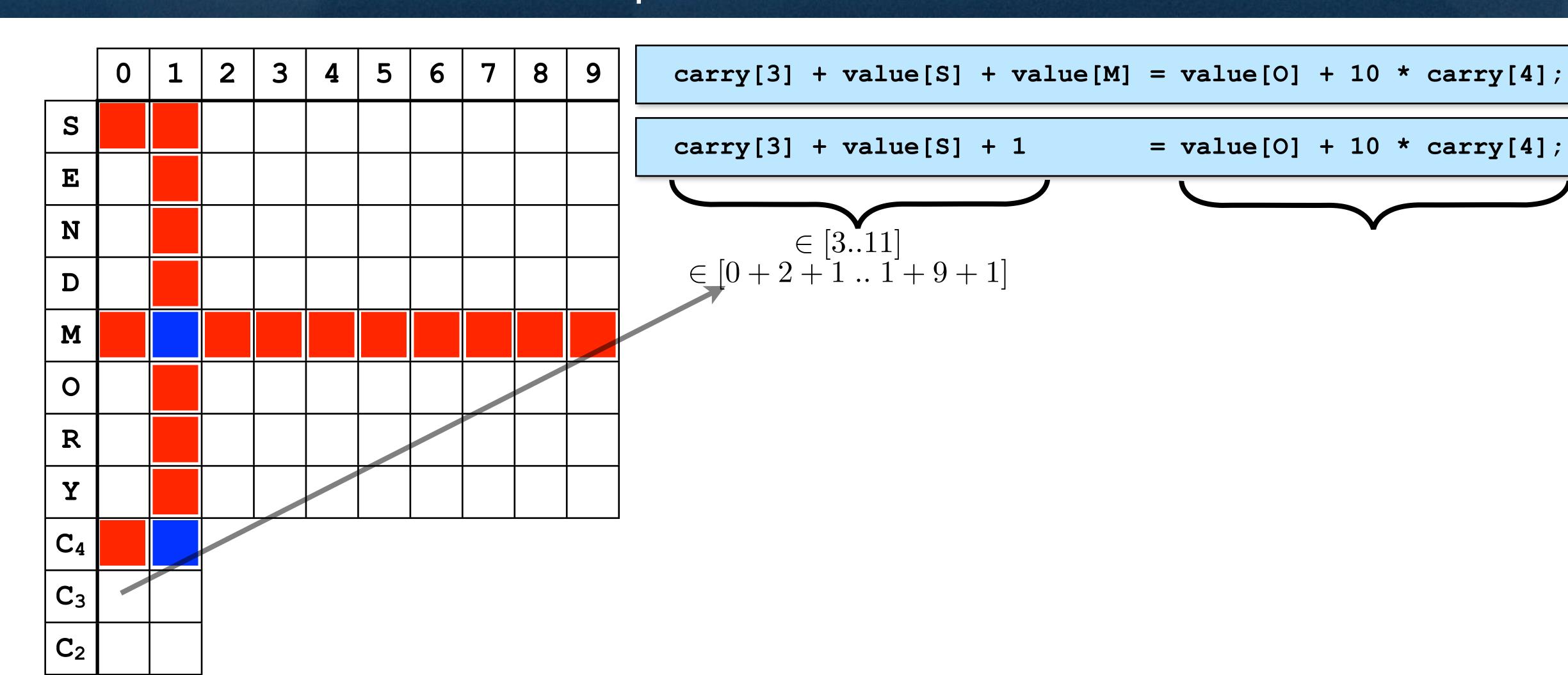
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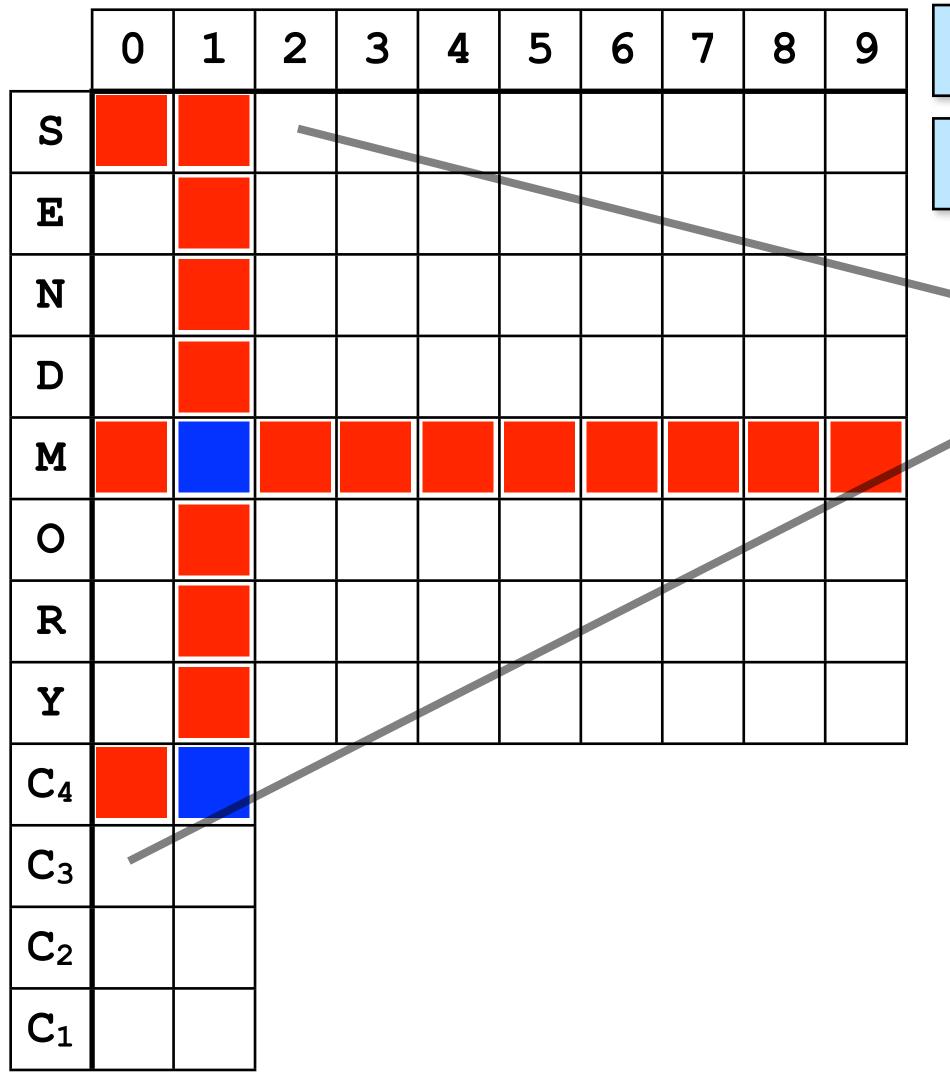


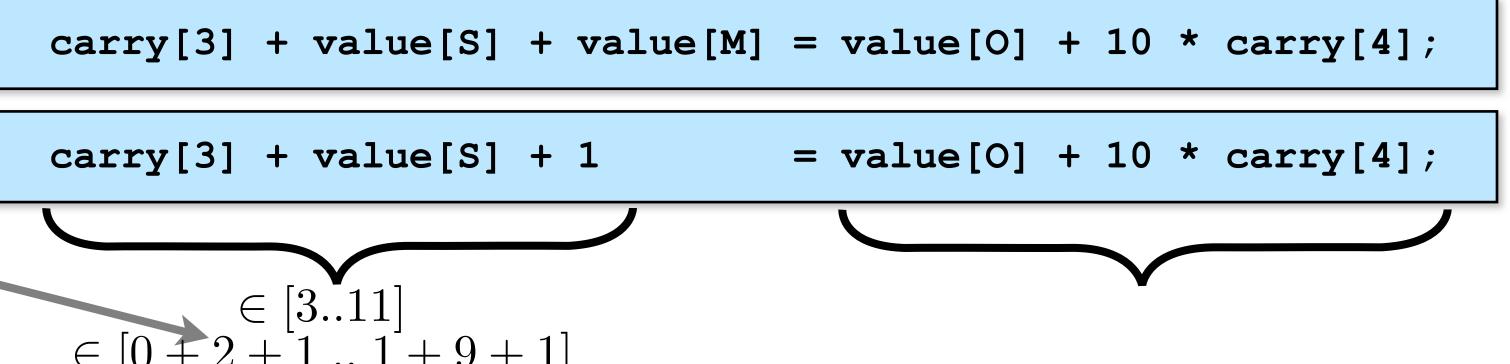


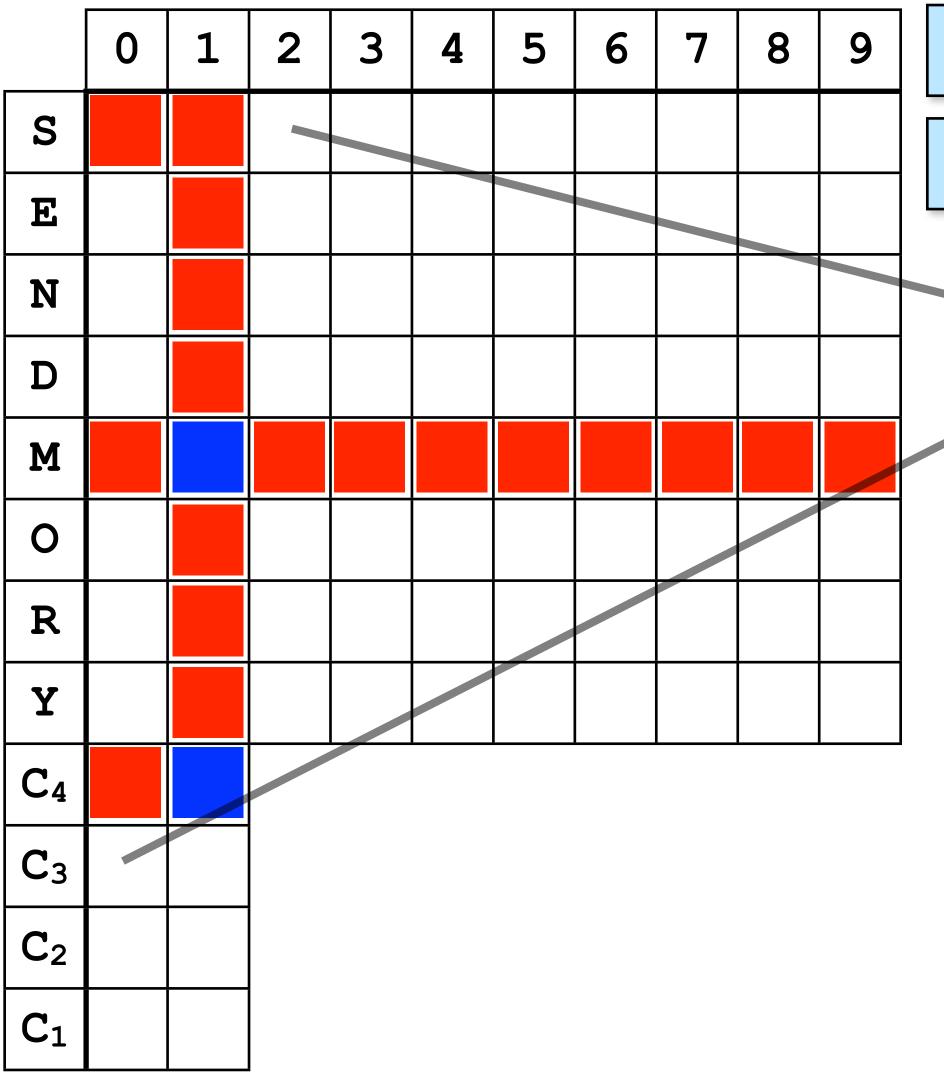
	0	1	2	3	4	5	6	7	8	9
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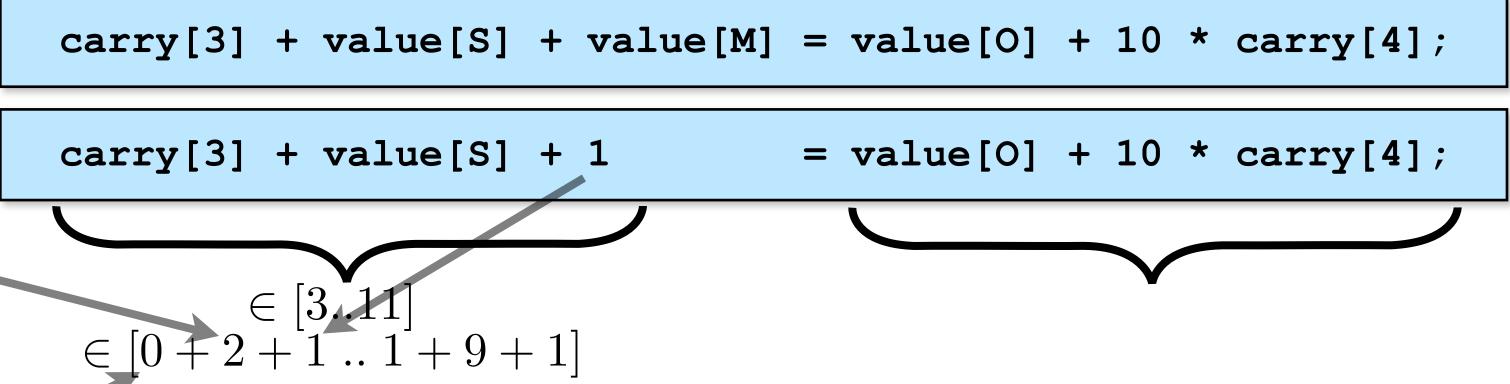


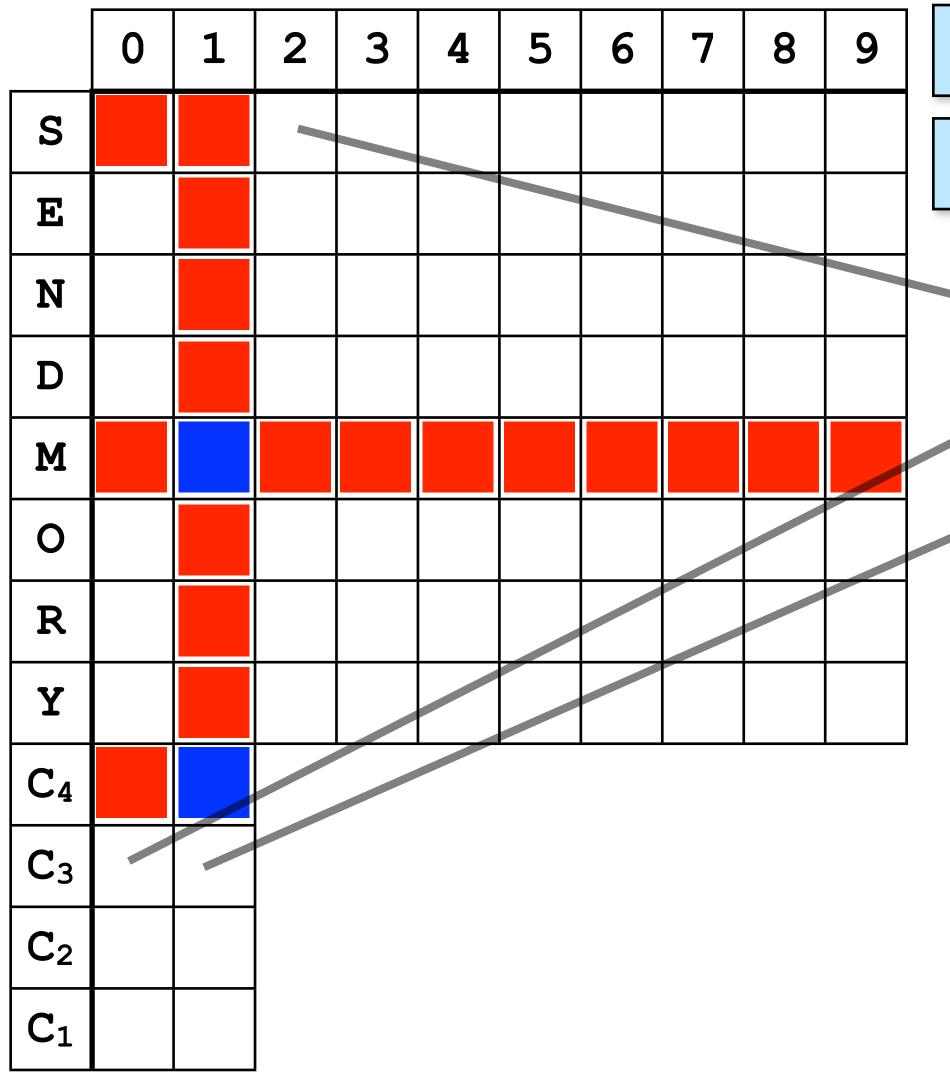


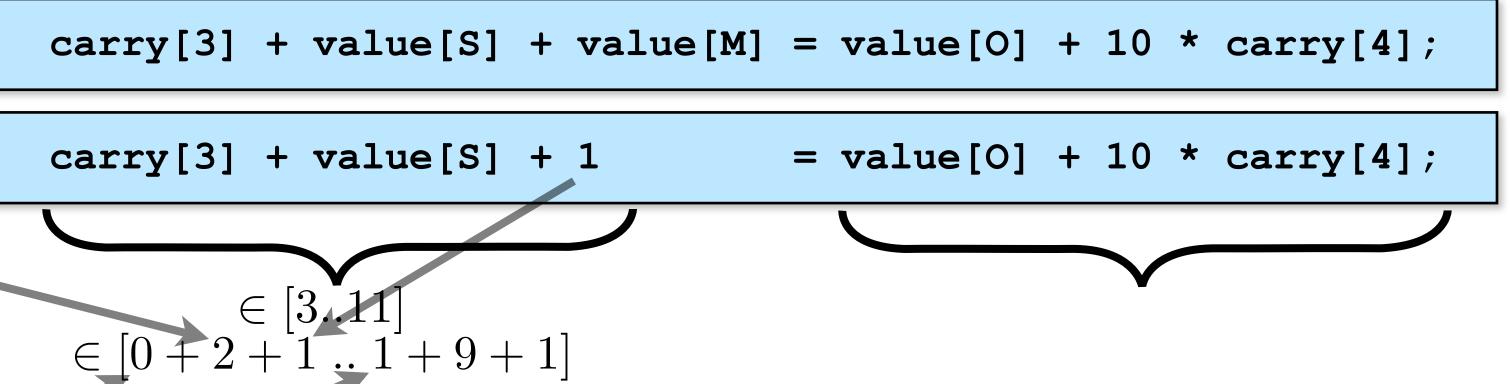


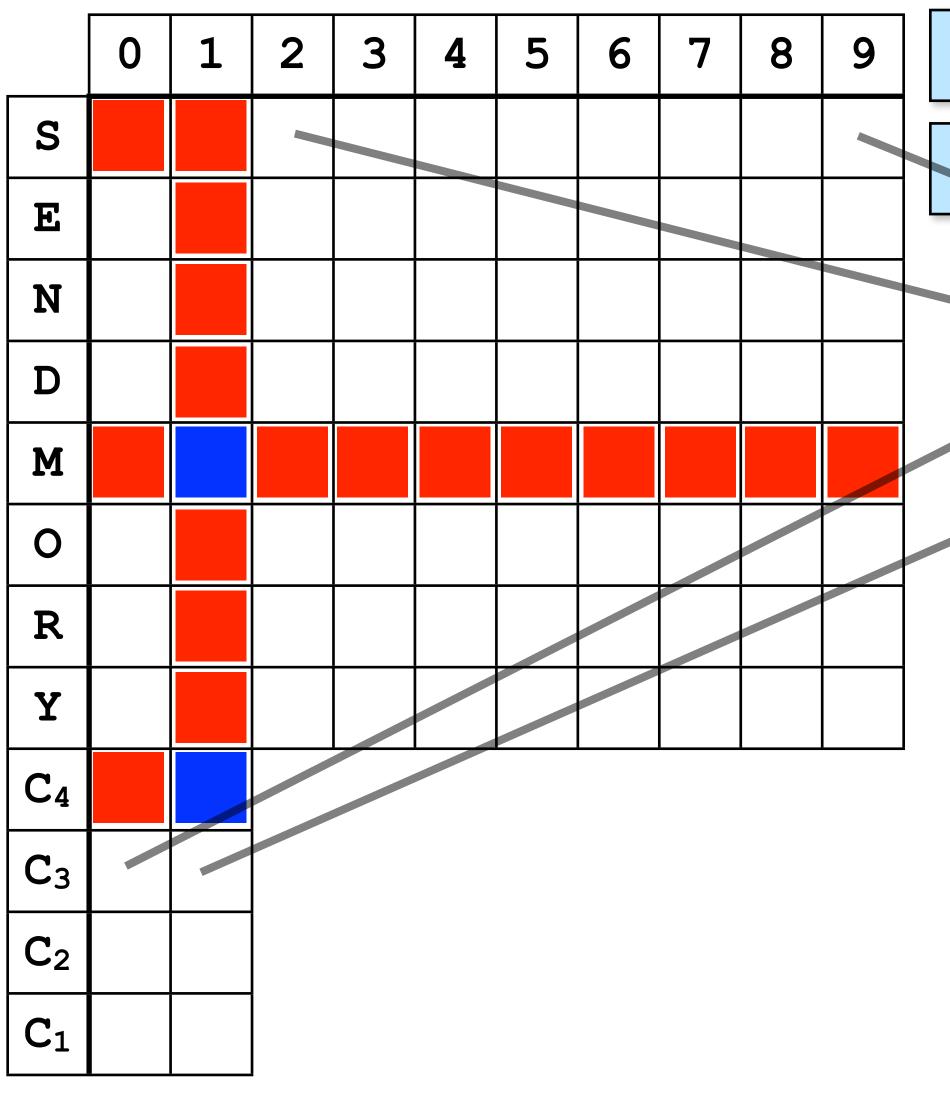




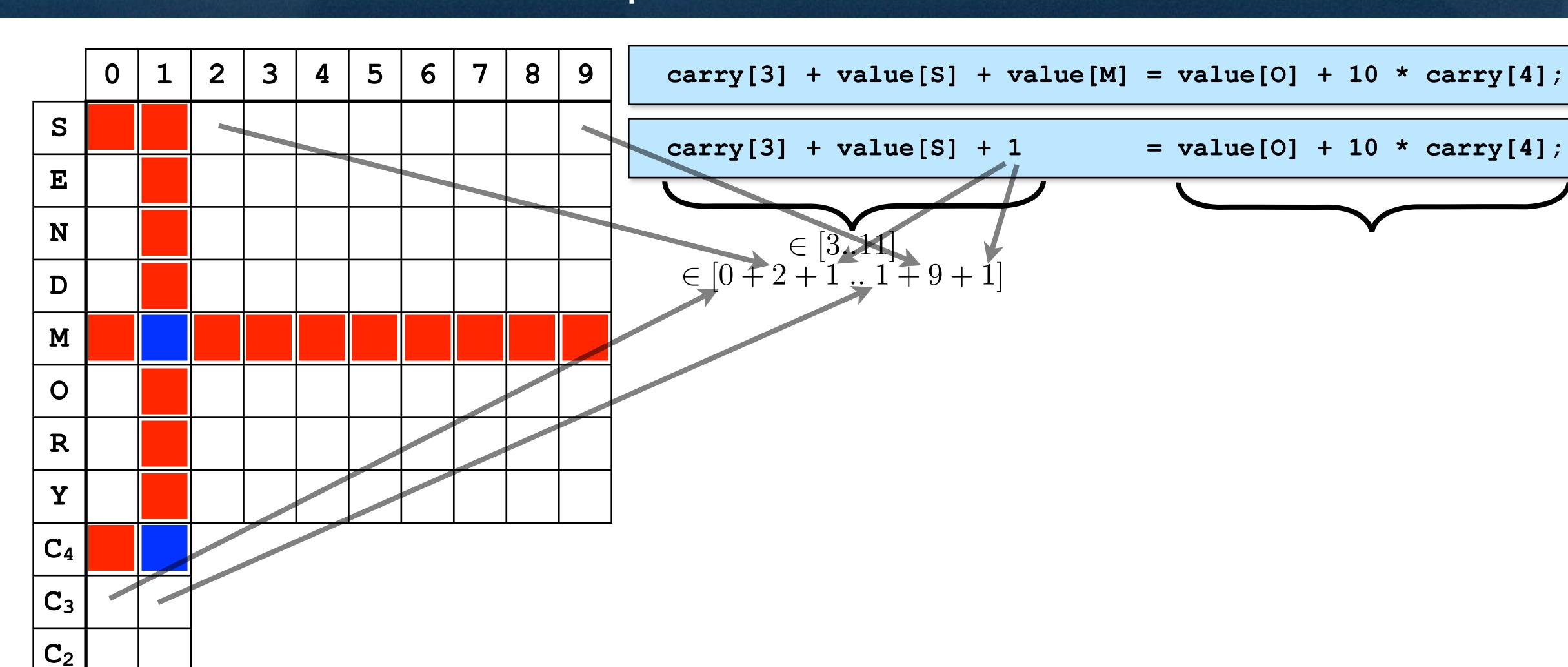


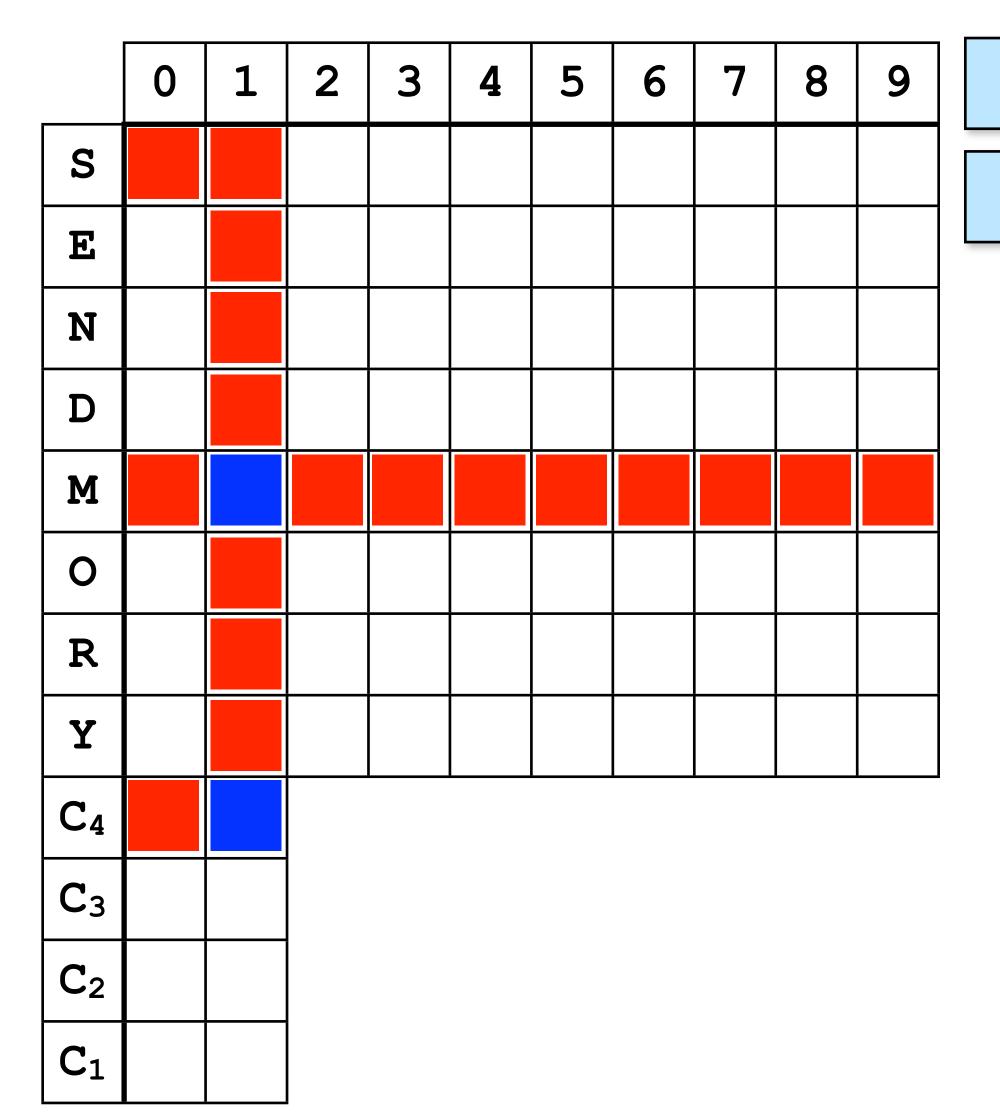


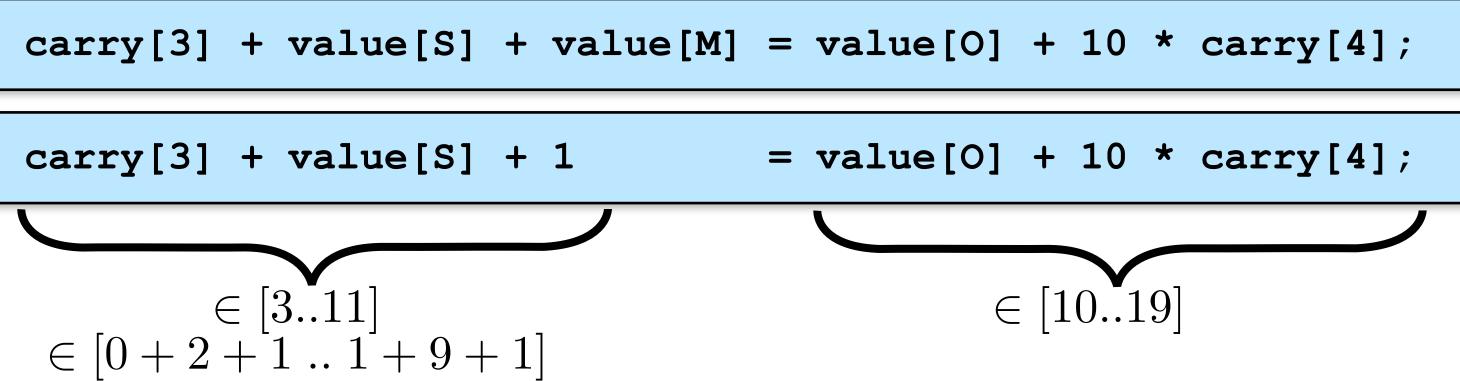


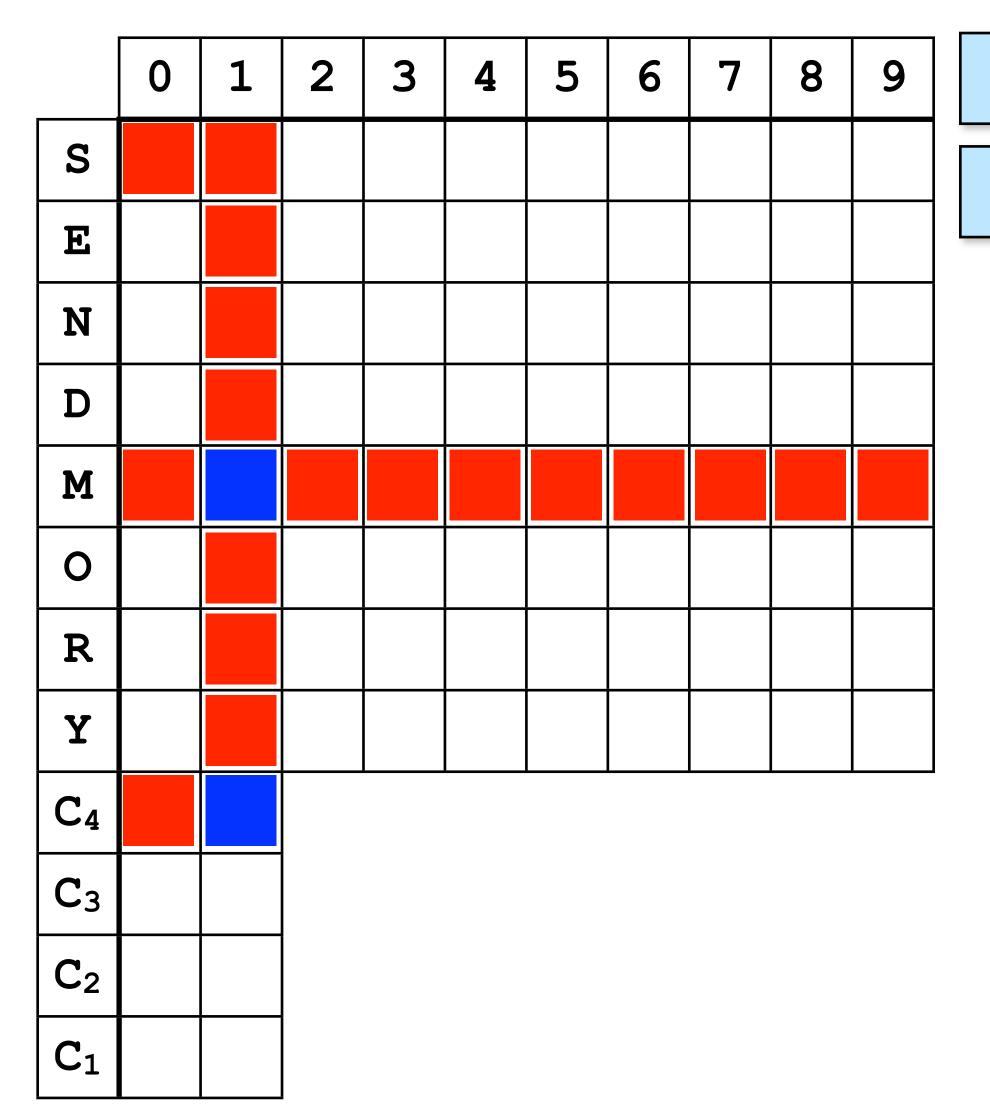


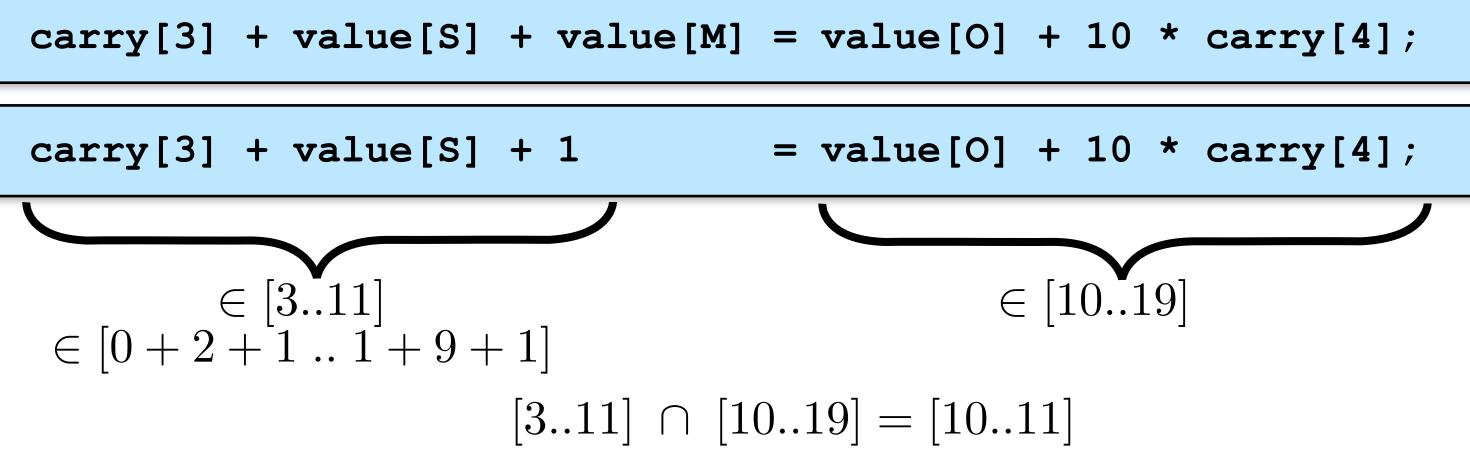
```
carry[3] + value[S] + value[M] = value[0] + 10 * carry[4];
carry[3] + value[S] + 1 = value[0] + 10 * carry[4];
```

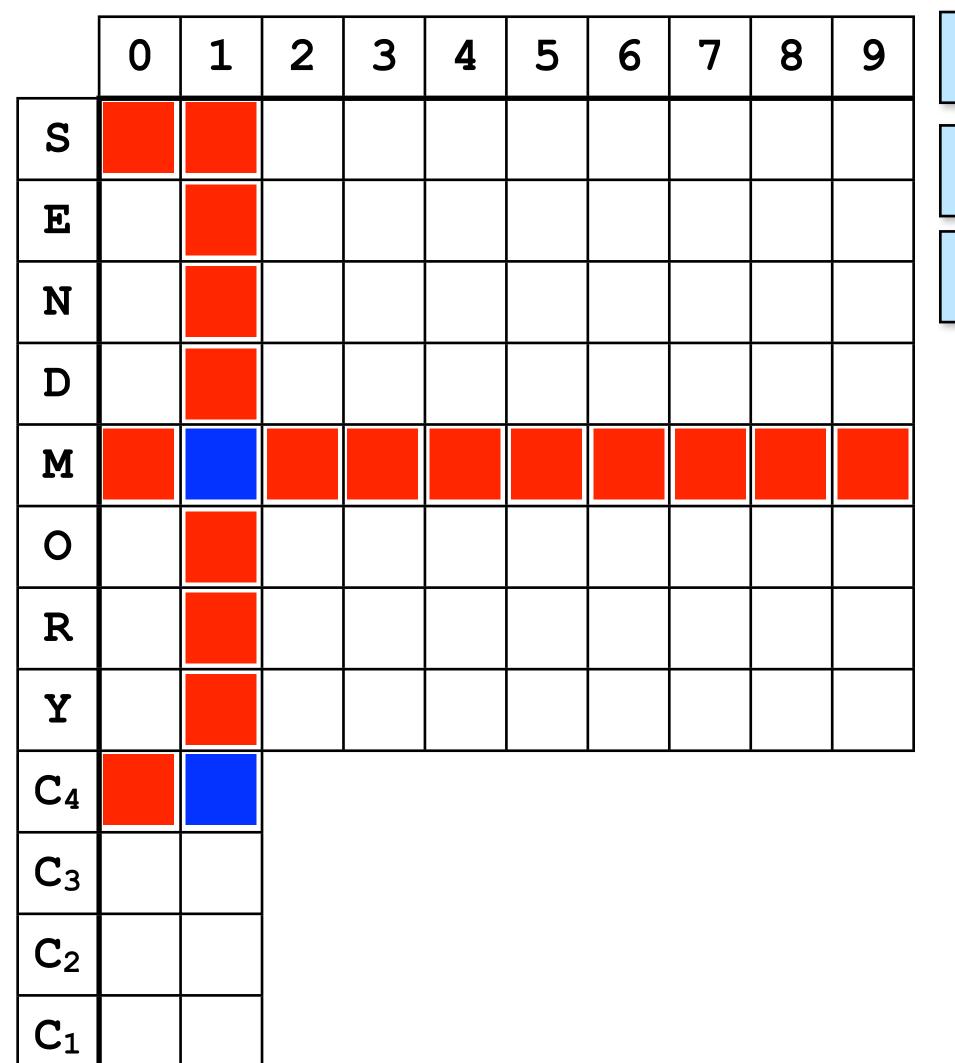












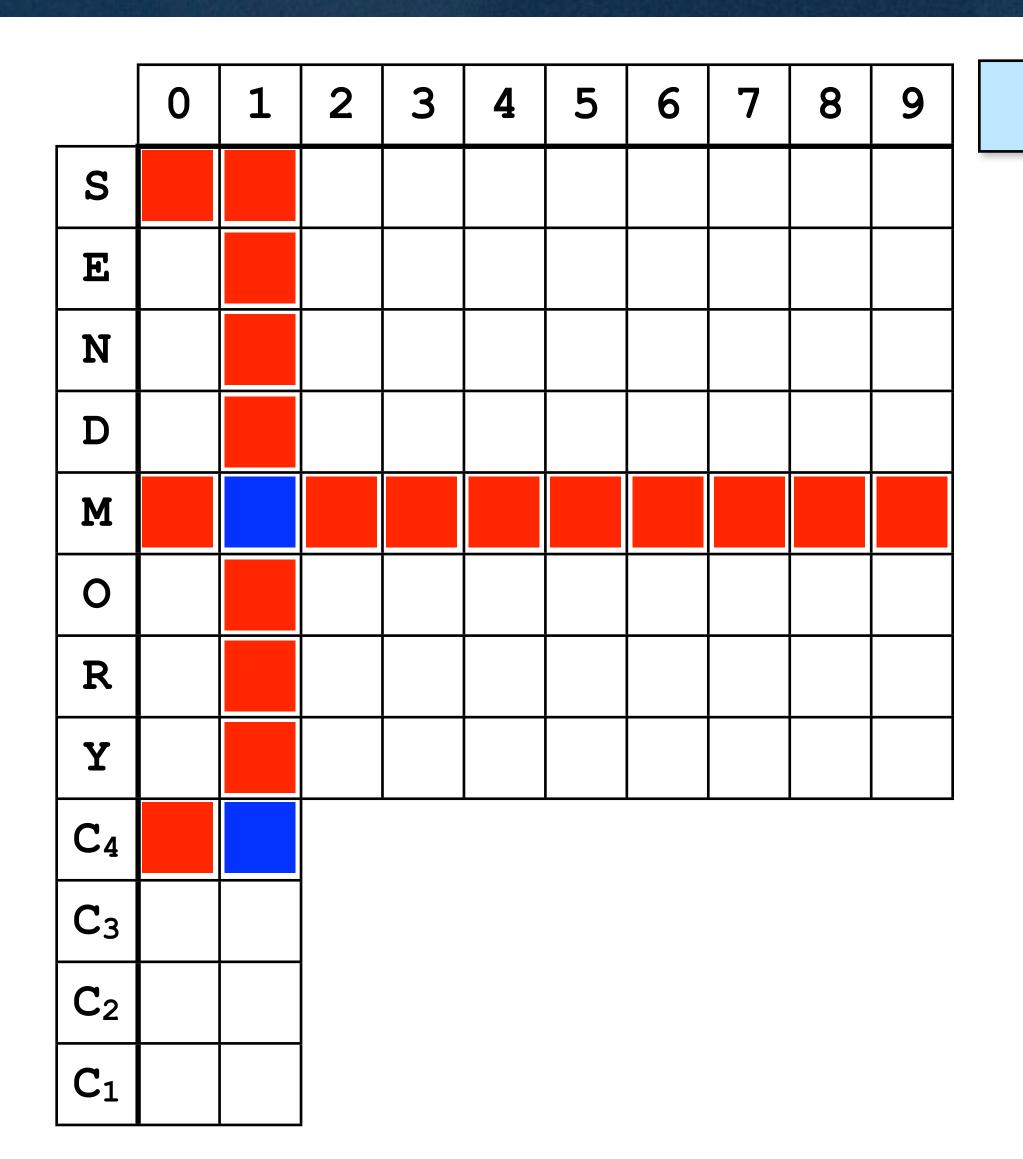
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carry[3] + value[S] + 1 = value[0] + 10 * carry[4];

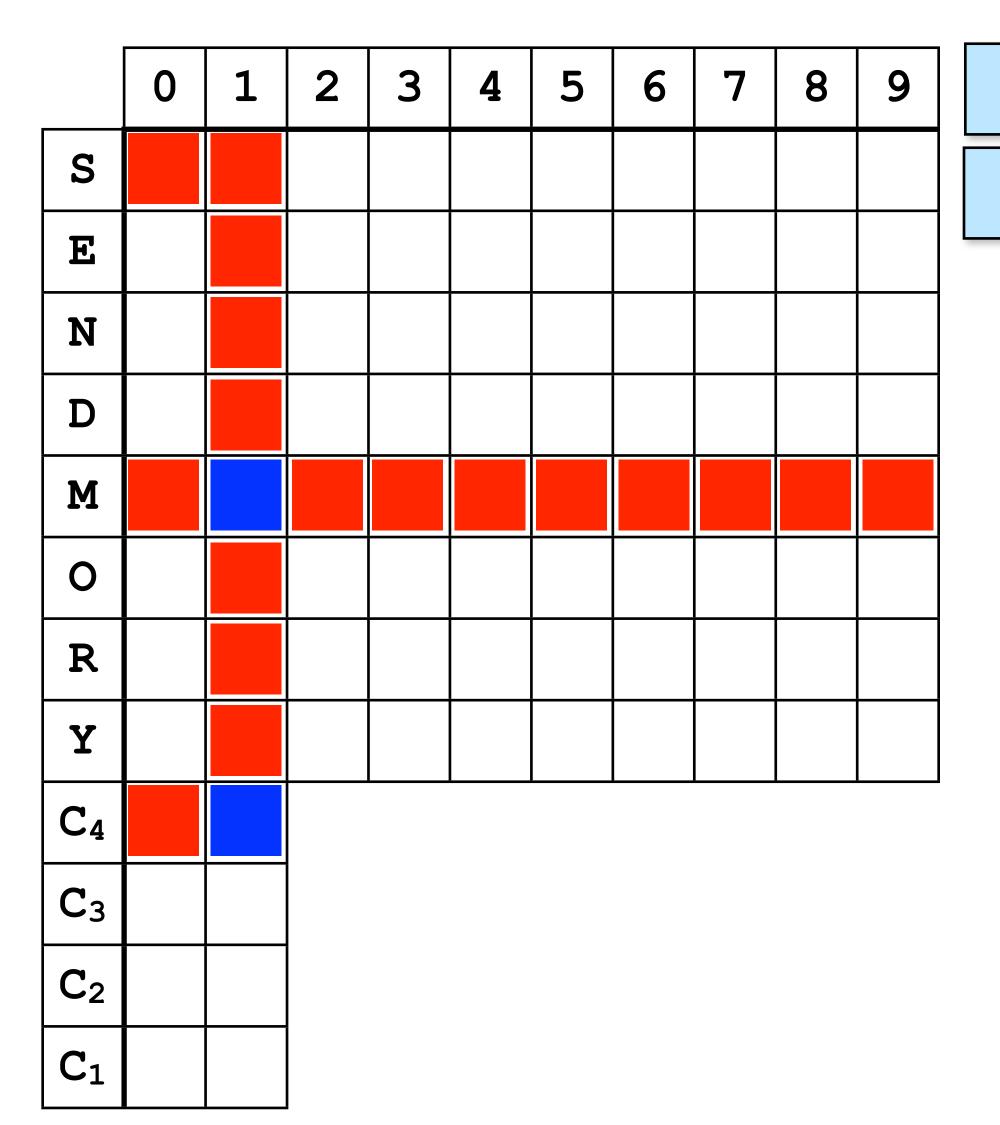
10 <= carry[3] + value[S] + 1 <= 11

\in [0+2+1..1+9+1]

= [3..11] \cap [10..19] = [10..11]
```

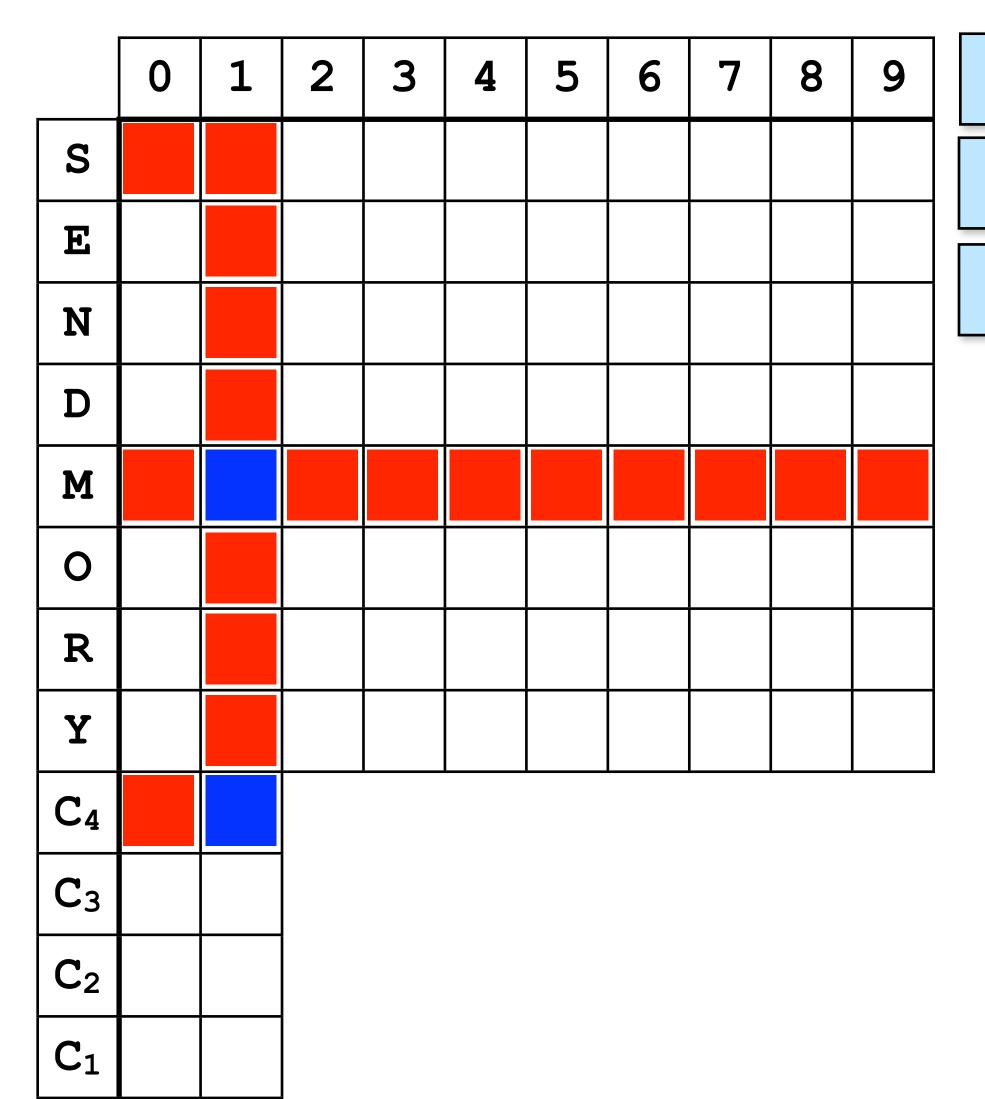


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```
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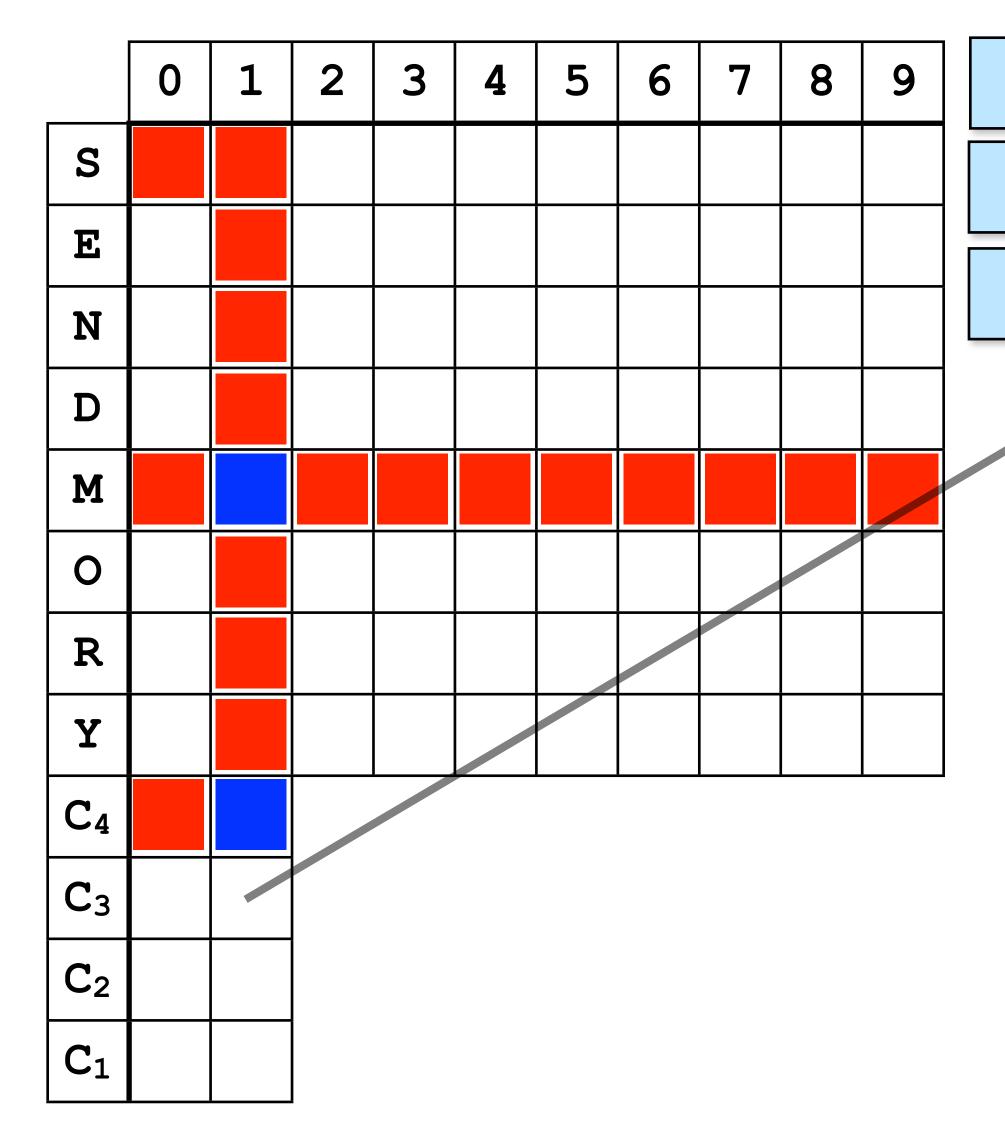
10 - 1 <= carry[3] + value[S] <= 11 - 1
```



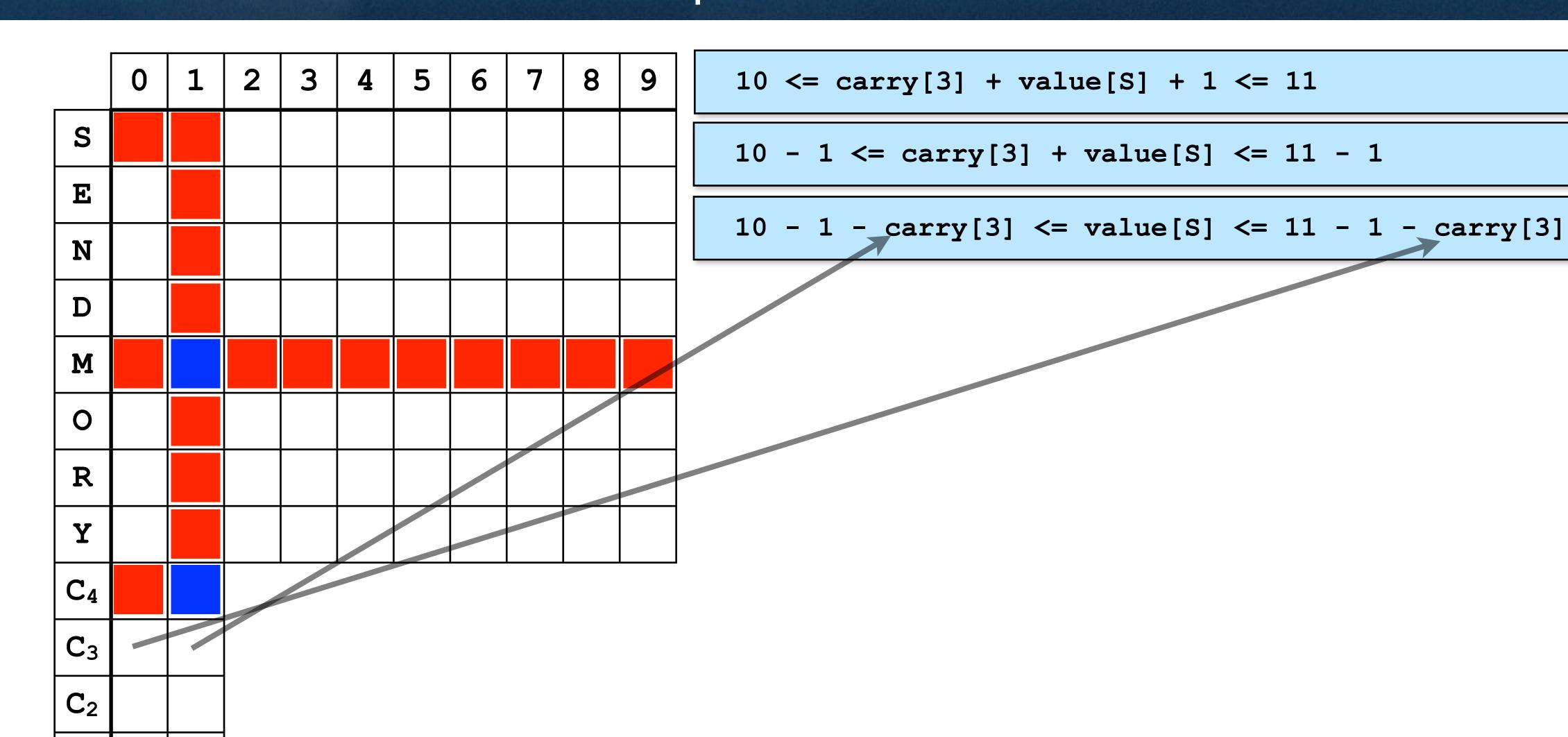
```
10 <= carry[3] + value[S] + 1 <= 11

10 - 1 <= carry[3] + value[S] <= 11 - 1

10 - 1 - carry[3] <= value[S] <= 11 - 1 - carry[3]
```



```
10 <= carry[3] + value[S] + 1 <= 11
10 - 1 <= carry[3] + value[S] <= 11 - 1
10 - 1 - carry[3] <= value[S] <= 11 - 1 - carry[3]</pre>
```



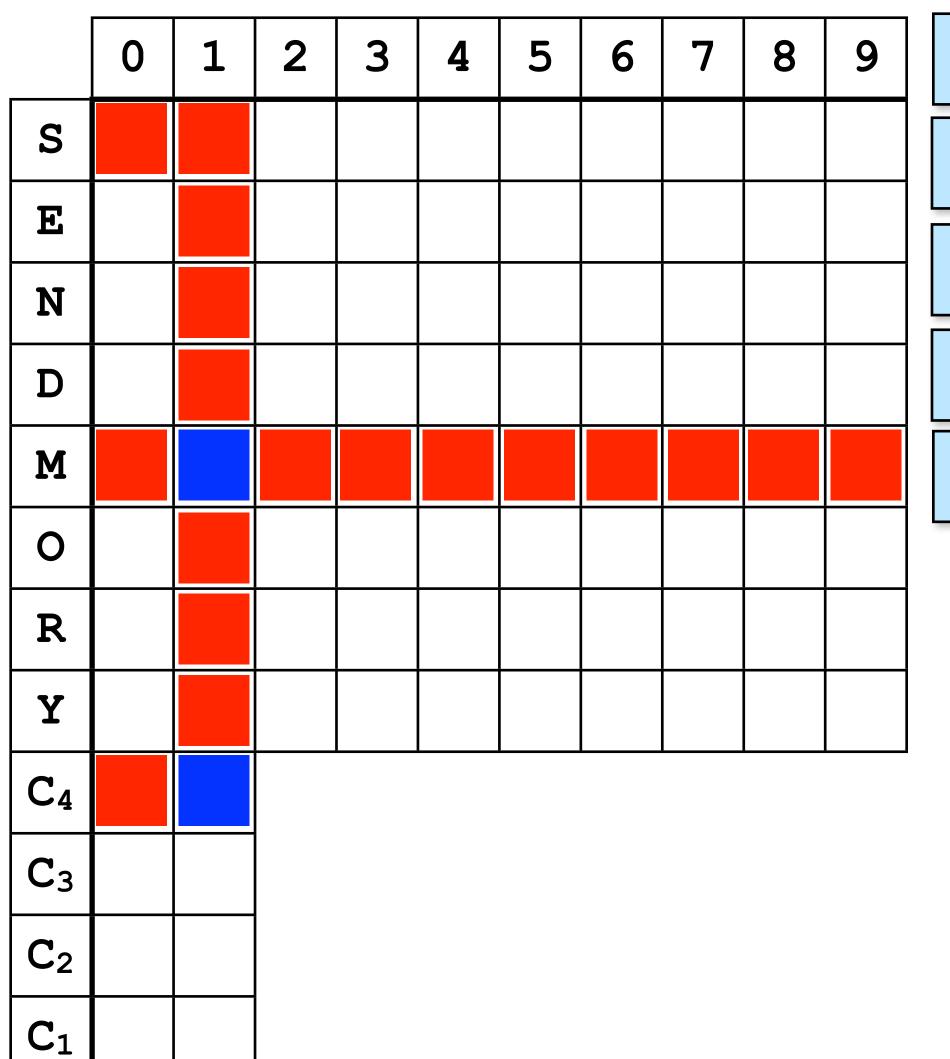
	0	1	2	3	4	5	6	7	8	9	
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C ₄											I
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```
10 <= carry[3] + value[S] + 1 <= 11

10 - 1 <= carry[3] + value[S] <= 11 - 1

10 - 1 - carry[3] <= value[S] <= 11 - 1 - carry[3]

10 - 1 - 1 <= value[S] <= 11 - 1 - 0
```



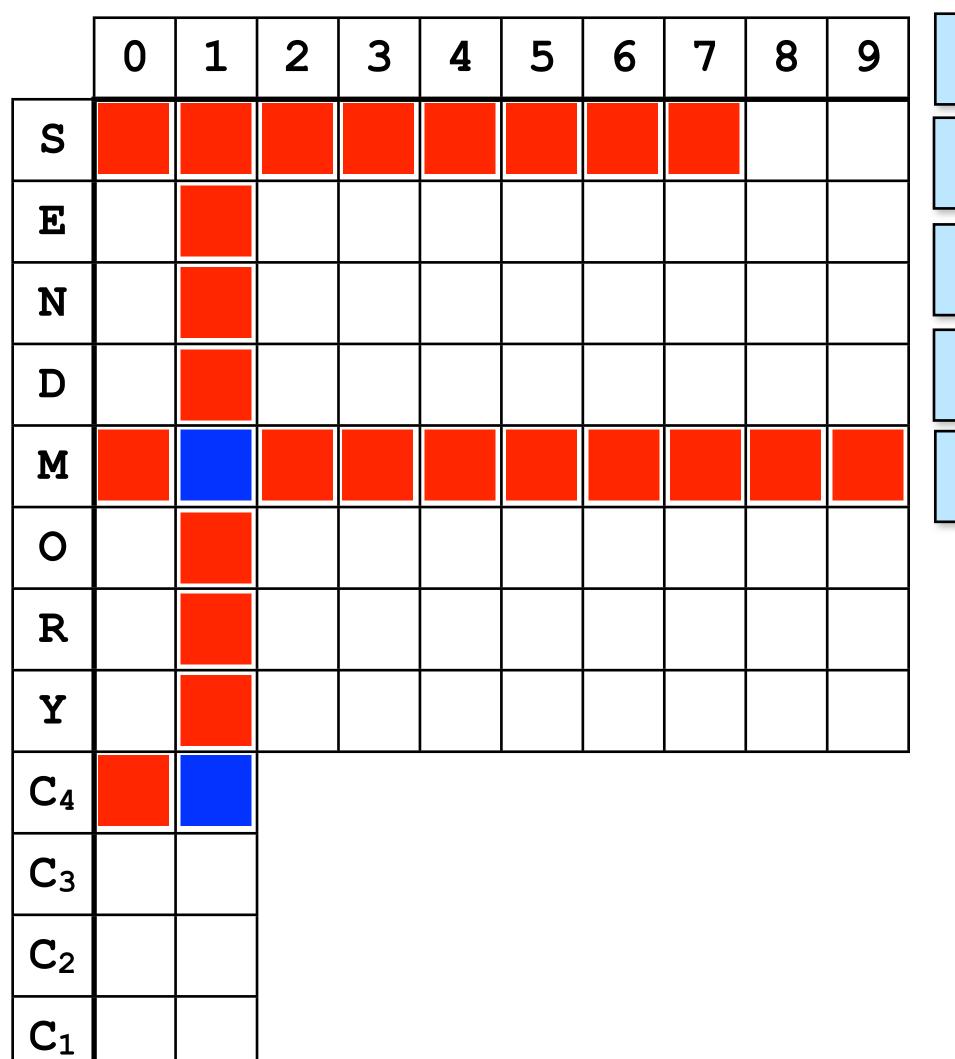
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10 <= carry[3] + value[S] + 1 <= 11

10 - 1 <= carry[3] + value[S] <= 11 - 1

10 - 1 - carry[3] <= value[S] <= 11 - 1 - carry[3]

10 - 1 - 1 <= value[S] <= 11 - 1 - 0

8 <= value[S] <= 10
```



```
10 <= carry[3] + value[S] + 1 <= 11

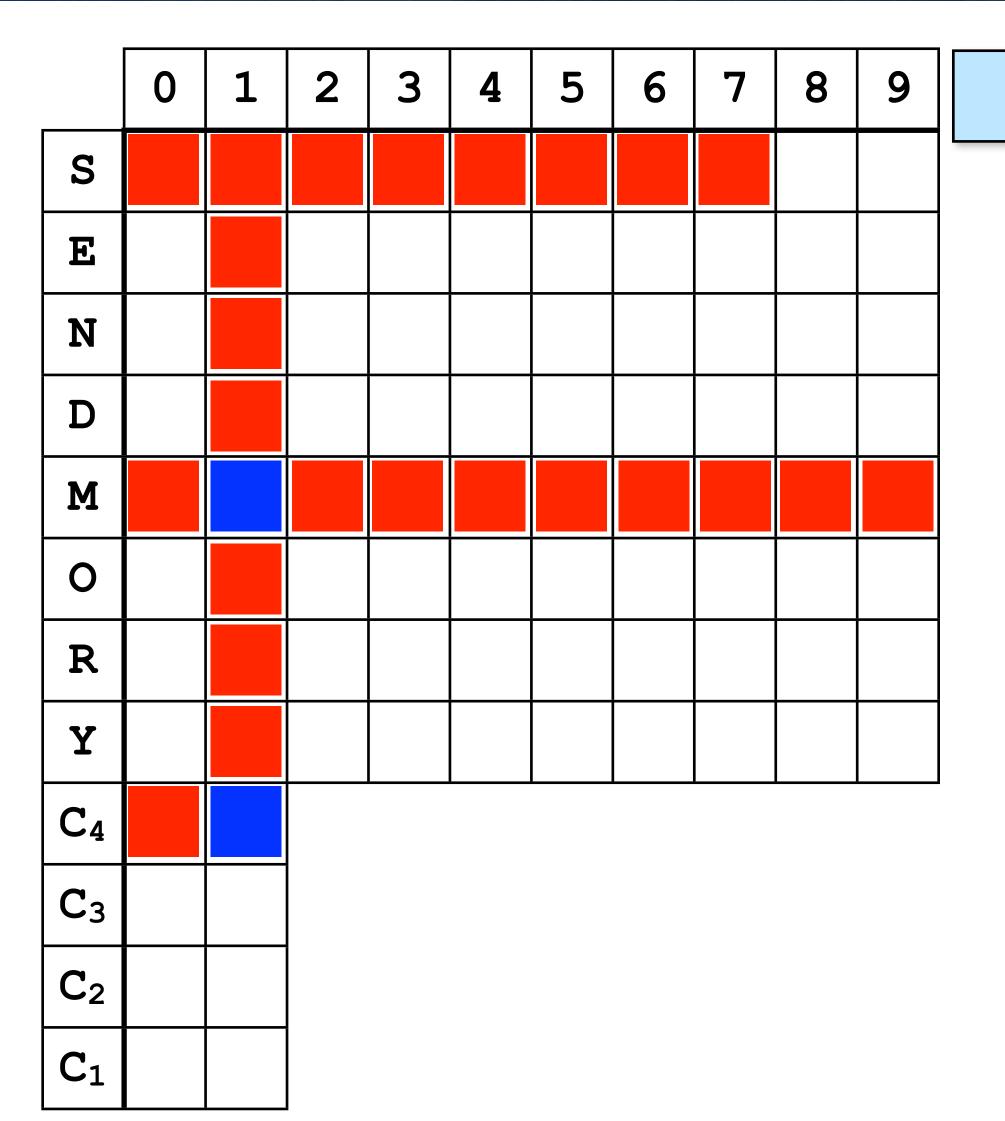
10 - 1 <= carry[3] + value[S] <= 11 - 1

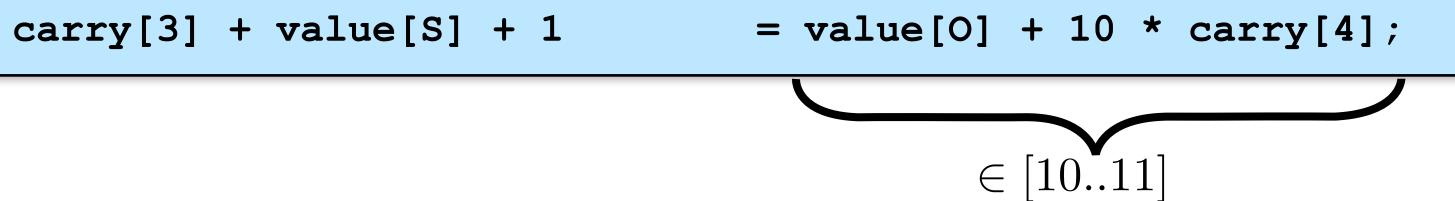
10 - 1 - carry[3] <= value[S] <= 11 - 1 - carry[3]

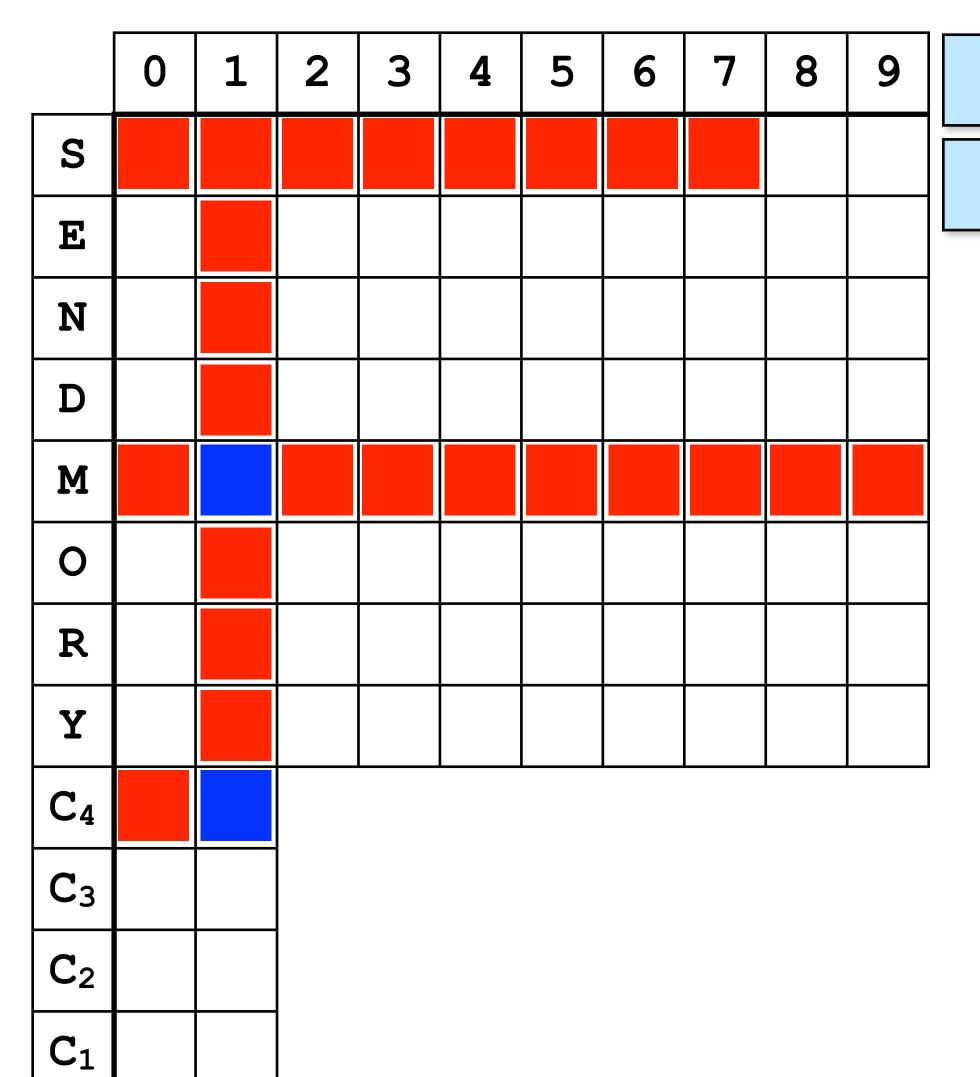
10 - 1 - 1 <= value[S] <= 11 - 1 - 0

8 <= value[S] <= 10
```

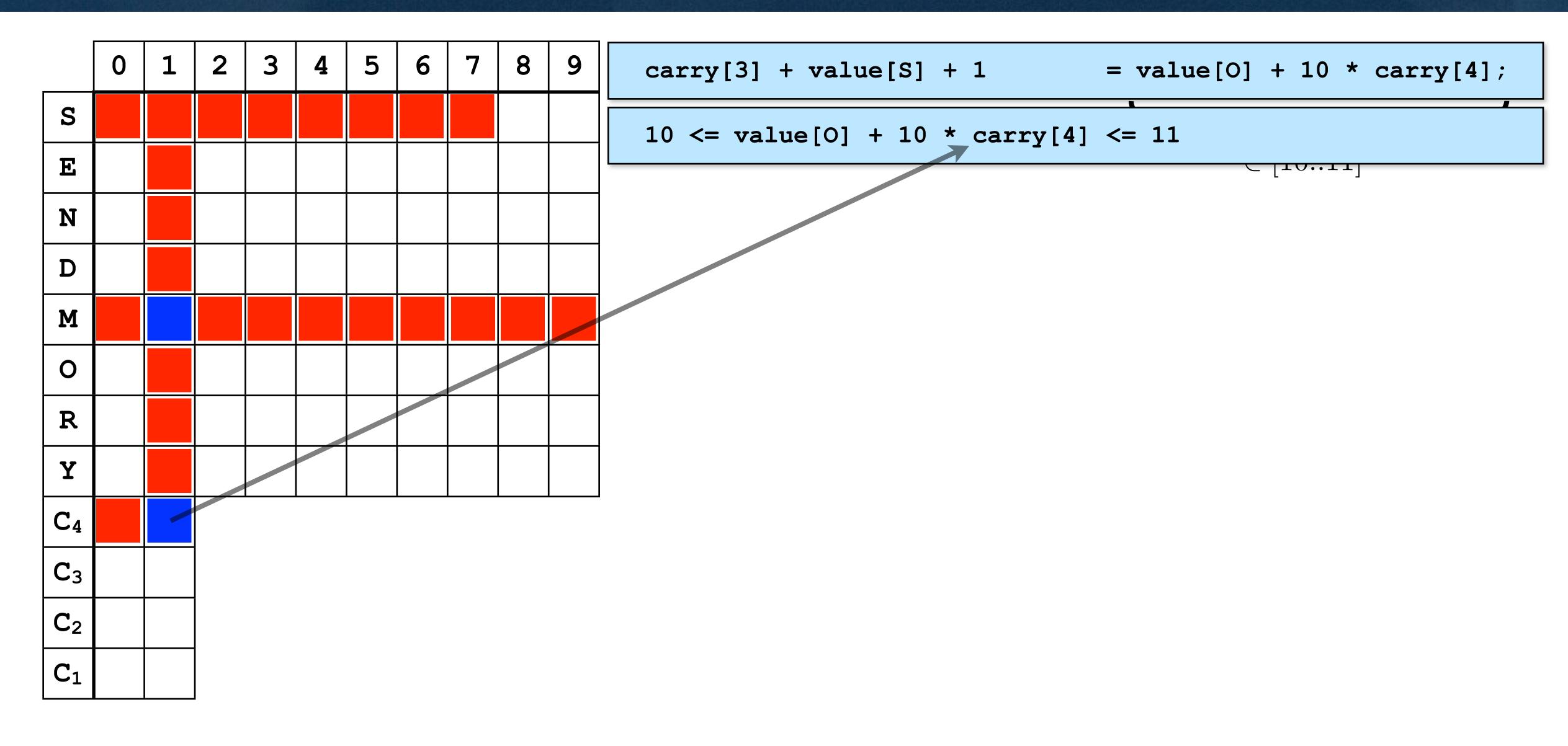
	0	1	2	3	4	5	6	7	8	9
S										
E										
N	-									
D										
M										
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R										
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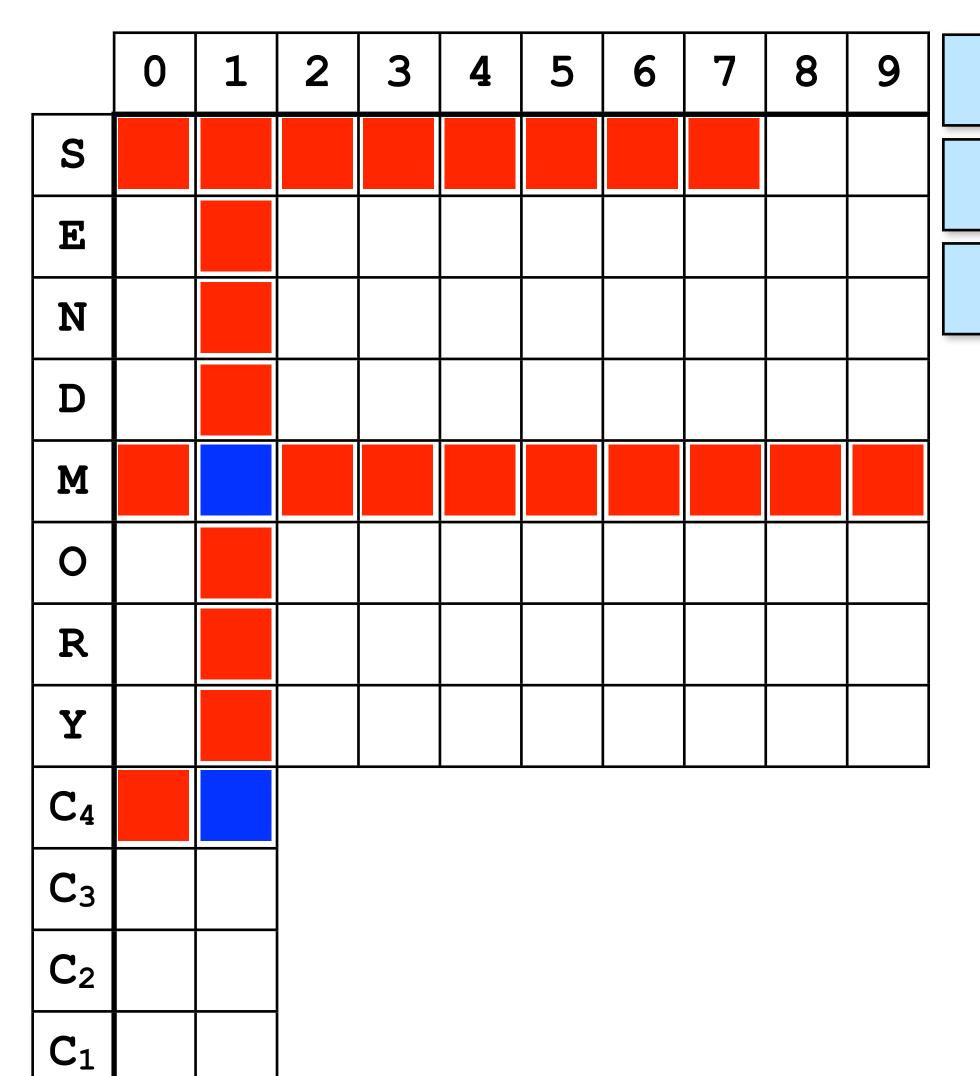


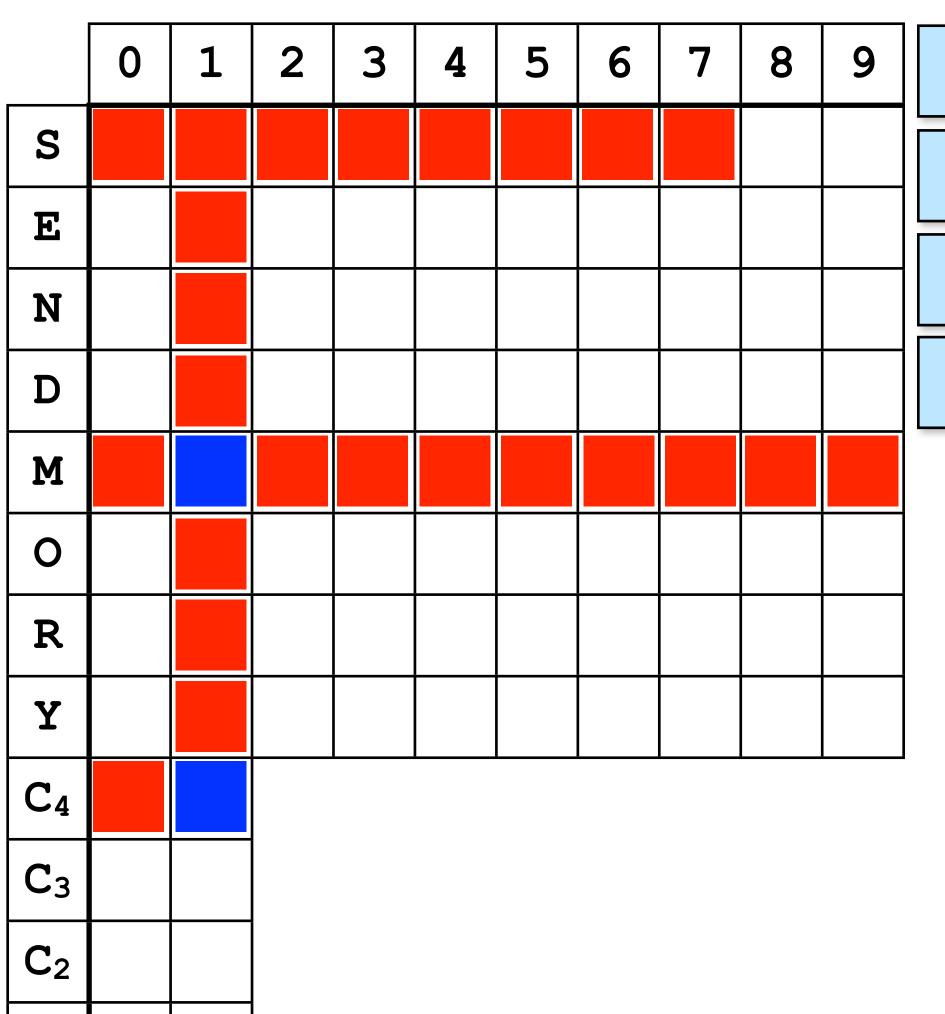


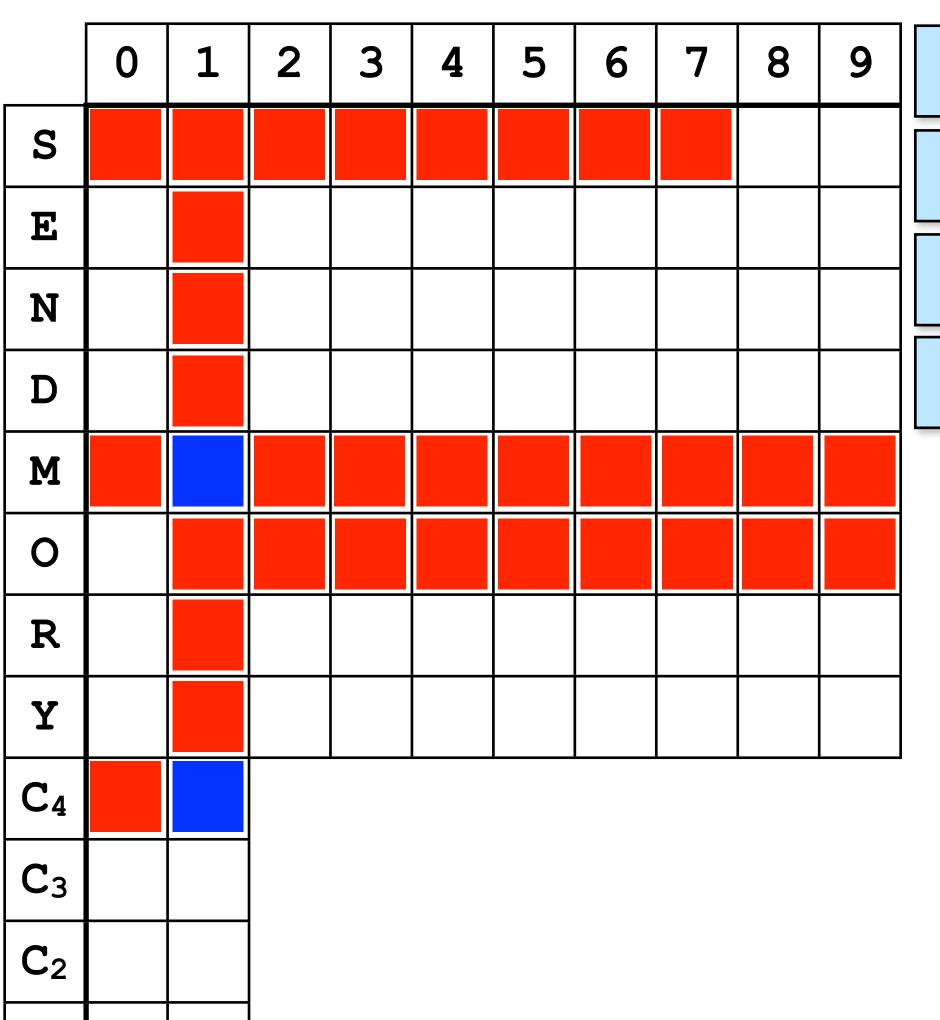


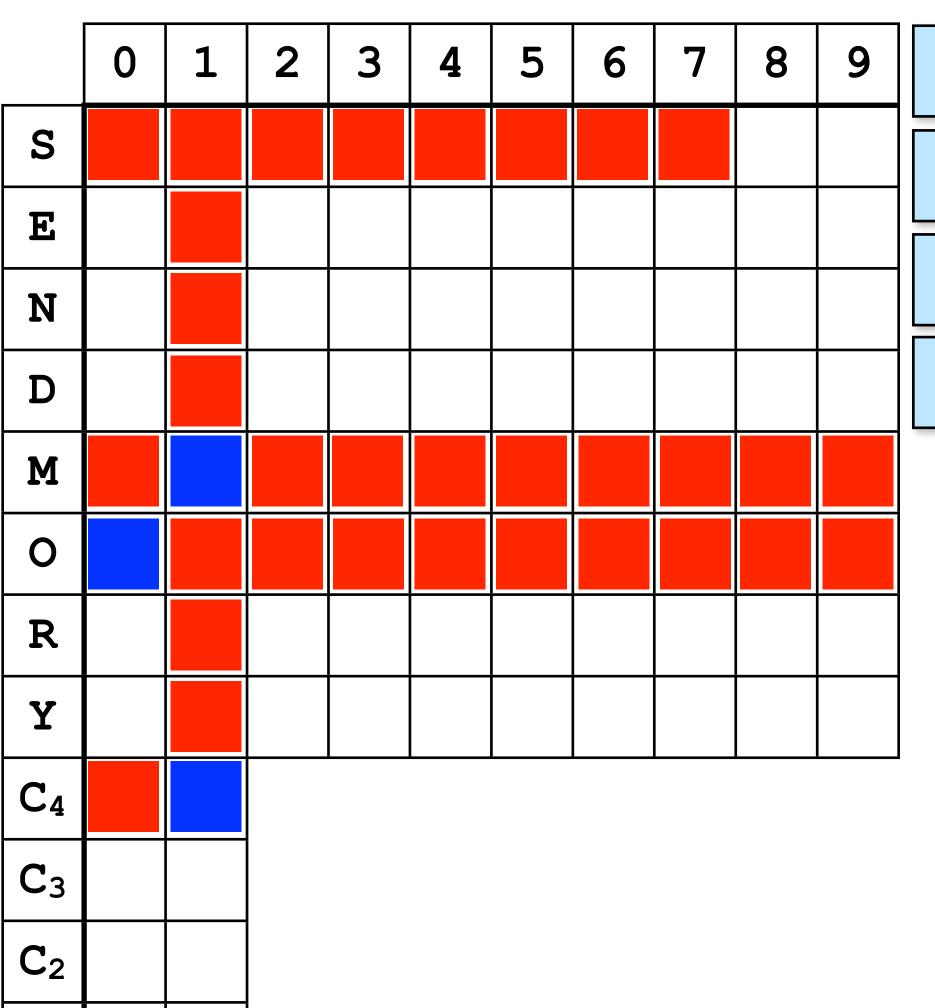
```
carry[3] + value[S] + 1 = value[O] + 10 * carry[4];
10 <= value[O] + 10 * carry[4] <= 11
```

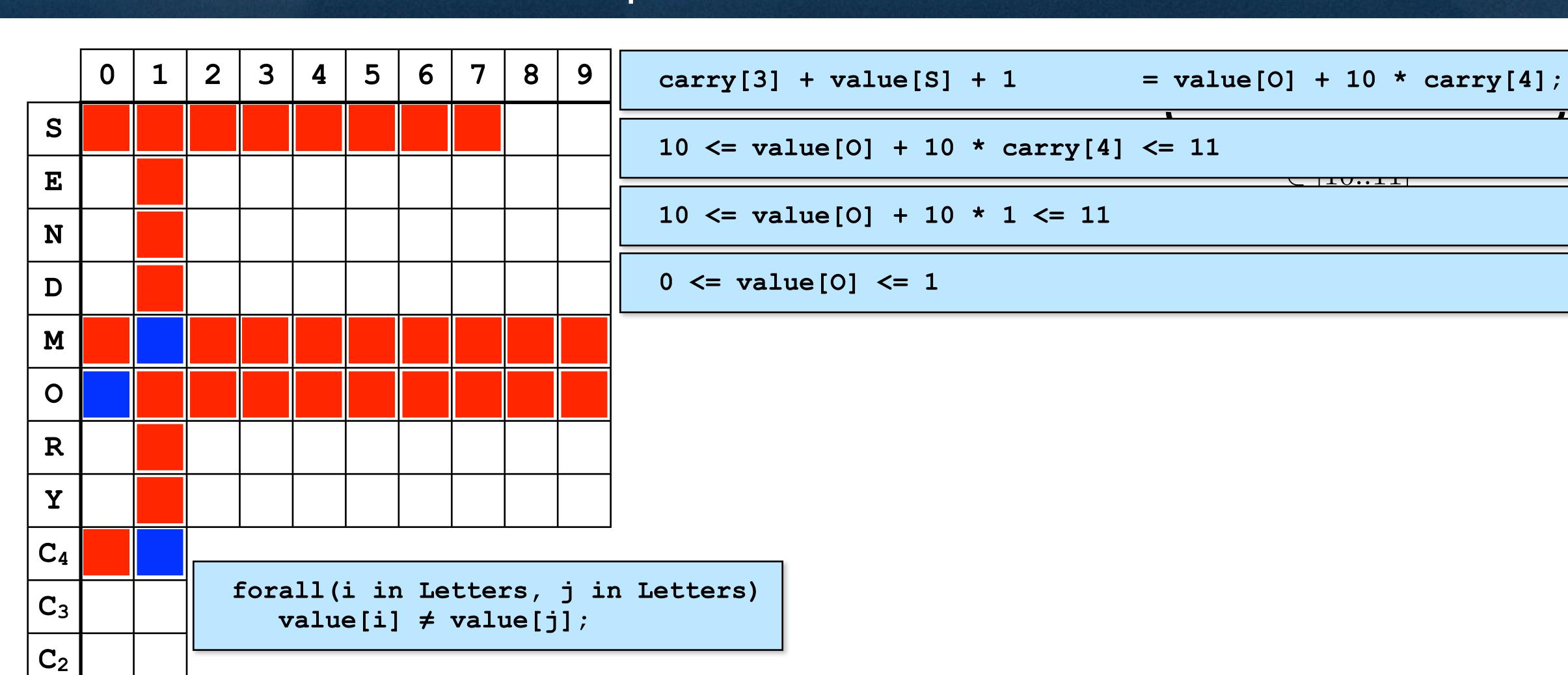


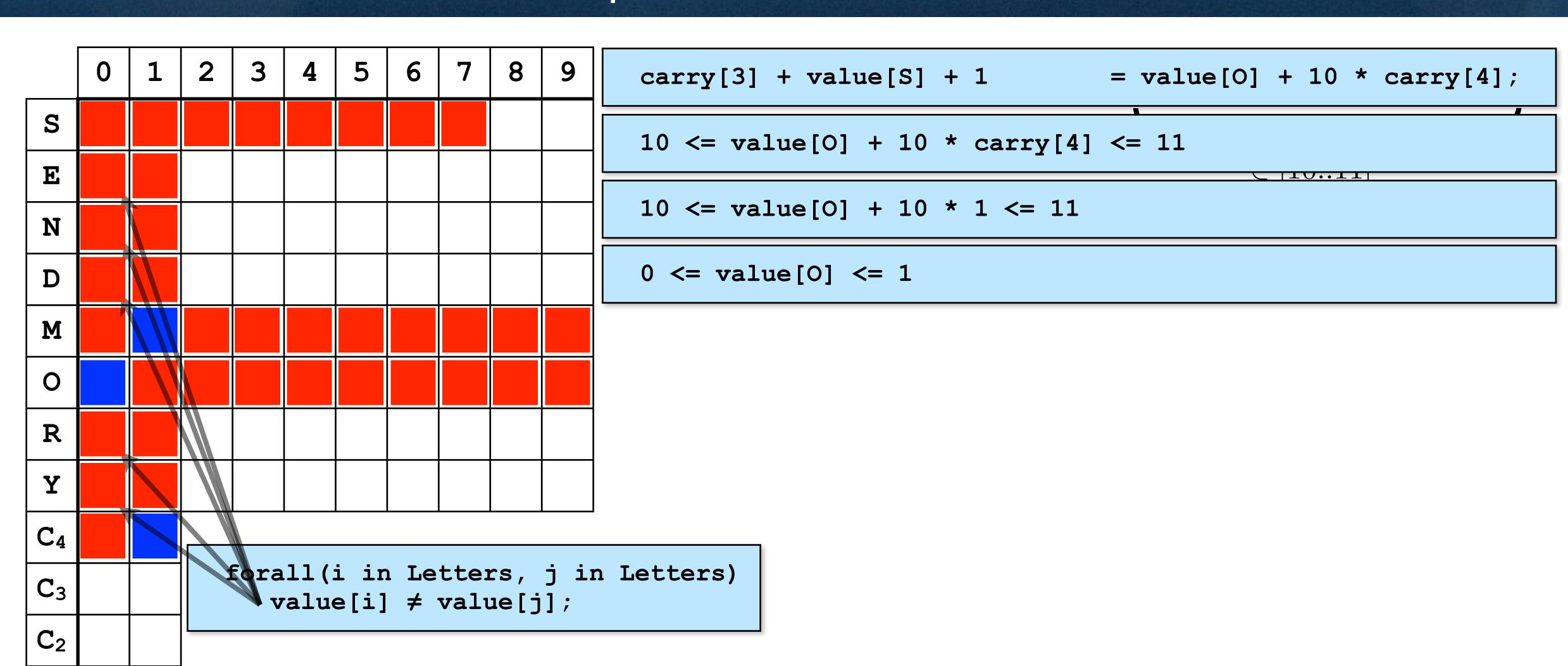










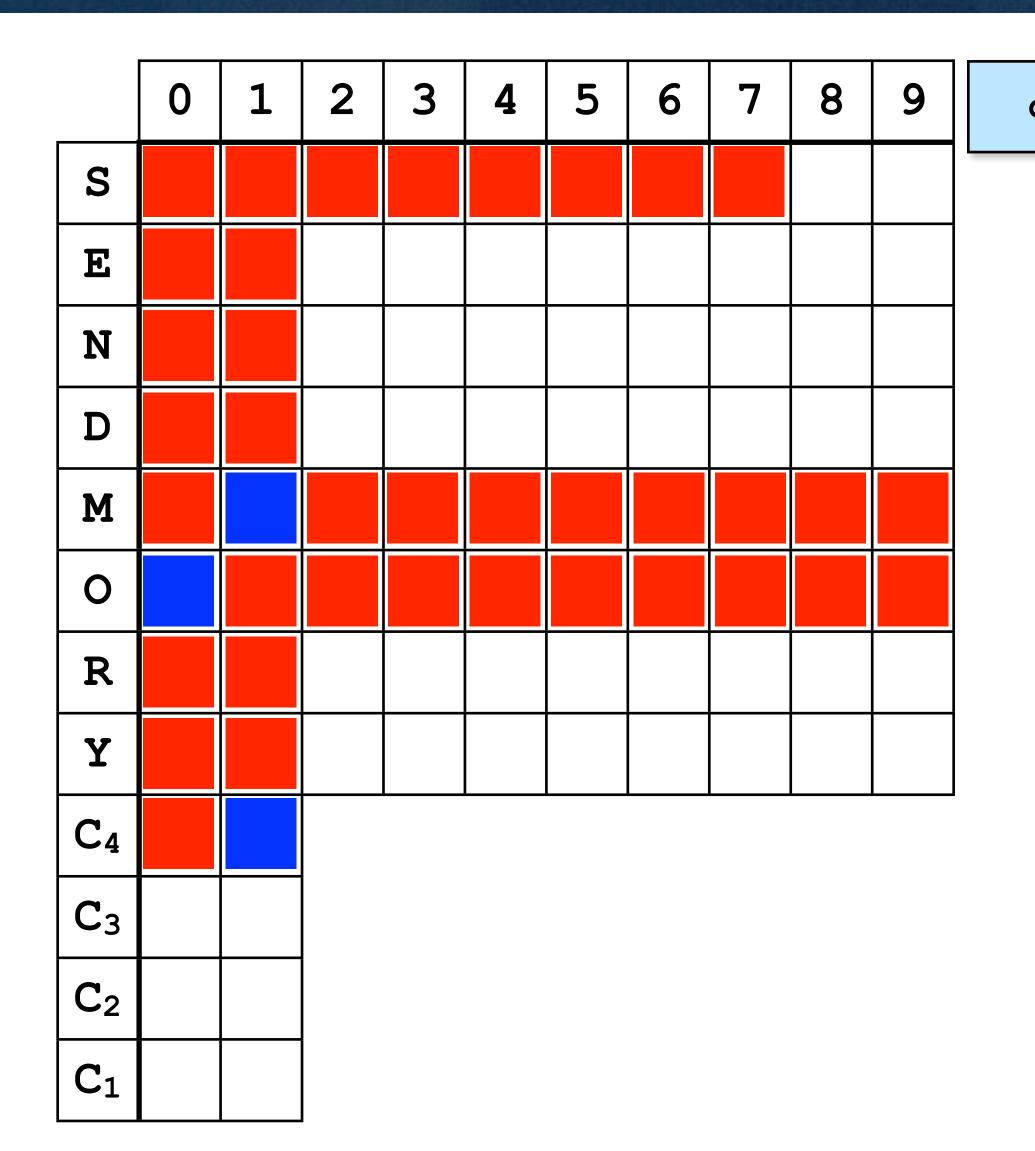


Send More Money

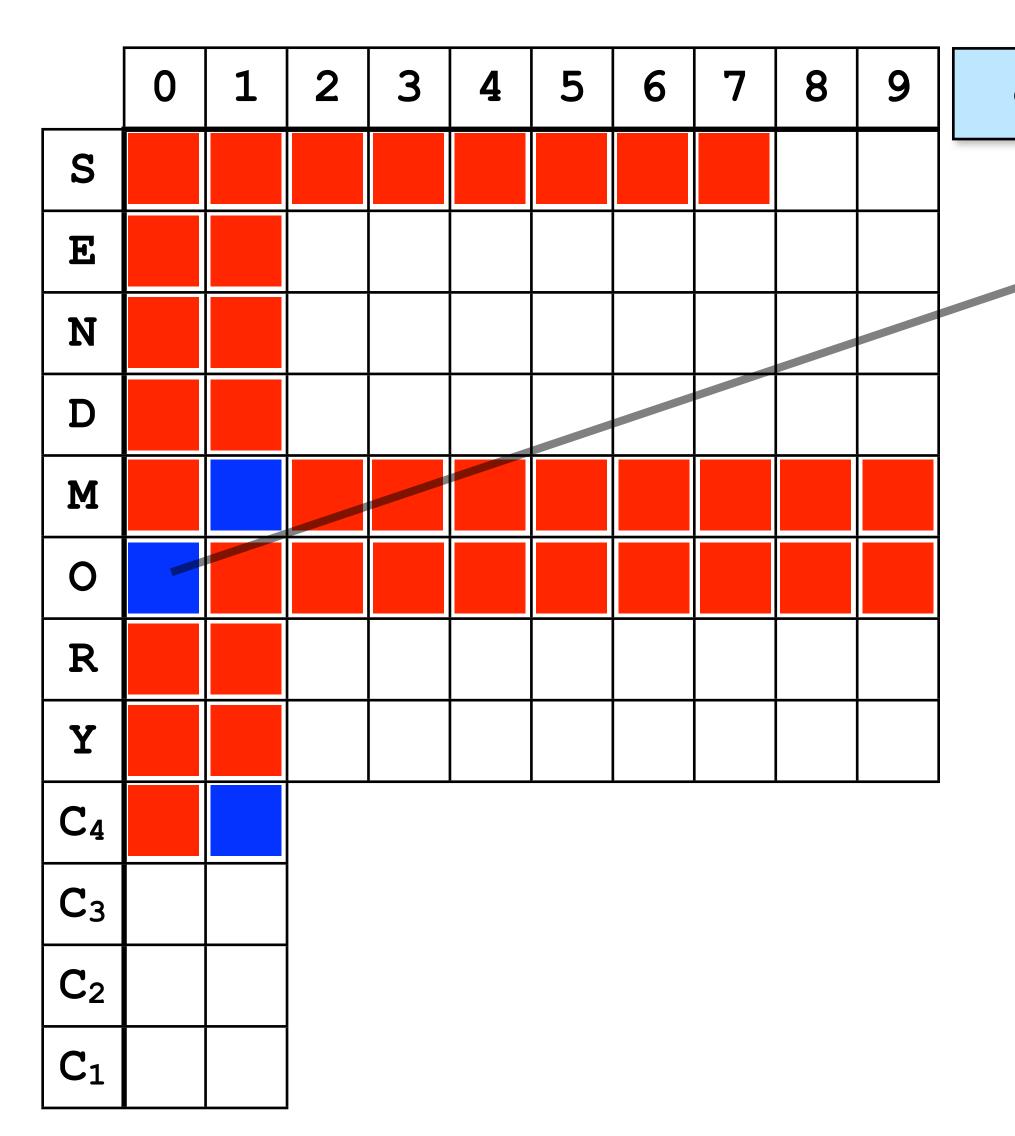
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range Digits = 0..9;
var{int} value[Letters] in Digits;
var{int} carry[1..4] in 0..1;
solve {
 forall(i in Letters, j in Letters: i < j)</pre>
   value[i] ≠ value[j];
 value[S] \neq 0;
 value[M] \neq 0;
                                 = value[M];
 carry[4]
 carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];
carry[2] + value[E] + value[O] = value[N] + 10 * carry[3];
 carry[1] + value[N] + value[R] = value[E] + 10 * carry[2];
            value[D] + value[E] = value[Y] + 10 * carry[1];
```

```
C<sub>4</sub> C<sub>3</sub> C<sub>2</sub> C<sub>1</sub>
S E N D
+ M O R E

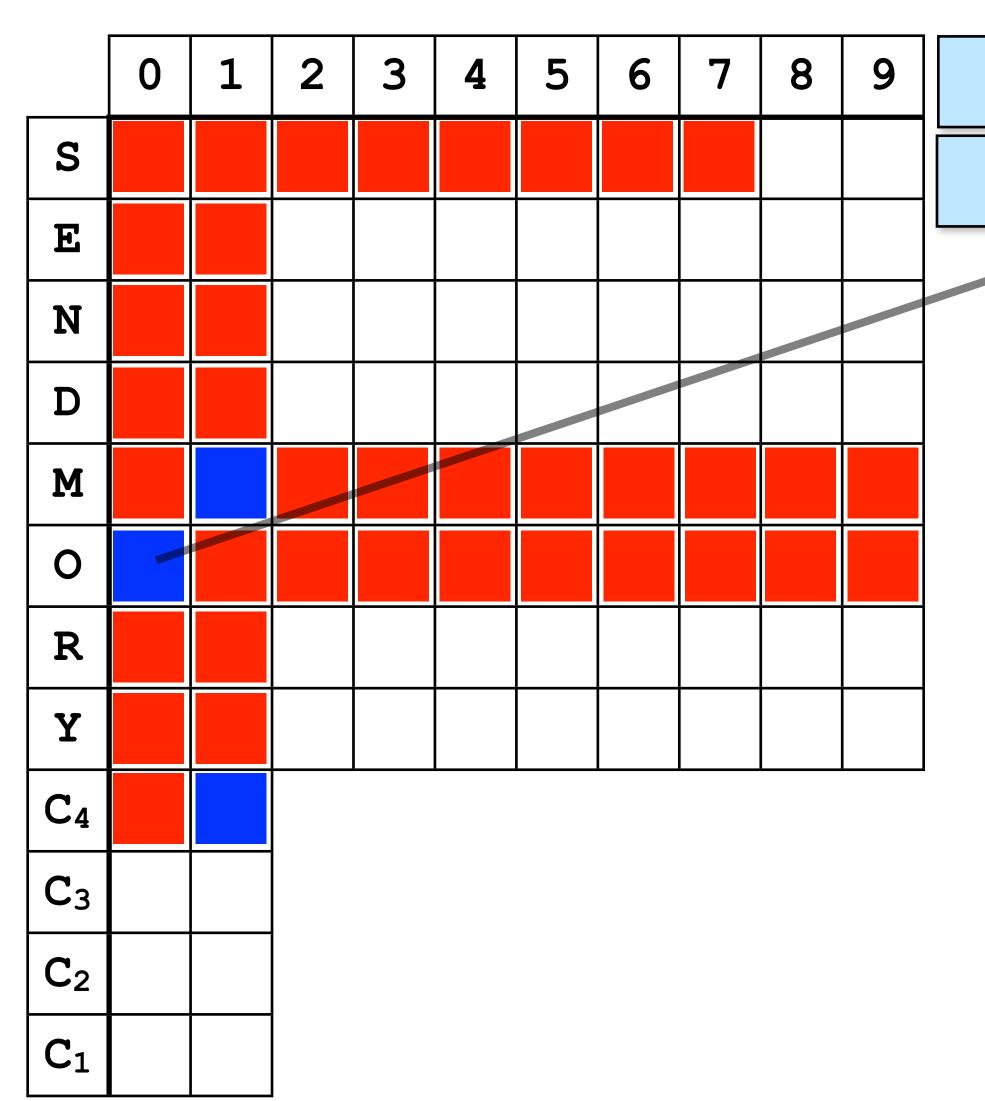
= M O N E Y
```



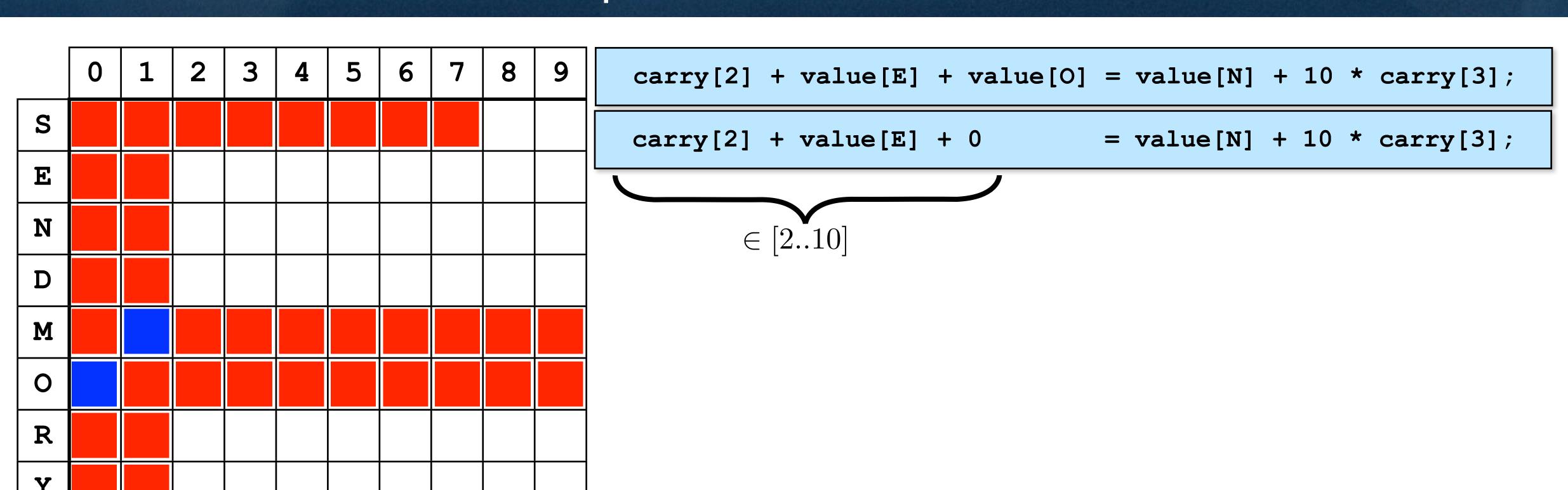
carry[2] + value[E] + value[O] = value[N] + 10 * carry[3];



carry[2] + value[E] + value[O] = value[N] + 10 * carry[3];



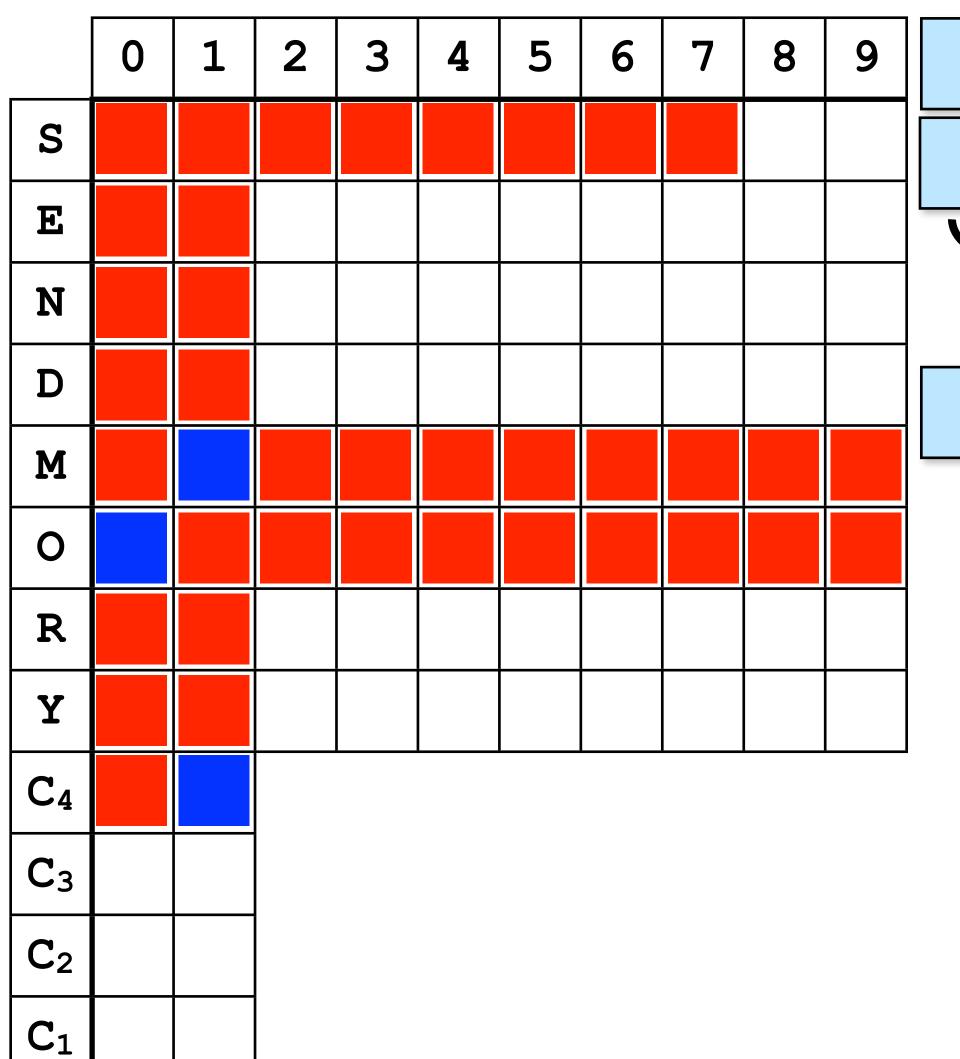
```
carry[2] + value[E] + value[O] = value[N] + 10 * carry[3];
carry[2] + value[E] + 0 = value[N] + 10 * carry[3];
```

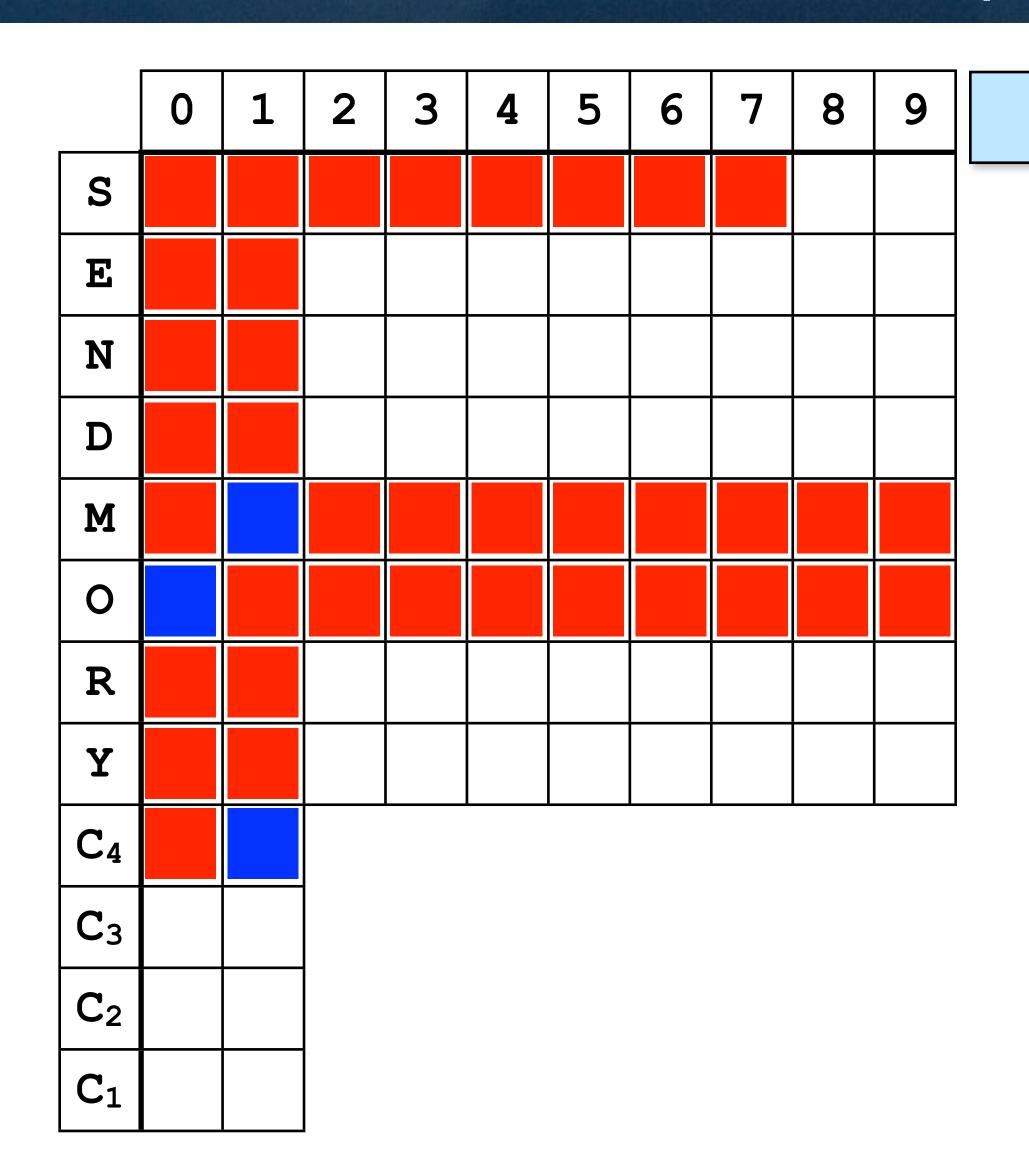


 C_4

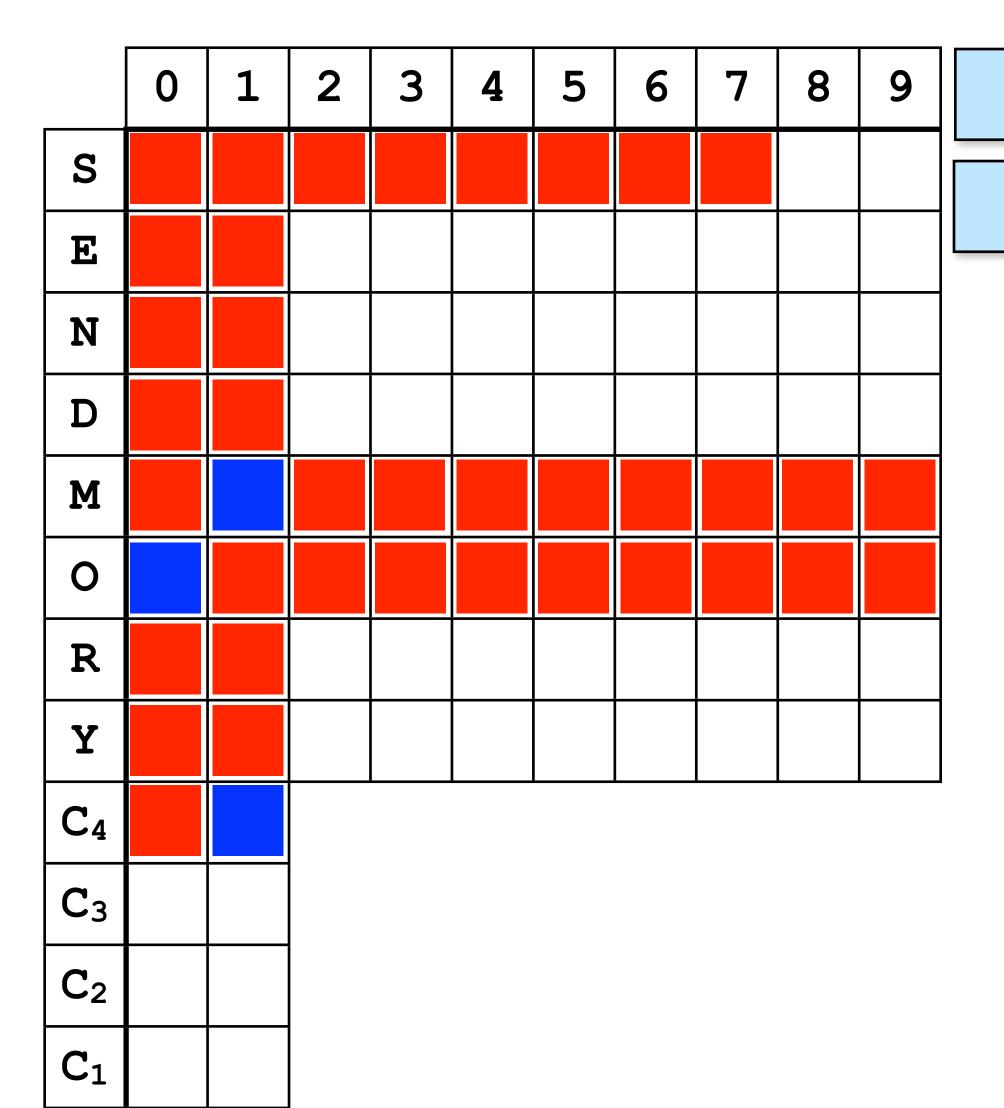
 C_3

 C_2

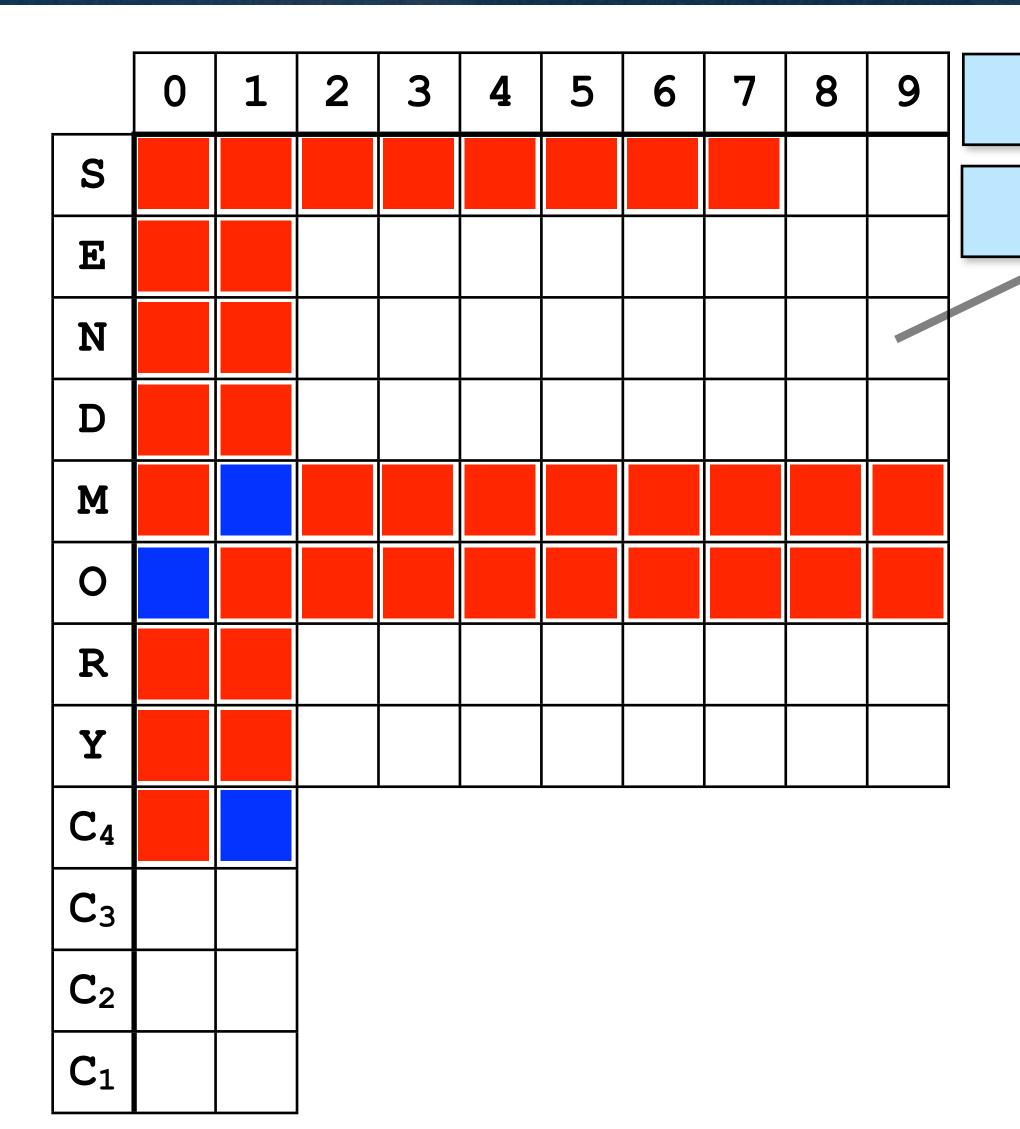




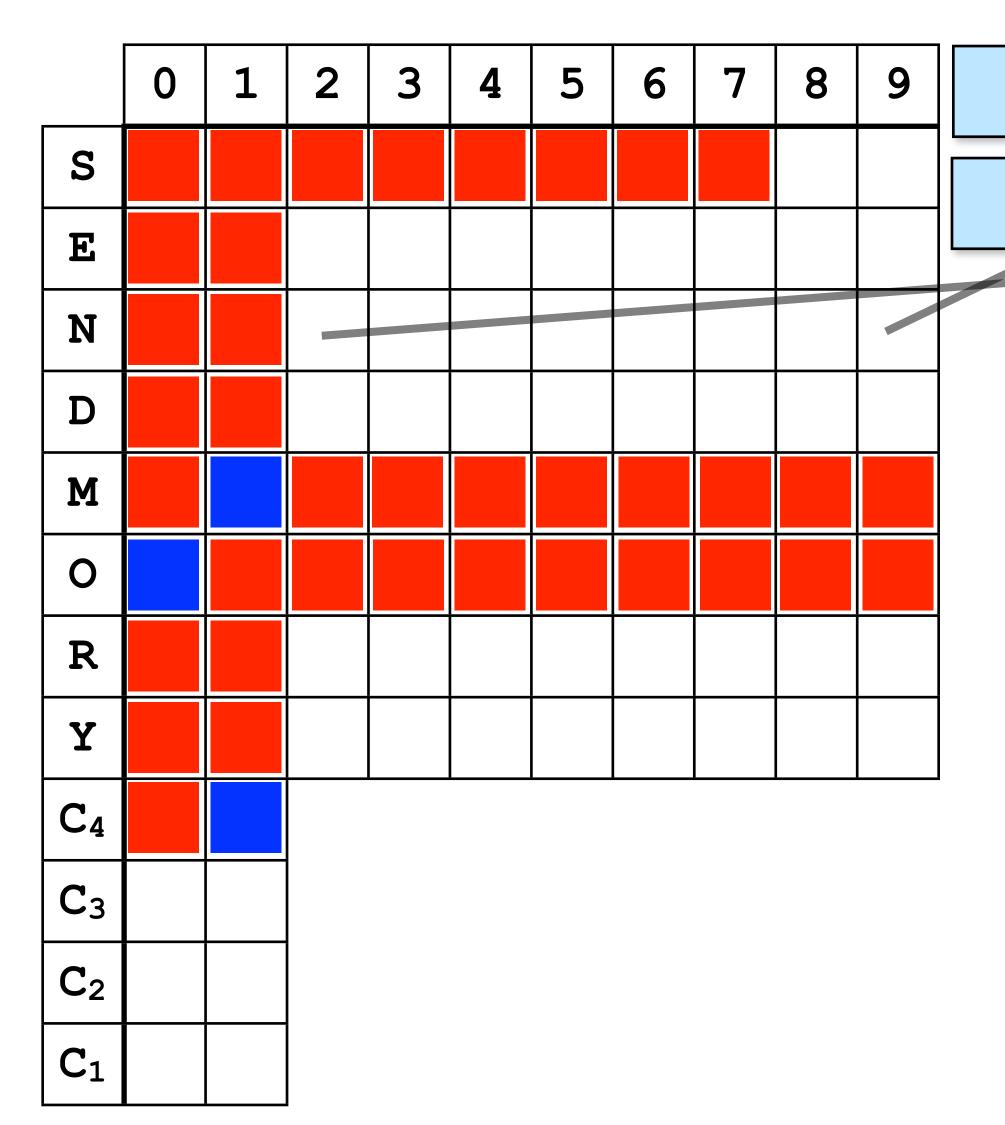
2 <= value[N] + 10 * carry[3] <= 10



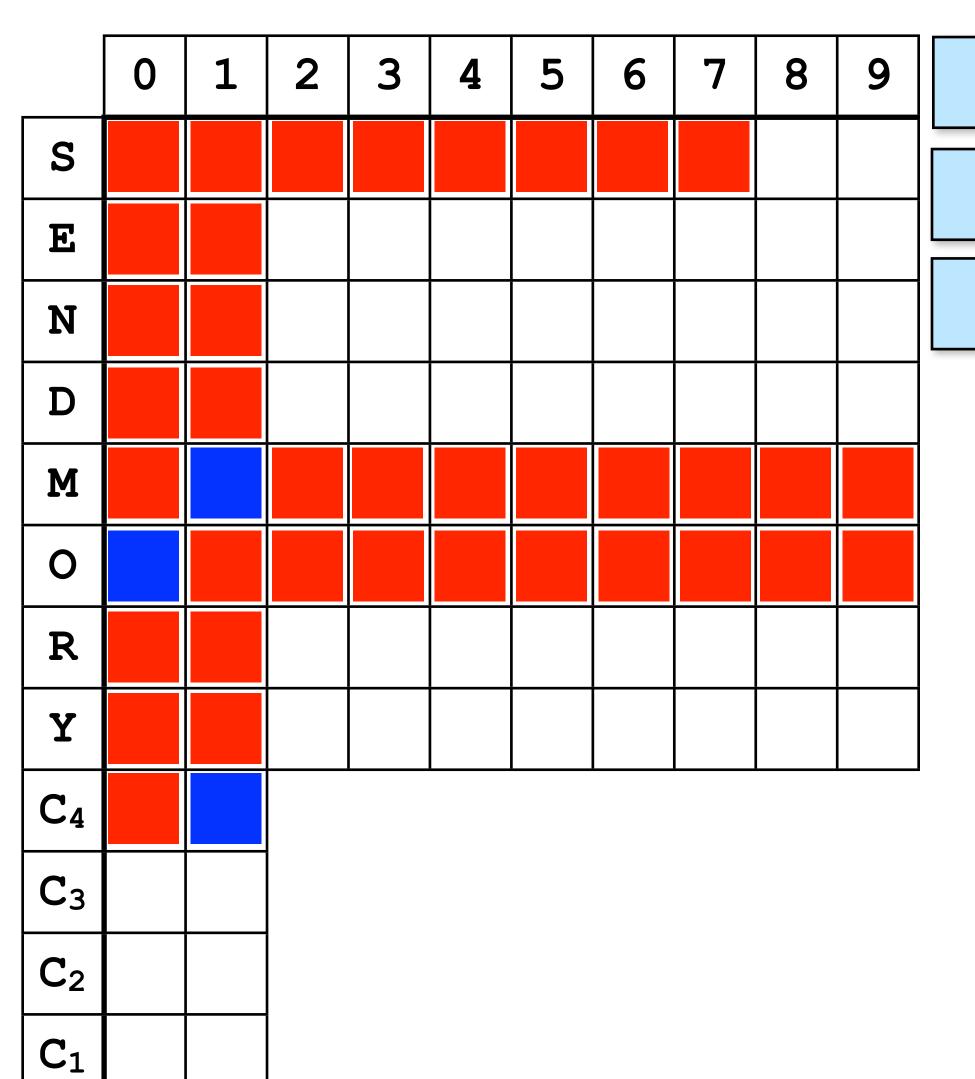
```
2 <= value[N] + 10 * carry[3] <= 10
```



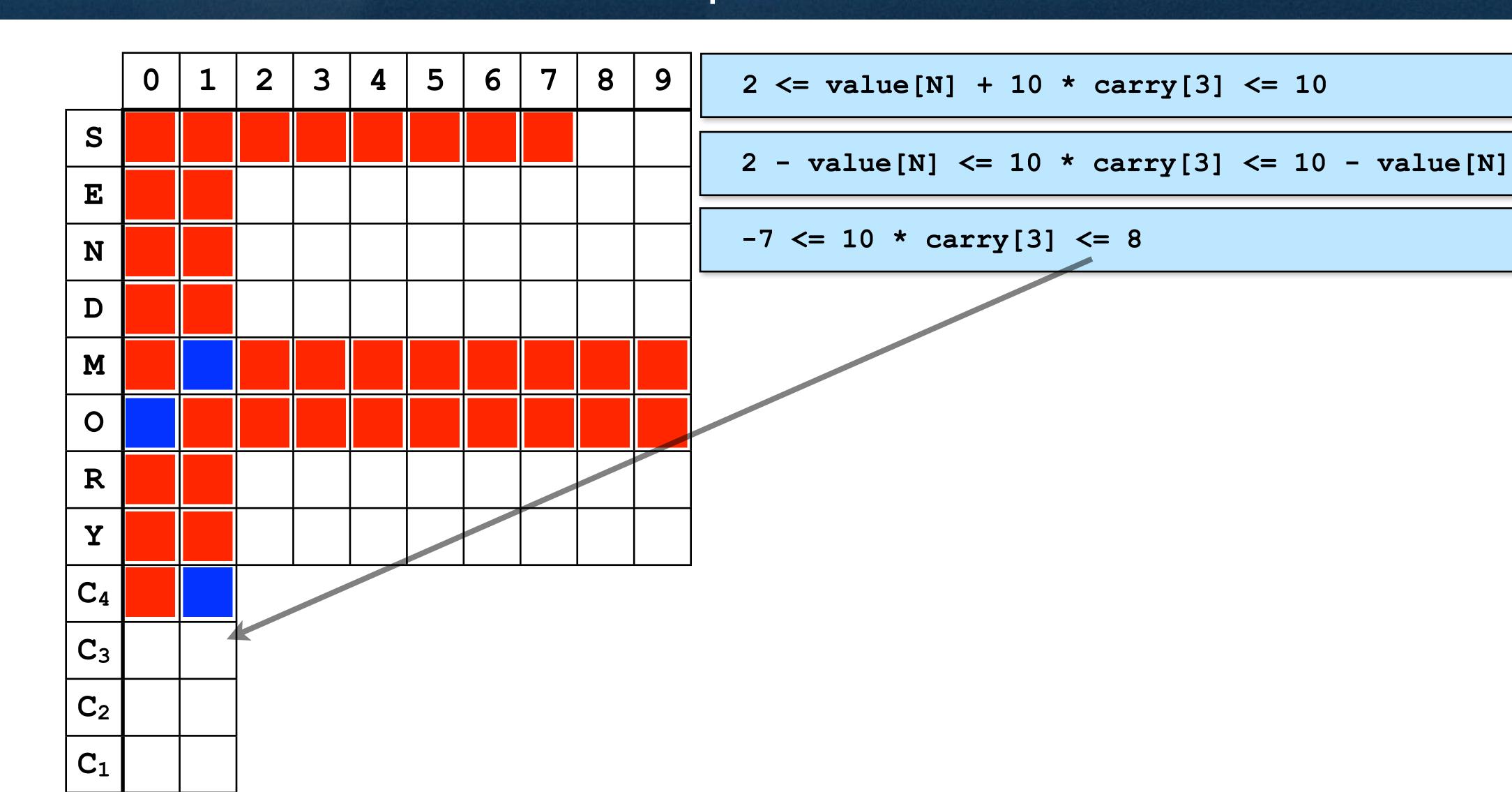
```
2 <= value[N] + 10 * carry[3] <= 10
```

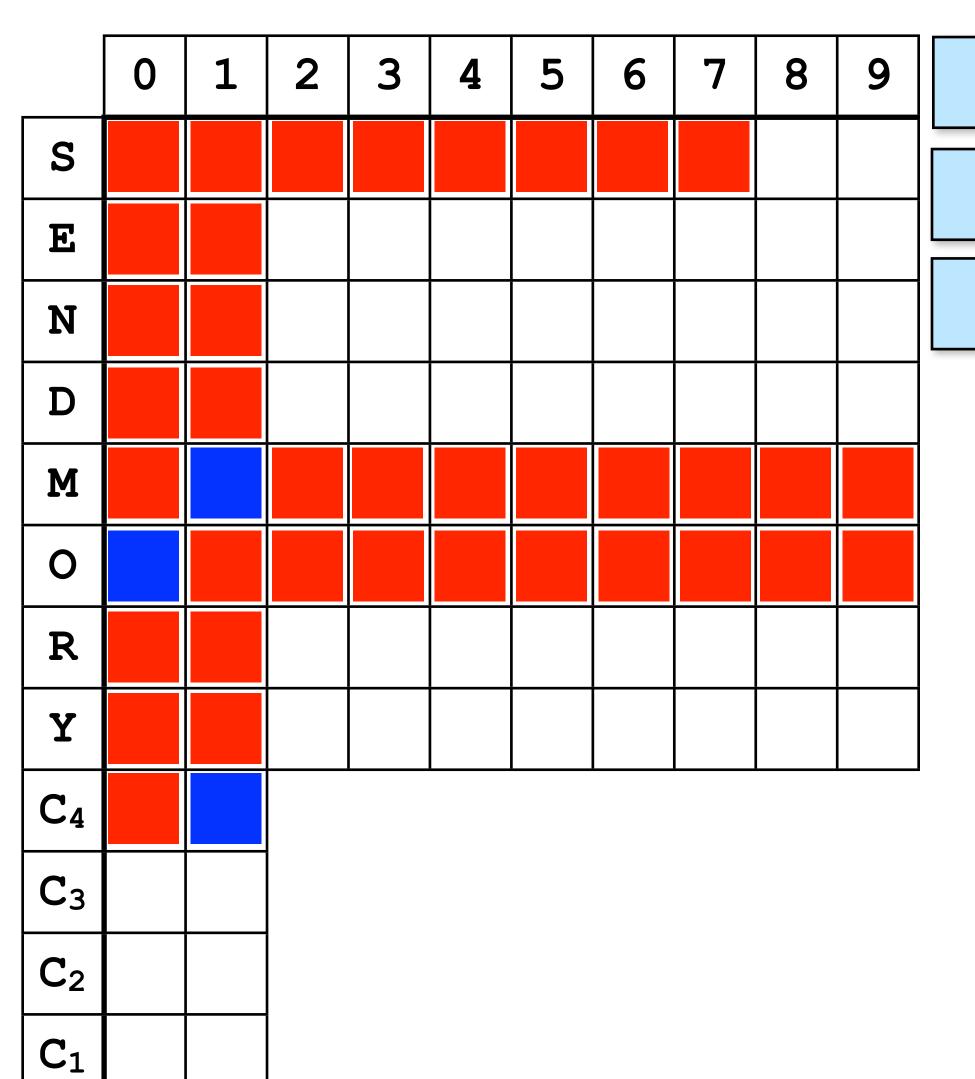


```
2 <= value[N] + 10 * carry[3] <= 10
```

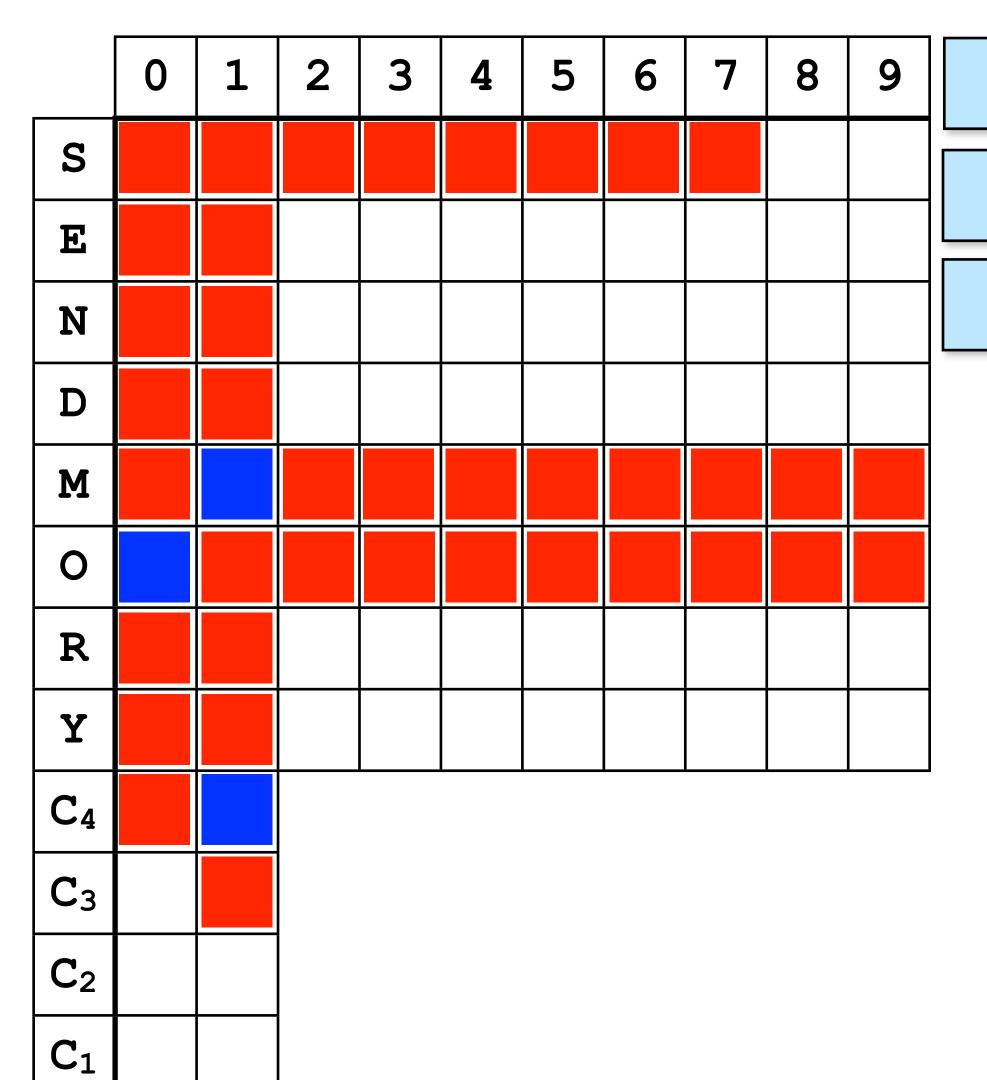


```
2 <= value[N] + 10 * carry[3] <= 10
```

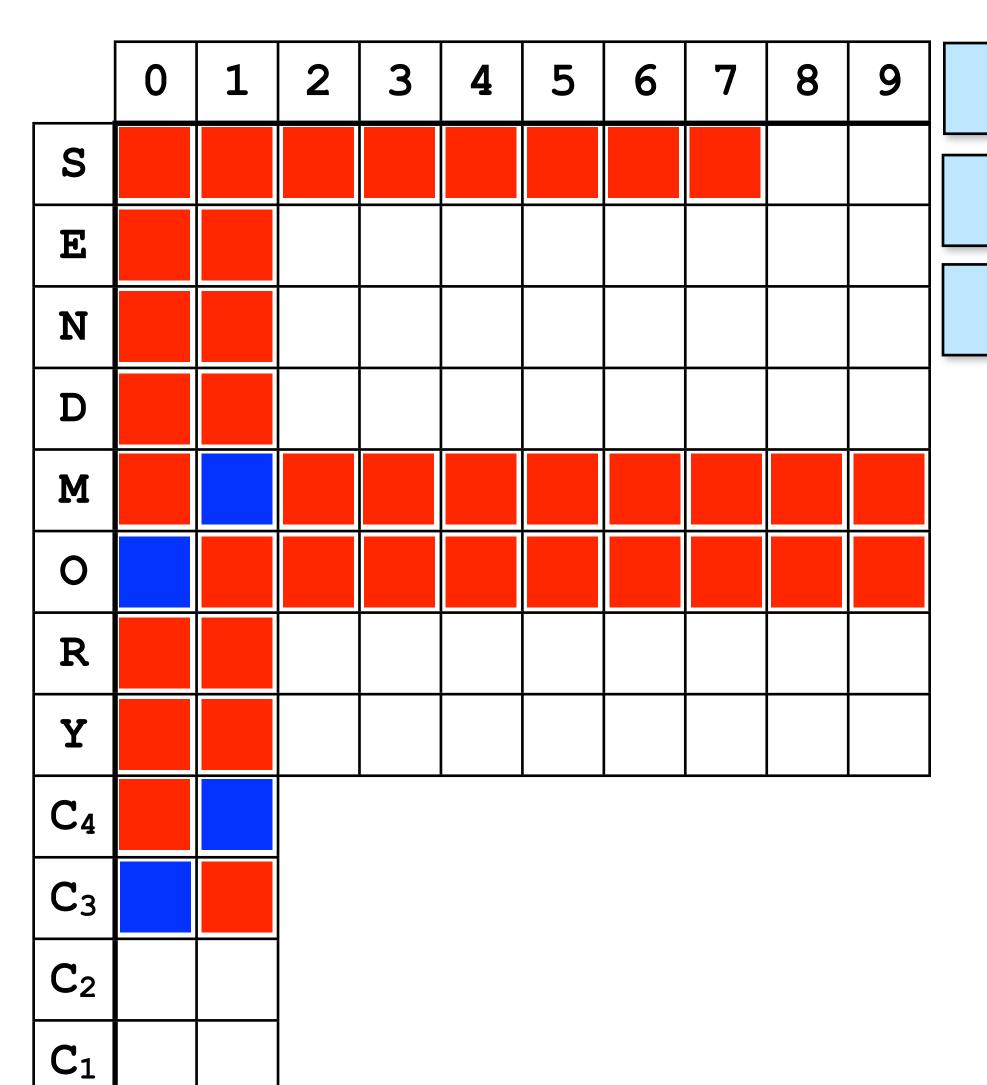




```
2 <= value[N] + 10 * carry[3] <= 10
```

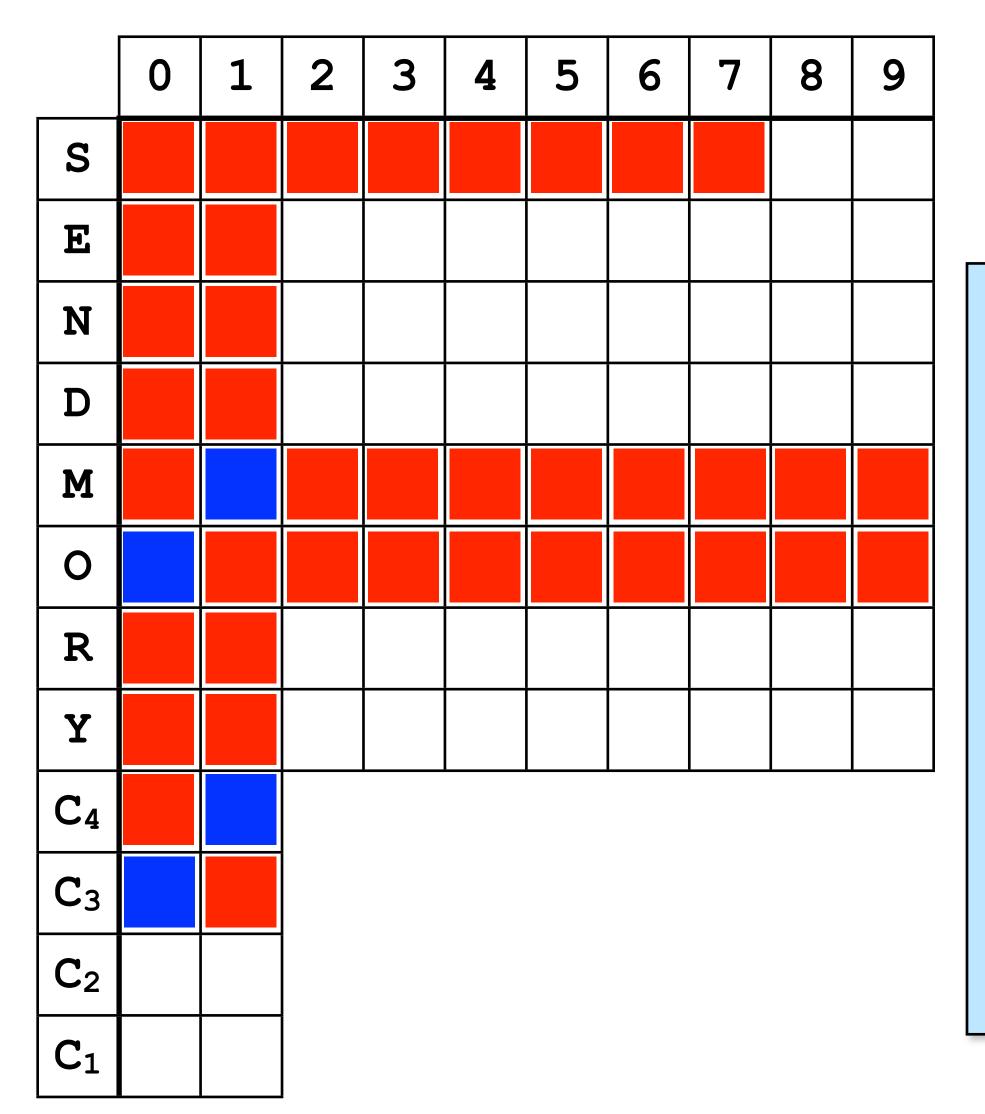


```
2 <= value[N] + 10 * carry[3] <= 10
```

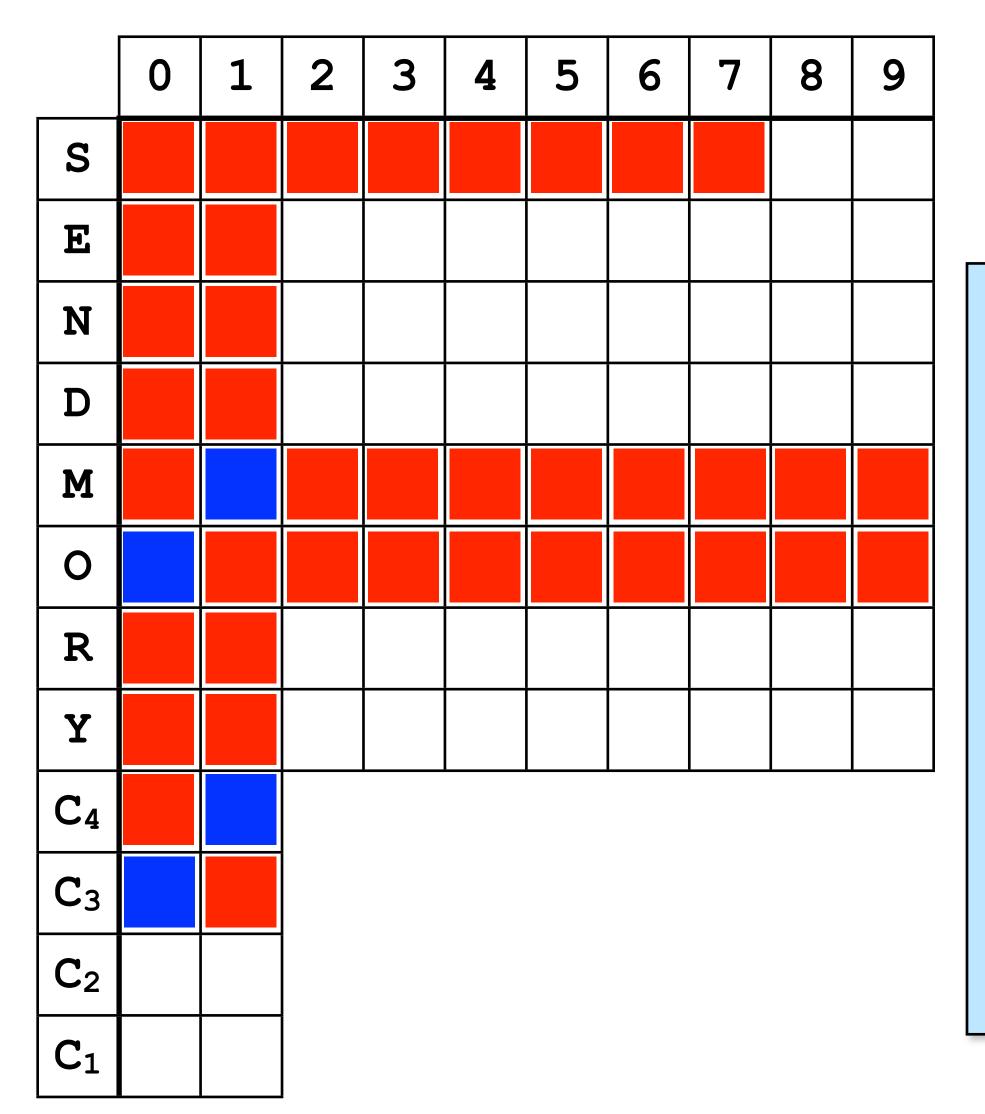


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2 <= value[N] + 10 * carry[3] <= 10
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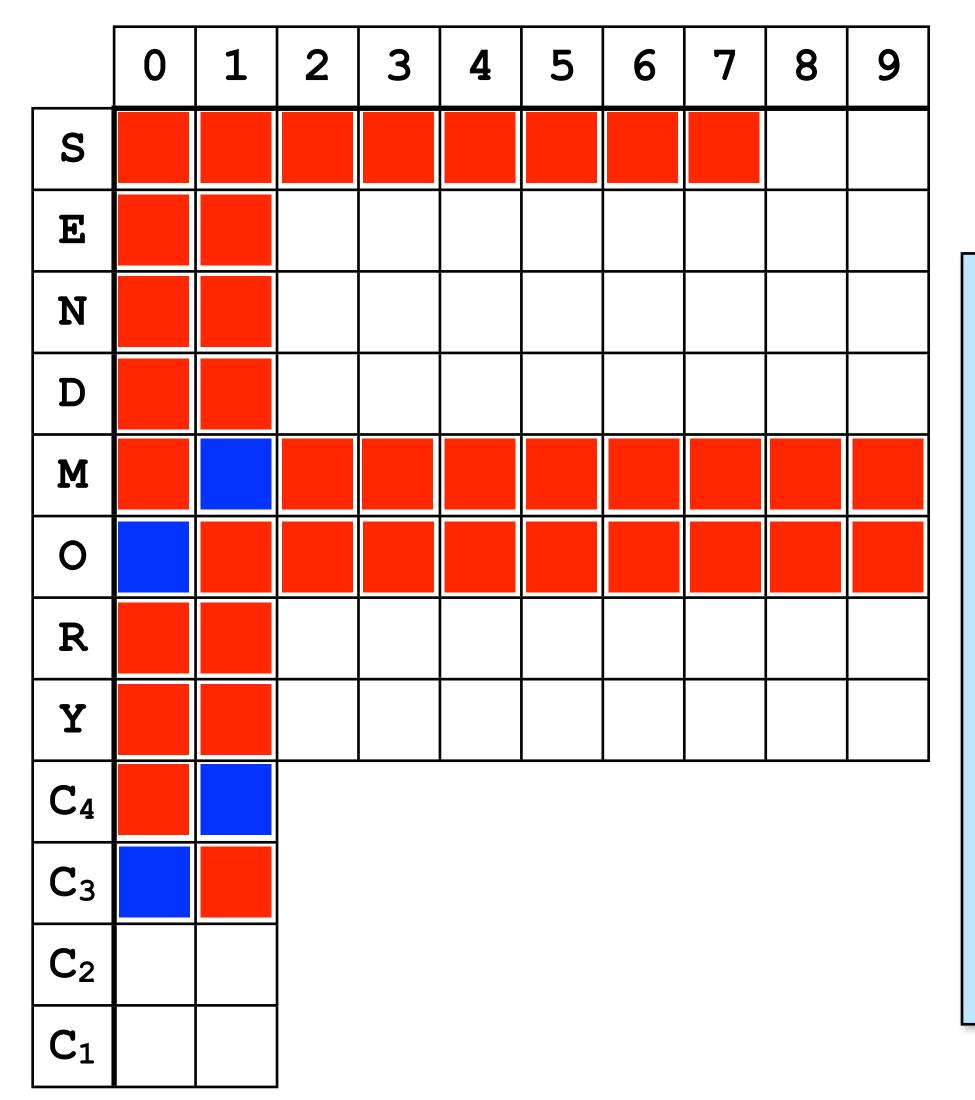
	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
0										
R										
Y										
C ₄										
C ₃										
C ₂										
C_1										



```
enum Letters = { S, E, N, D, M, O, R, Y};
range Digits = 0..9;
var{int} value[Letters] in Digits;
var{int} carry[1..4] in 0..1;
solve {
 forall(i in Letters, j in Letters: i < j)</pre>
    value[i] ≠ value[j];
 value[S] \neq 0;
 value[M] \neq 0;
                                 = value[M];
 carry[4]
 carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];
 carry[2] + value[E] + value[0] = value[N] + 10 * carry[3];
 carry[1] + value[N] + value[R] = value[E] + 10 * carry[2];
            value[D] + value[E] = value[Y] + 10 * carry[1];
```



```
enum Letters = { S, E, N, D, M, O, R, Y};
range Digits = 0..9;
var{int} value[Letters] in Digits;
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 value[S] \neq 0;
 value[M] \neq 0;
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 carry[4]
 carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];
 carry[2] + value[E] + value[0] = value[N] + 10 *: carry[3];;
 carry[1] + value[N] + value[R] = value[E] + 10 * carry[2];
            value[D] + value[E] = value[Y] + 10 * carry[1];
```

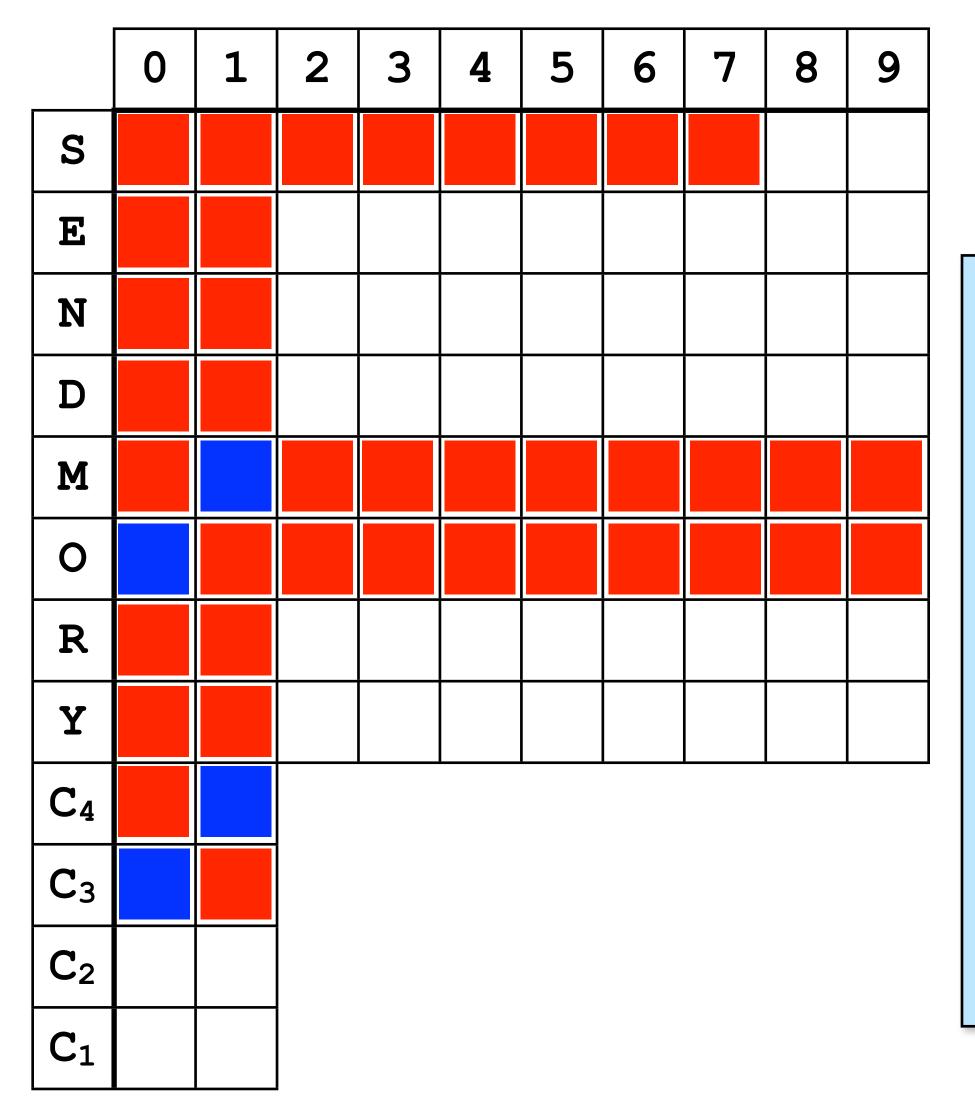


```
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                                 = value[M];
carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];
 carry[2] + value[E] + value[0] = value[N] + 10 * carry[3];
 carry[1] + value[N] + value[R] = value[E] + 10 * carry[2];
            value[D] + value[E] = value[Y] + 10 * carry[1];
```

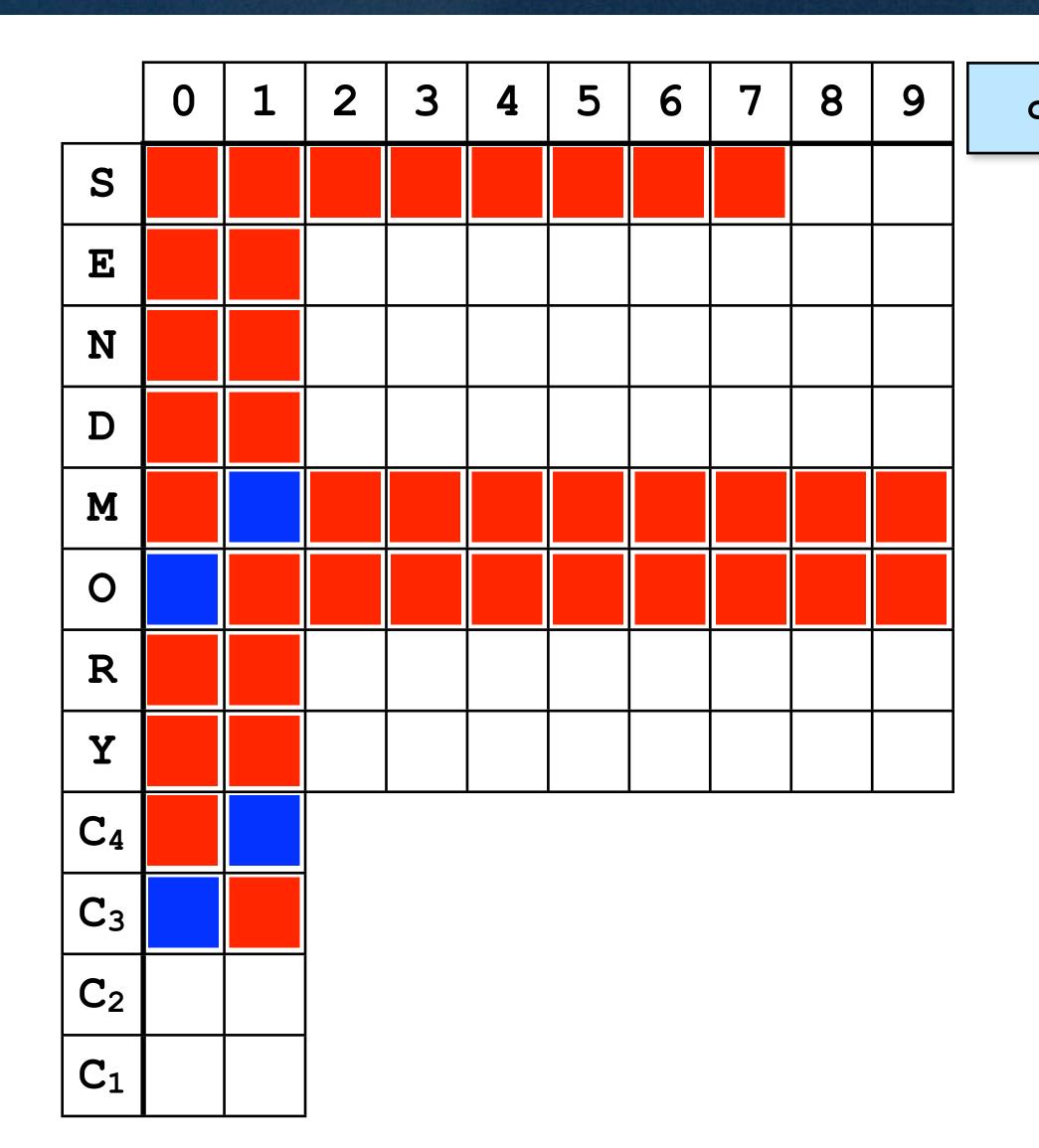
Computational Paradigm

- ► The propagation engine
 - this is the core of any constraint-programming solver
 - a simple (fixpoint) algorithm

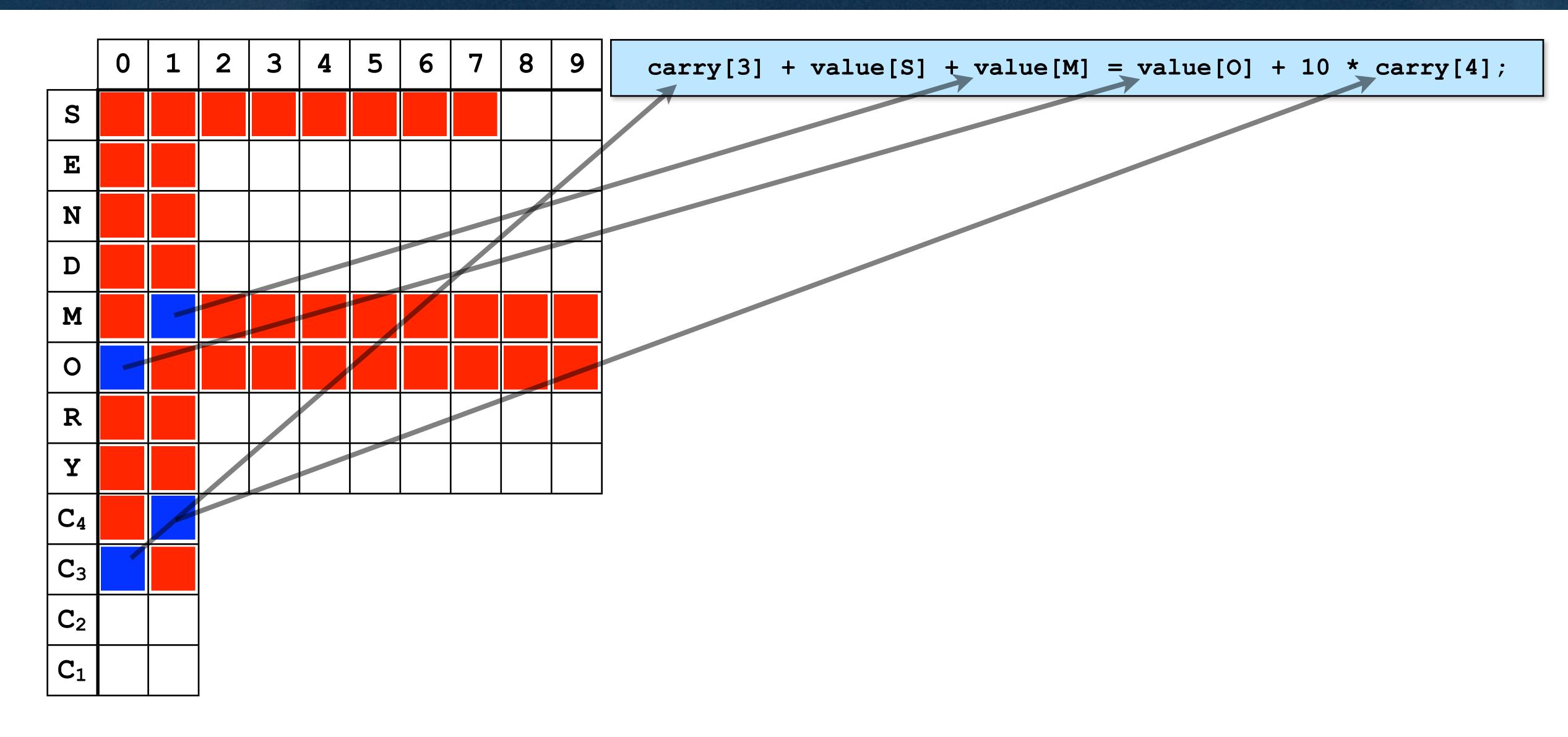
```
propagate()
{
  repeat
    select a constraint c;
    if c is infeasible given the domain store then
       return failure;
    else
       apply the pruning algorithm associated with c;
  (until no constraint can remove any value from the domain of its variables;
    return success;
}
```

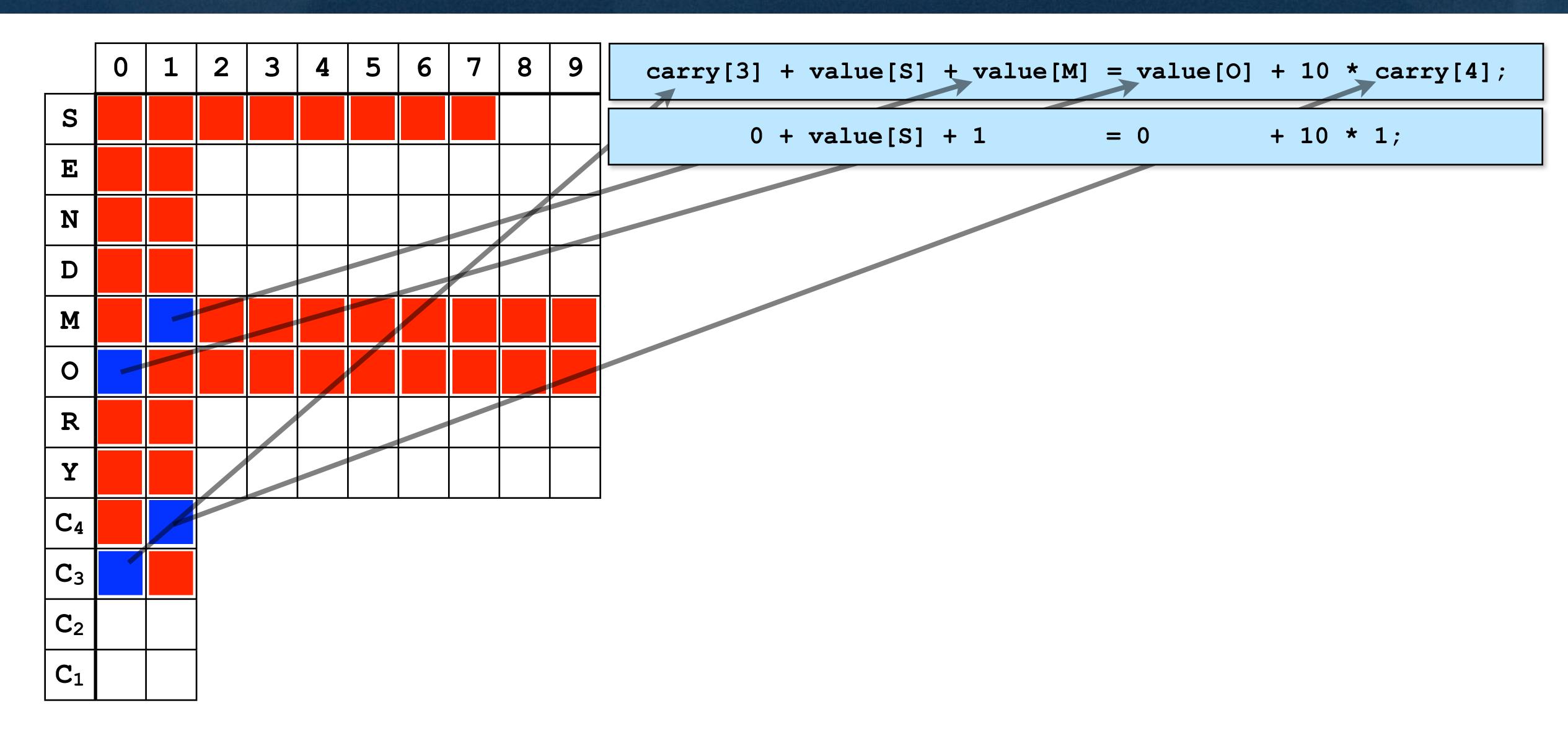


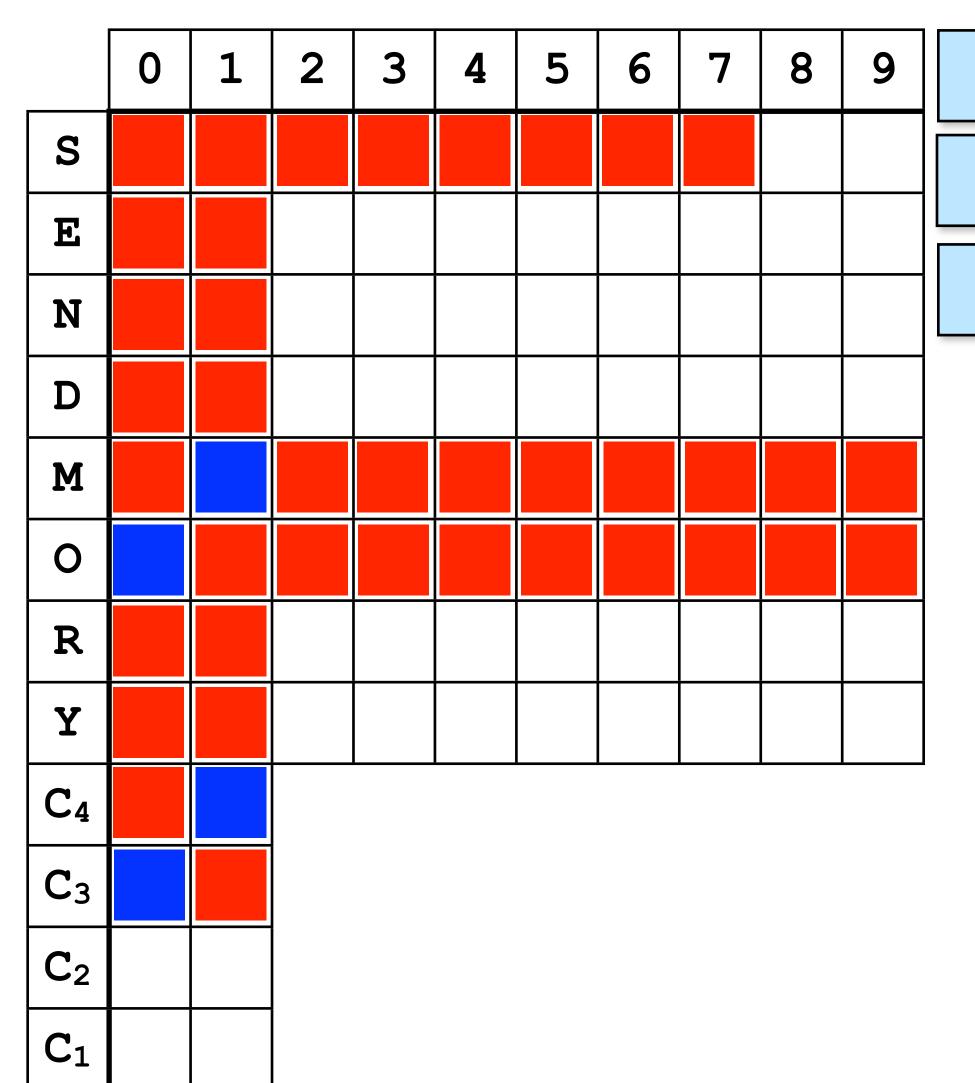
```
enum Letters = { S, E, N, D, M, O, R, Y};
range Digits = 0..9;
var{int} value[Letters] in Digits;
var{int} carry[1..4] in 0..1;
solve {
 forall(i in Letters, j in Letters: i < j)</pre>
    value[i] # value[j];
 value[S] \neq 0;
 value[M] \neq 0;
 carry[4]
                                 = value[M];
carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];
 carry[2] + value[E] + value[0] = value[N] + 10 * carry[3];
 carry[1] + value[N] + value[R] = value[E] + 10 * carry[2];
            value[D] + value[E] = value[Y] + 10 * carry[1];
```



carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];



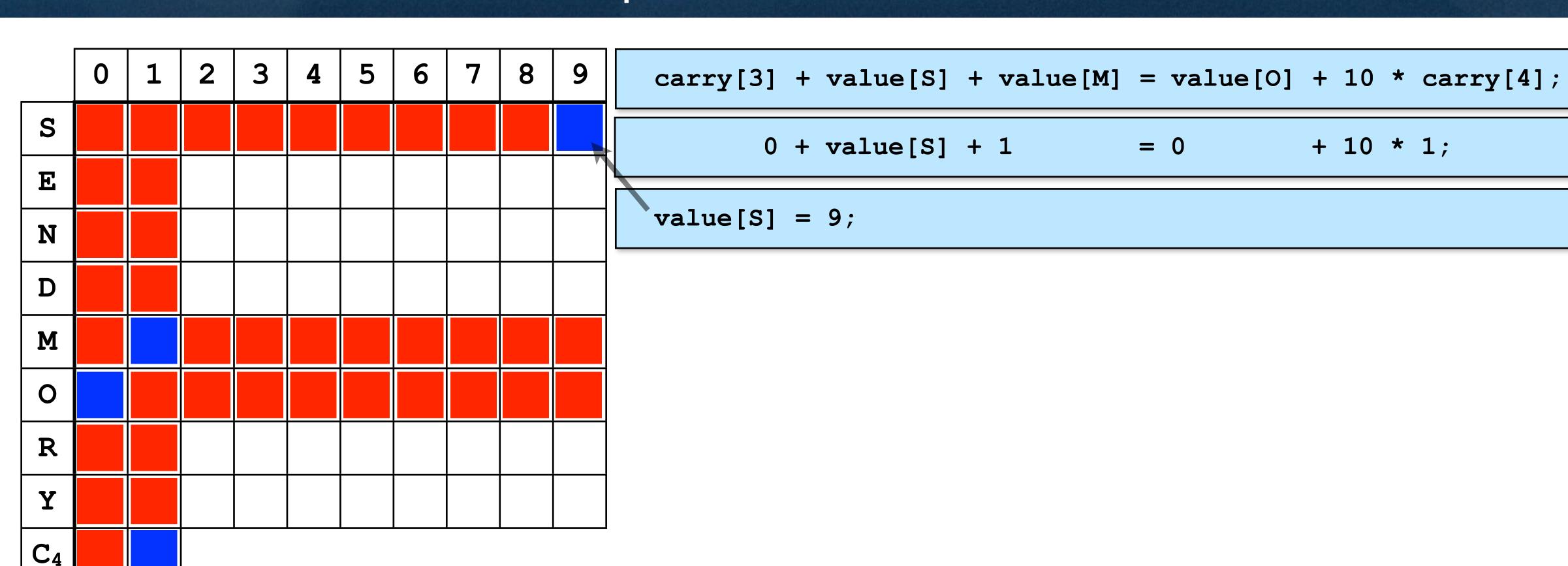




```
carry[3] + value[S] + value[M] = value[O] + 10 * carry[4];

0 + value[S] + 1 = 0 + 10 * 1;

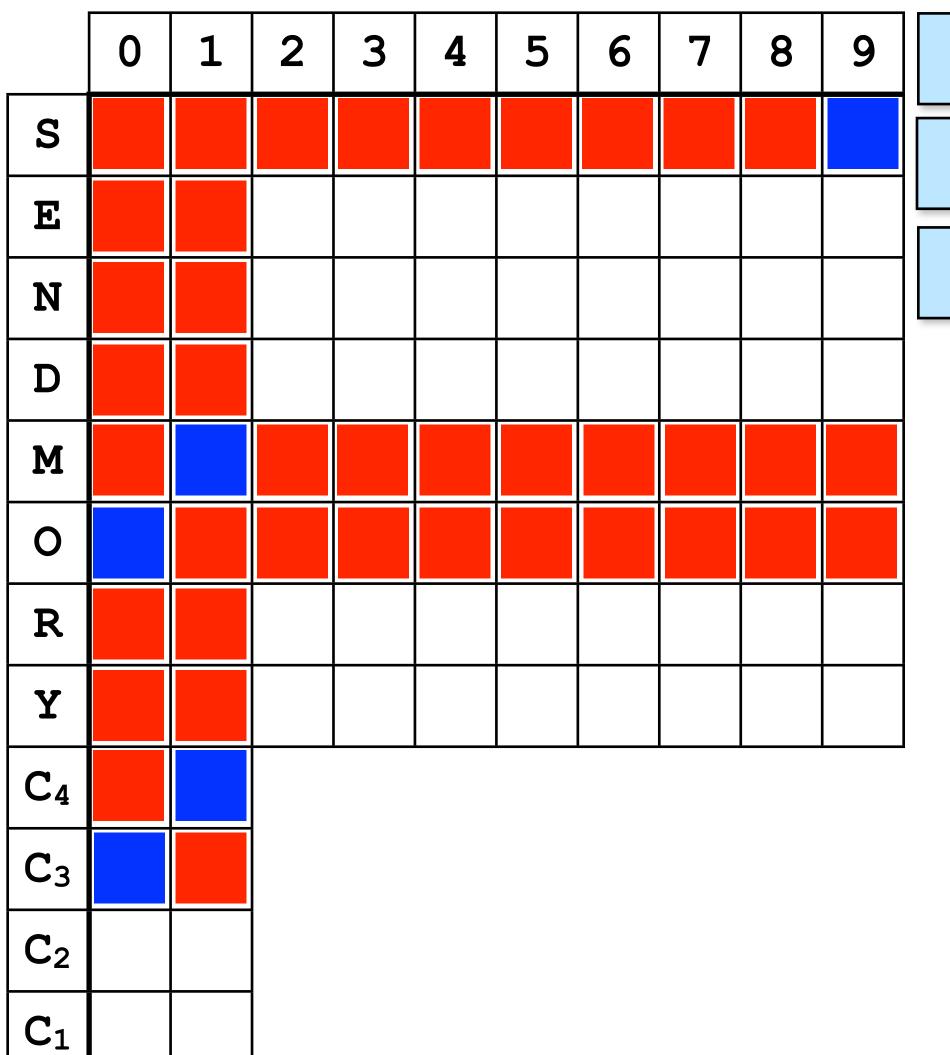
value[S] = 9;
```



 C_3

 C_2

 C_1

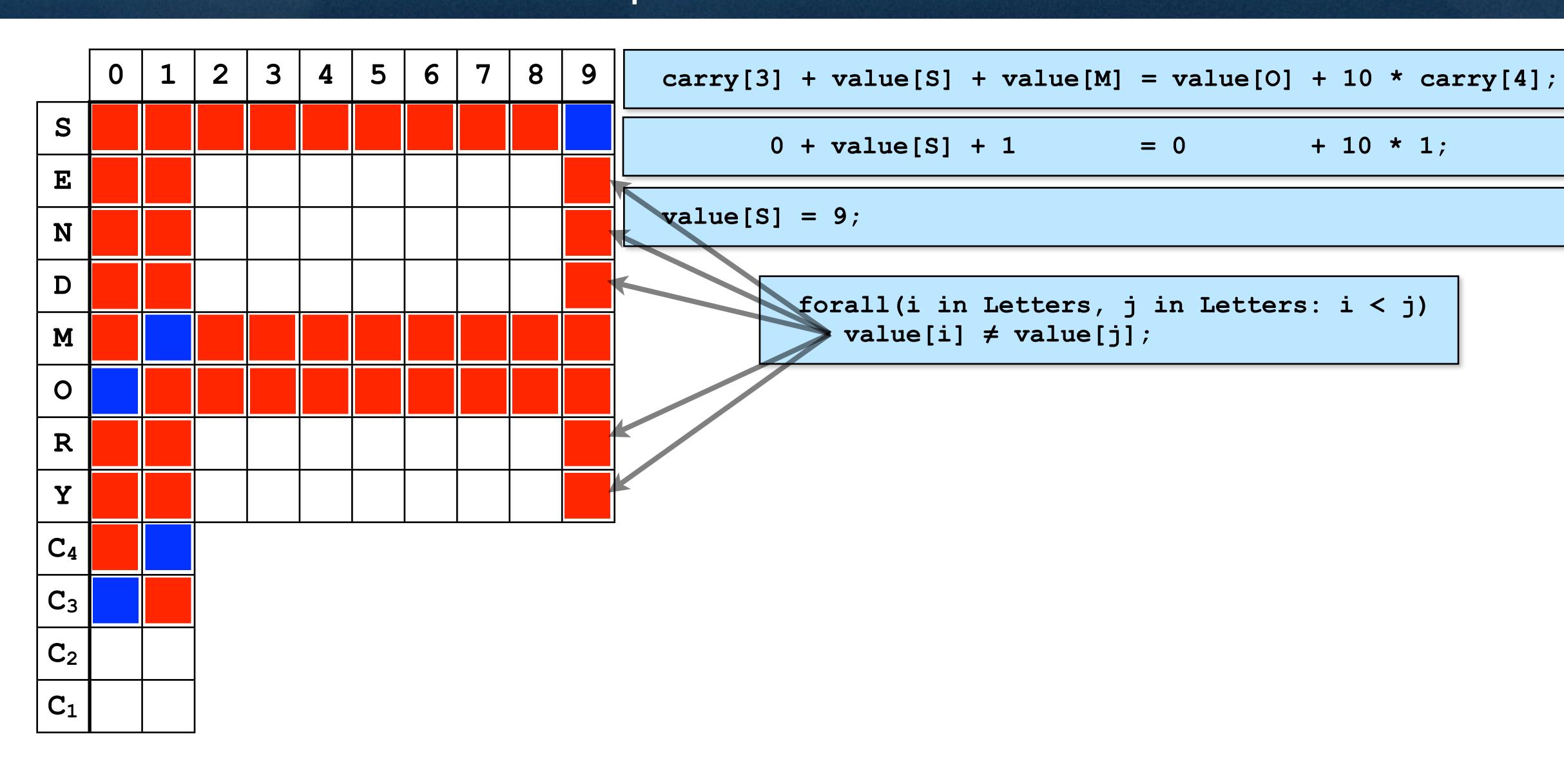


```
carry[3] + value[S] + value[M] = value[0] + 10 * carry[4];

0 + value[S] + 1 = 0 + 10 * 1;

value[S] = 9;
```

```
forall(i in Letters, j in Letters: i < j)
  value[i] ≠ value[j];</pre>
```



	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
0										
R										
Y										
C ₄										
C ₃										
C ₂										
C_1										

Consider a constraint

$$a_1x_1 + \ldots + a_nx_n \ge b_1y_1 + \ldots + b_my_m$$

 $a_i, b_j \geq 0$ are constants

 x_i, y_j are variables with domains $D(x_i), D(y_j)$

Consider a constraint

$$a_1x_1 + \ldots + a_nx_n \ge b_1y_1 + \ldots + b_my_m$$

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Feasibility test

Consider a constraint

$$a_1x_1 + \ldots + a_nx_n \ge b_1y_1 + \ldots + b_my_m$$

 $a_i, b_j \ge 0$ are constants
 x_i, y_j are variables with domains $D(x_i), D(y_j)$

Feasibility test

$$a_1 \max(D(x_1)) + \ldots + a_n \max(D(x_n)) \ge b_1 \min(D(y_1)) + \ldots + b_m \min(D(y_m))$$

Consider a constraint

$$a_1x_1 + \ldots + a_nx_n \ge b_1y_1 + \ldots + b_my_m$$

 $a_i, b_j \ge 0$ are constants
 x_i, y_j are variables with domains $D(x_i), D(y_j)$

Feasibility test

$$a_1 \max(D(x_1)) + \ldots + a_n \max(D(x_n)) \ge b_1 \min(D(y_1)) + \ldots + b_m \min(D(y_m))$$

$$l = a_1 \max(D(x_1)) + \ldots + a_n \max(D(x_n))$$

 $r = b_1 \min(D(y_1)) + \ldots + b_m \min(D(y_m))$

Consider a constraint

$$a_1x_1 + \ldots + a_nx_n \ge b_1y_1 + \ldots + b_my_m$$

 $a_i, b_j \ge 0$ are constants
 x_i, y_j are variables with domains $D(x_i), D(y_j)$

► Pruning

$$a_i x_i \ge r - (l - a_i \max(D(x_i)))$$

$$x_i \ge \left\lceil \frac{r - (l - a_i \max(D(x_i)))}{a_i} \right\rceil \qquad y_j \le \left\lfloor \frac{l - (r - b_j \min(D(y_j)))}{b_j} \right\rfloor$$

$$y_j \le \left\lfloor \frac{l - (r - b_j \min(D(y_j))}{b_j} \right\rfloor$$

Until Next Time